

Lecture 8 : Computational Approaches to Lyapunov functions

Recap :- DP view of the world - minimizing some cost-to-go.

$$\int_0^{\infty} L(x, u) dt \quad , s.t. \quad \dot{x} = f(x, u)$$

Cond. for Continuation

$$J^*(x) \Rightarrow 0 = \min_u \left[L(x, u) + \frac{\partial J}{\partial x} f(x, u) \right]$$

$$\frac{dJ^*}{dt} = -L(x, u^*) \quad (\text{cost-to-go must go down the hill at exactly the rate of the loss when we are taking optimal controller})$$

Lyapunov view of the world

→ Instead of find a J that's going down the hill at exactly the above rate $(-L(x, u^*))$

find a function $V(x) \geq 0$

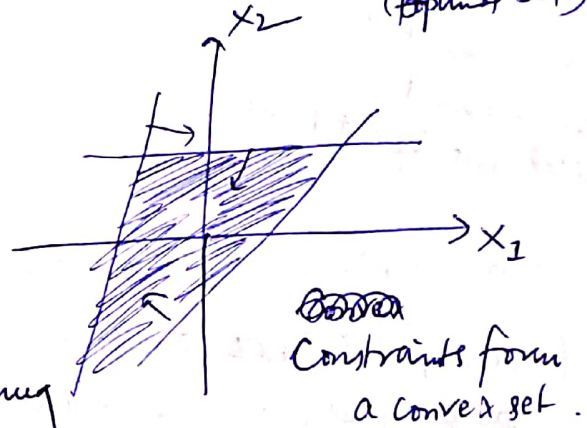
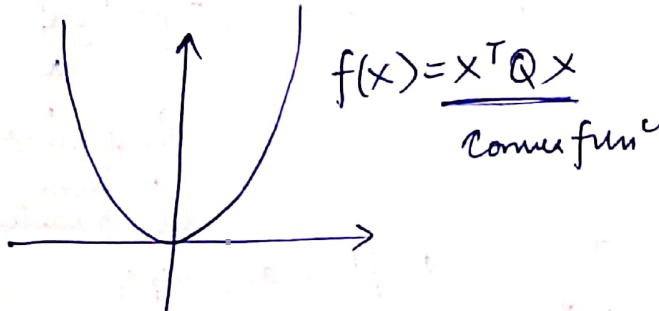
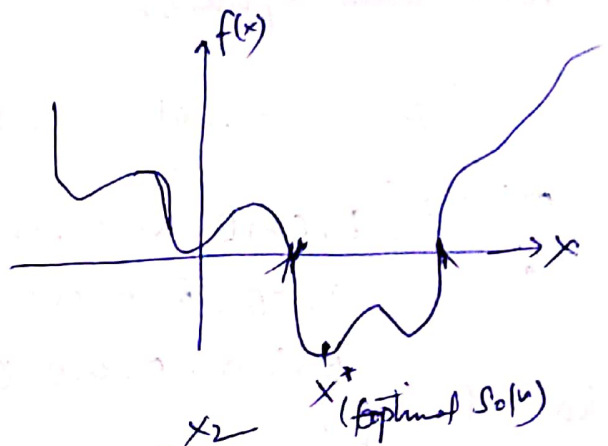
$$s.t. \quad \frac{dV}{dt} \leq 0 \quad (\text{going down the hill at all})$$

(There are more functions that will satisfy this inequality constraints and because of that it would be easier to search over them)

↓
eventually but not exactly

Optimization Crash Course

$\min_x f(x)$
decision variable x subject to $\forall x, g(x) \leq 0$
scalar objective



When the objective function is convex and constraints form a convex set, then we are in the land of convex optimization.

Linear Programming (LP)

Linear cost function, Linear constraints

$$\min_x C^T x$$
$$\text{s.t. } Ax \leq b$$

Second order cone programming (SOCP)

Quadratic Programming (QP)

$$\min_x \sqrt{Q} x$$
$$\text{s.t. } Ax \leq b$$

If Q_x is +ve definite, then it's a convex quadratic programming but still have linear constraints.

Semidefinite programming (SDP)

$$\min_x C^T x$$
$$Ax \leq b$$

$P(x) \geq 0 \leftarrow$ Semidefinite Constraint

Pendulum dynamics

$$ml^2 \ddot{\theta} + b \dot{\theta} + mgl \sin \theta = 0$$

$$\phi(x) = \begin{bmatrix} 1, \cos \theta, \sin \theta, \dot{\theta}, \\ \cos^2 \theta, \sin \theta \cos \theta, \\ \dot{\theta} \sin \theta, \dot{\theta} \cos \theta, \dot{\theta}^2 \end{bmatrix}$$

↑
basis of
the trigonometric
polynomial
that we expect to see

$$V(x) = \alpha_0 + \alpha_1 \cos \theta + \alpha_2 \sin \theta + \alpha_3 \dot{\theta} +$$

find α_i 's ~~such~~ (coeffs. of
above polynomials) such
that the conditions ~~are~~ hold.

$$V(x_i) > 0 \quad V(x_i) \geq \epsilon x_i^T x_i \quad \leftarrow \epsilon^{-3}$$

Drake Tutorial - for diff. kind
of programming

Search over a class of functions
that could be Lyapunov functions

Use ~~DP~~ linear function approximates

$$V(x) = \sum_i \alpha_i \phi_i(x) = \alpha^T \phi(x)$$

$$\dot{V}(x) = \frac{\partial V}{\partial x} f(x) = \left(\alpha^T \frac{\partial \phi}{\partial x} \right) f(x)$$

(still linear
in those
sample points)

Idea - Sample many x_i 's

find α , so that $\forall_i V(x_i) \geq 0$

$$\dot{V}(x_i) \leq 0$$

How can we certify that Lyapunov condition are true $\forall x$?
→ what if we choose a different basis function.

$$V(x) = \sum_i \alpha_i \phi_i^2(x), \quad \alpha \geq 0 \quad (\text{constraint})$$

$$\geq \phi^T(x) \begin{bmatrix} \alpha & & \\ & \alpha & \\ & & \alpha \end{bmatrix} \phi(x) \quad (\text{matrix form})$$

Nice generalization → $\phi^T(x) G \phi(x), G \geq 0$

(+ve semidefinite
matrix)
(convex set)

$$\forall x, x^T G x \geq 0$$

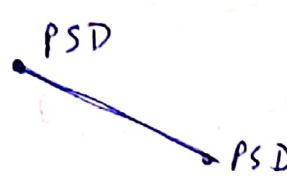
G is +ve semidefinite



PSD - positive semi-definite matrix

$$x^T G_1 x \geq 0 \Rightarrow G_1 \succeq 0$$

$$x^T G_2 x \geq G_{22} \Rightarrow G_2 \succeq 0$$



$$\alpha G_1 + (1-\alpha) G_2 \succeq 0$$

$$\forall \alpha \in [0, 1]$$

All of the positive semi-definite matrix that can be made by the convex combination of them (end points) are also PSD.

So, $\alpha x^T G_1 x$
the set of PSD is a convex set.

Note - If instead of writing a bunch of x_i 's $V(x_i) \geq 0$, we would have parametrised $V(x)$ like below

$$V(x) = \sum_i \alpha_i \phi_i^2(x), \alpha_i \geq 0$$

$$= \phi^T(x) \begin{bmatrix} \alpha & & \\ & \alpha & \\ & & \alpha \end{bmatrix} \phi(x) \quad \forall x, x^T G_1 x \geq 0$$

$\phi^T(x) G_1 \phi(x)$ and just add the constraint $G_1 \succeq 0$, then that guarantees for all x 's, V is positive definite function.

No sampling required.

What about $\dot{V}(x) = \frac{\partial V}{\partial x} f(x) = -P(x)$? → basis function

$$P(x) = m^T(x) G_2 m(x), G_2 \succeq 0$$

P for polynomial

when $P(x)$ is polynomial, this requires finite linear constraints.

$$P(x) = 1 + 2x^2 \geq 0$$

(Sum of squares decomposition of a polynomial)

$$P(x) = 2 - 4x + 5x^2 = \begin{bmatrix} 1 & x & x^2 \end{bmatrix} \begin{bmatrix} G_{11} & G_{12} & G_{13} \\ -G_{12} & G_{22} & G_{23} \\ -G_{13} & -G_{23} & G_{33} \end{bmatrix} \begin{bmatrix} 1 \\ x \\ x^2 \end{bmatrix}$$

Is $P(x)$ positive for all x ?

If $G \geq 0$ (Positive definite) and if above equality holds, then $P(x)$ is positive $\forall x$.

For polynomial, coeff. matching is linear and can be written as linear constraints & is sufficient

$$G_{11} = 2$$

$$G_{12} + G_{21} = -4$$

(square terms = 5)
coeff.

If we can find a matrix G constrained to match above polynomial $P(x) = 2 - 4x + 5x^2$ such that $G \geq 0$ (positive ^{semi}definite), then $2 - 4x + 5x^2$ must have been positive.

→ General optimizing over polynomials using 'sum of squares' decomposition trick is called 'sum of squares optimization' (SOS).

$$\dot{x} = Ax \quad \text{is it stable?}$$

$$V(x) = x^T P x, \quad P > 0$$

$$\dot{V}(x) = x^T A^T P x + x^T P A x < 0$$

$$x^T (A^T P + P A) x \quad A^T P + P A < 0$$

find a P ~~matrix~~ such that P is positive definite and $\dot{V}(x)$, i.e., $A^T P + P A$ is negative definite

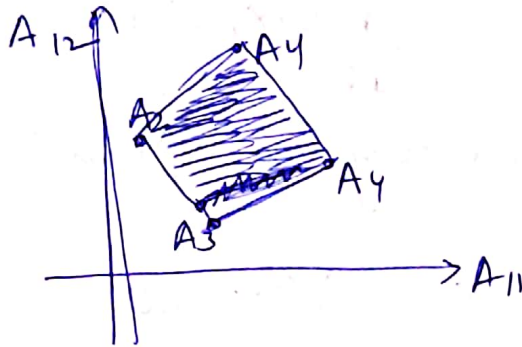
$$\underline{A^T P + P A = -P_2}, \quad P_2 > 0 \quad (\text{Positive semidefinite Constraint})$$

(Linear constraint)

Robust stability (for linear systems) via Common Lyapunov func

$$\dot{X} = AX$$

$$A = \sum_i \beta_i A_i, \quad \beta_i \geq 0, \quad \sum_i \beta_i = 1$$



Find a P

s.t. $P > 0$

$$V(x) = x^T P x$$

$\forall i,$

$$P A_i + A_i^T P < 0$$

when the optimisation fails in the land of convex optimisation it means there is no solution to the problem. There is no quadratic common Lyapunov function for the system

if this is true, then

$$P(\sum \beta_i A_i) + (\sum \beta_i A_i)^T P < 0$$

(anything in the convex hull is also going downhill)

$$\dot{X}_1 = -X_1 - 2X_1^2$$

$$\dot{X}_2 = -X_2 - X_1 X_2 - 2X_2^3$$

$$V(x) = \phi^T(x) G \phi(x)$$

$$\dot{V}(x) < 0$$

$$G = \begin{bmatrix} x_1^4 \\ x_2^4 \\ x_1^2 x_2^2 \\ x_1^2 x_2^2 \\ x_2^4 \end{bmatrix} \text{ (symmetric)}$$

find a G such that so that

$$V(x) = \phi^T(x) G \phi(x) \text{ is true}$$

$$\text{and } \dot{V}(x) < 0$$

find a Lyapunov func that

satisfies any optimality

certify that optimality we can do this