

Talos

PROTECTING YOUR NETWORK

Tyler Bohan - @1blankwall1 PacSec 2016



In the Zone: OS X Heap Exploitation





Introduction

- Tyler Bohan
 - Senior Research Engineer
 - Cisco Talos
- Team
 - Aleksandar Nikolich
 - Ali Rizvi-Santiago
 - Cory Duplantis
 - Marcin Noga
 - Piotr Bania
 - Richard Johnson
 - Yves Younan

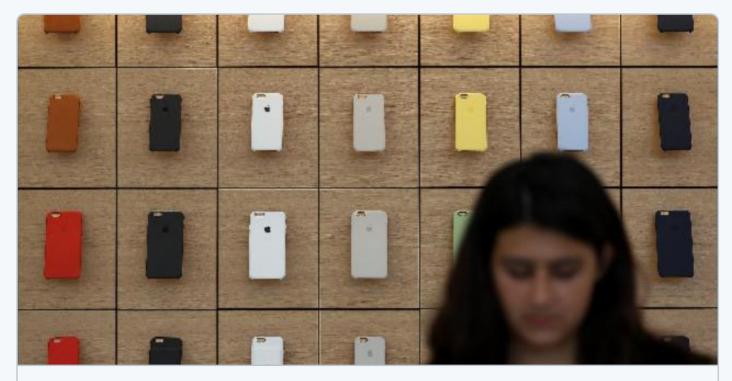
- Talos Vulndev
 - Third party vulnerability research
 - 170 bug finds in last 12 months
 - Microsoft
 - Apple
 - Oracle
 - Adobe
 - Google
 - IBM, HP, Intel
 - 7zip, libarchive, NTP
 - Security tool development
 - Fuzzers, Crash Triage
 - Mitigation development
 - FreeSentry





Richard Johnson @richinseattle · Jul 19

Update iPhones/OSX ASAP Apple patched critical vulns found by @1blankwall1 of our @TalosSecurity vulndev team today!



Apple Lovers Get Patching -- Just One Text Can Steal Your iPhone Pa...

It's the Android Stagefright bugs but for all Apple devices: just one multimedia message could leak passwords.

forbes.com









ImageIO

Available for: iPhone 4s and later, iPod touch (5th generation) and later, iPad 2 and later

Impact: A remote attacker may be able to execute arbitrary code

Description: Multiple memory corruption issues were addressed through improved memory handling.

CVE-2016-4631: Tyler Bohan of Cisco Talos (talosintel.com/vulnerability-reports)

CoreGraphics

Available for: iPhone 4s and later, iPod touch (5th generation) and later, iPad 2 and later

Impact: A remote attacker may be able to execute arbitrary code

Description: A memory corruption issue was addressed through improved memory handling.

CVE-2016-4637: Tyler Bohan of Cisco Talos (talosintel.com/vulnerability-reports)

Why?

- Discovered an image based heap overflow and moved forward with exploitation
- Not much recent documentation about the default malloc implementation on OS X
- Needed to ensure that overflow does not cause any unintentional side effects



Why?

- Heap spraying not reliable or elegant soon to be obsolete
- In depth knowledge of the heap algorithm is important to consistently position chunks relative to one another
- Constantly running into already freed blocks corrupting headers and failing



Prior Work

- Insomnia
 - https://www.insomniasec.com/downloads/ publications/Heaps_About_Heaps.ppt
- Immunity
 - http://www.slideshare.net/seguridadapple/attackingthe-webkit-heap-or-how-to-write-safari-exploits
- Phantasmal Phantasmagoria
- Chris Valasek
 - https://media.blackhat.com/bh-us-12/Briefings/
 ValasekBH_US_12_Valasek_Windows_8_Heap_Internals_Slides.pdf
- Etc.

Prior Work (OS X)

- OS X Heap Exploitation Techniques 2005
 - Nemo
- Mac OS Xploitation (and others) 2009
 - Dino A. Dai Zovi
- Cocoa with Love How Malloc Works on Mac 2010
 - Matt Gallagher



OS X

- Research done on El Capitan version 10.11.5
 - minor changes easy to implement into tooling
 - Currently 10.11.6 so far so good (Sierra beta)
- Scalable Malloc Documented in OS X Hackers Handbook
- Magazine Malloc Default allocator in El Capitan



Roadmap

- Heap Structures Boring (important)
- Heap Strategies Super Fun (unique techniques)
- Exploit Strategies Fun (putting it all together)



Heap Primer(?)







Heap Implementations

- Magazine allocator This one
- Scalable allocator Older OSX allocator
- JEmalloc Jason Evans (bsd)
- DLmalloc Doug Lea (glibc)
- PTmalloc Wolfram Gloger (newer-glibc)
- TCmalloc Webkit's General allocator (google)
- Lookaside List -- Windows (earlier)
- LF Heap -- Windows (recent)
- Etc.

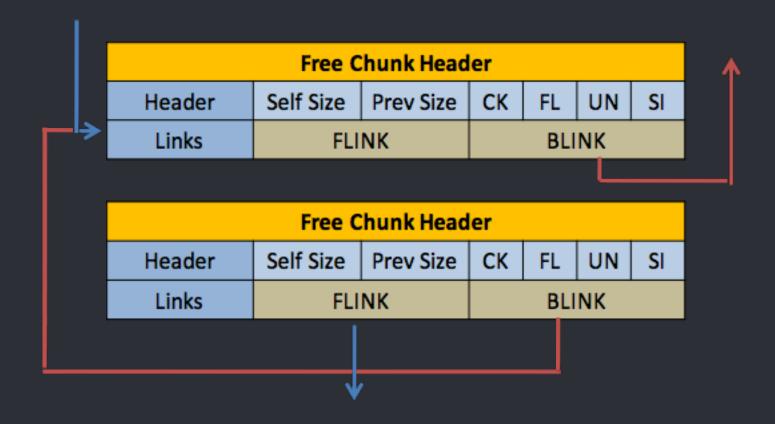


Heap Implementations

- DL Malloc, TC Malloc, JE Malloc, etc...
 - Used to cache or sort chunks that were recently freed
 - Generally a hybrid between a linked-list/metadata and an arena
 - Chunk-based allocators prefix a chunk with metadata headers that allow an allocator to divide or navigate through the different chunks that are available.

Heap Implementation Schemes

Linked List Allocator:



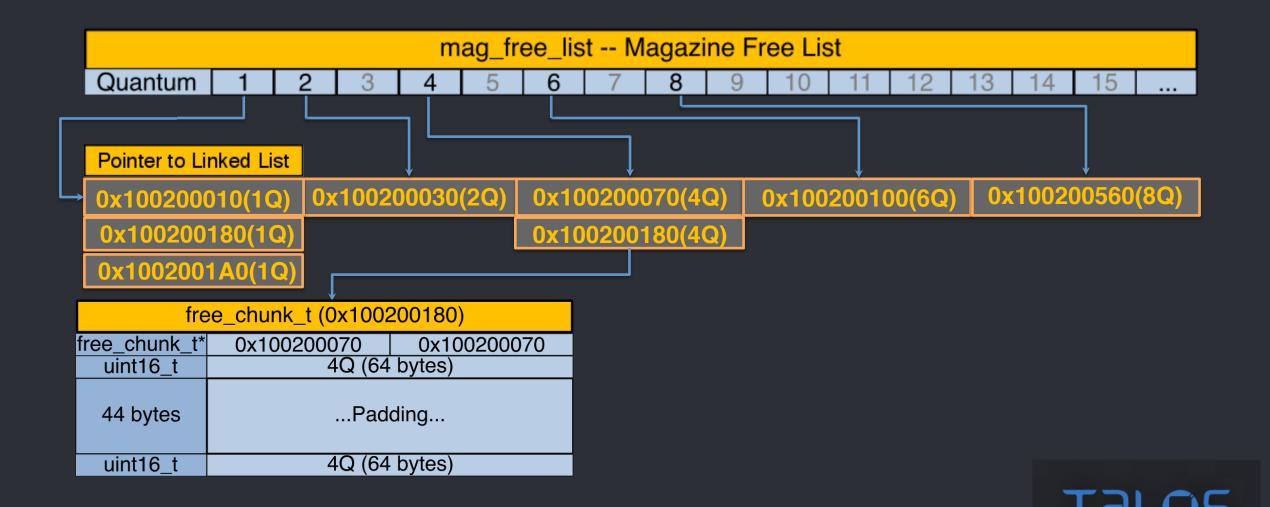


Heap Implementations

- Magazine Allocator, LFH, BSD malloc etc...
 - Arena-based allocators tend to group allocations of the same size on each page
 - Allocated objects efficiently deallocated



Heap Implementation Schemes



OS X Divergence







OS X - CoreFoundation

- Contains various higher-level types for passing arguments to lower-level api.
- These include things such as CFDictionary, CFSet, CFString, etc.
- Referenced and garbage collected pointers: CFRetain/ CFRelease



OS X - CoreFoundation

- Allocates things such as hash table using the system malloc.
- Initialization code allocates onto default heap for setup changing heap state — more on this later
- Heap objects available for overwrite coming from CoreFoundation
- Used by WebKit and Safari for various allocations

- Provided heap utilities
- Guard Malloc
 - Page heap like
 - OpenBSD Allocator original

libgmalloc is used in place of the standard system malloc, and uses the virtual memory system to identify memory access bugs. Each malloc allocation is placed on its own virtual memory page (or pages). By default, the returned address for the allocation is positioned such that the end of the allocated buffer is at the end of the last page, and the next page after that is kept unallocated. Thus, accesses beyond the end of the buffer cause a bad access error immediately. When memory is freed, libgmalloc deallocates its virtual memory, so reads or writes to the freed buffer cause a bad access error. Bugs which had been difficult to isolate become immediately obvious, and you'll know exactly which code is causing the problem.



- Malloc Stack Logging:
 - Requires debug flags MallocStackLogging
 - Alters heap layout
 - malloc_info in LLDB
 - command script import IIdb.macosx.heap
 - malloc_info -s 0x1080715e0

```
(lldb) malloc_info -s 0x00000001080715e0
0x00000001080715e0: malloc( 160) -> 0x1080715e0
stack[0]: addr = 0x1080715e0, type=malloc, frames:
```



- Malloc Stack Logging:
 - malloc_history from the terminal
 - more detailed output
 - all stack frames for malloc and free throughout program history
 - not just current allocation stack frame



Malloc Stack Logging:

```
→ ~ malloc_history 48863 0x00000001080715e0
malloc_history Report Version: 2.0
ALLOC 0x108071550-0x10807160f [size=192]: thread_7fff7a2a5000 | start | NSApplicationMain | -[NSApplication run] | -[NSApplication _nextEventMatchingEventMask:untilDate:inMode:dequeue:] | _DPSNextEvent | _BlockUntilNextEventMatchingListInModeWithFilter | ReceiveNextEventCommon | RunCurrentEventLoopInMode | CFRunLoopRunSpecific | _CFRunLoopRun | _CFRUNLOOP_IS_SERVICING_THE_MAIN_DISPATCH_QUEUE__ | _dispatch_main_queue_callback_4CF | _dispatch_client_callout | _dispatch_call_block_and_release | 0x1000f4dfe | 0x100009513 | 0x100000b0df | -[NSObject(NSThreadPerformAdditions) performSelectorOnMainThread:withObject:waitUntilDone:] | -[NSObject(NSThreadPerformAdditions) performSelector:onThread:withObject:waitUntilDone:modes:] | 0x100000b19e | -[NSWindowController showWindow:] | 0x100009d9a | -[NSWindowController window] | -[NSWindowController windowController windowController windowController windowController windowController windowCont
```

FREE 0x108071550-0x10807160f [size=192]: thread_7fff7a2a5000 | start | NSApplicationMain | -[NSApplication run] | -[NSApplication _nextEventMatchingEventMask:untilDate:inMode:dequeue:] | _DPSNextEvent | _BlockUntilNextEventMatchingListInModeWithFilter | ReceiveNextEventCommon | RunCurrentEventLoopInMode | CFRunLoopRunSpecific | __CFRunLoopRun | __CFRUNLOOP_IS_SERVICING_THE_MAIN_DISPATCH_QUEUE__ | _dispatch_main_queue_callback_4CF | _dispatch_client_callout | _dispatch_call_block_and_release | 0x10000f4dfe | 0x1000009f13 | 0x100000b0df | -[NSObject(NSThreadPerformAdditions) performSelectorOnMainThread:withObject:waitUntilDone: | | -[NSObject(NSThreadPerformAdditions) performSelector:onThread:withObject:waitUntilDone:modes:] | 0x100000b19e | -[NSWindowController showWindow:] | 0x100009d9a | -[NSWindowController _windowDidLoad] | 0x100000b47 | 0x1000014e5d | 0x100001b162 | -[IKImageContentView setImageURL:imageAtIndex:withDisplayProperti









OS X - Allocators

- Malloc Zones
- Manages multiple heap implementations used by an application
- Each zone contains a version and name for identification
- Scalable Allocator (version=5)
- Magazine Allocator (version=8)
- Mapped/Released via mmap/munmap.



- Used by:
 - Anything calling the default CRT malloc
 - libsystem_malloc.dylib, CoreFoundation, etc.
 - Reachable through components such as Preview,
 Color Sync, Safari, Keynote, etc...
- Not Used by:
 - Webkit (minus lower-level things like Objective-c, libxpc, renderers)
 - ImageIO copies implementation internally



- Free list compared to a bucket list on other allocators
- Contains a cookie/checksum for pointers

```
// Because the free list entries are previously freed objects, a misbehaved 
// program may write to a pointer after it has called free() on that pointer, 
// either by dereference or buffer overflow from an adjacent pointer. This write 
// would then corrupt the free list's previous and next pointers, leading to a 
// crash. In order to detect this case, we take advantage of the fact that 
// malloc'd pointers are known to be at least 16 byte aligned, and thus have 
// at least 4 trailing zero bits. 
// 
// When an entry is added to the free list, a checksum of the previous and next 
// pointers is calculated and written to the low four bits of the respective 
// pointers. Upon detection of an invalid checksum, an error is logged and NULL
```



- Each region/arena is bound to a core or "magazine"
- One Magazine per core. One special magazine for free'd regions known as the "depot".
- Allocation sizes are divided into 3 types (tiny, small, large)



- Tiny and Small allocations are managed by a unit of measurement known as a "Quantum" (Q)
- Each magazine has a single-allocation cache for shortlived memory allocations

 i.e.

NSNumber* myNumber = [NSNumber alloc]



- Tiny allocations (Q of 16) maximum size of 1008
- Small allocations (Q of 512) minimum size of 1009
- Large allocations (size defined in zone and is defined based on "hw.memsize" sysctl)



- All allocations come from an arena known as a region
- Free-lists are linked-lists indexed by number of Q
- When a chunk is moved into free-list it is coalesced
- When a region in the depot is empty, region is relinquished back to the OS



Magazine Allocator – Tiny Allocations

- Chunks are handed out by a tiny region.
 - Allocations are from 1Q (16 bytes) to 63Q (1008 bytes)
 - Region is always 0x100000 bytes (0x100 pages) in size
 - Contains 64520 Q-sized blocks for allocations
 - Tiny metadata is described with two bitmaps:
 One describing which block starts a chunk
 Another describing whether the chunk is busy/free



Magazine Allocator – Small Allocations

- Chunks handed out from a small region (Q = 512b)
 - Allocations are from 1Q (1008 / 512b) to szone.large_threshold
 - Region is always 0x800000 bytes (0x800 pages) in size
 - Contains 16320 Q-sized blocks
 - Metadata is a simple list of uint16_t each representing a block and the number of Q to get to the next one
 - The high-bit of each uint16_t describes free/busy



Magazine Allocator – Large Allocations

- Anything larger than szone.large_threshold.
 - * Determined by sysctl hw.memsize being > 230
- Managed by a circular linked list + cache
- Cached until a certain fragmentation threshold is reached.
- When fragmentation threshold is reached, returned back to OS.
- Not elaborated on in this presentation not practical for exploitation purposes



- Thread Migration
 - Threads migrate between cores when user space yields to kernel
 - Kernel can wake up thread on a different core!



- Thread Migration
 - Core determines magazine
 - magazine determines heap layout
 - Precise heap manipulation difficult or impossible if constantly migrating cores



- Thread Migration
 - Patent pending information here



- Thread Migration How to deal?
 - Try to prevent core switching via occupying threads
 - i.e. tight loops etc...



- Thread Migration How to deal?
- 5 iterations available vs blocking
 - available = ~460,000 iterations
 - blocking = ~860,000 iterations



Magazine Allocator Structures





Magazine Allocator

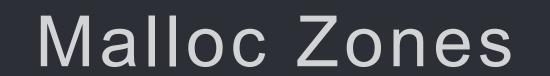
- Tied together by various structures to manage the three different allocation types
- Attempts to keep the regions associated with a magazine tied to it's respective core for any allocations
- Keeps track of an arena and region emptiness for release back to OS
- Tracks large allocations and fragmentation due to reuse of cached large allocations



Magazine Allocator

- Each magazine maintains a single-allocation cache that is used for very short-lived allocations
- Each magazine contains a free-list indexed by Q for quickly identifying free chunks in a region
- Free chunks contain metadata to aid with forwards or backwards coalescing
- Regions contain metadata used to describe how much of the region is in use, and how it's in use









Magazine Allocator – malloc_zones

Zones:

- Version types
 - built in support for multiple different mallocs in one container
- Major feature is extensibility
 - can handle many allocation schemes
 - makes platforms like Safari possible



Magazine Allocator – malloc_zone_t

```
struct malloc_introspection_t {
struct malloc_zone_t {
                                                       funcptr t* check;
   funcptr_t* malloc;
                                                       funcptr_t* print;
                                                       funcptr_t* log;
   funcptr_t* calloc;
                                                       funcptr_t* statistics;
. . .
  const char* zone_name;
   struct malloc_introspection_t* introspect;-
   unsigned version;
```



Magazine Allocator – malloc_zone_t

- General structure used to identify each malloc
- Zone_Name pointer to a string with an identifier
- Version uint32_t identifying allocator type
- Contains function pointers for the different
 - entry points: malloc, free, calloc, etc
 - introspection: size, memalign, pressure_relief, etc
- Szone_t attacked by Nemo in Phrack article



Magazine Allocator – malloc_zone_t

- Phrack digression
 - Szone_t overwrite not practical would need new szone creation
 - Wrap around bug squashed with move to 64 bit systems
 - Double-Free now checked against



Magazine Allocator – szone_t

```
struct szone_t {
  malloc zone t basic zone;
  unsigned debug_flags;
  ...tiny allocations...
  ...small allocations...
  unsigned num_small_slots;
  ...large allocations...
  unsigned is_largemem;
  unsigned large_threshold;
  uintptr_t cookie; _
```

 pointer-sized and is xor'd against pointers that are intended to be obfuscated



Magazine Allocator – szone_t

- Encoded pointers contain a 4-bit checksum within the top 4-bits of the pointer.
- Obfuscated ptr is converted into a checksum ptr via: csum := pointer ^ cookie
- Checksum extracted from pointer via: rol(csum,4) & 0xF



Magazine Allocator – szone_t

```
static INLINE uintptr t
free list checksum ptr(szone t *szone, void *ptr)
    uintptr t p = (uintptr t)ptr;
    return (p >> NYBBLE)
           (free list gen checksum(p ^ szone->cookie)
            << ANTI NYBBLE); // compiles to rotate instruction
  0x100103310 \(^0x233688 = 0x100330598
                                            \rightarrow Checksum(0x100330598) = 1
     ror(0x100103310,4) | Checksum(...) = 100000010010331
```



Magazine Allocator – tiny szone_t

```
struct szone_t {
  size_t num_tiny_regions;
  region_hash_generation_t* tiny_region_generation;
  region_hash_generation_t[2] trg;
                                                struct region_hash_generation_t {
  int num_tiny_magazines;
                                                  size_t num_regions;
                                                  size_t num_regions_shift;
                                                  region_t*[num_regions] hash_regions;
  magazine_t* tiny_magazines;
                                                  region_hash_generation_t* nextgen;
  unsigned num_small_slots; \rightarrow 64
  region_t[64] initial_tiny_regions;
```



Magazine Allocator – small szone_t

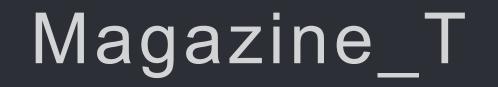
```
struct szone_t {
. . .
  size_t num_small_regions;
region_hash_generation_t* small_region_generation; -
  region_hash_generation_t[2] srg;
                                        struct region_hash_generation_t {
                                         size_t num_regions;
  int num_small_magazines;
                                         size_t num_regions_shift;
                                         region_t*[num_regions] hash_regions;
  magazine_t* small_magazines;
                                        region_hash_generation_t* nextgen;
  unsigned large_threshold;
  region_t[64] initial_small_regions;
```



Magazine Allocator – large szone_t

```
struct szone_t {
unsigned num_large_objects_in_use;
  unsigned num_large_entries;
  large_entry_t* large_entries;
  size_t num_bytes_in_large_objects;
  int large_entry_cache_oldest, large_entry_cache_newest;
  large_entry_t[16] large_entry_cache;
unsigned large_threshold;
```









- One magazine allocated per core. +1 for the "depot".
 - depot used as place holder magazine
 - indexed at slot -1
 - never gets used, only for caching freed regions
- Used to bind regions locally to a core
 - Regions may migrate between magazines only after being relinquished to the depot



- Contains a chunk cache, a free list with bitmap, last region used for an allocation, linked list of all regions
- Monitors some usage stats which are used to manage allocations from the current region



```
struct magazine_t {
  void* mag_last_free;-
                                       Local magazine cache
  region_t mag_last_free_rgn;
  free_list_t*[256] mag_free_list;
                                                1000000001100111000110000111
  unsigned[8] mag_bitmap;
  size_t mag_bytes_free_at_end, mag_bytes_free_at_start;
  region_t mag_last_region;
  unsigned mag_num_objects;
  size_t mag_num_bytes_in_objects, num_bytes_in_magazine;
  unsigned recirculation_entries;
  region_trailer_t* firstNode, *lastNode;
                                                        Linked list of all regions
```

- Statistics
 - Mag_bytes_free_at_Start
 - Mag_num_objects
 - number of contiguous slots available free and busy

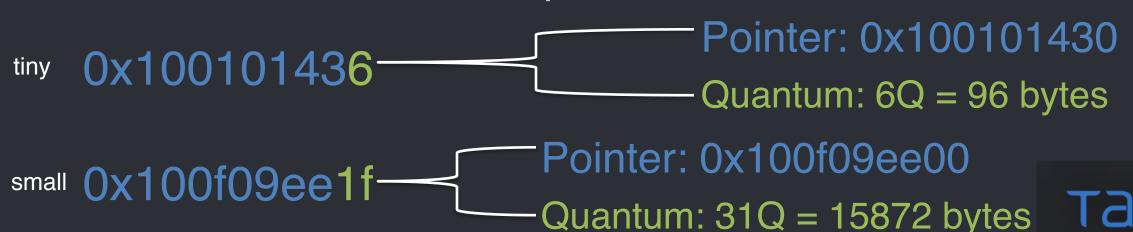
Region

mag_bytes_free_at_start

mag_bytes_free_at_end

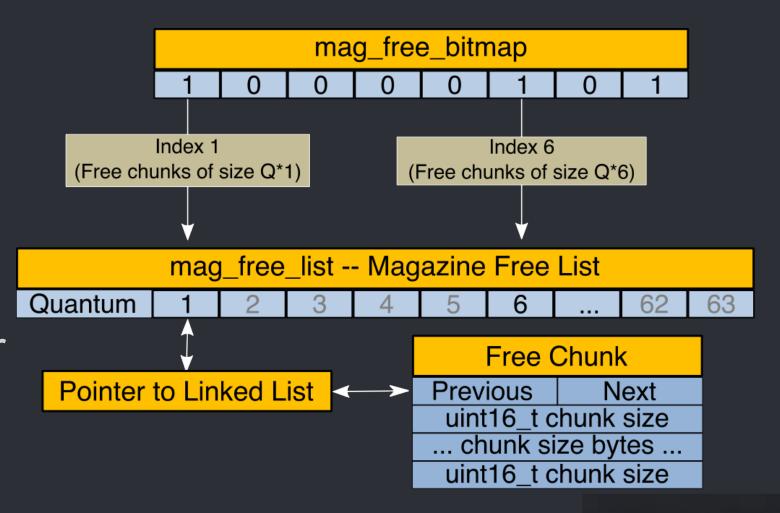


- mag_last_free cache Short-lived allocations
 - Quantum size encoded along with pointer.
 - Tiny Q=16, lower 4-bits represent quantum size
 - Small Q=512, lower 9-bits represent quantum size
 - Rest of bits contain pointer with chunk.



Mag_free_bitmap
Represents which
free-list entries are
populated

Mag_free_list
Pointers to a circular
linked list of a free
chunk.





- Magazine Free-list
 - Points to a circular-linked list of all free-chunks.
 - Next and Free pointers are checksummed.
 - Checksum is encoded in the top 4-bits of pointers.
 - Pointer is bit-rotated to the right by 4-bits, and then Or'd with the 4-bit checksum.
 - Checksum not checked if chunk is being coalesced with.

Next Pointer in Chunk

Actual Pointer







- When a chunk moves from cache to free-list
 - Circular linked list written to chunk.
 - The size (in Q) is written at the end of chunk if chunk is > 1Q.
 - Used for forwards and backwards coalescing

free_chunk_t			
free_chunk_t*	Previous Pointer	Next Pointer	
uint16_t	Forwards Coalesce Size		
	Padding		Size*Q bytes
uint16_t	Backwards C	oalesce Size	

- Free list contains 256 slots only 64 in use.
 - Defined by szone.num_small_slots
- Leads to some dead-space due to 1008 maximum.
 - NUM_SMALL_SLOTS = 64



- Free-list uses all 256 slots available for it
- based on large threshold
- typically 0x1fc00 / 512 == 254









Magazine Allocator – region_t

- Two different types for "tiny" and "small" allocations.
- Chunks are composed of a number of chunks split on
 Q: Tiny = 16 bytes, Small = 512 bytes.
- Contains a number of Q-sized blocks that is calculated based on the size of the region + metadata + region trailer



Magazine Allocator – region_t

- Metadata contains information describing chunk sizes and whether chunk is busy or free.
- Tiny metadata maps each bit in header/in-use to the region's blocks
- Small metadata maps each uint16_t directly to a block



Magazine Allocator – region_t

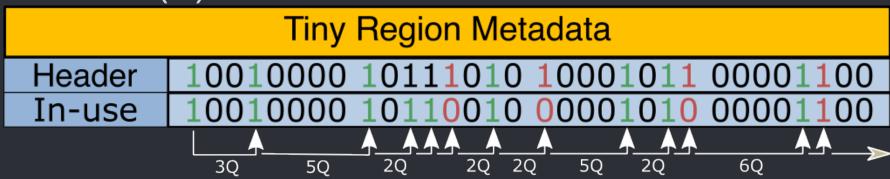
```
typedef uint8_t[Q] Quantum_t;
                                        struct small_metadata_t {
                                          uint16_t[NUM_BLOCKS] msize;
struct region_t {
 Quantum_t[NUM_BLOCKS] blocks;
 region_trailer_t trailer;
                                        struct tiny_header_inuse_pair_t {
 metadata_t metadata;
                                           uint32 t header, inuse;
struct tiny_metadata_t {
tiny_header_inuse_pair_t[NUM_BLOCKS / sizeof(uint32_t)] blocks;
```

* Metadata correlates directly with blocks in region_t



Magazine Allocator – tiny region_t

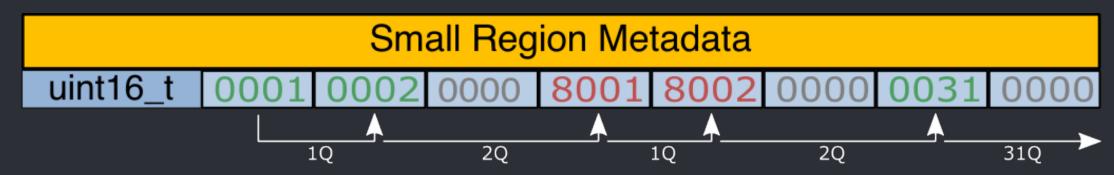
- Tiny Metadata Array of structures with two uint32_t
 - 1. Bit represents whether block is beginning of a chunk. (Used to calculate chunk sizes in Q).
 - 2. Bit represents whether that entire chunk is busy(1) or free(0)





Magazine Allocator – small region_t

- Small Metadata Array of encoded uint16_t
 - High-bit represents whether busy(0) or free(1)
 - Rest of bits describe the number of Q for chunk
 - Q represents how many elements to skip to get to next chunk.





Magazine Allocator – large regions

- Large Region
- Simple. Just an address and it's size.
- doesn't get unmapped until heavily fragmented
- stored in cache

```
typedef struct large_entry_t {
     vm_address_t address;
     vm_size_t size;
     boolean_t did_madvise_reusable;
}
```



Magazine Allocator – region_trailer_t

- Each region knows which magazine it's associated with
 - To navigate, a magazine_t points to the trailer, allows a magazine to iterate through all regions











Strategies – for cache

- Tool we call Mac Heap
- LLDB python script allowing access to all this information
- Open Source





Strategies

- Positioning
 - Strategic block placement, cache management
- Overwriting
 - Block placement before or after free and busy chunks
- Metadata
 - Unlinking attacks



Strategies – for cache

- Cache
 - Emptying cache
 - Malloc exact size as cached object
 - Cache to free list
 - free another object less than 256b(tiny) in size
 - moves new object to cache and cache object to free list



Strategies – for free list

- Free-List
 - Linked List
 - Indexed by Quantum Size
 - Tiny 64 slots Small 256 slots
 - Small minimum size 1009 cant use 1 Q in size due to inability to allocate
 - Q = 512 512 goes into tiny region 1 Q allocations are dead



Strategies - mag_free_list

- Coalescing When chunk is moved to free-list
 - ForwardQ/BackwardQ Only written to chunks >1Q
 - If overwriting ForwardQ or BackwardQ, need to ensure that specified count is pointing to a free chunk.
 - This will add entire chunk (including any used chunks in between) to free list.



Strategies - mag_free_list - Coalesce

Overwrite chunk with 2Q + 3Q bytes data. Set BackwardQ to 12 (3+2+7).

Free Chunk
Previous Next
Forward 3Q

Backward 3Q

Busy Chunk
5Q

Freeing Busy chunk will coalesce with 3Q Chunk and use overwritten BackwardQ of 12 joining 2Q chunk with 7Q chunk.

Due to BackwardQ being 12. Freeing 5Q chunk will read BackwardQ and add 7Q+2Q+3Q (12Q) to free list whilst chunk still being by program.



Strategies - mag_free_list

- Forwards Coalesce When realloc() or moving from cache to free list.
 - Looks in ForwardQ of free-chunk to determine how far to coalesce or how far to realloc in-place...
 - When done, gets added to free list.



Strategies - mag_free_list - Coalesce

Overwrite Busy with 2Q data to get to Free. Write 4+4 (8Q) into Free's ForwardQ.

If Busy Chunk gets free'd, 2Q+4Q+4Q (10Q) gets added to free list. 3Q still in use by program.

Free Chunk
Previous Next
Forward 4Q

Backward 4Q

Similarly if Busy Chunk gets realloc'd, The different of 8Q and the new allocation size gets added to the free list where 3Q chunk is still in use.



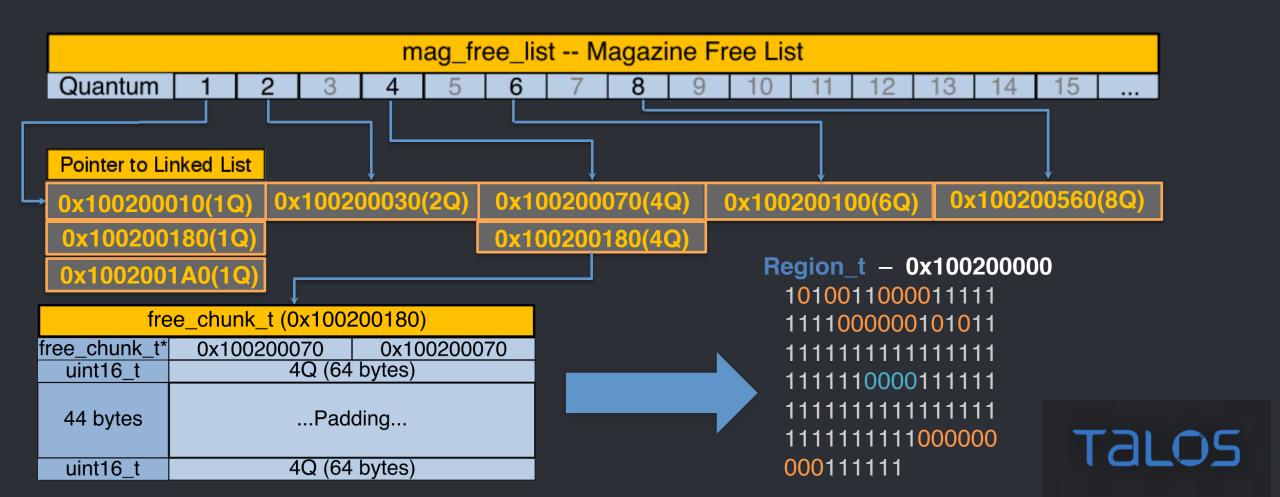
Strategies - mag_free_list

- Overwrite linked list at beginning of freed chunk
 - Write 4 with unlink not super useful ASLR 7% odds of getting correct check sum
 - If linked list pointers are NULL, then checksum evaluates to 0. Unfortunately, this clears the bit in mag_bitmap.
 - Need to move another chunk to free-list to set bit in bitmap, to re-gain access to coalesced entry.

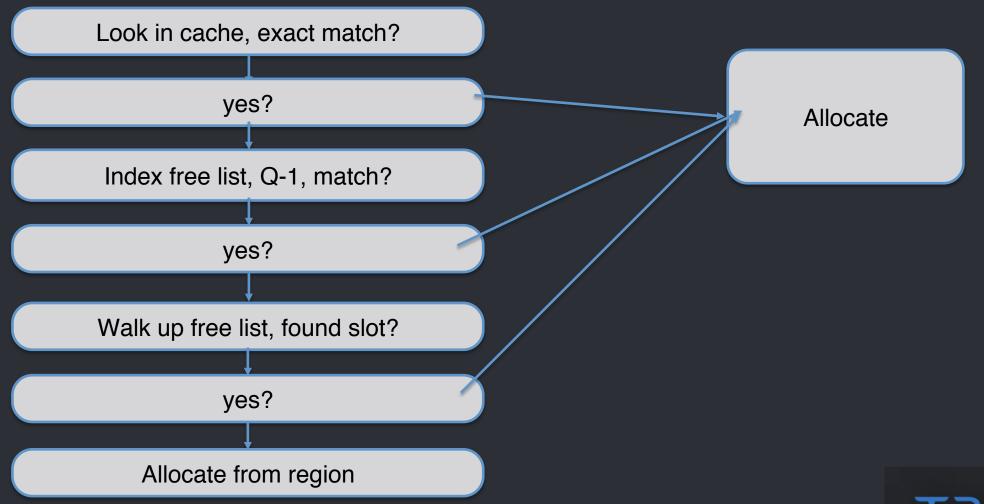


Strategies - mag_free_list

Allocating from the free list - 4Q

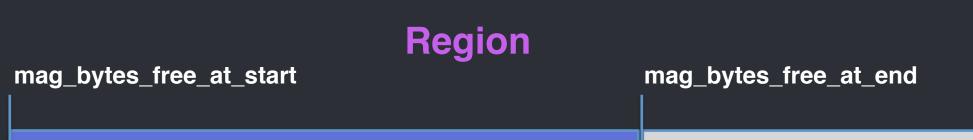


Strategies – for free list





- Allocate from region
 - calculate bytes currently in use
 - add to region base address
 - allocate space needed
 - free'd goes back on the free list





Demo

Single Core Script

```
script import onemag
breakpoint set -N singlemag1 -s libsystem_malloc.dylib -a 0x262a
breakpoint set -N singlemag2 -s libsystem_malloc.dylib -a 0x2a6a
breakpoint set -N singlemag3 -s libsystem_malloc.dylib -a 0xb95d
breakpoint set -N singlemag4 -s libsystem_malloc.dylib -a 0xbd43
breakpoint command add -F onemag.cont singlemag1
breakpoint command add -F onemag.cont singlemag2
breakpoint command add -F onemag.cont2 singlemag3
breakpoint command add -F onemag.cont2 singlemag4
```

- Region Bitmap
 - Bitmap fitting
 - fill up single quanta slots inside of region bitmap to grow region bitmap
 - slot next allocation at end of in use region
 - allocations slotted together

```
>>>region.getoffset()
0x100200000
>>>region.bitmap()
```

```
Fragmented Bitmap
In Use = 1 Free = 0
```

```
>>> print magazine.dump()
[1] 0x100202eb0
[2] 0x100202ed0
[4] 0x100202f10
[6] 0x100202fb0
[8] 0x100203070
[13] 0x100203550
[22] 0x1002032b0
```

```
>>>region.getoffset()
0x100200000
>>>region.bitmap()
```

```
>>>t['mag_free_list'].quantum()
56
```

make 56 single quantum allocations to flatten bitmap

Fragmented Bitmap In Use = 1 Free = 0

region.bitmap()

>>> len(region.bitmap())
901

next allocation will grow the bitmap (subsequently the region) the necessary quantum

Fragmented Bitmap In Use = 1 Free = 0

region.bitmap()

>>> len(region.bitmap())

Strategies - Scoping bugs(?)

- Double Free
 - Not possible when freeing due to explicit checks against metadata and cache



Strategies - Scoping bugs(?)

- Heap Spraying
 - Blocks at beginning of boundary (0x100000 tiny) (0x800000 - small)
 - 0xXXXFC080 through 0xXXXFFFFF tiny structures
 - 0xXX800000 through 0xXXFF8000 small structures



Debugging Strategies







- LLDB lacks functionality
 - viewing page protections
 - dumping memory
 - expression handling
 - etc



- Overly verbose commands (limited functionality)
 - breakpoint set -a `(void()) main`
 - breakpoint command add -o "x/x \$rax" 1



- LLDB lacks functionality
- Difficult or cumbersome for certain commands
- Python functionality inside of LLDB init is quite cumbersome



- Created generic way of adding functions
- Alias and breakpoint managers
- Full expression parser built in to allow more WinDBG like commands, breakpoint management etc...
- ie. poi(main+\$rax+local_1)



Exploit Strategies







Safari

- 5 malloc zones
 - Default Malloc version 8 (magazine alloc)
 - GFX Malloc version 8 (magazine alloc)
 - WebKit Malloc version 4 (bmalloc)
 - Quartz Core Malloc version 5(scalable alloc)
 - Default Purgeable Malloc version 8(magazine alloc)



Safari

• 5 malloc zones

MALLOC ZONE	VIRTUAL SIZE	ALLOCATION COUNT	BYTES ALLOCATED	% FULL	REGION COUNT
DefaultMallocZone_0x102226000	72.0M	16490	3466K	4%	5
WebKit Malloc_0x7fff7bf4ce68	12.0M	3	12.0M	100%	7
QuartzCore_0x7fa735009600	100K	767	46K	46%	10
GFXMallocZone_0x10245c000	Ø K	0	ØK		0
DefaultPurgeableMallocZone_0x108036000	0K	0	0K		0
					
TOTAL	84.1M	17260	15.4M	18%	22



Vulnerability Details

- Image based heap overflow
- Core Foundation -> ImageIO
- Default malloc zone
- ImageIO has its own built in malloc (ImageIOMalloc) it's barely used :)



- Heap based overflow
- Primitives needed:
 - Information leak
 - Object to overwrite
- And here comes the zone problem...



- Majority of controllable allocations fall into the WebKit malloc
- However, calls to New go into the default zone
 - opens up quite a few objects to overwrite
- Information Leak still needed



- Heap massaging primitives found via Javascript's interaction with CoreFoundation
- AudioContext object makes reliable allocations in default malloc zone
- Multiple images needed to properly position heap, due to CoreFoundations and initialization and setup allocating to the default heap - CSS used for proper loading

More reversing reveals one string allocated into our default zone

- Date.prototype.ToLocale
- Overwrite able via our overflow and still accessible via JavaScript



- Date.prototype.ToLocale ... 3-quantum allocation
- Empty a quantum after our locale
- Metadata overwrite -> backwards coalesce
- Get chunk freed -> clear the region bitmap
- Next allocation positions us in the middle of the locale string
- Remove the null terminator
- Read the date



- Information leak obtained
- Overwrite object
- Code execution





Conclusion

- https://github.com/blankwall/MacHeap
 - MacHeap tool
 - libsystem_malloc.idb
 - https://github.com/blankwall/LLDBInit
- @1blankwall1

