P	$::=$	$R\ \overline{C}\ e$	*program*
R	$::=$	$\overline{\mathtt{m} \leq \mathtt{m}'}$	*mode order*
C	$::=$	**class** $\mathtt{c}\ \Delta$ **extends** $\mathtt{c}'\{\overline{F}\ \overline{M}\ A\}$	*class*
F	$::=$	$T\ \mathtt{fd} = e$	*field*
M	$::=$	$T\ \mathtt{md}(\overline{T}\ \overline{\mathtt{x}})\{e\}$	*method*
A	$::=$	e	*attributor*
e	$::=$	$\mathtt{x} \mid e.\mathtt{fd} \mid$ **new** $\mathtt{c}\langle\iota\rangle \mid e.\mathtt{md}(\overline{e})$	*expressions*
	$\mid$	$(T)e \mid$ **snapshot** $e\ [\eta, \eta] \mid e \triangleright \eta$	
	$\mid$	$\{\overline{\mathtt{m} : e}\}^T$	

Figure 1. Syntax

T	$::=$	$\mathtt{c}\langle\iota\rangle \mid$ **mcase**$\langle T\rangle$	*programmer type*
ι	$::=$	$\overline{\eta} \mid ?, \overline{\eta}$	*object mode parameter list*
η	$::=$	$\mathtt{m} \mid \mathtt{mt} \mid \top \mid \bot$	*static mode*
μ	$::=$	$\eta \mid ?$	*mode*
$\mathtt{mt}$			*mode type variable*
$?$			*dynamic mode type*
ω	$::=$	$\eta \leq \mathtt{mt} \leq \eta'$	*constrained mode*
Δ	$::=$	$? \to \omega, \Omega \mid \Omega$	*class mode parameter list*
Ω	$::=$	$\overline{\omega}$	*constrained mode list*
τ	$::=$	$T \mid \exists\omega.\tau \mid \mathtt{modev} \mid \mathtt{mt}$	*type*
$\mathtt{K}$	$::=$	$\overline{\eta \leq \eta'}$	*constraints*

Figure 2. Type Elements

$$\text{(WF-Class)}\ \frac{\textbf{class}\ \mathtt{c}\ \Delta\ \textbf{extends}\ \mathtt{c}'\ \cdots \in P \qquad \mathtt{iparam}(\Delta) = \iota' \quad \mathtt{cons}(\Delta) = \mathtt{K}' \quad \mathtt{K} \models \mathtt{K}'\{\overline{\eta}/\iota'\}}{\mathtt{K} \vdash_{\mathtt{wft}} \mathtt{c}\langle\overline{\eta}\rangle}$$

$$\text{(WF-ClassDyn)}\ \frac{\textbf{class}\ \mathtt{c}\ ? \to \omega, \Omega\ \textbf{extends}\ \mathtt{c}' \cdots \in P \qquad \mathtt{iparam}(? \to \omega, \Omega) = \iota' \qquad \mathtt{cons}(\Omega) = \mathtt{K}' \quad \mathtt{K} \models \mathtt{K}'\{\overline{\eta}/\iota'\}}{\mathtt{K} \vdash_{\mathtt{wft}} \mathtt{c}\langle ?, \overline{\eta}\rangle}$$

$$\text{(WF-Top)}\ \mathtt{K} \vdash_{\mathtt{wft}} \mathtt{Object}\langle\eta\rangle$$

$$\text{(WF-Exist)}\ \frac{\omega = \eta \leq \mathtt{mt} \leq \eta' \qquad \mathtt{K} = \mathtt{K}' \cup \{\eta \leq \mathtt{mt}, \mathtt{mt} \leq \eta'\} \quad \mathtt{mt} \notin \mathtt{K}' \quad \mathtt{K} \vdash_{\mathtt{wft}} \tau}{\mathtt{K} \vdash_{\mathtt{wft}} \exists\omega.\tau}$$

$$\text{(WF-MCase)}\ \frac{\mathtt{K} \vdash_{\mathtt{wft}} T}{\mathtt{K} \vdash_{\mathtt{wft}} \textbf{mcase}\langle T\rangle}$$

Figure 3. Type Well-Formedness

$$\text{(WF-Empty)}\ P \vdash_{\mathtt{wfe}} \epsilon$$

$$\text{(WF-ESpec)}\ \frac{P \vdash_{\mathtt{wfe}} \Omega \qquad \eta \leq \eta'}{P \vdash_{\mathtt{wfe}} \Omega, \eta \leq \mathtt{mt} \leq \eta'}$$

$$\text{(WF-TSpec)}\ \frac{P \vdash_{\mathtt{wfe}} \omega, \Omega}{P \vdash_{\mathtt{wfe}} ? \to \omega, \Omega}$$

Figure 4. Environment Well-Formedness

$$\text{(FD-Object)}\ \mathtt{fields}(\mathtt{Object}\langle\eta\rangle) = \bullet$$

$$\text{(FD-Class)}\ \frac{\textbf{class}\ \mathtt{c}\ \Delta\ \textbf{extends}\ \mathtt{c}'\{\overline{T\ \mathtt{fd} = e}\ \overline{M}\ A\} \qquad \mathtt{eparam}(\Delta) = \iota' \quad \mathtt{fields}(\mathtt{c}'\langle\iota\rangle) = \overline{T_0\ \mathtt{fd}_0 = e_0}}{\mathtt{fields}(\mathtt{c}\langle\iota\rangle) = \overline{T_0\ \mathtt{fd}_0 = e_0}, \overline{T\{\iota/\iota'\}\ \mathtt{fd} = e\{\iota/\iota'\}}}$$

$$\text{(MT-Class)}\ \frac{\textbf{class}\ \mathtt{c}\ \Delta\ \textbf{extends}\ \mathtt{c}'\{\overline{F}\ \overline{M}\ A\} \qquad T\ \mathtt{md}(\overline{T}\ \overline{\mathtt{x}})\ \{\ e\ \} \in \overline{M} \quad \mathtt{eparam}(\Delta) = \iota'}{\mathtt{mtype}(\mathtt{md}, \mathtt{c}\langle\iota\rangle) = (\overline{T} \to T)\{\iota/\iota'\}}$$

$$\text{(MT-Super)}\ \frac{\textbf{class}\ \mathtt{c}\ \Delta\ \textbf{extends}\ \mathtt{c}'\{\overline{F}\ \overline{M}\ A\} \quad \mathtt{md} \notin \overline{M}}{\mathtt{mtype}(\mathtt{md}, \mathtt{c}\langle\iota\rangle) = \mathtt{mtype}(\mathtt{md}, \mathtt{c}'\langle\iota\rangle)}$$

$$\text{Override}\ \frac{\mathtt{mtype}(\mathtt{md}, T) = \overline{T'} \to T_0' \quad \mathtt{K} \vdash T_0 <: T_0'}{\mathtt{override}(\mathtt{md}, T, \mathtt{K}, \overline{T} \to T_0)}$$

$$\text{(MB-Class)}\ \frac{\textbf{class}\ \mathtt{c}\ \Delta\ \textbf{extends}\ \mathtt{c}'\{\overline{F}\ \overline{M}\ A\} \qquad \mathtt{iparam}(\Delta) = \iota' \quad T\ \mathtt{md}(\overline{T}\ \overline{\mathtt{x}})\{e\} \in \overline{M}}{\mathtt{mbody}(\mathtt{md}, \mathtt{c}\langle\iota\rangle) = \overline{\mathtt{x}}.e\{\iota/\iota'\}}$$

$$\text{(MB-Super)}\ \frac{\textbf{class}\ \mathtt{c}\ \Delta\ \textbf{extends}\ \mathtt{c}'\{\overline{F}\ \overline{M}\ A\} \quad \mathtt{md} \notin \overline{M}}{\mathtt{mbody}(\mathtt{md}, \mathtt{c}\langle\iota\rangle) = \mathtt{mbody}(\mathtt{md}, \mathtt{c}'\langle\iota\rangle)}$$

$$\text{(AB-Class)}\ \frac{\textbf{class}\ \mathtt{c}\ \Delta\ \textbf{extends}\ \mathtt{c}'\{\overline{F}\ \overline{M}\ A\} \qquad \mathtt{eparam}(\Delta) = \iota' \quad A = e}{\mathtt{abody}(\mathtt{c}\langle\iota\rangle) = e\{\iota/\iota'\}}$$

Figure 5. FJ Functions

$$\text{(T-Program)}\ \frac{R \text{ form a latice} \quad \varnothing \vdash e \quad \overline{C}\ \text{OK}}{R\ \overline{C}\ e\ \text{OK}}$$

$$\text{(T-Class)}\ \frac{\mathtt{K} = \mathtt{cons}(\Delta) \qquad \overline{M}\ \text{OK IN}\ \mathtt{c}, \Delta \quad A\ \text{OK IN}\ \mathtt{c}, \Delta \quad \overline{F} = \overline{T}\ \overline{\mathtt{fd}} = \overline{e} \qquad \varnothing; \mathtt{K} \vdash \overline{e} : \overline{T} \quad \textbf{class}\ \mathtt{c}'\ \Delta\ \textbf{extends}\ \mathtt{c}''\{\ldots\}\ \text{FJ OK}}{\textbf{class}\ \mathtt{c}\ \Delta\ \textbf{extends}\ \mathtt{c}'\{\overline{F}\ \overline{M}\ A\}\ \text{OK}}$$

$$\text{(T-Attributor)}\ \frac{\iota = \mathtt{iparam}(\Delta) \quad \mathtt{K} = \mathtt{cons}(\Delta) \qquad A = e \quad \Delta; \textbf{this} : \mathtt{c}\langle\iota\rangle \vdash e : \mathtt{modev} \quad \mathtt{K} \vdash_{\mathtt{wft}} \mathtt{c}\langle\iota\rangle}{A\ \text{OK IN}\ \mathtt{c}, \Delta}$$

$$\text{(T-Method)}\ \frac{\iota = \mathtt{iparam}(\Delta) \qquad \mathtt{K} = \mathtt{cons}(\Delta) \quad \overline{x} : \overline{T}; \textbf{this} : \mathtt{c}\langle\iota\rangle; \mathtt{K} \vdash e : T \qquad \mathtt{override}(\mathtt{md}, \mathtt{c}\langle\iota\rangle, \mathtt{K}, \overline{T} \to T) \quad \mathtt{K} \vdash_{\mathtt{wft}} \mathtt{c}\langle\iota\rangle}{T\ \mathtt{md}(\overline{T}\ \overline{x})\{\ e\ \}\ \text{OK IN}\ \mathtt{c}\ \Delta}$$

Figure 6. Class Typing

$$\text{(T-Var)}\ \frac{}{\Gamma;\mathtt{K}\vdash \mathtt{x} : \Gamma(\mathtt{x})}$$

$$\text{(T-New)}\ \frac{\begin{array}{c}\iota = ?, \iota' \text{ iff } \textbf{class}\ \mathtt{c}\ \Delta \cdots \in P \text{ and } \mathtt{ethis}(\Delta) = ?\\ \iota \neq ?, \iota' \text{ iff } \textbf{class}\ \mathtt{c}\ \Delta \cdots \in P \text{ and } \mathtt{ethis}(\Delta) \neq ? \qquad \mathtt{K} \models \mathtt{cons}(\Delta)\end{array}}{\Gamma;\mathtt{K}\vdash\ \textbf{new}\ \mathtt{c}\langle\iota\rangle : \mathtt{c}\langle\iota\rangle}$$

$$\text{(T-Cast)}\ \frac{\Gamma;\mathtt{K}\vdash e : T'}{\Gamma;\mathtt{K}\vdash (T)e : T}$$

$$\text{(T-Msg)}\ \frac{\begin{array}{c}\Gamma;\mathtt{K}\vdash\ e : \mathtt{c}\langle\iota\rangle \qquad \mathtt{mtype}(\mathtt{md},\mathtt{c}\langle\iota\rangle) = \overline{T}\to T\\ \Gamma;\mathtt{K}\vdash \overline{e} : \overline{T} \qquad \Gamma;\mathtt{K}\vdash\ \textbf{this} : T_{this}\\ \mathtt{K}\models \{\mathtt{mode}(\mathtt{c}\langle\iota\rangle)\leq \mathtt{mode}(T_{this})\} \qquad \mathtt{mode}(\mathtt{c}\langle\iota\rangle)\neq ?\end{array}}{\Gamma;\mathtt{K}\vdash e.\mathtt{md}(\overline{e}) : T}$$

$$\text{(T-Field)}\ \frac{\begin{array}{c}\Gamma;\mathtt{K}\vdash e : \mathtt{c}\langle\iota\rangle \qquad \Gamma;\mathtt{K}\vdash \textbf{this} : T_{this} \qquad \mathtt{fields}(\mathtt{c}\langle\iota\rangle) = \overline{T}\ \overline{\mathtt{fd}}\\ \mathtt{K}\models \{\mathtt{mode}(\mathtt{c}\langle\iota\rangle)\leq \mathtt{mode}(T_{this})\} \qquad \mathtt{mode}(\mathtt{c}\langle\iota\rangle)\neq ?\end{array}}{\Gamma;\mathtt{K}\vdash e.\mathtt{fd}_i : T_i}$$

$$\text{(T-Snapshot)}\ \frac{\Gamma;\mathtt{K}\vdash e : \mathtt{c}\langle ?,\iota\rangle \qquad \omega = \eta_1 \leq \mathtt{mt} \leq \eta_2}{\Gamma;\mathtt{K}\vdash \textbf{snapshot}\ e\,[\eta_1,\eta_2] : \exists\omega.\mathtt{c}\langle\mathtt{mt},\iota\rangle}$$

$$\text{(T-MCase)}\ \frac{\overline{\mathtt{m}} = \mathtt{modes}(P) \qquad \Gamma;\mathtt{K}\vdash e_i : T \text{ for all } i}{\Gamma;\mathtt{K}\vdash \{\overline{\mathtt{m}:e}\}^T : \textbf{mcase}\langle T\rangle}$$

$$\text{(T-ElimCase)}\ \frac{\Gamma;\mathtt{K}\vdash e : \textbf{mcase}\langle T\rangle \qquad \eta\in\mathtt{modes}(P) \text{ or } \eta \text{ appears in } \mathtt{K}}{\Gamma;\mathtt{K}\vdash e \rhd \eta : T}$$

$$\text{(T-ModeValue)}\ \frac{\mathtt{m}\in\mathtt{modes}(P)}{\Gamma;\mathtt{K}\vdash \mathtt{m} : \mathtt{modev}}$$

$$\text{(T-Sub)}\ \frac{\Gamma;\mathtt{K}\vdash e : \tau \qquad \mathtt{K}\vdash \tau <: \tau'}{\Gamma;\mathtt{K}\vdash e : \tau'}$$

Figure 7. Expression Typing

$$\text{(S-Dynamic)}\ \mathtt{K}\vdash \mathtt{c}\langle\mu;\overline{\eta}\rangle <: \mathtt{c}\langle ?;\overline{\eta}\rangle$$

$$\text{(S-Mcase)}\ \frac{\mathtt{K}\vdash\tau <: \tau'}{\mathtt{K}\vdash \textbf{mcase}\langle\tau\rangle <: \textbf{mcase}\langle\tau'\rangle}$$

$$\text{(S-ExistOpen)}\ \frac{\begin{array}{c}\omega = \eta_1\leq\mathtt{mt}\leq\eta_2 \qquad \mathtt{K}' = \mathtt{K}\cup\{\eta_1\leq\mathtt{mt},\mathtt{mt}\leq\eta_2\}\\ \mathtt{mt}\notin\mathtt{K}' \qquad \mathtt{K}'\vdash\tau <: \tau'\end{array}}{\mathtt{K}\vdash\exists\omega.\tau <: \tau'}$$

$$\text{(S-ExistAbstract)}\ \frac{\omega = \eta_1\leq\mathtt{mt}\leq\eta_2 \qquad \mathtt{K}\models\{\eta_1\leq\mathtt{mt},\mathtt{mt}\leq\eta_2\}}{\mathtt{K}\vdash\tau <: \exists\omega.\tau}$$

$$\text{(S-Class)}\ \frac{\textbf{class}\ \mathtt{c}\ \Delta\ \textbf{extends}\ \mathtt{c}'\cdots\in P \qquad \mathtt{K}\models\mathtt{cons}(\Delta)}{\mathtt{K}\vdash\mathtt{c}\langle\iota\rangle <: \mathtt{c}'\langle\iota\rangle}$$

Figure 8. Subtyping (reflexivity and transitivity rules are omitted.)

$$\text{(M-Sub)}\ \frac{\{\eta\leq\mathtt{mt},\mathtt{mt}\leq\eta',\eta\leq\eta'\}\in\mathtt{K}}{\mathtt{K}\models\{\eta\leq\mathtt{mt},\mathtt{mt}\leq\eta',\eta\leq\eta'\}}$$

Figure 9. Submoding

$$\begin{array}{lcll}
e & ::= & \ldots & \textit{runtime expressions}\\
 & | & \textbf{check}(e,\mathtt{m},\mathtt{m}') & \\
 & | & \mathtt{obj}(\alpha,\mathtt{c}\langle\iota\rangle,\overline{e}) & \\
 & | & \mathtt{cl}(\mathtt{m},e) & \\
\\
\mathbf{E} & ::= & \odot & \textit{evaluation context}\\
 & | & \mathbf{E}.\mathtt{md}(\overline{e}) & \\
 & | & o.\mathtt{md}(\ldots,o,\mathbf{E},e,\ldots) & \\
 & | & (T)\mathbf{E} \mid \mathbf{E}.\mathtt{fd} & \\
 & | & \textbf{snapshot}\ \mathbf{E}\,[\mathtt{m}_1,\mathtt{m}_2] & \\
 & | & \{\ldots\mathtt{m}:v;\mathtt{m}:\mathbf{E};\mathtt{m}:e\ldots\} \mid \mathbf{E}\rhd\mu & \\
 & | & \textbf{check}(\mathbf{E},\mathtt{m},\mathtt{m}') & \\
 & | & \mathtt{obj}(\alpha,\mathtt{c}\langle\mathbf{E},\iota\rangle,\overline{e}) & \\
 & | & \mathtt{obj}(\alpha,\mathtt{c}\langle\iota\rangle,\ldots v,\mathbf{E},e\ldots) & \\
 & | & \mathtt{cl}(\mathtt{m},\mathbf{E}) &
\end{array}$$

Figure 10. Run-Time Elements

$$\text{(T-Obj)}\ \frac{\begin{array}{c}\Gamma;\mathtt{K}\vdash e : \mathtt{mt}\\ \Gamma;\mathtt{K}\vdash\overline{e} : \overline{T} \qquad \mathtt{fields}(\mathtt{c}\langle\mathtt{mt},\iota\rangle) = (\overline{T}\ \overline{\mathtt{fd}} = \overline{e})\end{array}}{\Gamma;\mathtt{K}\vdash\mathtt{obj}(\alpha,\mathtt{c}\langle e,\iota\rangle,\overline{e}) : \mathtt{c}\langle\mathtt{mt},\iota\rangle}$$

$$\text{(T-Check)}\ \frac{\Gamma;\mathtt{K}\vdash e_1 : \mathtt{modev} \qquad \mathtt{mt}\ \mathit{fresh}}{\Gamma;\mathtt{K}\vdash\textbf{check}(e_1,\mathtt{m}_1,\mathtt{m}_2) : \exists\mathtt{m}_1\leq\mathtt{mt}\leq\mathtt{m}_2.\mathtt{mt}}$$

$$\text{(T-Closure)}\ \frac{\Gamma;\mathtt{K}\vdash e : \tau}{\Gamma;\mathtt{K}\vdash\mathtt{cl}(\mathtt{m},e) : \tau}$$

Figure 11. Auxiliary Run-time Expression Typing

(R-New)	**new** $\mathtt{c}\langle\iota\rangle$	$\stackrel{\mathtt{m}}{\Longrightarrow}$	$\mathtt{obj}(\alpha, \mathtt{c}\langle\iota\rangle, \mathtt{init}(P, \mathtt{c}))$	if α is $fresh$
(R-Cast)	$(\tau_0)\,o$	$\stackrel{\mathtt{m}}{\Longrightarrow}$	o	if $\tau <: \tau_0$
(R-Msg)	$o.\mathtt{md}(\overline{v}')$	$\stackrel{\mathtt{m}}{\Longrightarrow}$	$\mathtt{cl}(\mathtt{m}', e\{\overline{v}'/\overline{\mathtt{x}}\}\{o/\textbf{this}\})$	if $\mu \leq \mathtt{m}, \mathtt{m}' = \mathtt{emode}(o)$
(R-Field)	$o.\mathtt{fd}_i$	$\stackrel{\mathtt{m}}{\Longrightarrow}$	v_i	if $\mu \leq \mathtt{m}$
(R-Snapshot1)	**snapshot** o $[\mathtt{m}_1, \mathtt{m}_2]$	$\stackrel{\mathtt{m}}{\Longrightarrow}$	$\mathtt{obj}(\alpha', \mathtt{c}\langle\textbf{check}(e_a\{o/\textbf{this}\}, \mathtt{m}_1, \mathtt{m}_2), \iota\rangle, \overline{v})$	if $\mu = ?$, **class** $\mathtt{c}\ \cdots\ \{\ \cdots\ A\ \} \in P$, α' is $fresh$, $\mathtt{abody}(\mathtt{c}\langle ?, \iota\rangle) = e_a$
(R-Snapshot2)	**snapshot** o $[\mathtt{m}_1, \mathtt{m}_2]$	$\stackrel{\mathtt{m}}{\Longrightarrow}$	o	if $\mu = \mathtt{m}'$, **class** $\mathtt{c}\ \cdots\ \{\ \cdots\ A\ \} \in P, \mathtt{m}_1 \leq \mathtt{m}' \leq \mathtt{m}_2$
(R-Check)	**check**$(\mathtt{m}', \mathtt{m}_1, \mathtt{m}_2)$	$\stackrel{\mathtt{m}}{\Longrightarrow}$	$\mathtt{m}'$	if $\mathtt{m}_1 \leq \mathtt{m}' \leq \mathtt{m}_2$
(R-McaseProj)	$\{\overline{\mathtt{m} : v}\}^T \rhd \mathtt{m}_j$	$\stackrel{\mathtt{m}}{\Longrightarrow}$	v_j	
(R-Closure1)	$\mathtt{cl}(\mathtt{m}', e)$	$\stackrel{\mathtt{m}}{\Longrightarrow}$	$\mathtt{cl}(\mathtt{m}', e')$	if $e \stackrel{\mathtt{m}'}{\Longrightarrow} e'$
(R-Closure2)	$\mathtt{cl}(\mathtt{m}', v)$	$\stackrel{\mathtt{m}}{\Longrightarrow}$	v	
(R-Context)	$\mathbf{E}[\,e_1\,]$	$\stackrel{\mathtt{m}}{\Longrightarrow}$	$\mathbf{E}[\,e_2\,]$	if $e_1 \stackrel{\mathtt{m}}{\Longrightarrow} e_2$

for all rules: $o = \mathtt{obj}(\alpha,\, T, \overline{v},), \mathtt{mbody}(\mathtt{md},\, T) = \overline{\mathtt{x}}.e,\, T = \mathtt{c}\langle\mu, \iota\rangle$

Figure 12. Reduction Rules

$\mathtt{modes}(P)$	$\triangleq$	$\overline{\mathtt{m} \leq \mathtt{m}'}$	
$\mathtt{mode}(\mathtt{c}\langle\overline{\iota}\rangle)$	$\triangleq$	μ	if $\iota = \mu, \overline{\eta}$
$\mathtt{attr}(\mathtt{c}\langle\iota\rangle)$	$\triangleq$	$A\{\iota/\mathtt{eparam}(\Delta)\}$	if **class** $\mathtt{c}\ \Delta$ **extends** $\tau\ \{\ \overline{F}\ \overline{M}\ A\}\ \in\ P$
$\mathtt{eparam}(\overline{\eta \leq \mathtt{mt} \leq \eta'})$	$\triangleq$	$\overline{\mathtt{mt}}$	
$\mathtt{eparam}(? \rightarrow \omega, \Omega)$	$\triangleq$	$\mathtt{mt} \cup \mathtt{eparam}(\Omega)$	if $\omega = \eta \leq \mathtt{mt} \leq \eta'$
$\mathtt{ethis}(\Omega)$	$\triangleq$	$\mathtt{mt}$	if $\mathtt{eparam}(\Omega) = \mathtt{mt}$
$\mathtt{init}(P, \mathtt{c})$	$\triangleq$	$\mathtt{init}(\mathtt{c}') \cup \overline{e\{\iota/\mathtt{eparam}(\Delta)\}}$	if **class** $\Delta\ \mathtt{c}$ **extends** $\mathtt{c}'\ \overline{\tau\ \mathtt{fd} = e} \in P$
$\mathtt{init}(P, \mathtt{c})$	$\triangleq$	ϵ	if $\mathtt{c} = \mathtt{Object}$
$\mathtt{eargs}(\mathtt{c}\langle\iota\rangle)$	$\triangleq$	ι	
$\mathtt{eargs}(\exists\omega.\tau)$	$\triangleq$	$\mathtt{eargs}(\tau)$	
$\mathtt{cons}(\eta \leq \mathtt{mt} \leq \eta')$	$\triangleq$	$\bigcup\{\eta \leq \mathtt{mt}, \mathtt{mt} \leq \eta'\}$	
$\mathtt{cons}(? \rightarrow \omega, \Omega)$	$\triangleq$	$\{\eta \leq \mathtt{mt}, \mathtt{mt} \leq \eta'\} \cup \mathtt{cons}(\Omega)$	if $\omega = \eta \leq \mathtt{mt} \leq \eta'$

We require $\overline{\mathtt{m}}$ as a lattice. We use $\perp$ and $\top$ to represent the bottom and top of $\overline{\mathtt{m}}$ respectively.
We define $\mathtt{init}(P, \mathtt{c})$ as $\mathtt{init}(P, \mathtt{c}') \cup \overline{e}$ if **class** $\mathtt{c}$ **extends** $\mathtt{c}'\ \overline{\tau\ \mathtt{fd} = e} \in P$ or ϵ if $\mathtt{c} = \mathtt{Object}$.

Figure 13. Compile Functions

$\mathtt{emode}(\mathtt{m})$	$\triangleq$	$\mathtt{m}$
$\mathtt{emode}(\mathtt{obj}(\mathtt{c}\langle\iota\rangle, \overline{v},))$	$\triangleq$	$\mathtt{mode}(\mathtt{c}\langle\iota\rangle)$

Figure 14. Runtime Functions