SOME VC-DENSITY COMPUTATIONS IN SHELAH-SPENCER GRAPHS

ANTON BOBKOV

Abstract. We investigate vc-density in Shelah-Spencer graphs. We provide an upper bound on a formula-by-formula basis and show that there isn't a uniform lower bound, forcing the vc-function to be infinite. In addition we show that Shelah-Spencer graphs do not have a finite dp-rank, in particular they are not dp-minimal.

VC-density was studied in [1] by Aschenbrenner, Dolich, Haskell, MacPherson, and Starchenko as a natural notion of dimension for NIP theories. In a complete NIP theory T we can define the vc-function

$$\mathrm{vc}^T = \mathrm{vc} : \mathbb{N} \longrightarrow \mathbb{R} \cup \{\infty\}$$

where $\mathrm{vc}(n)$ measures the worst-case complexity of families of definable sets in an n-fold Cartesian power of the underlying set of a model of T (see 1.13 below for a precise definition of vc^T). We always have $\mathrm{vc}(n) \geq n$ for each n, and the simplest possible behavior is $\mathrm{vc}(n) = n$ for all n. Theories with the property that $\mathrm{vc}(1) = 1$ are known to be dp-minimal, i.e., having the smallest possible dp-rank (see Definition 1.15). It is not known whether there can be a dp-minimal theory which doesn't satisfy $\mathrm{vc}(n) = n$ (see [1], diagram in section 5.3).

In this paper, we investigate vc-density of definable sets in Shelah-Spencer graphs. First major model-theoretic breakthrough for these structures was made in [8]. In our description of Shelah-Spencer graphs we follow closely the treatment in [2]. A Shelah-Spencer graph is a limit of random structures $G(n, n^{-\alpha})$ for an irrational $\alpha \in (0, 1)$. Here $G(n, n^{-\alpha})$ is a random graph on n vertices with edge probability $n^{-\alpha}$.

Our first result is that in Shelah-Spencer graphs

$$\mathrm{vc}(n) = \infty \text{ for each } n.$$

We also show that Shelah-Spencer graphs don't have a finite dp-rank, which in particular implies that they are not dp-minimal. Our second result provides an upper bound on the vc-density of a given formula $\phi(x, y)$:

$$\mathrm{vc}(\phi) \leq D(\phi)$$

where $D(\phi)$ is an expression involving $|y|$ and number of vertices and edges defined by ϕ.

Section 1 introduces basic facts about VC-dimension and vc-density. More can be found in [1]. Section 2 summarizes notation and basic facts concerning Shelah-Spencer graphs. We direct the reader to [2] for a more in-depth treatment. In section 3 we introduce key lemmas that will be useful in our proofs. Section 4 computes a lower bound for vc-density to demonstrate that $\mathrm{vc}(n) = \infty$. Here we also do computations involving dp-rank. Section 5 computes an upper bound for vc-density on a formula-by-formula basis.

1. VC-dimension and vc-density

Throughout this section we work with a collection $\mathcal{F}$ of subsets of an infinite set X. We call the pair $(X, \mathcal{F})$ a <u>set system</u>.

Definition 1.1.

- Given a subset A of X, we define the set system $(A, A \cap \mathcal{F})$ where $A \cap \mathcal{F} = \{A \cap F \mid F \in \mathcal{F}\}$.
- For $A \subseteq X$ we say that $\mathcal{F}$ <u>shatters</u> A if $A \cap \mathcal{F} = \mathcal{P}(A)$ (the power set of A).

Definition 1.2. We say $(X, \mathcal{F})$ has <u>VC-dimension</u> n if the largest subset of X shattered by $\mathcal{F}$ is of size n. If $\mathcal{F}$ shatters arbitrarily large subsets of X, we say that $(X, \mathcal{F})$ has infinite VC-dimension. We denote the VC-dimension of $(X, \mathcal{F})$ by $\mathrm{VC}(X, \mathcal{F})$.

Note 1.3. We may drop X from the notation $\mathrm{VC}(X, \mathcal{F})$, as the VC-dimension doesn't depend on the base set and is determined by $(\bigcup \mathcal{F}, \mathcal{F})$.

Set systems of finite VC-dimension tend to have good combinatorial properties, and we consider set systems with infinite VC-dimension to be poorly behaved.

Another natural combinatorial notion is that of the dual system of a set system:

Definition 1.4. For $a \in X$ define $X_a = \{F \in \mathcal{F} \mid a \in F\}$. Let $\mathcal{F}^* = \{X_a \mid a \in X\}$. We call $(\mathcal{F}, \mathcal{F}^*)$ the dual system of $(X, \mathcal{F})$. The VC-dimension of the dual system of $(X, \mathcal{F})$ is referred to as the dual VC-dimension of $(X, \mathcal{F})$ and denoted by $\mathrm{VC}^*(\mathcal{F})$. (As before, this notion doesn't depend on X.)

Lemma 1.5 (see 2.13b in [3])**.** *A set system $(X, \mathcal{F})$ has finite VC-dimension if and only if its dual system has finite VC-dimension. More precisely*

$$\mathrm{VC}^*(\mathcal{F}) \leq 2^{1+\mathrm{VC}(\mathcal{F})}.$$

For a more refined notion of complexity of $(X, \mathcal{F})$ we look at the traces of our family on finite sets:

Definition 1.6. Define the shatter function $\pi_{\mathcal{F}} \colon \mathbb{N} \longrightarrow \mathbb{N}$ of $\mathcal{F}$ and the dual shatter function $\pi^*_{\mathcal{F}} \colon \mathbb{N} \longrightarrow \mathbb{N}$ of $\mathcal{F}$ by

$$\pi_{\mathcal{F}}(n) = \max \{|A \cap \mathcal{F}| \mid A \subseteq X \text{ and } |A| = n\}$$

$$\pi^*_{\mathcal{F}}(n) = \max \{\mathrm{atoms}(B) \mid B \subseteq \mathcal{F}, |B| = n\}$$

where $\mathrm{atoms}(B)$ = number of atoms in the boolean algebra of sets generated by B. Note that the dual shatter function is precisely the shatter function of the dual system: $\pi^*_{\mathcal{F}} = \pi_{\mathcal{F}^*}$.

A simple upper bound is $\pi_{\mathcal{F}}(n) \leq 2^n$ (same for the dual). If the VC-dimension of $\mathcal{F}$ is infinite then clearly $\pi_{\mathcal{F}}(n) = 2^n$ for all n. Conversely we have the following remarkable fact:

Theorem 1.7 (Sauer-Shelah '72, see [5], [6]). *If the set system $(X, \mathcal{F})$ has finite VC-dimension d then $\pi_{\mathcal{F}}(n) \leq \binom{n}{\leq d}$ for all n, where $\binom{n}{\leq d} = \binom{n}{d} + \binom{n}{d-1} + \ldots + \binom{n}{1}$.*

Thus the systems with a finite VC-dimension are precisely the systems where the shatter function grows polynomially. The vc-density of $\mathcal{F}$ quantifies the growth of the shatter function of $\mathcal{F}$:

Definition 1.8. Define the vc-density and dual vc-density of $\mathcal{F}$ as

$$\mathrm{vc}(\mathcal{F}) = \limsup_{n\to\infty} \frac{\log \pi_{\mathcal{F}}(n)}{\log n} \in \mathbb{R}^{\geq 0} \cup \{+\infty\},$$
$$\mathrm{vc}^*(\mathcal{F}) = \limsup_{n\to\infty} \frac{\log \pi^*_{\mathcal{F}}(n)}{\log n} \in \mathbb{R}^{\geq 0} \cup \{+\infty\}.$$

Generally speaking a shatter function that is bounded by a polynomial doesn't itself have to be a polynomial. Proposition 4.12 in [1] gives an example of a shatter function that grows like $n \log n$ (so it has vc-density 1).

So far the notions that we have defined are purely combinatorial. We now adapt VC-dimension and vc-density to the model theoretic context.

Definition 1.9. Work in a first-order structure M. Fix a finite collection of formulas $\Phi(x, y)$ in the language $\mathcal{L}(M)$ of M.

- For $\phi(x, y) \in \mathcal{L}(M)$ and $b \in M^{|y|}$ let

$$\phi(M^{|x|}, b) = \{a \in M^{|x|} \mid \phi(a, b)\} \subseteq M^{|x|}.$$

- Let $\Phi(M^{|x|}, M^{|y|}) = \{\phi(M^{|x|}, b) \mid \phi_i \in \Phi, b \in M^{|y|}\} \subseteq \mathcal{P}(M^{|x|})$.
- Let $\mathcal{F}_\Phi = \Phi(M^{|x|}, M^{|y|})$, giving rise to a set system $(M^{|x|}, \mathcal{F}_\Phi)$.
- Define the VC-dimension $\mathrm{VC}(\Phi)$ of Φ to be the VC-dimension of $(M^{|x|}, \mathcal{F}_\Phi)$, similarly for the dual.

- Define the vc-density $\mathrm{vc}(\Phi)$ of Φ to be the vc-density of $(M^{|x|}, \mathcal{F}_\Phi)$, similarly for the dual.

We will also refer to the vc-density and VC-dimension of a single formula ϕ viewing it as a one element collection $\Phi = \{\phi\}$.

Counting atoms of a boolean algebra in a model theoretic setting corresponds to counting types, so it is instructive to rewrite the shatter function in terms of types.

Definition 1.10.

$$\pi^*_\Phi(n) = \max\left\{\text{number of } \Phi\text{-types over } B \mid B \subseteq M, |B| = n\right\}.$$

Here a Φ-type over B is a maximal consistent collection of formulas of the form $\phi(x, b)$ or $\neg\phi(x, b)$ where $\phi \in \Phi$ and $b \in B$.

The functions π^*_Φ and $\pi^*_{\mathcal{F}_\Phi}$ do not have to agree, as one fixes the number of generators of a boolean algebra of sets and the other fixes the size of the parameter set. However, as the following lemma demonstrates, they both give the same asymptotic definition of dual vc-density.

Lemma 1.11.

$$\mathrm{vc}^*(\Phi) = \textit{degree of polynomial growth of } \pi^*_\Phi(n) = \limsup_{n\to\infty} \frac{\log \pi^*_\Phi(n)}{\log n}.$$

Proof. With a parameter set B of size n, we get at most $|\Phi|n$ sets $\phi(M^{|x|}, b)$ with $\phi \in \Phi, b \in B$. We check that asymptotically it doesn't matter whether we look at growth of boolean algebra of sets generated by n or by $|\Phi|n$ many sets. We have:

$$\pi^*_{\mathcal{F}_\Phi}(n) \le \pi^*_\Phi(n) \le \pi^*_{\mathcal{F}_\Phi}(|\Phi|n).$$

Hence:

$$\begin{aligned}
\mathrm{vc}^*(\Phi) &\le \limsup_{n\to\infty} \frac{\log \pi^*_\Phi(n)}{\log n} \le \limsup_{n\to\infty} \frac{\log \pi^*_{\mathcal{F}_\Phi}(|\Phi|n)}{\log n} = \\
&= \limsup_{n\to\infty} \frac{\log \pi^*_{\mathcal{F}_\Phi}(|\Phi|n)}{\log |\Phi|n} \frac{\log |\Phi|n}{\log n} = \limsup_{n\to\infty} \frac{\log \pi^*_{\mathcal{F}_\Phi}(|\Phi|n)}{\log |\Phi|n} \le \\
&\le \limsup_{n\to\infty} \frac{\log \pi^*_{\mathcal{F}_\Phi}(n)}{\log n} = \mathrm{vc}^*(\Phi).
\end{aligned}$$

□

One can check that the shatter function and hence VC-dimension and vc-density of a formula are elementary notions, so they only depend on the first-order theory of the structure M.

NIP theories are a natural context for studying vc-density. In fact we can take the following as the definition of NIP:

Definition 1.12. Define ϕ to be NIP if it has finite VC-dimension in a theory T. A theory T is NIP if all the formulas in T are NIP.

In a general combinatorial context (for arbitrary set systems), vc-density can be any real number in $0 \cup [1, \infty)$ (see [4]). Less is known if we restrict our attention to NIP theories. Proposition 4.6 in [1] gives examples of formulas that have non-integer rational vc-density in an NIP theory, however it is open whether one can get an irrational vc-density in this model-theoretic setting.

Instead of working with a theory formula by formula, we can look for a uniform bound for all formulas:

Definition 1.13. For a given NIP structure M, define the vc-function

$$\begin{aligned}
\mathrm{vc}^M(n) &= \sup\{\mathrm{vc}^*(\phi(x,y)) \mid \phi \in \mathcal{L}(M), |x| = n\} \\
&= \sup\{\mathrm{vc}(\phi(x,y)) \mid \phi \in \mathcal{L}(M), |y| = n\} \in \mathbb{R}^{\geq 0} \cup \{+\infty\}.
\end{aligned}$$

As before this definition is elementary, so it only depends on the theory of M. We omit the superscript M if it is understood from the context. One can easily check the following bounds:

Lemma 1.14 (Lemma 3.22 in [1]). *We have* $\mathrm{vc}(1) \geq 1$ *and* $\mathrm{vc}(n) \geq n\,\mathrm{vc}(1)$.

However, it is not known whether the second inequality can be strict or even whether $\mathrm{vc}(1) < \infty$ implies $\mathrm{vc}(n) < \infty$.

Dp-rank is a common measure used in study of NIP theories, with dp-minimality being a special case. Those notions originated in [7], and further studied in [10], showing, for example, that dp-rank is additive. Here it is easiest for us to define dp-rank in terms of vc-density over indiscernible sequences.

Definition 1.15. Work in a first-order structure M. Fix a finite collection of formulas $\Phi(x, y)$ in the language $\mathcal{L}(M)$ of M.

- Suppose $A = (a_i)_{i\in\mathbb{N}}$ is an indiscernible sequence with each $a_i \in M^{|x|}$. Let

$$\mathscr{I}(A, \Phi) = \{\phi(\bigcup_{i\in\mathbb{N}} a_i, b) \mid \phi \in \Phi, b \in M^{|y|}\} \subseteq \mathcal{P}(M^{|x|}).$$

  This gives rise to a set system $(M^{|x|}, \mathscr{I}(A, \Phi))$.
- Define

$$\mathrm{vc}_{\mathrm{ind}}(\Phi) = \sup\left\{\mathrm{vc}(\mathscr{I}(A, \Phi)) \mid A = (a_i)_{i\in\mathbb{N}} \text{ is indiscernible}\right\}.$$

- Dp-rank of a theory T is $\leq n$ if $\mathrm{vc}_{\mathrm{ind}}(\phi) \leq n$ for all formulas ϕ.
- A theory T is said to have finite dp-rank if its dp-rank is $\leq n$ for some n.
- A theory T is dp-minimal if its dp-rank ≤ 1.

Refer to [9] for the connection between to the classical definition of dp-rank and the definition given here.

2. Graph Combinatorics

Throughout this paper A, B, C, M will denote finite graphs, and $\mathbb{D}$ will be used to denote potentially infinite graphs. For a graph A the set of its vertices is denoted by $v(A)$, and the set of its edges by $e(A)$. Number of vertices of A will be denoted as $|A|$. Subgraph always means induced subgraph and $A \subseteq B$ means that A is a subgraph of B. For two subgraphs A, B of a larger graph, the union $A \cup B$ denotes the graph induced by $v(A) \cup v(B)$. Similarly, $A - B$ means a subgraph of A induced by the vertices of $v(A) - v(B)$. For $A \subseteq B \subseteq D$ and $A \subseteq C \subseteq D$, graphs B, C are said to be disjoint over A if $v(B) - v(A)$ is disjoint from $v(C) - v(A)$ and there are no edges from $v(B) - v(A)$ to $v(C) - v(A)$ in D.

For the remainder of the paper fix $\alpha \in (0, 1)$, irrational.

Definition 2.1.

- For a graph A let $\dim(A) = |A| - \alpha|e(A)|$.
- For A, B with $A \subseteq B$ define $\dim(B/A) = \dim(B) - \dim(A)$.
- We say that $A \leq B$ if $A \subseteq B$ and $\dim(A'/A) > 0$ for all $A \subsetneq A' \subseteq B$.
- Define A to be positive if for all $A' \subseteq A$ we have $\dim(A') \geq 0$.
- We work in theory S_α in the language of graphs axiomatized by:
  - Every finite substructure is positive.
  - Given a model $\mathbb{G}$ and graphs $A \leq B$, every embedding $f : A \longrightarrow \mathbb{G}$ extends to an embedding $g : B \longrightarrow \mathbb{G}$.

  (Here an embedding maps edges to edges and nonedges to nonedges.) This theory is complete and stable (see 5.7 and 7.1 in [2]). From now on fix an ambient model $\mathbb{G} \models S_\alpha$. This will be the only infinite graph we work with.
- For A, B positive, (A, B) is called a minimal pair if $A \subseteq B$, $\dim(B/A) < 0$ but $\dim(A'/A) \geq 0$ for all proper $A \subseteq A' \subsetneq B$. We call B a minimal extension of A. The dimension of a minimal pair is defined as $|\dim(B/A)|$.
- A sequence $\langle M_i \rangle_{0 \leq i \leq n}$ is called a minimal chain if (M_i, M_{i+1}) is a minimal pair for all $0 \leq i < n$.
- For a graph A with the tuple of vertices x let $\mathrm{diag}_A(x)$ be the atomic diagram of A, i.e. the first-order formula recording whether there is an edge between every pair of vertices.

- Given $A \subseteq B$ let

$$\phi_{A,B}(x) = \mathrm{diag}_A(x) \wedge \exists z \ \mathrm{diag}_B(x,z).$$

  Any graph isomorphic to B is called a witness of $\phi_{A,B}$.

- A formula $\phi_{A,B}$ is called a basic formula if there is a minimal chain $\langle M_i \rangle_{0 \le i \le n}$ such that $A = M_0$ and $B = M_n$.

Theorem 2.2 (Quantifier elimination, 5.6 in [2])**.** *In theory S_α every formula is equivalent to a boolean combination of basic formulas.*

Definition 2.3. A graph $S \subseteq \mathbb{D}$ is called N-strong if for any $S \subseteq T \subseteq D$ with $|T| - |S| \le N$ we have $S \le T$.

3. Basic Definitions and Lemmas

Definition 3.1. Suppose $\phi(x,y)$ is a basic formula. Define X to be the graph on vertices x with edges defined by ϕ. Similarly define Y. Note that X, Y are positive. Additionally, let Y' be a subgraph of Y induced by vertices of Y that are connected to $W - (X \cup Y)$, where W is a witness of ϕ.

Definition 3.2. Suppose A, B are subgraphs of $\mathcal{D}$ such that $v(A), v(B)$ are disjoint. Then define $\mathscr{E}(A,B)$ to be the number of edges between the vertices in $v(A)$ and the vertices in $v(B)$.

We will require the following lemmas from [2]:

Lemma 3.3. *[See 2.3 in [2]] Let $A, B \subseteq \mathbb{D}$. Then*

$$\dim(A \cup B/A) \le \dim(B/A \cap B).$$

Moreover,

$$\dim(A \cup B/A) = \dim(B/A \cap B) - \alpha E,$$

where E is the number of edges connecting the vertices of $B - A$ to the vertices of $A - B$.

Lemma 3.4. *[See 4.1 in [2]] Suppose A is a positive graph, with at least $1/\alpha + 2$ vertices. Then for any $\epsilon > 0$ there exists a graph B such that (A, B) is a minimal pair with dimension $\leq \epsilon$. Moreover, every vertex in A is connected to a vertex in $B - A$.*

Lemma 3.5. *[See 4.4 in [2]] Suppose A is a positive graph, and $\mathcal{G}$ a model of S_α. Then for any integer S there exists an embedding $f \colon A \longrightarrow \mathcal{G}$ such that $f(A)$ is S-strong in $\mathcal{G}$.*

Lemma 3.6. *[See 3.8 in [2]] For all $S > 0$ there exists $M = M(S, \alpha) \in \mathbb{N}$ with the following property. Suppose $A \subseteq \mathcal{G}$ where $\mathcal{G}$ is a model of S_α. Then there exists B with $A \subseteq B \subseteq \mathcal{G}$ such that B is S-strong in $\mathbb{G}$ and $|B| \leq M|A|$.*

We conclude this section by stating a couple of technical lemmas that will be useful in our proofs later.

Lemma 3.7. *Work in an ambient graph $\mathbb{D}$. Suppose we have a set B and a minimal pair (A, M) with $A \subseteq B$ and $\dim(M/A) = -\epsilon$. Then either $M \subseteq B$ or $\dim(M \cup B/B) < -\epsilon$.*

Proof. By Lemma 3.3

$$\dim(M \cup B/B) \leq \dim(M/M \cap B),$$

and as $A \subseteq M \cap B \subseteq M$

$$\dim(M/A) = \dim(M/M \cap B) + \dim(M \cap B/A).$$

In addition we are given $\dim(M/A) = -\epsilon$. If $M \nsubseteq B$ then $A \subseteq M \cap B \subsetneq M$ and by minimality $\dim(M \cap B/A) > 0$. Combining the inequalities above we obtain the desired result:

$$\dim(M \cup B/B) \leq \dim(M/M \cap B) = \dim(M/A) - \dim(M \cap B/A) < -\epsilon.$$

$\square$

Lemma 3.8. *Work in an ambient graph $\mathbb{D}$. Suppose we have a set B and a minimal chain $\langle M_i \rangle_{0 \leq i \leq n}$ with dimensions*

$$\dim(M_{i+1}/M_i) = -\epsilon_i$$

and $M_0 \subseteq B$. Let $\epsilon = \min_{0 \leq i \leq n} \epsilon_i$. Then either $M_n \subseteq B$ or $\dim((M_n \cup B)/B) < -\epsilon$.

Proof. Let $\bar{M}_i = M_i \cup B$. Then:

$$\dim(\bar{M}_n/B) = \dim(\bar{M}_n/\bar{M}_{n-1}) + \ldots + \dim(\bar{M}_2/\bar{M}_1) + \dim(\bar{M}_1/B).$$

Either $M_n \subseteq B$ or at least one of the summands above is nonzero. Apply previous lemma. □

Lemma 3.9. *Suppose we have a minimal pair (A, M) with dimension ϵ. Suppose we have some $B \subseteq M$. Then $\dim B/(A \cap B) \geq -\epsilon$. Moreover if $B \cup A \neq M$ then $\dim B/(A \cap B) \geq 0$.*

Proof. We have $\dim(B \cup A/A) \leq \dim B/(A \cap B)$ by Lemma 3.3. As $A \subseteq B \cup A \subseteq M$ we have $\dim(B \cup A/A) \geq -\epsilon$ by minimality. Moreover, minimality implies that it is positive if $B \cup A \neq M$. □

Lemma 3.10. *Suppose we have a minimal chain $\langle M_i \rangle_{0 \leq i \leq n}$ with dimensions*

$$\dim(M_{i+1}/M_i) = -\epsilon_i.$$

Let ϵ be the sum of all ϵ_i. Suppose we have a graph B with $B \subseteq M_n$. Then $\dim B/(M_0 \cap B) \geq -\epsilon$.

Proof. Let $B_i = B \cap M_i$. We have $\dim B_{i+1}/B_i \geq \dim M_{i+1}/M_i$ by the previous lemma. Thus

$$\dim B/(M_0 \cap B) = \dim B_n/B_0 = \sum \dim B_{i+1}/B_i \geq -\epsilon.$$

□

4. Lower bound

In this section we restrict our attention to the following family of basic formulas $\phi(x, y)$:

- All formulas have $Y' = Y$ (see Definition 3.1).
- All formulas define no edges between X and Y.
- Minimal chain of $\phi(x, y)$ consists of one step, that is we only have one minimal extension as opposed to a chain of minimal extensions.
- The dimension of that minimal extension is smaller than α.

We obtain a lower bound for the formulas that are boolean combinations of basic formulas written in the disjunctive-conjunctive form. First, define $\epsilon_L(\phi)$.

Definition 4.1. For a basic formula $\phi = \phi_{\langle M_i \rangle_{0 \le i \le n}}(x, y)$ let

- $\epsilon_i(\phi) = -\dim(M_i/M_{i-1})$.
- $\epsilon_L(\phi) = \sum_1^n \epsilon_i(\phi)$.

Definition 4.2 (Negation)**.** If ϕ is a basic formula, then define

$$\epsilon_L(\neg\phi) = \epsilon_L(\phi).$$

Definition 4.3 (Conjunction)**.** Take a collection of formulas $\phi_i(x, y)$ where each ϕ_i is a positive or a negative basic formula. If both positive and negative formulas are present then $\epsilon_L(\phi) = \infty$. We don't have a lower bound for that case. If different formulas define X or Y differently then $\epsilon_L(\phi) = \infty$. In the case of conflicting definitions the formula would have no realizations. Otherwise let

$$\epsilon_L\left(\bigwedge \phi_i\right) = \sum \epsilon_L(\phi_i).$$

Definition 4.4 (Disjunction)**.** Take a collection of formulas ψ_i where each instance is a conjunction as above all agreeing on X and Y. Then

$$\epsilon_L\left(\bigvee \psi_i\right) = \min \epsilon_L(\psi_i).$$

Theorem 4.5. *For a formula ψ as above we have*

$$\operatorname{vc}\psi \geq \left\lfloor \frac{Y(\psi)}{\epsilon_L(\psi)} \right\rfloor,$$

where $Y(\psi)$ is $\dim(Y)$ (as all basic componenets agree on Y).

Proof. First, work with a formula that is a conjunction of positive basic formulas $\psi = \bigwedge_{i\in I} \phi_i$. Then as we have defined above

$$\epsilon_L(\psi) = \sum_{i\in I} \epsilon_L(\phi_i).$$

If W_i is a witness of ϕ_i, let $S_i = |W_i|$. Let n_1 be the largest natural number such that

$$n_1\epsilon_L(\psi) < Y(\psi).$$

Let ϵ' be the smallest value among $\epsilon_L(\phi_i)$. Suppose it corresponds to the formula ϕ'. Let n_2 be the largest natural number such that

$$n_1\epsilon_L(\psi) + n_2\epsilon' < Y(\psi).$$

Fix some $N > n_1 + n_2$. Let

$$J = \{0 \leq j < N\} \subseteq \mathbb{N}.$$

Let a_j be a graph isomorphic to X for each $j \in J$, pairwise disjoint. Let $A = \bigcup_{1\leq j\leq N} a_j$. Let

$$S = |Y| + (n_1 + n_2 + 1)\sum_{i\in I} S_i.$$

By Lemma 3.5 the graph A can be embedded into $\mathbb{G}$ as an S-strong graph. Abusing notation, we identify A with this embedding. Thus we have $A \subseteq \mathbb{G}$, S-strong.

Let J_1, J_2 be disjoint subsets of J, of sizes n_1, n_2 respectively. Let b be a graph isomorphic to Y. For each $i \in I, j \in J_1$ let W_{ij} be a witness of $\phi_i(a_j, b)$. (Note that then $(a_j \cup b, W_{ij})$ is a minimal pair.) For each $j \in J_1$ let W_j be a union of $\{W_{ij}\}_{i\in I}$ disjoint over $a_j \cup b$. For each $j \in J_2$ let W_j be a witness of $\phi'(a_j, b)$. Let W' be a union of $\{W_j\}_{j\in J_1\cup J_2}$ disjoint over b. Let W be a union of W' and A disjoint over $\{a_j\}_{j\in J_1\cup J_2}$.

Claim 4.6. *We have $A \leq W$.*

Proof. Consider some $A \subsetneq B \subseteq W$. We need to show $\dim(B/A) > 0$. Let $\bar{A} = A\cup b$. We have

$$\dim(B/A) = \dim(B/B \cap \bar{A}) + \dim(B \cap \bar{A}/A).$$

Let $B_{ij} = B\cap W_{ij}$. Let $B_j = B\cap W_j$. To unify indices, relabel all the graphs above as $\{B_k\}_{k\in K}$ for some index set K. By the construction of W we have

$$\dim(B/B \cap \bar{A}) = \sum_{k\in K} \dim(B_k/B_k \cap \bar{A}).$$

Fix k. We have $B_k \subseteq W_k$, where W_k is a minimal extension of $M_0^k = a \cup b$ for some $a \in A$. Let ϵ_k be the dimension of this minimal extension. We have $\dim(B_k/B_k \cap \bar{A}) = \dim(B_k/a \cup (B \cap b))$.

Case 1: $B \cap b = b$. Then $M_0^k \subseteq B_k \subseteq W_k$ and

$$\dim(B_k/a \cup (B \cap b)) = \dim(B_k/M_0^k).$$

By minimality of (M_0^k, B_k) we have $\dim(B_k/M_0^k) \geq -\epsilon_k$. Thus

$$\dim(B/B \cap \bar{A}) \geq -\sum_{k\in K} \epsilon_k = -\left(n_1\epsilon_L(\psi) + n_2\epsilon'\right).$$

In addition

$$\dim(B \cap \bar{A}/A) = \dim(b) = Y(\psi).$$

Combining the two, we get

$$\dim(B/A) \geq Y(\psi) - \left(n_1\epsilon_L(\psi) + n_2\epsilon'\right),$$

which is positive by the construction of n_1, n_2 as needed.

Case 2: $B \cap b \subsetneq b$.

Claim 4.7. *We have* $\dim(B_k/B_k \cap \bar{A}) > 0$.

Proof. Recall that $\dim(B_k/B_k \cap \bar{A}) = \dim(B_k/a \cup (B \cap b))$. First, suppose that $B_k \cup M_0^k \neq W_k$. Then by Lemma 3.9 we get the required inequality. Thus we may assume that $B_k \cup M_0^k = W_k$. By Lemma 3.3 we have

$$\dim(B_k \cup M_0^k/M_0^k) = \dim(B_k/B_k \cap M_0^k) - \alpha E,$$

where E is the number of edges connecting the vertices of $B_k - M_0^k = B_k \cup M_0^k - M_0^k$ to the vertices of $M_0^k - B_k = M_0^k - B_k \cap M_0^k$. Noting that $B_k \cup M_0^k = W_k$, $\dim W_k/M_0^k = -\epsilon_k$, and $B_k \cap M_0^k = a \cup (B \cap b)$ we may rewrite the equality above as

$$\dim(B_k/a \cup (B \cap b)) = \alpha E - \epsilon,$$

and E is the number of edges connecting the vertices of $W_k - M_0^k$ to the vertices of $M_0^k - a \cup (B \cap b)$. As $Y = Y'$ and $B \cap b \subsetneq b$ we must have $E \geq 1$. But then as $\alpha > \epsilon$ we have $\dim(B_k/a \cup (B \cap b)) > 0$ as needed. $\square$

Now, recall that

$$\dim(B/A) = \dim(B \cap \bar{A}/A) + \sum_{k \in K} \dim(B_k/B_k \cap \bar{A}).$$

By the claim above each of $\dim(B_k/B_k \cap \bar{A}) > 0$, thus

$$\dim(B/A) > \dim(B \cap \bar{A}/A).$$

In addition

$$\dim(B \cap \bar{A}/A) = \dim(B \cap b) \geq 0,$$

as b is postive. Thus $\dim(B/A) > 0$ as needed. $\square$

As $A \leq W$ and $A \subseteq \mathbb{G}$, we can embed W into $\mathbb{G}$ over A. Abusing notation again, we identify W with its embedding $A \leq W \subseteq \mathbb{G}$. In particular, now we have $b \in \mathbb{G}$. Also note that

$$\dim(W/A) = Y(\psi) - (n_1\epsilon_L(\psi) + n_2\epsilon'),$$
$$|W| - |A| \leq |b| + (n_1 + n_2)\sum_{i\in I} S_i.$$

Lemma 4.8. *We have*

$$\{a_j\}_{j\in J_1} \subseteq \psi(A, b) \subseteq \{a_j\}_{j\in J_1\cup J_2}.$$

Proof. First inclusion $\{a_j\}_{j\in J_1} \subseteq \psi(A, b)$ is immediate from the construction of W, as W_{ij} witnesses that $\phi_i(a_j, b)$ holds. For the second inclusion, suppose that there is $a \in A - \{a_j\}_{j\in J_1\cup J_2}$ such that $\psi(a, b)$ holds. Let $W' \subseteq \mathbb{G}$ be a witness of $\phi_1(a, b)$. First, note that the case $W' \subseteq W$ is impossible as there are no edges between a and $W - a$, but there are edges between a and $W' - a$. Thus assume $W' \nsubseteq W$. As $(a \cup b, W')$ is minimal, by Lemma 3.7 we have $\dim(W' \cup W/W) < -\epsilon_1$. Therefore

$$\dim(W' \cup W/A) = \dim(W' \cup W/W) + \dim(W/A) < Y(\psi) - (n_1\epsilon_L(\psi) + n_2\epsilon') - \epsilon_1,$$

which is negative by the construction of n_1, n_2. Thus $A \not\leq W \cup W'$, as then it would have a positive dimension. Additionally,

$$|W' \cup W| - |A| \leq |W' - W| + |W| - |A| \leq S_1 + |b| + (n_1 + n_2)\sum_{i\in I} S_i \leq S,$$

but then this contradicts that A is S-strong, as then we would have $A \leq W \cup W'$. $\square$

In the construction of W we have chosen indices J_1, J_2 arbitrarily. In particular, suppose we let J_2 to be the last n_2 indices of J and J_1 an arbitrary n_1-element subset of the first $N - n_2$ elements of J. Each of those choices would then yield a different trace $\psi(A, b)$ by the lemma above. Thus $\psi(A, M^{|y|}) \geq \binom{N-n_2}{n_1}$ and therefore $\mathrm{vc}(\psi) \geq n_1$. By the definition of n_1 we have $n_1 = \left\lfloor \frac{Y(\psi)}{\epsilon_L(\psi)} \right\rfloor$, so this proves the theorem for ψ.

Now consider a formula which is a conjunction consisting of negative basic formulas $\psi = \bigwedge_{i \in I} \neg\phi_i$. Let $\bar{\psi} = \bigwedge_{i \in I} \phi_i$. Do the construction above for $\bar{\psi}$ and suppose its trace is $X \subseteq A$ for some b. Then over b the same construction gives trace $(A - X)$ for ψ. Thus we get as many traces as above, and the same bound.

Finally consider a formula which is a disjunction of formulas considered above $\theta = \bigvee_{k \in K} \psi_k$. Choose the one with the smallest ϵ_L, say ψ_k, and repeat the construction above for ψ_k. Any trace we obtain is automatically a trace for θ, and thus we get as many traces as above, and the same bound. □

Corollary 4.9. *VC-function is infinite in Shelah-Spencer random graphs:*

$$\mathrm{vc}(n) = \infty.$$

Proof. Let A be a graph consisting of $1/\alpha + 2 + n$ disconnected vertices. Fix $\epsilon > 0$. By Lemma 3.4, there exists B such that (A, B) is minimal with dimension $\leq \epsilon$. Consider a basic formula $\psi_{A,B}(x, y)$ where $|x| = 1/\alpha + 2$ and $|y| = n$. Then by the theorem above $\mathrm{vc}(n) \geq \mathrm{vc}(\psi_{A,B}) \geq \frac{n}{\epsilon}$. As ϵ was arbitrary, this number can be made arbitrarily large, giving $vc(n) = \infty$ as needed. □

Corollary 4.10. *Shelah-Spencer random graphs don't have finite dp-rank. In particular they are not dp-minimal.*

Proof. We would like to modify the proof of Theorem 4.5 such that A is indiscernible. Note that in the proof we can construct sets $A = \{a_j\}_{j \in J}$ of arbitrary length. Moreover for every finite $J' \subseteq J$, the set $A = \{a_j\}_{j \in J'}$ is still S-strong. Thus we can find an infinite set $A = \{a_j\}_{j \in \mathbb{N}}$ indiscernible and S-strong. Repeating the construction of the corollary above, we can obtain a formula with an arbitrarily large vc-density over the indiscernible sequence A. □

5. Upper bound

Consider a basic formula $\phi(x, y)$ with a minimal chain $\langle M_i \rangle_{0 \le i \le n_\phi}$ with dimensions $\dim(M_{i+1}/M_i) = -\epsilon_i$. Define

$$\epsilon(\phi) = \min \{\epsilon_i\}_{0 \le i \le n_\phi}$$
$$K(\phi) = |M_{n_\phi}|.$$

Now consider a finite collection of basic formulas

$$\Phi = \Phi(\vec{x}, \vec{y}) = \{\phi_i(\vec{x}, \vec{y})\}_{i \in I}\,.$$

Define

$$\epsilon(\Phi) = \min \{\epsilon(\phi_i)\}_{i \in I} \cup \{\alpha\}\,,$$
$$K(\Phi) = \max \{K(\phi_i)\}_{i \in I}\,,$$
$$D_1(\Phi) = \frac{K(\Phi)}{\epsilon(\Phi)},$$
$$D(\Phi) = |y| D_1(\Phi).$$

We claim that

Theorem 5.1. *If ϕ is a boolean combination of formulas from Φ, then* $\mathrm{vc}(\phi) \le D(\Phi)$.

Let

$$S = \left\lceil \left(\frac{|y|}{\epsilon(\phi)} + 1 \right) K(\phi) \right\rceil.$$

Suppose we have a finite parameter set $A_0 \subseteq \mathbb{G}^{|x|}$ with $|A_0| = N_0$. We would like to bound $\phi(A_0, \mathbb{G}^{|y|})$. Let $A_1 \subseteq \mathbb{G}$ consist of the components of the elements of A_0. Then $|A_1| \le |x| N_0$. Using Lemma 3.6 let A be a graph $A_0 \subseteq A \subseteq \mathbb{G}$, S-strong in $\mathbb{G}$. Let $N = |A|$. We have $N \le |x| N_0 M$ (where M is the constant from the Lemma

3.6). As $A_0 \subseteq A^{|x|}$ we have

$$\left|\phi(A_0, \mathbb{G}^{|y|})\right| \leq \left|\phi(A^{|x|}, \mathbb{G}^{|y|})\right|.$$

Therefore it suffices to bound $\left|\phi(A^{|x|}, \mathbb{G}^{|y|})\right|$.

Definition 5.2. For $A \subseteq \mathbb{G}^{|x|}, B \subseteq \mathbb{G}^{|y|}$ define

$$\Phi(A, B) = \{(a, i) \in A \times I \mid \mathbb{G} \models \phi_i(a, b)\} \subseteq A \times I$$

Lemma 5.3. *For $A \subseteq \mathbb{G}^{|x|}, B \subseteq \mathbb{G}^{|y|}$ if ϕ is a boolean combination of formulas from Φ then*

$$|\phi(A, B)| \leq |\Phi(A, B)|$$

Proof. Clear, as for all $a \in A, b \in B$ the set $\Phi(a, b)$ determines the truth value of $\phi(a, b)$. $\square$

Thus it suffices to bound $\left|\Phi(A^{|x|}, \mathbb{G}^{|y|})\right|$.

Definition 5.4.

- For all $i \in I, a \in A^{|x|}, b \in \mathbb{G}^{|y|}$ if $\phi_i(a, b)$ holds fix $W^i_{a,b} \subseteq \mathbb{G}$, a witness of this formula.
- For $b \in \mathbb{G}^{|y|}$ let

$$W_b = \bigcup \left\{ W^i_{a,b} \mid a \in A^{|x|}, i \in I, \mathbb{G} \models \phi_i(a, b) \right\}.$$

- For sets $C, B \subset \mathbb{G}$ define the <u>boundary</u> of C over B

$$\partial(C, B) = \{b \in B \mid \mathscr{E}(b, C - B) \neq \emptyset\}$$

  (see Definition 3.2 for $\mathscr{E}$).
- For $b \in \mathbb{G}^{|y|}$ let $\partial_b \subseteq A$ be the boundary $\partial(W_b, A)$.
- For $b \in \mathbb{G}^{|y|}$ let $\bar{W}_b = (W_b - A) \cup \partial_b$.

- For $b_1, b_2 \in \mathbb{G}^{|y|}$ we say that $b_1 \sim b_2$ if $\partial_{b_1} = \partial_{b_2}$, $b_1 \cap A = b_2 \cap A$, and there exists a graph isomorphism from $\bar{W}_{b_1} \cup b_1$ to $\bar{W}_{b_2} \cup b_2$ that fixes ∂_{b_1} and maps b_1 to b_2. One easily checks that this defines an equivalence relation.
- For $b \in \mathbb{G}^{|y|}$ define $\mathscr{I}_b$ to be the $\sim$-equivalence class of b.

Lemma 5.5. *For $b_1, b_2 \in \mathbb{G}^{|y|}$ if $b_1 \sim b_2$ then $\Phi(A^{|x|}, b_1) = \Phi(A^{|x|}, b_2)$.*

Proof. Fix a graph isomorphism $\bar{f} \colon \bar{W}_{b_1} \cup b_1 \longrightarrow \bar{W}_{b_2} \cup b_2$. Extend it to an isomorphism $f \colon W_{b_1} \cup A \longrightarrow W_{b_2} \cup A$ by being an identity map on the new vertices. Suppose $\mathbb{G} \models \phi_i(a, b_1)$ for some $a \in A^{|x|}$. Then $f(W^i_{a,b_1})$ is a witness for $\phi_i(a, b_2)$ (though not necessarily equal to W^i_{a,b_2}) and thus $\mathbb{G} \models \phi_i(a, b_2)$. As a was arbitrary, this proves the equality of the traces. $\square$

Thus to bound the number of traces it is sufficient to bound the number of possibilities for $\mathscr{I}_b$.

Theorem 5.6. *Suppose we have $b \in \mathbb{G}^{|y|}$. Let $Y = |b - A|$. Then*

$$|\partial_b| \leq Y D_1(\phi)$$

$$|\bar{W}_b| \leq 3Y D_1(\phi)$$

From this theorem we get the desired result:

Corollary 5.7. *(Theorem 5.1) If ϕ is a boolean combination of formulas from Φ, then $\mathrm{vc}(\phi) \leq D(\Phi)$.*

Proof. We count possible equivalence classes of $\sim$. This essentially boils down to bounding possibilities for ∂_b, $b \cap A$, and number of isomorphism classes of W_b. Fix $b \in \mathbb{G}^{|y|}$ and let $Y = |b - A|$. Let $D = Y D_1(\Phi)$. By the first part of Theorem 5.6 there are $\binom{N}{D}$ choices for boundary ∂_b. By the second part of Theorem 5.6 there are at most $3D$ vertices in $\bar{W}_b$. Thus to determine the graph $\bar{W}_b$ we need to choose how many vertices it has and then decide where edges go. This amounts to at most $3D2^{(3D)^2}$ choices. Additionally there are $\binom{N}{|y|-Y}$ choices for $b \cap A$. In total this gives us at most

$$\binom{N}{D} \cdot \binom{N}{|y|-Y} \cdot 3D2^{(3D)^2} = O(N^{D+|y|-Y}) =$$
$$= O(N^{YD_1(\Phi)+|y|-Y}) = O(N^{|y|D_1(\Phi)}) = O(N^{D(\Phi)})$$

choices (second to last inequality uses $D_1(\Phi) \geq 1$). By Lemma 5.5 we have $\left|\Phi(A^{|x|}, \mathbb{G}^{|y|})\right| = O(N^{D(\Phi)})$. Recall that

$$\left|\phi(A_0, \mathbb{G}^{|y|})\right| \leq \left|\Phi(A^{|x|}, \mathbb{G}^{|y|})\right|.$$

Therefore we have

$$\left|\phi(A_0, \mathbb{G}^{|y|})\right| = O(N^{D(\Phi)}) = O((|x|N_0M)^{D(\Phi)}) =$$
$$= O((|x|M)^{D(\Phi)}\, N_0^{D(\Phi)}) = O(N_0^{D(\Phi)}).$$

As A_0 was arbitrary, this shows that $\mathrm{vc}(\phi) \leq D(\Phi)$ as needed. (Note that throughout this proof we are using the fact that quantities $D_1(\Phi), D(\Phi), M$ are completely determined by Φ, thus independent from A_0.) $\square$

Proof. (of Theorem 5.6)

The graph W_b is a union of witnesses of the from $W_{a,b}$ for some $a \in A^{|x|}, b \in \mathbb{G}^{|y|}$. Enumerate all of them as $\{W_j\}_{1\leq j\leq J}$. Define $M_j = \bigcup_1^j W_{j'}$ for $1 \leq j \leq J$ and let $M_0 = b$. Let $\bar{A} = A \cup b$.

Definition 5.8. For $0 \leq j \leq J$ define:

- Let $v_j = 1$ if new vertices are added to M_j outside of $\bar{A}$, that is if $M_j - \bar{A} \neq M_{j-1} - B$, and let it be 0 otherwise.
- Let $E_j = \partial(A - W_j, M_j - A)$.
- Let

$$m_j = \sum_{j'=0}^{j} (v_j + |E_j|).$$

(Here assume $M_{-1} = \emptyset$.)

Lemma 5.9. *For $0 \leq j \leq J$ we have*

$$|\partial(M_j, A)| \leq |E_0| + m_j K(\Phi)$$

Proof. Proceed by induction. The base case $j = 0$ is clear. For an induction step suppose that

$$|\partial(M_{j-1}, A)| \leq m_{j-1} K(\Phi)$$

holds. Let

$$\begin{aligned} \delta_1 &= \partial(M_j, A) - \partial(M_{j-1}, A) = \\ &= \{a \in A \mid \mathscr{E}(a, M_j - A) \neq \emptyset \text{ and } \mathscr{E}(a, M_{j-1} - A) = \emptyset\}. \end{aligned}$$

If $M_j - A = M_{j-1} - A$ then $\delta_1 = \emptyset$ and we are done as m_j is increasing. Suppose not. We have $|\delta_1| = |\delta_1 \cap W_j| + |\delta_1 - W_j|$, and

$$\delta_1 - W_j = \{a \in A - W_j \mid \mathscr{E}(a, M_j - A) \neq \emptyset \text{ and } \mathscr{E}(a, M_{j-1} - A) = \emptyset\}.$$

But then it's clear that $\delta_1 - W_j \subseteq E_j$ as

$$\begin{aligned} &W_j - M_{j-1} - A \subseteq M_j - A, \\ &(W_j - M_{j-1} - A) \cap (M_{j-1} - A) = \emptyset. \end{aligned}$$

As $b \in M_{j-1}$ and $M_j - A \neq M_{j-1} - A$, then $M_j - \bar{A} \neq M_{j-1} - \bar{A}$, and thus $v_j = 1$. Therefore we have

$$\begin{aligned} |\delta_1| &= |\delta_1 \cap W_j| + |\delta_1 - W_j| \leq |W_j| + |E_j| \leq \\ &\leq K(\Phi) + |E_j| \leq (v_j + |E_j|) K(\Phi) \leq (m_j - m_{j-1}) K(\Phi), \end{aligned}$$

as needed. $\square$

Lemma 5.10. *For $0 \leq j \leq J$ we have*

$$|M_j - \bar{A}| \leq \sum_{j'=0}^{j} v_{j'} K(\Phi)$$

Proof. Proceed by induction. The base case $j = 0$ is clear. For an induction step suppose that

$$|M_{j-1} - \bar{A}| \leq \sum_{j'=0}^{j-1} v_{j'} K(\Phi)$$

holds. If $M_j - \bar{A} = M_{j-1} - \bar{A}$ then the inequality is immediate as $v_j \geq 0$. Therefore assume this is not the case, so $v_j = 1$ and $|M_j - A| - |M_{j-1} - A| \leq |W_j| \leq v_j K(\Phi)$, and so we get the required inequality.

□

Lemma 5.11. *For $0 \leq j \leq J$ we have*

$$\dim(M_j \cup \bar{A}/\bar{A}) \leq -m_j \epsilon(\Phi),$$

Proof. Proceed by induction. Base case $j = 0$ is clear. For an induction step suppose that

$$\dim(M_{j-1} \cup \bar{A}/\bar{A}) \leq -m_{j-1} \epsilon(\Phi)$$

holds. We have

$$\dim(M_j \cup \bar{A}/\bar{A}) = \dim(M_j \cup \bar{A}/M_{j-1} \cup \bar{A}) + \dim(M_{j-1} \cup \bar{A}/\bar{A}) \leq$$

$$\leq \dim(M_j \cup \bar{A}/M_{j-1} \cup \bar{A}) - m_{j-1} \epsilon(\Phi).$$

Let $\bar{M}_{j-1} = M_{j-1} \cup \bar{A}$. By Lemma 3.3

$$\dim(M_j \cup \bar{A}/M_{j-1} \cup \bar{A}) = \dim(W_j \cup \bar{M}_{j-1}/\bar{M}_{j-1}) = \dim(W_j/W_j \cap \bar{M}_{j-1}) - e\alpha$$

where e is the number of edges connecting the vertices of $\bar{M}_{j-1} - W_j$ to the vertices of $W_j - \bar{M}_{j-1}$. Recall that $E_j = \partial(A - W_j, M_j - A)$. We have $A - W_j \subseteq \bar{M}_{j-1} - W_j$ (as $A \subseteq \bar{M}_{j-1}$) and $W_j - M_{j-1} - A = W_j - \bar{M}_{j-1}$ (as for $j > 1$, we have $b \subseteq M_{j-1}$). Thus $|E_j| \leq e$, and we get

$$\dim(M_j \cup \bar{A}/M_{j-1} \cup \bar{A}) \leq \dim(W_j/W_j \cap \bar{M}_{j-1}) - |E_j|\alpha.$$

If $W_j \subseteq \bar{M}_{j-1}$ then $\dim(W_j/W_j \cap \bar{M}_{j-1}) = 0$. If not, then by Lemma 3.8 we have $\dim(W_j/W_j \cap \bar{M}_{j-1}) \leq -\epsilon(\Phi)$. Either way, we have $\dim(W_j/W_j \cap \bar{M}_{j-1}) \leq -v_j\epsilon(\Phi)$. Using this and the fact that $\epsilon(\Phi) \leq \alpha$, we obtain

$$\dim(M_j \cup \bar{A}/M_{j-1} \cup \bar{A}) \leq -v_j\epsilon(\Phi) - |E_j|\epsilon(\Phi) = -(m_j - m_{j-1})\epsilon(\Phi).$$

Finally,

$$\begin{aligned}\dim(M_j \cup \bar{A}/\bar{A}) &\leq \dim(M_j \cup \bar{A}/M_{j-1} \cup \bar{A}) - m_{j-1}\epsilon(\Phi) \leq \\ &\leq -(m_j - m_{j-1})\epsilon(\Phi) - m_{j-1}\epsilon(\Phi) = -m_j\epsilon(\Phi),\end{aligned}$$

as needed. □

(Proof of Theorem 5.6 continued) For any $0 \leq j \leq J$ we have

$$\begin{aligned}\dim(M_j \cup A/A) &= \dim(\bar{A}/A) + \dim(M_j \cup \bar{A}/\bar{A}) \\ &\leq Y - |E_0|\alpha + \dim(M_j \cup \bar{A}/\bar{A}).\end{aligned}$$

Lemma 5.11 gives us

$$\dim(M_j \cup \bar{A}/\bar{A}) \leq -m_j\epsilon(\Phi).$$

Thus

$$\dim(M_j \cup A/A) \leq Y - |E_0|\alpha - m_j\epsilon(\Phi).$$

Suppose j is an index such that

$$Y - |E_0|\alpha - m_j\epsilon(\Phi) \geq 0,$$

$$Y - |E_0|\alpha - m_{j+1}\epsilon(\Phi) < 0$$

if one exists. Then

$$m_j \leq \frac{Y - |E_0|\alpha}{\epsilon(\Phi)}.$$

By Lemma 5.10 we have

$$\begin{aligned} |M_{j+1} - A| &\leq \left(\sum_{j'=1}^{j+1} v_{j'}\right) K(\Phi) \leq (m_j + 1)K(\Phi) \\ &\leq \left(\frac{Y - |E_0|\alpha}{\epsilon(\Phi)} + 1\right) K(\Phi) \leq S. \end{aligned}$$

This is a contradiction, as A is S-strong and $\dim(M_{j+1} \cup A/A)$ is negative. Thus $Y - |E_0|\alpha - m_j\epsilon(\Phi) \geq 0$ for all $j \leq J$. In particular $Y - |E_0|\alpha - m_J\epsilon(\Phi) \geq 0$, so $m_J \leq \frac{Y - |E_0|\alpha}{\epsilon(\Phi)}$. Noting that $M_J = W_b$, Lemma 5.9 gives us

$$|\partial_b| = |\partial(W_b, A)| \leq |E_0| + m_J K(\Phi) \leq |E_0| + K(\Phi)\frac{Y - |E_0|\alpha}{\epsilon(\Phi)}.$$

As $K(\Phi) \geq 1$ and $\epsilon(\Phi) \geq \alpha$, we get

$$|\partial_b| \leq K(\Phi)\frac{Y}{\epsilon(\Phi)} = YD_1(\Phi).$$

But this is precisely the first inequality we need to prove. For the second inequality, Lemma 5.10 gives us

$$\begin{aligned} \left|W_b - \bar{A}\right| &\leq Y + \left(\sum_{j'=0}^{J} v_{j'}\right) K(\Phi) \leq Y + m_J K(\Phi) \leq \\ &\leq Y + K(\Phi)\frac{Y}{\epsilon(\Phi)} \leq 2YD_1(\Phi). \end{aligned}$$

Thus we have

$$|\bar{W}_b| \leq |W_b - A| + |\partial_b| \leq 3YD_1(\Phi),$$

as needed. This ends the proof for Theorem 5.6. □

6. Conclusion

This paper computes upper and lower bounds for certain types of formulas in Shelah-Spencer graphs. The bounds are not tight: in the best case scenario for a basic formula $\phi(x, y)$ defining a minimal extension of dimension ϵ we have

$$\frac{|y|}{\epsilon} \leq \mathrm{vc}(\phi) \leq K\frac{|y|}{\epsilon},$$

where K is the number of vertices in the minimal extension. Thus there is a multiple of K gap between lower and upper bounds. It is this author's hope that a refinement of presented techniques can yield better estimates of the vc-density. One potential direction towards this goal is to have a closer study on how multiple minimal extensions can intersect without increasing overall dimension.

Note that this paper doesn't answer the question whether there can be exotic values for vc-density of individual formulas, such as non-integer or irrational values. A better bound can help address this question.

Another observation is that while $\mathrm{vc}(n) = \infty$ there seems to be a good structural behavior of the vc-density for individual formulas. This perhaps suggests that the vc-function is not the best tool to describe behaviour of the definable sets in Shelah-Spencer graphs, and some more refined measure might be required. One potential way to do this is to separate the formulas based on values of $K(\phi), \epsilon(\phi)$. Once those are bounded, vc-density seems to be well-behaved. This author hopes to explore this further in his future work.

References

[1] M. Aschenbrenner, A. Dolich, D. Haskell, D. Macpherson, S. Starchenko, *Vapnik-Chervonenkis density in some theories without the independence property*, I, Trans. Amer. Math. Soc. 368 (2016), 5889-5949

[2] Michael C. Laskowski, *A simpler axiomatization of the Shelah-Spencer almost sure theories*, Israel J. Math. **161** (2007), 157186. MR MR2350161

[3] P. Assouad, *Densite et dimension*, Ann. Inst. Fourier (Grenoble) 33 (1983), no. 3, 233-282.

[4] P. Assouad, *Observations sur les classes de Vapnik-Cervonenkis et la dimension combinatoire de Blei*, in: Seminaire dAnalyse Harmonique, 1983-1984, pp. 92-112, Publications Mathematiques dOrsay, vol. 85-2, Universite de Paris-Sud, Departement de Mathematiques, Orsay, 1985.

[5] N. Sauer, *On the density of families of sets*, J. Combinatorial Theory Ser. A 13 (1972), 145-147.

[6] S. Shelah, *A combinatorial problem; stability and order for models and theories in infinitary languages*, Pacific J. Math. 41 (1972), 247-261.

[7] S. Shelah, *Strongly dependent theories* (2009), preprint.

[8] S. Shelah and J. Spencer, *Zero-one laws for sparse random graphs*, Journal of AMS 1 (1988) 97-115.

[9] V. Guingona, C. Hill, *Local dp-rank and VC-density over indiscernible sequences*, Math. Log. Quart., 60: 59-65

[10] I. Kaplan, A. Onshuus, and A. Usvyatsov, *Additivity of the dp-rank*, Trans. Amer. Math. Soc. 365 (2013), 5783-5804

E-mail address: `bobkov@math.ucla.edu`