

# VC-DENSITY IN AN ADDITIVE REDUCT OF P-ADIC NUMBERS

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ABSTRACT. [1] computed a bound  $2n + 1$  for the VC function in p-adic numbers, but it is not known to be optimal. I investigate a C-minimal additive reduct of p-adic numbers and using techniques of [2] I compute an optimal bound  $n$  for that structure.

VC density was introduced in [1] by Aschenbrenner, Dolich, Haskell, MacPherson, and Starchenko as a natural notion of dimension for NIP theories. In a NIP theory we can define the VC function

$$\text{vc} : \mathbb{N} \longrightarrow \mathbb{N}$$

Where  $\text{vc}(n)$  measures complexity of definable sets in an  $n$ -dimensional space. Simplest possible behavior is  $\text{vc}(n) = n$  for all  $n$ . [1] computes an upper bound for this function to be  $2n + 1$ , and it's not known whether it's optimal. This same bound would hold in any reduct of p-adic numbers, so one may hope that the simplified structure of the reduct would allow a better bound. In [2], Leenknegt provides a cell decomposition result for the C-minimal additive reduct of p-adic numbers. Using that I'm able to improve the bound for the VC function, showing that  $\text{vc}(n) = n$ .

## 1. CELL DECOMPOSITION

We work with the reduct of p-adic numbers in the language  $L = \{\mathbb{Q}_p, Q_{n,m}, +, -, \{\bar{c}\}_{c \in K}\}$ , where  $\bar{c}$  is a scalar multiplication by  $c$ , and  $Q_{n,m}$  is a unary predicate

$$Q_{n,m} = \left\{ \bigcup_{k \in \mathbb{Z}} p^{kn}(1 + p^m \mathbb{Z}_p) \right\}$$

[2] provides a cell decomposition result for this structure. Any formula  $\phi(t, x)$  with  $t$  singleton decomposes as the union of the following cells:

$$\{(x, t) \in D \times K \mid \text{val } a_1(x) \square_1 \text{val}(t - c(x)) \square_2 \text{val } a_2(x), t - c(x) \in \lambda Q_{n,m}\}$$

where  $D$  is a cell of a smaller dimension,  $a_1, a_2, c$  are linear polynomials in  $x$ ,  $\square$  is  $<$  or no condition,  $\lambda \in \mathbb{Q}_p$ .

We analyze a formula  $\phi(x, y)$  to find an upper bound of its VC-density. Using cell decomposition, without loss of generality we may assume that we only need to bound the following family of formulas  $\Psi(x, y)$

$$\begin{aligned} \text{val } p_i(x) - c_i(y) &< \text{val } p_j(x) - c_j(y) & i, j \in I \\ \text{val } p_i(x) - c_i(y) &\in \lambda_k Q & i \in I, k \in K \end{aligned}$$

where  $I, K$  some finite index sets,  $p_i$  is linear in  $x$ ,  $c_i$  is a linear polynomial in  $y$ ,  $\lambda_k \in \mathbb{Q}_p$ , and  $Q = Q_{n,m}$  for some  $n', m'$ .

To see that apply cell decomposition theorem to  $\phi(x_1, \bar{x}; y)$ . Extract from the cells all the polynomials  $a_1(\bar{x}, y), a_2(\bar{x}, y), x_1 - c(\bar{x}, y)$ , and separate  $x$  and  $y$  parts into  $p_i(x) - c_i(y)$ . Choose  $n', m'$  large enough to cover all  $n, m$  that come up in the cells. Finally choose  $\lambda_k$  to go over all cosets of  $Q$ .

Then  $(x, y), (x', y')$  agreeing on  $\Psi$ , will agree on being contained in those cells, and thus will agree on satisfying  $\phi$ .

## 2. KEY LEMMAS AND DEFINITIONS

**Definition 2.1.** A tuple  $p \in \mathbb{Q}_p^m$  can be viewed as a vector  $\vec{p}$ , treating  $\mathbb{Q}_p^m$  as a vector space over  $\mathbb{Q}_p$ .

**Lemma 2.2.** Suppose we have a collection of vectors  $\{\vec{p}_i\}_{i \in I}$  with each  $\vec{p}_i \in \mathbb{Q}_p^m$ . Pick a subset  $J \subset I$  and  $j \in I$  such that

$$\vec{p}_j \in \text{span} \{\vec{p}_i\}_{i \in J}$$

Suppose we have  $\vec{x} \in \mathbb{Q}_p^m, \alpha \in \mathbb{Z}$  with

$$\text{val}(\vec{p}_i \cdot \vec{x}) > \alpha \text{ for all } i \in J$$

Then

$$\text{val}(\vec{p}_j \cdot \vec{x}) > \alpha - \gamma$$

for some  $\gamma \in \mathbb{Z}^{\geq 0}$ . Moreover  $\gamma$  can be chosen independent of choice of  $J, j, \vec{x}, \alpha$  depending only on  $\{\vec{p}_i\}_{i \in I}$ .

The following lemma is an adaptation of lemma 7.4 in [1].

**Lemma 2.3.** For  $n, m$  there exists  $D = D(n, m) \in \mathbb{Z}$  such that for any  $x, y, a \in \mathbb{Q}_p$  if

$$\text{val } x - a = \text{val } y - a < \text{val}(x - y) - D$$

then  $x - a, y - a$  are in the same coset of  $Q_{n,m}$ .

**Definition 2.4.** For  $c \in \mathbb{Q}_p, \alpha \in \mathbb{Z}$  we define an open ball

$$B(c, \alpha) = \{c' \in \mathbb{Q}_p \mid \text{val}(c' - c) \leq \alpha\}$$

Suppose we have a finite  $T \subset \mathbb{Q}_p$ . We view it as a tree as follows. Branches through the tree are elements of  $T$ . With this tree we associate open balls  $B(t_1, \text{val}(t_1 - t_2))$  for all  $t_1, t_2 \in T$ . An interval is two balls  $B(t_1, v_1) \supset B(t_2, v_2)$  with no balls in between. An element  $a \in \mathbb{Q}_p$  belongs to this interval if  $a \in B(t_1, v_1) \setminus B(t_2, v_2)$ . There are at most  $2|T|$  different intervals and they partition the entire space.

We may rewrite our collection of formulas  $\Psi(x, y)$  as

$$\begin{aligned} \text{val } \vec{p}_i \cdot \vec{x} - c_i(y) &< \text{val } \vec{p}_j \cdot \vec{x} - c_j(y) & i, j \in I \\ \text{val } \vec{p}_i \cdot \vec{x} - c_i(y) &\in \lambda_k Q & i \in I, k \in K \end{aligned}$$

Fix a parameter set  $B$  of size  $N$ .

Consider a tree  $T = \{c_i(b) \mid b \in B, i \in I\}$  It has at most  $O(N) = N \cdot |I|$  many intervals. For the remainder of the paper we work with this tree.

For some  $c, c' \in \mathbb{Q}_p^m$  we say they have the same  $\Psi$ -type if they have the same  $\Psi$  type over  $B$ .

For some  $x, x' \in \mathbb{Q}_p$  we say they have the same tree type if

- $x + c_i(b)$  is in the same  $Q$ -coset as  $x' + c_i(b)$  for all  $i \in I, b \in B$
- $\text{val}(x + c_i(b)) < \text{val}(x + c_j(b))$  iff  $\text{val}(x' + c_i(b)) < \text{val}(x' + c_j(b))$  for all  $i, j \in I, b \in B$

**Lemma 2.5.** Let  $c, c' \in \mathbb{Q}_p^m$ . If  $p_i(c), p_i(c')$  have the same tree type for all  $i \in I$ , then  $c, c'$  have the same  $\Psi$ -type.

**Lemma 2.6.** For any  $Q = Q_{n,m}$  there exists  $\theta_Q$  such that for all  $\theta \geq \theta_Q$  the following holds. Suppose we have  $x, y, c \in \mathbb{Q}_p$  such that

$$\text{val}(x - y) - \theta > \text{val}(x - c) = \text{val}(y - c)$$

Then  $x - c$  and  $y - c$  lie in the same coset of  $Q$ .

Next lemma is along the lines of lemma 7.5 of [1].

**Lemma 2.7.** Fix  $\theta$  sufficiently large to satisfy previous lemma for all  $Q_i$ . Define an enumeration of near balls

$$B_1(c, \alpha), B_2(c, \alpha), \dots, B_{N_\theta}(c, \alpha)$$

**Definition 2.8.** Let  $c \in \mathbb{Q}_p$ . It lies in our tree in one of the intervals  $B(c_L, \alpha_L) \setminus B(c_U, \alpha_U)$ . Suppose  $c$  lies in one of the near balls corresponding to  $B(c_L, \alpha_L)$  or  $B(c_U, \alpha_U)$ . Then define its interval type to be the index of that near ball. Otherwise define its interval type to be the coset of  $c - c_U$  of  $Q$ . Denote the space of all the possible branch types Bt. We have

$$|\text{Bt}| = N_\theta + \text{number of cosets of } Q$$

depending only on  $\Psi$ , independent from  $B$ .

**Lemma 2.9.** *If  $c, c'$  are in the same interval and have the same interval type then they have the same tree type.*

**Definition 2.10.** For  $c \in \mathbb{Q}_p$  and  $\alpha, \beta \in \mathbb{Z}$  let  $c \upharpoonright [\alpha, \beta] \in \mathbb{Z}_p^{\beta-\alpha}$  be the record of coefficients of  $c$  for valuations between  $\alpha, \beta$ . More precisely write  $c$  in its power series form

$$c = \sum_{\gamma \in \mathbb{Z}} c_\gamma p^\gamma \text{ with } c_\gamma \in \mathbb{Z}/p\mathbb{Z}$$

Then  $c \upharpoonright [\alpha, \beta]$  is just  $(c_\alpha, c_{\alpha+1}, \dots, c_\beta)$ .

**Definition 2.11.** Let  $c \in \mathbb{Q}_p$ . It lies in our tree in one of the intervals  $B(c_L, \alpha_L) \setminus B(c_U, \alpha_U)$ . Define  $F(c)$ , the floor of  $c$  to be  $\alpha_L$ .

Let  $f : \mathbb{Q}_p^n \rightarrow \mathbb{Q}_p^I$  with  $f(\vec{c}) = (p_i(\vec{c}))_{i \in I}$ . Define segment space Sg to be the image of  $f$ .

For some element  $(a_i)$  in segment space look at floors  $F(a_i)$ . Partition the segment space by order type of  $\{F(a_i)\}$ . Work in a fixed partition Sg'. After relabeling we may assume that

$$F(a_1) \geq F(a_2) \geq \dots$$

Consider (relabelled) sequence of vectors  $\vec{p}_1, \vec{p}_2, \dots, \vec{p}_I$ . Choose the unique subset of linearly independent vectors  $J \subset I$ . For any index  $i \in I$  we call it independent if  $i \in J$  and we call it dependent otherwise.

For all  $a_i$  record its interval type.

For  $a_i$  with  $i$  independent, record the interval of  $a_i$ .

Pick  $a_i$  with  $i$  dependent. Let  $j$  be the largest independent index with  $j < i$ . Record  $a_i \upharpoonright [F(a_j) - \gamma, F(a_j)]$ .

Combining all the records defines a function

$$g : \text{Sg}' \rightarrow \text{Bt}^I \times \text{Pt}^m \times \text{Ct}^I$$

We claim that for  $\vec{a}, \vec{a}' \in \text{Sg}'$  if we have  $g(\vec{a}) = g(\vec{a}')$  then all  $a_i, a'_i$  have the same tree type.

*Proof.* Suppose we have  $\vec{a}, \vec{a}' \in \text{Sg}'$  that map to the same image by  $g$ . Suppose  $i$  is independent. Then by construction,  $a_i, a'_i$  map to the same interval of the tree and have the same interval type. Thus they have the same tree type. Otherwise, suppose  $i$  is dependent. Pick largest  $j < i$  such that  $j$  is independent. We have  $F(a_i) \leq F(a_j)$  and  $F(a'_i) \leq F(a'_j)$ . Moreover  $F(a_j) = F(a'_j)$  as they are mapped to the same interval (as  $j$  is independent).

**Claim 2.12.**  $\text{val}(a_i - a'_i) > F(a_j) - \gamma$

*Proof.* Let  $\bar{x}, \bar{x}' \in \mathbb{Q}_p^m$  be some elements with

$$\vec{p}_k \cdot \bar{x} = a_k$$

$$\vec{p}_k \cdot \bar{x}' = a'_k \text{ for all } k \in I$$

Let  $J$  be the set of independent indices less than  $i$ . We have

$$\text{val}(a_k - a'_k) > F(a_k) \text{ for all } k \leq J$$

as for independent indices  $a_k, a'_k$  lie in the same interval.

$$\text{val}(a_k - a'_k) > F(a_j) \text{ for all } k \leq J \text{ by monotonicity of } F(a_k)$$

$$\text{val}(\vec{p}_k \cdot \bar{x} - \vec{p}_k \cdot \bar{x}') > F(a_j) \text{ for all } k \leq J$$

$$\text{val}(\vec{p}_k \cdot (\bar{x} - \bar{x}')) > F(a_j) \text{ for all } k \leq J$$

$J$  and  $i$  match the requirements of the claim above by independence so we conclude

$$\text{val}(\vec{p}_i \cdot (\bar{x} - \bar{x}')) > F(a_j) - \gamma$$

$$\text{val}(\vec{p}_i \cdot \bar{x} - \vec{p}_i \cdot \bar{x}') > F(a_j) - \gamma$$

$$\text{val}(a_i - a'_i) > F(a_j) - \gamma$$

as needed.  $\square$

By record of continuations (which  $a_i, a'_i$  agree on) we have

$$a_i = a'_i \upharpoonright F(a_j)$$

As  $F(a_i) \leq F(a_j)$ ,  $a_i, a'_i$  have to lie in the same interval. They also agree on interval type. Thus they have the same tree type.  $\square$

Now suppose we have  $c, c' \in \mathbb{Q}_p^m$  such that  $g(f(c)) = g(f(c'))$ . Then  $f(c)$  components have the same tree type as  $f(c')$  components. Then  $c, c'$  have the same  $\Psi$ -type. Thus the number of possible  $\Psi$ -types is bound by the size of the range of  $g$ .

$$|\text{Ct}| = p^\gamma$$

$$|\text{Pt}| \leq N \cdot I^2 \text{ (the only component dependent on } N)$$

Moreover we need at most  $I!$  many partitions of Sg. This gives us

$$I! \cdot |\text{Bt}|^I \cdot (N \cdot I^2)^m \cdot p^{\gamma I} = O(N^m)$$

upper bound for the possible number of  $\Psi$ -types.

## REFERENCES

- [1] M. Aschenbrenner, A. Dolich, D. Haskell, D. Macpherson, S. Starchenko, *Vapnik-Chervonenkis density in some theories without the independence property*, I, preprint (2011)
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