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Template hashing

```
| md5sum | cut -c-6
```

The cli command md5sum is used to validate if the template was written cor-

Write the one liner above without line breaks and save it as hash.sh, and run

chmod +x hash.sh

to make it executable, vai luan

To get the hash in file.cpp between lines <11> and <12>, run

Hash is made by scope, delimited by the curly brackets ("{", "}").

The empty hash is d41d8 while the m5dsum of hash.sh is 9cd12.

Matrix Inverse 2x2

$$A = \begin{pmatrix} a & b \\ c & d \end{pmatrix}, A^{-1} = \frac{1}{ad-bc} \begin{pmatrix} d & -b \\ -c & a \end{pmatrix}$$

Pick Theorem

 $A = i + \frac{b}{2} - 1$, where i is the number of lattice points inside the polygon and b is the number on boundary. Work for simple polygon.

Burnside Lemma

Let G be a finite group that acts on set X. Let $X^g := \{x \in X \mid g.x = x\}$. The number of orbits |X/G| is

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

Lucas Theorem

For prime p,
$$\binom{n}{m} = \prod_{i=0}^{n} \binom{n_i}{m_i} \pmod{p}$$

Catalan

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \frac{(2n)!}{n!*(n+1)!} \mid C_n^k = \frac{k+1}{n+k+1} \binom{2n+k}{n}$$

Stirling numbers of the first kind

These are the number of permutation of size n with exactly k cycles

Stirling numbers of the second kind

These are the number of ways to partition n into exactly k non-empty sets.

$$\begin{bmatrix} n \\ k \end{bmatrix} = (k) \begin{bmatrix} n-1 \\ k \end{bmatrix} + \begin{bmatrix} n-1 \\ k-1 \end{bmatrix}$$

$$\begin{bmatrix} n \\ k \end{bmatrix} = \frac{1}{k!} \sum_{j=0}^{k} (-1)^{k-j} {k \choose j} j^n$$

$$\begin{bmatrix} 0 \\ 0 \end{bmatrix} = 1, \begin{bmatrix} n \\ 0 \end{bmatrix} = \begin{bmatrix} 0 \\ n \end{bmatrix} = 1$$

Planar Graph

If G has k connected components, then n - m + f = k + 1.

m < 3n - 6. If G has no triangles, m < 2n - 4.

The minimum degree is less of equal 5. And can be 6 colored in O(n+m)/

Spherical coordinates

$$\begin{split} x &= r \sin(\theta) \cos(\varphi) \mid r = \sqrt{x^2 + y^2 + z} \\ y &= r \sin(\theta) \sin(\varphi) \mid (\theta) = \arccos\left(\frac{z}{r}\right) \\ z &= r \cos(\theta) \mid (\varphi) = \operatorname{atan2}\left(y, x\right) \end{split}$$

series and sums

$$\begin{split} &\frac{1}{1-x} = \sum_{k=0} x^k \mid e^x = \sum_{k=0} \frac{x^k}{k!} \mid \ln(1+x) = \sum_{k=1} \left(-1\right)^{k+1} \frac{x^k}{k!} \\ &\sqrt{1+x} = 1 + \sum_{k=1} \left(-1\right)^{k+1} \frac{x^k}{2^{2k-1}} \mid \frac{1}{\sqrt{1-x}} = \sum_{k=0} \frac{x^k}{4^k} \binom{2k}{k} \\ &\sin x = \sum_{k=0} \left(-1\right)^k \frac{(x)^{2k+1}}{(2k+1)!} \mid \cos x = \sum_{k=0} \left(-1\right)^k \frac{x^{2k}}{(2k)!} \\ &1+2+\ldots = \frac{n(n+1)}{2} \mid 1^2+2^2+\ldots = \frac{n(2n+1)(n+1)}{6} \\ &1^3+2^3+\ldots = \frac{n^2(n+1)^2}{4} \mid 1^4+2^4+\ldots = \frac{n(n+1)(2n+1)(3n^2+3n+1)}{30} \end{split}$$

Code

divide and conquer

```
d41 // Divide and Conquer DP
d41 //
d41 // A dp of the form
d41 // dp[i][j] = min_{k < j}(dp[i - 1][k] + cost(k, j))
d41 // can be solved in O(m n logn) with divide and conquer
    optimization if we have that
          opt[i][j] <= opt[i][j + 1]
d41 // where
         dp[i][j] = dp[i - 1][opt[i][j]] + cost(opt[i][j], j).
d41 // Complexity: O(m n logn) time (for a partition in m
   subarrays of an array of size n)
                  O(n) memory
d41 //
d41
d41
547 ll dp[MAX][2];
94b void solve(int k, int l, int r, int lk, int rk) {
```

```
if (l > r) return;
109
        int m = (l+r)/2, p = -1;
d2b
        auto& ans = dp[m][k\&1] = LINF;
6e2
        for (int i = max(m, lk); i \le rk; i++) {
d73
            ll at = dp[i+1][\sim k\&1] + cost(m, i);
57d
            if (at < ans) ans = at, p = i;
5ed
1ee
        solve(k, l, m-1, lk, p), solve(k, m+1, r, p, rk);
b62 }
d41
cf1 ll DC(int n, int k) {
        dp[n][0] = dp[n][1] = 0;
f27
        for (int i = 0; i < n; i++) dp[i][0] = LINF;
        for (int i = 1; i \le k; i++) solve(i, 0, n-i, 0, n-i);
b76
8e7
        return dp[0][k&1];
5e9 }
```

1

dynamic cht

d41 // Dynamic Convex Hull Trick

```
d41 //
d41 // Description:
            Maintains the convex hull of some functions.
           Copied from github.com/brunomaletta/Biblioteca/blob/
    master/Codigo/Estruturas/chtDinamico.cpp.
d41 //
d41 // Functions:
d41 //
            add(a,b); adds line (ax+b) to the convex hull.
            query(x): returns the maximum value of any line on
   point x.
d41 //
d41 // Complexity:
d41 //
                    O(logn)
d41 //
            query: O(logn)
d41 //
d41 // Details:
           If you want to maintain the bottom convex hull, it is
            easier to just change the sign. Be careful with
d41 //
            on guery. Can use int128 to avoid.
72c struct Line {
        mutable ll a, b, p;
        bool operator<(const Line& o) const { return a < o.a; }</pre>
8e3
abf
        bool operator<(ll x) const { return p < x; }</pre>
469 };
326 struct dynamic hull : multiset<Line, less<>>> {
33a
       ll div(ll a, ll b) {
a20
            return a / b - ((a ^ b) < 0 and a % b);
a8a
       }
d41
bbb
        void update(iterator x) {
459
            if (next(x) == end()) x -> p = INF;
            else if (x->a == next(x)->a) x->p = x->b >= next(x)-
eec >b ? INF : -INF;
424
          else x->p = div(next(x)->b - x->b, x->a - next(x)->a);
d37
d41
71c
        bool overlap(iterator x) {
```

```
f18
            update(x);
cfa
            if (next(x) == end()) return 0;
a4a
            if (x->a == next(x)->a) return x->b >= next(x)->b;
d40
            return x \rightarrow p >= next(x) \rightarrow p;
901
       }
d41
176
        void add(ll a, ll b) {
1c7
            auto x = insert({a, b, 0});
4ab
            while (overlap(x)) erase(next(x)), update(x);
           if (x != begin() and !overlap(prev(x))) x = prev(x),
dbc
    update(x);
0fc
            while (x != begin() and overlap(prev(x)))
4d2
                x = prev(x), erase(next(x)), update(x);
48f
       }
d41
4ad
        ll query(ll x) {
229
            assert(!emptv()):
7d1
            auto l = *lower bound(x):
d41
            // if(l.a and abs(x) >= abs(INF/l.a)) return INF/2;
aba
            return l.a * x + l.b:
3f5
      }
905 }:
```

knuth

```
d41 // Knuth DP
d41
d41 // A dp of the form
d41 // dp[l][r] = min_{l} < m < r_{l} (dp[l][m] + dp[m][r]) +
_{
m d41} // can be solved in O(n^2) with Knuth opmitization if we
    have that
d41 // opt[l][r - 1] <= opt[l][r] <= opt[l + 1][r]</pre>
d41 // where
d41 // dp[l][r] = dp[l][opt[l][r]] + dp[opt[l][r]][r] +
    cost(l, r).
d41 // Other sufficient condition (that implies the previous
d41 // given a <= b <= c <= d, we have:
    // - quadrangle inequality: cost(a, c) + cost(b, d)
    \leq cost(a, d) + cost(b, c)
d41 //
            - monotonicity: cost(b, c) <= cost(a, d)</pre>
d41 //
d41 // Complexity: O(n^2) time
d41 //
                  0(n^2) memory
297 ll knuth(int n) {
       vector<vector<ll>>> dp = vector<vector<ll>>>(n,
    vector<ll>(n));
       vector<vector<int>> opt = vector<vector<int>>>(n,
    vector<int>(n));
83e
       for(int k = 0; k \le n; k++) {
831
            for(int l = 0; l + k \le n; l++) {
cdd
               int r = l + k;
2b2
               if(k < 2) {
653
                    dp[l][r] = 0; // base case
358
                    opt[l][r] = l;
2ad
740
                dp[l][r] = INF;
```

```
for(int m = opt[l][r - 1]; m <= opt[l + 1][r];</pre>
cba <sub>m++</sub>) {
                   ll cur = dp[l][m] + dp[m][r] + cost(l, r); //
elb must define 0(1) cost function
                     if(cur < dp[l][r]) {
3b8
                         dp[l][r] = cur;
f17
                         opt[l][r] = m;
9e4
                     }
5ee
                 }
22f
            }
39e
ec9
         return dp[0][n];
124 }
```

extra

fastIO

```
d41 // fast io
d41
7a5 int read int() {
32h
        bool minus = false:
c18
       int result = 0;
d28
       char ch:
ca3
        ch = getchar();
31e
       while (1) {
b92
           if (ch == '-') break;
736
           if (ch >= '0' && ch <= '9') break;
ca3
            ch = getchar();
771
       }
822
        if (ch == '-') minus = true;
62e
        else result = ch-'0';
31e
       while (1) {
ca3
            ch = getchar();
a46
            if (ch < '0' || ch > '9') break;
f79
            result = result*10 + (ch - '0');
722
28e
        if (minus) return -result;
98e
        else return result;
9b5 }
```

pragmas

```
d41 // Pragmas
881 #pragma GCC optimize("03,unroll-loops")
827 #pragma GCC target("avx2,bmi,bmi2,lzcnt,popcnt")
```

random

```
d41 // Random
798 mt19937_64 rng(chrono::steady_clock::now().time_since_epoch().
836 shuffle(permutation.begin(), permutation.end(), rng);
lee uniform_int_distribution<int>(a,b)(rng);
```

template

```
d41 // template
d41
2b7 #include <bits/stdc++.h>
d41
```

```
ca4 using namespace std;
f51 #define fastio ios base::sync with stdio(0); cin.tie(0)
90e using ll = long long;
0e0 #define int ll
efe #define pb push_back
d69 #define all(a) a.begin(), a.end()
21e void dbg out() { cerr << endl; }</pre>
2dc template <typename H, typename... T>
f62 void dbg out(H h, T... t) { cerr << ' ' << h; dbg out(t...); }
   #define dbg(...) { cerr << # VA ARGS << ':';</pre>
89a #@eiille dbg(..., dbg_out(__VA_ARGS__); }
d41
63d void solve(){
fa9 }
0dd signed main(){
         fastio:
0f9
        solve();
78d }
```

geometry

circle

```
d41 // Circle
d41 //
d41 // Complexity: O(1) unless stated
d39 const Point<int> ccw(1,0);
4fc template <class T>
d2f struct Circle {
       using P = Point<T>;
20e
       using C = Circle;
d41
543
        Pc; Tr;
ca1
        Circle() {}
439
        Circle(P c, T r) : c(c), r(r) {}
        bool in(C rhs) { return (c-rhs.c).norm2() <= (rhs.r-</pre>
d4d
   ' r)*(rhs.r-r): }
        bool has(P p) { return (c-p).norm2() \le r*r; }
296
d41
d41
        // returns intersection points between line and circle
a07
        vector<P> intersect(Line<T> l) {
oAifOt());
            T h2 = r * r - l.dist2(c):
95d
            if (h2 < -EPS) return {};</pre>
d41
dc8
            P p = l.proj(c);
b23
            P h = l.v * sqrt(h2) / l.v.norm();
0c8
            if (h.norm() < EPS) return {p};</pre>
49c
            return {p + h, p - h};
96e
d41
d41
        // returns intersection points between two circles
```

```
vector<P> intersect(const C& rhs) const {
5e1
            vector<P> inter:
2b3
            C d = (c - rhs.c).norm();
            if (d > r + rhs.r + EPS or d + min(r, rhs.r) + EPS <
    max(r,rhs.r)) {
c17
                return inter;
800
d41
2a0
            T x = (d * d - rhs.r * rhs.r + r * r) / (d*2),
653
             y = sqrt(r*r - x*x);
485
            P v = (rhs.c - c) / d;
451
            inter.push_back(c + v * x - v.rot(ccw) * y);
            if (y > EPS) inter.push back(c + v * x + v.rot(ccw)
a29
c17
            return inter;
a0a
d41
d41
        // returns the tangents between two circles
834
        vector<pair<P,P>> tangents(C rhs, bool inner = false) {
153
            if (inner) rhs.r = -rhs.r:
260
            P d = (rhs.c - c):
            double dr = (r - rhs, r), d2 = d.norm2(), h2 = d2 - rhs
d7e dr * dr;
fe3
            if (d2 < EPS or h2 < -EPS) return {}:
d41
002
            vector<pair<P. P>> ret:
092
            for (double sign : {-1, 1}) {
042
              P v = (d * dr - d.perp() * sqrt(h2) * sign) / d2;
1da
                ret.emplace back(c + v * r, rhs.c + v * rhs.r);
19b
            }
edf
            return ret;
33b
        }
d41
d41
        // returns the tangent between a point and a circle
        vector<pair<P,P>>> tangents(P p, bool inner = false) {
6e8
            return tangents(C(p, 0), inner);
9c1
310
       }
d85 };
d41 // Given a circle defined by points (a, b, c)
d41 // checks if d is contained in that circle.
199 bool in circle(Pa, Pb, Pc, Pd) {
421 P ad = (a-d), bd = (b-d), cd = (c-d);
11f
b51
            ad.norm2() * cross(bd, cd) -
            bd.norm2() * cross(ad, cd) +
            cd.norm2() * cross(ad, bd)
271
871
      ) >= 0;
573 };
```

convex hull

```
d41 // Convex hull
d41 //
d41 // Description: Monotone chain algorithm.
d41 //
d41 // Details: Border represents the inclusion of collinear points
d41 //
d41 // Complexity: O(nlgn)
```

```
d41
67a template <typename T>
57f vector<T> convex hull(vector<T> p, bool border=0){
        sort(p.begin(), p.end());
d41
        // removes duplicated points
582
        p.erase(unique(p.begin(), p.end()), p.end());
d41
d55
        if (p.size() == 1) return p;
f3d
        int n = p.size(), m = 0;
1d4
        vector<T> ch;
620
        for (int i = 0; i < n; i++, m++) {
           while (m > 1 \text{ and } sign(cross(ch[m-2], ch[m-1], p[i]))
    <= -border) {
                ch.pop back(), m--;
cc3
2h0
3db
            ch.push back(p[i]);
458
        }
d41
4f9
        for (int i = n - 2, t = m; i \ge 0; i - -, m + +) {
           while (m > t and sign(cross(ch[m-2], ch[m-1], p[i]))
61e
    <= -border) {
cc3
                ch.pop back(), m--;
2e4
3db
            ch.push back(p[i]);
902
9d9
        ch.pop back(): // last element is duplicated
d41
66c
        return ch;
30e }
```

halfplane

```
d41 // Halfplane
d41 // Complexity: 0(1) unless stated
d41
4fc template <class T>
c27 struct Halfplane {
h38
       using P = Point<T>;
728
        using H = Halfplane;
d41
812
        P p, pq;
fe9
        double ang;
d41
175
        Halfplane() {}
58b
        Halfplane(P p, P q) : p(p), pq(q-p) {
2a5
            ang = atan2l(pg.y, pg.x);
973
        }
d41
bc2
        bool operator<(H h) const { return ang < h.ang; }</pre>
d41
        bool operator==(H h) const { return fabsl(ang-h.ang)
062
    < EPS; }
d41
fb6
        bool out(P r) { return (pq^(r-p)) < -EPS; }</pre>
d41
3a2
        friend P inter(H s, H t) {
fd0
            double alpha = ((t.p - s.p)^t.pq) / (s.pq^t.pq);
453
            return s.p + s.pq*alpha;
```

```
5e7 }
bc0 };
```

halfplane intersection

```
d41 // Halfplane intersection algorithm
d41 // Description: Finds a convex polygon that is covered by
d41 // If does not exist, then returns an empty vector.
d41 //
d41 // Complexity: O(nlgn)
b37 vector<P> hp intersect(vector<H>& h) {
2f9
        P box[4] = \{
            P(INF, INF),
948
            P(-INF, INF),
c2b
            P(-INF, -INF),
            P(INF, -INF)
c4c
63b
        };
d41
1cd
        for (int i = 0; i < 4; i++) {
7d4
            H = aux(box[i], box[(i+1)%4]);
cd1
            h.push_back(aux);
80f
d41
d77
        sort(all(h));
17a
        deaue<H> da:
486
        int len = 0:
306
        for (int i = 0; i < h.size(); i++) {</pre>
            while (len > 1 and h[i].out(inter(dg[len-1].
    da[len-2]))) {
                dq.pop back(), --len;
cfb
6cb
d41
cdf
            while (len > 1 and h[i].out(inter(dq[0], dq[1]))) {
21b
                 dq.pop front(), --len;
edf
d41
26b
            if (len > 0 and fabsl(h[i].pq^dq[len-1].pq) < EPS) {</pre>
82a
             if (h[i].pg*dg[len-1].pg < 0.0) return vector<P>();
d41
8f5
                if (h[i].out(dq[len-1].p)) {
cfb
                     dq.pop_back(), --len;
a50
                }
64e
                else continue;
bee
d41
395
            dq.push_back(h[i]);
250
            ++len;
e45
d41
        while (len > 2 and dq[0].out(inter(dq[len-1],
    dg[len-2]))) {
cfb
            dq.pop back(), --len;
d81
d41
2e3
        while (len > 2 and dq[len-1].out(inter(dq[0], dq[1]))) {
21h
            dq.pop front(), --len;
365
```

```
d41
1a3
       if (len < 3) return vector<P>();
d41
7e7
        vector<P> ret(len);
cc7
        for (int i = 0; i+1 < len; i++) {
01e
            ret[i] = inter(dq[i], dq[i+1]);
00 f
4fd
        ret.back() = inter(dq[len-1], dq[0]);
edf
        return ret;
ced }
```

intercircles

```
d41 // Circle intersection area algorithm
d41 //
d41 // Description: finds the area covered by at least k circles.
d41 // Complexity: O(n^2 lgn) with high constant.
659 using P = Point<long long>;
3dc using C = Circle<long long>;
d41
397 bool cmp(P a, P b, P o) {
     a = a - 0, b = b - 0:
       return make tuple(-a.side(), 0) < make tuple(-b.side(),</pre>
318
2ce }
d41
clb vector<double> intercircles(vector<C> &c) {
       // area covered by at least k circles
       vector<double> cover(c.size()+1);
018
       for (int i = 0; i < c.size(); i++) {
hfe
            int k = 1:
            vector<pair<P, int>> p = \{\{c[i].c + P(1,0)*c[i].r,
    0}, {c[i],c - P(1.0)*c[i],r, 0}};
d41
226
            for (int j = 0; j < c.size(); j++) {</pre>
bd4
                bool b0 = c[i].in(c[i]), b1 = c[i].in(c[i]):
abe
                if(b0 and (!b1 or i<j)) k++;</pre>
83e
                else if (!b0 and !b1) {
793
                    auto v = c[i].intersect(c[i]);
00b
                    if (v.size() == 2) {
                     p.push_back({v[0], 1}); p.push_back({v[1],
530
    -1}):
99a
                        if (cmp(v[1], v[0], c[i].c)) k++;
7d8
                    }
b50
                }
e05
            }
d41
80e
            P \circ = c[i].c;
381
            sort(p.begin(), p.end(),[&o](auto a, auto b) {
328
                return cmp(a.first,b.first,o);
c2e
            });
d41
b80
            for (int j = 0; j < p.size(); j++) {</pre>
             P a = p[j? j-1: p.size()-1].first, // (j-1+p.size())
bdb
    %p.size()
ff1
                      b = p[i].first:
3c9
                double ang = (a - c[i].c).angle(b-c[i].c);
```

line

```
d41 // Line
d41 //
d41 // Complexity: O(1) unless stated
4fc template <class T>
72c struct Line {
        using P = Point<T>;
ebc
        using L = Line;
d41
b79
        P v; T c;
369
        Line() {};
c30
        Line(P p, T c) : v(v), c(c) {}
193
        Line(P p, P q) : v(q - p), c(v ^ p) {}
d41
5c3
        T side(P p) { return (v^p)-c; }
1a3
        double dist(P p) { return sqrt(dist2(p)); }
8d1
        T dist2(P p) { return side(p) * side(p) / v.norm2(); }
d41
        L perp(P p) { return L(p, p + v.perp()); }
98d
d41
d41
       // orthogonal projection and reflection
efc
       P proj(P p) { return p - v.perp() * side(p) / v.norm2(); }
        P refl(P p) { return p - v.perp() * 2 * side(p) /
    v.norm2(); }
d41
e24
        bool parallel(L l) { return (v ^ l,v) == 0; }
d41
ce0
        L translate(P t) { return L(v, c + (v ^ t)); }
d41
54c
        friend bool inter(L l1, L l2, P &q) {
71b
            if (l1.parallel(l2)) return false;
dd6
            q = (l2.v*l1.c - l1.v*l2.c)/(l1.v^l2.v);
8a6
            return true:
472
        }
d41
db1
        friend L bisector(L l1, L l2, bool interior) {
456
            assert((l1.v^l2.v) != 0);
7e7
            double sign = (interior? 1: -1);
            return L(l2.v / l2.v.norm() + l1.v / l1.v.norm()
b93 * sign,
                     12.c / 12.v.norm() + 11.c / 11.v.norm()
80b * sign);
304
494 };
d41
```

point

```
d41 // Point
d41 //
```

```
67a template <typename T>
afl int sign(T x) { return (x > 0) - (x < 0); }
4fc template <class T>
f26 struct Point {
645
       Tx,y;
11b
        using P = Point;
        Point(T x=0, T y=0) : x(x), y(y) {}
0fa
d41
b3a
        double norm() { return sqrt(norm2()); }
2c3
        T norm2() { return (*this) * (*this): }
d41
286
        P operator+(P p) { return P(x + p.x, y + p.y); }
f4b
        P operator-(P p) { return P(x - p.x, y - p.y); }
a5b
        P operator*(T t) { return P(x * t, y * t); }
da9
        P operator/(T t) { return P(x / t, v / t); }
d41
301
        T operator*(P p) { return x * p.x + y * p.y; }
a91
        T operator^(P p) { return x * p.y - y * p.x; }
d41
163
        bool operator<(P p) { return tie(x,y) < tie(p.x,p.y); }</pre>
eda
        bool operator>(P p) { return tie(x,v) > tie(p,x,p,v); }
cd3
       bool operator==(P p) { return tie(x,y) == tie(p.x,p.y); }
854
       bool operator!=(P p) { return tie(x,v) != tie(p,x,p,v); }
d41
d41
        // rotations are counter-clockwise
785
        P rot(P p) { return P((*this)^p, (*this)*p); }
db2
        P rot(double ang) { return rot(P(sin(ang), cos(ang))); }
d41
889
        double angle() { return atan2(y,x); }
d48
        double angle(P p) {
            return acos(max(-1.0, min(1.0, (*this)*p/
    (norm()*p.norm())));
69b
       }
d41
56d
        P perp() { return P(-y,x); }
6eb
        bool is_perp(P p) { return ((*this) * p) == 0; }
d41
356
        bool side() { return (y > 0) or (y == 0 \text{ and } x < 0); }
d41
0e9
        friend T cross(P a, P b, P c) { return (b-a)^(c-a); }
        friend bool left(P a, P b, P c) { return cross(a,b,c)
cae
        friend bool right(P a, P b, P c) { return cross(a,b,c)
9b8
d41
902
        friend ostream& operator<<(ostream& os, P p) {</pre>
989
            return os<<"("<<p.x<<", "<<p.v<<")":
9a9
6e1 }:
67a template <typename T>
823 void polarsort(vector<T> &v) {
      sort(v.begin(), v.end(), [](T a, T b) {
          return make_tuple(a.side(), 0) < make_tuple(b.side(),</pre>
c05 a ^ b);
f41 });
d35 }
```

d41 // Complexity: 0(1) unless stated

polygon

```
d41 // Polygon structure
d41 //
d41 // Complexity: O(1) unless stated
4fc template <class T>
666 struct Polygon {
        using P = Point<T>;
d41
1a8
        int n;
573
        vector<P> p;
d73
        Polygon() {}
121
        Polygon(vector<P> q) : n(q.size()), p(q) {}
d41
218
        int prev(int i) { return (i-1+n)%n; }
88 f
        int next(int i) { return (i+1)%n; }
d41
       // Positive represents counter-clockwise order. Negative
d41
    represent clockwise order.
d41
       // Complexity: 0(n)
320
       T orientation() {
a1d
            T acum=0:
603
            for (int i = 0; i < n; i++) {
733
                acum += p[i] ^ p[next(i)];
6f0
a13
            return acum;
763
        }
d41
d41
        // Complexity: 0(n)
e04
        double area() { return double(area2())/2; }
        T area2() { return abs(orientation()); }
19e
d41
        // Return true if point q is inside the polygon. The
    polygon must be convex and in ccw order.
d41
        // Complexity: O(lgn)
df7
        bool contains(P q) {
          if (right(p[0],p[1],q) or left(p[0],p[n-1],q)) return
e86
9b1
            int l=1, r=n;
219
            while (r-l > 1) {
ee4
                int m = (l+r)/2;
9eb
                if (!right(p[0],p[m],q)) l=m;
ef3
                else r=m;
638
            }
            if (r == n) return (q-p[0]).norm2() \leftarrow= (p[n-1]-
    p[0]).norm2()+EPS:
e10
            return !right(p[l],p[r],q);
f59
       }
d41
        // Returns true if point q is inside the polygon (does
    not require to be convex).
d41
        // Complexity: 0(n)
35c
        bool rayCasting(P q, bool strict = true) {
3dc
            function<bool(P,P)> above = [](Pa, Pp) {
d3e
                return p.y >= a.y;
940
d41
            function<bool(P,P,P)> crossesRay = [&above](P a, P
    p, P q) {
```

```
return (above(a,q) - above(a,p)) * cross(a,p,q)
eba > 0;
2b7
            };
d41
537
             int numCrossing = 0;
603
             for (int i = 0; i < n; i++) {
6b5
                 if (Segment<T>(p[i], p[next(i)]).has(q)) {
468
                     return !strict;
78a
8c9
                numCrossing += crossesRay(q, p[i], p[next(i)]);
496
6be
             return (numCrossing&1);
83b
        }
d41
d41
        // Returns the squared diameter (uses rotating calipers).
d41
        // Complexity: 0(n)
d5f
        T diameter() {
966
            T ans=0:
1ed
             for (int i = 0, j = 1; i < n; i++) {
                 while (((p[next(i)]-p[i]) ^ (p[next(j)]-p[j]))
^{465} > 0) j = next(j);
                 ans = max(ans, (p[i]-p[j]).norm2());
dcb
37e
ba7
             return ans:
db8
d41
d41
        // Returns the index of an extreme point
d41
        // Complexity: O(lgn)
23b
        int extreme(const function<bool(P, P)> &cmp) {
9b0
            auto isExtreme = [&](int i, bool& curDir) -> bool {
b36
                 curDir = cmp(p[next(i)], p[i]);
                 return !cmp(p[prev(i)], p[i]) && !curDir;
224
2b9
            };
d41
1a0
            bool lastDir, curDir;
c7c
            if (isExtreme(0, lastDir)) return 0;
993
            int l = 0, r = n;
ead
            while (l + 1 < r) {
27e
                 int m = (l + r) >> 1;
b60
                 if (isExtreme(m, curDir)) return m;
ac7
                 bool relDir = cmp(p[m], p[l]);
                if ((!lastDir && curDir) || (lastDir == curDir
     % relDir == curDir)) {
8a6
                    l = m:
986
                    lastDir = curDir;
3dd
                else r = m;
e12
792
             return l;
34c
d41
        // Finds the indices of the two tangents to an external
d41
    point a
d41
        // Complexity: 0(lgn)
349
        pair<int, int> tangent(P q) {
7fc
             auto leftTangent = [\&](P r, P s) \rightarrow bool {
60a
                 return cross(q, r, s) > 0;
403
            };
             auto rightTangent = [\&](P r, P s) \rightarrow bool {
702
e5d
                 return cross(q, r, s) < 0;
ded
            };
```

```
f49
          return {extreme(leftTangent), extreme(rightTangent)};
d35
       }
d41
d41
        // Rotates the polygon points such that
d41
        // p[0] is the lowest leftmost point
d41
        // Complexity: 0(n)
df5
        void normalize() {
            rotate(p.begin(), min_element(p.begin(), p.end()),
   p.end());
7c6
d41
d41
        // Minkowsky sum
d41
        // Complexity: 0(n)
        Polygon operator+(Polygon& other) {
47c
a98
            vector<P> sum;
335
            normalize();
c4f
            other.normalize();
f2b
            for(int i = 0, j = 0; i < n \mid\mid j < other.n; i += dir
   >= 0, j += dir <= 0) {
                sum.push back(p[i % n] + other.p[j % other.n]);
de3
                dir = (p[(i + 1) % n] - p[i % n])
                       ^ (other.p[(j + 1) % other.n] - other.p[j
37e
      other.nl):
090
28c
            return Polygon(sum);
e7e
      }
60a };
```

segment

```
d41 // Segment
d41 //
d41 // Complexity: 0(1) unless stated
4fc template <class T>
fb3 struct Segment {
        using P = Point<T>;
        using S = Segment;
cf0
d41
2ff
        P p,q;
9e4
         Segment() {}
        Segment(P p, P q) : p(p), q(q) {}
07f
d41
2b3
        bool inDisk(P r) { return (p-r)*(q-r) \le 0; }
a04
        bool has(P r) { return cross(p,q,r) == 0 and inDisk(r); }
d41
259
        T side(P r) { return sign((r-p)^(q-p)); }
        double len() { return (q-p).norm(); }
31f
882
        T len2() { return (q-p).norm2(); }
d41
        // for Shamos-Hoey
    #warning caution in the case p == s.q (ex. checking if a
    polygon is simple)
787
        bool operator<(S s) {</pre>
347
            if (p == s.p) return cross(p, q, s.q) > 0;
a30
            if (p.x != q.x \text{ and } (s.p.x == s.q.x \text{ or } p.x < s.p.x))
e68
                 return cross(p, q, s.p) > 0;
d57
            return cross(p, s.q, s.p) > 0;
6c3
```

```
d41
d41
        // checks if two segments intersect
7fb
        friend bool inter(S a, S b) {
            if (a.has(b.p) or a.has(b.q) or b.has(a.p) or
e1e
    b.has(a.q)) {
8a6
                return true;
            }
2a6
cba
            P out;
0d1
            return properInter(a, b, out);
501
d41
bfd
        friend bool properInter(S s1, S s2, P &out) {
3c1
            P = s1.p. b = s1.q. c = s2.p. d = s2.q;
d41
8cf
            double oa = cross(c, d, a),
392
                   ob = cross(c, d, b),
d14
                   oc = cross(a, b, c).
fd6
                   od = cross(a, b, d):
d41
b70
            if (oa*ob < 0 \text{ and } oc*od < 0) {
947
                out = (a*ob - b*oa) / (ob-oa);
8a6
                return true:
787
d41
d1f
            return false:
8ce
        }
d41
a35
        Line<T> getLine() { return Line(p,q); }
cd5 };
d41
```

shamos hoev

```
d41 // Shamos-hoey algorithm
d41 //
    // Description: given a set of segments finds if any two
    segments intersect
d41 // Complexity: O(nlgn)
d41
4cb using S = Segment<ll>;
d41
a21 enum {
83e
       ADD.
f5b
        REM
405 };
d41
779 bool shamos_hoey(vector<S> seg) {
        vector<array<ll, 3> > ev;
071
        for (int i = 0; i < seg.size(); i++) {</pre>
ea5
            if(seg[i].q < seg[i].p) swap(seg[i].p, seg[i].q);</pre>
            ev.push_back({seg[i].p.x, ADD, i});
ee1
9f3
            ev.push_back({seg[i].q.x, REM, i});
2cf
49e
       sort(ev.begin(), ev.end());
e27
        set<S> s:
d9h
        for (const auto &e : ev) {
         S at = seg[e[2]];
dcb
d7d
           if (e[1] == ADD) {
                auto nxt = s.lower bound(at);
```

```
f67
                if((nxt != s.end() && inter(*nxt, at))
                    || (nxt != s.begin() && inter(*prev(nxt),
   at))) {
6a5
                        return 1;
e3e
9be
                s.insert(at);
851
           } else {
                auto nxt = s.upper_bound(at), cur = nxt, prv =
eb2
b75
                if(nxt != s.end() && prv != s.begin()
6d0
                    && inter(*nxt, *(prev(prv)))) return 1;
005
                s.erase(cur);
fe5
           }
d26
       }
hh3
        return 0;
174 }
```

graphs

2sat

```
d41 // 2-SAT algorithm
d41
d41 // Description:
d41 //
           Decideds if a 2-SAT instance has a solution and
d41 //
            finds a valid solution (if possible).
d41 //
d41 // Details:
d41 //
           'sz' in the constructor and reset function represents
d41 //
           the number of clausures.
d41 //
d41 // Complexity:
d41 //
           solve() and reset(n) are O(n).
d41
417 struct TwoSAT {
97b
       int n, c, t;
61a
        vector<vector<int>> adj, adjr;
b57
        vector<int> topo, vis, comp, value;
d41
        TwoSAT(int sz) : n(2*sz), adj(n), adjr(n), topo(n),
    vis(n), comp(n), value(n) {
8bd
            reset(sz):
083
       }
d41
a64
        void reset(int sz) {
686
           n = 2*sz:
603
            for(int i = 0; i < n; i++){
eab
                topo[i] = vis[i] = comp[i] = value[i] = 0;
b1d
                adj[i].clear();
7aa
                adir[i].clear():
245
a34
           t = 0:
809
           c = 1;
e96
       }
d41
00c
        int to vertex(int i) {
9a6
            if(i > 0) return (i - 1) << 1;
            return to_vertex(-i) + 1;
ccb
955
       }
d41
```

```
697
        void add(int i, int j) { // add clause (i or j)
904
            int vi = to vertex(i);
d57
            int vj = to vertex(j);
2c9
            adj[vi^1].push back(vj);
5b0
            adj[vj^1].push back(vi);
d3b
            adjr[vj].push_back(vi^1);
fb2
            adjr[vi].push_back(vj^1);
4d9
       }
d41
aa1
        void add true(int i) { // add clause (i)
7b0
            add(i, i);
a72
       }
d41
6a7
        void add false(int i) { // add clause (neg i)
            add(-i, -i):
6cf
a24
d41
21f
        void add xor(int i, int i) {    // add clause (i xor i)
50e
            add(i, i):
3d2
            add(-i, -i):
def
       }
d41
af5
        void add xnor(int i, int j) { // add clause (i xnor j)
de7
            add(i, -i):
dfb
            add(-i, i):
ece
       }
d41
839
        void dfs1(int v) {
cca
            vis[v] = 1;
5cf
            for(auto x : adj[v]) if(vis[x] == 0) dfs1(x);
11f
            topo[t++] = v;
413
       }
d41
4b2
        void dfs2(int v) {
d5f
            comp[v] = c;
ee1
            for(auto x : adjr[v]) if(comp[x] == 0) dfs2(x);
1d5
       }
d41
6c0
        bool solve() {
830
            for(int i = 0; i < n; i++)
205
                if(vis[i] == 0) dfs1(i);
45b
            for(int i = n - 1; i \ge 0; i - -) {
                if(comp[topo[i]] == 0) {
b4b
256
                    dfs2(topo[i]):
6a4
                    C++;
616
                }
20f
e29
            for(int i = 0; i < n; i+=2) {
5cb
                if(comp[i] == comp[i^1])
d1f
                    return false;
4e6
df1
                    value[i] = (comp[i] < comp[i^1]) ? 0 : 1;
bd8
                    value[i^1] = 1 - value[i]:
f18
               }
е3е
8a6
            return true;
edf
       }
d41
9ca
        bool get_value(int i) {
af7
            return value[to_vertex(i)];
```

```
843
d41
30f
        void print() {
e29
            for(int i = 0; i < n; i += 2) {
b46
                 if(i != 0) cout << ' ';
2e8
                 cout << value[i];</pre>
222
            }
1fb
             cout << endl;</pre>
e99
       }
926 };
```

```
f0d
                     }
863
                 }
7f8
                 if(root && arb <= 1) gart[v].clear();</pre>
5ee
            };
d41
778
             for(int v=0; v<n; v++) {</pre>
d0d
                 if(!vis[v]) dfs(v, 1);
1b3
            }
6d6
        }
c2c };
```

block cut

```
d41 // Block Cut Tree
d41 //
    // Descrition: builds the forest of bluck cut trees for an
    UNDIRECTED graph
d41 //
d41 // Complexity: O(N+M)
d41
142 struct BlockCutTree {
8d3
        int ncomp; // number of components
f7a
        vector<int> comp; // comp[e]: component of edge e
       vector<vector<int>>> gart; // gart[v]: list of components
    an articulation point v is adjacent to
                                 // if v is NOT an articulation
    point, then gart[v] is empty
d41
        // assumes auto [neighbor vertex, edge id] =
d41
    a[current vertex][i]
        BlockCutTree(int n, int m, vector<pair<int,int>>> g[]):
    ncomp(0), comp(m), gart(n) {
6bc
            vector<bool> vis(n), vise(m);
594
            vector<int> low(n), prof(n);
46e
            stack<pair<int,int>> st;
d41
45f
          function<void(int,bool)> dfs = [\&](int v, bool root) {
cca
                vis[v] = 1:
dc9
                int arb = 0; // arborescences
e8a
                for(auto [p, e]: g[v]) if(!vise[e]) {
c8a
                    vise[e] = 1;
934
                    int in = st.size();
20c
                    st.emplace(e, vis[p] ? -1 : p);
137
                    if(!vis[p]) {
f07
                        arb++;
690
                        low[p] = prof[p] = prof[v] + 1;
397
                        dfs(p, 0);
de7
                        low[v] = min(low[v], low[p]);
23d
                    } else low[v] = min(low[v], prof[p]);
d41
c52
                    if(low[p] >= prof[v]) {
c80
                        gart[v].push_back(ncomp);
080
                        while(st.size() > in) {
2b5
                            auto [es, ps] = st.top();
8b3
                            comp[es] = ncomp;
81d
                            if(ps != -1 && !gart[ps].empty())
746
                                gart[ps].push back(ncomp);
25a
                            st.pop();
229
a8f
                        ncomp++;
```

centroid

```
d41 // Centroid decomposition
d41
195 namespace centroid {
a2e
        const int MAXN = 2e5+10;
d41
e81
        vector<int> sz(MAXN, 1);
46f
        vector<bool> removed(MAXN);
def
        vector<vector<int>> adj(MAXN);
d41
765
        int set sizes(int u, int p = -1) {
267
            sz[u] = 1;
087
            for (int v: adj[u]) if (v != p and not removed[v]) {
8d0
                sz[u] += set sizes(v, u);
9c6
            }
f93
            return sz[u];
54d
        }
d41
f47
        int get centroid(int u, int size) {
f8d
            for (int v: adj[u]) if (sz[v] < sz[u]) {
869
              if (sz[v] > size/2) return get centroid(v, size);
7d9
            }
03f
            return u;
1ed
        }
d41
d41
        // For example
846
        vector<int> dists(MAXN);
d41
8a5
        int solve(int root, int k) {
f1e
            int c = get_centroid(root, set_sizes(root));
6f4
            removed[c] = true;
d41
d41
fd0
             * Answer in linear time here
f24
             * Example: number of paths of length k
b44
d41
11e
            int res = 0;
ce6
            queue<int> add, rem;
            function<void(int, int, int)> dfs = [&](int u, int
873 dep, int p) {
926
                if (k >= dep) res += dists[k-dep];
4de
                add.push(dep); rem.push(dep);
087
              for (int v: adi[u]) if (v != p and not removed[v]) {
455
                    dfs(v, dep+1, u);
166
f9f
            };
d41
```

```
3e2
            dists[0] = 1;
            for (int v: adj[c]) if (not removed[v]) {
3c3
05d
                dfs(v, 1, -1);
ce8
                while (add.size()) {
021
                    int d = add.front(); add.pop();
a17
                    dists[d] += 1;
4bf
                }
1e9
            }
d41
51c
            while (rem.size()) {
7ae
                int d = rem.front(); rem.pop();
93b
                dists[d] -= 1;
da7
            }
d41
d41
bfc
             * Example end.
b44
d41
3c3
            for (int v: adj[c]) if (not removed[v]) {
2e4
                res += solve(v, k):
d2b
b50
            return res;
149
       }
071 }:
```

dinic

```
d41 // Dinitz algorithm for finding maximum flows
d41 // Complexity: O(n^2m) (O(m sqrt(n)) on unit networks)
67a template <typename T>
14d struct Dinic {
       struct Edge {
791
            int to:
d90
            T cap, flow:
63d
            Edge(int to, T cap) :
e02
                to(to), cap(cap), flow(0) {}
ffe
       };
d41
327
        int src, snk;
6ff
        vector<int> nxt, ptr, frst, dist, q;
321
        vector<Edge> edges;
d41
7ab
        Dinic(int n) :
02b
            dist(n), frst(n,-1), q(n) {}
d41
4c7
        void add(int u, int v, T cap){
2b3
            edges.emplace back(v,cap);
05a
            nxt.push_back(frst[u]);
197
            frst[u] = edges.size()-1;
d41
265
            edges.emplace_back(u,0);
741
            nxt.push back(frst[v]);
67b
            frst[v] = edges.size()-1;
6e7
       }
d41
838
        bool bfs() {
4f7
            fill(dist.begin(), dist.end(), -1);
f3c
            int st=0, ed=0;
```

```
af1
            q[ed++] = src;
e13
            dist[src] = 0;
d88
            while(st<ed){</pre>
eff
                int u = q[st++];
7f0
                for(int e=frst[u]; e != -1; e = nxt[e]) {
9ac
                    int v = edges[e].to;
                    if(dist[v] == -1 and edges[e].cap >
1b4
    edges[e].flow){
8a0
                        dist[v] = dist[u]+1;
a38
                        q[ed++] = v;
1df
                    }
6ff
               }
c9d
            }
27b
            return dist[snk] != -1;
a67
       }
d41
013
       T dfs(int u, T f=INF){
9b2
            if(u == snk or f == 0) return f:
dca
            for(int &e=ptr[u]; e != -1; e=nxt[e]) {
9ac
                int v = edges[e].to:
                if(dist[u]+1 == dist[v] and edges[e].cap >
    edges[e].flow){
                    int df = dfs(v, min(f,edges[e].cap-
    edges[e].flow)):
f2d
                    if(df>0){
4bc
                        edges[e].flow += df;
1ef
                        edges[e^1].flow -= df;
4ad
                        return df;
e90
                    }
eb8
                }
1c4
            }
bb3
            return 0;
6ac
d41
dfd
       T flow(int _src, int _snk) {
401
            src= src, snk= snk;
8f6
            T flow=0;
6d4
            while(bfs()){
89d
                ptr=frst;
ee5
                while(T add_flow=dfs(src)) flow += add_flow;
b15
            }
99d
            return flow;
59a
971 }:
```

```
24e
        vector<int> ebcc, sz;
d41
       tarjan ebcc(int n, vector<int> adj[]) : nebcc(0), ebcc(n)
5f7
aa0
            int t=0, sn=0;
9a9
            vector<int> disc(n,-1), low(n,-1), stk(n);
d41
87f
            function<void(int,int)> dfs = [&](int u, int p) {
782
                disc[u] = low[u] = t++;
de4
                stk[sn++] = u;
d1c
                for(auto v: adj[u]) {
a30
                   if(disc[v] == -1) {
95e
                        dfs(v.u):
ab6
                        low[u] = min(low[u], low[v]);
h83
f9b
                 else if(v != p) low[u] = min(low[u], disc[v]);
263
c23
                if(low[u] == disc[u]) {
493
                   int v:
f41
                    sz.emplace back(0):
016
                    do {
1ac
                        v = stk[--sn]:
5a5
                        ebcc[v] = nebcc;
1d5
                       sz[nebccl++:
759
                   } while(u != v):
e59
                    nebcc++:
030
               }
e07
           };
d41
603
            for(int i=0; i<n; i++) {</pre>
f2f
                if(disc[i] == -1) dfs(i,i);
527
           }
daf
        }
d41
        vector<vector<int>> condensed graph(int n, vector<int>
a68
    adi[]) {
246
            vector<vector<int>> adj ebcc(nebcc);
687
            for(int u=0; u<n; u++) {
372
                for(int v: adj[u]) {
                   if(ebcc[v] != ebcc[u])
    adj_ebcc[ebcc[u]].push_back(ebcc[v]);
b05
665
ddb
            return adj ebcc;
874
5eb }:
```

ebcc

```
d41 // Tarjan algorithm for edge-biconnected components
d41 //
d41 // Description:
d41 //
           Builds an array ebcc such that ebcc[u] = ebcc[v] iff
d41 //
           u and v are in the same component.
d41 //
d41 //
           Builds an array sz such that sz[ebcc[u]] is the size
d41 //
           of u's component.
d41 //
d41 // Complexity: O(n+m)
d41
cc5 struct tarjan ebcc {
       int nebcc:
```

hld vertex

```
d41 // Heavy-light decomposition for vertices
d41 //
d41 // Description:
d41 // Allows range queries and updates on trees. Template
uses

d41 // lazy segment tree, keep in mind that if structure
is changed
d41 // then the complexities change.
d41 //
d41 // Functions:
```

```
query_path(a,b):
                                  sum of elements in the path a
d41
            update path(a,b,x): sums x to elements in path a -
d41
    > b
            query subtree(a):
                                 sum of elements in the subtree
d41
    of a
d41 ^{\prime}_{\text{of a}}
           update subtree(a,x): sums x to elements in the subtree
d41 //
d41 // Complexity:
d41 //
            query path: 0(lg^2 n)
            update path: O(\lg^2 n)
d41 //
d41 //
            query subtree: O(lgn)
            update subtree: O(lgn)
d41 //
d41
833 const int MAXN = 6e5;
826 namespace hld {
        segtree<long long> seg;
        vector<int> h(MAXN), v(MAXN), p(MAXN), pos(MAXN),
    sz(MAXN), peso(MAXN);
d41
baf
        void build(int root, int n, vector<int> adj[]) {
6bb
d41
5ae
            function<void(int,bool)> dfs = [\&](int u, bool f) {
f1a
                 v[pos[u]=t++] = peso[u];
267
                 sz[u] = 1;
cab
                 for(auto &v: adj[u]) {
567
                     if(v == p[u]) continue;
                     p[v] = u;
                    h[v] = (v == adj[u][0]? h[u]: v);
e60
1ae
                     dfs(v,f);
                     sz[u] += sz[v];
                     if(sz[v] > sz[adj[u][0]] \text{ or } adj[u][0] ==
    } ([v]q
d67
                         swap(v, adj[u][0]);
bbf
                    }
155
94c
                if(u == root and f) dfs(h[u]=u,t=false);
f60
            };
d41
09a
            dfs(root.true);
0da
            seg = segtree<long long>(t,v);
e94
d41
4f6
        long long query path(int a, int b) {
aa1
            if(pos[a] < pos[b]) swap(a,b);</pre>
a42
            if(h[a] == h[b]) return seg.query(pos[b],pos[a]);
            return seg.query(pos[h[a]], pos[a]) +
    query path(p[h[a]], b);
567
        }
d41
920
        void update path(int a, int b, int x) {
aa1
            if(pos[a] < pos[b]) swap(a,b);</pre>
            if(h[a] == h[b]) return seq.update(pos[b],pos[a],x);
370
            seg.update(pos[h[a]], pos[a], x);
            update_path(p[h[a]], b, x);
7c6
1db
        }
d41
```

```
665    long long query_subtree(int a) {
        return seg.query(pos[a], pos[a]+sz[a]-1);
7f8    }
d41
acc    void update_subtree(int a, int x) {
        seg.update(pos[a], pos[a]+sz[a]-1, x);
002    }
827 };
```

kuhn

```
d41 // Kuhn algorithm
d41 //
    // Description: computes a maximum matching in a bipartite
    graph
d41 //
d41 // Complexity: O(VE)
    // Details: if not using add, the adjacency of only one side
    is needed
d41
d41
    chrono::steady_clock::now().time_since_epoch().count());
d41
el2 struct Kuhn {
14e
        int n.m:
903
        vector<vector<int>> adi:
d23
        vector<int> mu. mv:
d41
0f2
        Kuhn(int n, int m) :
fb4
            n(n), m(m), adj(n), mu(n,-1), mv(m,-1) {}
d41
0a9
        void add(int u, int v) {
cc9
            adj[u].push back(v);
451
d41
bf7
        int matching() {
6a9
            vector<bool> vis(n+m, false);
d41
6a4
            function<br/><br/>bool(int)> dfs = [\&](int u) {
b9c
                vis[u] = true;
f07
                for(int v: adj[u]) if(!vis[n+v]) {
a46
                    vis[n+v] = true;
ea2
                    if(mv[v] == -1 \text{ or } dfs(mv[v]))  {
12d
                        mu[u] = v, mv[v] = u;
8a6
                         return true;
32a
                    }
b31
                }
d1f
                return false;
0a2
            };
d41
1ae
            int ret=0, aum=1;
92b
            for(auto& u: adj) shuffle(u.begin(), u.end(), rng);
392
            while(aum) {
0fa
                fill(vis.begin()+n, vis.end(), false);
c5d
                aum=0:
687
                for(int u=0: u<n: u++) {
                    if(mu[u] == -1 and dfs(u)) ret++, aum=1;
5da
5e9
```

```
891 }
edf return ret;
140 }
2f5 };
```

lca-rmg

```
d41 // Lowest common ancestor using RMQ
d41 // Description: finds the LCA in O(1) using RMQ
d41
74b namespace LCA {
adc
        vector<int> dep, pos, inv;
e16
        RMQ<pair<int,int>> rmg;
d41
7a2
        void build(int root, const vector<vector<int>>& adj) {
19f
            dep.resize(2 * adj.size());
940
            inv.resize(2 * adj.size());
be5
            pos.resize(adj.size());
d41
6bb
            int t = 0;
            function<void(int, int, int)> dfs = [&](int u, int
    d, int p) {
a86
                inv[t] = u, pos[u] = t, dep[t++] = d;
914
                 for (int v: adj[u]) if (v != p) {
1fe
                    dfs(v, d+1, u);
678
                    inv[t] = u, dep[t++] = d;
043
18e
            }:
d08
            dfs(root, 0, -1);
904
             rma = RMO(dep):
d41
8a4
            vector<pair<int,int>> vals(t);
960
             for (int i = 0; i < t; i++) {
f29
                vals[i] = {dep[i], i};
7ba
d9d
             rmq = RMQ(vals);
a67
        }
d41
310
        int lca(int u, int v) {
bdc
            if (pos[u] > pos[v]) swap(u, v);
e91
             return inv[rmq.query(pos[u], pos[v])];
f02
        }
d41
c11
        int dist(int u, int v) {
           return dep[pos[u]] + dep[pos[v]] - 2 * dep[pos[lca(u,
bb1 v)]];
10d
340 }
```

scc

```
d41 // Tarjan algorithm for strongly connected components
d41 //
d41 // Description:
d41 // Builds an array scc such that scc[u] = scc[v]
d41 // iff u and v are in the same component. And build an
d41 // array sz such that sz[scc[u]] is the size of u's
component.
d41 //
```

```
d41 // Complexity: O(n+m)
24b struct tarjan_scc {
f43
        int nscc;
35f
        vector<int> scc, sz;
d41
11b
       tarjan_scc(int n, vector<int> adj[]) : nscc(0), scc(n) {
aa0
            int t=0, sn=0;
9a9
            vector<int> disc(n,-1), low(n,-1), stk(n);
d41
214
            function<void(int)> dfs = [\&](int u) {
782
                disc[u] = low[u] = t++;
de4
                stk[sn++1 = u:
d1c
                for(auto v: adj[u]) {
0h4
                    if(disc[v] == -1) dfs(v);
                    low[u] = min(low[u], low[v]);
ab6
8a2
c23
                if(low[u] == disc[u]) {
d93
                    int v:
f41
                    sz.emplace back(0):
016
                    do {
1ac
                        v = stk[--sn]:
2fh
                        scc[v] = nscc;
2a5
                        sz[nscc]++:
b62
                        low[v] = INF:
db0
                    } while(u != v):
983
                    nscc++;
3cd
                }
47b
            };
d41
603
            for(int i=0; i<n; i++) {
162
                if(disc[i] == -1) dfs(i);
5cd
9fc
       }
d41
d41
        // Complexity: O(n+m)
        vector<vector<int>> condensed graph(int n, vector<int>
a68
2a3
            vector<vector<int>> adj_scc(nscc);
687
            for(int u=0; u<n; u++) {
372
                for(int v: adj[u]) {
                    if(scc[v] != scc[u])
    adj scc[scc[u]].push back(scc[v]);
6e8
cc7
6f4
            return adj scc;
f3d
ecd };
```

math

bit iterator

d41

```
5f7
             Mask(uint32_t _msk): msk(_msk) {}
             bool operator!=(const Mask& rhs) const { return msk
22e < rhs.msk; };
                  void operator++(){const uint32 t t=msk|
     (msk-1); msk=(t+1)|(((\sim t\&-\sim t)-1)>> builtin ffs(msk));}
600
            uint32_t operator*() const { return msk; }
cc7
        };
1dc
        uint32_t n, m;
75a
        BitIterator(uint32_t _n, uint32_t _m): n(_n), m(_m) {}
17a
        Mask begin() const { return Mask((1<<m)-1); }</pre>
d0d
        Mask end() const { return Mask((1<<n)); }</pre>
8ca };
```

convolutions

```
d41 // Convolutions
d41 // Description: Multiply two polinomial
d41 // Complexity: O(N logN)
d41 // Functions:
d41 // multiply(a, b)
          multiply mod(a, b, m) - return answer modulo m
d41
d41 // Details:
d41 // For function multiply mod, any modulo can be used.
d41 //
          It is implemented using the technique of dividing
          in sart to use less fft. Function multiply may have
d41 //
d41 //
          precision problems.
          This code is faster than normal. So you may use it
d41 //
d41 //
          if TL e tight.
d32 const double PI=acos(-1.0):
35b namespace fft {
3b2 struct num {
662
           double x.v:
c0a
           num() \{x = v = 0; \}
6da
         num(double x, double y): x(x), y(y){}
cd4
      }:
       inline num operator+(num a, num b) {return num(a.x +
    b.x, a.y + b.y);}
       inline num operator-(num a, num b) {return num(a.x -
    b.x, a.y - b.y);}
       inline num operator*(num a, num b) {
            return num(a.x * b.x - a.y * b.y, a.x * b.y + a.y
9f0
d63
db0
       inline num conj(num a) {return num(a.x, -a.y);}
d41
b58
        int base = 1;
e47
        vector<num> roots=\{\{0,0\}, \{1,0\}\};
8a4
        vector<ll> rev={0, 1};
148
        const double PI=acosl(-1.0);
d41
d41
       // always try to increase the base
d50
       void ensure base(int nbase) {
11e
         if(nbase <= base) return:</pre>
49f
            rev.resize(1 << nbase);</pre>
55a
            for (int i = 0; i < (1 << nbase); i++)
19e
             rev[i] = (rev[i>>1] >> 1) + ((i&1) << (nbase-1));
            roots.resize(1<<nbase);</pre>
```

```
775
            while(base<nbase) {</pre>
21f
                 double angle = 2*PI / (1<<(base+1));
8cf
                 for(int i = 1 << (base-1); i < (1 << base); i++) {
52a
                     roots[i<<1] = roots[i];
aef
                    double angle i = angle * (2*i+1-(1<<base));
                     roots[(i << 1)+1] =
     num(cos(angle_i),sin(angle_i));
958
                 base++;
af4
            }
ae9
        }
d41
        void fft(vector<num> &a,int n=-1) {
b94
05e
            if(n==-1) n=a.size();
421
             assert((n&(n-1)) == 0);
2fd
            int zeros = builtin ctz(n);
a02
            ensure base(zeros):
4fa
            int shift = base - zeros:
603
             for (int i = 0: i < n: i++) {
3fc
                if(i < (rev[i] >> shift)) {
b8b
                     swap(a[i],a[rev[i] >> shift]);
9ac
h97
7cd
             for(int k = 1: k < n: k <<= 1) {
                for(int i = 0: i < n: i += 2*k) {
cda
0c2
                    for(int j = 0; j < k; j++) {
d85
                         num z = a[i+j+k] * roots[j+k];
20a
                         a[i+j+k] = a[i+j] - z;
c9a
                         a[i+j] = a[i+j] + z;
ee1
                    }
804
                }
b62
            }
382
        }
d41
ba5
        vector<num> fa, fb;
d41
        // multiply with less fft by using complex numbers.
318
        vector<ll> multiply(vector<ll> &a, vector<ll> &b);
d41
        // using the technique of dividing in sgrt to use less
d41 <sub>fft.</sub>
        vector<ll> multiply_mod(vector<ll> &a, vector<ll> &b,
    ll m. ll ea=0):
        vector<ll> square mod(vector<ll>&a. ll m):
7a3 }:
d41
7b3 vector<ll> fft::multiply(vector<ll> &a, vector<ll> &b) {
        int need = a.size() + b.size() - 1:
d41
217
        int nbase = 0:
8da
        while((1 << nbase) < need) nbase++;</pre>
4e5
        ensure base(nbase):
d41
729
        int sz = 1 << nbase;</pre>
9dh
        if(sz > (int)fa.size()) fa.resize(sz);
887
        for(int i = 0; i < sz; i++) {</pre>
422
            ll x = (i < (int)a.size() ? a[i] : 0);
435
            ll y = (i < (int)b.size() ? b[i] : 0);
685
             fa[i] = num(x, y);
3e3
        }
```

```
fft(fa, sz);
650
4db
        num r(0, -0.25/sz);
        for(int i = 0; i \le (sz >> 1); i++) {
3ec
b13
            int j = (sz-i) & (sz-1);
afc
            num z = (fa[j] * fa[j] - conj(fa[i] * fa[i])) * r;
            if(i != j) fa[j] = (fa[i] * fa[i] - conj(fa[j] *
    fa[j])) * r;
            fa[i] = z;
        }
488
d41
650
        fft(fa, sz);
5e0
        vector<ll> res(need):
07f
        for(int i = 0; i < need; i++) res[i] = fa[i].x + 0.5;
h50
        return res:
16b }
d41
d41 // 4eb347
    vector<ll> fft::multiply mod(vector<ll> &a, vector<ll> &b,
    ll m, ll eq) {
        int need = a.size() + b.size() - 1;
217
        int nbase = 0:
Sh8
        while((1 << nbase) < need) nbase++;</pre>
        ensure base(nbase):
729
        int sz = 1 << nbase:</pre>
9db
        if(sz > (int)fa.size()) fa.resize(sz);
        for(int i = 0; i < (int)a.size(); i++) {</pre>
538
            ll x = (a[i] % m + m) % m;
7e5
            fa[i] = num(x & ((1 << 15) - 1), x >> 15);
b60
26e
        fill(fa.begin() + a.size(), fa.begin() + sz, num{0,0});
650
        fft(fa, sz);
32a
        if(sz > (int)fb.size()) fb.resize(sz);
b19
        if(eq) copy(fa.begin(), fa.begin() + sz, fb.begin());
4e6
            for(int i = 0; i < (int)b.size(); i++) {</pre>
1da
044
                ll x = (b[i] % m + m) % m;
418
                fb[i] = num(x & ((1 << 15) - 1), x >> 15);
9f0
           fill(fb.begin() + b.size(), fb.begin() + sz, num{0,0}
535 );
07e
            fft(fb.sz):
59c
df3
        double ratio = 0.25 / sz:
dc2
        num r2(0, -1), r3(ratio, 0), r4(0, -ratio), r5(0, 1);
3ec
        for(int i = 0; i \le (sz >> 1); i++) {
            int j = (sz - i) & (sz - 1);
b13
d96
            num al = (fa[i] + conj(fa[j]));
6c3
            num a2 = (fa[i] - conj(fa[j])) * r2;
712
            num b1 = (fb[i] + coni(fb[i])) * r3:
e45
            num b2 = (fb[i] - coni(fb[i])) * r4:
41e
            if(i != j) {
123
                num c1 = (fa[j] + conj(fa[i]));
7ce
                num c2 = (fa[j] - conj(fa[i])) * r2;
92b
                num d1 = (fb[j] + conj(fb[i])) * r3;
a76
                num d2 = (fb[i] - conj(fb[i])) * r4;
35f
                fa[i] = c1 * d1 + c2 * d2 * r5;
525
                fb[i] = c1 * d2 + c2 * d1;
55a
dc5
            fa[j] = a1 * b1 + a2 * b2 * r5;
```

```
fb[i] = a1 * b2 + a2 * b1;
d92
dc3
f68
       fft(fa, sz); fft(fb, sz);
5e0
        vector<ll> res(need);
ae6
        for(int i = 0; i < need; i++) {
6e0
           ll aa = fa[i].x + 0.5;
bb7
            ll\ bb = fb[i].x + 0.5;
0ee
            ll cc = fa[i].y + 0.5;
407
            res[i] = (aa + ((bb%m) << 15) + ((cc%m) << 30))%m;
0d6
b50
       return res;
ca4 }
d41
b86 vector<ll> fft::square mod(vector<ll> &a, ll m) {
        return multiply mod(a, a, m, 1);
dde }
```

crt

```
d41 // Chinese Remainder Theorem
d41 // Joins modular linear equations: x = a \pmod{m}
d41 // if there are no solutions, a = -1
153 template<typename T> tuple<T, T, T> ext_gcd(T a, T b) {
3bd
      if (!a) return {b, 0, 1};
550
       auto [g, x, y] = ext gcd(b%a, a);
c59
       return \{g, y - b/a*x, x\};
537 }
d41
bfe template<typename T = ll> struct crt {
627
       Ta.m:
d41
5f3
       crt() : a(0), m(1) {}
7eb
       crt(T a , T m ) : a(a ), m(m ) {}
       crt operator * (crt C) {
911
238
           auto [q, x, v] = ext qcd(m, C.m):
dc0
           if ((a - C.a) % g) a = -1;
4f9
           if (a == -1 or C.a == -1) return crt(-1, 0):
d09
           T lcm = m/g*C.m;
eb2
           T ans = a + (x*(C.a-a)/q % (C.m/q))*m;
d8d
           return crt((ans % lcm + lcm) % lcm, lcm);
1f2 }
0d9 };
```

division trick

```
d41 // Division trick
d41 //
d41 // complexity: 0(sqrt(n))
d41
79c for(int l = 1, r; l <= n; l = r + 1) {
746     r = n / (n / l);
d41     // n / i has the same value for l <= i <= r
5bf }</pre>
```

ext_gcd

```
d41 // Extended GCD
d41 //
d41 // Descrption:
```

```
d41 // Let g = gcd(a,b)
d41 // Returns {g,x,y} such that ax+by = g
d41 //
d41 // Complexity: O(log(min(a,b)))
d41
6f9 tuple<int,int,int> ext_gcd(int a, int b) {
3bd         if (!a) return {b, 0, 1};
550         auto [g, x, y] = ext_gcd(b%a, a);
c59         return {g, y - b/a*x, x};
596 }
```

fwht

```
d41 // Fast Walsh-Hadamard trasform
d41 // Description: Multiply two polynomials such that x^a * x^b
d41 // - op(a, b) = a "xor" b, a "or" b, a "and" b
d41 // Complexity: O(n log n)
29a const ll N = 1<<20;
d41
67a template <typename T>
372 struct FWHT {
        void fwht(T io[], ll n) {
495
            for (ll d = 1; d < n; d <<= 1) {
h98
                for (ll i = 0, m = d << 1; i < n; i += m) {
                  for (ll j = 0; j < d; j++) { /// Don't forget
499 modulo if required
bdc
                        T x = io[i+j], y = io[i+j+d];
703
                     io[i+j] = (x+y), io[i+j+d] = (x-y); // xor
d41
                       // io[i+j] = x+y; // and
d41
                        // io[i+j+d] = x+y; // or
afa
7f3
               }
           }
cf6
fe6
1f8
        void ufwht(T io[], ll n) {
495
            for (ll d = 1: d < n: d <<= 1) {
b98
                for (ll i = 0, m = d << 1; i < n; i += m) {
                  for (ll j = 0; j < d; j++) { /// Don't forget
499 modulo if required
bdc
                        T x = io[i+j], y = io[i+j+d];
d41
                        /// Modular inverse if required here
                        io[i+j] = (x+y)>>1, io[i+j+d] = (x-i)
   v)>>1; // xor
d41
                        // io[i+j] = x-y; // and
d41
                        // io[i+j+d] = v-x; // or
027
                   }
711
               }
fd4
           }
e69
       }
        // a, b are two polynomials and n is size which is power
d41
683
        void convolution(T a[], T b[], ll n) {
26b
            fwht(a, n), fwht(b, n);
e95
            for (ll i = 0: i < n: i++)
33e
                a[i] = a[i]*b[i];
23c
            ufwht(a, n):
5b4
d41
        // for a*a
```

11

integration

matrix

```
d41 // Matrix structure
d41 //
d41 // Complexity:
d41 //
            operator*: 0(n^3)
d41 //
            operator^: O(n^3 lgn)
67a template<typename T>
953 struct Matrix : vector<vector<T>>> {
      int n,m;
               Matrix(const vector<vector<T>>& c) :
36b vector<vector<T>>(c), n(c.size()), m(c[0].size()) {}
       Matrix(int n, int m) : vector<vector<T>>(n, vector<T>(m,
cf8 0)),n(n),m(m){}
      Matrix(const initializer list<initializer list<T>>& c) {
bd0
f7e
            vector<vector<T>> val:
212
            for(auto& i:c) val.push back(i):
5a8
            *this = Matrix(val):
ad3
       }
254
        Matrix operator*(Matrix& r) {
1e2
            assert(m == r.n):
0f7
            Matrix M(n, r,m):
830
            for(int i=0:i<n:i++)</pre>
800
                for(int k=0:k<m:k++)</pre>
418
                    for(int j=0;j<r.m;j++)</pre>
47e
                        M[i][j] += (*this)[i][k]*r[k][j];
474
            return M;
658
       }
7cf
        Matrix operator^(ll e){
0d6
            Matrix M(n.n):
40a
            for(int i=0;i<n;i++) M[i][i] = 1;</pre>
13c
            for(Matrix at=*this;e;e/=2,at=at*at)
2e2
                if(e\&1) M = M*at;
474
            return M;
9ee
      }
b3e };
```

mint

```
d41 // Modular integer structure
9a3 template <const int MOD>
e54 struct mint {
3ec
       int x;
9f5
        mint(): x(0) \{
        mint(int _x): x(_x%MOD<0?_x%MOD+MOD:_x%MOD) {}</pre>
5b8 =MOD; }
        void operator+=(mint rhs) { x+=rhs.x; if(x>=MOD) x-
        void operator-=(mint rhs) { x-=rhs.x; if(x<0)x+=MOD; }</pre>
a9a
d08
        void operator*=(mint rhs) { x*=rhs.x; x%=MOD; }
        void operator/=(mint rhs) { *this *= rhs.inv(); }
        mint operator+(mint rhs) { mint res=*this; res+=rhs;
     return res; }
        mint operator-(mint rhs) { mint res=*this; res-=rhs;
ee4
     return res: }
        mint operator*(mint rhs) { mint res=*this; res*=rhs;
384
     return res: }
        mint operator/(mint rhs) { mint res=*this; res/=rhs;
dd6
     return res: }
7ea
        mint inv() { return this->pow(MOD-2); }
714
        mint pow(int e) {
30h
            mint res(1):
65a
            for(mint p=*this; e>0; e/=2, p*=p) if(e%2)
bbc
                res*=p:
b50
            return res;
f35
     }
b64 }:
7e5 using Z = mint<998244353>;
```

xor_basis

```
d41 // XOR basis
d41 //
d41 // Description:
d41 //
            builds a basis in \Integers 2 that spans the
d41 //
            vectors inserted.
d41 //
d41 //
            think of integers as vectors in \Integers 2,
d41 //
            i.e., 4 \rightarrow (1.0.0).
d41 //
d41 // Complexity: O(SZ), unless stated otherwise
722 template <const int SZ>
b5d struct XorBasis {
727
       int base[SZ];
d41
79f
       XorBasis() { memset(base, 0, sizeof base); }
d41
1f5
       bool bit on(int x, int i) {
2ab
            return (x&(1LL<<i));</pre>
627
       }
d41
1b3
        void insert(int v) {
da0
            for (int i = SZ-1; i \ge 0; i--) {
0b3
                if (bit_on(v, i)) {
7e8
                   if (!base[i]) {
e02
                        base[i] ^= v:
```

```
return;
7ed
                    }
0d7
                    v ^= base[i];
1e7
9b5
           }
736
        }
d41
d41
        // merges two basis in O(SZ^2)
92f
        void merge(XorBasis &b) {
da0
            for (int i = SZ-1; i \ge 0; i--) {
f59
                if (b.base[i]) insert(b.base[i]);
019
           }
b87
        }
d41
d41
        // checks if v can be achieved
3e7
        bool check(int v) {
c67
            int x = 0:
592
            for (int i = SZ-1: i >= 0: i--) if (base[i]) {
5aa
                if (bit on(x, i) != bit on(v, i)) {
25a
                    x ^= base[i]:
2af
fe8
            }
57d
            return (x == v);
8d8
d41
d41
        // returns the maximum value
e92
        int max() {
404
            int value = 0;
            for (int i = SZ-1; i \ge 0; i--) if (!bit on(value,
781 i)) {
cca
                value ^= base[i];
214
af2
            return value;
e92 }
2b4 };
```

strings

aho

```
d41 // Aho-Corasick
d41 // Description: data structure that allows search
d41 // for multiple patterns in text.
d41 //
d41 // Details: K is the alphabet size, in template
d41 // it is using alphabet from 'a' to 'z' with offset 'a'
f12 template <const int K = 26>
d3c struct aho_corasick {
3c9
      struct node {
714
           int next[K], len,
78 f
               link, // suffix link
a92
             term link; // suffix link to "first" terminal state
           bool is term; // terminal state (i.e., a word from
    dictionary ends here)
           node() : link(-1), term_link(-1), is_term(false) {
cf9
               fill(next, next+K, -1);
237
           }
e2c
a46
      };
```

```
d41
bbb
        vector<node> t;
d41
968
        aho corasick() : t(1) {}
        aho corasick(int size) : t(1) { t.reserve(size); }
cbd
d41
498
        int new node() {
83d
            t.emplace_back();
b18
            return t.size()-1;
1ae
d41
d41
        // 0(|s|)
d45
        void add string(const string& s) {
            int u = 0:
400
aec
            for (const char& ch: s) {
15f
                int c = ch - 'a';
              if (t[u].next[c] == -1) t[u].next[c] = new node();
a1c
2c2
                u = t[u].next[c]:
01d
d34
            t[u].is term = true, t[u].len = s.size();
842
        }
d41
d41
        // 0(|dictionary size|)
        void build() {
0bb
            queue<int> q; q.push(0);
402
            while (q.size()) {
                int u = q.front(); q.pop();
                for (int c = 0; c < K; c++) if (t[u].next[c] !=</pre>
176
                    int v = t[u].next[c], l = t[u].link;
0d6
ceb
                    while (l != -1 \text{ and } t[l].next[c] == -1) {
a96
                        l = t[l].link;
                    t[v].link = (l == -1? 0: t[l].next[c]);
                    t[v].term link = (t[t[v].link].is term?
f42
t[v].link: t[t[v].link].term_link);
                    q.push(v);
                }
ff7
            }
227
        }
d41
        // 0(|dictionary size|)
c3b
        void add dictionary(const vector<string>& dict) {
            for (const string& s: dict) add string(s):
4bb
6f2
            build():
dd5
        }
d41
d41
61c
         * Example: CSES Word Combinations.
2fa
         * dp[i] is the # of ways to generate string [1...i]
6a0
         * (where s[1...0] = emptv string)
h44
b2f
        mint query(const string& s) {
b49
            int at = 0, i = 0;
0d7
            vector<mint> dp(s.size()+1); dp[0] = 1;
            for (const char& ch: s) {
15f
                int c = ch - 'a';
d41
3e4
                if (t[at].term) dp[i] += dp[i-t[at].len];
```

12

```
for (int e = t[at].term_link; e != -1; e =
759 t[e].term_link) {
                   dp[i] += dp[i-t[e].len];
eee
d41
385
                while (at != -1 and t[at].next[c] == -1) {
054
                    at = t[at].link;
409
2bd
                at = (at == -1? 0: t[at].next[c]);
0dd
db2
d41
            if (t[at].term) dp[i] += dp[i-t[at].len]:
            for (int e = t[at].term link; e != -1; e =
    t[e].term link) {
d70
               dp[i] += dp[i-t[e].len];
eee
d41
648
            return dp[s.size()];
9a3
c42 }:
```

hash

```
d41 // String Hashing
   // Description: A data structure that transforms a string
d41
d41 // Functions:
d41 // str hash - Builds the hash in O(|S|)
d41 // operator() - Gives the number representing substring
    s[l,r] in O(1)
d41
d41 // Details:
d41 // - To use more than one prime, you may use long
    long, __int128 or array<int>
    // - You may easily change it to handle vector<int>
    instead of string
    // - Other large primes: 1000041323, 100663319,
    201326611
               - If smaller primes are needed(For instance,
    need to store the mods in an array):
              - 50331653, 12582917, 6291469, 3145739, 1572869
d41 //
d41
4ba const long long mod1 = 1000015553, mod2 = 1000028537;
    mt19937 rng((int) chrono::steady_clock::now().time_since_epocl
    random number generator
d41
463 int uniform(int l, int r) {
       uniform_int_distribution<int> uid(l, r);
f54
       return uid(rng);
d9e }
d41
3fb template<int MOD>
d7d struct str hash {
      static int P:
c63
dcf
       vector<ll> h, p;
      str hash () {}
```

```
str_hash(string s) : h(s.size()), p(s.size()) {
7a2
            p[0] = 1, h[0] = s[0];
ad7
            for (int i = 1; i < s.size(); i++)</pre>
                p[i] = p[i - 1]*P%MOD, h[i] = (h[i - 1]*P +
84c s[i])%MOD;
af7
        ll operator()(int l, int r) { // retorna hash s[l...r]
749
           ll hash = h[r] - (l?h[l-1]*p[r-l+1]%MOD:0);
dfd
            return hash < 0 ? hash + MOD : hash;</pre>
      }
3ba
977 };
    template<int MOD> int str_hash<MOD>::P = uniform(256, MOD = 0.000)
    1): // l > |sigma|
d41
61c struct Hash {
d41
       // Uses 2 primes to better avoid colisions
3h6
      str hash<mod1> H1:
b36
        str hash<mod2> H2:
d41
        Hash (string s) : H1(str hash<mod1>(s)),
e3d
    H2(str hash<mod2>(s)) {}
d41
af7
        ll operator()(int l, int r) {
f6f
            ll ret1 = H1(l, r), ret2 = H2(l, r);
742
            return (ret1 << 30) ^ (ret2);</pre>
d2e
      }
b31 };
```

kmp

```
d41 // Knuth-Morris-Pratt algorithm (KMP)
 d41 //
 d41 // Description:
 d41 // . The prefix function, p[i] is the longest suffix that
     ends at i
 d41 //
           that is also a suffix.
 d41 //
 d41 // Complexity: O(n)
 67a template<typename T>
 29c vector<int> kmp(int sz, const T s[]) {
         vector<int> p(sz);
 e8d
         for(int i=1;i<sz;i++) {</pre>
 42d
             int \&j = p[i];
 084
             for(j=p[i-1];j>0 && s[i]!=s[j];j=p[j-1]);
 04b
             if(s[i] == s[i]) i++;
 b13
()74eunt())et/u/rn p;
 8af }
 d41
 f42 vector<vector<int>> automaton(string s) {
 46e
         s += '#';
 c73
         int n = (int)s.size();
 cf1
        auto p = kmp(s.size(), s.data());
 c3d
         vector<vector<int>> aut(n, vector<int>(26,0));
 603
         for(int i=0:i<n:i++) {
 1e1
             for(int c=0;c<26;c++) {
                if(i>0 \&\& 'a'+c != s[i]) aut[i][c] = aut[p[i-1]]
42c [c];
                 else aut[i][c] = i + ('a' + c == s[i]);
```

```
237 }
40c }
8a7 return aut;
593 }
```

manacher

```
d41 // Manacher algorithm
d41 //
d41 // Description:
d41 //
            Finds a vector with palindromes size.
d41 // in i
            ret[2*i] = size of largest palindrome with center
d41 ^{\prime\prime} ret[2*i+1] = size of largest palindrome with center in i and i+1
d41 //
d41 // Complexity: O(n)
67a template <typename T>
44d vector<int> manacher(const T& s) {
        int l = 0, r = -1, n = s.size();
fc9
        vector<int> d1(n), d2(n);
        for (int i = 0; i < n; i++) {
821
           int k = i > r ? 1 : min(d1[l+r-i], r-i);
            while (i+k < n \&\& i-k >= 0 \&\& s[i+k] == s[i-k]) k++;
61a
61e
            d1[i] = k--:
9f6
            if (i+k > r) l = i-k, r = i+k:
950
e03
       l = 0. r = -1:
603
        for (int i = 0: i < n: i++) {
a64
            int k = i > r ? 0 : min(d2[l+r-i+1], r-i+1): k++:
            while (i+k \le n \&\& i-k \ge 0 \&\& s[i+k-1] == s[i-k])
2c6 k++;
            d2[i] = --k;
            if (i+k-1 > r) l = i-k, r = i+k-1;
26d
4fe
c41
        vector<int> ret(2*n-1);
        for (int i = 0; i < n; i++) ret[2*i] = 2*d1[i]-1;
e1d
        for (int i = 0; i < n-1; i++) ret[2*i+1] = 2*d2[i+1];
edf
9ca }
```

minimum rotation

palindromic tree

```
d41 // Palindromic tree (eertree)
d41 // Complexity: O(|s|) amortized
c71 struct palindromic_tree {
3c9
       struct node {
ff6
            int length, link;
3a4
            map<char, int> to;
          node(int length, int link): length(length), link(link)
697
45d
       };
b9e
       vector<node> nodes;
9fc
        int current:
d36
        palindromic tree(): current(1) {
31a
            nodes.push back(node(-1, 0));
5ec
            nodes.push back(node(0, 0));
2fa
       }
ea5
       void add(int i, string& s) {
            int parent = nodes[current].length == i ?
    nodes(current).link : current:
f2b
            while (s[i - nodes[parent].length - 1] != s[i])
c9c
                parent = nodes[parent].link;
            if(nodes[parent].to.find(s[i]) !=
    nodes[parent].to.end()) {
b6c
                current = nodes[parent].to[s[i]];
9a4
4e6
            else {
490
                int link = nodes[parent].link;
0bb
                while (s[i - nodes[link].length - 1] != s[i])
304
                    link = nodes[link].link;
569
                link = max(1, nodes[link].to[s[i]]);
b39
               current = nodes[parent].to[s[i]] = nodes.size();
                nodes.push_back(node(nodes[parent].length +
    2, link));
739
          }
ec3
bfc
       void insert(string& s) {
dd9
            current = 1;
9a8
            for (int i = 0; i < int(s.size()); i++)</pre>
df9
                add(i, s);
c50
      }
0d1 };
d41
```

suffix array

```
d41 // Suffix array
d41 //
d41 // Description:
d41 //
            Builds a sorted array of all suffixes of a string.
d41 //
d41 // Complexity:
           build: O(nlgn)
d41 //
3f4 struct SuffixArray {
19b
       vector<int> v. inv. lcp:
d41
        template <tvpename T>
        SuffixArray(int sz, const T s[]): v(sz+1), inv(sz+1),
292 lcp(sz) {
```

```
5b8
             iota(v.begin()+1, v.end(), 0);
            sort(v.begin()+1, v.end(), [&](int i, int j) { return
d5f s[i] < s[j]; });
b8a
             v[0] = sz++;
311
            vector<int> ra(sz), newra(sz), newv(sz), cnt(sz+1);
941
             ra[v[1]] = 1;
             for(int i=2;i<sz;i++) ra[v[i]] = ra[v[i-1]] +</pre>
     (s[v[i]] != s[v[i-1]]);
             for(int k=1; k<sz; k*=2) {</pre>
22f
                 for(int i=0; i < sz; i++) v[i] = (v[i]-k+sz) %sz;
0ca
35b
                 fill(cnt.begin(), cnt.end(), 0);
1b1
                 for(int x: ra) cnt[x+1]++;
6db
               partial sum(cnt.begin(), cnt.end(), cnt.begin());
                 for(int i=0;i<sz;i++) newv[cnt[ra[v[i]]]++]</pre>
eb3
    = v[i];
2de
                 v.swap(newv);
e8d
                 for(int i=1;i<sz;i++) {</pre>
                   int diff = ra[v[i]] != ra[v[i-1]] || ra[(v[i]]
     +k)%sz] != ra[(v[i-1]+k)%sz]:
0ec
                     newra[v[i]] = newra[v[i-1]] + diff;
729
694
                 ra.swap(newra);
dh2
f28
             for(int i=0:i<sz:i++) inv[v[i]] = i:</pre>
f7a
             for(int l=0.i=0:i<sz-1:i++) {</pre>
e3e
                 int i = v[inv[i]-1]:
040
                 while(\max(i,j)+l < sz-1 \&\& s[i+l] == s[j+l]) l++;
f58
                 lcp[inv[i]-1] = l;
89d
                 if(l) l--;
a89
            }
afc
       }
901 };
```

suffix automaton

```
d41 // Suffix Automaton
d41 //
d41 // Description:
d41 //
           Builds a minimal DFA that accepts all suffixes
d41 //
           of a string
d41 //
d41 // Complexity:
           build: O(n)
d41 //
d41
92e struct SuffixAutomaton {
       struct Node {
bf2
a46
            int len, link;
3a4
            map<char, int> to;
            Node() : len(0), link(-1) {}
ad3
a4d
       };
d41
0f1
        int n, sz, last;
f5b
        vector<Node> t;
d41
244
       void add(char c) {
37a
            int cur = sz++:
24a
            t[cur].len = t[last].len + 1;
           int p = last:
81e
8d7
            while (p != -1 \&\& !t[p].to.count(c)) {
                t[p].to[c] = cur;
```

```
e06
                p = t[p].link;
2a9
924
            if (p == -1) {
                t[cur].link = 0;
b64
c44
            } else {
a7d
                int q = t[p].to[c];
9a9
                if (t[p].len + 1 == t[q].len) {
cdd
                    t[cur].link = q;
913
               } else {
90a
                    int clone = sz++;
477
                    t[clone].len = t[p].len + 1;
46d
                    t[clone].to = t[q].to;
b6f
                    t[clone].link = t[q].link;
520
                    while (p != -1 \&\& t[p].to[c] == q) {
                        t[p].to[c] = clone;
h1e
e06
                        p = t[p].link;
b85
762
                    t[a].link = t[cur].link = clone:
914
a9c
691
            last = cur;
607
        }
d41
       SuffixAutomaton(string &s): n(s.size()), sz(1), last(0),
    t(2*n) {
217
            for (auto c: s) add(c):
      };
99b };
```

tri

```
d41 // Trie data structure
d41 //
d41 // Description:
d41 //
           Data structure for strings.
d41 //
d41 // Complexity:
d41 //
           add: 0(|s|)
d41 //
           erase: 0(|s|)
d41 //
           find: O(|s|)
d41
ab5 struct trie {
        int sigma, offset;
5e5
        vector<int> wrd, cnt;
        vector<vector<int>> to;
d41
498
       int new node() {
a2e
           wrd.push back(0);
a02
           cnt.push back(0);
526
           to.push_back(vector<int>(sigma));
501
           return to.size()-1;
a02
       }
d41
30f
        trie(int sigma=26, int offset='a'):
bc6
           sigma( sigma), offset( offset) {
839
           new node():
819
       }
d41
d41
        // adds s to the trie
22d
        void add(string &s) {
```

```
181
           int i = 0;
b4f
           for(char c: s) {
2a6
               int &nxt = to[i][c-offset];
29d
               if(!nxt) nxt = new node();
b70
               cnt[i=nxt]++;
f14
           }
c8f
           wrd[i]++;
333
       }
d41
d41
       // if s is in the trie removes it and returns true,
d41
       // else returns false
5c9
       bool erase(string &s) {
3d0
           if(!find(s)) return false:
181
           int i=0:
h4f
           for(char c: s) {
2a6
               int &nxt = to[i][c-offset];
d25
               cnt[i=nxt]--:
4d9
               if(!cnt[nxt]) nxt=0:
678
           }
8cc
           wrd[i]--:
8a6
           return true:
       }
65e
d41
       // if s is in the trie returns true.
d41
       // else returns false
9b8
       bool find(string &s) {
181
        int i=0;
b4f
           for(char c: s) {
2a6
               int &nxt = to[i][c-offset];
59c
               if(!nxt) return false;
799
               i = nxt;
ff6
           }
243
           return (wrd[i] > 0);
bbd
     }
789 };
```

Z

```
d41 // Z-function
d41 //
d41 // Description:
d41 // z[i]: is the longest suffix that starts in i that is also
d41 // a suffix that starts at 0.
d41 //
d41 // Complexity: O(n)
67a template<typename T>
7f4 vector<int> z function(int n, const T s[]) {
       vector<int> z(n);
2c3
       int x = 0, y = 0;
6f5
       for (int i = 1; i < n; i++) {
063
        z[i] = max(0, min(z[i-x], y-i+1));
0d8
           while (i+z[i] < n \&\& s[z[i]] == s[i+z[i]]) {
654
               x = i; y = i+z[i]; z[i]++;
54c
           }
a64
       }
070
       return z;
944 }
```

structures

```
d41 // BIT (Fenwick tree)
d41 //
d41 // Functions:
d41 //
           query(l,r): sum of elements in range [l,r]
d41 //
           update(i,x): sums x to value in position i
d41 //
d41 // Details: guery and update are 1-indexed.
d41 //
d41 // Complexity:
d41 //
           build: O(n)
d41 //
            query: O(lqn)
d41 //
           update: O(lgn)
67a template <typename T>
714 struct BIT{
1a8
        int n:
678
        vector<T> bit:
d41
03f
        BIT(int n) : n(n+1), bit(n+1){}
6c8
        BIT(vector<int> &v) : n(ssize(v)+1), bit(ssize(v)+1) {
465
            for(int i=1:i<n:i++) bit[i] = v[i-1]:</pre>
6f5
            for(int i=1:i<n:i++) {</pre>
edf
               int i = i+(i\&-i):
2h8
                if(j<n) bit[j] += bit[i];</pre>
604
           }
       }
64b
d41
b13
        T pref querv(int i) {
501
           T sum = 0:
185
            for(;i;i-=i&-i) sum+=bit[i];
e66
            return sum;
       }
bea
d41
b7a
        T query(int l, int r) {
6b8
            return pref guery(r)-pref guery(l-1);
1d4
d41
2ac
        void update(int i, T x) {
4a3
            for(;i<n;i+=i&-i) bit[i]+=x;</pre>
0e8
      }
099 };
```

lazy seg

```
d41 // Segment tree with lazy propagation
d41 //
d41 // Functions:
d41 //
          update(l,r,x): sums x to all elements in range [l,r]
d41 //
           query(l,r): sum of all elements in range [l,r]
d41 //
d41 // Complexity:
           build: O(n)
d41 //
d41 //
           update: O(lgn)
d41 //
           query: O(lqn)
d41
d41 // neutral element
cad const int NEUT = 0;
67a template <typename T>
```

```
07c struct segtree {
1a8
        int n:
243
        vector<T> seq, lazy;
d41
ff0
        T build(int p, int l, int r, vector<T> &v) {
3c7
            lazv[p] = 0:
6cd
            if(l == r) return seq[p] = v[l];
ee4
            int m = (l+r)/2;
         return seg[p] = build(2*p,l,m,v)+build(2*p+1,m+1,r,v);
06d
d41
        segtree(vector<T> &v) : n(v.size()), seg(4*n), lazy(4*n)
      build(1,0,n-1,v); }
        segtree(int n) : n(n), seg(4*n), lazy(4*n) {}
9a8
d41
ceb
        void prop(int p, int l, int r) {
            seg[p] += lazy[p]*(r-l+1);
            if(l != r) lazv[2*p] += lazv[p], lazv[2*p+1] +=
2c9 lazy[p];
            lazv[p] = NEUT:
7fd
       }
d41
69e
        T update(int a, int b, int x, int p, int l, int r) {
6h9
            prop(p.l.r):
e9f
            if(b < l or r < a) return seg[p]:</pre>
9a3
            if(a <= l and r <= b) {
b94
                lazy[p] += x;
6h9
                prop(p,l,r);
534
                return seq[p];
821
            int m=(l+r)/2;
            return seg[p] = update(a,b,x,2*p,l,m)+update(a,b,x,
    2*p+1.m+1.r):
9d4
       }
d41
88b
        T query(int a, int b, int p, int l, int r){
6b9
            prop(p,l,r);
527
            if(a <= l and r <= b) return seq[p];</pre>
            if(b < l or r < a) return NEUT;</pre>
            int m = (l+r)/2;
ee4
b1f
            return query(a,b,2*p,l,m)+query(a,b,2*p+1,m+1,r);
98b
d41
        void update(int l, int r, int x) { update(l,r,x,
    1.0.n-1): }
       T query(int l, int r) { return query(l,r,1,0,n-1); }
a37 }:
d41
```

15

median

```
6c5
               ll x = *(qt.begin());
9ad
               qt.erase(qt.begin());
2da
               sm.insert(x);
e12
           } else if (sm.size() > (size + 1) / 2) {
e79
               ll x = *(sm.begin());
440
               sm.erase(sm.begin());
27d
               qt.insert(x);
b89
           }
b39
       }
d41
e1b
       void insert(ll v) {
           if (size == 0 || v <= median_first()) sm.insert(v);</pre>
8ab
005
           else at.insert(v):
b56
           size++:
da6
           balance():
ef3
       }
d41
32d
       bool remove(ll v) {
b47
        if (size == 0) return false;
4b9
           if (v <= median first()) {</pre>
8f2
               auto it = sm.find(v):
               if (it == sm.end()) return false:
5e2
570
               sm.erase(it);
5cc
           } else {
ba5
               auto it = gt.find(v);
5dc
               if (it == qt.end()) return false:
59d
               qt.erase(it);
e93
           }
1ec
            size--;
da6
           balance();
8a6
           return true;
579
       }
d41
fcd
       ll median_first() {
4fb
           if (size % 2 == 1) return *(sm.begin());
e77
            return *(sm.begin());
187
       }
d41
23d
       double median() {
4fb
        if (size % 2 == 1) return *(sm.begin());
a29
           return (*(sm.begin()) + *(gt.begin())) / 2.0;
9ab
       }
d41
393
       string median string() {
           if (size % 2 == 1) return to string(*(sm.begin()));
bdc
bf7
           ll a = *(sm.begin()) + *(gt.begin());
ac0
           strina s:
9c6
           if (a < 0) s += '-':
           s += to string(abs(a) / 2);
5e9
           if (abs(a) % 2 == 1) s += ".5":
013
047
           return s:
2f1 }
55f }:
```

mo

```
d41 // MOs algorithm
d41 //
d41 // Description:
d41 // Answers queries offline with sqrt decomposition
```

```
d41 //
d41 // Complexity:
d41 // exec: O(nsgrt(n)(O(remove)+O(add)))
678 \text{ const int } SZ = 230;
d41
670 struct Query {
738
        int l, r, idx;
9a6
        Query () {}
       Query (int _l, int _r, int _idx) : l(_l), r(_r), idx(_idx)
        bool operator < (const Query &o) const {</pre>
9ae
04d
            return {l / SZ, r} < {o.l / SZ, o.r};</pre>
4a3
d0f };
d41
67a template <typename T>
5ce struct MO {
8d9
        int sum:
b61
        MO(\text{vector} < T > \& v) : sum(0), v(v), cnt(MAXN), C(MAXN) {}
d41
f5d
        void exec(vector<Query> &queries, vector<T> &answers) {
14d
            answers.resize(queries.size());
bfa
            sort(queries.begin(), queries.end());
d41
3df
            int cur l = 0:
cf5
            int cur r = -1;
d41
275
            for (Query q : queries) {
71e
                while (cur l > q.l) {
ec6
                    cur l--;
939
                    add(cur_l);
60c
294
                while (cur_r < q.r) {</pre>
bda
                    cur r++;
d95
                    add(cur r);
c3b
b32
                while (cur l < q.l) {
631
                     remove(cur_l);
cf9
                    cur l++;
ddf
6eb
                while (cur r > q.r) {
198
                    remove(cur r);
99e
                    cur r--:
d76
553
                answers[q.idx] = get answer(cur l, cur r);
8bc
c0e
        }
d41
        void add(int i) {
c96
683
            sum += v[i]:
0c3
d41
17e
        void remove(int i) {
f2f
            sum -= v[i]:
9a0
d41
83a
        T get_answer(int l, int r) {
e66
             return sum;
```

```
42d }
d84 };
```

ordered set

```
d41 // Ordered set (0-indexed)
d41 //
d41 // Description:
^{\prime\prime} It has the same properties and functions of the
    regular set and
d41 //
            two extra functions.
d41 //
d41 // Functions:
d41 ^{//}_{\text{index k}}
            find by order(k) - iterator to the element with
d41 //
            order of key(k) - index of the element k in the set
d41 //
d41 // Complexity:
d41 //
            find_by_order: O(logn)
d41 //
            order_of_key: O(logn)
774 #include <ext/pb ds/assoc container.hpp>
30f #include <ext/pb ds/tree policy.hpp>
0d7 using namespace __gnu_pbds;
4fc template <class T>
def using ord set = tree<T, null_type, less<T>, rb_tree_tag,
       tree order statistics node update>;
```

persistent seg

```
d41 // Persistent Segment Tree
d41 //
d41 // Functions:
d41 // update(
             update(p,x,t): updates the value at position p to
\mbox{\bf d41} \begin{tabular}{ll} // & \mbox{update(p,x):} & \mbox{updates the value at position p to} \\ x \mbox{ of the last version} & \end{tabular}
d41 // query(l,r,t): sum of all elements in range [l,r]
    at version t
d41 //
d41 // Complexity:
             build: O(n)
d41 //
d41 //
             update: O(lan)
d41 //
             query: O(lgn)
67a template <typename T>
e22 struct pseaf
        const int LIM = 1e7:
        int n. rt. nxt:
130
         vector<T> sea:
8h1
         vector<int> L. R:
d41
af3
         pseq(int n): seq(LIM), L(LIM), R(LIM), n(n), nxt(0) {}
d41
188
        int new_node(T x, int l=0, int r=0) {
2f6
             int nn = nxt++;
21d
             seq[nn] = x; L[nn] = l; R[nn] = r;
7c3
             return nn;
605
        }
d41
```

```
7e7
        int build(vector<T> &v, int l, int r) {
ed7
            if(l==r) return new node(v[l]);
            int m=(l+r)/2, nl=build(v,l,m), nr=build(v,m+1,r);
565
e0f
            return new_node(seg[nl]+seg[nr],nl,nr);
ccb
        }
d41
291
        int build(vector<T> &v) { return rt = build(v,0,n-1); }
d41
a15
        int update(int p, T x, int k, int l, int r) {
4ca
            int nn = new node(seg[k],L[k],R[k]);
e88
            if(l == r){seq[nn]=x;return nn;}
            int m = (l+r)/2;
ee4
bd6
            if(p <= m) L[nn] = update(p,x,L[nn],l,m);</pre>
            else R[nn] = update(p,x,R[nn],m+1,r);
c4e
fff
            seg[nn] = seg[L[nn]] + seg[R[nn]];
7c3
            return nn:
0a1
        }
d41
d99
        T query(int a, int b, int k, int l, int r) {
a2c
            if (r < a \text{ or } b < l) return 0:
f49
            if(a <= l and r <= b) return seg[k];</pre>
ee4
            int m = (l+r)/2:
            return query(a,b,L[k],l,m) + query(a,b,R[k],m+1,r);
84e
a22
        }
d41
        int update(int p, T x, int t) { return rt=update(p,x,t,
ab2
    0, n-1); }
2b0
        int update(int p, T x) { return update(p,x,rt); }
        T query(int l, int r, int t) { return query(l,r,t,
c72 0,n-1); }
1ad };
```

seg_iterative

```
d41 // Seament tree
d41
4fc template <class T>
9cc struct segment tree {
374
       T NEUT = T():
d41
1a8
        int n;
130
       vector<T> seq;
d41
d41
       // O(n)
409
        segment tree(vector<T> &v) : n(v.size()), seg(2*n) {
            for (int i = 0; i < n; i++) seq[i+n] = v[i];
1a1
            for (int i = n-1; i; i--) seq[i] = seq[2*i] +
d15
    seq[2*i+1];
d13
       }
d41
d41
       // O(lgn)
6a3
        void update(int pos, T val) {
            for (seg[pos += n] = val; pos >>= 1;)
fa3
385
                seq[pos] = seq[2*pos] + seq[2*pos+1];
1d1
       }
d41
d41
       // O(lgn)
       T query(int l, int r) {
b7a
66d
           T tl = NEUT, tr = NEUT;
            for (l += n, r += n; l < r; l >>= 1, r >>= 1) {
```

segtree

```
d41 // Seament tree
d41 //
d41 // Functions:
d41 //
            update(p,x): updates the value at position p to x
d41 //
           query(l,r): the sum of all elements in range [l,r]
d41 //
d41 // Complexity:
d41 //
           build: O(n)
d41 //
           update: O(lgn)
d41 //
           query: O(lgn)
d41
67a template <typename T>
07c struct segtree {
1a8
        int n;
130
        vector<T> seq;
d41
ff0
        T build(int p, int l, int r, vector<T> &v) {
           if(l == r) return seg[p] = v[l];
6cd
ee4
            int m = (l+r)/2:
         return seg[p] = build(2*p,l,m,v)+build(2*p+1,m+1,r,v);
e2a
94d
d41
f87
        segtree(vector<T> &v) : n(v.size()), seg(4*n) {
7ed
           build(1,0,n-1,v);
811
4f2
        segtree(int n) : n(n), seg(4*n) {}
d41
6c6
        T update(int a, T x, int p, int l, int r) {
fa6
            if(a < l or r < a) return seg[p]:</pre>
b0e
            if(l == r) return seg[p] = x;
ee4
            int m=(l+r)/2;
            return seq[p] = update(a,x,2*p,l,m)+update(a,x,
9bd
   2*p+1,m+1,r);
      }
394
d41
88b
        T query(int a, int b, int p, int l, int r){
527
            if(a <= l and r <= b) return seg[p];</pre>
786
            if(b < l or r < a) return 0;</pre>
ee4
            int m = (l+r)/2;
b1f
            return guery(a,b,2*p,l,m)+guery(a,b,2*p+1,m+1,r);
85d
       }
d41
7d4
        void update(int p, T x) { update(p,x,1,0,n-1); }
190
        T query(int l, int r) { return query(l,r,1,0,n-1); }
2e5 };
```

slope_trick

```
d41 // SlopeTrick
d41 //
d41 // Armazena uma estrutura convexa piecewise linear
```

```
// Comentarios acima das funcoes para explicar o que cada
67a template<typename T>
406 struct SlopeTrick {
       T inf = numeric_limits<T>::max() / 3;
        T min f;
f32
        priority queue<T, vector<T>, less<>> L;
6ef
        priority gueue<T, vector<T>, greater<>> R;
        T add l, add r;
a20
d41
055
        T top R() {
            if (R.empty()) return inf;
a34
ffe
            else return R.top() + add r;
074
       }
d41
70c
       T pop R() {
            T val = top_R();
66f
8e0
            if (R.size()) R.pop();
d94
            return val:
21d
       }
d41
d9d
        T top L() {
b7b
            if (L.empty()) return -inf;
470
            else return L.top() + add_l;
31d
       }
d41
821
       T pop L() {
66a
            T val = top L();
1e0
            if (L.size()) L.pop();
d94
            return val;
dfd
       }
d41
86d
        size t size() {
7ff
            return L.size() + R.size();
c4b
d41
0e8
        SlopeTrick() : min_f(0), add_l(0), add_r(0) {};
d41
d41
        // return {min f(x), lx, rx}
d41
        // Em que [lx, rx] eh o intervalo que atinge o minimo
5ee
        arrav<T, 3> querv() {
e8a
            return {min f, top L(), top R()};
14f
d41
d41
        // f(x) += a
ad4
        void add all(T a) {
            min_f += a;
f8c
78a
d41
d41
        // add \
d41
        // f(x) += max(a - x, 0)
60a
        void add a minus x(T a) {
            min f += max(T(0), a - top R());
8c6
            R.push(a - add_r);
cdb
            L.push(pop_R() - add_l);
416
44c
       }
d41
```

// Permite adicionar slopes sem peso e realizar query de

17

```
d41
        // add /
d41
        // f(x) += max(x - a, 0)
7a9
        void add x minus a(T a) {
b36
            min_f += max(T(0), top_L() - a);
988
            L.push(a - add l);
e5a
            R.push(pop_L() - add_r);
f3a
        }
d41
d41
        // add \/
d41
        // f(x) += abs(x - a)
825
        void add abs(T a) {
9cc
            add_a_minus_x(a);
e55
            add x minus a(a);
639
        }
d41
d41
        // \/ -> \
d41
        // f {new} (x) = min f(y) (y <= x)
73b
        void clear right() {
b8e
            while (R.size()) R.pop();
2b3
        }
d41
d41
        // \/ -> /
d41
        // f \{new\} (x) = min f(y) (y >= x)
fd5
        void clear left() {
e21
            while (L.size()) L.pop();
bc4
        }
d41
d41
        // \/ -> \ /
d41
        // f_{\text{new}} (x) = \min f(y) (x-b \le y \le x-a)
564
        void shift(T a, T b) {
25b
            assert(a <= b);</pre>
b95
            add_l += a;
165
            add r += b;
29a
       }
d41
d41
        // \/. -> .\/
d41
        // f_{\text{new}} (x) = f(x - a)
5d6
        void shift(T a) {
a77
            shift(a, a);
af1
d41
d41
       // Retorna f(x)
d41
       // 0(size)
c2a
       T get(T x) {
7ce
            auto L2 = L;
202
            auto R2 = R;
bf4
            T ret = min f;
6a9
            while (L.size()) {
efd
                ret += \max(T(0), pop L() - x);
e50
            }
886
            while (R.size()) {
97b
                ret += \max(T(0), x - pop_R());
8ef
98a
            L = L2, R = R2;
edf
            return ret:
093
        }
d41
d41
        // O(min(size, st.size))
9e9
        void merge(SlopeTrick &st) {
f68
            if (st.size() > size()) {
```

```
079
                swap(*this, st);
788
           }
1a3
           while (st.R.size()) {
85b
                add_x_minus_a(st.pop_R());
2c5
8c6
            while (st.L.size()) {
897
                add_a_minus_x(st.pop_L());
b31
           }
eaf
            min_f += st.min_f;
       }
3df
f24 };
```

sparse seg

```
d41 // Sparse segment tree
d41 //
d41 // Functions:
d41 //
            query(l,r): sum in range [l,r]
d41 //
           update(l,r,x): sums x to all elements in [l,r]
d41 //
d41 // Complexity:
d41 //
           build: 0(1)
d41 //
           query: O(lgn)
d41 //
           update: O(lgn)
d41
67a template <typename T>
75a struct sseg{
86f
        const int LIM = 1e7;
e42
        const int N = 1e9:
d41
2a7
        int nxt:
        vector<T> seg, lazy;
243
8h1
        vector<int> L.R:
d41
9a3
        sseg() : seg(LIM), lazy(LIM), L(LIM), R(LIM), nxt(2) {}
d41
e9a
        int get l(int i) {
ba7
           if(!L[i]) L[i] = nxt++;
a96
            return L[i];
a5a
       }
d41
943
        int get_r(int i) {
2e9
           if(!R[i]) R[i] = nxt++;
283
            return R[i];
cb5
       }
d41
ceb
        void prop(int p, int l, int r) {
cdf
            seg[p] += lazy[p]*(r-l+1);
            if(l!=r) lazy[get_l(p)]+=lazy[p], lazy[get_r(p)]
821
    +=lazy[p];
3c7
           lazy[p] = 0;
       }
cae
d41
183
        T update(int a, int b, T x, int p, int l, int r) {
6b9
           prop(p,l,r);
9a3
           if(a \le l and r \le b)
h94
                lazy[p] += x;
6b9
                prop(p,l,r);
534
                return seg[p];
821
```

```
e9f
            if(b < l or r < a) return seg[p];</pre>
ee4
            int m = (l+r)/2;
            return seg[p] = update(a,b,x,get_l(p),l,m) +
    update(a,b,x,get_r(p),m+1,r);
       }
d41
88b
        T query(int a, int b, int p, int l, int r) {
6b9
            prop(p,l,r);
527
            if(a <= l and r <= b) return seg[p];</pre>
786
            if(b < l or r < a) return 0;
ee4
            int m = (l+r)/2;
            return query(a,b,get_l(p),l,m)
    +query(a,b,get r(p),m+1,r);
ca4
d41
6b7
        void update(int l, int r, T x){update(l,r,x,1,0,N-1);};
fab
        T query(int l, int r){return query(l,r,1,0,N-1);};
f87 }:
```

18

sparse table

```
d41 // Sparse table
d41 //
d41 // Functions:
d41 //
            guery(l,r): minimum element in range [l,r]
d41 //
d41 // Complexity:
d41 //
            build: O(nlan)
d41 //
            query: 0(1)
67a template <typename T>
991 struct sparse {
1a8
        int n:
5a6
        vector<vector<T>> m;
d41
5c0
        sparse(vector<T> &v) : n(v.size()) {
020
            int lg = 32 - builtin clz(n);
60a
            m.assign(lg,vector<T>(n));
78e
            for(int i=0;i<n;i++) m[0][i] = v[i];</pre>
05c
            for(int j=1;(1<<j)<=n;j++) {
ed9
                for(int i=0;i+(1<<j)<=n;i++) {
5d5
                m[j][i] = min(m[j-1][i], m[j-1][i+(1<<(j-1))]);
642
b3a
            }
7d1
       }
d41
b7a
       T query(int l, int r) {
133
            int j = 31 - __builtin_clz(r-l+1);
56f
            return min(m[j][l], m[j][r-(1<<j)+1]);</pre>
7b9
       }
87a };
```