# PSET5

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1

 $\mathbf{a}$ 

We can first begin calculating MRS for each individual. Without a loss of generality, we can begin to note the following:

$$MRS_{i} = \frac{\alpha(x_{1}^{i})^{\alpha - 1}(x_{2}^{i})^{\beta}}{\beta(x_{2}^{i})^{\beta - 1}(x_{1}^{i})^{\alpha}}$$

This above quantity implies that  $\frac{x_2^1}{x_1^1} = \frac{x_2^2}{x_1^2}$ . Let  $e_1 = x_1^1 + x_1^2$  and  $e_2 = x_2^1 + x_2^2$ . Using this equalities as well as the implication derived from the MRSes, we can find that we get:

$$x_2^1 = \frac{e_2}{e_1} x_1^1$$
  $x_2^2 = \frac{e_2}{e_1} x_1^2$ 

Since both the above quantities are linear in nature with no intercept, this implies that indeed the contact curve is that of connecting endpoints.

b

Since we we are working with different utility functions, we can find that after similar caluclations to above that:

$$MRS_1 = \frac{\alpha(x_1^1)^{\alpha - 1}(x_2^1)^{1 - \alpha}}{(1 - \alpha)(x_1^1)^{\alpha}(x_2^1)^{-\alpha}} = \frac{\alpha x_2^1}{(1 - \alpha)x_1^1}$$

and using similar calculations, we find that:

$$MRS_2 = \frac{\beta x_2^2}{(1 - \beta)x_1^2}$$

Since we know that  $1 > \alpha > \beta > 0$ , we find that:

$$\frac{\alpha}{1-\alpha} > \frac{\beta}{1-\beta}$$

Thus, we can see that for MRS to equal to each other, we know that:

$$\frac{x_2^1}{x_1^1} < \frac{x_2^2}{x_1^2}$$

Using the equations derived above, we can find that:

$$x_2^1 < \frac{e_2}{e_1} x_1^1$$

this implies that the graph still intersects the origins, but now  $x_2^1 < x_1^1$ , where we have all a curve that will be strictly below that of the original line derived in **a**. The contract curve is seen as below.

 $\mathbf{c}$ 

For the contract curve to exist, we want  $MRS_1 = MRS_2$ . Let  $e_1 = x_1^1 + x_1^2$ . We can see that

$$MRS_1 = MRS_2 \implies \alpha(x_1^1)^{\alpha - 1} = \beta(x_1^2)^{\beta - 1}$$

Thus, substuting the endowment, we find that

$$\alpha(x_1^1)^{\alpha-1} = \beta(e_1 - x_1^1)^{\beta-1}$$

So we see that as  $x_1^1 \to e_1^1$ , we find that the consumers will not consume any  $x_2$ , and consume only  $x_1$  However, since we know that  $\alpha > \beta$ , this implies that consumer  $x_1^1$  has greater value on  $x_1^1$ , which implies that  $x_1^1 > x_1^2$ . Note that any level of  $x_2$  sastifies the MRS equality argument, thus we are only concerned about when  $e_1 = x_1^1 + x_1^2$  and one of these inputs are the utility maximizing solution. The contract curve will like the one below:

 $\mathbf{2}$ 

 $\mathbf{a}$ 

See Graph

b

See Graph

 $\mathbf{c}$ 

See Graph

#### $\mathbf{d}$

The core would just be the point  $(e_1^1, e_2^1)$  as any trade would make this individual worse off, hence no trade would not be blocked by any coalition.

 $\mathbf{e}$ 

We first derive the general Marshallian Demand function for each individual. We are interested in the following optimization problem:

$$\max x_1 x_2$$
s.t  $p_1 e_1 + p_2 e_2 \ge p_1 x_1 + p_2 x_2$ 

Where the langrangian is:

$$L = x_1 x_2 - \lambda (p_1 e_1 + p_2 e_2 - p_1 x_1 - p_2 x_2)$$

where we see that the FOCs are:

$$[x_1]$$
  $x_2 + \lambda p_1 \le 0$  and  $x_1 \ge 0$   
 $[x_2]$   $x_1 + \lambda p_2 \le 0$  and  $x_2 \ge 0$   
 $[\lambda]$   $p_1e_1 + p_2e_2 \le p_1x_1 + p_2x_2$ 

We can see all FOCs must be strict equality, as if that is not the case, then markets will fail to clear  $([x_1], [x_2])$  and by the AU assumption that we want to use all of our endowment to maximize utility. Using these FOCs, we can find that:

$$p_1x_1 = p_2x_2$$

Thus, using this equation and the constraint, we find that:

$$x_1^m = \frac{e_1}{2} + \frac{p_2}{2p_1}e_2$$
  $x_2^m = \frac{p_1}{2p_2}e_1 + \frac{e_2}{2}$ 

Thus, we find that:

$$x_1^1 = \frac{p_2}{2p_1} + 1$$
  $x_2^1 = \frac{p_1}{p_2} + \frac{1}{2}$   $x_1^2 = \frac{3p_2}{2p_1} + 1$   $x_2^2 = \frac{p_1}{p_2} + \frac{3}{2}$ 

Thus we can find that where we let  $\mathbf{p} = (p_1, p_2)$ :

$$z_1(\mathbf{p}) = 2 + 2\left(\frac{p_2}{p_1}\right) - 4 = -2 + 2\left(\frac{p_2}{p_1}\right)$$

$$z_2(\mathbf{p}) = 2 + 2\left(\frac{p_1}{p_2}\right) - 4 = -2 + 2\left(\frac{p_1}{p_2}\right)$$

and we can verify that:

$$\mathbf{p} \cdot \mathbf{z} = \begin{bmatrix} p_1 & p_2 \end{bmatrix} \cdot \begin{bmatrix} 2 \begin{pmatrix} \frac{p_2}{p_1} \end{pmatrix} - 2 \\ 2 \begin{pmatrix} \frac{p_1}{p_2} \end{pmatrix} - 2 \end{bmatrix} = 2p_2 - 2p_1 + 2p_1 - 2p_2 = 0$$

 $\mathbf{f}$ 

We can see if  $p_2 = p_1$ , we find that obviously, **z** goes to 0. Thus, we find that the set of Walrasian equilibria is  $p^*(\mathcal{E}) = \{p_1, p_2\} | p_1 = p_2\}$  We find that a Walrasian equilibrium allocation is

$$x^W = \{(x_1, x_2)\} = \{(1.5, 1.5), (2.5, 2.5)\}$$

since there is one relative price, we can find that the set of Walrasian Equilbirum Allocations is:

$$W(\mathcal{E}) = \bigcup_{p^*} x^W(p^*(\mathcal{E}), \mathcal{E}) = \{(1.5, 1.5), (2.5, 2.5)\}$$

 $\mathbf{g}$ 

Note that  $e_1 = e_2$ , this implies that the slope of the line is 1, which means that all allocations that have  $x_1 = x_2$  will be in the core. Thus, we can see that the above set is a subset of the core.

#### h

Switching the utility function, we find that we are interested in the following optimization problem:

$$\max x_1^{\frac{2}{3}} x_2^{\frac{1}{3}}$$
s.t  $p_1 e_1 + p_2 e_2 \le p_1 x_1 + p_2 x_2$ 

the constraint remains the same, but with the following FOCs.

$$[x_1] \quad \left(\frac{2}{3}x_1^{-\frac{1}{3}}\right)x_2^{\frac{1}{3}} = p_1\lambda$$
$$[x_2] \quad \left(\frac{1}{3}x_1^{-\frac{1}{3}}\right)x_1^{\frac{2}{3}} = p_2\lambda$$

We know that these FOCs must have strict equality due to the same reasons as stated above. Using these FOCs, the following can be derived:

$$2x_2p_2 = x_1p_1$$

which implies that:

$$x_1^2 = \frac{2}{3p_1}(p_1e_1 + p_2e_2)$$
  $x_2^2 = \frac{1}{3p_1}(p_1e_1 + p_2e_2)$ 

Using, previous results, we can find that:

$$z_1(\mathbf{p}) = 1 + \frac{p_2}{2p_1} + \frac{4}{3} + \frac{2p_2}{p_1} - 4 = \frac{5p_2}{2p_1} - \frac{5}{3}$$

$$z_2(\mathbf{p}) = \frac{1}{2} + \frac{p_1}{p_2} + \frac{2p_1}{3p_2} + 1 - 4 = \frac{5p_1}{3p_2} - \frac{5}{2}$$

Note that this implies that  $3p_2 = 2p_1$ , as this is the only relative price that makes Walras' law hold. Thus, we see that:

$$x^{W} = \{(x_1, x_2)\} = \left\{ \left(\frac{4}{3}, 2\right), \left(\frac{8}{3}, 2\right) \right\}$$

and since we have only one relative price, we can find that the set of Walrasian equilbirum is

$$W(\mathcal{E}) = \bigcup_{p^*} x^W(p^*(\mathcal{E}), \mathcal{E}) = \left\{ \left(\frac{4}{3}, 2\right), \left(\frac{8}{3}, 2\right) \right\}$$

i

Consider the general form of the equation,

3

4

a

We can note that as we are all using the minimum function, we can see that  $x_1^1 = x_2^1 = x_2^2 = x_3^2 = x_1^3 = x_3^3$ , which implies that we are interested in the allocation:

$$x_1^1 = x_2^1 = x_2^2 = x_3^2 = x_1^3 = x_3^3 = 0.5$$

which is optimal as moving away would cause utilty lose and maximizes utilty for all and markets clear. Thus, this allocation is in the core, and let this allocation equal  $\overline{x}$ 

## $\mathbf{b}$

Consider the UMP for consumer 1. We are interested in the following optimization problem:

$$\max \min\{x_1^1, x_2^1\}$$
  
s.t  $x_1^1 p_1 + p_1 x_2^1 \le p e_1$ 

note that we want  $x_1^1 = x_2^1$ , which implies using the strict equality of the constraint (which we know is true by the AU assumption) that

$$x_1^1 = x_2^1 = \frac{p_1}{p_1 + p_2}$$

using a similar logic, we can see that:

$$x_2^2 = x_3^2 = \frac{p_2}{p_2 + p_3}$$

and

$$x_1^3 = x_3^3 = \frac{p_3}{p_1 + p_3}$$

Thus, we can now analyze the aggegate demand functions. We find that  $z_i = x_i^1 + x_i^2 + x_i^3 - 1$ 

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