**CS 314: Operating Systems lab** 

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Lab: 4

Part: 1) SJF – shortest job first

#### Introduction:

The class-driven shortest-job-first scheduler, which was implemented in C++, treats each process as a process object. The SJF method makes use of multiple data structures, such as priority queues for CPU bursts and process vectors. Parallel increments are made to a clock variable to maintain keeping tabs on the simulation's time.

The job in the ready queue with the lowest CPU burst time is selected for execution by the SJF scheduler while the CPU is idle. It is also important to remember that the SJF scheduler is not a pre-emptive scheduler, meaning that one job executing on the CPU must finish before another can start.

### **Execution:**

Run: g++ sjf.cpp -o sjf

./sjf process1.dat

# **Analysis:**

 The "convoy effect," in which lengthier processes could make shorter ones wait, can affect non-preemptive SJF and lengthen turnaround times overall.

- Penalty ratio is the ratio of turnaround time to the burst time. In nonpreemptive SJF, shorter processes may wait for longer ones, leading to potentially higher penalty ratios.
- Non-preemptive SJF can result in longer waiting times for shorter processes if they arrive while a longer process is already executing.

```
Process ID=2
Wait Time = 370
Turn Around Time = 661
Penalty Ratio = 2.32143
Process ID=3
Wait Time = 349
Turn Around Time = 550
Penalty Ratio = 2.83684
Process ID=4
Wait Time = 106
Turn Around Time = 119
Penalty Ratio = 12.7778
Process ID=5
Wait Time = 97
Turn Around Time = 102
Penalty Ratio = 20.4
Process ID=6
Wait Time = 1214
Turn Around Time = 1419
Penalty Ratio = 6.9803
Overall Performance
Avg. turnaround time: 755.286
Avg. Waiting time : 544.286
Avg. penalty ratio : 7.41632
CPU Throughput
                    : 0.0049054
```

### Part 2) Shortest Remaining Time First

#### Introduction:

Operating systems use the preemptive scheduling method known as SRTF (Shortest Remaining Time First) to schedule processes. It chooses which process to run next, the one with the shortest burst time left. Reducing the overall processing time and enhancing system responsiveness are the main goals.

Since SRTF is a preemptive algorithm, if a new process with a shorter burst duration comes, the scheduler has the ability to halt the execution of a running process. A preemptive variant of the SJF (Shortest Job First) algorithm is how SRTF is frequently referred to. SRTF makes advantage of the remaining time, whereas SJF schedules depending on the overall burst time of processes. Each process's remaining burst duration dynamically determines its priority. Priority increases with a reduction in remaining time. Because the method is preemptive, it frequently switches contexts, which could result in overhead. Shorter jobs are completed quickly with SRTF, which reduces response time for interactive procedures.

### **Execution:**

Run: g++ srtf.cpp -o srtf

.\srtf process1.dat

## **Analysis:**

- Preemptive SRTF aims to minimize turnaround time by allowing shorter processes to interrupt longer ones, reducing waiting times.
- Penalty ratios can be lower in SRTF because it allows shorter processes to execute promptly, minimizing waiting times.
- Preemptive SRTF reduces waiting times by allowing the scheduler to switch to shorter processes when they arrive, even if longer processes are running.
- SRTF attempts to mitigate the convoy effect by preempting longer processes when shorter ones arrive, minimizing the impact on overall system responsiveness.

Penalty Ratio : 1.70103 Process ID : 3 Turnaround Time: 278 Waiting Time : 77 Penalty Ratio : 1.38308 Process ID : 4 Turnaround Time: 18 Waiting Time : 5 Penalty Ratio : 1.38462 Process ID : 5 Turnaround Time: 8 Waiting Time : 3 Penalty Ratio : 1.6 Process ID : 6 Turnaround Time: 1420 Waiting Time : 1215 Penalty Ratio : 6.92683 Overall Performance Avg. Turnaround Time: 611.571 Avg. Waiting Time : 400.571 Avg. Penalty Ratio : 2.62196 CPU Throughput : 0.00489168

# **Graphs:**







