

# 143A: Principles of Operating Systems

## Lecture 7: System boot

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October, 2017

# Outline for today

## Boot operating system

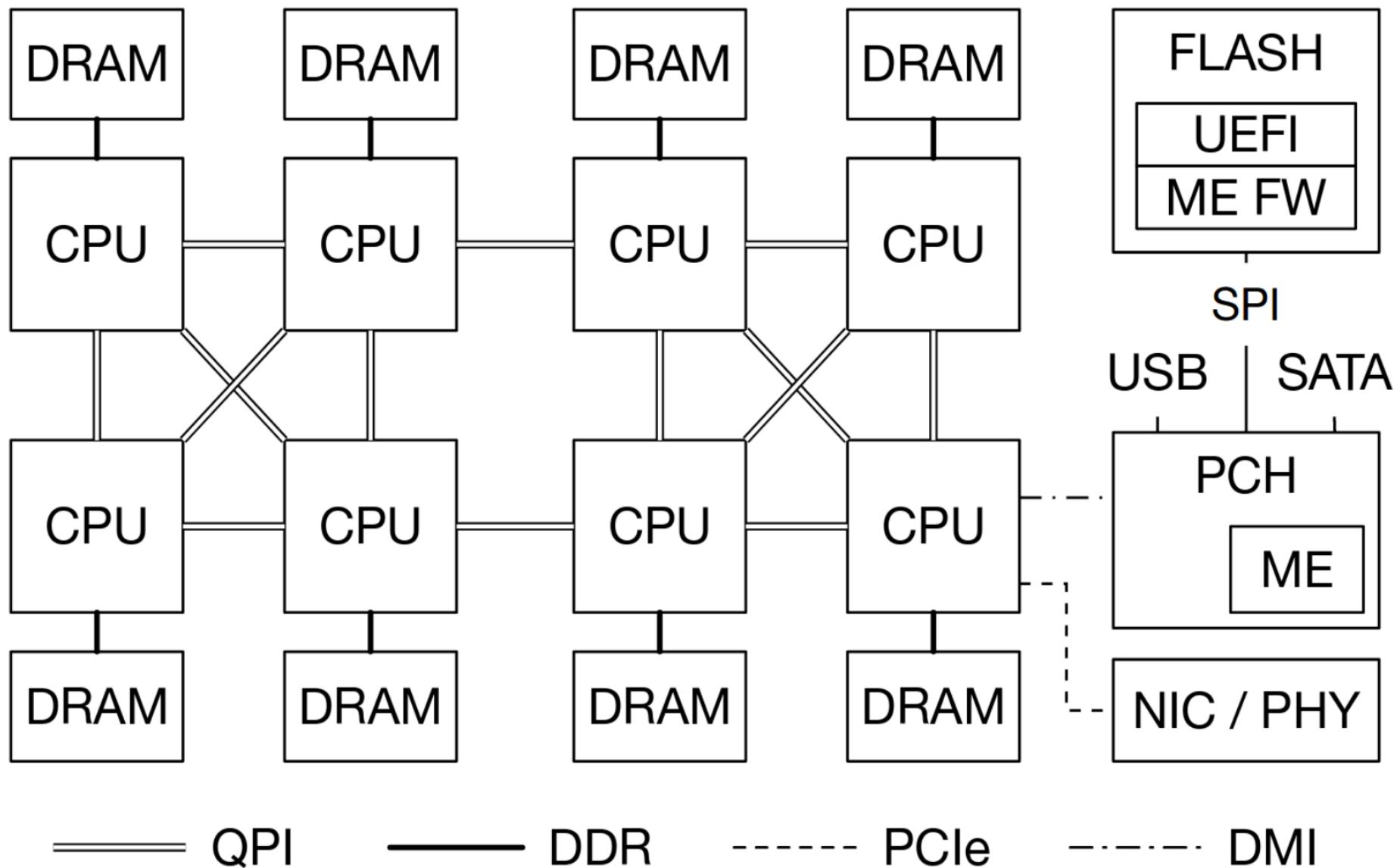
- Setup segments (data and code)
- Switch to protected mode
  - Load GDT (turn segmentation on)
- Setup stack (needed to call C functions)
- Load the kernel from disk into memory
- Setup first page table
  - 2 entries [ 0 : 4MB ] and [ 2GB : (2GB + 4MB) ]
- Setup high-address stack
- Jump to main()
  - Start executing kernel code

# What happens when we turn on the power?

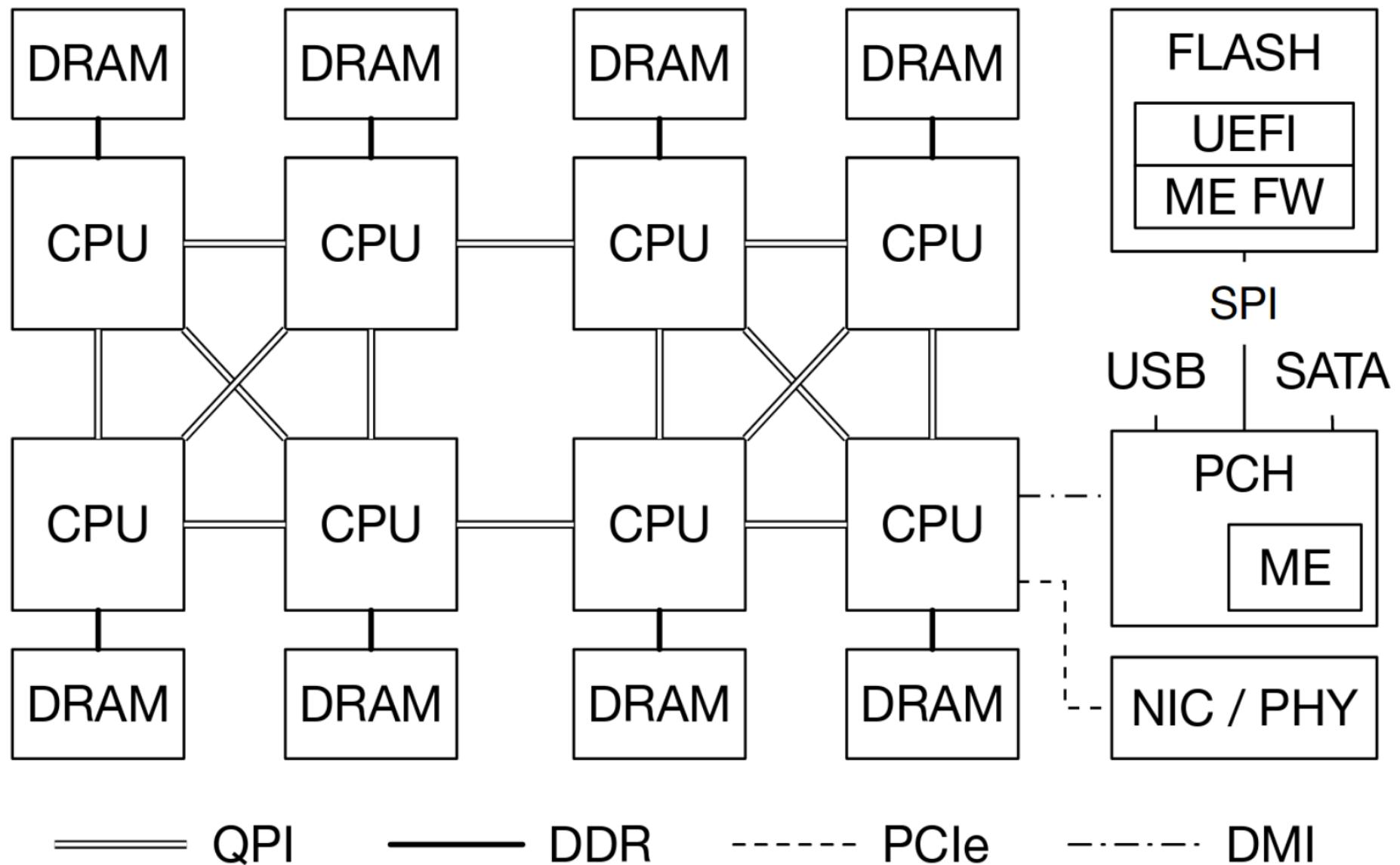
- Well it's complicated
  - Intel SGX Explained is a good start (Section 2.13 [1])
- At a high-level a sequence of software pieces initializes the platform
  - Microcode, firmware (BIOS), bootloader

- The most important thing, the OS is not the only software running on the machine
  - And not the most privileged
- Today, at least two layers sit underneath the OS/hypervisor
  - System Management Mode (SMM) (ring -2)
    - Runs below the hypervisor/OS
  - Intel Management Engine (ring -3)
    - Runs on a separate CPU

# PC motherboard components



# ME gets power before CPUs



# Intel Management Engine (ME)

- Full-featured computer
  - Argonaut RISC Core (ARC), 200-400MHz
  - Internal RAM (640KB)
  - Can access all DRAM via DMA
  - Can control boot chain
  - Can access network interface (NIC) on the motherboard
    - Via System Management Bus (SMBus)
    - Or an ATM compatible NIC
  - Connected to the power supply
    - Stays on as long as power is provided to power supply

# ME: Theft prevention use-case

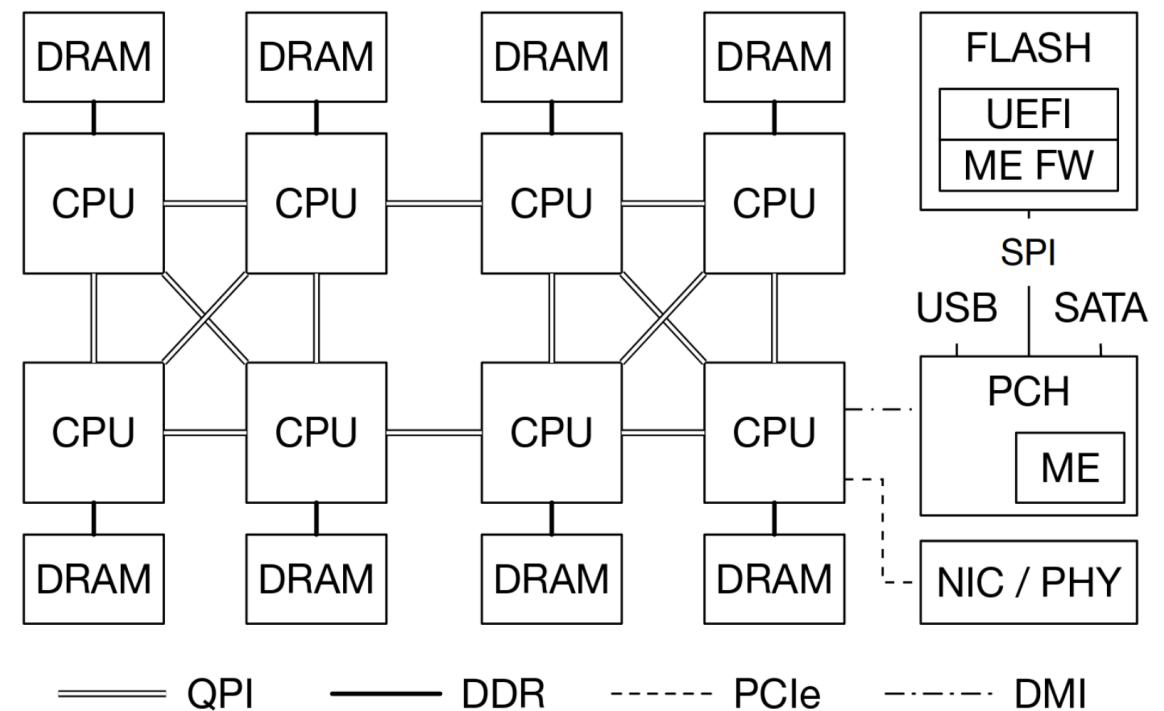
- In S5 (computer off) ME cannot access DRAM
  - DRAM is off
  - But ME can use its internal memory
  - ME can disable a stolen laptop equipped with cellular modem remotely
    - As long as power is connected
    - And cell network has signal

# Intel Management Engine (ME)

- All modern Intel CPUs contain ME
  - Part of Active Management Technology (AMT)
  - Convenient way for administrators to fix your machine remotely
    - Obviously a huge opportunity for an attack

# ME starts first

- Reads its initialization code from the BIOS chip
  - Via the SPI bus



# Bootstrap processor (BSP)

- One of the logical processors is chosen as bootstrap processor (BSP)
  - Will start initialization
- Others become “application processors” (AP)
  - Waiting for a special interrupt from the BSP

# BSP starts reading BIOS

- Executes instructions stored in the BIOS chip
- An interesting detail is that BSP starts with DRAM disabled
  - Hence there is no stack to call functions
  - What can be done?

# BSP starts without DRAM

- Custom-written assembly code that uses no stack
- Or a ROMCC compiler
  - Generates code from C that uses no stack
  - Used in the coreboot project

# Cache-as-RAM

- Use CPU caches as temporary replacement for RAM
  - Initialize DRAM
  - Copy BIOS firmware into DRAM and continue

# BIOS frimware

- Initialize
  - Interrupt controllers
  - Devices, e.g., network interfaces
  - If one of PCI devices contains “option ROM” load and execute it
    - Network cards may contain iPXE ROM
    - Implement boot from the network host

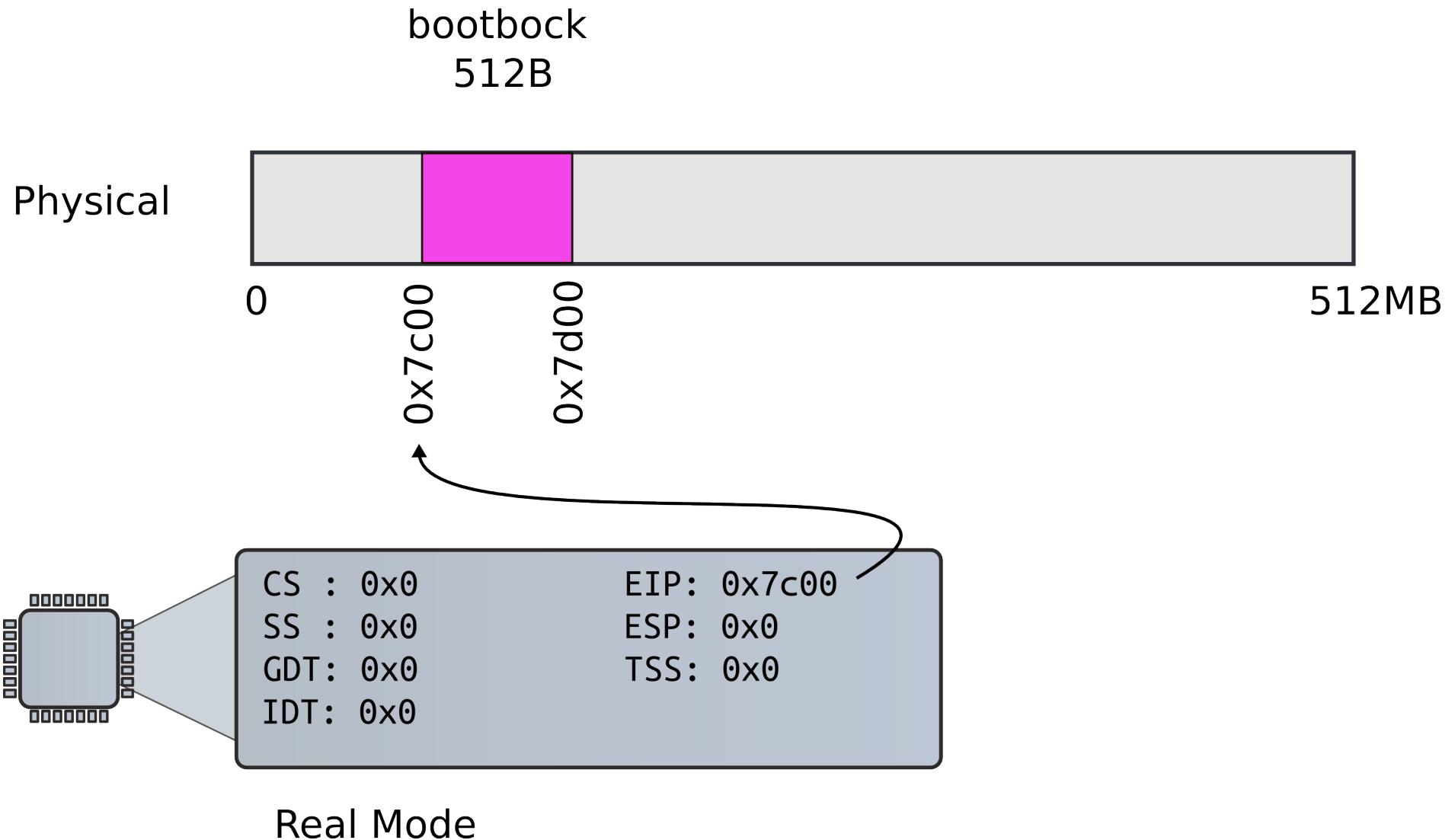
# System Management Mode

- Another compartment that runs underneath your OS or hypervisor
- Initialized by BIOS
- Protected with hardware memory mechanisms
  - OS cannot access this region of memory
- Runs under your OS or hypervisor
  - Receives interrupts periodically, can take over the entire system any time
  - Impossible to disable

# BIOS loads the boot loader

- BIOS ends by loading a boot loader
  - Modern BIOSes can load the boot loader from a variety of sources (hard disks, USB drives, optical disks)
  - Default way is to load the first sector (512 bytes) from disk into the memory location at 0x7c00
  - BIOS then starts executing instructions at the address 0x7c00
    - This is exactly what we see when we run xv6 under QEMU
    - QEMU emulates hardware: runs BIOS, follows the same protocol

# BIOS loads bootloader



# Bootloader starts

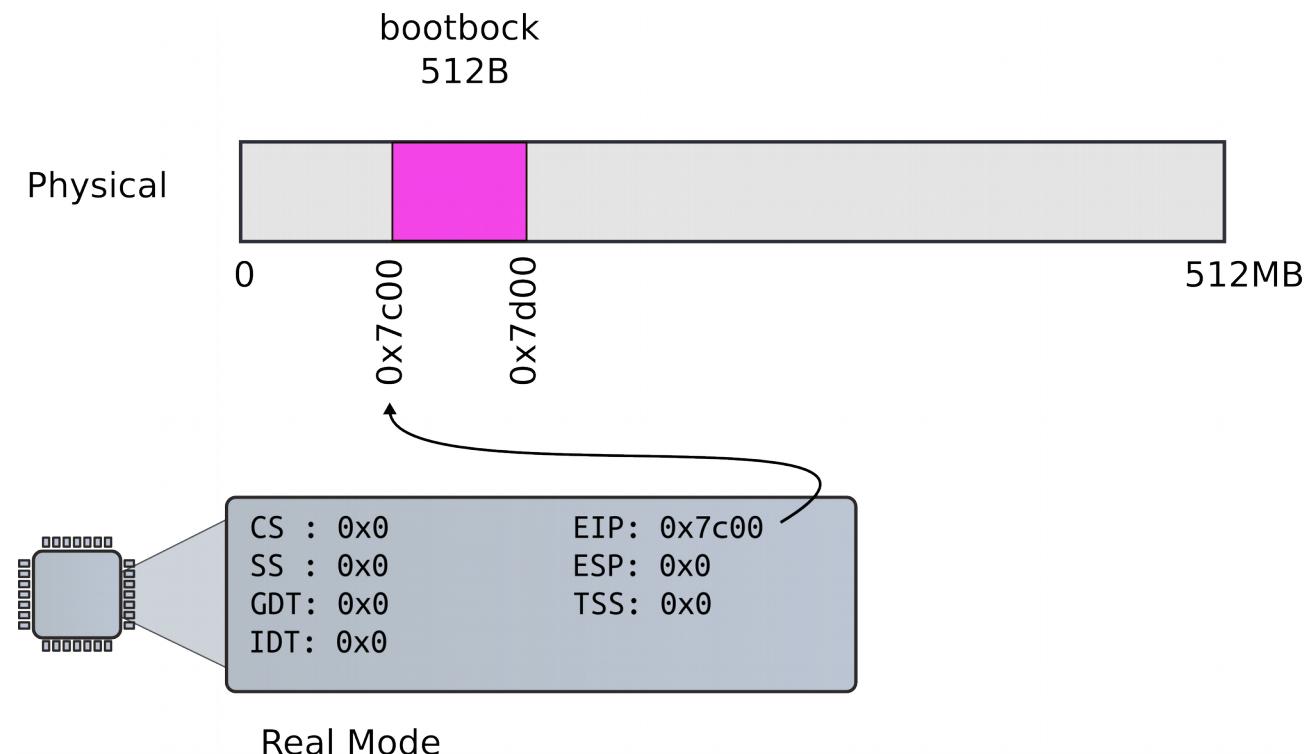
```
9111 start:  
9112     cli # BIOS enabled interrupts; disable  
9113  
9114     # Zero data segment registers DS,ES, and  
              SS.  
9115     xorw %ax,%ax # Set %ax to zero  
9116     movw %ax,%ds # -> Data Segment  
9117     movw %ax,%es # -> Extra Segment  
9118     movw %ax,%ss # -> Stack Segment
```

# Why start happens to be 0x7c00?

9111 start:

9112 cli # BIOS enabled interrupts; disable

9113



# Linker is instructed to link the boot block code in the Makefile

```
9111 start:  
9112     cli # BIOS enabled interrupts; disable  
9113  
  
bootblock: bootasm.S bootmain.c  
           $(CC) $(CFLAGS) -fno-pic -O -nostdinc -I. -c bootmain.c  
           $(CC) $(CFLAGS) -fno-pic -nostdinc -I. -c bootasm.S  
           $(LD) $(LDFLAGS) -N -e start -Ttext 0x7C00 -o bootblock.o  
bootasm.o bootmain.o  
           $(OBJDUMP) -S bootblock.o > bootblock.asm  
           $(OBJCOPY) -S -O binary -j .text bootblock.o bootblock  
           ./sign.pl bootblock
```

# Switch to protected mode

- Switch from real to protected mode
  - Use a bootstrap GDT that makes virtual addresses map directly to physical addresses so that the effective memory map doesn't change during the transition.

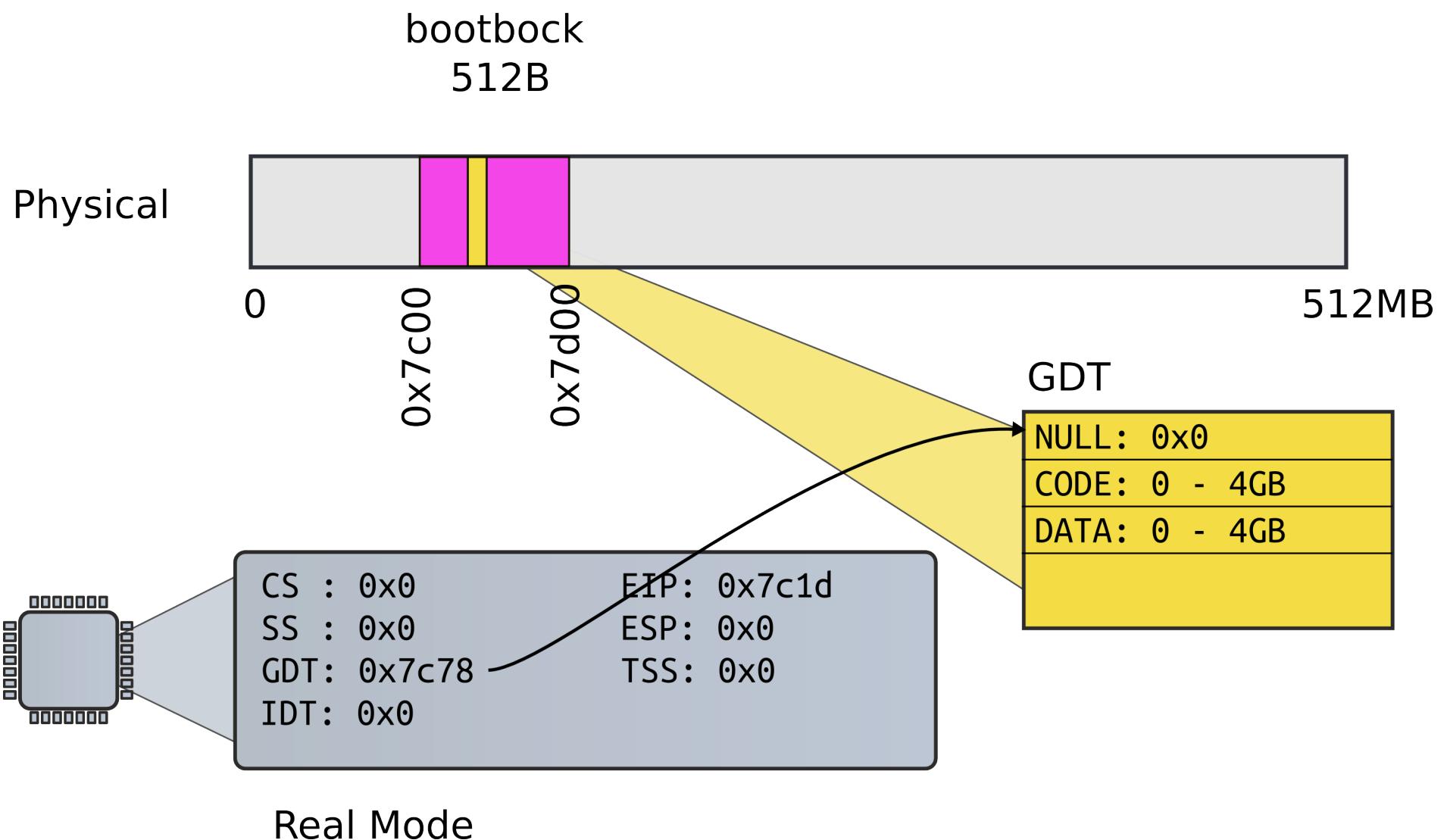
```
9141 lgdt gdtdesc
```

```
9142 movl %cr0, %eax
```

```
9143 orl $CR0_PE, %eax
```

```
9144 movl %eax, %cr0
```

# Load GDT



# How GDT is defined

```
9180 # Bootstrap GDT  
  
9181 .p2align 2 # force 4 byte alignment  
  
9182 gdt:  
  
9183     SEG_NULLASM # null seg  
  
9184     SEG_ASM(STA_X|STA_R, 0x0, 0xffffffff) # code seg  
  
9185     SEG_ASM(STA_W, 0x0, 0xffffffff) # data seg  
  
9186  
  
9187 gdtdesc:  
  
9188     .word (gdtdesc - gdt - 1) # sizeof(gdt) - 1  
  
9189     .long gdt
```

# How GDT is defined

```
9180 # Bootstrap GDT  
  
9181 .p2align 2 # force 4 byte alignment  
  
9182 gdt:  
  
9183     SEG_NULLASM # null seg  
  
9184     SEG_ASM(STA_X|STA_R, 0x0, 0xffffffff) # code seg  
  
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9186  
  
9187 gdtdesc:  
  
9188     .word (gdtdesc - gdt - 1) # sizeof(gdt) - 1  
  
9189     .long gdt
```

# Actual switch

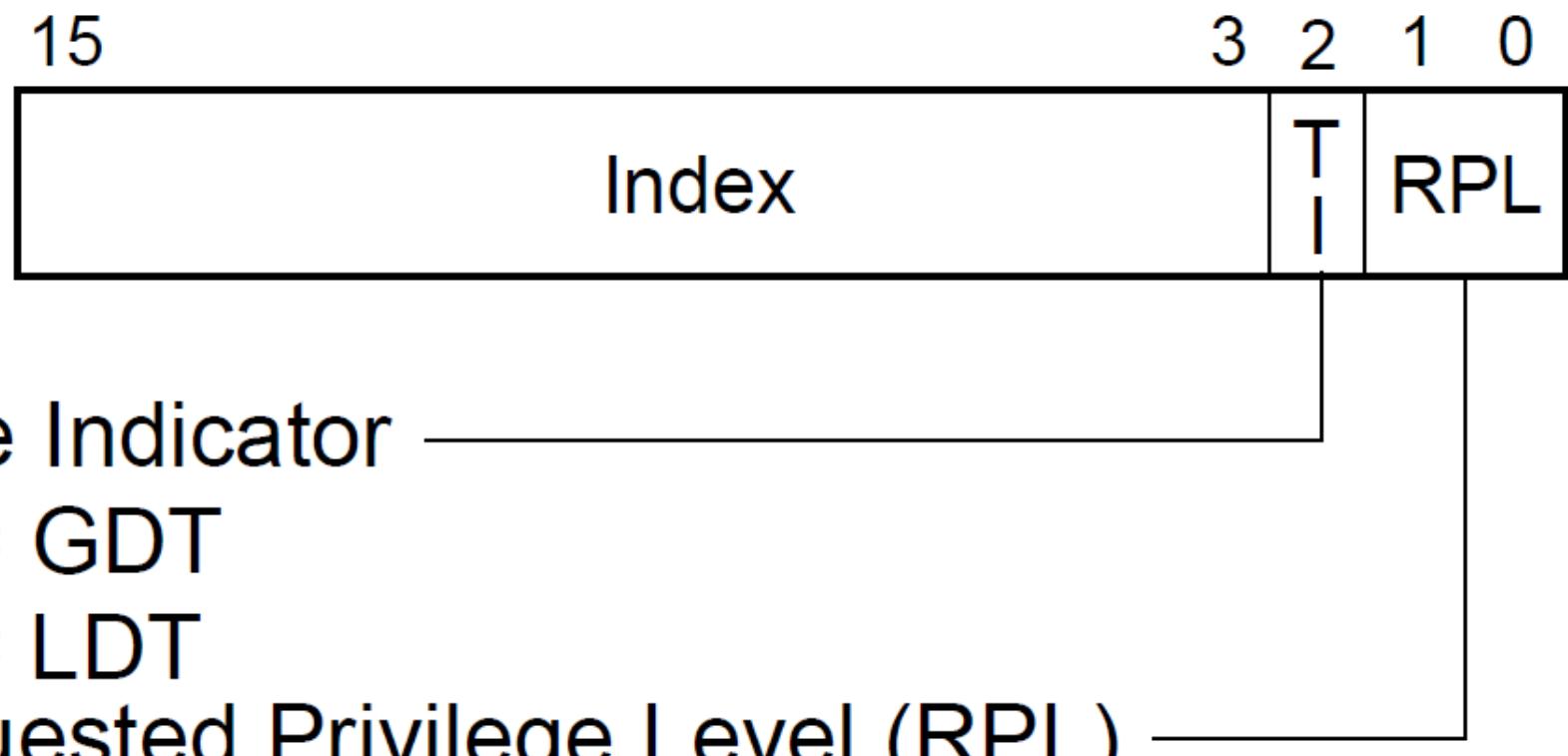
- Use long jump to change code segment

9153 1jmp \$(SEG\_KCODE<<3), \$start32

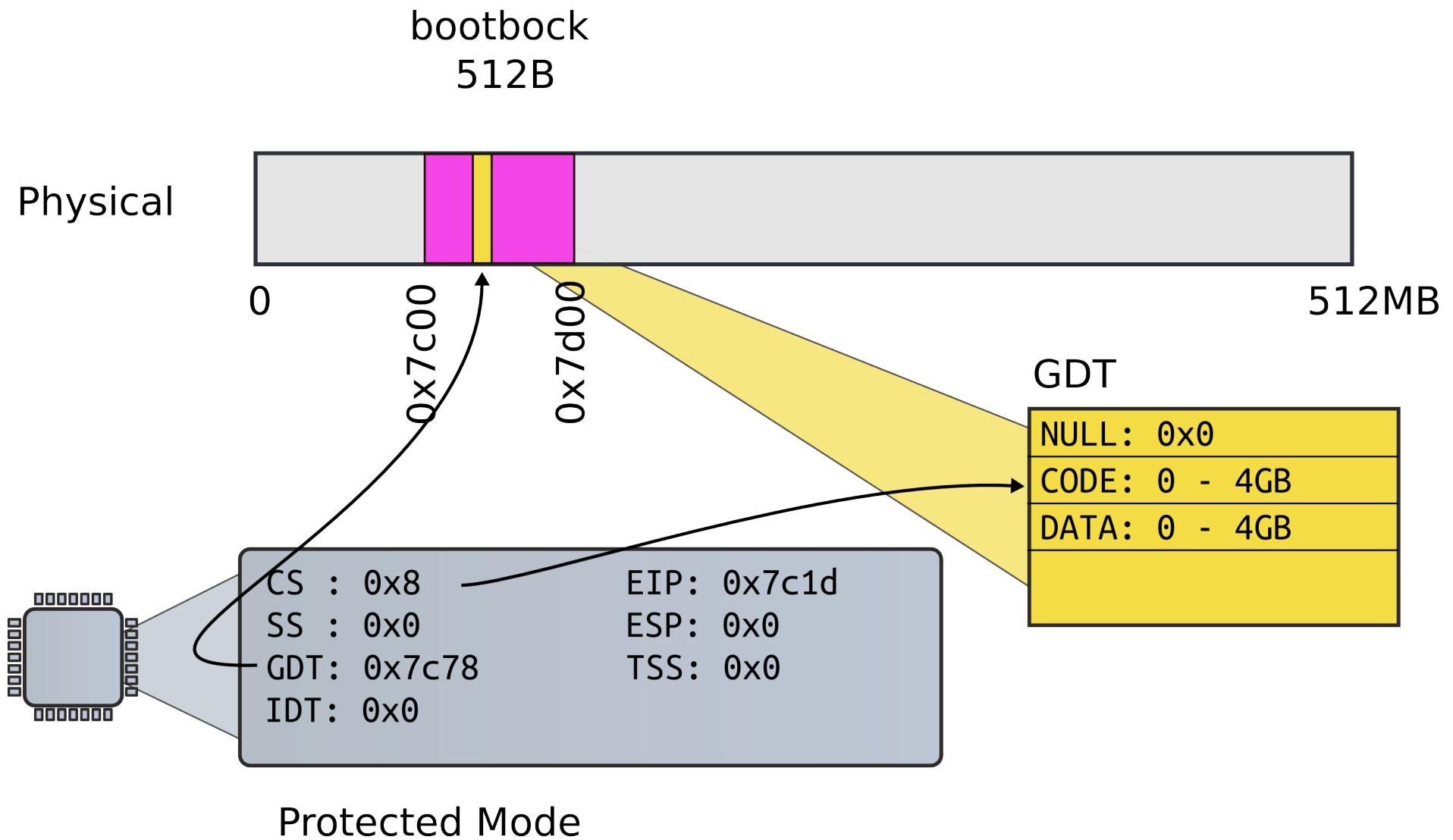
- Explicitly specify code segment, and address
- Segment is 0b1000 (0x8)

# Why CS is 0x8, not 0x1?

- Segment selector:



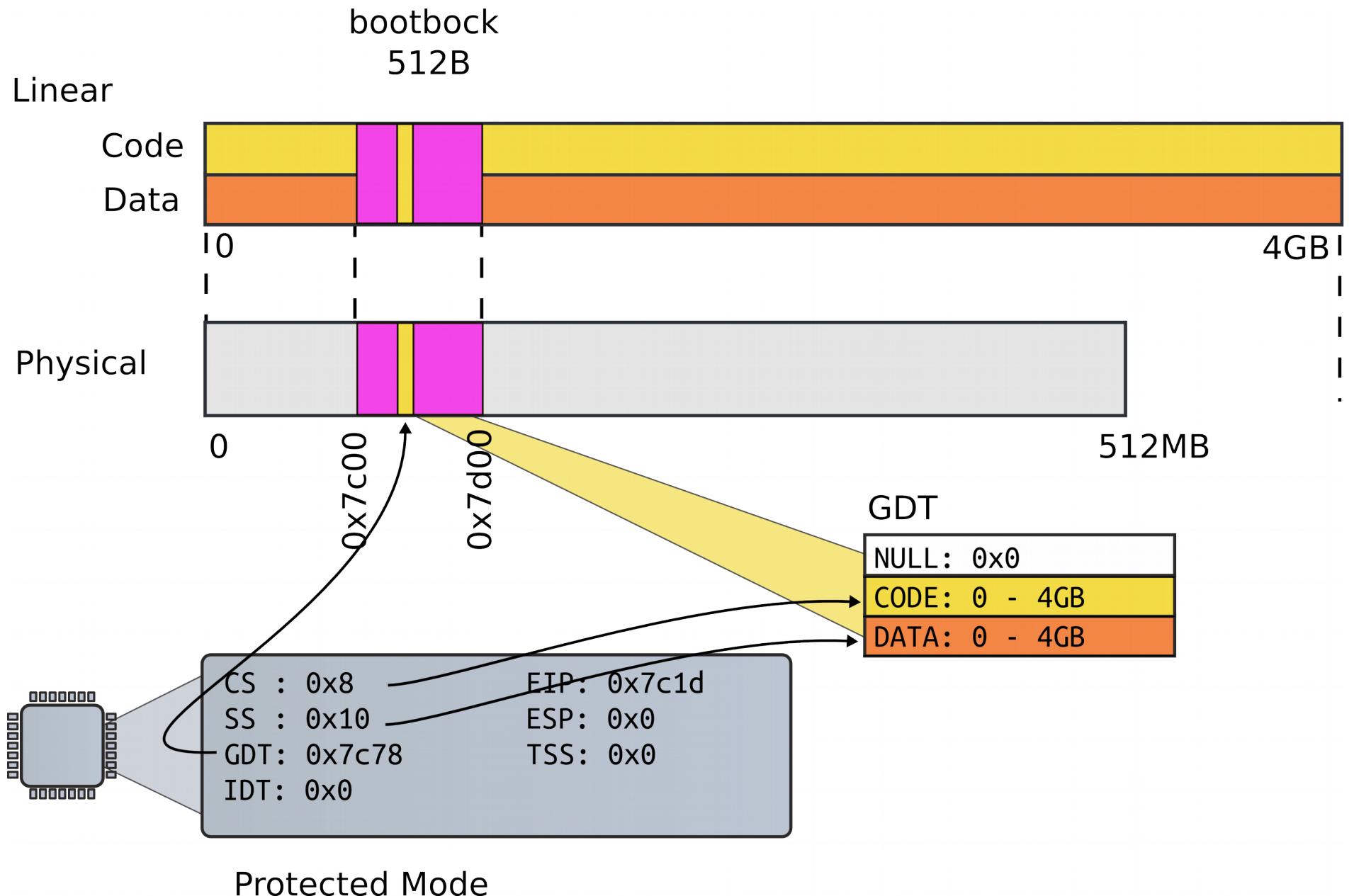
# Long jump



# Segments

```
9155 .code32 # Tell assembler to generate 32-bit code now.  
9156 start32:  
  
9157     # Set up the protected-mode data segment registers  
9158     movw $(SEG_KDATA<<3), %ax # Our data segment selector  
9159     movw %ax, %ds # -> DS: Data Segment  
9160     movw %ax, %es # -> ES: Extra Segment  
9161     movw %ax, %ss # -> SS: Stack Segment  
9162     movw $0, %ax # Zero segments not ready for use  
9163     movw %ax, %fs # -> FS  
9164     movw %ax, %gs # -> GS
```

# Segments



# Setup stack

- Why do we need a stack?

9166 movl \$start, %esp

9167 call bootmain

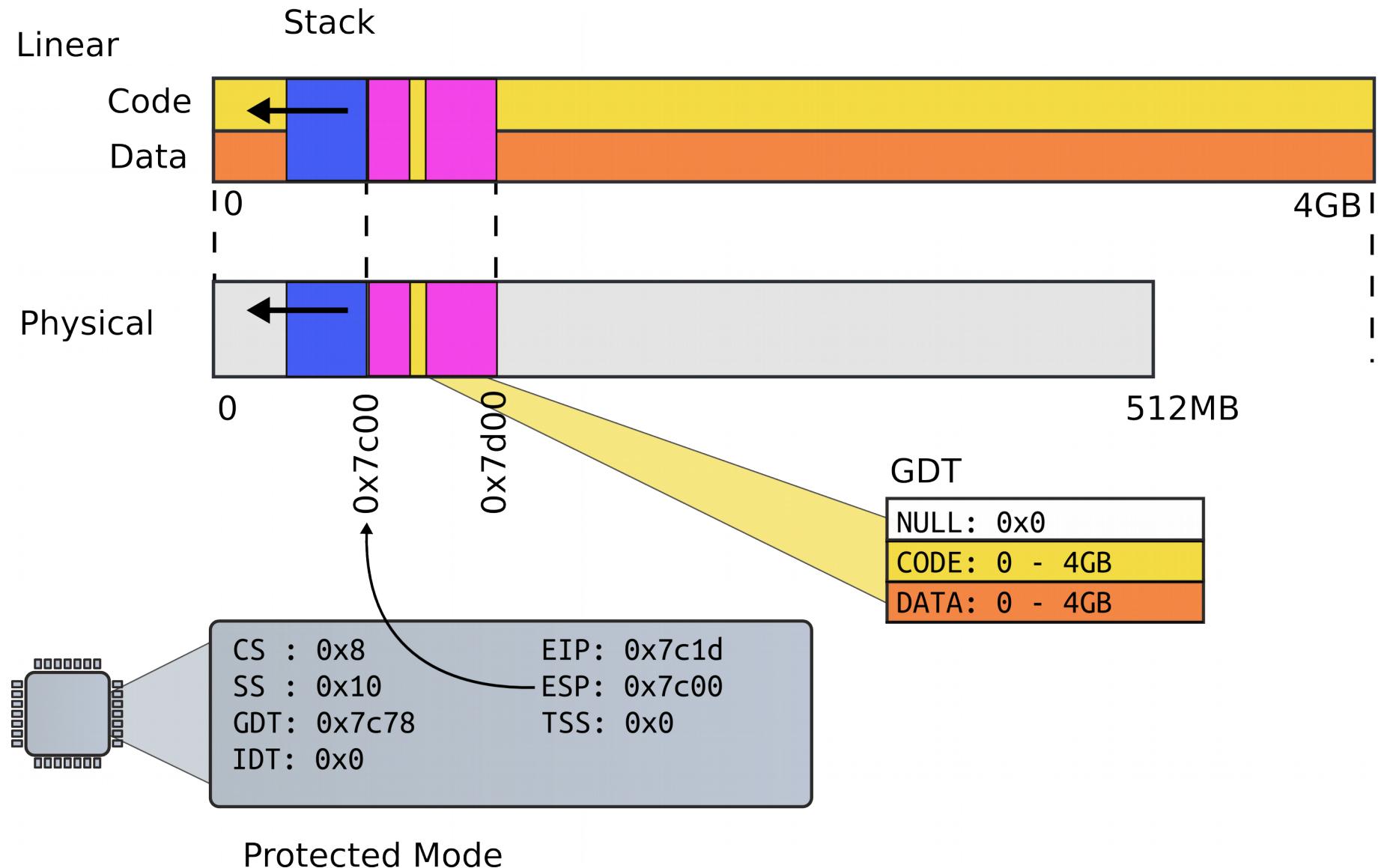
# Setup stack

- Need stack to use C
  - Function invocations
  - Note, there were no stack instructions before that

9166 movl \$start, %esp

9167 call bootmain

# First stack



# Invoke first C function

9166 movl \$start, %esp

9167 call bootmain

# bootmain(): read kernel from disk

```
9216 void
9217 bootmain(void)
9218 {
9219     struct elfhdr *elf;
9220     struct proghdr *ph, *eph;
9221     void (*entry)(void);
9222     uchar* pa;
9223
9224     elf = (struct elfhdr*)0x10000; // scratch space
9225
9226     // Read 1st page off disk
9227     readseg((uchar*)elf, 4096, 0);
9228
9229     // Is this an ELF executable?
9230     if(elf->magic != ELF_MAGIC)
9231         return; // let bootasm.S handle error
9232
```

```
9232
9233     // Load each program segment (ignores ph flags).
9234     ph = (struct proghdr*)((uchar*)elf + elf->phoff);
9235     eph = ph + elf->phnum;
9236     for(; ph < eph; ph++){
9237         pa = (uchar*)ph->paddr;
9238         readseg(pa, ph->filesz, ph->off);
9239         if(ph->memsz > ph->filesz)
9240             stosb(pa + ph->filesz, 0, ph->memsz - ph->filesz);
9241     }
9242
9243     // Call the entry point from the ELF header.
9244     // Does not return!
9245     entry = (void(*)(void))(elf->entry);
9246     entry();
9247 }
```

# bootmain(): read kernel from disk

# How do we read disk?

```
9257
9258 // Read a single sector at offset into dst.
9259 void
9260 readsect(void *dst, uint offset)
9261 {
9262     // Issue command.
9263     waitdisk();
9264     outb(0x1F2, 1); // count = 1
9265     outb(0x1F3, offset);
9266     outb(0x1F4, offset >> 8);
9267     outb(0x1F5, offset >> 16);
9268     outb(0x1F6, (offset >> 24) | 0xE0);
9269     outb(0x1F7, 0x20); // cmd 0x20 - read sectors
9270
9271     // Read data.
9272     waitdisk();
9273     insl(0x1F0, dst, SECTSIZE/4);
9274 }
```

xv6/bootmain.c

# How do we read disk (cont)?

```
9250 void  
9251 waitdisk(void)  
9252 {  
9253     // Wait for disk ready.  
9254     while((inb(0x1F7) & 0xC0) != 0x40)  
9255     ;  
9256 }  
9257
```

```
9232
9233     // Load each program segment (ignores ph flags).
9234     ph = (struct proghdr*)((uchar*)elf + elf->phoff);
9235     eph = ph + elf->phnum;
9236     for(; ph < eph; ph++){
9237         pa = (uchar*)ph->paddr;
9238         readseg(pa, ph->filesz, ph->off);
9239         if(ph->memsz > ph->filesz)
9240             stosb(pa + ph->filesz, 0, ph->memsz - ph->filesz);
9241     }
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9243     // Call the entry point from the ELF header.
9244     // Does not return!
9245     entry = (void(*)(void))(elf->entry);
9246     entry();
9247 }
```

Call kernel entry

```
1. #include <stdio.h>
2.
3. void func_a(void){
4.     printf("func_a\n");
5.     return;
6. }
7.
8. void func_b(void) {
9.     printf("func_b\n");
10.    return;
11. }
12.
13. int main(int ac, char **av)
14. {
15.     void (*fp)(void);
16.
17.     fp = func_b;
18.     fp();
19.     return;
20. }
```

# Function pointers

08048432 <func\_b>:

8048432:	55	push %ebp
8048433:	89 e5	mov %esp,%ebp
8048435:	83 ec 18	sub \$0x18,%esp
8048438:	c7 04 24 07 85 04 08	movl \$0x8048507,(%esp)
804843f:	e8 ac fe ff ff	call 80482f0 <puts@plt>
8048444:	90	nop
8048445:	c9	leave
8048446:	c3	ret

08048447 <main>:

8048447:	55	push %ebp
8048448:	89 e5	mov %esp,%ebp
804844a:	83 e4 f0	and \$0xfffffffff0,%esp
804844d:	83 ec 10	sub \$0x10,%esp
		# Load pointer to func_p on the stack
8048450:	c7 44 24 0c 32 84 04	movl \$0x8048432,0xc(%esp)
8048457:	08	
8048458:	8b 44 24 0c	mov 0xc(%esp),%eax
804845c:	ff d0	call *%eax
804845e:	90	nop
804845f:	c9	leave
8048460:	c3	ret

# Function pointers

# Function pointers

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8048433:	89 e5	mov %esp,%ebp
8048435:	83 ec 18	sub \$0x18,%esp
8048438:	c7 04 24 07 85 04 08	movl \$0x8048507,(%esp)
804843f:	e8 ac fe ff ff	call 80482f0 <puts@plt>
8048444:	90	nop
8048445:	c9	leave
8048446:	c3	ret

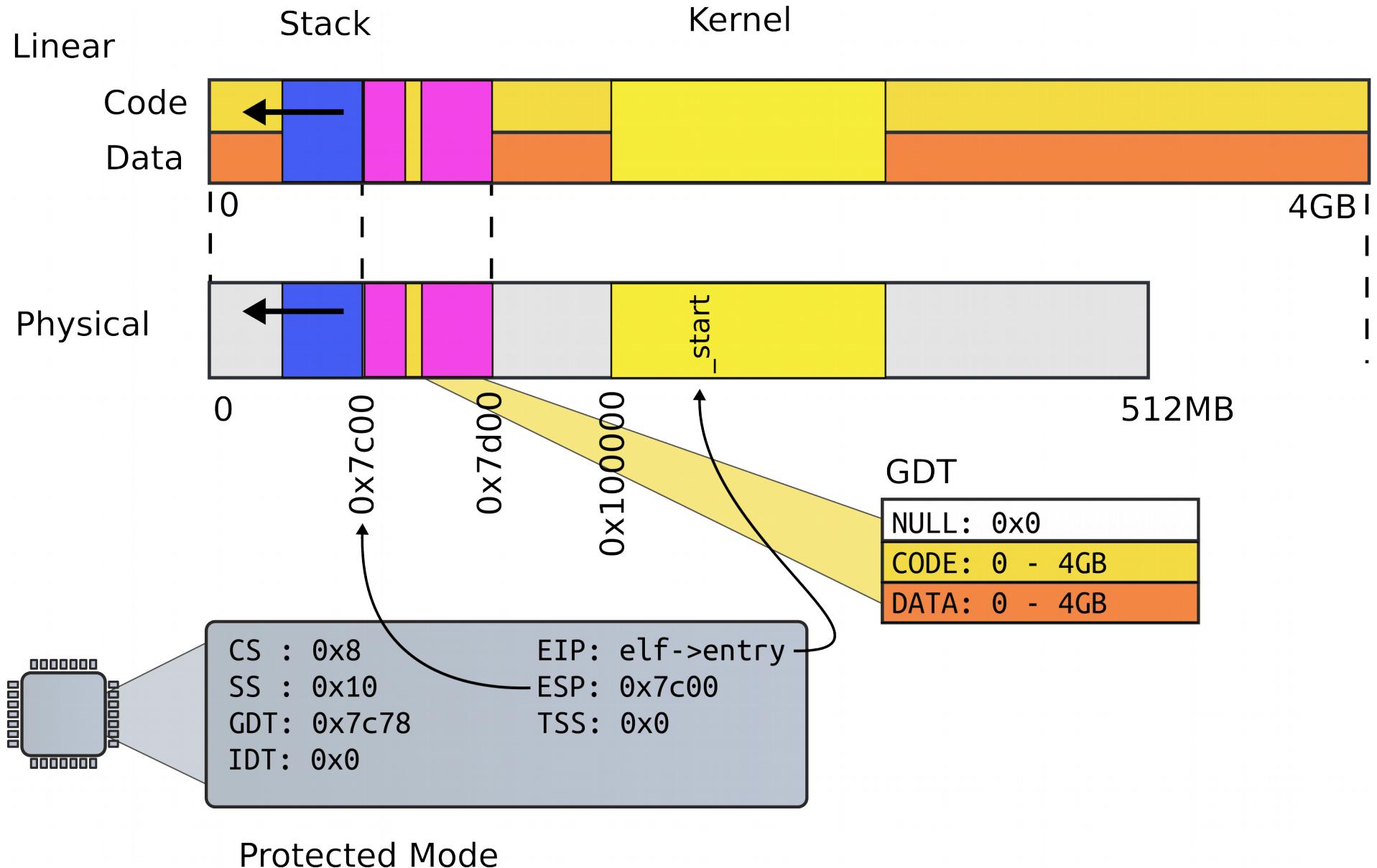
08048447 <main>:

8048447:	55	push %ebp
8048448:	89 e5	mov %esp,%ebp
804844a:	83 e4 f0	and \$0xffffffff0,%esp
804844d:	83 ec 10	sub \$0x10,%esp
		# Load pointer to func_p on the stack
8048450:	c7 44 24 0c 32 84 04	movl \$0x8048432,0xc(%esp)
8048457:	08	# Move func_b into %eax
8048458:	8b 44 24 0c	mov 0xc(%esp),%eax
804845c:	ff d0	call *%eax # Call %eax
804845e:	90	nop
804845f:	c9	leave
8048460:	c3	ret

```
1039 .globl entry
1136 # By convention, the _start symbol specifies the ELF entry point.
1137 # Since we haven't set up virtual memory yet, our entry point is
1138 # the physical address of 'entry'.
1139 .globl _start
1140 _start = V2P_W0(entry)
1141
1142 # Entering xv6 on boot processor, with paging off.
1143 .globl entry
1144 entry:
1145 # Turn on page size extension for 4Mbyte pages
1146     movl %cr4, %eax
1147     orl $(CR4_PSE), %eax
1148     movl %eax, %cr4
```

# entry(): kernel ELF entry

# Kernel



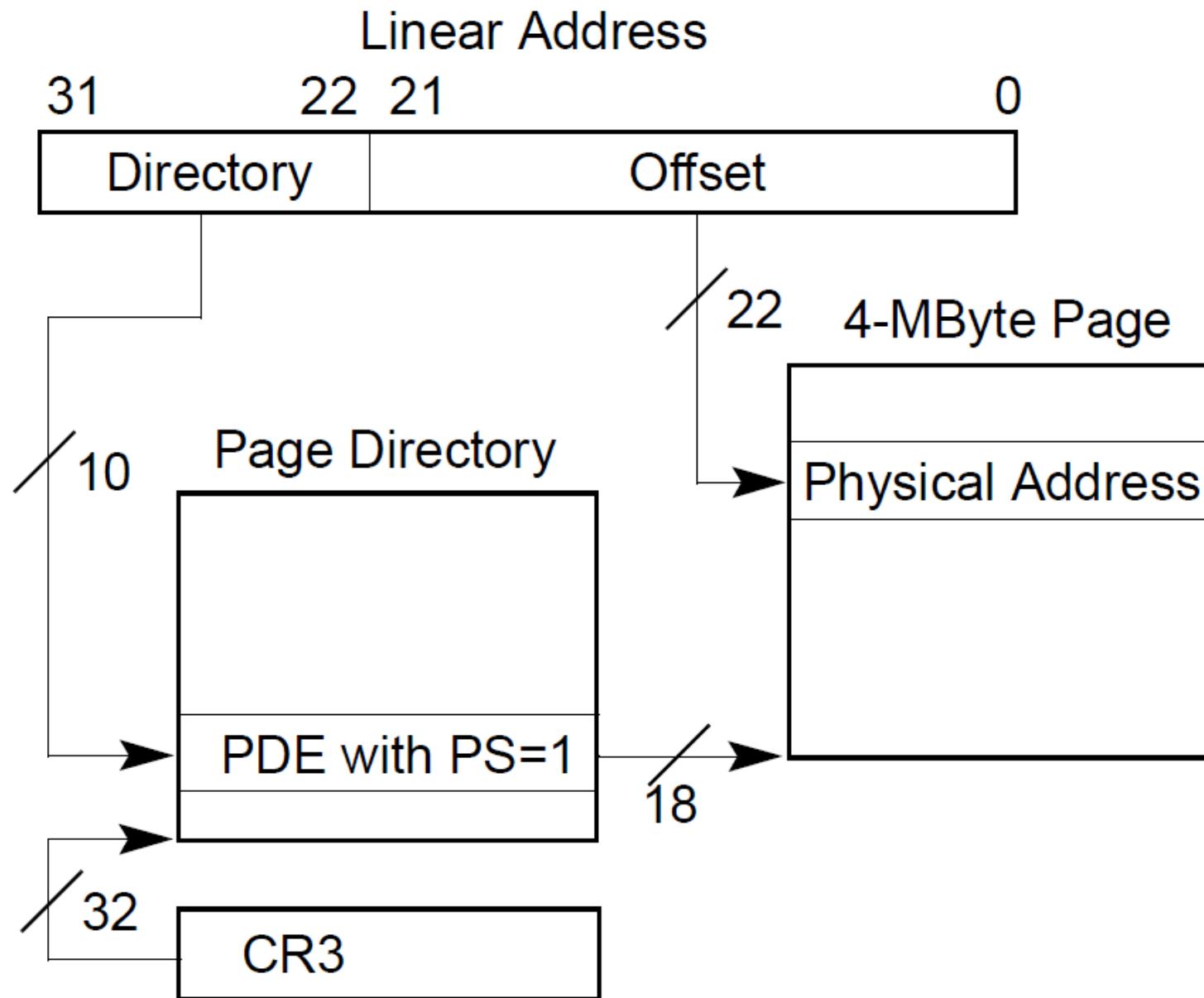
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xv6/entry.S

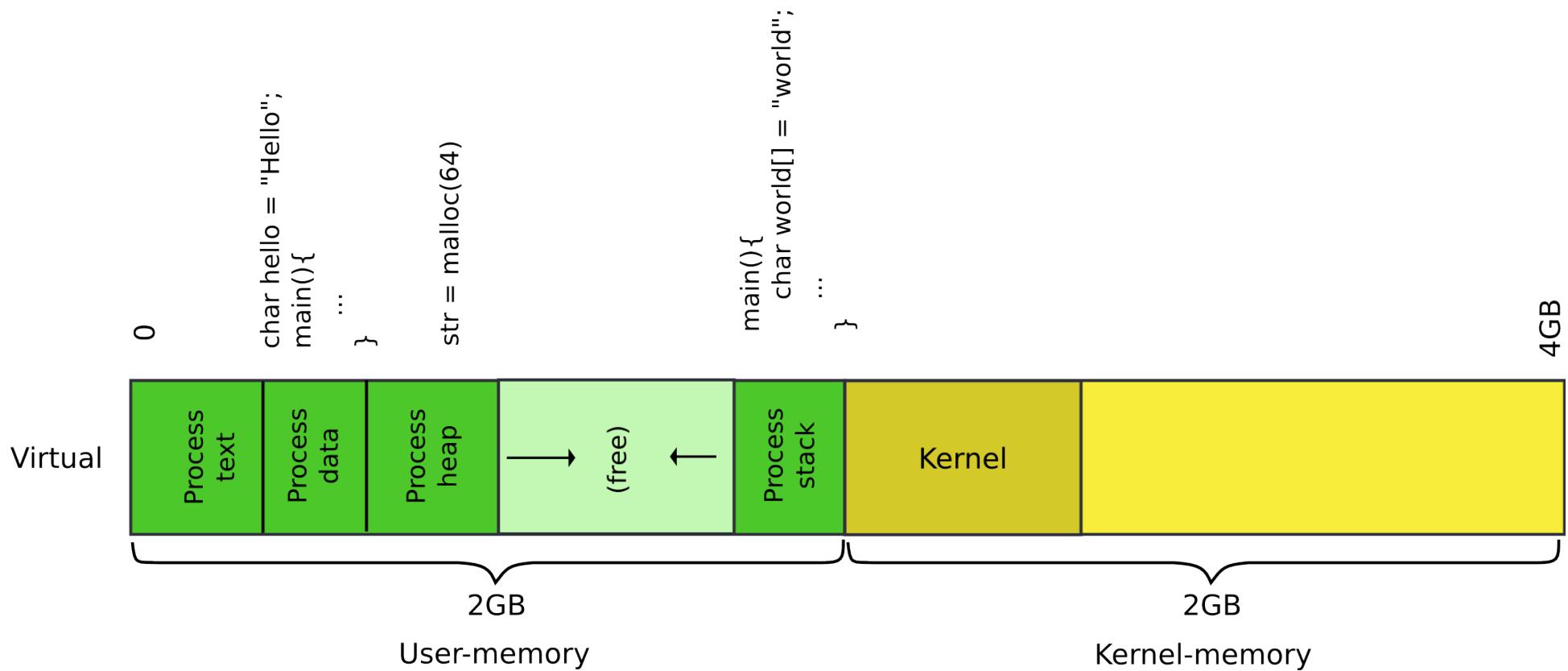
entry(): kernel ELF entry

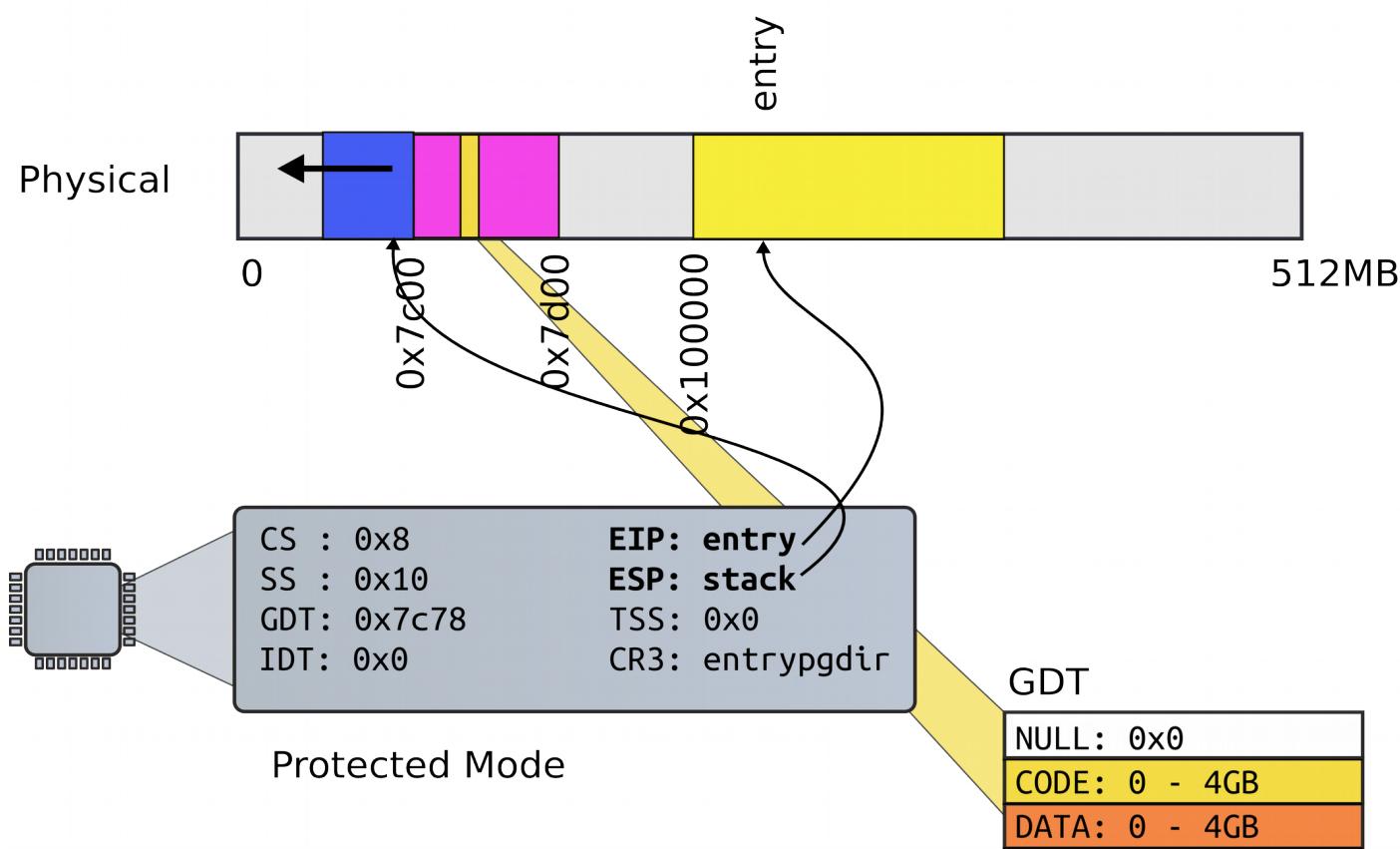
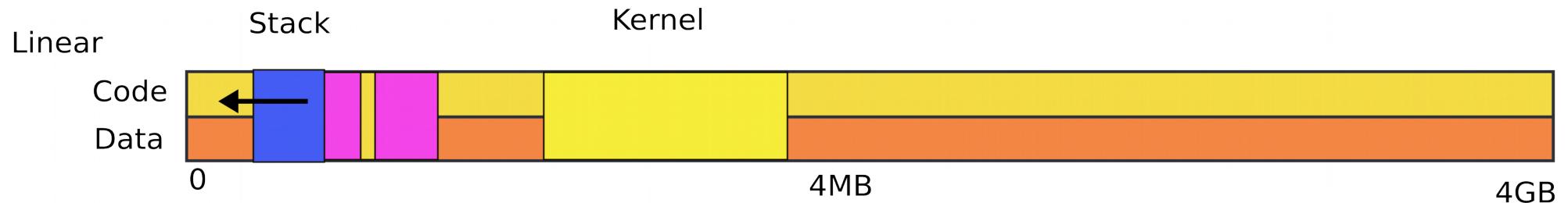
# Set up page directory

```
1149 # Set page directory  
1150 movl $(V2P_W0(entrypgdir)), %eax  
1151 movl %eax, %cr3
```



# Our goal: 2GB/2GB address space





# First page table

- Two 4MB entries (large pages)
- Entry #0
  - $0x0 - 4MB \rightarrow 0x0:0x400000$
- Entry #512
  - $0x0 - 4MB \rightarrow 0x8000000:0x80400000$

```
1406 // The boot page table used in entry.S and entryother.S.  
1407 // Page directories (and page tables) must start on page  
     boundaries,  
1408 // hence the __aligned__ attribute.  
1409 // PTE_PS in a page directory entry enables 4Mbyte  
pages.  
1410  
1411 __attribute__((__aligned__(PGSIZE)))  
1412 pde_t entrypgdir[NPDENTRIES] = {  
1413     // Map VA's [0, 4MB) to PA's [0, 4MB)  
1414     [0] = (0) | PTE_P | PTE_W | PTE_PS,  
1415     // Map VA's [KERNBASE, KERNBASE+4MB) to PA's [0, 4MB)  
1416     [KERNBASE>>PDXSHIFT] = (0) | PTE_P | PTE_W | PTE_PS,  
1417 };
```

## First page table

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```

## First page table

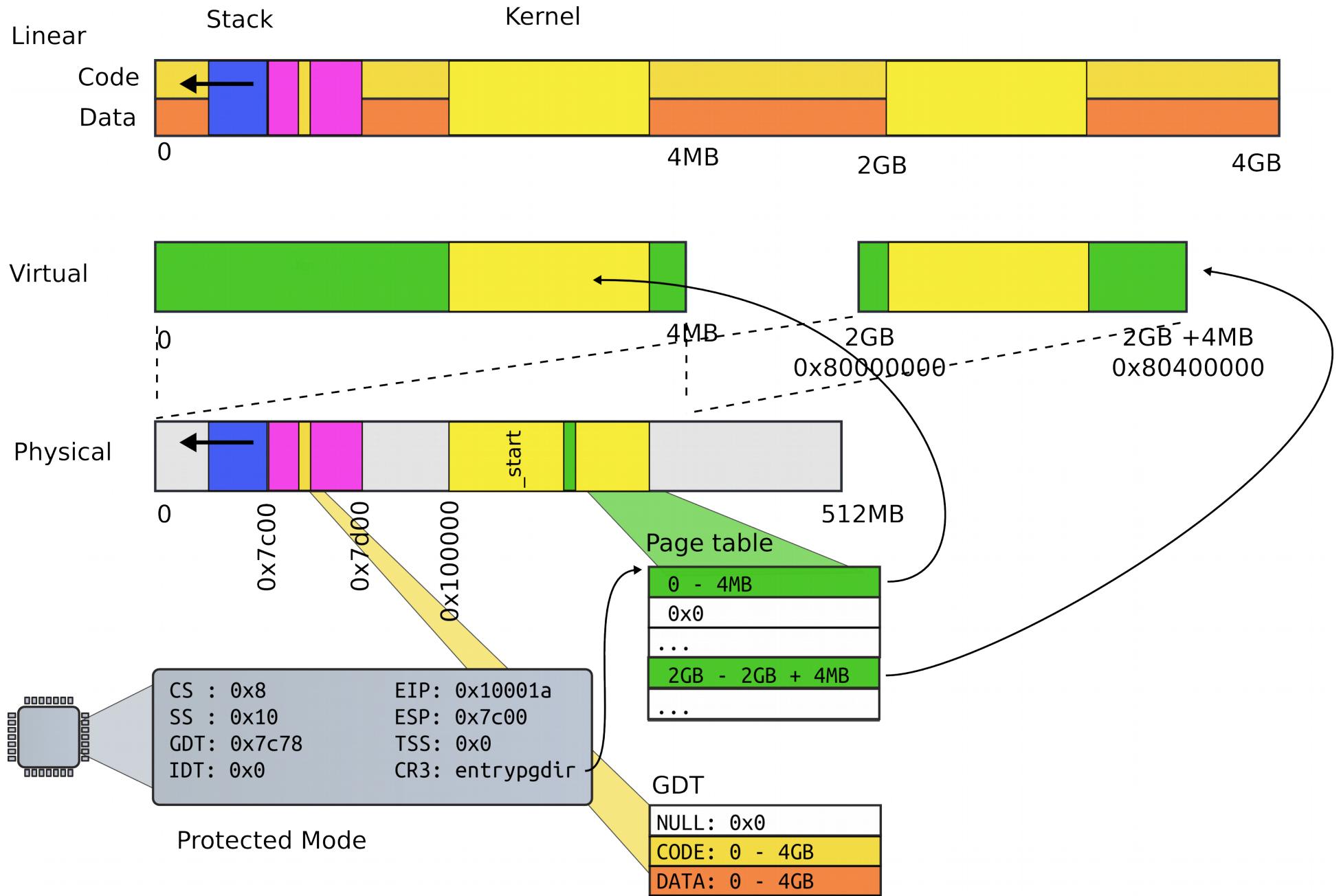
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1414     [0] = (0) | PTE_P | PTE_W | PTE_PS,  
1415     // Map VA's [KERNBASE, KERNBASE+4MB) to PA's [0, 4MB)  
1416     [KERNBASE>>PDXSHIFT] = (0) | PTE_P | PTE_W | PTE_PS,  
1417 };
```

## First page table

# First page table (cont)

```
0870 // Page directory and page table constants.  
0871 #define NPDENTRIES 1024
```

# First page table



# Turn on paging

```
1152 # Turn on paging.
```

```
1153 movl %cr0, %eax
```

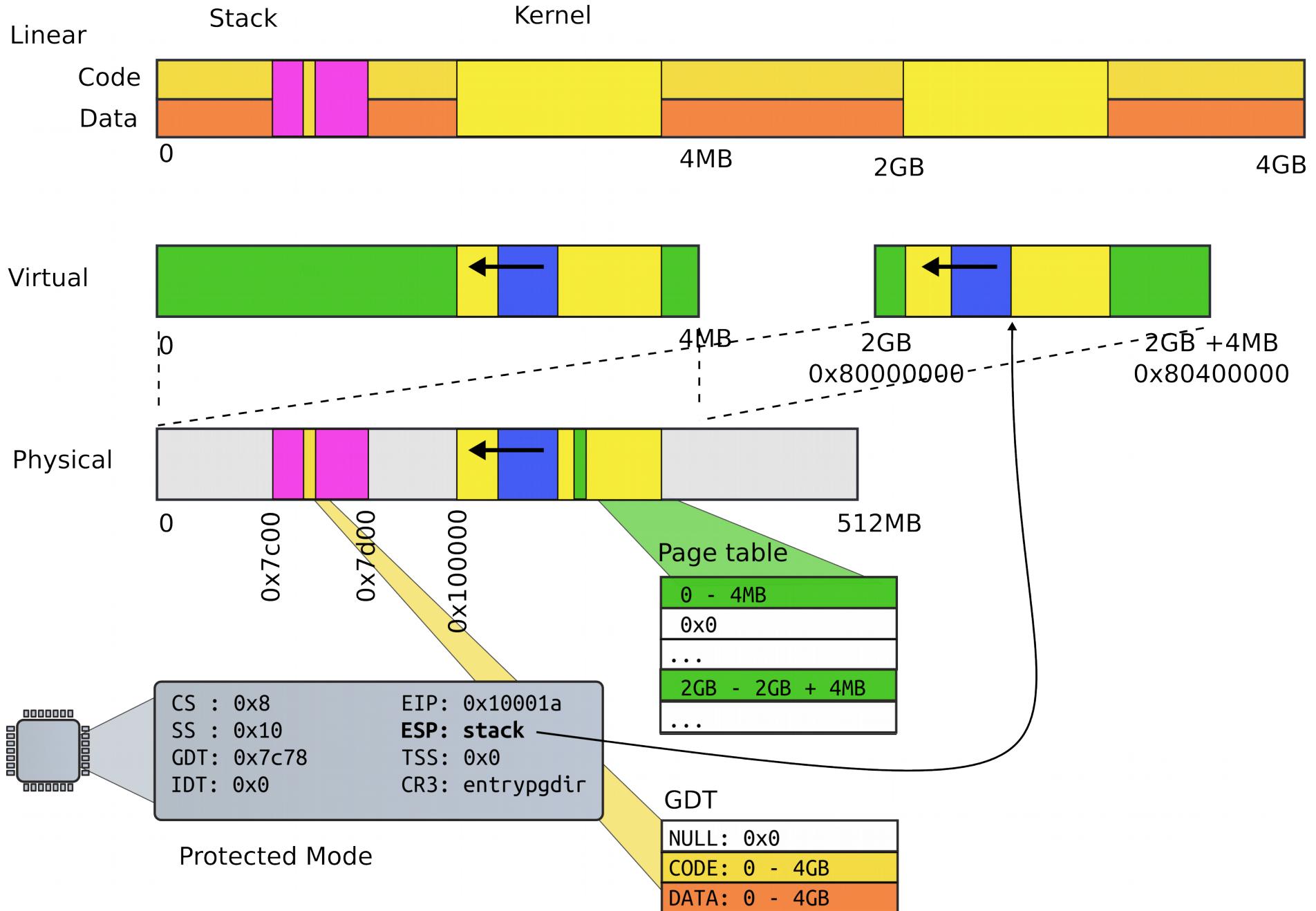
```
1154 orl $(CRO_PG|CRO_WP), %eax
```

```
1155 movl %eax, %cr0
```

# High address stack (4K)

```
1157 # Set up the stack pointer.  
1158 movl $(stack + KSTACKSIZE), %esp  
1159  
...  
1167 .comm stack, KSTACKSIZE  
  
0151 #define KSTACKSIZE 4096 // size of  
           per-process kernel stack
```

# High address stack (4K)



# Jump to main()

```
1160 # Jump to main(), and switch to executing at  
1161 # high addresses. The indirect call is  
     needed because  
1162 # the assembler produces a PC-relative  
     instruction  
1163 # for a direct jump.  
1164 mov $main, %eax  
1165 jmp *%eax  
1166
```

# Running in main()

```
1313 // Bootstrap processor starts running C code here.  
1314 // Allocate a real stack and switch to it, first  
1315 // doing some setup required for memory allocator to work.  
1316 int  
1317 main(void)  
1318 {  
1319     kinit1(end, P2V(4*1024*1024)); // phys page allocator  
1320     kvmalloc(); // kernel page table  
1321     mpinit(); // detect other processors  
1322     lapicinit(); // interrupt controller  
1323     seginit(); // segment descriptors  
1324     cprintf("\ncpu%d: starting xv6\n\n", cpunum());  
...  
1340 }
```

# Recap of the boot sequence

- Setup segments (data and code)
- Switched to protected mode
  - Loaded GDT (segmentation is on)
- Setup stack (to call C functions)
- Loaded kernel from disk
- Setup first page table
  - 2 entries [ 0 : 4MB ] and [ 2GB : (2GB + 4MB) ]
- Setup high-address stack
- Jumped to main()

# Conclusion

- We've booted
  - We're running in main()
- Next time:
  - Process and kernel address space

Thank you!

# References

- [1] Costan, Victor, and Srinivas Devadas. "Intel SGX Explained." IACR Cryptology ePrint Archive 2016 (2016): 86.  
<https://eprint.iacr.org/2016/086.pdf>