

Consultation restricted to the "Alloy quick reference". Duration: 50 minutes.

Student name: _____ Number _____

1. [11.2 points] For each question, mark with a cross (X) the correct answer
Each correct answer is graded 1.6 points. Each incorrect answer is graded -0.4 points.

a) Pixels always have three components: red, green and blue. A possible translation to Alloy would be:
<pre> sig Pixel { } sig Red, Green, Blue extends Pixel { } enum Component { Red, Green, Blue } sig Pixel { components: Component -> Int } sig Pixel { Red : Int, Green: Int, Blue: Int } All the previous answers are correct </pre>
b) In an University, a School has several Departments, but a Department belongs to just one school. How to translate that in Alloy?
<pre> sig University { schools: some School, deps : Department some -> one schools } sig University { schools: some School, deps : schools -> Department } sig University { schools: some School, deps : schools some -> one Department } sig University { schools: some School, deps : schools 1..* -> 1 Department } </pre>
c) What is the converse ($\sim R$) of the binary relation $R = \{(a,b), (b,c), (c,b)\}$?
<pre> ~R = { (a,b), (b,a), (b,c), (c,b) } ~R = { (a,b), (b,c), (c,b), (a,a), (b,b), (c,c) } ~R = { (b,a), (c,b), (b,c) } ~R = { (b,c), (c,b) } </pre>
d) Given $R1 = \{(a,a), (a,b), (b,c)\}$ and $R2 = \{(a,c)\}$ what is the value of the restriction $R1 \rightarrow R2$?
<pre> R1 :> R2 = {(a,a)} R1 :> R2 = {(a,a,a), (a,a,c)} R1 :> R2 = {(a,a), (b,c)} R1 :> R2 = {(a,a), (a,b)} </pre>
e) Given $R1 = \{(a,b), (b,b)\}$, $R2 = \{(a,a)\}$ and $R3 = \{(b,a)\}$ what is the value of $(R1 \cup R2) \cup R3$?
<pre> {(a,b), (b,b), (a,a), (b,a)} {(a,a), (b,b), (b,a)} {(a,a), (b,a)} None of the previous answers is correct </pre>
f) Consider a graph definition where each node has a set of adjacent nodes. A graph is connected if there is a path from every node to any other node. How to express that constraint in Alloy?
<pre> fact connected { iden in ^ adjacent } fact connected { all disj n1, n2: Node n2 in n1.^ adjacent } fact connected { all n1: Node n1 in n1.* adjacent } All the previous answers are correct </pre>
g) Given $\text{sig Exam}(\text{grades: Student} \rightarrow \text{Int})$, how can we obtain the pairs (exam, student) that received a certain grade?
<pre> fun results [g : Int]: Exam -> Student { g <: grades } fun results [g : Int]: Exam -> Student { grades <: g } fun results [g : Int]: Exam -> Student { grades :> g } None of the previous answers are correct </pre>

2. [8.8 points] Fill in the empty blocks.

```
sig Account {}
```

```
abstract sig Transaction { amount: Int }
```

```
0.2 sig Deposit, Withdrawal extends Transaction -- A transaction is either a deposit
                                         or a withdrawal
```

```
sig Client {
```

```
0.4   accounts: some Account, -- a client can access several accounts (1 or more)
```

```
0.4   withdrawPrivileges: set Account, -- but can't withdraw from all of them (0..*)
```

```
0.6   balance: Account set -> lone Int -- the amount each account currently has
```

```
0.6   transactions: Account one -> set Transaction -- a list of all account movements
}
```

```
pred invariants[c: Client] {
```

```
-- the balance of an account should never be lower than 0
```

```
0.6   all n : Account.balance | n >= 0
```

```
-- a client can only withdraw from accounts she has access to
```

```
0.6   withdrawPrivileges in accounts
```

```
-- a client only has balance from accounts she has access to
```

```
0.6   balance.Int in accounts
}
```

```
-- transaction t withdraws quantity q from account a of client c,
-- resulting in a new state c'
```

```
0.2 pred withdraw[c, c': Client, a: Account, qty: Int, t: Transaction] {
    -- pre-conditions (without using predicate invariants)
```

```
(TODO)
```

```
1.4
```

```
-- post-conditions (without using predicate invariants)
```

```
(TODO)
```

```
1.6
```

```
-- gives the total balance of a client c
```

```
fun totalBalance[c: Client] : Int {
```

```
sum a: c.accounts | c.balance[a]
```

```
0.8
```

```
}
```

```
assert withdraw_preserves_invariants {
```

```
all c, c': Client, account: Account, qty: Int, t: Transaction |
```

```
-- if one withdraws from a consistent client
```

```
0.4   (invariants[c] and withdraw[c, c', a, qty, t]) =>
```

```
-- one ends up with a new consistent client state
```

```
0.4   ( invariants[c'] )
```

```
}
```

```
check withdraw_preserves_invariants
```

Good luck!