Page tables

xv6 Runs on SV39 RISC.

PT → 4096 byte 2^{12}.

RISV page table u logically an array of 2^{27}. PTE.

Context Switch

↓

happens during I/O or Interrupt

If Timer interrupt, give up CPU.

if (dev == 2) – yield

it is timer interrupt.

yield → proc.c

↓

takes current process state

saves it & resumes it later.

save all registers callee register

ra → return address

sp → stack point

process is running, CPU wants to switch to another process2 – then....

→ there will be a timer interrupt

→ call yield() in trap.c

↳ turns a RUNNING Process to

RUNNABLE

yield is in proc.

↓

yield calls sched

↓

double check all conditions

↓

calls switch to switch the context.

sched()

↳ calls swtch (&p→ context,

↓

it is an assembly code

ld , sd

it dosn't know anything about which process is called etc

it just takes a new context & an old one & save old context registers and load from new.

```
sd , ra , 0(a0)
sd , sp , 8(a0)
.
.
.
ld ra 0(a1)
ld , sp , 8(a1)
```

now CPU will start at address at indicated at new ra

↓ scheduler → proc.c

each CPU calls scheduler after setting itself up, it is never ending.

→ keep interrupt ON

```
for (p = proc ; p < &proc[NPROC] ; p++)
    acquire(&p->lock)
```

uthreads.c

in thread_init (which is called first) we initialise all_thread[0] with RUNNING ~~RUNNABLE~~ state & current_thread = &all_thread[0],

then control goes to create_thread. then schedule. "we make threads RUNNABLE

In schedule

1st fn t = current_thread + 1 &all_thread[1]

```
for (i < MAX_THREADS)
    if (t >= all_threads + MAX_THREAD)
        t = all_thread
    if (t->state = RUNNABLE)
        next_thread = t;   // here we get our next Runnable thread
        break
    t = t+1;
```

(2) struct entry *table [NBUCKET]

array of pointers that point to structures [pointer to array].

```
struct entry {
    int key;
    int value;
    struct entry *next ;
};
```

now entry has a pointa next, so. *table [NBUCKET], each element can point to either one struct or even a list.head.