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<b>Experiment</b>	2
<b>Aim</b>	Understand sorting algorithms based on Divide and Conquer approach
<b>Objective</b>	1) Learn Divide and Conquer strategy 2) Find any one local minima in a given array
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<b>Class</b>	CSE AIML
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<b>Date of Submission</b>	05/02/2024

<b>Algorithm and Explanation of the technique used</b>	<ol style="list-style-type: none"><li>1. Accept the size of the array from the user.</li><li>2. Accept elements for the array.</li><li>3. The localMinima function takes the array, its size, start index, and end index as parameters.</li><li>4. Compute the middle index <math>m</math> as <math>(start + end) / 2</math>.</li><li>5. Check if the middle element is a local minima:     If <math>m</math> is at the beginning (<math>m == 0</math>) and the element at <math>m</math> is less than the one following it (<math>arr[m + 1]</math>), return <math>arr[m]</math>.     If <math>m</math> is at the end (<math>m == size - 1</math>) and the element at <math>m</math> is less than the one preceding it (<math>arr[m - 1]</math>), return <math>arr[m]</math>.     If <math>m</math> is neither at the beginning nor at the end, and the element at <math>m</math> is less than both its neighbors, return <math>arr[m]</math>.</li><li>6. If the middle element is not a local minima:     If the element to the left of <math>m</math> is smaller, recursively search the left half (start to <math>m - 1</math>).     If the element to the right of <math>m</math> is smaller, recursively search the right half (<math>m + 1</math> to end).</li><li>7. If no local minima is found, return -1.</li><li>8. End</li></ol> <p><b>Divide-and-Conquer Approach:</b> The function localMinima takes an array, its size, and the start and end indices. It calculates the middle index and checks if the middle element is a local minima.</p>
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	<p>If the middle element is not a local minima, it recursively searches in the left or right half of the array based on which side has a smaller neighboring element.</p> <p>The recursion terminates when a local minima is found or when the search space reduces to a single element.</p>
<b>Program(Code)</b>	<pre> import java.util.Scanner; public class Main {     static int localMinima(int[] arr, int size, int start, int         end) {         int m=(start+end)/2;         if(m==0 &amp;&amp; arr[m]&lt;arr[m+1]){             return arr[m];         }         else if(m==size-1 &amp;&amp; arr[m-1]&lt;arr[m]){             return arr[m];         }         else if((m!=0 &amp;&amp; m!=size-1) &amp;&amp; (arr[m]&lt;arr[m-1] &amp;&amp; arr[m]&lt;arr[m+1] ) ){             return arr[m];         }         else if (m!=0 &amp;&amp; arr[m-1]&lt;arr[m]) {             return localMinima(arr, size,0,m-1);         }         else if (m!=size-1){             return localMinima(arr, size,m+1,size-1);         }         return -1;     }     public static void main(String[] args) {         Scanner sc = new Scanner(System.in);         System.out.print("Enter the size of array: ");         int size = sc.nextInt();         int arr[]=new int[size];         System.out.println("Enter elements for the array:");         for (int i = 0; i &lt; size; i++) {             arr[i] = sc.nextInt();         }         System.out.println("Local minima is :             "+localMinima(arr, size,0,size-1));     } } </pre>

<b>Output</b>	<pre> Enter the size of array: 7 Enter elements for the array: 2 1 5 3 4 7 9 Local minima is : 3  Enter the size of array: 8 Enter elements for the array: 2 1 4 7 5 7 9 8 Local minima is : 1  Enter the size of array: 7 Enter elements for the array: 1 2 3 4 5 6 7 Local minima is : 1 </pre>
<b>Justification of the complexity calculated</b>	<p><b>TIME COMPLEXITY:</b></p> <p>At each step of the recursion, the array size reduces by half: <math>n/(2^k)</math>, where <math>k</math> is the depth of recursion.</p> <p>After <math>k</math> steps of halving, the size of the array becomes 1: <math>n/(2^k)=1</math></p> <p>Solving for <math>k</math>, <math>n=2^k</math>, that is <math>k=\log(n)</math> to base 2. Therefore, the time complexity of the algorithm is <math>O(\log n)</math>.</p> <p>Best Case Time Complexity: <math>O(1)</math> :- Constant time when the local minima is found in the first comparison.</p> <p>Worst Case Time Complexity: <math>O(\log n)</math> :- Logarithmic time when the local minima is located at one of the ends, requiring a complete binary search.</p> <p>Average Case Time Complexity: <math>O(\log n)</math> :- Logarithmic time when the local minima is found within the array not necessarily at the ends.</p>
<b>Conclusion</b>	<p>I learned about the divide and conquer algorithm. The code implements an algorithm to find a local minima within an array using a binary search approach. The algorithm divides the array into halves at each step of the recursion, efficiently reducing the search space until it identifies a local minima. Overall, the code demonstrates an effective implementation of the local minima search algorithm.</p>