

## Page tables

xv6 runs on SV39 Risc.

PT  $\rightarrow$  4096 byte  $2^{12}$

RISV page table is logically an array  
of  $2^{24}$  PTE

## Context Switch

↓  
happens during I/O or Interrupt

If timer interrupt, give up CPU.  
if (dev == 2) - yield  
it is timer interrupt.

yield  $\rightarrow$  proc.c  
↓

takes current process state  
saves it & resumes it later.

save all registers callee registers

ra  $\rightarrow$  return address

sp  $\rightarrow$  stack point

process is running, CPU wants to switch  
to another process then...  
 $\rightarrow$  there will be a timer interrupt  
 $\rightarrow$  call yield() in trap.c  
 $\rightarrow$  turns a RUNNING process to

yield  
double  
calls

context.

yield()

$\rightarrow$  calls

sw

it does  
process is called  
one it just  
and load

RUNNABLE  
yield is in proc.

↓  
yield calls sched

↓  
double check all conditions

↓  
calls switch to switch the

context.

switch()

↳ calls switch (sp → context,

↓  
it is an assembly code

ld, sd

it doesn't know anything about which  
process is called etc

it just takes a new context & an old  
one & save old context registers  
and load from new.

sd, ra, 0(a0)

sd, sp, 8(a0)

ld ra, 0(a1)

ld sp, 8(a1)

now  
CPU  
will start  
at address  
indicated at  
new ra



• scheduler → proc.c  
each CPU calls scheduler after setting up, it is never ending.

→ keep interrupt ON

```
for (p = proc; p < bproc [NPROC]; p++)  
    acquire (&p → lock)
```

### Uthreads.c

in thread-init (which is called first)  
we initialise all-thread[0] with  
~~RUNNABLE~~ <sup>RUNNING</sup> state & current-thread =  
&all-threads[0],

then  
then control goes to create-thread.  
then schedule. "we make threads ~~RUNNABLE~~

in schedule

1st if  $t = \text{current-thread} + 1$  & all-thread[1]

for (i < MAX-THREADS)

if (t >= all-threads + MAX-THREAD)  
t = all-thread

if (t → state = RUNNABLE)

next-thread = t; // here we  
get our  
next runnable  
thread  
break

t = t + 1;

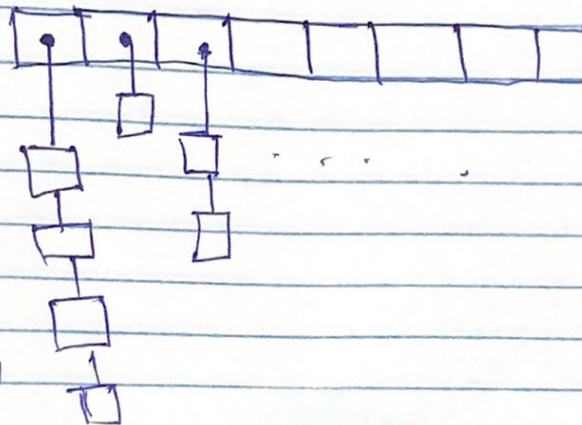


② struct entry \*table[NBUCKET]

array of pointers that point to structures [pointers to array].

```
struct entry {
    int key;
    int value;
    struct entry *next;
};
```

now entry has a pointer next, so, \*table[NBUCKET], each array element can point to either one struct or even a list head.



20, 50  
20, 50

1, 20

table[0] → [ ] [ ] [ ] [ ]  
again points to list

as we  
get our  
at Runnable  
thread