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ANUBHAV SINGH

IC-2KL6 - 54

IC - 703

ARTIFICIAL INTELLIGENCE LAB

MCA 9th Sem

Anubhan
29th Jan. 2021

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Record of Work Done



1. Represent the following facts in Prolog :

a. Tom is a singer.
singer (Tom).

b. Mia and John are married.
married (Mia, John).

c. Rani is dead.
dead (Rani).

d. John kills everyone who gives Mia a footmassage.
kills (John, X) :-
footmassage (X, mia).

e. Mia leaves everyone who is a good dancer.
leaves (Mia, X) :-
gooddancer (X).

f. Rahul eats anything that is nutritious or tasty.
eats (Rahul (X)) :-
nutritious (X),
tasty (X).



2. Suppose we are working with the following knowledge base:

hasWand(harry).

quidditchPlayer(harry).

wizard(ron).

wizard(X) :- hasBroom(X), hasWand(X).

hasBroom(X) :- quidditchPlayer(X).

How does Prolog respond to the following queries?

wizard(ron)

true

wizard(hermione)

false

wizard(harry)

true

wizard(Y)

Y = ron ;

Y = harry.

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Q. Define and describe the difference b/w knowledge, belief, hypothesis and data.

i. Knowledge: It can be defined as the body of facts and principles accumulated by humankind or fact or state of knowing.

ii. Belief: It is defined as essentially any meaningful and coherent expression that can be represented. Thus, a belief can be true or false.

iii. Hypothesis: It is defined as justified belief that is not known to be true. Thus a hypothesis is a belief that is backed up with some supporting evidence, but it may still be false.

In other words, it is preliminary assumption or tentative explanation that accounts for a set of facts, taken to be true for the purpose of investigating and testing.

iv. Data: Data in computer terminology means new facts and figures for eg. 'Anubhan', '201654', etc are data. Data are processed to form information.

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4. Write Prolog clauses for the following mgt system.

i. Edward manages Smith & Jackson.

manages (Edward , Smith),
manages (Edward , Jackson).

ii. Smith manages Jones , Harris & Peters.

manages (Smith , Jones),
manages (Smith , Harris),
manages (Smith , Peters).

iii. Harris manages Carter & Fletcher.

manages (Harris , Carter),
manages (Harris , Fletcher).

iv. Jackson manages Harper, Pitchard & Glover.

manages (Jackson , Harper),
manages (Jackson , Pitchard),
manages (Jackson , Glover).



	PERSON	RESPONSIBLE FOR
a.	Jones	Finance
b.	Carter	Sales
c.	Fletcher	Purchases
d.	Peters	Personnel
e.	Harper	Transport
f.	Pritchard	Communications
g.	Glover	Industrial Relations
responsible (Jones, finance)		
responsible (Carter, Sales)		
responsible (Fletcher, Purchase)		
responsible (Peters, Personnel)		
responsible (Harper, Transport)		
responsible (Pritchard, Communication)		
responsible (Glover, Industrial-Relations)		
A person responsible for a work if he manages someone who is responsible for that work.		
responsible (A, B) :-		
manages (A, A),		
responsible (A, B).		

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Answer the following question:

1. Is Smith responsible for transport.
responsible (Smith, transport).
False

2. Who are responsible for communication.
responsible (X, communication).
 $X =$ Pitchard
 $X =$ Edward
 $X =$ Jackson
False

3. Who are responsible for sales & personal.
responsible (X, sales).
 $X =$ Carter
 $X =$ Edward
 $X =$ Smith
 $X =$ Harris
False
responsible (X, personal).
 $X =$ Peters
 $X =$ Edward
 $X =$ Smith
False

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5. Implement the following problems in Prolog.

a. The Farmer, Wolf, Goat & Cabbage.

$\Rightarrow \text{change}(e, w)$

$\Rightarrow \text{change}(w, e)$

$\text{move}([x; x, Goat, Cabbage], wolf [y; y, Goat, Cabbage]) :-$
 $\quad \text{change}(x, y).$

$\text{move}([x, wolf, x, Cabbage], goat [y, wolf, y, Cabbage]) :-$
 $\quad \text{change}(x, y).$

$\text{move}([x, wolf, goat, x], cabbage, [y, wolf, goat, y]) :-$
 $\quad \text{change}(x, y).$

$\text{move}([x, wolf, goat, cabbage], nothing, [y, wolf, goat, cabbage]) :-$
 $\quad \text{change}(x, y).$

$\text{oneEq}(x, x, -).$

$\text{oneEq}(x, -, x).$

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safe([Man, Wolf, Goat, Cabbage]) :-

oneEq(Man, Goat, Wolf),

oneEq(Man, Goat, Cabbage).

solution([e, e, e, e], []).

solution(config, [FirstMove | otherMoves]) :-

move(config, Move, NextConfig),

safe(NextConfig),

solution(NextConfig, otherMoves).

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2. The Monkey & the Banana :-

domains :

$x = \text{string}$

predicates

$\text{take}(x, x)$

$\text{move}(x, x)$

$\text{get-on}(x, x)$

$\text{hit}(x, x, x)$

/* go clause */

go :-

$\text{take}(\text{"Monkey"}, \text{"Stick"}),$

$\text{move}(\text{"Monkey"}, \text{"Chair"}),$

$\text{get-on}(\text{"Monkey"}, \text{"Chair"}),$

$\text{hit}(\text{"Monkey"}, \text{"Stick"}, \text{"Banana"}),$

$\text{write}(\text{"The monkey has hit the banana"}).$

go :-

$\text{write}(\text{"The monkey could not reach to the banana"}).$

$\text{take}(\text{Animal}, \text{Object}) :-$

$\text{write}(\text{"Does the, "Animal", take the, " Object, ?"} (\text{Y/N})),$

$\text{readchar}(\text{Reply}),$

$\text{Reply} = \text{'Y'}$.

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move (Animal, Object) :-

write ("Does the", Animal, " moves the", Object, "? (Y/N)"),

readchar (Reply),

Reply = 'y'.

geton (Animal, Object) :-

write ("Does the", Animal, "Get on", Object, "? (Y/N)'),

readchar (Reply),

Reply = 'y'.

hit (Animal, Object, Fruit) :-

write ("Does the", "Animal", "hit the", "Fruit", "with",
"the", Object, "? (Y/N)'),

readchar (Reply),

Reply = 'y'.

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c. Waterjug problem.

domain

$x, y = \text{Integers}$.

predicates

$f(x, y)$

clauses.

$f(2, 4) :- \text{write}(\text{"\\n } 2 \text{ o \\n"}),$
 $f(x, 2) :- \text{write}(\text{"\\n } 0 \text{ o } 2 \text{ \\n"}).$

$f(x, 4) :- y=0, x < 4!,$
 $\text{write}(\text{"\\n"}, x, \text{'3'}),$
 $f(x, 3).$

$f(x, y) :- x=0, y < 3!,$
 $\text{write}(\text{"\\n } 4 \text{ ", } y),$
 $f(4, y).$

$f(x, y) :- x+y >= 4; y > 0, x < 4, !,$
 $y_1 = y - 4 + x.$

$\text{write}(\text{"\\n } 4 \text{ ", } y_2),$
 $f(4, y_2); x+y >= 4, y > 0, y < 3, !;$
 $y_2 = y - 4 + x,$
 $\text{write}(\text{"\\n } 4 \text{ ", } y_1),$
 $f(4, y_1).$

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$f(x,y) :- x+y >= 3, x > 0, x < 4, !,$

$$X_1 = X - 3 + Y,$$

write ("\\n", X₂, " 3"),

$f(X_2, 3), X+Y >= 3, X > 0, Y < 3, !,$

$$X_1 = X - 3 + Y$$

write ("\\n", X₂, " 3"),

$f(x, 3).$

$f(x, y) :- x+y <= 4, x > 0, !,$

$$X_1 = X + Y,$$

write ("\\n 0", X₁, " 0"),

$f(x, 0).$

$f(X_2, y) :- X+Y <= 3, X > 0, !,$

$$Y_1 = X + Y$$

write ("\\n 0", Y₁),

$f(0, Y_1)$

$f(X_3, Y) :- X > 0, !,$

write ("\\n 0", Y)

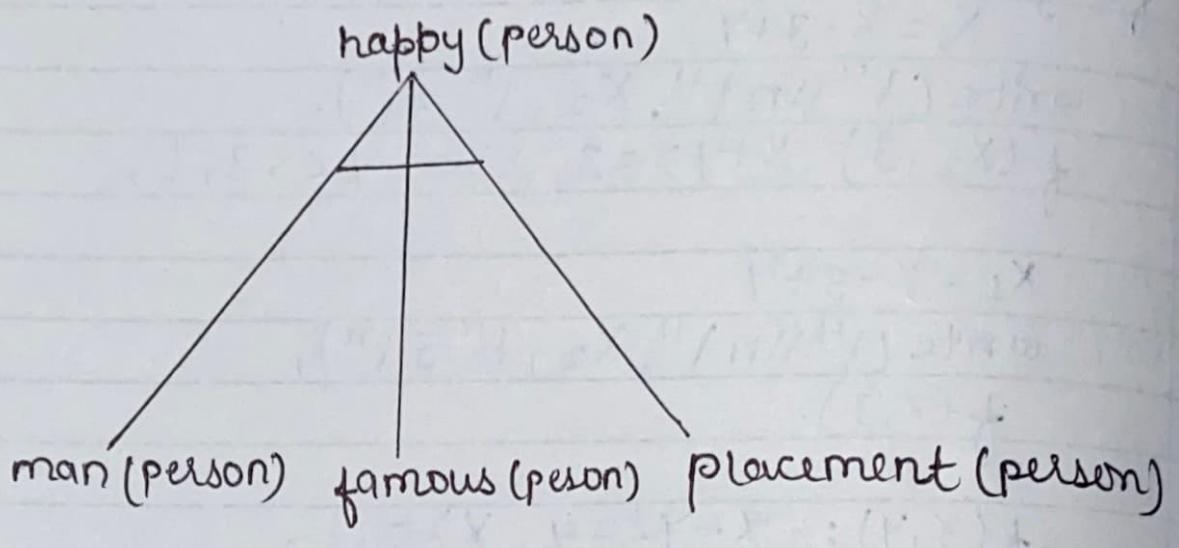
$f(0, Y)$

$f(x, y) :- Y > 0, !,$

write ("\\n", X, " 0"),

$f(x, 0) -$

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AND Tree

(fig 1)

1. Explain the rules of conjunction & disjunction with a example & draw AND/OR tree.

Rule of conjunction:

A man is happy if he is rich & famous might be transtated to :-

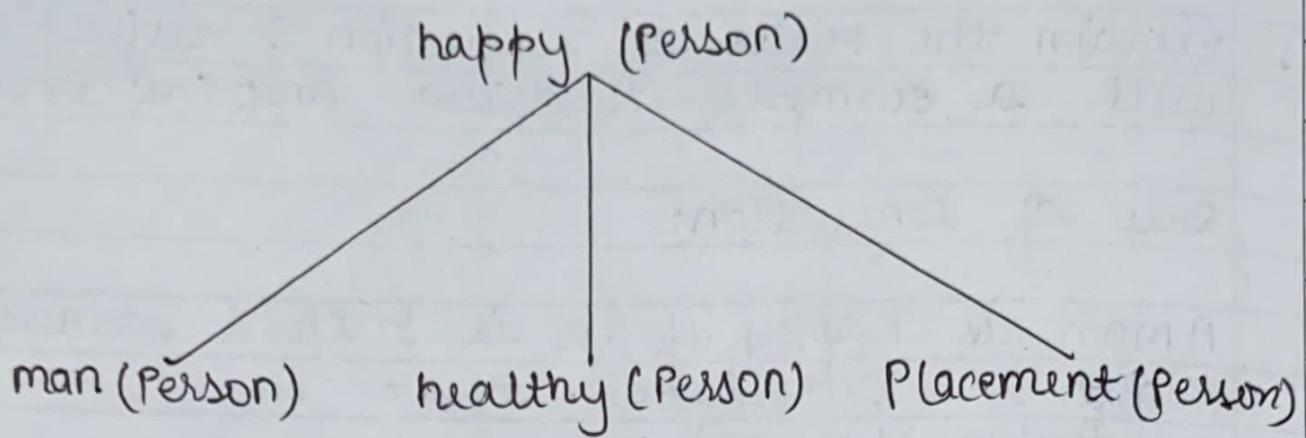
happy (Person) :-

man (Person),
famous (Person),
placement (Person).

The ',' indicates the conjunction & is roughly equivalent to a predicate calculus.

Therefore read ',' as 'AND'. Subgoals are cojoined.

We can represent conjunction using an AND tree. Here is an AND tree that represents above example. The way in which we delineate b/w OR tree & AND tree, is the use of horizontal bar to link the rule goal. We need this distinction because we are going to separate the structure of a program using a compiled AND/OR tree.



OR Tree

(fig 2)



Rule of Disjunction

Someone is happy if he is man, healthy or has placement.

Translate to:

happy (Person) :- man (Person)

happy (Person) :- healthy (Person)

happy (Person) :- Placement (Person)

Note that we have rewrite these stnts.

Someone is happy if they are men. OR

Someone is happy if they are wealthy OR

Someone is happy if they have placements.

The predicate name "Happy" is known as function.

The function happy has one argument.

We describe a predicate with name

"predname" with argument 'n' as predname

In. It has one argument so we can say that, its arity is 1.

The predicate happy| \perp is defined by 3 clauses.

The AND/OR tree is represented as given



2. Explain logical variables. Explain the diff b/w logical variable & variables used in procedural language.

In prolog, logical variable enable you to write general facts & ask general questions. To represent a variable in prolog, captitiathe first letter of it.

Ex:- likes ("Rama", x) if likes ("bill", x).
The object x is a logical variable.
It might be able to match anything that Bill likes.

Logical variable v/s variable

Prolog has no assignment attachment, this is a significant distinction b/w them.

- i. Variables in prolog get their value by being matched to constants in facts or rules, you can't store information in logical variables, unlike procedural language.
- ii. Anomous variables can be used in place of any other logical variables.

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3. Write short note on following :-

a. First order predicate logic :-

First order predicate logic is a powerful language that develops info about the object in a more easy way & can also express relationship b/w the objects. first order logic didn't assume that the world contains facts like propositional logic. world contains facts the following:-

i. Object

ii. Relations

iii. Functions

As a natural language, first order logic also has two main parts:-

i. Syntax

ii. Semantics.

Basic elements of first order logic are:-

i. Constant

ii. Variable

iii. Predicate

iv. Connecting equality.

b. Propositional & predicate logic:-

Propositional logic is the study of propositions where a proposition is a statement that is either true or false. Proposition logic may be used to encode simple arguments that are expressed in natural language. It uses truth table & inference rule.

Predicate rules logic allows complex facts world to be represented & new fact may be determined via deduction reasoning predicate calculus include predicate variable & quantities.

c. Forward chaining vs backward chaining

Forward chaining starts from the known fact & move forward by applying inference rule to extract more data & it continues until it reaches to the goal, whereas backward by using inference rule to determine the fact that satisfy the goal.

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- Forward chaining is data driven whereas backward chaining is goal driven.
- Forward chaining uses BFS strategy whereas backward chaining used DFS.
- d. Unification : Unification is a general method of matching a query to a fact each of which may contain variables. In order to determine contradictions, we need a matching procedure that compare 2 literals & discourses whether there exists a set of substitutions that makes them identical. This is unification Algorithm.
- e. Fail Predicate : If we want to force a rule to fail under certain condition Then we built in predicate called failure used in prolog. Predicate fail takes prolog interpreter to fail a predicate goal & subsequently force batch tracking.



f. Input Predicate: Prolog can read & write terms from one or several files. These thus form the stream of data that are either input or output.

The input predicates are:

read(x) : read one clause from the current input & unify it with x . If there is no further input, x is unified with end of file.

see (file) : Open file as the current input file.

seen : Close the current input file.

g. Backtracking: Backtracking is basically a form of searching. Prolog handles the computation that may have more than one results by backtracking i.e. undoes all work done since a tentative choice was made as an alternative choice can be made or tried. If a particular clause can be satisfied more than once, we know that prolog will go back & try to find all of those solutions.

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h. Computable functions & predicates:

It may be necessary to compute functions as part of fact. In this case a computable predicate is used. A computable predicate may include computable function such as +, -, *, etc. For ex:-

$gt(x-y, 10) \rightarrow bigger(x)$ contains the computable predicates gt which perform the greater than function. Note this computable predicates uses the computational function subtraction.

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4. Assume you have database --- ex :-

beat (rami, rakesh)

beat (tom, rami)

beat (jerry, tom)

You define a relation

class (Player, category)

that ranks player into categories,
you just have three categories.

winner: Every player who all his
games is a ^{won} winner.

fighter: Any player that won some
games & lost her games.

sportsman: Every player that has
lost all his matches.

Qd: To specify the rules of a fighter.

x is a fighter if

There is some y such that x
beats y &

There is some z such that z
beats x.

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New rules for the winner:-

x is a winner if

x beat some y &

y was not beaten by anybody.

The rules for sportsman are:-

x is a sportman if

there is a y who beats x &

no one is beaten by x

The formulation contains not which can't be directly expressed with our present protog facilities, so the formulation of winner appears trickster. The same problem occurs with sportsman. The problem can be circumvented by combining the definition of winner with that of fighter & we the otherwise, connective. Such a formation is:

If x beats somebody & x was beaten by somebody, then x is a fighter.

otherwise, if x beats somebody then x is a winner

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otherwise if x got beaten by somebody
then x is sportman.

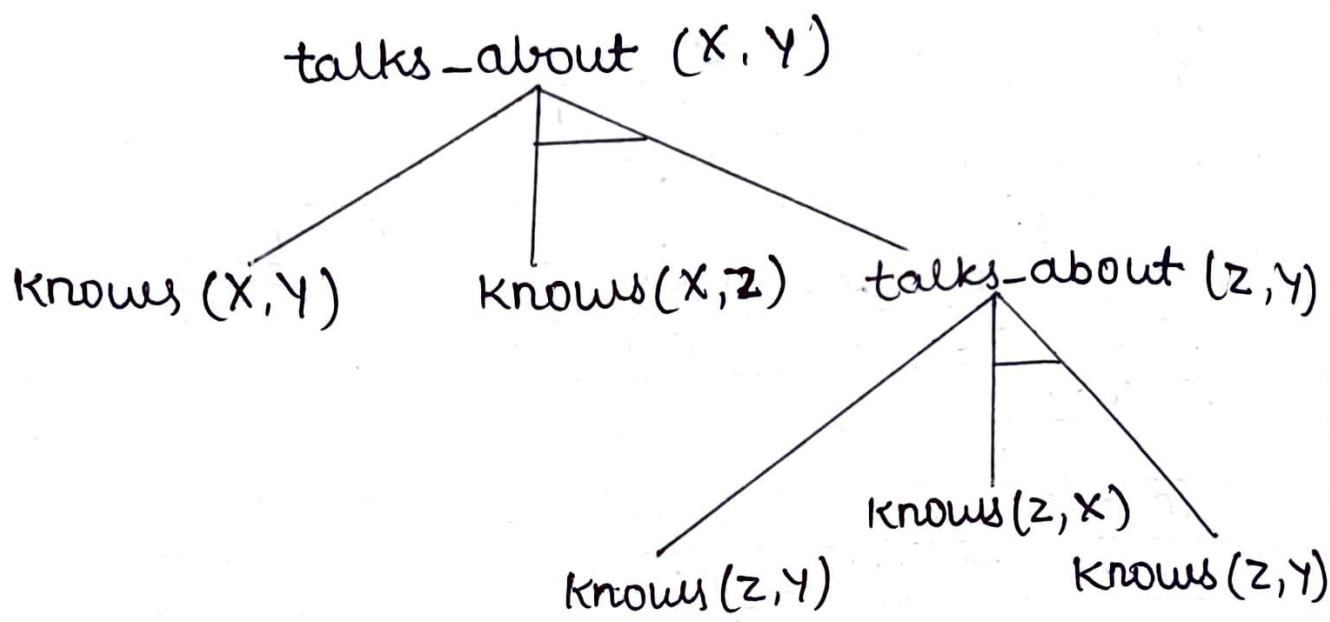
The formulation can be readily
translated into prolog. The mutual
exclusion of the three alternative
categories is indicated by the
cuts.

class (X , fighter):-
beat ($X, -$),
beat ($-, X$).

class (X , winner):-
beat ($X, -$).

class (X , sportsman):-
beat ($-, X$).

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AND/OR Tree for the program

(fig 3)

5. Write output & draw AND/OR tree for the following program:-

`talks-about (A, B) :-`
`know (A, B).`

`talks-about (P, R) :-`
`know (P, Q),`
`talks-about (Q, R).`

Consider the following facts

`know (rathi, ram).`

`know (ram, pat).`

`know (rami, fred).`

`know (fred, rami).`

Explanation:

You talk about someone either you know them or you know someone who talks about them.

If you look AND/OR tree you see that :-

- There is a such tree which is the same shape as the whole tree reflecting the simple recursion.

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2. The solution of given problem depend on being able to stop recursion at same point. Because the left most path down the tree is not infinite in length. It is responsible to help for a solution.

How the prolog respond to the following queries.

a. goal : talks-about (X,Y) :

X = Rani.

Y = Ram.

b. goal : talks-about (Fred,X) :

X = Rani.

c. goal : talks-about (Ram,-) :

True

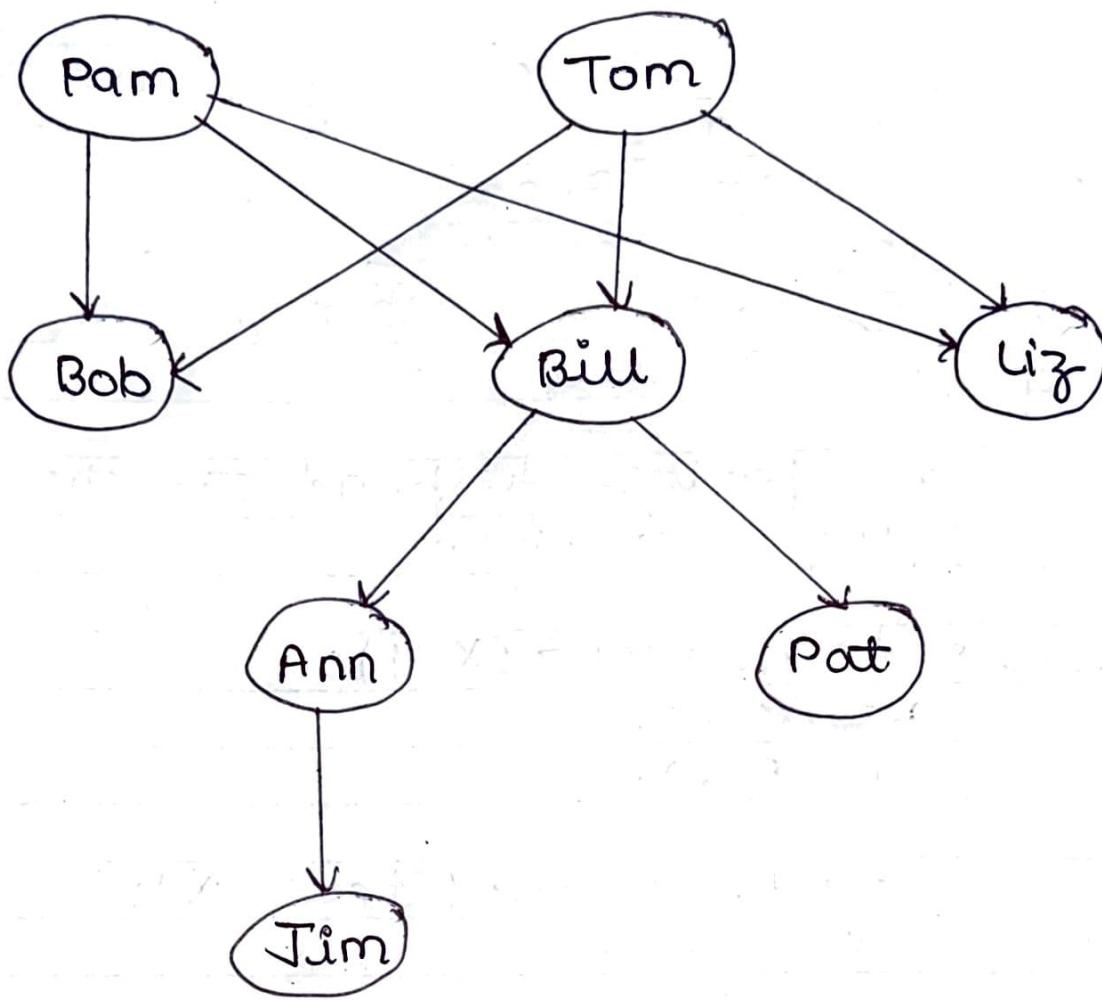
d. goal : talks-about (-,Rani).

True

e. goal : talks-about (-,-)

True

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(fig 4)

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LAB ASSIGNMENT #3



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1. Genealogy problem

- 1.1 Build a knowledge base to represent the parent relationship can be deducted from the tree.

parent (child, parent)

parent (bob, bdm)

parent (bob, tom)

parent (bill, pam)

parent (bill, tom)

parent (liz, pam)

parent (liz, tom)

parent (ann, bill)

parent (pat, bill)

parent (jim, ann)

- 1.2 Build predicates that describes the following family relationship.

a. grandparent

grandparent (grandchild, grandparent) :-

parent (grandchild, parent),

parent (parent, grandparent).

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b. Mother

mother (child, mother) :-

female (mother),

parent (child, mother).

c. Father

father (child, father) :-

male (father),

parent (child, father).

d. Brother

brother (sibling, bro) :-

male (bro),

father (sibling, father),

father (bro, father),

bro = sibling,

mother (sibling, mother),

mother (bro, mother).

e. Sister

sister (sibling, sis) :-

female (sis),

father (sibling, father),

father (sis, father)

sis = sibling,

mother (sibling, mother),

mother (sis, mother).

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f. Aunt

i. Sister of once father or another.

aunt (kid, auntie) :-

female (Auntie),

parent (kid, parent),

sister (parent, auntie).

ii. The wife of once uncle.

aunt (kid, Auntie) :-

female (Auntie),

parent (kid, person),

brother (person, brother),

married (Auntie, brother).

g. Uncle :-

i. The brother of one's mother or father.

uncle (kid, uncle) :-

male (uncle),

parent (kid, person),

brother (parent, uncle).

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ii. The husband of one's aunt

uncle (kid, uncle) :-

male (uncle),

parent (kid, person),

sister (person, sister),

married (uncle, sister).

h. Ancestor :-

ancestor (Person, Ancestor) :-

parent (person, parent),

ancestor (parent, ancestor).

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3. Consider the following sentence:-

a. Translate these sentences into formula in predicate logic & prolog clause form

- John likes all kind of food.

$$\forall x: \text{food}(x) \rightarrow \text{likes}(\text{Jones}, x) \quad /* \text{ Predicate logic */}$$

- Apples are food.

$\text{food}(\text{apples})$.

- Dosa is food.

$\text{food}(\text{dosa})$.

- Anything anyone eats & not killed by its food.

$$\forall x: \forall y: \text{eats}(x, y) \wedge \neg \text{killed}(x, y) \rightarrow \text{food}(y)$$

- Bill eats peanut & still alive.

$\text{eats}(\text{bill}, \text{peanut}) \wedge \neg \text{killed}(\text{bill}, \text{peanut})$.

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f. Liz eats everything Bill eats.

$\forall x, \text{food}(x) \wedge \text{eats}(\text{Bill}, x) \rightarrow \text{eats}(\text{Liz}, x)$.

CNF form

Rule $P \rightarrow Q = \neg P \vee Q$

i. $\neg \text{food}(x_1) \vee \text{likes}(\text{John}, x_1)$

ii. $\text{food}(\text{apple})$

iii. $\text{food}(\text{dosa})$

iv. $\neg \text{eats}(x_2, y) \wedge \neg \text{killed}(x_2, y) \vee \text{food}(y)$
 $= \neg \text{eats}(x_2, y) \vee \neg \text{killed}(x_2, y) \vee \text{food}(y)$.

Rule used $\neg(P \wedge \neg Q) = \neg P \vee Q$

v.a. $\text{eats}(\text{Bill}, \text{peanut})$

b. $\neg \text{killed}(\text{Bill}, \text{peanut})$

vi. $\neg \text{food}(x_3) \vee \neg \text{eats}(\text{Bill}, x_3) \vee \text{eats}(\text{Liz}, x_3)$

Rules:- $(A \cap B) \rightarrow C$
 $\Rightarrow \neg(A \cap B) \vee C$
 $\Rightarrow \neg A \vee \neg B \vee C$

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b. To prove:-

John likes dosa.

$\text{likes}(\text{John}, \text{dosa})$

Let us assume that John doesn't like dosa.

$\neg \text{like}(\text{John}, \text{dosa})$

$\neg \text{like}(\text{John}, \text{dosa})$

(1)

$\neg \text{food}(\text{dosa})$

(1+1)

E

Empty clause

Since, an empty clause occur means assumption is wrong. Therefore

John likes dosa Proved.

Liz like peanuts.

Suppose Liz doesn't like peanuts.

$\neg \text{likes}(\text{liz}, \text{peanuts})$

As. no CNF can be used. So. we can say the information is insufficient.

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4. consider the following sentences
- a. Translate the sentence into formulas in predicate logic & prolog clause form.
- Tom only likes hard courses.
 $\text{hardcourses}(x) \rightarrow \text{likes}(\text{Tom}, x)$
 - Management courses are easy.
Easy can be converted into not hard
 $\neg \text{hardcourse}(\text{management})$
 - All courses in Engineering department are hard.
 $\text{eng-dept}(x) \rightarrow \text{hardcourse}(x)$.
 - Computer engineering is an engg. dept course.
 $\text{Engg-dept}(\text{computer science})$

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CNF FORM:-

Rules used:-

$$P \rightarrow Q = \neg P \vee Q$$

- i. Thardcourse (x) \vee likes (Tom, x)
- ii. Thardcourse (management)
- iii. \neg Engg-dept (x_1) \vee hard courses (x_2)
- iv. Engg-dept (computer science)

b. Prove that "what course would Tom like".

To prove likes (Tom, x) = ?

$$\neg \text{likes} (\text{Tom}, x) \vee \text{likes} (\text{Tom}, x)$$

(i)

$$\neg \text{hardcourse} (x) \vee \text{likes} (\text{Tom}, x_1)$$

(ii)

$$\neg \text{Engg-dept} (x_2) \vee \text{likes} (\text{Tom}, x_2)$$

(iv)

$$x_2 = \text{Computer Science}$$

likes (Tom, Computer Science)

Hence, Tome would like Computer Science

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