

DCS Assignment 1

(1)

It is given that most of the user queries involve the Author-Name and/or Book Name.

So, we can use first 3 letters of the Book Name to design hash function.

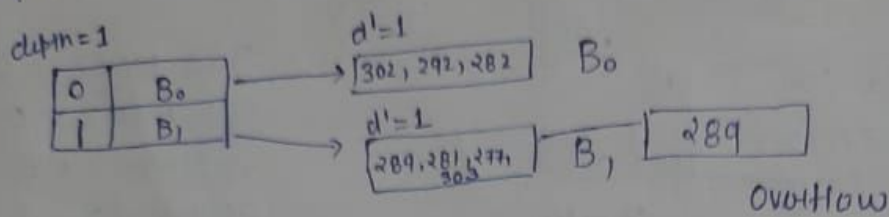
Hash function used Sum of ASCII Value of first three characters.

Book first 3 characters	ASCII Value (sum)	Binary
How	302 - - - -	100101110
Sei	292 - - - -	1100010000
Mal	282 - - - -	100010000
Emo	289 - - - -	100100001
Pha	281 - - - -	100011001
Fan	277 - - - -	100010101
Gui	303 - - - -	100101111
The	289 - - - -	100100001
Wod	298 - - - -	100101010

for finding Sum of ASCII Value we consider first letter Capital ~~and~~ uppercase next two in lower case.

(a) Bucket size : 4

Step 1: mod. function used : $\text{key mod } 2$



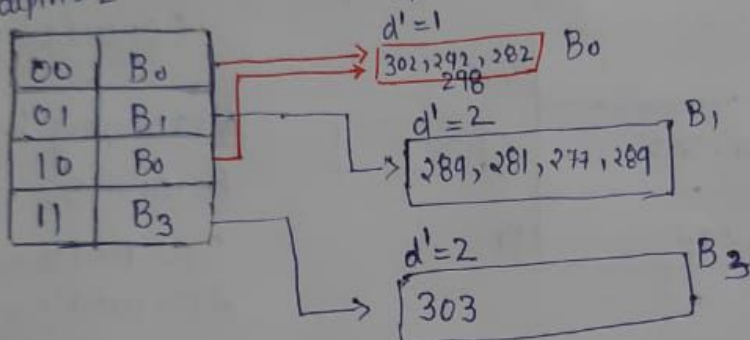
Value inserted

$302 \text{ mod } 2 = 0 \rightarrow B_0$
 $292 \text{ mod } 2 = 0 \rightarrow B_0$
 $282 \text{ mod } 2 = 0 \rightarrow B_0$
 $289 \text{ mod } 2 = 1 \rightarrow B_1$
 $281 \text{ mod } 2 = 1 \rightarrow B_1$
 $277 \text{ mod } 2 = 1 \rightarrow B_1$
 $303 \text{ mod } 2 = 1 \rightarrow B_1$
 $289 \text{ mod } 2 = 1 \rightarrow B_1$
 (Overflow)
 Now have to split

Step 2

depth=2

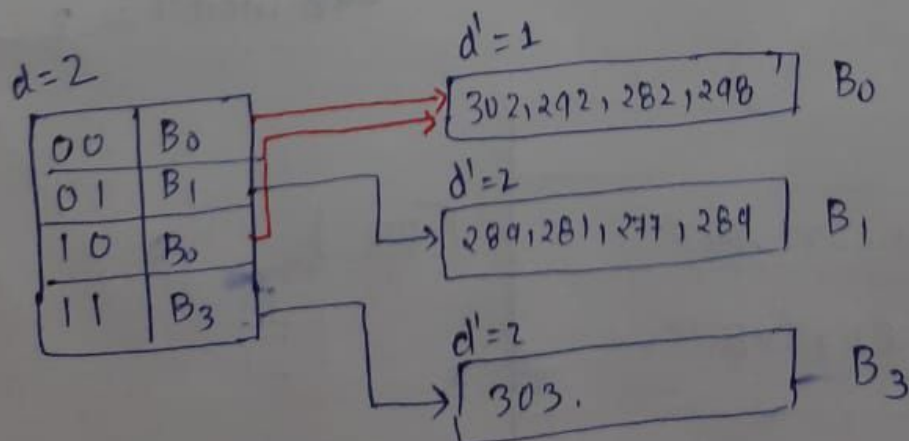
hash function : $(\text{key}) \text{ mod } 4$



Value inserted after splitting

$289 \text{ mod } 4 = 1 = B_1$
 $281 \text{ mod } 4 = 1 = B_1$
 $277 \text{ mod } 4 = 1 = B_1$
 $298 \text{ mod } 4 = 2 = B_0$
 $303 \text{ mod } 4 = 3 = B_3$
 $289 \text{ mod } 4 = 1 = B_1$
 Now New value
 $298 \text{ mod } 4 = 2 = B_0$
 since B_2 is pointing to B_0

Final structure.



(b) To be inserted 302, 292, 282, 289, 281, 277, 303, 289, 298 (3)

① Number of bucket = 2

Bucket size = 4

$$H_0 = \text{key} \bmod 2$$

It denotes
which bucket
to be split

B_0	302, 292, 282,	H_0
B_1	289, 281, 277, 303	H_0

Overflow

So split B_0 to B_0 and B_3

New Hash function $H_1 = \text{key} \bmod 4$

B_0	292	H_1
B_1	289, 281, 277, 303	H_0
B_2	302, 282, 298	H_1

$$302 \bmod 2 = 0 \quad B_0$$

$$292 \bmod 2 = 0 \quad B_0$$

$$282 \bmod 2 = 0 \quad B_0$$

$$289 \bmod 2 = 1 \quad B_1$$

$$281 \bmod 2 = 1 \quad B_1$$

$$277 \bmod 2 = 1 \quad B_1$$

$$303 \bmod 2 = 1 \quad B_1$$

$$289 \bmod 2 = 1 \quad B_1$$

(Overflow)

Arranging keys of
 B_0 and B_2 using H_1

$$302 \bmod 4 = 2 \quad B_2$$

$$292 \bmod 4 = 0 \quad B_0$$

$$282 \bmod 4 = 2 \quad B_2$$

Now insert 298

$$298 \bmod 2 = 0$$

\therefore It is directed B_0
which is using H_1

So again to check between
 B_0 and B_2

$$298 \bmod 4 = 2 = B_2$$

final

B_0	292	
B_1	289, 281, 277, 303	289
B_2	302, 282, 298	

(C) record < MMI-009; Mi009-Soc, Marvin Minsky,
Society of Mind >

Book first 3 char ASCII value

Binary.

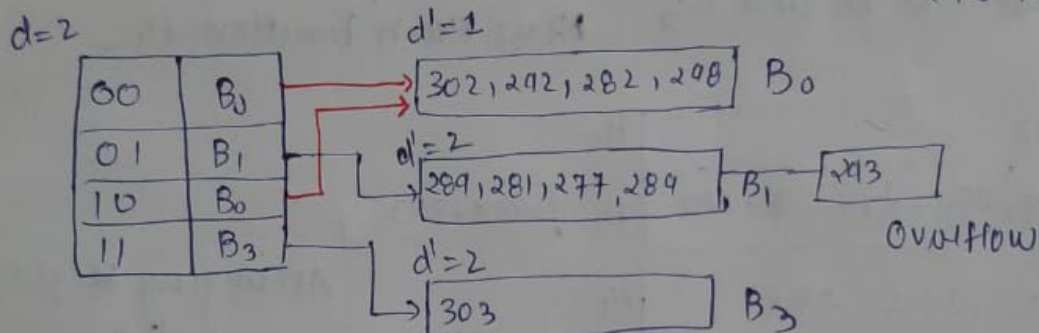
Soc

$$83 + 111 + 99 = 293$$

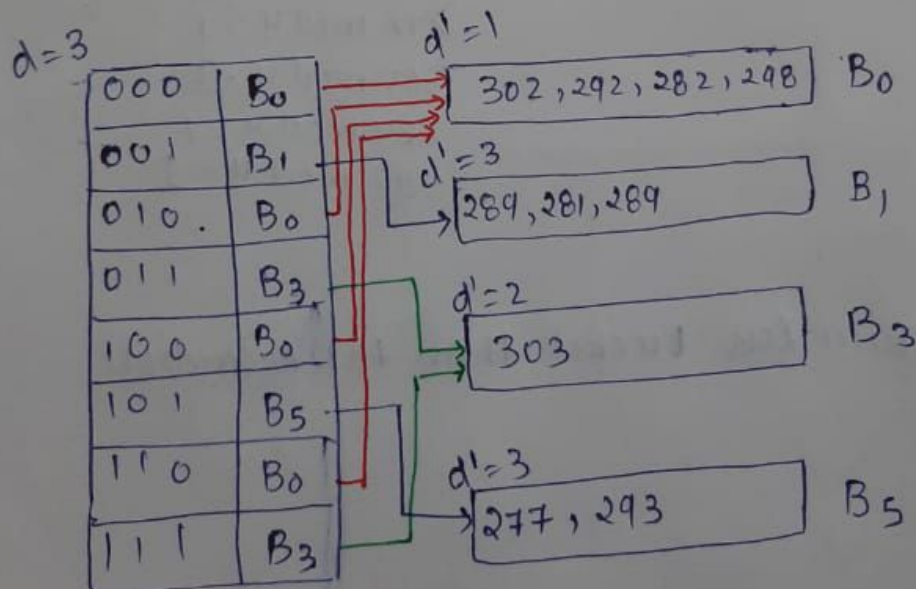
100100101

Adding into structure of Question (a)

$$293 \bmod 4 = 1 \quad B_1$$



Hash function: $(key) \bmod 8$



$$\begin{aligned} 289 \bmod 8 &= 1 \quad B_1 \\ 281 \bmod 8 &= 1 \quad B_1 \\ 277 \bmod 8 &= 5 \quad B_5 \\ 289 \bmod 8 &= 1 \quad B_1 \\ 293 \bmod 8 &= 5 \quad B_5 \end{aligned}$$

Redistribution
b/w B_1 and B_5

(5)

Adding 293 in structure of Question 6.

B_0	292	H_1
B_1	289, 281, 277, 303	H_2
B_2	302, 282, 298	H_1

→ 289, 293 overflow

So split B_1 to B_1 and B_3 New Hash function H_1

B_0	292	H_1
B_1	289, 281, 277, 303 289	H_1
B_2	302, 282, 298	H_1
B_3	303	H_1

289, 293

Arranging keys of B_1 and B_3

$$289 \bmod 4 = 1 \quad B_1$$

$$281 \bmod 4 = 1$$

$$277 \bmod 4 = 1$$

$$303 \bmod 4 = 3$$

$$289 \bmod 4 = 1$$

$$293 \bmod 4 = 1$$

Comparison:

- * Linear hashing requires less bucket than the extendible hashing. only bucket needed
- * No directory needed in linear hashing, whereas in both directory and bucket needed in extendible hashing.