

Homework 1

CIS-623 STRUCTURED PROGRAMMING & FORMAL METHODS

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Question 1:

Consider the following logic:

Syntax: A wff can be defined using the following BNF rules.

- $\Phi ::= p \mid (\Phi + \Phi) \mid (\neg \Phi)$
- p is propositional symbol representing a proposition from the set of propositions $\{p_1, p_2, p_3, \dots\}$

Semantics:

- A proposition can map to “true” or “false”
- “ \neg ” is interpreted as “negation” ($\neg \text{true} = \text{false}$; $\neg \text{false} = \text{true}$)
- “ $+$ ” is interpreted as “exclusive or”
 - $\text{true} + \text{true} = \text{false}$
 - $\text{false} + \text{false} = \text{false}$
 - $\text{true} + \text{false} = \text{true}$
 - $\text{false} + \text{true} = \text{true}$

Which strings are well-formed formulas (wff's) in this logic?

1. p_1 is a well-formed formula according to the BNF rules.
2. $(p_1 + (\neg p_2))$ is a wff.
3. $+p_1$ is not a wff because the exclusive or operator is defined as a binary operator in the BNF rules. The proposition on the left-hand side is missing.
4. $(\neg \Phi)$ is a wff.
5. $p_1 + p_2$ is not a wff because there is a missing set of parentheses.
6. $\neg(p_1)$ is not a wff. The negation symbol should be inside the parentheses.
7. $\neg(p_1$ is not a wff. There is a missing closed parenthesis.
8. $(p_1 + (p_2 + p_3))$ is a wff.

Question 2:

Let us assume that unary operator \neg has higher precedence than $+$ operator, and $+$ is a left associative operator. What are the fully parenthesized versions of the following wffs?

1. $((\neg p_1) + (\neg p_2))$
2. $((((\neg(\neg p_1)) + (\neg p_2)) + p_3))$
3. $((((\neg p_1) + p_2) + (\neg p_3))$

Question 3:

Axioms are the formulas that are regarded to be true in the logic. Can you give three axioms in this logic? (using schemas - generic versions of formulas)

The following three axioms are formulas that are valid (always true).

1. $((\neg p_1) + p_1)$

2. $((\neg p_1) + p_1) + (p_1 + p_1)$
3. $((((\neg p_1) + p_1) + (p_1 + p_1)) + (p_1 + p_1))$

Question 4:

Can you give some inference rules in this logic? (using schemas generic versions of formulas)

Rule 1: $(\Phi + \Psi), \Phi \vdash (\neg\Psi)$

Rule 2: $(\Phi + \Psi), \Psi \vdash (\neg\Phi)$

Rule 3: $(\Phi + \Psi), \neg\Phi \vdash \Psi$

Rule 4: $\Phi, \neg\Psi \vdash (\Phi + \Psi)$

Rule 5: $\neg\Phi, \Psi \vdash (\Phi + \Psi)$

Rule 6: $(\Phi + \Psi), \neg\Psi \vdash \Phi$

Rule 7: $\neg\neg\Phi \vdash \Phi$

Question 5:

Can you give a proof p_3 from p_1, p_1+p_2, p_2+p_3 using the inference rules from question 4?

$p_1, (p_1 + p_2), (p_2 + p_3) \vdash p_3$

1. p_1 (premise)
2. $(p_1 + p_2)$ (premise)
3. $(p_2 + p_3)$ (premise)
4. $(\neg p_2)$ [Rule 1], 1, 2
5. p_3 [Rule 3], 3, 4

Question 6:

Is $\neg p_2$ a logical consequence of p_1 and $p_1 + p_2$?

$p_1, (p_1 + p_2) \models \neg p_2$?

Model Number	p_1	p_2	$p_1 + p_2$
1	F	F	F
2	F	T	T
3	T	F	T
4	T	T	F

From the above truth table, only Model 3 satisfies the given premises that p_1 and $p_1 + p_2$ are both true. In this model, p_2 is false. Therefore, $\neg p_2$ is true.

$\neg p_2$ is a logical consequence of p_1 and $p_1 + p_2$.

$p_1, (p_1 + p_2) \models \neg p_2$