

Lecture 10: Lyapunov/Sum-of-Squares for Control.

$$V(x) > 0 \quad \dot{V}(x) < 0$$

$$\frac{\partial V}{\partial x} f(x)$$

If we choose Lyapunov candidates which are polynomials and our dynamics is polynomial (i.e., if $V(x)$ polynomial $f(x)$ polynomial)

Then Sum-of-squares (SOS) optimization.

But, Eqn of motion

Simple pendulum

$m l^2 \ddot{\theta} + b \dot{\theta} + m g l \sin \theta = 0$ is not polynomial in the original coordinates.

$\sin \theta_i \rightarrow s_i$ (new decision variable)

$\cos \theta_i \rightarrow c_i$

In order to make sure, they are related, add another constraint

$$s_i^2 + c_i^2 = 1$$

they are polynomials but they are not actually polynomials in original coordinates

Note
The eqn we get out of our mechanical systems have sines and cosines that enter in a particular way. we can get $\sin(\theta_1 + \theta_2)$ which we can break

into polys. of \sin & \cos of θ_1 or θ_2 via trigonometry.

But, we will never get something like $\sin \theta_1$ or $\sin \theta_2$, etc. then trigonometric terms are direct functions of one variable which is θ .

These special structures allows us to do above substitution $\sin \theta_i \rightarrow s_i$ $\cos \theta_i \rightarrow c_i$

In particular
Coordinate System $x = \begin{bmatrix} s \\ c \\ \dot{\theta} \end{bmatrix}$

$$\dot{x} = f(x) = \begin{bmatrix} c \dot{\theta} \\ -s \dot{\theta} \\ \frac{b \dot{\theta} + m g l s}{m l^2} \end{bmatrix} \text{ (still polynomial in } x \text{ coordinates)}$$

Screw joints are not polynomial because they mix translation and rotation in the same coordinate.

find (P, α) s.t. \rightarrow decision variables

$$V = m^T(x) P m(x)$$

$$-\dot{V}(x) + \lambda(x)(s^2 + c^2 - 1)$$

monomial is a basis of $[1, s, c, \dots]$

is SoS.

x are indeterminant

another polynomial

$V(0) = 0$ (to set the scaling).

S-procedure

$$g(x) \leq 0 \Rightarrow \text{---}$$

$$g(x) \geq 0 \Rightarrow \text{---}$$

(s, c) are not decision variable, they are indeterminant (they are never actually handed to the solver)

* A decision variable has to take on a particular value at the optimal solution whereas an indeterminant is a quantity we want to hold for all values — undetermined value / indeterminant value.

$s^2 + c^2 = 1$ is not a constraint. It ~~do~~ makes the problem easier in some sense but it's a language that says when $s^2 + c^2 = 1$ that implies $\dot{V}(x)$ is -ve

$$M(q)\ddot{q} + C(q, \dot{q})\dot{q} = \tau_g(q) + \beta \dot{u}^0$$

$M(q), C(q, \dot{q}), \tau_g$ are polynomial in \dot{q}, s, c

$$\dot{x} = f(x) = \begin{bmatrix} \dot{q} \\ M^{-1}(q)[\tau_g(q) - C \dots] \end{bmatrix}$$

inverse makes it a rational polynomial

on the scalar can

$$\frac{P_1(x)}{P_2(x)}$$



$$\dot{x} = f(x)$$

more general description
of system of equations

an explicit way to write eqn of motion
when the derivation is an explicit
function of the state.

$$g(x, \dot{x}) = 0 \Rightarrow f(x) - \dot{x} \quad (\text{implicit form of dynamics})$$

$$M(q)\ddot{q} + C(q, \dot{q})\dot{q} = \tau_j(q) + B\ddot{u}^0$$

(implicit form of dynamics)

$$\dot{x} = f(x) = \begin{bmatrix} \dot{q} \\ M(q)[\tau_j(q) - C \dots] \end{bmatrix}$$

explicit form
of dynamics

$$g(q, \dot{q}, \ddot{q}) - M(q)\ddot{q} - C - \tau_j - B\ddot{u}^0 = 0 \quad (\text{implicit form of dynamics})$$

Lyapunov cond. in implicit form

$$V(x) \text{ is SoS}$$

Indeterminates x, z $-\frac{\partial V}{\partial x} z$ is SoS
↑
placeholder for \dot{x}

Quotient Ring

finally
reveals the
fundamental
structure of the eqn
of motion in our
optimization.

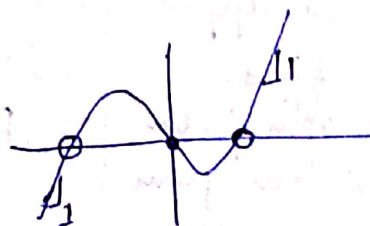
$$g(x, z) = 0 \Rightarrow -\frac{\partial V}{\partial x} z \text{ is SoS.}$$

$$-\frac{\partial V}{\partial x} z + \lambda^T(x) g(x, z) \text{ is SoS.}$$

polynomial
of same
fixed degree

Example $\dot{x} = \frac{-x + x^3}{1 + x^2}$

Rational polynomial
form.



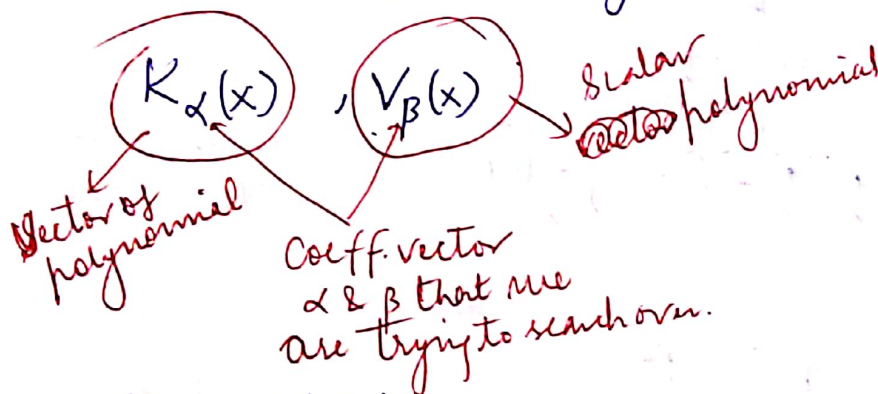
when $g(x, z) = 0$
 $\Rightarrow \frac{\partial V}{\partial x} z \geq 0$

Lyapunov Control (we are going to design controllers and Lyapunov function that proves it's stable)

$$\dot{X} = f_1(x) + f_2(x)u \quad (\text{dynamics are control affine})$$

find a controller $u = K(x)$ (polynomial in x), $V(x)$

that makes above system satisfy the Lyapunov ~~good~~ conditions in closed loop.



find $V_\beta(x) > 0$
 α, β

$$\dot{V} = \frac{\partial V_\beta}{\partial x} [f_1(x) + f_2(x)K_\alpha(x)] < 0 \quad \text{is SOS requires linear in decision variable.}$$

Non-Convex Constraint (bad)
since this will have terms like $\alpha_1 \beta_2, \dots$ which ~~are not~~ are not jointly convex in α & β .

we want this to be a polynomial linear in decision variables α & β .

Above problem is Bilinear in decision variables

$$\alpha_i \beta_j (\dots)$$

↑
indeterminants

A standard approach would be ~~to~~ alternation. (if one controller is fixed, we know that the optimization is convex for looking for Lyapunov functions and similarly if one Lyapunov func is fixed, the problem finding good K 's that satisfy that is convex in parameters of K .)

$\left[\begin{array}{l} \text{Fix } V_\beta(x), \text{ optimize } K_\alpha(x) \\ \text{Fix } K_\alpha(x), \text{ optimize } V_\beta(x) \end{array} \right]$
 (repeat until converge)

if α is fixed, then remaining term is convex in β and vice-versa.

- can have local minima.
- Recursive feasibility (once we find a Lyapunov/Controller, they will only improve).
- Monotonic improvement

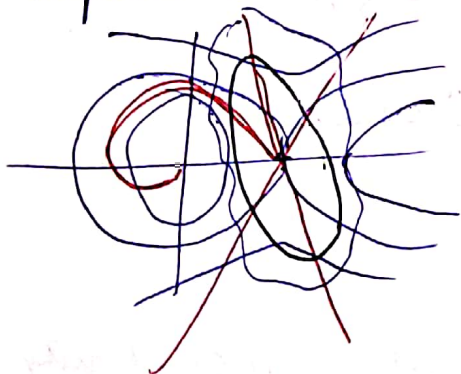
If we start with an initial K and optimize K ,

one initial guess at K is already a feasible solution

So, optimizing we can only do better at each of these steps.

✓ we can get a guess at K , try to find a better Lyapunov function maybe one that goes downhill faster, for instance we can only do better at each of these steps.

★ Hard part is we have to find an initial Lyapunov function that is feasible. once we have that we can expect monotonic improvement until convergence -



powerful recipe

- Linearize nonlinear dynamics at some fixed point x^*, u^*
- Solve LQR: gives $u = -Kx$,

If we can compute $\hat{J} = x^T S x$ (cost-to-go)

RoA using SoS of $u = -Kx$ Controller with $\hat{V} = x^T S x$ Lyapunov guess

Now start Alternation to search $K(x), V(x)$ to prove bigger RoA.

Offline step

$$u = Kx + \alpha x^2 + \beta x^3$$

Sum-of-squares DP

Recall L.P. DP (Linear Programming for DP)

Discrete setting: $\forall s \quad J(s) = \min_a [l(s, a) + J(f(s, a))]$

\uparrow
discrete state

\Downarrow L.P. formulation

$$\forall s, \forall a \quad J(s) \leq l(s, a) + J(f(s, a))$$

$$\max_s \sum_s J(s)$$

New in Continuous time & state action

$$\forall s \quad 0 = \min_u \left[l(s, u) + \frac{\partial J}{\partial x} f(x, u) \right]$$

$$\forall x, \forall u \quad 0 \leq l(x, u) + \frac{\partial \hat{J}}{\partial x} f(x, u)$$

$$\Rightarrow l(x, u) + \frac{\partial \hat{J}}{\partial x} f(x, u) \text{ is SOS.}$$

$$\max \int_x \hat{J}(x) dx \quad \hat{J} \Rightarrow \text{estimate}$$

$$\begin{cases} \dot{x} = -x + x^3 \text{ (origin is already stable, so not an interesting control problem)} \\ \dot{x} = x - x^3 \text{ (origin is unstable)} \end{cases}$$

HJB vs Lyapunov

$$\dot{J} = l(x, u^*)$$

$$\dot{V} \leq 0$$

now,

$$\dot{J} \leq -l(x, u^*) \quad \text{(Not a convex optimisation)} \quad \text{upper bound on cost-to-go and certify Lyapunov}$$

But, what we did above $\dot{J} \geq -l(x, u^*)$ — (cont.)

(Convex optimisation) minimised u can be replaced by all u . (lower bound)

