

Quantifying the Noise Tolerance of Quantum Factoring: A Study of Shor's Algorithm

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This project investigates the impact of noise on Shor's algorithm, a quantum algorithm that could break RSA cryptography. Through software simulations on IBM quantum simulator, it aims to provide insights into the challenges of using quantum computers to threaten RSA security.

Introduction

Motivation

RSA encryption, widely used for online security, faces a threat that Quantum Computers (QC) could theoretically use Shor's algorithm to break RSA.

Objectives

- Model the imperfections, i.e. noises of superconducting architectures
 - Explore Shor's algorithm performance under imperfections
- ### Hypothesis
- Success rate ↓ as noise ↑
 - Readout error least sensitive ∴ only occur at measurement
 - Very low success rate factoring numbers larger than 15

How factoring crack RSA?

$$p \cdot q = N$$

- Step 1: Euler's totient function $\phi(n) = (p-1)(q-1)$
Step 2: Generate random e that $\text{GCD}(e, \phi(n)) = 1$
Step 3: Create d by $e \cdot d = 1 \bmod \phi(n)$

Public key (e, N) & Private key (d, N)

- Step 4: Encrypt message M by public key
 $C = M^e \bmod N$
Step 5: Decrypt message by private key
 $M = C^d \bmod N$

Know $e, N, C \rightarrow$ factor $N \rightarrow$ find out p & $q \rightarrow$ deduce value $d \rightarrow$ crack message M

Example to factor $N=15$ [1,2]

- Step 1: Guess any numbers "a". Let $a = 2$.
Step 2: Find order of $a^r = 1 \bmod N$

$$\begin{aligned} 2^0 \% 15 &= 1 \\ 2^1 \% 15 &= 2 \\ 2^2 \% 15 &= 4 \\ 2^3 \% 15 &= 8 \\ 2^4 \% 15 &= 1 \rightarrow \therefore r = 4 \end{aligned}$$

- Step 3: $\text{GCD}(n, a^{(r/2)+1}) \rightarrow \text{GCD}(15, 5) = 5$
 $\text{GCD}(n, a^{(r/2)-1}) \rightarrow \text{GCD}(15, 3) = 3$



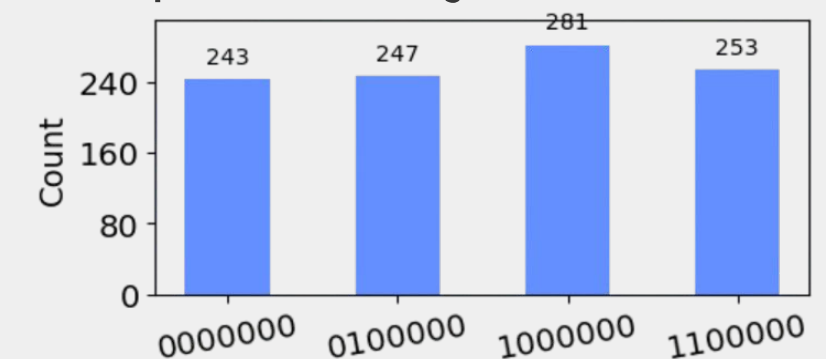
- Step 1: m qubits for control register
 n qubits for work register stores $|\Psi\rangle$. $N=1111 \rightarrow n=4$
Circuit uses " $2n+3$ " qubits $\rightarrow m=7$

- Step 2: Superposition (H)
 $[H|0\rangle]^{\otimes 7} |\Psi\rangle = \frac{1}{2^{n/2}} [|0\rangle + |1\rangle + |2\rangle + \dots + |127\rangle] |1\rangle$

- Step 3: Modular Exponentiation Function
 $\frac{1}{4} [|0\rangle|1\rangle + |1\rangle|2\rangle + |2\rangle|4\rangle + |3\rangle|8\rangle + |4\rangle|1\rangle + \dots + |124\rangle|1\rangle + |125\rangle|2\rangle + |126\rangle|4\rangle + |127\rangle|8\rangle]$
 $= \frac{1}{4} ([|0\rangle + |4\rangle + \dots + |120\rangle + |124\rangle] \otimes |1\rangle + \dots + [|3\rangle + |7\rangle + \dots + |123\rangle + |127\rangle] \otimes |8\rangle)$

- Step 4: $\text{QFT}^\dagger \rightarrow$ measure
Inverse Quantum Fourier Transform (QFT^\dagger) to perform Quantum Phase Estimation (QPE) \rightarrow determine the phase of eigenvalues of operators by $\phi = \ell/2^m$

- Peaks ℓ occur at
 $0, 32, 64, 96$
Phases ϕ (s/r) are
 $0, 1/4, 1/2, 3/4$



Methodology

- 1 Implement Shor's algorithm
Develop quantum circuit & post-quantum processing procedure

Build 3 types of noise models

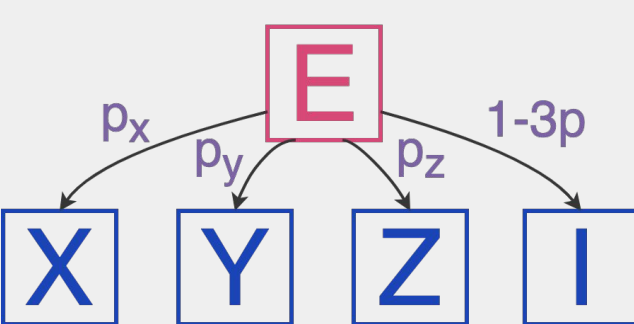
Sample IBM 15 real QC error data [3] & research findings [4]

- 3 Vary error rates
Vary error rates between min. & max. to implement circuit

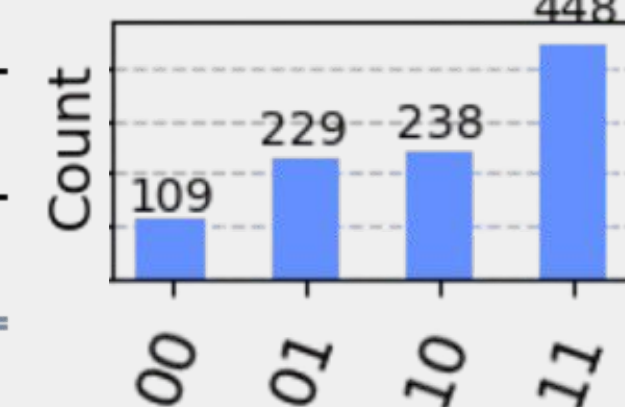
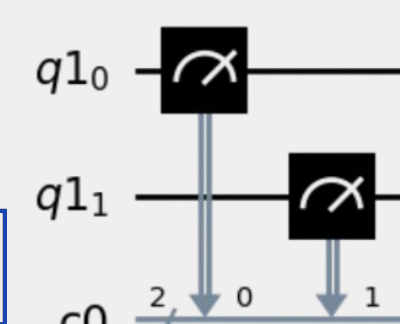
- Analysis
Analyse factoring success rate and compare to noise-free results

Types of noise

Depolarising



Bit-flip (X), phase-flip (Y), combined flip (Z) occurs with same probability



State	Theoretical	Simulated
00>	$\frac{1}{3} \cdot \frac{1}{3} = 0.111$	$109/1024 = 0.106$
01>	$\frac{1}{3} \cdot \frac{2}{3} = 0.222$	$229/1024 = 0.224$
10>	$\frac{2}{3} \cdot \frac{1}{3} = 0.222$	$238/1024 = 0.232$
11>	$\frac{2}{3} \cdot \frac{2}{3} = 0.444$	$448/1024 = 0.438$

Thermal Relaxation

- (i) Thermal relaxation occurs over time in the form of excitation / de-excitation (T_1 = relaxation constant)
(ii) Dephasing of qubit overtime (T_2 = dephasing constant)

Readout

Error in measuring qubits

Results

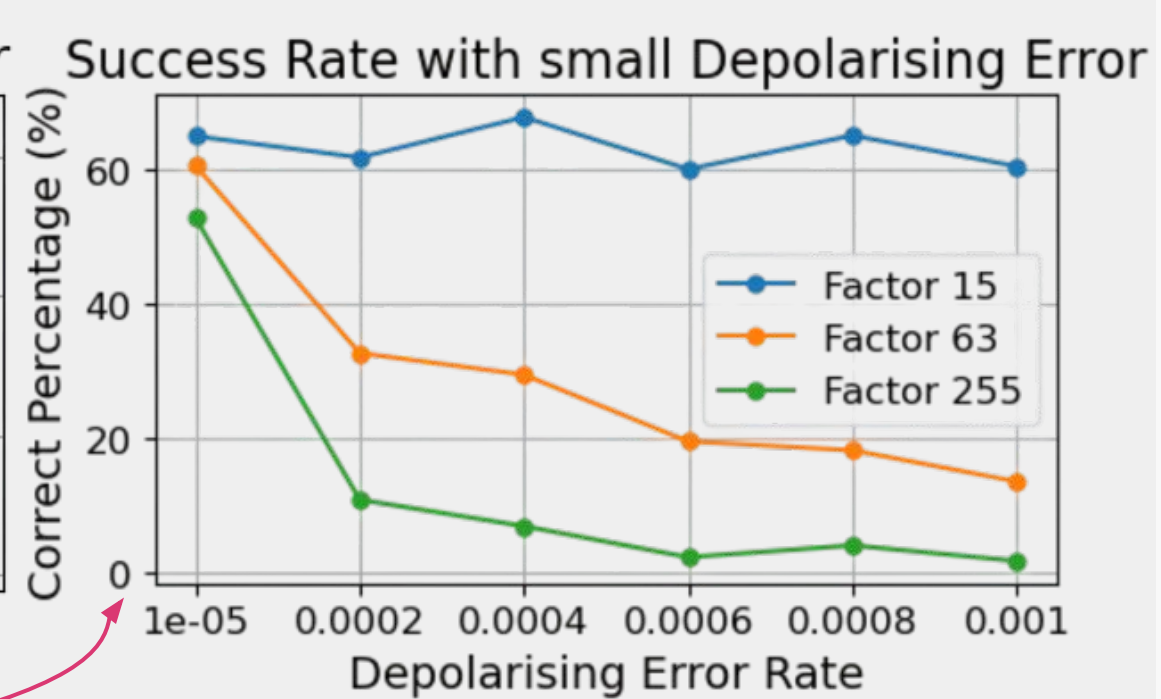
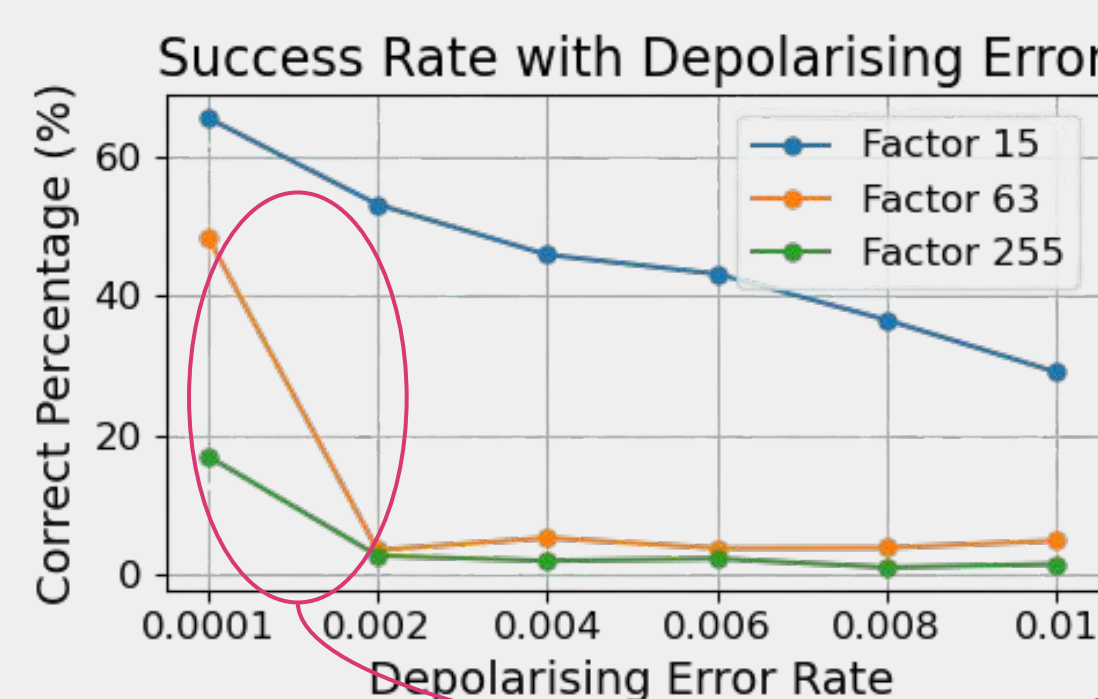
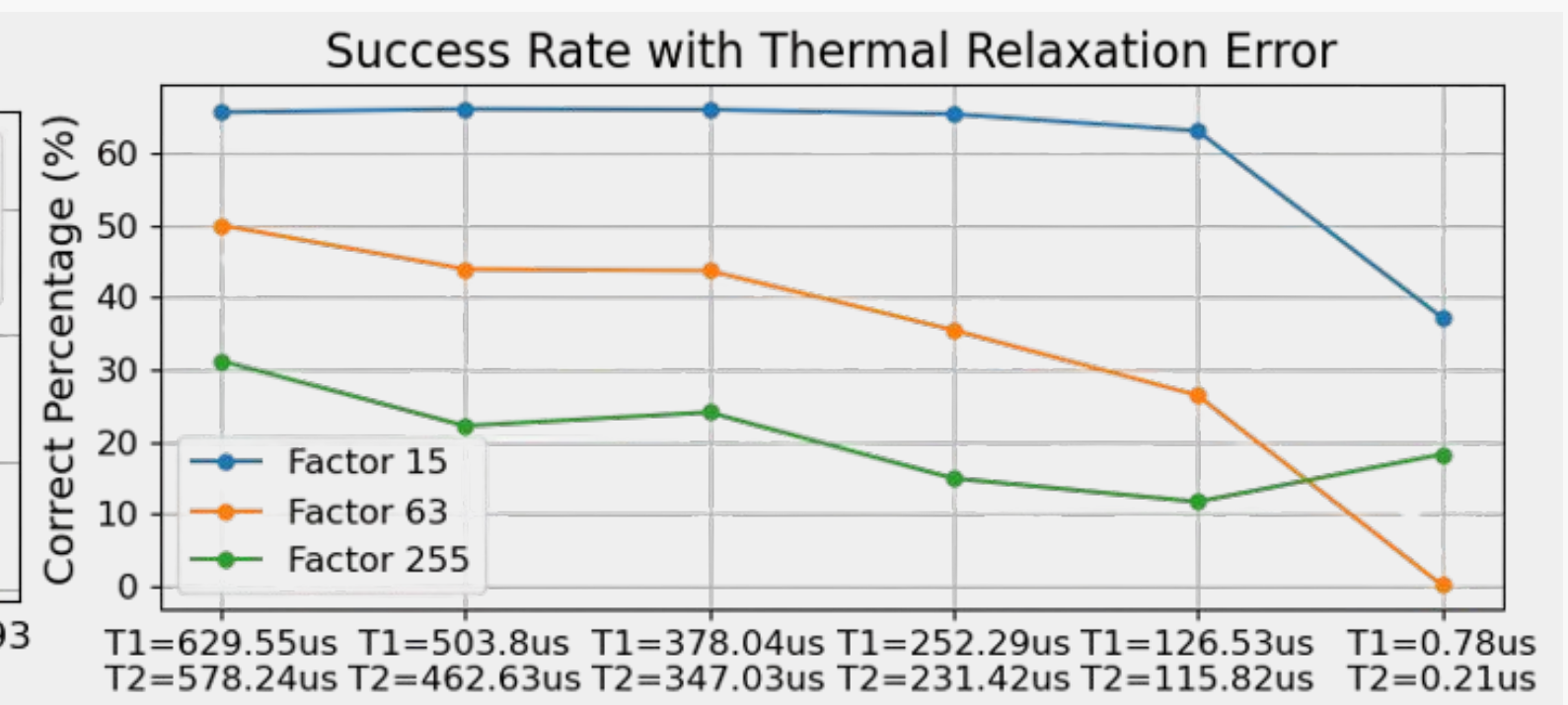
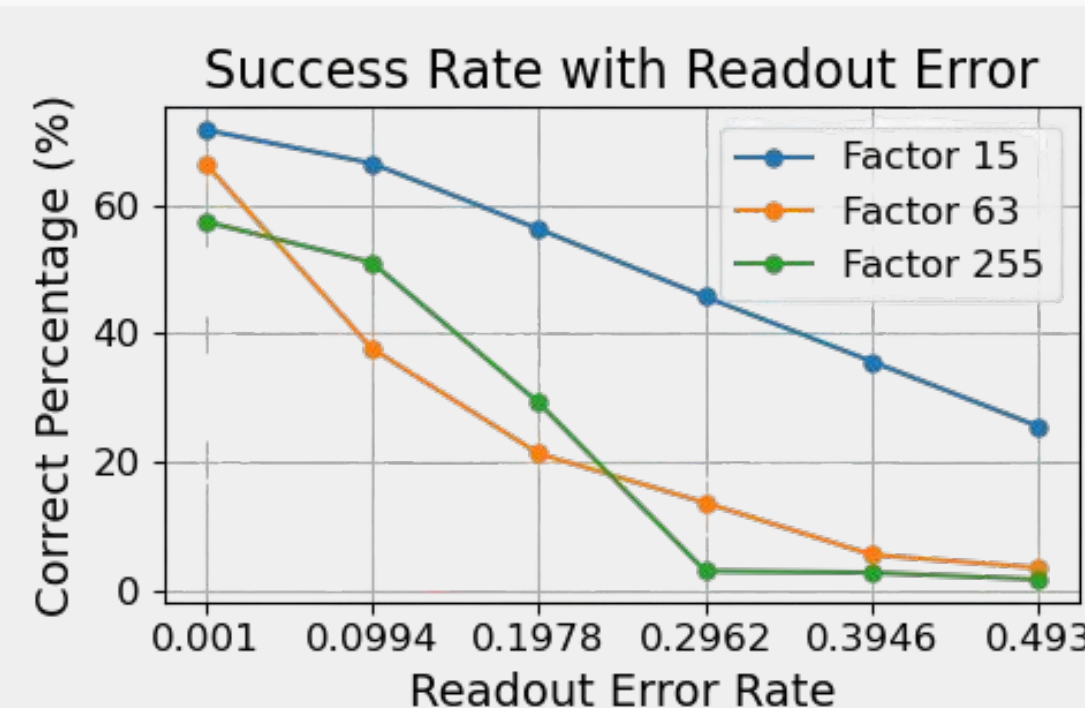
Line Graphs Analysis

- Readout Error:
 - Significantly drop close to 0 since median error rate for factoring 255
 - Thermal Relaxation Error:
 - Non sensitive to small circuit
 - Depolarising Error:
 - Biggest impact among the three errors
 - Sensitive between error = $1e-5$ to $2e-3$
- ### Table Analysis
- Readout is the least impact among the three errors
 - Real noisy QC has <5% success rate for factoring 63, 255.
 - Experimental noise models reveals depolarising error is the major challenge to Shor's algorithm

Discussion

- Fluctuation see in line graphs could be due to the nature of gates are not perfect, e.g. Hadamard gates \times produce 50/50 outcomes.
- Still far away from breaking RSA-2048.
- Future plan: error mitigation for depolarising error.

I would like to thank Dr. Imad Faruque and Dr. Jorge Barreto for their support throughout this research.



* Real QC backend ibm_brisbane	Factor success rate	Factor success rate		
		15	63	255
15 (1111)	70.65%	66.67%	56.44%	59.71%
63 (111111)	4.96%	47.57%	3.47%	2.40%
255 (11111111)	2.55%	64.14%	26.99%	10.49%
		65.05%	56.40%	52.92%