

**Assignment 01**  
**EECS 3101**

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## **Problem 1**

(a) Find the proper loop invariant for the while-loop

**Inv(p1, p2, p3):**

At the start of each iteration of the **while** loop of lines 7–18, the following conjuncts are true:

1. If  $p3 > 0$ , then the sub-array of `lst` from 0 to  $p3-1$  (including both indices) consists of the elements originally in `lst` from 0 to  $p3-1$ , but in sorted order.
2.  $p1 \leq p2 \leq p3$
3. if  $p1 > 0$ , then the sub-array of `lst` from 0 to  $p1-1$  contains only 'A'
4. if  $p2 > 0$ , then the sub-array of `lst` from  $p1$  to  $p2-1$  contains only 'B'
5. if  $p3 > 0$ , then the sub-array of `lst` from  $p2$  to  $p3-1$  contains only 'C'
6.  $p3 \leq \text{lst.length}$

(b) Prove the invariant that you find in (a), using induction.

Let's consider the base case, which is before the start of the while loop.  $p1=p2=p3=0$ . It is easy to see that all the above 6 conditions are true. Hence the invariant  $\text{Inv}(p1, p2, p3)$  is true for the base case.

Now let's assume that the invariant is true at an arbitrary iteration of the while loop. Then we have to show that the invariant will be true for the next iteration as well.

So at iteration  $n$ , we know the sub-array of `lst` from 0 to  $p3-1$  consists of the elements originally in `lst` from 0 to  $p3-1$ , but in sorted order. Then for iteration  $n+1$ , we have to consider several cases;

Case 1:  $\text{lst}[p3] == \text{'A'}$ : if the character at position  $p3$  is 'A', then we swap  $\text{lst}[p3]$  with the  $\text{lst}[p1]$ . There are three sub cases.  $\text{lst}[p1]$  can be either 'A', 'B' or 'C'

Case 1.1:  $\text{lst}[p1] == \text{'A'}$ : This can happen only if  $p1=p2=p3$  (by condition 3, 'A' appears only up to  $p1-1$ , so if 'A' appears again at  $p1$ , that would mean there are no 'B's and 'C's in the `lst` from 0 to  $p3$ ). So swapping changes nothing,  $p1$  is incremented by 1,  $p2$  is also incremented by 1 and  $p3$  is also incremented by 1. So  $p1 \leq p2 \leq p3$  also holds true. All conditions hold true & the  $\text{inv}(p1, p2, p3)$  holds true.

Case 1.2:  $\text{lst}[p1] == \text{'B'}$ : After swapping  $\text{lst}[p1]$  with  $\text{lst}[p3]$ ,  $\text{lst}[p3]$  becomes 'B' and  $\text{lst}[p1]$  becomes 'A'.  $p1$  is incremented. The condition that the sub-array of `lst` from 0 to  $p1-1$  contains only 'A' holds true again. Earlier we had  $\text{lst}[p1:p2-1]$  to be 'B', but after  $p1$  is incremented,  $p2$  is set as  $\max(p1, p2)$  which means that  $p2$  is at least  $p1$ . So again,  $\text{lst}[p1:p2-1]$  are all 'B' holds true again.

But we have a 'B' at  $\text{lst}[p3]$  and there could be 'C' before that (from  $\text{lst}[p2:p3-1]$ ). This breaks the sort order (condition 1), but it is taken care of by the second if condition inside the while loop.  $\text{lst}[p3]$  &  $\text{lst}[p2]$  are swapped. That is,  $\text{lst}[p2]$  is set as 'B' and  $p2$

is incremented. Hence the condition that sub-array of `lst` from `p1` to `p2-1` contains only 'B' holds true once again. If `p2` was equal to `p3`, then the swapping is b/w the same position, `p2` incremented so condition 4 holds true. Otherwise (if `p3` was greater than `p2`), then 'C' was at `p2` (by condition 5), and it goes to `p3`. `p3` is incremented by 1. The condition 5 holds true again.  $p1 \leq p2 \leq p3$  also holds true. All the 6 conditions will hold true and hence the invariant is true.

Case 1.3: `lst[p1] == 'C'`: This would mean that there are no 'B' from 0 to `p3-1` & `p1` is equal to `p2`. (by conditions 4 & 5). Then swapping `p1` & `p3`, means 'A' comes to `p1` (& `p1` is incremented means condition 3 holds true), and 'C' comes to `p3`. Since there are no 'B's, condition 4 is trivially true, and condition 5 also holds true because `p3` will be incremented by 1. The `inv(p1, p2, p3)` holds true.

Case 2: `lst[p3] == 'B'`: if the character at position `p3` is 'B', then it is swapped with character at position `p2`. `p2` can be equal to `p3`, in which case `p2` is simply incremented (so condition 5 holds true). `p3` is also incremented so condition 5 is also true trivially.

On the other hand, if  $p2 < p3$ , then the character at `lst[p2]` is 'C' (by condition 5). Hence C goes to `p3` and B goes to `p2`. `p2` is incremented, so condition 4 holds true again. Also `p3` is incremented, so condition 5 also holds true. Condition 3 will also hold true, because we haven't changed any element from 0 to `p1-1`.

Case 3: `lst[p3] == 'C'`: Both the if conditions are false. The `p3` is incremented by 1. So the condition 5 that 'C' appears from `p2` to `p3-1` holds true again. The other conditions also hold true, since we haven't touched `p1` or `p2`. The `inv(p1, p2, p3)` holds true.

Hence by the principle of induction, the `inv(p1, p2, p3)` is true for any arbitrary iteration `n`.

(c) If the loop terminates, then `p3 = lst.length`. Hence by condition 1, the `lst` contains all original elements in `lst` but in sorted order. Hence `SortABC` is proven to work correctly (partially correct) if it terminates.

(d) It is easy to see that the loop terminate after finite number of iterations. In fact the loop terminates after exactly `n` number of iterations (where  $n = \text{lst.length}$ ). This is because the variable `p3` is incremented by 1 in each iteration and as soon as `p3` becomes equal to `n`, the loop terminates.

## **Problem 2**

(a) Find the proper loop invariant  $\text{Inv}(x, y)$

The loop invariant contains the following conjuncts;

$\text{Inv}(x, y)$ :

Before the start of the while loop from lines 6-13, the following conditions should hold true;

1.  $x \leq \text{LCM}(x, y)$
2.  $y \leq \text{LCM}(x, y)$
3.  $x > 0$
4.  $y > 0$

(b) Prove the invariant by induction.

Base Case: Initially  $x=a, y=b$ , if the condition inside the while loop becomes false, i.e.  $x=y$ , then  $a=b \Rightarrow \text{LCM}(a,b) = a$ . Hence (return  $x$ ) will return  $a$ , which is the correct solution. So all the above conditions of the invariant are true. (by preconditions,  $a \geq 1$  &  $b \geq 1$ , so  $x > 0$  &  $y > 0$  hold true. Also  $x \leq \text{LCM}(x,y)$  and  $y \leq \text{LCM}(x,y)$  is also true. So the invariant is true). Otherwise, if  $x \neq y$ , then also the invariant is true trivially.

Induction: Assume the  $\text{inv}(x, y)$  is true at step  $n$ , then we have to show that the invariant is true for step  $n+1$  as well.

Case 1:  $y = \text{LCM}(x, y)$ : In this case  $x < y$  (otherwise the if  $x = \text{LCM}(x,y)$  then the loop would have exited in the previous iteration itself, or if  $x > \text{LCM}(x,y)$ , then the condition 1 will be violated, which we assumed to be true). Then we increment  $x$  by  $a$ .  $x$  cannot exceed  $\text{LCM}(x,y)$  because we increment  $x$  by  $a$  every step, and by the definition of LCM,  $\text{LCM}(x,y)$  is an integer multiple of  $a$ . So  $x$  has to be  $\leq \text{LCM}(x,y)$  after update. Hence the condition holds true.

Case 2:  $x = \text{LCM}(x, y)$ : (same argument with  $x$  &  $y$  interchanged)  $y < x$ .  $y$  is an integer multiple of  $b$  and  $\text{LCM}(x, y)$  is also an integer multiple of  $b$ . So  $y \leq \text{LCM}(x, y)$

Case 3: Both  $x$  &  $y < \text{LCM}(x, y)$ : None of  $x$  or  $y$  can exceed  $\text{LCM}(x, y)$  since  $\text{LCM}(x, y)$  is an integer multiple of  $a$  &  $b$ .  $x$  is an integer multiple of  $a$  and  $y$  is an integer multiple of  $b$ .

Hence the invariant  $\text{inv}(x, y)$  holds true at iteration  $n+1$  as well.

(c) If the loop terminates, the  $x = y$ . But we know by the definition of  $\text{LCM}(x, y)$ , it is the smallest number where an integer multiple of  $a$  becomes equal to an integer multiple of  $b$ . Coupled with the invariants that  $x \leq \text{LCM}(x,y)$  &  $y \leq \text{LCM}(x,y)$ ,  $x = y = \text{LCM}(x, y)$ .

Hence the proof.