

Transactions

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Database Systems I. Seminar

2018 Spring Semester



Recap

- Transaction properties:
 Atomicity, Consistency,
 Isolation, Durability;
 (ACID)
- Scheduling
- Conflicting transactions
- Conflict equivalent schedules
- View equivalent schedules
- Conflict serializable

- precedence graph
- locking techniques
- two phase protocol



 We have these two transaction, what is the result if we run: S1(T1,T2), and what is the result if we run S2(T2,T1). The initial values are: x=15; y=5; N=5; M=10

T1	T2
read(x)	read(x)
x:=x-N	x:=x+M
write(x)	write(x)
read(y)	
y:=y+N	
write(y)	



 Given the following scheduling, what will be the result with the same initial values: x=15; y=5; N=5; M=10

T1	T2
read(x)	
x:=x-N	
	read(x)
	x:=x+M
write(x)	
read(y)	
	write(x)
y:=y+N	
write(y)	



• Is the following schedule conflict serializable?

T1	T2
read(x)	
write(x)	
	read(x)
read(y)	
	write(x)
write(y)	



Are the following two schedules view equivalent?

T1	T2	T1	T2
read(X) 1		read(X) 1	
X:=X-N 2		X:=X+M 2	
	3 read(X)		3 read(X)
	4 X:=X+M		4 X:=X-N
write(X) 5			5 write(X)
Read(Y) 6		write(X) 6	
	7 write(X)	read(Y) 7	
Y:=Y+N 8		Y:=Y+N 8	
write(Y) 9		write(Y) 9	



 Is the following schedule conflict serializable? (draw the precedence graph)

T1	T2	T3	T4
	read(A)		
read(B)			
write(A)			
		read(C)	
		write(B)	
read(C)			
			read(B)



• What will be the result of S1(T1,T2) and S2(T2,T1), if the initial values are: X=20, Y=10?

T1	T2	
s-lock(y)	s-lock(x)	
Read(y)	Read(x)	
Unlock(y)	Unlock(x)	
x-lock(x)	s-lock(y)	
read(x)	Read(y)	
X:=X-Y/2	Unlock(y)	
Write(x)	x-lock(y)	
Unlock(x)	read(y)	
s-lock(y)	Y:=Y+X	
read(y)	Write(y)	
unlock(y)	unlock(y)	



Mark the two phases in the following scheduling:

```
T1
            T2
s-lock(Y); s-lock(X);
read(Y); read(X);
x-lock(X); x-lock(Y);
unlock(Y); unlock(X);
read(X); read(Y);
X:=X+Y; Y:=X+Y;
write(X); write(Y);
unlock(X); unlock(Y);
```