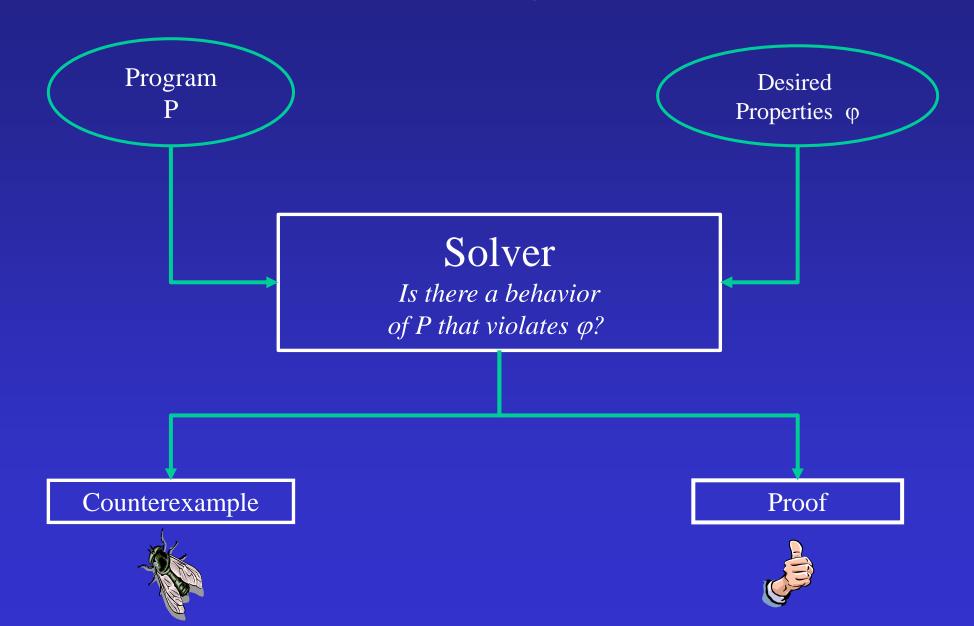
Introduction to Software Synthesis

Mooly Sagiv
Ideas and Slides taken from Sumit Gulwani, Armando *Solar-Lezama*Eran Yahav, Emain Torlak,

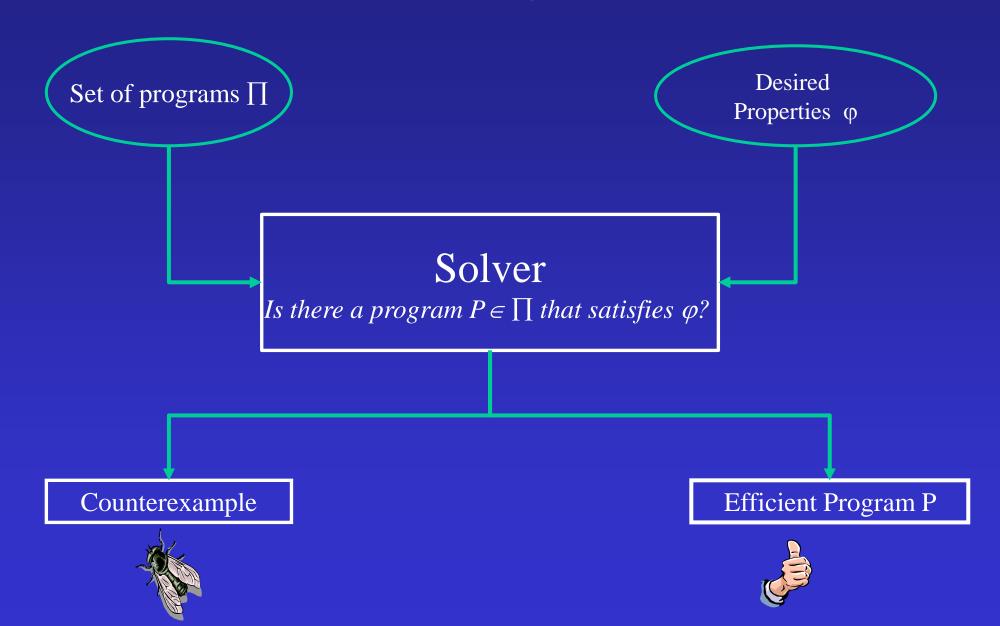
Recap

Problem	Tools
Propositional SAT solving	MiniSat, Z3
First order solving with theories (SMT)	Z3, CVC3
Bounded Model Checking	CBMC, JBMC
Concolic Execution	DART, KLEE, SAGE, Cloud9, Mayhem
Static analysis	SLAM(SDV), Astrée, TVLA, CSSV
Testing	PITTEST, AFL
Program Synthesis	SKETCH(MIT), Rosettee(UWASH)

Verification vs. Synthesis



Verification vs. Synthesis



Potential Applications

- Low level programming
 - Configuration
 - Bit manipulation
- Programming for end-users
 - Excel
 - Spreadsheet
 - Spark

What is software synthesis?

IEEE TRANSACTIONS ON SOFTWARE ENGINEERING, VOL. SE-5, NO. 4, JULY 1979

Synthesis: Dreams \Longrightarrow Programs

ZOHAR MANNA AND RICHARD WALDINGER

techniques are presented for deriving programs iven specifications. The specifications express the program without giving any hint of the algol. The basic approach is to transform the specificording to certain rules, until a satisfactory pro-

Introduction

N RECENT years there has been increasing activity in the field of program verification. The goal of these efforts is to construct computer systems for determining whether a



Zohar Manna



Richard Waldinger

Synthesis Methods

- Deductive Synthesis
 - Derive the low level implementation from high level implementation from high level specificaion
- Inductive Synthesis
 - Synthesize a program whose behavior satisfies a set of input/output examples

Why now?

- Computer programs are every ware
- Small programs are tricky
- Maturity of underlying technology
 - Machine learning algorithms
 - -SMT
 - Powerful hardware

Motivation



99% of computer users cannot program! They struggle with simple repetitive tasks.

Spreadsheet help forums

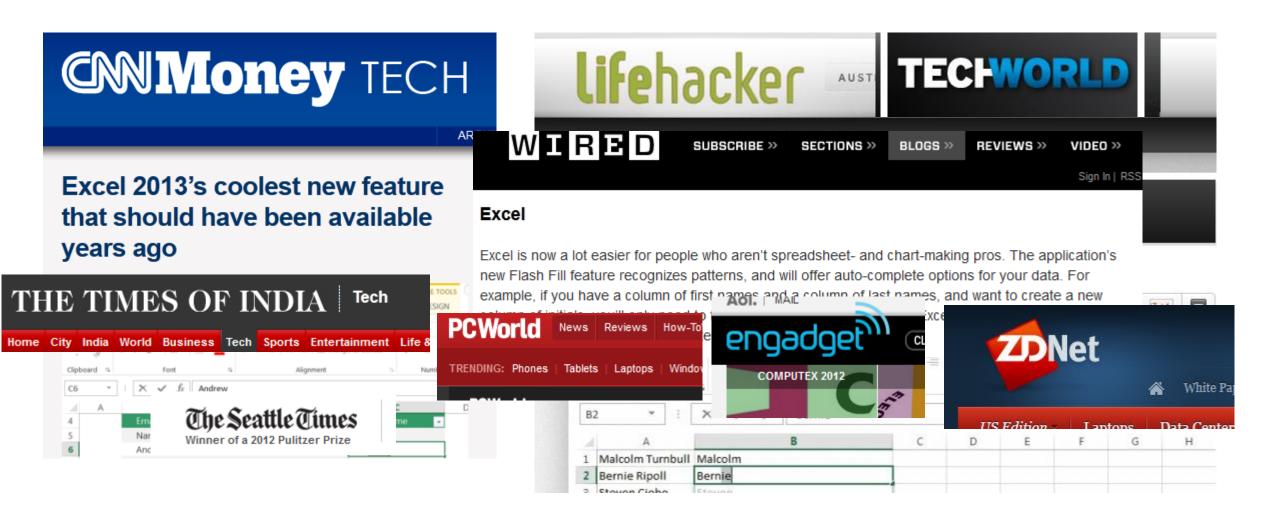




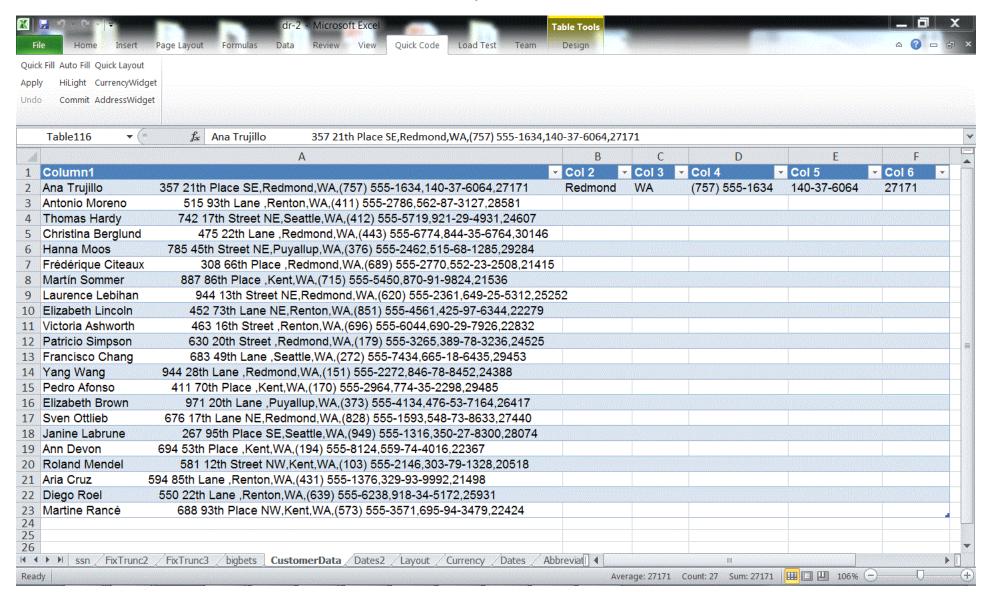




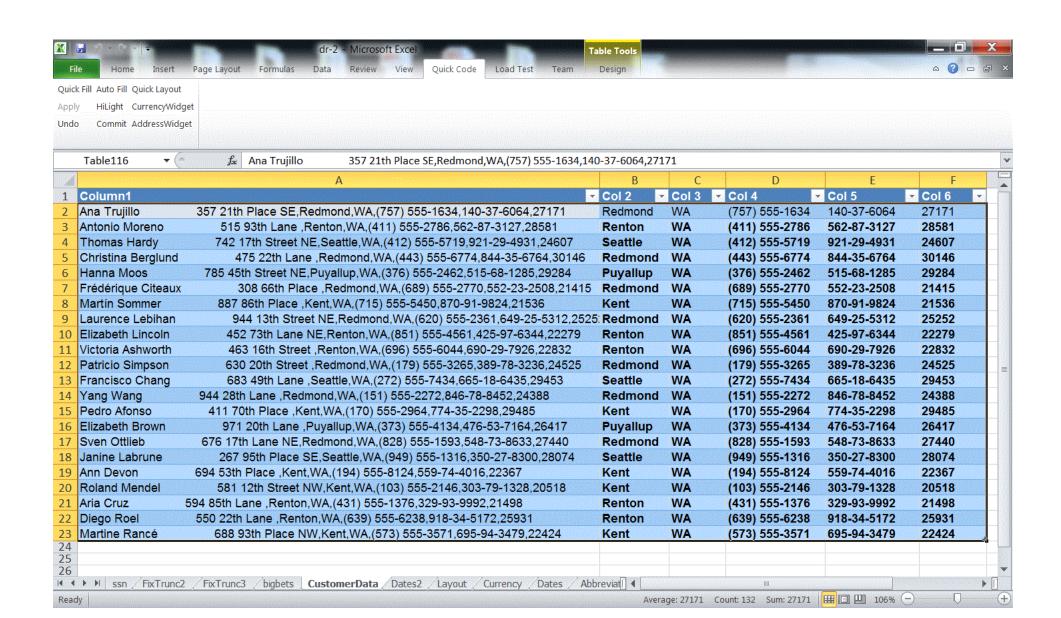
Real world application of synthesis



FlashFill: a feature of Excel 2013 (Sumit Gulwani POPL'11)



FlashFill: a feature of Excel 2013



Inductive Synthesis

Synthesize a program whose behavior satisfies a set of examples

Doesn't machine learning do that?

Traditional Bayesian Machine Learning

- Learn a function from a set of examples
- Scalability is very important, algorithms must scale to millions of data points
- Data is assumed to be noisy;
 - need to avoid overfitting
- Space of possible functions is highly stylized
- Background knowledge incorporated as preprocessing and feature selection

Inductive Synthesis

- Learn a function from a set of examples
- Scalability is not so important, usually we are dealing with small numbers of examples
- Data is assumed to be clean
 - It's annoying when user says f(x)=y and the system assumes the user is wrong and decides that f(x)=z
- Space of possible functions can be arbitrary
- Background knowledge encoded in the description of the space and in the search itself

Functional Synthesis

Goal: Synthesize a function that satisfies a specification

- How do we know the specification has been satisfied? Isn't verification itself already quite hard?
- Can we leverage inductive synthesis machinery for this problem?
- What is the relevant space of functions?
- How do we explore this space efficiently?

The Sketch Synthesis System

Armando Solar-Lezama bit.ly/iptutorial2015



Language Design Strategy

Extend base language with one construct

Constant hole: ??

```
int bar (int x)
{
   int t = x * ??;
   assert t == x + x;
   return t;
}
int bar (int x)
{
   int t = x * 2;
   assert t == x + x;
   return t;
}
```

Synthesizer replaces ?? with a constant High-level constructs defined in terms of ??

Integer Generator >> Sets of Expressions

```
Expressions with ?? == sets of expressions
```

- linear expressions
- polynomials
- sets of variables

```
x*?? + y*??
```

```
x*x*?? + x*?? + ??
```

?? ? x : y

Example: Registerless Swap

Swap two words without an extra temporary

```
int W = 32;
void swap(ref bit[W] x, ref bit[W] y) {
    if(??) \{ x = x ^ y; \} else \{ y = x ^ y; \}
    if(??) \{ x = x ^ y; \} else \{ y = x ^ y; \}
    if(??) { x = x ^ y; }else{ y = x ^ y; }
harness void main(bit[W] x, bit[W] y) {
    bit[W] tx = x; bit[W] ty = y;
    swap(x, y);
    assert x==ty && y == tx;
```

From simple to complex holes

We need to compose ?? to form complex holes

Borrow ideas from generative programming

- Define generators to produce families of functions
- Use partial evaluation aggressively

Generators

Look like a function

but are partially evaluated into their calling context

Key feature:

- Different invocations → Different code
- Can recursively define arbitrary families of programs

Example: Least Significant Zero Bit

0010 0101 → 0000 0010

Trick:

- Adding 1 to a string of ones turns the next zero to a 1
- i.e. 000111 + 1 = 001000

Sample Generator

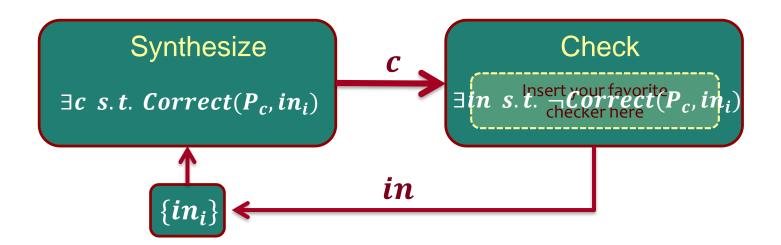
```
/**
 * Generate the set of all bit-vector expressions
 * involving +, &, xor and bitwise negation (~).
 * the bnd param limits the size of the generated expression.
 */
generator bit[W] gen(bit[W] x, int bnd) {
    assert bnd > 0;
    if(??) return x;
    if(??) return ??;
    if(??) return ~gen(x, bnd-1);
    if(??){
        return {| gen(x, bnd-1) (+ | & | ^) gen(x, bnd-1) |};
```

Example: Least Significant Zero Bit

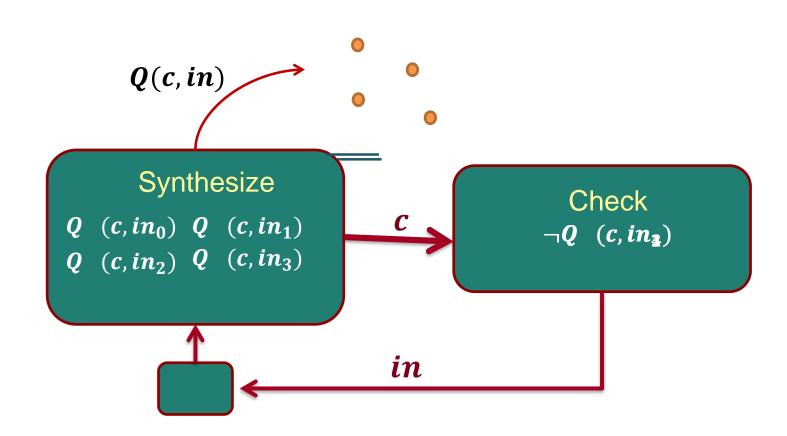
```
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    if(??) return ~gen(x, bnd-1);
    if(??){
        return {| gen(x, bnd-1) (+ | & | ^) gen(x, bnd-1) |};
    }
}
bit[W] isolate0sk (bit[W] x) implements isolate0 {
    return gen(x, 3);
}
```

HOW DOES IT WORK?

CEGIS Synthesis algorithm



CEGIS



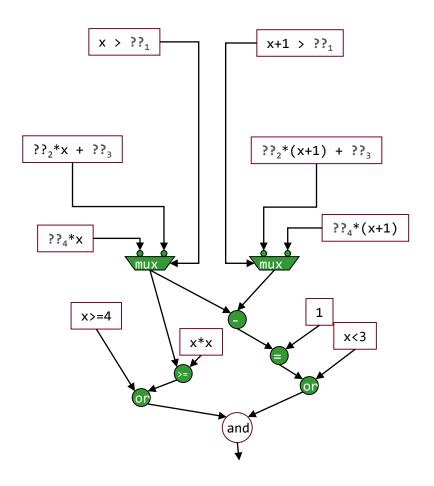
A sketch as a constraint system

```
int lin(int x){
   if(x > ??₁)
     return ??₂*x + ??₃;
   else
     return ??₄*x;
}

void main(int x){
   int t1 = lin(x);
   int t2 = lin(x+1);

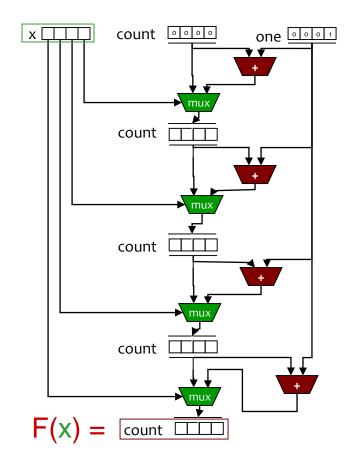
   if(x<4) assert t1 >= x*x;

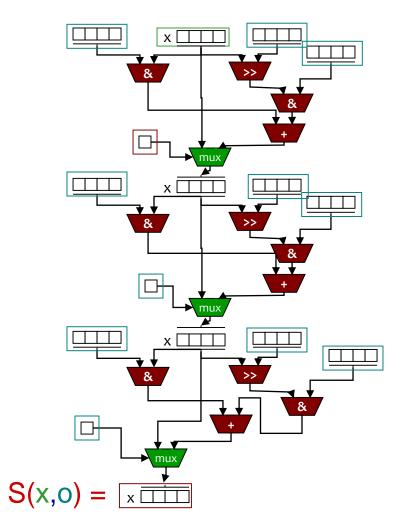
   if(x>=3) assert t2-t1 == 1;
}
```



Ex: Population count. 0010 0110 \rightarrow 3

```
int pop (bit[W] x)
{
    in count = 0;
    for int i = 0; i < W; i++) {
        in (x[i]) count++;
    }
    retarn count;
}</pre>
```





Synthesizing Locks and Fences

Input:

A concurrent program

Assertions

Output:

Locks that guarantee that the assertions hold

```
x = 1;

x = x + 1;

assert x = 3 x = x + 1;
```

Veselin Raychev, Martin T. Vechev, Eran Yahav: Automatic Synthesis of Deterministic Concurrency. SAS 2013: 283-303

Michael Kuperstein, Martin T. Vechev, Eran Yahav: Automatic inference of memory fences. SIGACT News 43(2): 108-123 (2012)

Synthesizing API calls

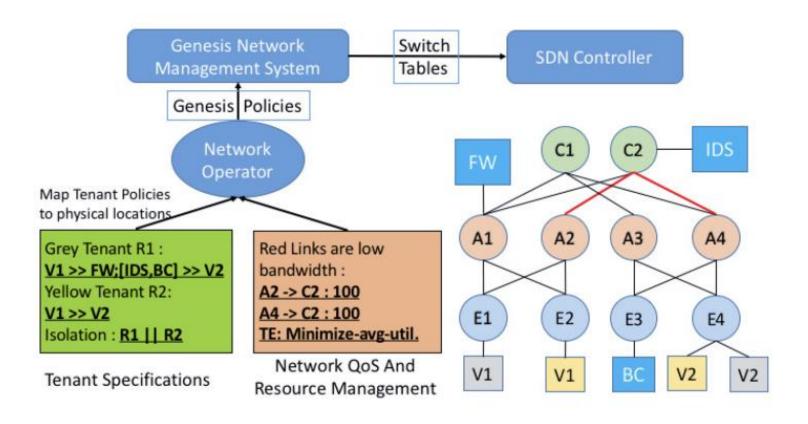
- Modern libraries provide a lot of functionality
- But hard to figure out the right sequence of calls
 - A lot of boilerplate code

```
public void test1()
{ Area a1 = new Area(new Rectangle(0, 0, 10, 2));
  Area a2 = new Area(new Rectangle(-2, 0, 2,
10));
  Point2D p = new Point2D.Double(0, 0);
  assertTrue(a2.equals(rotate(a1, p, Math.PI/2))); }
```

```
Area rotate(Area obj, Point2D pt, double angle)
{ AffineTransform at = new AffineTransform();
  double x = pt.getX();
  double y = pt.getY();
  at.setToRotation(angle, x, y);
  Area obj2 = obj.createTransformedArea(at);
  return obj2;
}
```

Yu Feng, Ruben Martins, Yuepeng Wang, Isil Dillig, Thomas W. Reps: Component-based synthesis for complex APIs. POPL 2017: 599-612 John K. Feser, Swarat Chaudhuri, Isil Dillig: Synthesizing data structure transformations from input-output examples. PLDI 2015: 229-239

Synthesizing Network Policies



Kausik Subramanian, Loris D'Antoni, Aditya Akella:

Genesis: synthesizing forwarding tables in multi-tenant networks. POPL 2017: 572-585

Synthesizing Cloud Configurations

- Number of machines
- Network
- Maximize throughputs
- Minimize cost