# OPERATING SYSTEMS SEMAPHORES SYNCHRONISATION

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## 1 Goal:

• Using semaphores to synchronize the *P* passenger processes and the *C* car processes and analyzing average time of riding for each person based on changing values of *P*, *C*.

# 2 Input:

- Input is the set of numbers:
  - 1. Number of Passenger threads (P)
  - 2. Number of Cars (C)
  - 3. Mean value of exponential wait time between 2 successive ride requests made by the passenger  $(\lambda_P)$
  - 4. Mean value of exponential wait time between 2 successive ride request accepted by a car  $(\lambda_C)$

# 3 Output:

• The output data is written to output file *output.txt* which contains all log for all the rides.

## 4 Procedure:

- Input is read from the file **input.txt** and stored in variables:
  - 1. P stores Passengers value
  - 2. C stores Consumers value
  - 3.  $lamb_1$  stores value of  $\lambda_1$
  - 4.  $lamb_2$  stores value of  $\lambda_2$
  - 5. K stores value of number of ride requests by each passenger

```
int P = 0;
int C = 0;
int lamb1 = 0;
int lamb2 = 0;
int K = 0;
```

- The mutex variables:
  - 1. **ret** This vector array stores all the strings outputted from the program which are needed to be written to the **output.txt** file.
  - 2. cars\_availability The array which keeps track of which car is free, and which car is already in ride.
  - 3. mtx\_writing is used for locking when there is a new string addition to the vector array ret is done.
  - 4. **sem\_array** This is a counting semaphore, which works based on the number of cars free for ride.
  - 5. **total\_time\_passengers** This adds up the total time for the whole process execution for each passenger, then we use it for calculating the average time for each set of  $(P, C, \lambda_1, \lambda_2, K)$

6. **total\_time\_cars** - This adds up the total time for the whole process execution for each car, then we use it for calculating the average time for each set of  $(P, C, \lambda_1, \lambda_2, K)$ 

```
vector<string> ret;
int *cars_availability;
pthread_mutex_t mtx_writing;
sem_t *sem_array;
double total_time_passengers = 0;
double total_time_cars = 0;
```

- This structure is used to send arugument to the passenger() whose parameters are:
  - 1. passenger number This variable stores the serial number of passenger
  - 2. **cars** This is the array which has the thread ID's of the cars threads.
  - 3. attr The default paramter for thread creation

```
struct data_to_passenger_func
{
   int passenger_number;
   pthread_t *cars;
   pthread_attr_t attr;
};
typedef struct data_to_passenger_func parampass;
```

- This structure is used to send arugument to the *cars*() whose parameters are :
  - 1. **car\_number** This variable stores the serial number of car in use by the current passenger.
  - 2. passenger\_number This variable stores the serial number of passenger
  - 3. **time\_sleep** This is the time for which the car thread sleeps so that the given constraint of not accepting new request is satisfied

```
struct data_to_car_func
{
   int car_number;
   int passenger_number;
   double time_sleep;
};
typedef struct data_to_car_func paramcar;
```

• The following code block corresponds to all the intialization of mutex, semaphore variable and the availability array for the cars.

```
// intializing the mutex variable
  pthread_mutex_init(&mtx_writing, nullptr);

// initializing semaphore variable to cars count
  sem_array = sem_open("/arraysem", O_CREAT | O_EXCL, 0644, C);
```

```
// initializing the cars availability array
  cars_availability = (int *)calloc(C, sizeof(int));
  for (int i = 0; i < C; i++)
    cars_availability[i] = 1;</pre>
```

- Creating passenger threads, with the pointer variable of **parampass struct** being passed as argument with suitable values loaded in.
- Control of program is now shifted from *main()* to *passengers()*.

```
pthread_t passengers[P]; // This is the thread pool of passengers
pthread_t cars[C]; // This is the thread pool of cars
pthread_attr_t attr;
pthread_attr_init(&attr);

for (int i = 0; i < P; i++)
{
    auto *dat = (parampass *)calloc(1, sizeof(parampass));
    dat->passenger_number = i + 1;
    dat->cars = cars;
    dat->attr = attr;
    pthread_create(&passengers[i], &attr, passenger, dat);
}
```

• The following code block is used for generation of the random times used in the program in the critical section of *passenger()* and *cars()* threads.

```
int seed = chrono::system_clock::now().time_since_epoch().count();
  default_random_engine generator(seed);
  exponential_distribution < double > t_1(1.0 / lamb1);
  exponential_distribution < double > t_2(1.0 / lamb2);

double t1 = t_1(generator);
  double t2 = t_2(generator);
```

• We use the mutex lock **mtx\_writing**, whenever we want to add a string value to the vector array **ret**.

```
pthread_mutex_lock(&mtx_writing);
    /*
    Adding a new string to ret vector array
    */
pthread_mutex_unlock(&mtx_writing);
```

• Whenever we want to find the local time the above code block is used. From the variable **timeinfo** we can get the required info about local time by accessing the struct data.

```
time_t raw_time;
struct tm *timeinfo;
  time(&raw_time);
timeinfo = localtime(&raw_time);
```

- The following code block is the entry section for the passenger thread.
- The semaphore variable **sem\_array** ensures that a passenger thread can only access the **car\_availability** array when there is a car available for riding i.e., **sem\_array**> 0.
- After crossing the *sem\_wait()* function search for a free car is done, after finding it the index of car it is stored in variable **index\_car** variable which is then passed on to thread function *cars()*.
- After loading required values to argument for passing we create a car thread with the **pthread\_t** variable **dat->cars[index\_car]**.

```
// Entry section of the passenger thread
 for (int j = 0; j < K; j++)
  {
    int index_car = 0;
// Blocked this section inorder to remove race conditions
// for array value choosing
    sem_wait(sem_array);
    for (int i = 0; i < C; i++)
      if (cars_availability[i] == 1)
      {
        index_car = i;
        cars_availability[i] = 0;
        break:
      }
    }
 paramcar *a = (paramcar *)calloc(1, sizeof(paramcar));
  a->passenger_number = dat->passenger_number;
 a->car_number = index_car + 1;
  a->time_sleep = t2;
 pthread_create(&dat->cars[index_car], &dat->attr, cars, a);
  pthread_join(dat->cars[index_car], nullptr]);
  cars_availability[index_car] = 1;
  sem_post(sem_array); // Incrementing the sem_array value
```

- Now here the program control shifts from *passengers*() to *cars*().
- The following code block is the Entry Section and CS of the car thread.
- The sleep time is in  $\mu$  sec.

```
// Entry section for car thread
  pthread_mutex_unlock(&mtx_writing);

// CS for car thread
  usleep(d->time_sleep*1000);
```

- We use the following code block to calculate time of execution of an action.
- The output time is in  $\mu$  seconds.

```
auto start = high_resolution_clock::now();
/*
  The required action is performed
*/
  auto stop = high_resolution_clock::now();
  auto timees = duration_cast<microseconds>(stop - start);
  cout << (double)timees.count();</pre>
```

- After execution of Remainder Section control of the program shifts from *cars*() to *passengers*().
- The function *pthread\_join()* ensures that the *car* thread is executed fully.
- Then the **cars\_availibity** arrays value is resette to 1. Then semaphore value is incremented using *sem\_post()*.

```
pthread_join(dat->cars[index_car], nullptr);
  cars_availability[index_car] = 1;
  sem_post(sem_array);
```

• The critical section of passenger depends on  $t_1$ ,  $t_2$  if  $t_1 > t_2$  then we need to still wait for a time of  $t_1 - t_2$ , for the passenger to pass new request to a car. But when  $t_1 < t_2$ , in this case as already the passenger has passed the time  $t_1$  passenger can put forth a new request as soon as the passeger exits the car.

```
// CS of passenger thread
if (t1 > t2)
usleep((int)(t1-t2)*1000);
```

The average Time of execution of Passengers and Cars is calculated like as shown below.

- File writing is done with the global vector array **ret**.
- The output to file is not in the order that was mentioned in the problem statement, this was done in purpose inorder to prove the concurrency of the program.
- In the output file the system time of the output can be seen it will not change this is because as we considered the sleep times in milli seconds.

```
ofstream op;
op.open("output.txt", ios_base::out);
  for (const auto & i : ret)
{
   op << i;
}</pre>
```

# 5 Analyzing Program Output:

- We analyze the program in two modes:
  - 1.  $T_{avg}$  for passengers to complete their tour with Total number of Cars as constant (C is constant).
  - 2.  $T_{avg}$  for Cars to complete their tour with Total number of Passengers as constant (P is constant).

## 5.1 C is constant:

- The value of *C* used by the program is fixed and is equal to 25.
- The value of *K* used by the program is equal to 5
- $\lambda_1$  is equal to 10
- $\lambda_2$  is equal to 10

## 5.1.1 Graph:

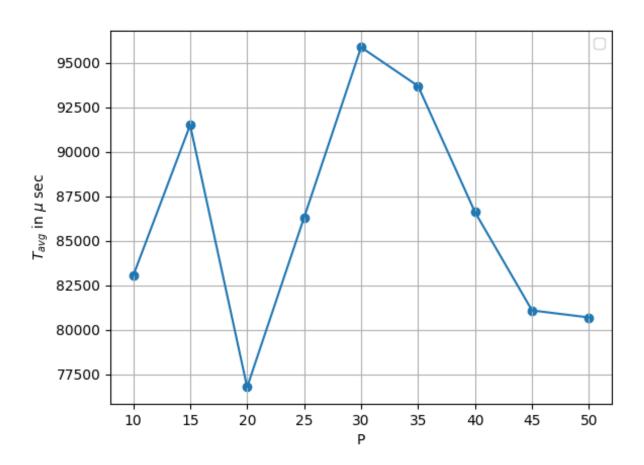


Figure 1:  $T_{avg}$  vs P(C is constant)

#### **5.1.2** Observations:

#### • Expected:

- As the value of P increases  $T_{avg}$  should increase as C is constant waiting time of each passenger increases.

#### • Observed:

- We observe some unusual trend in the  $T_{avg}$  of this case.
- We see that there isn't much difference between the average times between the values of P = 10 to P = 25 which is because of the fact that the value of C = 25 we have more cars than passengers, because of this the time differences are not much high in this range of P, i.e., the waiting time isn't that high.
- The value of  $T_{avg}$  increases as expected in P = 20 to P = 30.
- Except for the values at P=20 and P=30 all other  $T_{avg}$  are nearly at the same times, as we are measuring the time in  $\mu$  seconds.
- This irregular time variations can be attributed to hardware implementation of the PC.

#### 5.2 *P* is constant:

- The value of *P* used by the program is fixed and is equal to 50.
- The value of K used by the program is equal to 3
- $\lambda_1$  is equal to 10
- $\lambda_2$  is equal to 10

# 5.2.1 Graph:

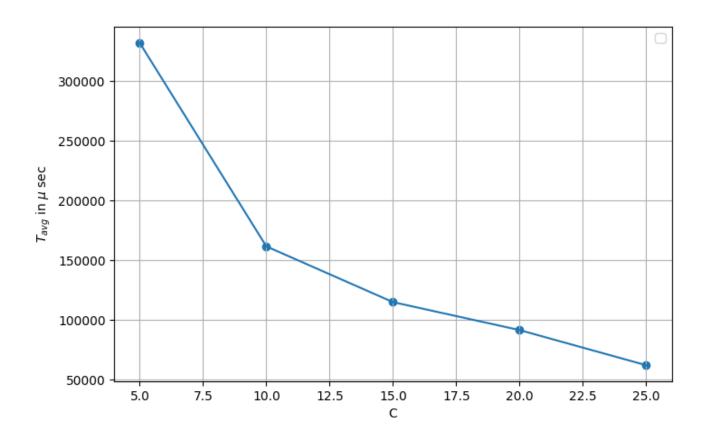


Figure 2:  $T_{avg}$  vs C (P is constant)

#### 5.2.2 Observations:

## • Expected:

- As the value of *C* increases the average working time decreases.

#### • Observed:

- As the value of C increases  $T_{avg}$  decreases, which follows the expected pattern.
- This is because as the number of cars increases work load on each car decreases which inturn decreases the time of ride for each car which results in decrease of  $T_{avg}$ .