

# VC-DENSITY IN AN ADDITIVE REDUCT OF $p$ -ADIC NUMBERS

ANTON BOBKOV

ABSTRACT. Aschenbrenner et. al. computed a bound  $\text{vc}(n) = 2n - 1$  for the VC density function in the field of  $p$ -adic numbers, but it is not known to be optimal. I investigate a certain  $P$ -minimal additive reduct of the field of  $p$ -adic numbers and using a cell decomposition result of Leenknegt I compute an optimal bound  $\text{vc}(n) = n$  for that structure.

VC density was introduced into model theory in [1] by Aschenbrenner, Dolich, Haskell, MacPherson, and Starchenko as a natural notion of dimension for NIP theories. In a NIP theory we can define the VC function

$$\text{vc} : \mathbb{N} \longrightarrow \mathbb{N}$$

where  $\text{vc}(n)$  measures complexity of the definable sets in an  $n$ -dimensional space. The simplest possible behavior is  $\text{vc}(n) = n$  for all  $n$ . [1] computes an upper bound for this function to be  $2n + 1$ , and it is not known whether it is optimal. This same bound would hold in any reduct of  $p$ -adic numbers, so one may hope that the simplified structure of the reduct would allow a better bound. In [2], Leenknegt provides a cell decomposition result for a certain  $P$ -minimal additive reduct of  $p$ -adic numbers. Using that I'm able to improve the bound for the VC function, showing that  $\text{vc}(n) = n$ .

## 1. CELL DECOMPOSITION

**Definition 1.1.** Let

$$Q_{n,m} = \bigcup_{k \in \mathbb{Z}} p^{kn} (1 + p^m \mathbb{Z}_p)$$

It is a subgroup of the multiplicative group of  $\mathbb{Q}_p$  with finitely many cosets.

We work with the reduct of  $p$ -adic numbers in the language  $\mathcal{L}_{aff} = \{\mathbb{Q}_p, \{R_{n,m}\}_{n,m \in \mathbb{N}}, +, -, \{\bar{c}\}_{c \in \mathbb{Q}_p}\}$ , where  $\bar{c}$  is a scalar multiplication by  $c$ , and  $R_{n,m}$  is a predicate for cosets of  $Q_{n,m}$

$$Q_{n,m} = \bigcup_{k \in \mathbb{Z}} p^{kn} (1 + p^m \mathbb{Z}_p)$$

In [2], Leenknegt provides a cell decomposition result for this structure. Any formula  $\phi(t, x)$  with  $t$  singleton decomposes as the union of the following cells:

$$\{(t, x) \in K \times D \mid \text{val } a_1(x) \square_1 \text{val}(t - c(x)) \square_2 \text{val } a_2(x), t - c(x) \in \lambda Q_{n',m'}\}$$

where  $D$  is a cell of a smaller dimension,  $a_1, a_2, c$  are linear polynomials in  $x$ ,  $\square$  is  $<$  or no condition,  $\lambda \in \mathbb{Q}_p$ .

**Lemma 1.2.** For a formula  $\phi(x)$  with  $x = (t, \bar{x})$  there exists a family of formulas  $\Psi'(x)$

$$\begin{aligned} \text{val}(q_i(x)) &< \text{val}(q_j(x)) & i, j \in I \\ \text{val}(q_i(x)) &\in \lambda_k Q_{n,m} & i \in I, k \in K \\ \bar{x} &\in D_l & l \in L \end{aligned}$$

with  $I, K, L$  finite,  $D_l$  cells,  $q_i$  linear polynomials,  $\lambda_k \in \mathbb{Q}_p$ , and  $Q = Q_{n,m}$  for some  $n, m$ . Moreover we have that if  $a, a' \in Q_p^{[x]}$  agree on all the formulas from  $\Psi'$  then they agree on  $\phi$ .

*Proof.* To see that, apply cell decomposition theorem to  $\phi(t, \bar{x})$ . Let  $q_i$  enumerate all of the polynomials  $a_1(\bar{x}), a_2(\bar{x}), t - c(\bar{x})$  that show up in the cells. Let  $D_l$  be the smaller cells for the  $\bar{x}$  components that appear in the cells. Choose  $n, m$  large enough to cover all  $n', m'$  that come up in the cells for  $Q_{n',m'}$ . Choose  $\lambda_k$  to go over all the cosets of  $Q_{n,m}$ .  $\square$

Applying this lemma inductively to smaller cells, we obtain a family  $\Psi(x)$

$$\begin{aligned} \text{val}(q_i(x)) &< \text{val}(q_j(x)) & i, j \in I \\ \text{val}(q_i(x)) &\in \lambda_k Q_{n,m} & i \in I, k \in K \end{aligned}$$

with  $I, K$  finite,  $q_i$  linear polynomials,  $\lambda_k \in \mathbb{Q}_p$ , and  $Q = Q_{n,m}$  for some  $n, m$ . Moreover whenever  $a, a' \in Q_p^{[x]}$  agree on all the formulas from  $\Psi$  then they agree on  $\phi$ .

Now fix a formula  $\phi(x; y)$  for finding an upper bound of its VC-density. Using the result above we can construct a family of formulas  $\Psi(x; y)$  which can be now written as

$$\begin{aligned} \text{val}(p_i(x) - c_i(y)) &< \text{val}(p_j(x) - c_j(y)) & i, j \in I \\ \text{val}(p_i(x) - c_i(y)) &\in \lambda_k Q & i \in I, k \in K \end{aligned}$$

where  $I, K$  finite,  $p_i$  a homogeneous linear polynomials in  $x$ ,  $c_i$  is a linear polynomial in  $y$ ,  $\lambda_k \in \mathbb{Q}_p$ , and  $Q = Q_{n,m}$  for some  $n, m$  (to do this we simply split the polynomial  $q_i$  into its  $x$  part and into its  $y$  part including the constant term). Now for any parameter set  $B$  we have that if  $a, a'$  have the same  $\Psi$ -type over  $B$  then they have the same  $\phi$ -type over  $B$ . Thus it suffices to bound VC-density for  $\Psi$ .

## 2. KEY LEMMAS AND DEFINITIONS

**Definition 2.1.** A tuple  $p \in \mathbb{Q}_p^{[x]}$  can be viewed as a vector  $\vec{p}$ , treating  $\mathbb{Q}_p^{[x]}$  as a vector space over  $\mathbb{Q}_p$ .

We may rewrite our collection of formulas  $\Psi(x, y)$  as

$$\begin{aligned} \text{val}(\vec{p}_i \cdot \vec{x}) - c_i(y) &< \text{val}(\vec{p}_j \cdot \vec{x}) - c_j(y) & i, j \in I \\ \text{val}(\vec{p}_i \cdot \vec{x}) - c_i(y) &\in \lambda_k Q & i \in I, k \in K \end{aligned}$$

**Lemma 2.2.** Suppose we have a collection of vectors  $\{\vec{p}_i\}_{i \in I}$  with each  $\vec{p}_i \in \mathbb{Q}_p^{[x]}$ . Pick a subset  $J \subset I$  and  $j \in I$  such that

$$\vec{p}_j \in \text{span}\{\vec{p}_i\}_{i \in J}$$

Suppose we have  $\vec{x} \in \mathbb{Q}_p^{[x]}, \alpha \in \mathbb{Z}$  with

$$\text{val}(\vec{p}_i \cdot \vec{x}) > \alpha \text{ for all } i \in J$$

Then

$$\text{val}(\vec{p}_j \cdot \vec{x}) > \alpha - \gamma$$

for some  $\gamma \in \mathbb{Z}^{\geq 0}$ . Moreover  $\gamma$  can be chosen independently from  $J, j, \vec{x}, \alpha$  depending only on  $\{\vec{p}_i\}_{i \in I}$ , independent of their order.

*Proof.* Fix some  $i, J$ . For some  $c_i$

$$\begin{aligned} \vec{p}_j &= \sum_{i \in J} c_i \vec{p}_i \\ \vec{p}_j \cdot \vec{x} &= \sum_{i \in J} c_i \vec{p}_i \cdot \vec{x} \end{aligned}$$

We have

$$\text{val}(c_i \vec{p}_i \cdot \vec{x}) = \text{val}(c_i) + \text{val}(\vec{p}_i \cdot \vec{x}) > \text{val}(c_i) + \alpha$$

Pick  $\gamma = -\max \text{val}(c_i)$  or 0 if all those values are positive. Then we have

$$\begin{aligned} \text{val}(c_i \vec{p}_i \cdot \vec{x}) &> \alpha - \gamma & \text{for all } i \in J \\ \sum_{i \in J} c_i \vec{p}_i \cdot \vec{x} &> \alpha - \gamma \end{aligned}$$

This shows that we can pick such  $\gamma$  for a given choice of  $i, J$ , but independent from  $\alpha, \vec{x}$ . To get a choice independent from  $i, J$ , go over all such eligible choices (of which there are finitely many as  $I$  is finite), pick  $\gamma$  for each, and then take the maximum of those values.  $\square$

**Definition 2.3.** For  $c \in \mathbb{Q}_p, \alpha \in \mathbb{Z}$  we define an open ball

$$B(c, \alpha) = \{c' \in \mathbb{Q}_p \mid \text{val}(c' - c) \leq \alpha\}$$

**Definition 2.4.** Suppose we have a finite  $T \subset \mathbb{Q}_p$ . We view it as a tree as follows. Branches through the tree are elements of  $T$ . With this tree we associate open balls  $B(t_1, \text{val}(t_1 - t_2))$  for all  $t_1, t_2 \in T$ . An interval is two balls  $B(t_1, v_1) \supset B(t_2, v_2)$  with no balls in between. An element  $a \in \mathbb{Q}_p$  belongs to this interval if  $a \in B(t_1, v_1) \setminus B(t_2, v_2)$ . There are at most  $2|T|$  different intervals and they partition the entire space.

Fix a parameter set  $B$  of size  $N$ .

Consider a tree  $T = \{c_i(b) \mid b \in B, i \in I\}$  It has at most  $O(N) = N \cdot |I|$  many intervals. Denote the set of all intervals as  $\text{Pt}$ . For the remainder of the paper we work with this tree.

**Definition 2.5.** Let  $c \in \mathbb{Q}_p$ . It lies in the tree in one of the unique intervals  $B(c_L, \alpha_L) \setminus B(c_U, \alpha_U)$ . Define  $F(c)$ , the floor of  $c$  to be  $\alpha_L$ .

**Definition 2.6.** We say  $x, x' \in \mathbb{Q}_p$  have the same tree type if

- $\text{val}(x - c_i(b)) < \text{val}(x - c_j(b))$  iff  $\text{val}(x' - c_i(b)) < \text{val}(x' - c_j(b))$  for all  $i, j \in I, b \in B$
- $x + c_i(b)$  is in the same  $Q$ -coset as  $x' + c_i(b)$  for all  $i \in I, b \in B$

**Lemma 2.7.** Let  $a, a' \in \mathbb{Q}_p^{|x|}$ . If  $p_i(a), p_i(a')$  have the same tree type for all  $i \in I$ , then  $a, a'$  have the same  $\Psi$ -type.

*Proof.* Clear from the construction.  $\square$

**Definition 2.8.** For  $c \in \mathbb{Q}_p$  and  $\alpha, \beta \in \mathbb{Z}$  let  $c \upharpoonright [\alpha, \beta] \in (\mathbb{Z}/p\mathbb{Z})^{\beta-\alpha}$  be the record of coefficients of  $c$  for the valuations between  $\alpha, \beta$ . More precisely write  $c$  in its power series form

$$c = \sum_{\gamma \in \mathbb{Z}} c_\gamma p^\gamma \text{ with } c_\gamma \in \mathbb{Z}/p\mathbb{Z}$$

Then  $c \upharpoonright [\alpha, \beta]$  is just  $(c_\alpha, c_{\alpha+1}, \dots, c_\beta)$ .

The following lemma is an adaptation of lemma 7.4 in [1].

**Lemma 2.9.** For  $n, m$  there exists  $D = D(n, m) \in \mathbb{Z}$  such that for any  $x, y, a \in \mathbb{Q}_p$  if

$$\text{val}(x - c) = \text{val}(y - c) < \text{val}(x - y) - D$$

then  $x - c, y - c$  are in the same coset of  $Q_{n,m}$ .

*Proof.* Define that  $a, b \in \mathbb{Q}_p$  are similar if  $\text{val } a = \text{val } b$  and

$$a \upharpoonright [\text{val } a, \text{val } a + (m + n)] = b \upharpoonright [\text{val } b, \text{val } b + (m + n)]$$

If  $a, b$  are similar then

$$a \in Q_{n,m} \leftrightarrow b \in Q_{n,m}$$

Moreover for any  $\lambda \in \mathbb{Q}_p$ , if  $a, b$  are similar we would also have  $a/\lambda, b/\lambda$  are similar. Thus if  $a, b$  are similar, then they belong in the same coset of  $Q_{n,m}$ . If we pick  $D = n + m$  then conditions of the lemma force  $x - c, y - c$  to be similar.  $\square$

The following construction is along the lines of lemmas 7.3, 7.5 of [1].

**Definition 2.10.** For two balls  $B(a, \alpha), B(b, \beta)$  let  $\gamma = \min(\alpha, \beta, \text{val}(a - b))$ . Define the distance between those two balls to be  $|\alpha - \gamma| + |\beta - \gamma|$ . In  $\mathbb{Q}_p$  value group is discrete and residue field is finite, so there are finitely many balls at a fixed distance from a given ball. Near balls of  $B(a, \alpha)$  are defined to be balls with distance  $\mathcal{D}$  from  $B(a, \alpha)$ . Enumerate those as:

$$B_1(a, \alpha), B_2(c, \alpha), \dots, B_{N_D}(a, \alpha)$$

Near balls partition the space

$$\{b \in \mathbb{Q}_p \mid |\text{val}(a - b) - \alpha| \leq D\}$$

**Definition 2.11.** Let  $c \in \mathbb{Q}_p$ . It lies in our tree in one of the intervals  $B(c_L, \alpha_L) \setminus B(c_U, \alpha_U)$ . Suppose  $c$  lies in one of the near balls of  $B(c_L, \alpha_L)$  or  $B(c_U, \alpha_U)$ . Then define its interval type to be the index of that near ball. Otherwise define its interval type to be the coset of  $c - c_U$  of  $Q$ . Denote the space of all the possible branch types  $\text{Bt}$ .

**Lemma 2.12.** If  $a, a'$  are in the same interval and have the same interval type then they have the same tree type.

*Proof.* First part of the tree type definition is satisfied as  $a, a'$  are in the same interval, so we only need to demonstrate that the corresponding  $Q$ -cosets match. Pick any element of our tree  $c_i(b)$ . We want to show that  $a - c_i(b), a' - c_i(b)$  are in the same  $Q$ -coset.

Suppose  $a$  is in one of the near balls. As  $a'$  has the same interval type, it has to be in the same near ball. By definition of the near ball we then have  $\text{val}(a - c_i(b)) = \text{val}(a' - c_i(b)) < \text{val}(a - a') - D$ . Thus by Lemma 2.10 we have  $a - c_i(b), a' - c_i(b)$  in the same  $Q$ -coset.

Now, suppose both  $a, a'$  aren't in any near balls. Label their interval as  $B(c_L, \alpha_L) \setminus B(c_U, \alpha_U)$ . Then we have

$$\alpha_L + D < \text{val}(a - c_U) < \alpha_U - D$$

$$\alpha_L + D < \text{val}(a' - c_U) < \alpha_U - D$$

as otherwise one (both) of them would be in one of the near balls. We have either  $\text{val}(c_U - c_i(b)) \geq \alpha_U$  or  $\text{val}(c_U - c_i(b)) \leq \alpha_L$  as otherwise it would contradict the definition of an interval.

Suppose it is the first case  $\text{val}(c_U - c_i(b)) \geq \alpha_U$ . Then

$$\text{val}(a - c_i(b)) = \text{val}(a - c_U) < \alpha_U - D \leq \text{val}(c_U - c_i(b)) - D$$

so by Lemma 2.10 we have  $a - c_i(b), a - c_U$  are in the same  $Q$ -coset. By a parallel argument we have  $a' - c_i(b), a' - c_U$  are in the same  $Q$ -coset. As we are assuming  $a, a'$  have the same tree type it implies that  $a - c_U, a' - c_U$  are in the same  $Q$ -coset. Thus by transitivity we get that  $a - c_i(b), a' - c_i(b)$  are in the same  $Q$ -coset.

For the second case, suppose  $\text{val}(c_U - c_i(b)) \leq \alpha_L$ . Then

$$\text{val}(a - c_i(b)) = \text{val}(c_U - c_i(b)) \leq \alpha_L < \text{val}(a - c_U) - D$$

so by Lemma 2.10 we have  $a - c_i(b), c_U - c_i(b)$  are in the same  $Q$ -coset. By a parallel argument we have  $a' - c_i(b), c_U - c_i(b)$  are in the same  $Q$ -coset. Thus by transitivity we get that  $a - c_i(b), a' - c_i(b)$  are in the same  $Q$ -coset.  $\square$

### 3. MAIN PROOF

Fix  $\gamma$  corresponding to  $\{\bar{p}_i\}_{i \in I}$  according to Lemma 2.2.

**Definition 3.1.** Denote  $\mathbb{Z}/p\mathbb{Z}^\gamma$  as  $\text{Ct}$ .

**Definition 3.2.** Let  $f : \mathbb{Q}_p^{|x|} \rightarrow \mathbb{Q}_p^I$  with  $f(\bar{c}) = (p_i(\bar{c}))_{i \in I}$ . Define the segment space  $\text{Sg}$  to be the image of  $f$ .

Given a tuple  $(a_i)_{i \in I}$  in the segment space look at the corresponding floors  $\{F(a_i)\}_{i \in I}$ . Those are ordered as elements of  $\mathbb{Z}$ . Partition the segment space by order type of  $\{F(a_i)\}$ . Work in a fixed partition  $\text{Sg}'$ . After relabeling we may assume that

$$F(a_1) \geq F(a_2) \geq \dots$$

Consider the (relabelled) sequence of vectors  $\bar{p}_1, \bar{p}_2, \dots, \bar{p}_I$ . There is a unique subset  $J \subset I$  such that all vectors with indices in  $J$  are linearly independent, and all vectors with indices outside of  $J$  are a linear combination of preceding vectors. For any index  $i \in I$  we call it independent if  $i \in J$  and we call it dependent otherwise.

Now, we define the following function

$$g : \text{Sg}' \rightarrow \text{Bt}^I \times \text{Pt}^J \times \text{Ct}^{I-J}$$

Let  $\bar{a} = (a_i)_{i \in I} \in \text{Sg}'$ . To define  $g(\bar{a})$  we need to specify where it maps  $\bar{a}$  in each individual component of the product.

For all  $a_i$  record its interval type  $\in \text{Bt}$ , giving the first component.

For  $a_j$  with  $j \in J$ , record the interval of  $a_j$ , giving the second component.

For the third component do the following computation. Pick  $a_i$  with  $i$  dependent. Let  $j$  be the largest independent index with  $j < i$ . Record  $a_i \upharpoonright [F(a_j) - \gamma, F(a_j)]$ .

**Lemma 3.3.** For  $\bar{a}, \bar{a}' \in \text{Sg}'$  if  $g(\bar{a}) = g(\bar{a}')$  then  $a_i, a'_i$  have the same tree type for all  $i \in I$ .

*Proof.* For each  $i$  we show that  $a_i, a'_i$  are in the same interval and have the same interval type, so the conclusion follows by Lemma 2.13.  $\text{Bt}$  records the interval type of each element, so if  $g(\bar{a}) = g(\bar{a}')$  then  $a_i, a'_i$  have the same interval type for all  $i \in I$ . Thus it remains to show that  $a_i, a'_i$  lie in the same interval for all  $i \in I$ . Suppose  $i$  is an independent index. Then by construction,  $\text{Pt}$  records the interval for  $a_i, a'_i$ , so those have to belong to the same interval. Now suppose  $i$  is dependent. Pick the largest  $j < i$  such that  $j$  is independent. We have  $F(a_i) \leq F(a_j)$  and  $F(a'_i) \leq F(a'_j)$ . Moreover  $F(a_j) = F(a'_j)$  as they are mapped to the same interval (using the earlier part of the argument as  $j$  is independent).

**Claim 3.4.**  $\text{val}(a_i - a'_i) > F(a_j) - \gamma$

*Proof.* Let  $\bar{x}, \bar{x}' \in \mathbb{Q}_p^{|x|}$  be some elements with

$$\begin{aligned}\bar{p}_k \cdot \bar{x} &= a_k \\ \bar{p}_k \cdot \bar{x}' &= a'_k \text{ for all } k \in I\end{aligned}$$

It is always possible to do that as  $\bar{a}, \bar{a}' \in \text{Sg}'$ . Let  $J'$  be the set of the independent indices less than  $i$ . We have

$$\text{val}(a_k - a'_k) > F(a_k) \text{ for all } k \in J'$$

as for the independent indices  $a_k, a'_k$  lie in the same interval.

$$\begin{aligned}\text{val}(a_k - a'_k) &> F(a_j) \text{ for all } k \in J' \text{ by monotonicity of } F(a_k) \\ \text{val}(\bar{p}_k \cdot \bar{x} - \bar{p}_k \cdot \bar{x}') &> F(a_j) \text{ for all } k \in J' \\ \text{val}(\bar{p}_k \cdot (\bar{x} - \bar{x}')) &> F(a_j) \text{ for all } k \in J'\end{aligned}$$

$J'$  and  $i$  match the requirements of Lemma 2.2 so we conclude

$$\begin{aligned}\text{val}(\bar{p}_i \cdot (\bar{x} - \bar{x}')) &> F(a_j) - \gamma \\ \text{val}(\bar{p}_i \cdot \bar{x} - \bar{p}_i \cdot \bar{x}') &> F(a_j) - \gamma \\ \text{val}(a_i - a'_i) &> F(a_j) - \gamma\end{aligned}$$

as needed, finishing the proof of the claim.  $\square$

Additionally  $a_i, a'_i$  have the same image in Ct component, so we have

$$\text{val}(a_i - a'_i) > F(a_j)$$

As  $F(a_i) \leq F(a_j)$ ,  $a_i, a'_i$  have to lie in the same interval.  $\square$

**Corollary 3.5.**  $\Psi(x, y)$  has VC-density  $\leq |x|$

*Proof.* Suppose we have  $c, c' \in \mathbb{Q}_p^{|x|}$  such that  $f(c), f(c')$  are in the same partition and  $g(f(c)) = g(f(c'))$ . Then by the previous lemma  $p_i(c)$  has the same tree type as  $p_i(c')$  for all  $i \in I$ . Then by Lemma 2.8  $c, c'$  have the same  $\Psi$ -type. Thus the number of possible  $\Psi$ -types is bounded by the size of the range of  $g$  times the number of possible partitions

$$(\text{number of partitions}) \cdot |Bt|^{|I|} \cdot |Pt|^{|J|} \cdot |Ct|^{|I-J|}$$

We have

$$\begin{aligned}|Bt| &= N_D + \text{number of cosets of } Q|Pt| \leq N \cdot I^2 \text{ (the only component dependent on } N) \\ |Ct| &= p^\gamma\end{aligned}$$

and there are at most  $|I|!$  many partitions of Sg. This gives us a bound

$$|I|! \cdot |Bt|^{|I|} \cdot (N \cdot |I|^2)^{|J|} \cdot p^{\gamma|I-J|} = O(N^{|J|})$$

Every  $p_i$  is an element of a  $|x|$ -dimensional vector space, so there can be at most  $|x|$  many independent vectors. Thus we have  $|J| \leq |x|$  and the bound follows.  $\square$

**Corollary 3.6.** In the language  $\mathcal{L}_{aff}$  we have  $\text{vc}(n) = n$ .

*Proof.* Previous lemma implies that  $\text{vc}(\phi) \leq \text{vc}(\Psi) \leq |x|$ . As choice of  $\phi$  was arbitrary, this implies that VC-density of any formula is bounded by the arity of  $x$ .  $\square$

#### REFERENCES

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E-mail address: bobkov@math.ucla.edu