

# CSE515 Advanced Algorithms

## Lecture 18

### Local Search and Scheduling Jobs on Identical Parallel Machines

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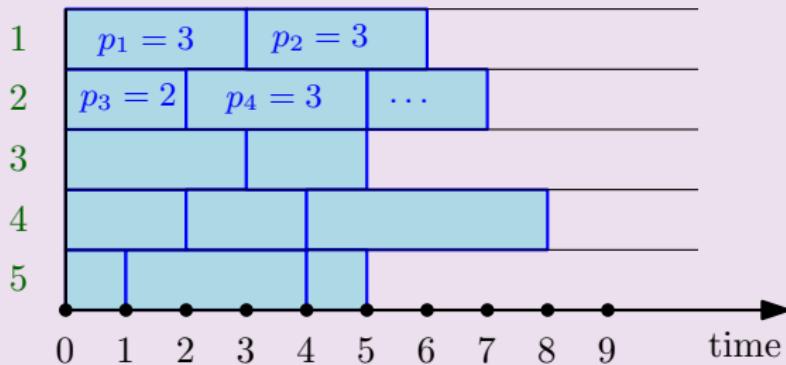
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# Introduction

- Assignment 3 is due tomorrow.
- Reference: Section 2.3 in [The design of approximation algorithms](#) by David P. Williamson and David B. Shmoys.

# Problem Statement

Machine



Here  $m = 5$ .

Here  $C_3 = 5$ ,  
 $C_{\max} = 8$ .

INPUT:

- Number of machines  $m$ .
- Processing times  $p_1, \dots, p_n$ .

OUTPUT:

- A schedule that minimizes the *makespan*, which is the maximum completion time  $C_{\max} = \max_{j=1, \dots, n} C_j$ .

# Local Search

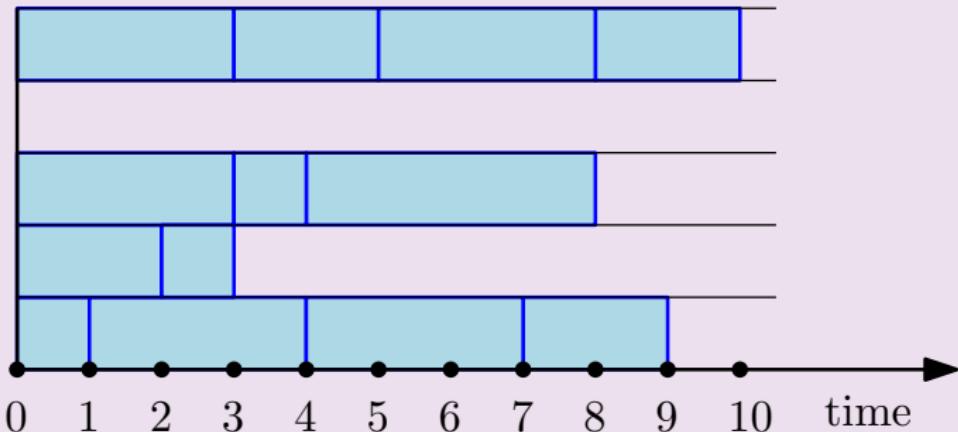
Simple local search:

- Start with a *feasible* solution.
- Apply a local change that improves the solution.
- Repeat until no local change gives an improvement.

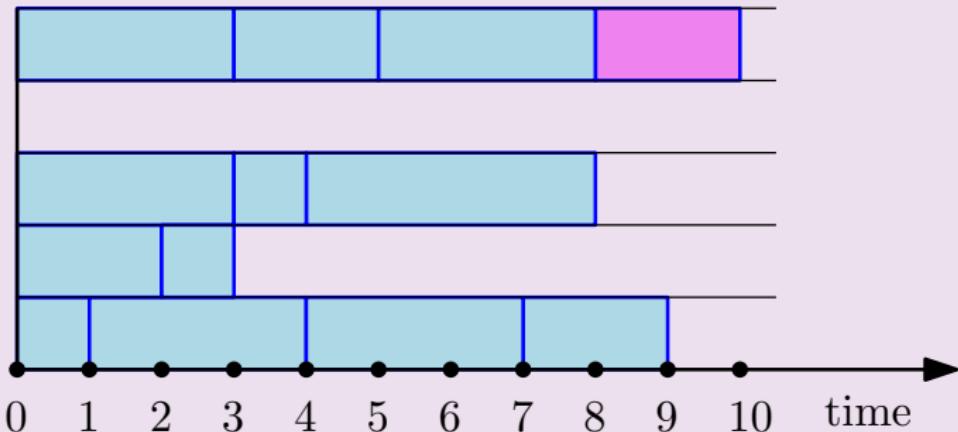
Difference with greedy algorithms:

- In a greedy algorithm, the initial solution may not be feasible.

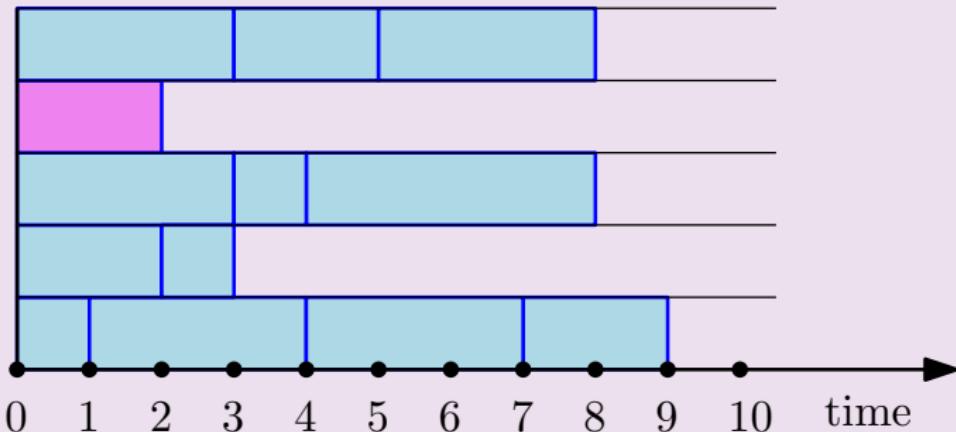
## Example of Local Search



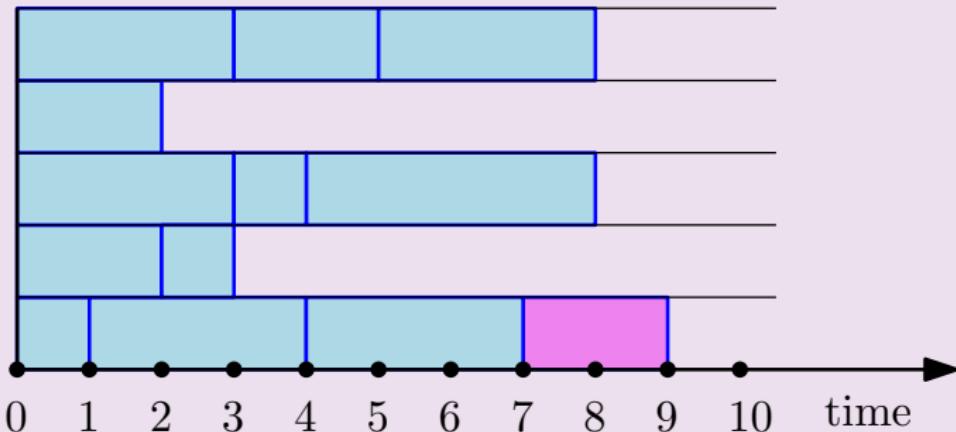
## Example of Local Search



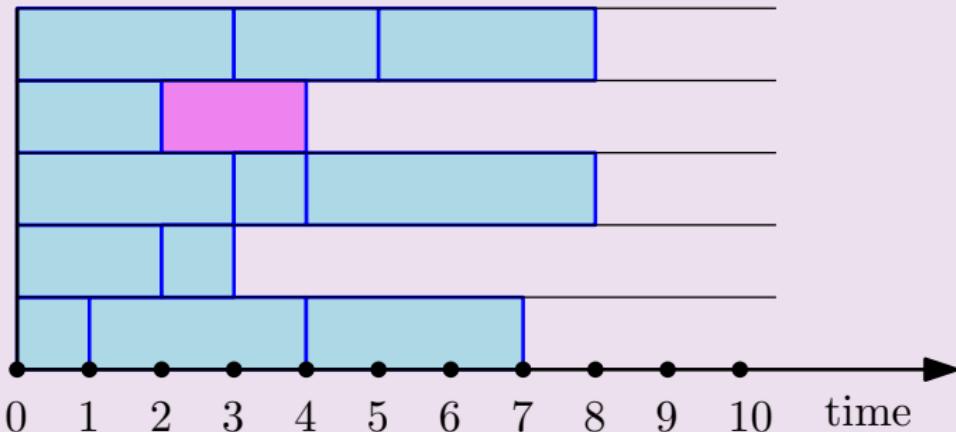
## Example of Local Search



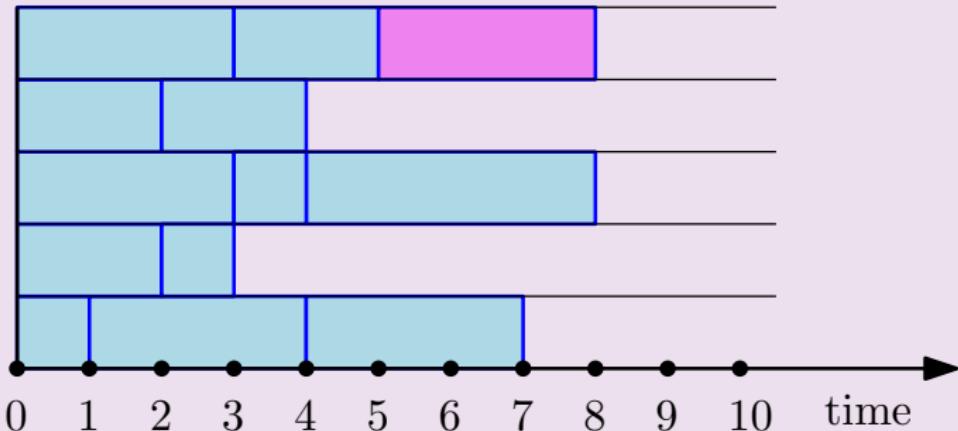
## Example of Local Search



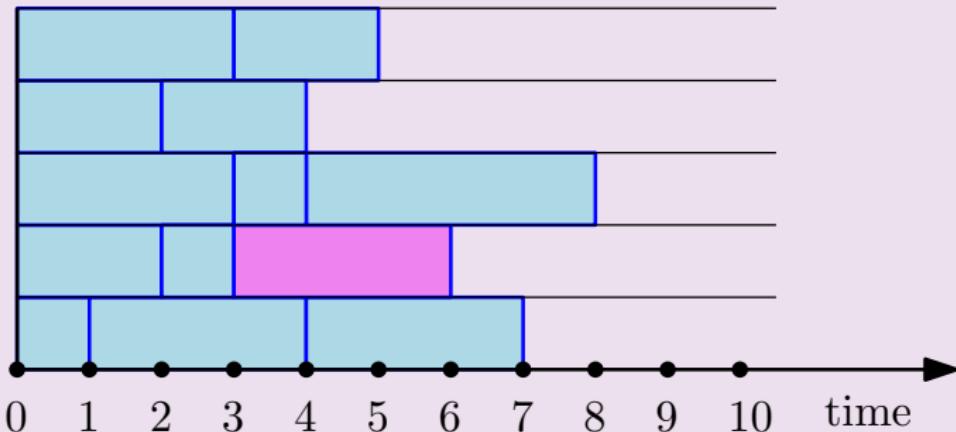
## Example of Local Search



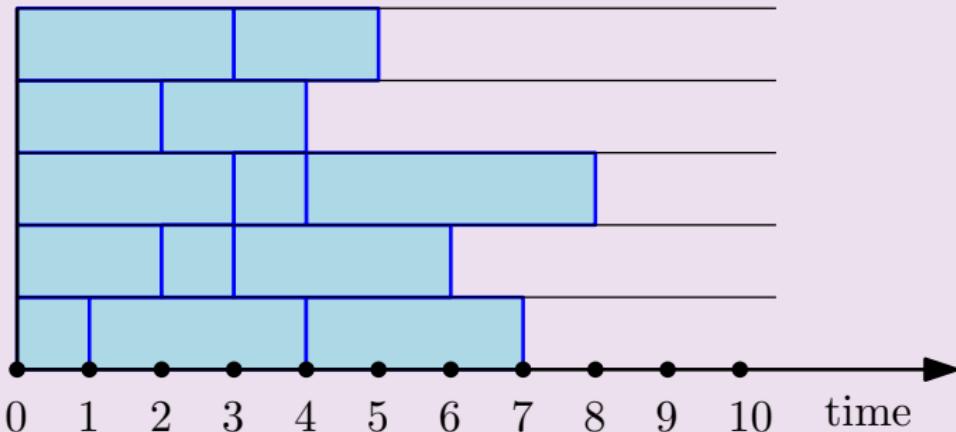
## Example of Local Search



## Example of Local Search



## Example of Local Search



# Pseudocode

Let  $C_{\min}$  be the completion time of a machine that completes all its jobs earliest.

## Local search algorithm for scheduling on parallel identical machines

- 1: start with an arbitrary feasible solution.
- 2: **while** there is a job  $j$  with  $C_j = C_{\max} > C_{\min} + p_j$  **do**
- 3:     transfer job  $j$  to a machine that completes at time  $C_{\min}$ .
- 4: **return** current solution

## Theorem

*This algorithm is a 2-factor approximation algorithm.*

# Proof

- We need to prove:
  - ▶ The local search algorithms runs in polynomial time.
  - ▶ Its approximation factor is 2.
- *Done in class. See textbook.*

# List Scheduling Algorithm

## List scheduling algorithm

- ① Build a list of all the jobs.
- ② Whenever a machine becomes idle, process the next job on the list with this machine.

## Theorem

*The list scheduling algorithm is a 2-approximation algorithm for scheduling jobs on identical parallel machines.*

## Proof.

The local search algorithm cannot improve the solution returned by this algorithm. □

# Longest Processing Time Rule

The *Longest Processing Time rule* (LPT) is to sort the jobs in the list by decreasing processing time.

## Theorem

*The list scheduling algorithm with LPT is a  $4/3$ -approximation algorithm.*

Proof: Partly done in class. See textbook and exercise set.