

→ do  
`pf("Main");`  
`while (i++ < 3);`

## Section 31

### Control statements

- Switch case
- break
- continue
- return

### Switch case:-

- Switch is a conditional non-iterative control statement.
- It is also called as multi-directional conditional control statement.
- Mainly, the switch case is used in a Menu based programs.

### Syntax :-

```
switch (expression)
{
    Case constant1: statement1
    Case constant2: statement2
    default : statement
}
```

```
switch (expression)
{
    Case label1:
        Statement
    Case label2:
        Statement
    default:
        Statement
}
```

- Note:- These labels should be integer family constants. Variables are not allowed.
- 'Break' is optional and 'default' is optional

Example:-

```
void main()
{
```

```
    int num;
    pf(("Enter the number\n"));
    sf("%d", &num);
```

With the labels

switch(num) ← The result of the third exp is  
compared.

```
{ case 1 : pf("1--\n");
    break;
```

Enter  
1  
1 --.

```
case 2 : pf("2--\n");
    break;
```

```
case 3 : pf("3--\n");
    break;
```

```
default : pf("Not 1234\n");
```

}

→ Default could be written anywhere in the middle.

If break statement is not given, the remaining code is also executed:-

eg:- switch(num)

{

```
case 1 : pf("1-\n");
```

Q/P

```
case 2 : pf("2-\n");
```

= Enter a num  
②

```
case 3 : pf("3-\n");
```

2

```
default : pf("Not 123\n");
```

3

Not 123

Valid :-

Case 1 : Case 'a'

Because break is  
not used so

Case 2 & 4 : Case 'a' & 'b'

after for 2 it

Invalid :-

Case "first"

will execute all  
as statements

# Vowels using switch case:-

```

char ch;
if (" Enter which character? ");
    sf ("%c", ch);
switch (ch)
{
    case 'a' :
    case 'A' :
    case 'e' :
    case 'E' :
    case 'i' :
    case 'I' :
    case 'o' :
    case 'O' :
    case 'u' :
    case 'U' :
        printf(" Yeah!");
        break;
    default : printf(" No! ");
}
printf(" Thank you! ");

```

Fig: b1 books

# Break :-

Break is an unconditional non-iterative.

Note:- Break should be within loop or switch case.

Eg:-

```
Void main()
```

```
{
```

```
    pf ("Hello! \n");
```

```
    break; → Treated as syntactical error
```

```
    p ("Bye! \n"); → error.
```

```
}
```

Syntax:-  
break;

## # return :-

- return belongs to unconditional non-iterative control statement.
- It is used to come out from the function. There are two types of return.
  1. Control return
  2. Value return <sup>more</sup> deals with function.

Statement below the keyword 'return' will not execute.

### // Control return example

```
→ {
    int i;
    printf("Hello\n");
    for(i=0; i<5; i++)
        printf("Hi\n");
    return;
}
printf("Bye\n");
```

## # Continue :-

- It is an unconditional <sup>non</sup> iterative control statement.

Cg :-

```
Void main()
{
    while(1)
        continue;
```

control goes from here to main()

Session - 39

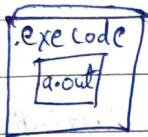
Date 23-7-2021  
Page

## → Secondary data types:

- Memory map
- Introduction to pointers

## Memory Map:-

On compiling a program, e.g. `cc p1.c` is given.  
 The executable file is created in the hard disk.



to execute the code type `./a.out`  
 represents present working directory

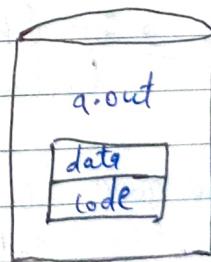
when `./a.out` is typed, this code is loaded into RAM for executions. At that time, OS allocates some memory.

→ The executable file contains 2 sections.

(1.) data section

(2.) code section

H/W

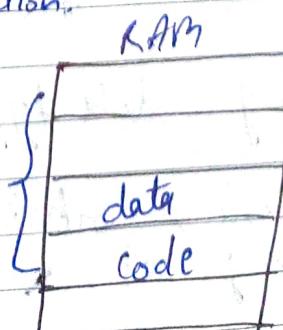


When the exe file is loaded into RAM, 2 more logical sections are allocated by the OS apart from code & data section.

(i) Stack section

(ii) Heap section

a.out



Stack  
heap

This memory is reserved until the program file

gets completed, (i.e. all instructions get executed)

→ In linux, stack section size is 8MB.

Memory section of one a.out file

↓ down growing  
stack.  
stack follows LIFO



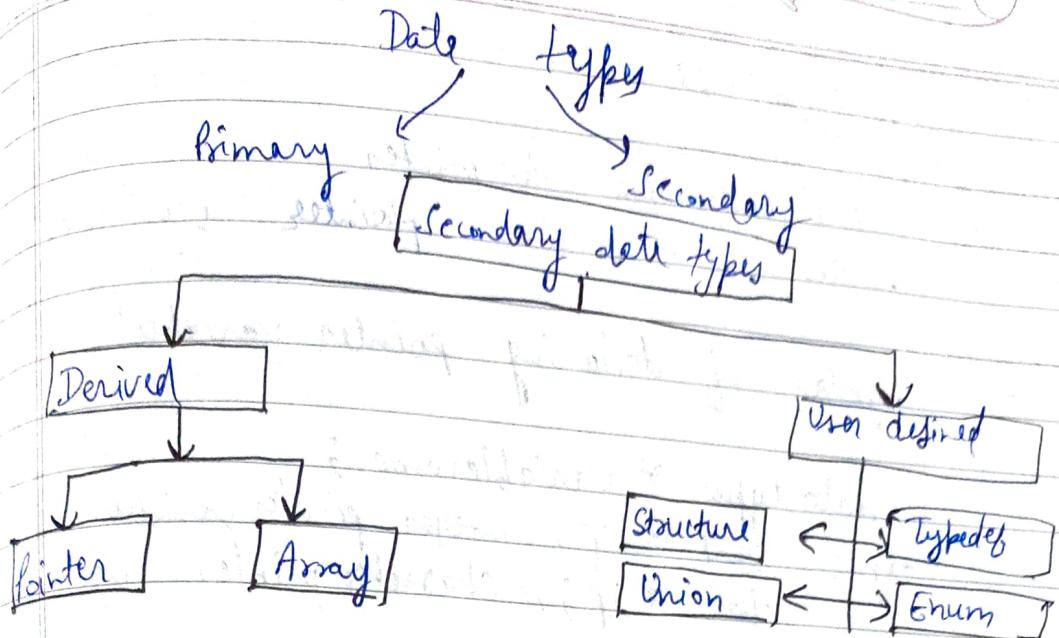
# BSS :- (Block started by symbol)  
= Uninitialised data section

Processors will fetch the instructions from the code section.

The above diagram shows the logical memory blocks created for an a.out file.

# Secondary data types:

Primary:- int, short int, long int, char, double float, string etc.



## # Pointers :- ( Derived Secondary data-types )

→ Pointer is a variable that represents the location of a data item. Pointer stores the address of the variable through which the data can be accessed indirectly.

i [ 10 ] ← Value

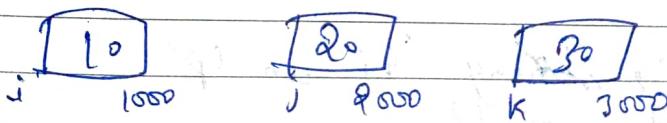
1000

This address can

← Address ↑ be stored in pointer.

→ Any value stored in pointers, it is treated as an address.

eg: int i=10, j=20, k=30



→ Through i, j and k can't be accessed. But there's a way for the above said access, using pointers.

eg:- p=&i → gives i's value using the same  
p= &j → gives j's value ptr variable 'p'  
p= &k → gives k's value all the data can be  
accessed.

→ Using pointers data can be accessed indirectly

## Session - 33

→ Introduction to pointers  
 → How to declare pointer variable

Syntax of declaring pointer variable:-

datatype \* variable name;  
 ex:- int \* ip; // integer pointer variable  
 char \* ip; // character pointer variables.

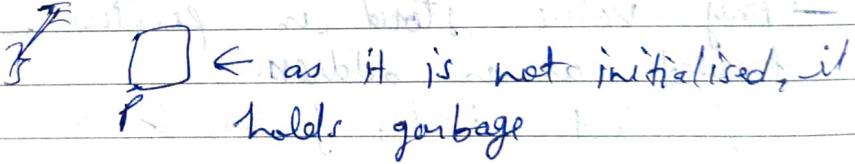
why we need to declare variable:-  
 To inform to the compiler what is what.

wild pointers:- The pointers which are not initialised or uninitialized are called as wild pointers.

e.g:- {

int i = 10, j = 20, k = 30;

int \* p; // p is a pointer variable

 ← as it is not initialised, it holds garbage

Now, storing 'i' in 'p'

→ Void main()

{

int i = 10

int \* p; // till here pointer is wild pointer

p = &i; // & is reference/starting address operator

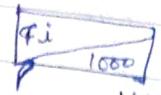
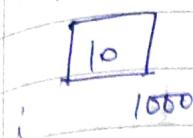
printf("i=%d %p=%d\n", i, p);

for get output dereference operator is

p = &i; ← To store the address in pointer

\*p ← to fetch the 'Data' from the add<sup>n</sup> (dereference operator)

$\%p \rightarrow i=10 \quad *p=10$



address of pointer

$i=10;$   
 $4i=1000$

$p=(4i)=1000$

$*p=4000$

$*p=10$  //  $*p$  will fetch the value inside the address value it holds

→ To access the value stored in an address, a dereference operator is used.

Here  $p$  holds the address of  $i$ ,  $*p$  will fetch the data present in address of  $i$ , which is the data stored in  $p$ .

[ $\%p$  is used to print the address]

Example:

$p=4i;$

`printf("%d %p", i=*p, p=%p\n", i, p);`

→ eg. {

int  $i=10;$

int  $*p;$

$p=4i$

`printf("%d %d", i=*p, p=%p);`

$*p=100$  // Here the value of  $i$  is indirectly changed.

`printf("%d %d", i=%d, *p=%d);`

}

→ Data could be read, written or modified using pointers

→ { int  $i=10; j=20, k=30;$

int  $*p;$

$p=4i;$

$*p=100$ ; // here pointer points to ' $i$ '

$p = \&j;$  // Here pointer points to  $j$

$\&p = \&i;$

$p = \&k;$

$\&p = \&v;$  // pointer points to  $k$

$pf("i=%d j=%d k=%d\n", i, j, k);$

$op = i=100 j=200 k=300$

$\rightarrow \{$

int  $i=10, j=20$

int  $\&p;$

int

$p = \&i;$

$\&p = 100;$  // here value of  $i$  is changed

$j = \&p;$  // now the value of  $i$  (which is 100) is stored

$pf("i=%d j=%d\n", i, j);$  // into  $j$

$\&p = j;$  // this is valid

$pf("i=%d j=%d\n", i, j);$

$i = 20 \quad j = 20;$

Size of pointers :-

{

char ch = 'a';  $\&(ch);$

int i = 10;  $\&i;$

float f = 23.5;  $\&f;$

\* is present here  
to inform the  
compiler that it  
is a pointer variable

$pf("%ld %ld\n", sizeof(ch), sizeof(ip));$

$pf("%ld %ld\n", sizeof(i), sizeof(ip));$

$pf("%ld %ld\n", sizeof(f), sizeof(fp));$

}

Opit 1      8      Same size (The addresses are same)  
                 8      In 64 bit OS, 8 bytes  
                 8      In 32 " " 4 "  
                 Size of pointer depends upon OS.

Session      34

25-7-20

## Secondary datatypes (pointer)

#Difference b/w char \* and int \*

#include < stdio.h >

'&' will give starting address

Void main()

{

int i=258 // size is 4

int \*ip; // size - 8

char \*cp; // size 8

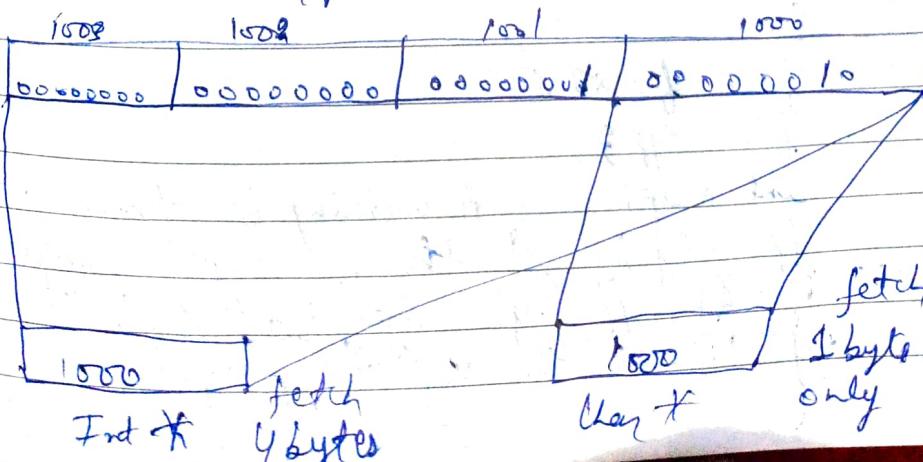
if = &i; // address of i is assigned to the int \*

cp = &i; // and char \*

if ('1' \* ip && \* cp = '1') {  
     if (\* ip == \* cp) {  
         x(ip);  
         x(cp);  
     }
 }

→ The compiler thrown a warning on the line  
 'cp = &i;' because the integer address is stored into char pointer. To avoid this warning, convert int to char explicitly  
 eg cp = (char \*) &i;

int 258      explicit conversion: Little endian



- char pointer → when dereferenced, will fetch only 1 byte  
 → int \* → will fetch 4 bytes, when dereferenced.

$$\text{off} = *i = 258 \quad *cp = 2.$$

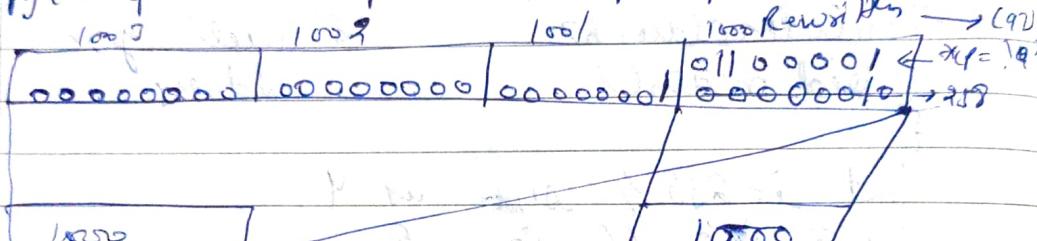
→ Modifying the pointer program a bit.

after  $(p = 'a');$

$\rightarrow cp = 'a';$

'a' is stored in ASCII value

$f(*ip = 7.0 + (p - 'd')n); *ip, *cp;$



So after  $*ip = 'a'$ ,  
 $*ip$  becomes  $\frac{97}{+97}$  Value of  $*ip = 97$ .

as the value of address 1000 is re-written  
 by  $*(p = 'a')$  value

o/p.

$*ip = 353 \rightarrow *ip = 97$

# Difference b/w int \* and float \*

→ void main()

{

float f = 23.5

float \*fp;

int \*ip;

$fp = &f;$

$if = fp;$

$if = fp; // eg:- ip = (int *) &f;$

$fp = (*ip = 7.0 + (p - 'd')n); *ip, *fp;$

}

A warning is thrown. To avoid this, do explicit typecast

e.g.:  $ip = (int *) &f;$

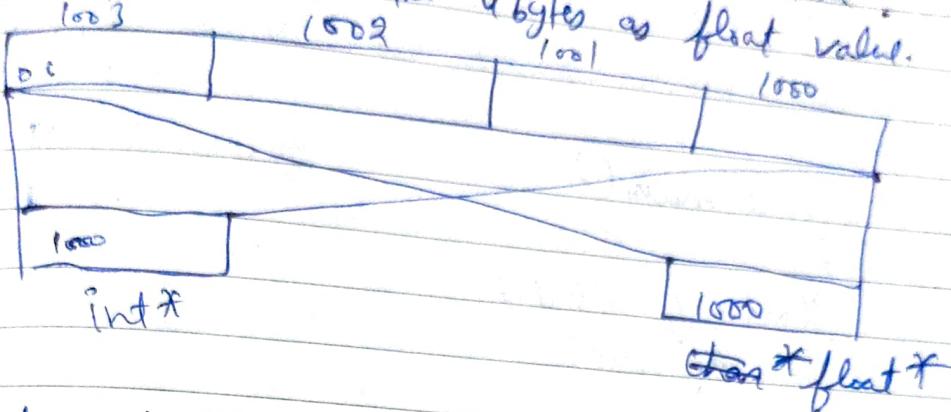
$fp = (float *) ip;$

affi-  
dijf = 1102839808

$$^*ff = 23.50000$$

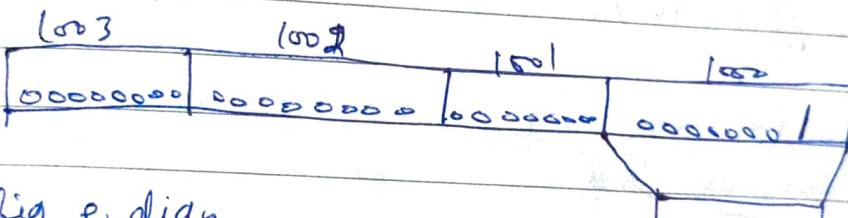
both are same

- int pointer treat the 4bytes as int value
- float pointer treat the 4bytes as float value.

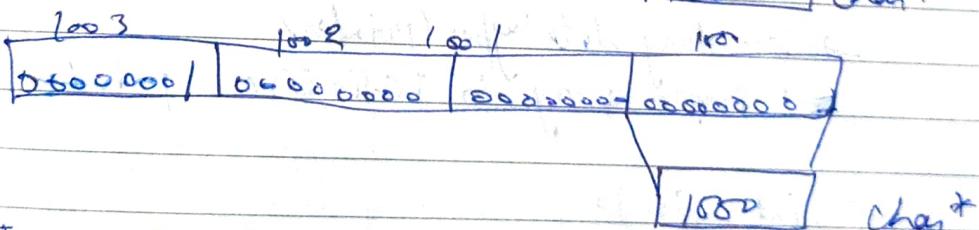


~~Let i=1;~~

little endian



Big endian



This char \* , on a big endian gives up off = 0 as the address 10002 gives us the result '0' char \* , on little endian gives the off = '1'

~~→ than program :-~~

Void main()

8

int i = 1;

Chan \* (p;)

$\{ p = \&i;$   
if ( $*p == 1$ )  
    printf("little \n");  
else  
    printf("big \n");

#

void main()  
{

float af = 23.5;

int \*p, pos;

p = &f;

for (pos = 31; pos >= 0; pos--)

    if ((\*p == 1) && \*p >= pos + 1);

    if ("1n");

}

Sessions 35

Session 38

27-7-20

→ Secondary datatype (pointer)

→ Proper way to initialize pointers

int i = 10

int \*p; // p is an integer pointer variable

p = &i; // the pointer here holds garbage value, as  
        // it is not initialised with any value  
        // & thus, it's a wild ptr.  
This is assigning a address to the  
pointer Variable

#

Syntax of initialise a pointer

datatype \* VariableName = address;

main()

{  
 int i; } These two can't be declared inside  
 int \*i; } the same block with the same  
 name. The compiler throws a syntactical  
 error.

when we write int  $\&p = &i$ ; //the address of i  
 stored into p

X  $\&p = &i$ ; // This is an example of wrong dealing  
 with the pointer  
 Here address of 'i' is stored in  $\&p$ , not in 'p'  
 But the approach →  $\boxed{\text{int } \&p = &i}$  is correct.

int i=10; int \*p;

p = &i; X ← wrong dealing with pointers

\* Address operator can't be written with a constant,  
 the compiler throws syntactical error.

Consider,

① int i=10; ✓

② int \*p; ✓

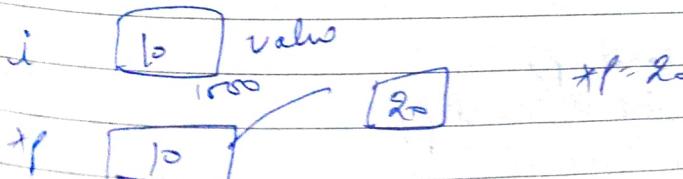
X ③  $\&p = &i$ ; X // pointer is dealing in wrong way!

Here,

① int i=10; ✓

② int  $\&p = &i$ ; X (wrong) ↗

③  $\&p = 20$ ; // ✓ (If warning is ignored, segmentation fault)



In RAM, lower addresses are allocated to OS.

Initially when a pointer has garbage value, it is permitted by the compiler.

But when it is trying to be accessed with a pointer, it generates error/segmentation fault.



{

int i=10;

int \*p;

int \*q;

p=q=&i;

p=q=&i;

\*p = 20;

if (\*p = 70) { q = &d; i = &d; \*p, \*q, i }

will give the  
latest value

\*p = 20; → will go to the address 'p' holds & changes the data there as '20'.

Consider,

\*p = 20;

\*q = 100;

pj("70 70 70", \*p, \*q, i);

pointer, the result is the modified data

# Applying constant term to the pointers:-

① const int \*p; } \*p is constant (i.e) read  
② int \* const p; } only

③ int \* const p;

④ const int \* const p;

⑤ const int \*p;

int i=10;

int \*p;

p = &i; → No error

\*p = 20; → No error

Now,

int i=10;

const int \*p; ← Read only pointer

p = &i;

\*p = 20; ← Error

← Read only

The data can't be modified through the pointer because  $\textcircled{P}$  has only READ permission on the data.

→ int i=10, j=20, k=30;

int \*p;

p = &i;

\*p = 100 ← No error will occur but when pointer is declared as const like below

① const int \*p;      ② p = &i;

\*p = 100; → Error

③ p = &j; ← This is allowed.

Here 'p' can be changed, but '\*p' can't be. 'p' can hold different address values but the address value can't be changed using the pointer.

④ int = const \*p; This is similar to the ① type. In both the cases, p is a read only pointer.

e.g.: { int i=10; int const \*p;

p = &i;

\*p = 20; ← Error

③ int \* const p;

This is a constant pointer.

→ {

int i = 10; → Here the \* points to 6 & which  
int \* const p; → it is trying to point to the diff.  
 p = &i; // error location, compiler generates error  
 \*p = 20;

→ This pointer when initialised to one location,  
 can't be pointed to other locations.

⇒ These type of pointers must be initialised.  
 A wild pointer of this type is no use.

→ e.g.: int \* const p = &i;

\*p = 20; ← This statement is permitted.

Here, the \* once initialised can't be made to point to any other locations, but the data can be changed using this pointer.

④ const int \* const p;

This should be initialised

Session 36

28-7-20

→ Secondary datatypes (pointer)

→ Inc and Dec pointer.

↑ Inc and Dec pointer :-

→ {

int i = 258;

int \*p;

char \*q;

$p = 4i;$ 

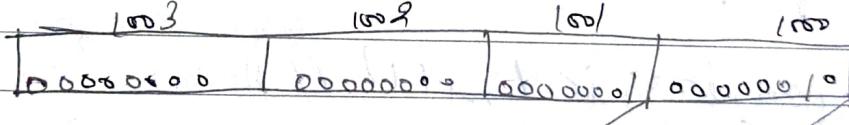
$q = 4i; // \text{warning} \rightarrow \text{int add. assign to char +}$   
 $f("xp\ln", 4i);$   
 $f("xp = \%p\ln", p, p+1); // \%u \text{ can also}$   
 $f("q = \%p\ln", q, q+1); \text{ be used}$

 $q = 0x - - (24)$  $p = 0x2 - - (24) \rightarrow p+1 = 0x2 - - (25) \text{ (4 byte inc)}$  $q = 0x2 - - (24) \rightarrow q+1 = 0x2 - - (25) \text{ (1 byte inc)}$ 

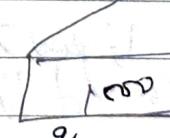
→ Integer pointer focuses 4 bytes

→ Char pointer focuses 1 bytes

Q: {  
 int i=258;  
 char \*q;

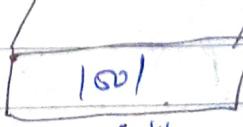
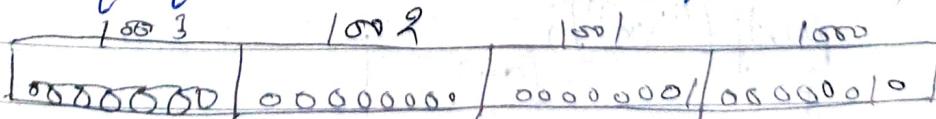
 $q = 4i;$  $f("%d\ln", *q);$  $q = q+1; // \text{the address is incremented}$  $f("%d\ln", *q);$  $j = 258$ 

address inc  
 $\underline{\quad}$



As q is char  
 fetch only 1 byte

when  $q = q+1 = 1000 + 1 = 1001 / \text{add increment by 1 byte}$



Ans = 1

 $00000001$ 

Instead of  $q = q+1 \rightarrow [xq = *q+1] = 259 \text{ Ans}$   
 Value inc

$q = q + 1; //$  address is incremented  
 $\underline{x}q = \underline{x}q + 1 //$  value is incremented.

After this statement executes, the address still remains the same as 1500.

$q + 1; //$   $q$  is not changed  
 $q = q + 1; //$  4 bytes incremented  
 $\underline{x}q = \underline{x}q + 1; //$  value is changed.

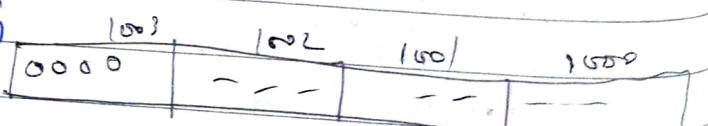
# with a program to display binary format of a float using character pointer.

→ #include <stdio.h>

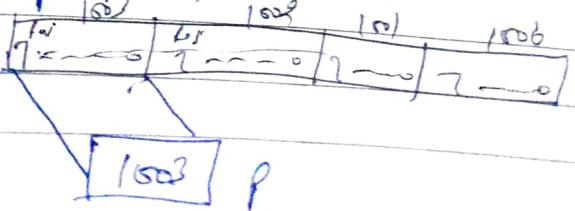
```
void main()
{
    float f = 23.5;
    char *p;
    int i, pos;
    f = 4f;
    f = p + 3;
```

for( $i = 0; i < 4; i++$ ) // As there <sup>outer</sup> lies 4 bytes we
 {
 for( $pos = 7; pos \geq 0; pos--$ ) // 8 times 8 bits
 printf("%d ", \*(p + pos));
 p = p - 1;
 }

printf("\n");



After  $p = p + 3;$



$p = p - 1$  → moves until false.

Now, how

The outer loop is written as  
starting address

for( $p = \&f$ ;  $p = p + 3$ ;  $p >= \&f$ ;  $p = p - 1$ )

}

→ if the char \* is declared as char \* const p;  
the  $p = p + 3$  // will throw an error

Example :-

$j = ++ * p;$

{

int i = 10, j;

int \*p = &i;

ff(" \*p = %p\n", p);

$j = ++ * p;$  //

← Associativity - right to left)

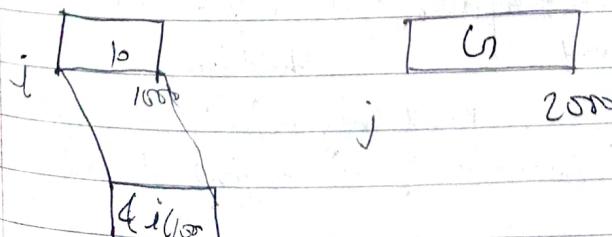
ff(" \*p = %p\n", p);

ff(" i = %d j = %d\n", i, j); // // //

$p = .ox--:-b18$

$p = 6x--b18$

$i = 11 \quad j = 11$



$j = ++ * p;$  there are 3 operators (2 binary  
1 unary) ↑  
unary has the higher priority      L ASS-R to L  
    + +

$j = ++xp;$  (so dereference is first) +  
 ←

$(xp)$  is the value of  $i$ )  
 $j = ++(xp);$  means  $++$  of  $10$  is  $11$   
 (value of  $i$  assigned to  $j$ )

→  $j = xp++;$  address  
 inc first  $p^1$

$x, +t$  same priority (R to L)  
 $j = xp$  after that is becomes  $1004$  in  
 case of int

e.g:-

int  $i = 10, j;$   
 int  $xp, *q;$   
 $q = &i$

suppose, ( $i = 1000$ , so  $f$ 's value is  $1000$ ).

Session - 37

29-7-2020

→ Secondary datatypes (pointer)

→ Inc & Dec pointers

→ Void pointers

→ Null pointers.

initial address $j = 10$ 1000	①	②	③	↓	④	↓	⑤	⑥
$q = p + i;$	$j = \cancel{xp}++;$ htoL pointing So $j$ gets $10$	$j = \cancel{xp}++;$ Data inc $(xp)$ $p$ is inc $f$ pointer goes to $1004$	$j = \cancel{x} + p$ $p$ is inc $f$ is still 1000, $j = 11$		$j = *f(p)$ $p$ becomes 1004, may hold garbage		$j = (\cancel{xp})++$ It is htoL Similar to ③	$j = (\cancel{xp})++$ htoL as Inc post inc

const int \*p; // Read only pointer.

(1)  $++*p; X$  (because  $*p$  is trying to inc.)

(2)  $*++p; \checkmark$

(3)  $*p++; \checkmark$

(4)  $((xp))++; X$  (because  $xp$  is trying to inc).

- int \* const p = &i; // Const pointer  
 (i) ++xp; ✓  
 (ii) xp++; X [Address can't be modified]  
 (iii) xp++ ; X const pointer  
 (iv) ++(xp); ✓  
 (v) (xp)++; ✓

Note:- If const int \* const p = &i; all are error.

# Void pointer or generic pointer:-

{

char ch = 'a';

int i = 10; It is pointer, but behaviour is not float f = 23.5; defined, so it can be used to void \*p; // point any datatype.

p = &ch; // Typecasting not req. when storing address.

printf("%d", \*(char \*)p); // Type cast necessary here

p = &i;

when dereferencing the pt.

((int \*)p = 10) \* (int \*)p; // After; once statement is over,

# p = &f;

it will changes into into

printf("%f", \*(float \*)p); original behaviour is void

}

(int \*)p; X wrong way of typecast.

# Null pointer:- (Null refers to zero)

If a ptr is holding '0'-zero as an address, that ptr is called as null ptr.

→ { int i = 10; :: To avoid wild pointer make it zero  
 int \*p = 0; // This is a null pointer

`pf("rd", *p);  
}`

- \* Note: when the work is finished with pointer, don't leave the pointer as it is. Make the pointer as Null pointer.
- To avoid null / dangling pointers we use null pointer.  
(Null pointer means resulting to zero)

## # Arithmetic operations on the pointer:-

$*p + q *$ ; ✓	Value $p+q$ ;	X	Address
$*p - q *$ ; ✓	are $q-p$ ; $p-q$ ;	✓	and
$xp * q *$ ; ✓	operated $p \times q$ ;	X	operator
$xp / q *$ ; ✓	$p/q$ ;	X	

When we subtract two pointers, we will get how many such type of data we can store in b/w those addresses.

→ {

```
int xp = (int *) 1016;
int *q = (int *) 1000;
pf("rd", p-q);
}
```

O/P:- 4 (1016 byte - 1000 byte gives result in bytes  
↓ format

4 int data type can be stored in b/w p and q.

→ {

```
char *p = (char *) 1016;
char *q = (char *) 1000;
pf("rd", p-q);
}
```

off = 16

16 bytes can be stored in b/w p and q.

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## Secondary data types (Array)

- int array
- array declaration

Array :- Array is a collection of similar types of data. or similar type of data can be put together in Array.  
In an array, similar elements can be put together, so that the processing becomes easy.

Syntax for array declaration :-

datatype VariableName [Elements];

Ex:-

Index operator

- (Ex) int a[5]; → a is an array of 5 int elements
- char ch[10]; → ch is an array of 10 char elements
- float f[5]; → f is an array of 5 float elements

→ If we want to access the elements of an array, we need to use the index operator [ ].

→ Array index starts with zero.

int a[5]	1000	1004	1008	1012	1016
	a[0]	a[1]	a[2]	a[3]	a[4]

Declaration:- a[0] a[1] a[2] a[3] a[4]

→ When an array is declared, it is filled with garbage values.

But we can initialise an array to avoid this  
Ex:- Cg:- int a[5] = {10, 20, 30, 40}