```
Pnt oci
       myPnt y;
       typodef mygnt smallant;
       smallint Z;
       Prantf ( "enter two Volues");
       Boanf (" %d &d", &x, &y);
       I=octy;
       Printf(" our value 10:0/d", x);
      getch();
      getuin o;
    output:
    enter two Values: 4, 2
    our Value &: 6
Otructure padding:
* Structure padding is a Concept in a that addo the
    one more empty bytes between the memory addresses to
    algon the data 9n momory
* Suppose we create a user defined structure. When we
   create an object of this otructure, then the contigious
   memory will be allocated to the structure member.
    Struct Student
     Chara;
    Chan b;
     Int C;
   4 stul;
```

- \* In 1HB. Student - structure name Otul - object of the structure
- \* After the creation of an object, a contiguous block of memory to allocated to 940 structuse members.
- \* Memory will be allocated to a then b' after that c' Now we calculate the give of the struct student struct student

chan a; 11 byte chas b; ///byte 9nt c; 1/4 byte.

In the above case, we calculate the size of the 'Student', 89ze come to 6 bytes. But this answer is wrong. Now, we will understand the concept of structure padding.

structure padding:

- \* Generally the processor does hot not read 1 syte at a 19me. It roads I word at a time.
- \* If we have 32 but processor, then the processor reads A byte at a time, which means that I word equal to Abytes. ge I word = 4 byte
  - I word = 8 byte [9n 64 691 processor] so spice of the word deponds anchatecture.

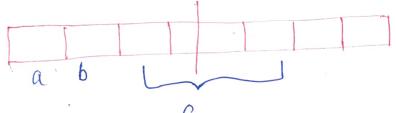
Why otructure padding:

Struct Student

Chool a; //1 byte chool b; // 1 byte

Chan c; 1 A byte

If we have a 82 best processor (A byte at a Ame) then the processor the memory for the cubive structure would be.



\* As we know that structure occupies the contiguous block of memory as shown in the above dragram. The I byte for chan'a, I byte for chan'b, and 4 byte for chan'c'

\* The 4 bytes can be accessed at a 19me as we are considering the 50 bit archatecture. The problem is that 90 90 one could cycle,

1 byte of chos a, 1 byte of chos b, and 2 byte of 9nt c can be accessed.

\* We will not face any problem during accepting of a and b Voorlables as both the voorlables can be accepted in one councycle, but into voorlable as a crossed in one countried to access the value of the c'voorlable.

- \* Suppose we do not Wart to access a so b' voorables, we only want to access the Voorable c' with regulares two cycles.

  This is an unnecessary wastage of CPU cycles.
- \* Due to the roason, the structure radding concert was introduced to save the number of CPV cycle.
- \* The structure padding 95 done by the compiler automatically



\* In order to action the otructure padding, an empty row is created on the left [ref picture] and it wo by the withch are occupied by the & voorable on the left cope possible it and the right so all the four bytes of & voorable one on the right.

\* Now the c variable can be accessed in a single con cycle.

After structure padding, the total momory occupied by the
atructure is 8 bytes

a-1 byte b-1 byte empty-2 bytes

C- A bytes

Which is greater than previous one. Although the memory by wasted in this case, the Vostable can be accessed in within a single CPU Eycle