

Sharing & Evaluating Networks

CPSC 433/533, Spring 2021

Anurag Khandelwal

Administrivia

- I hope by now, you have:
 - **All:** Checked Canvas for all the course policies and administrivia
 - **CPSC 533:** Found resources on how to read/critique papers

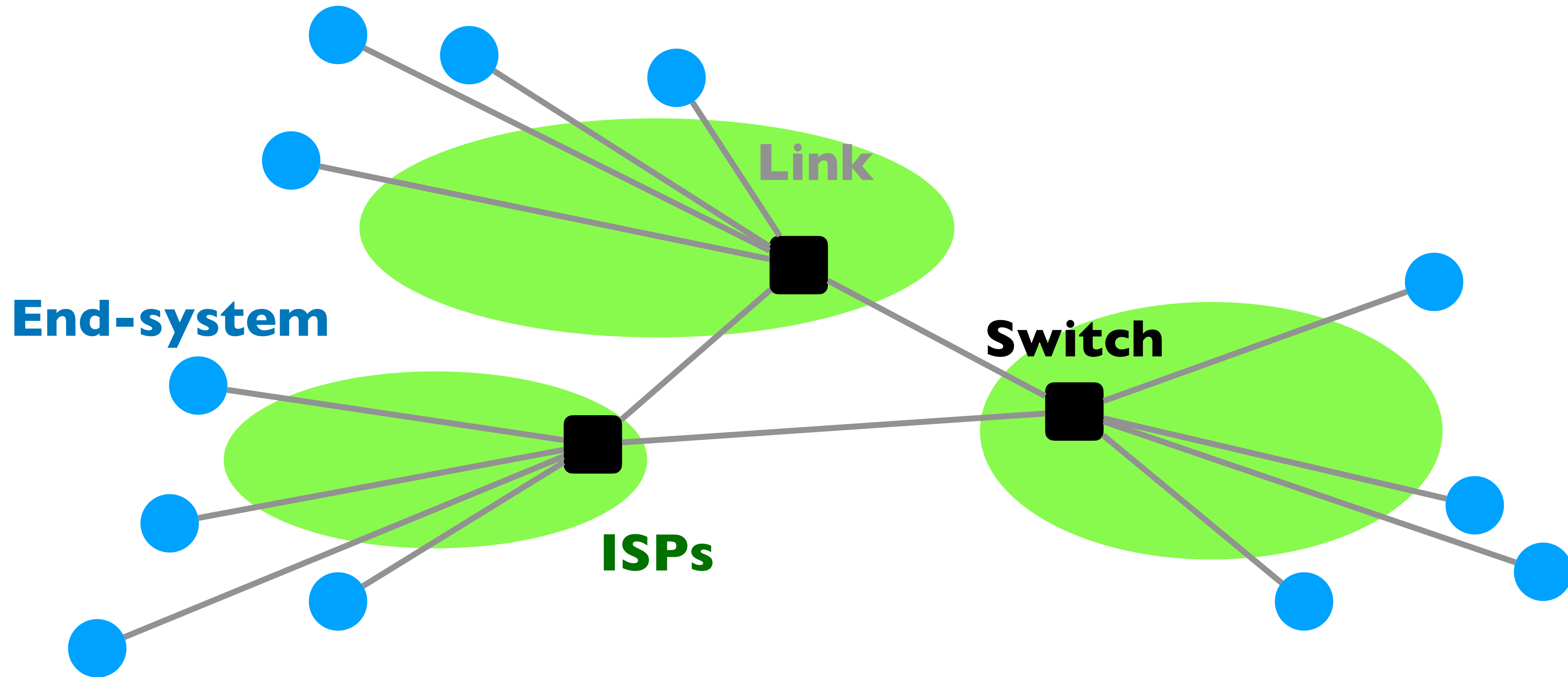
Today's Agenda

- More details on What & How of the Internet
- Sharing a network
- Evaluating a network

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- More details on What, How and Why of the Internet
- Sharing a network
- Evaluating a network

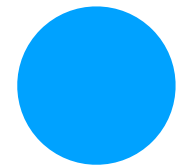
The What



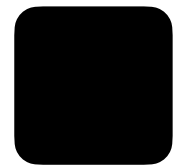
Digging Deeper into the “What”

The Last Mile

Home PC



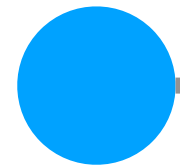
Switch



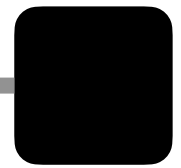
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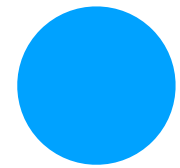
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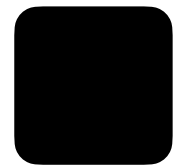
Edge Networks: The Last Mile

Telephony Access Technology

Home PC



Switch



Edge Networks: The Last Mile

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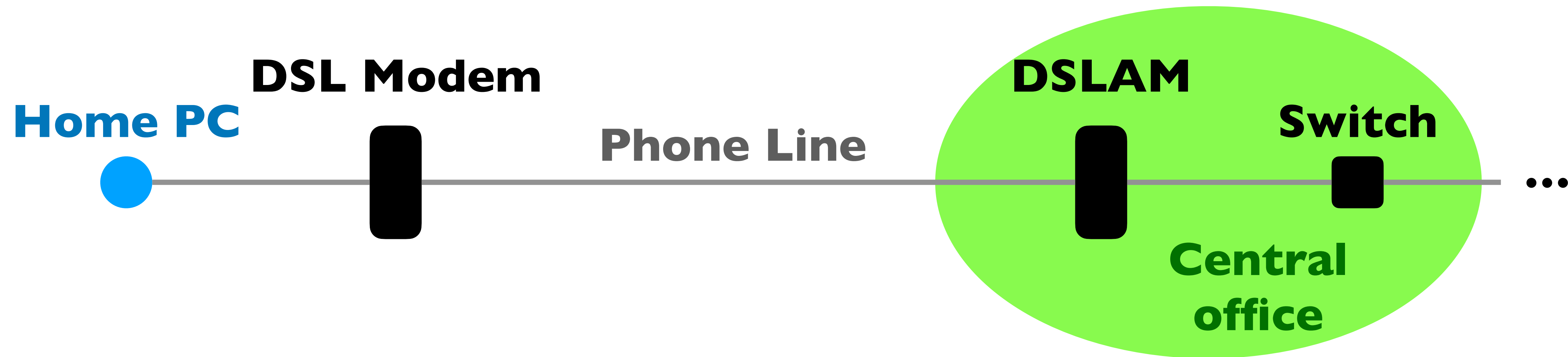
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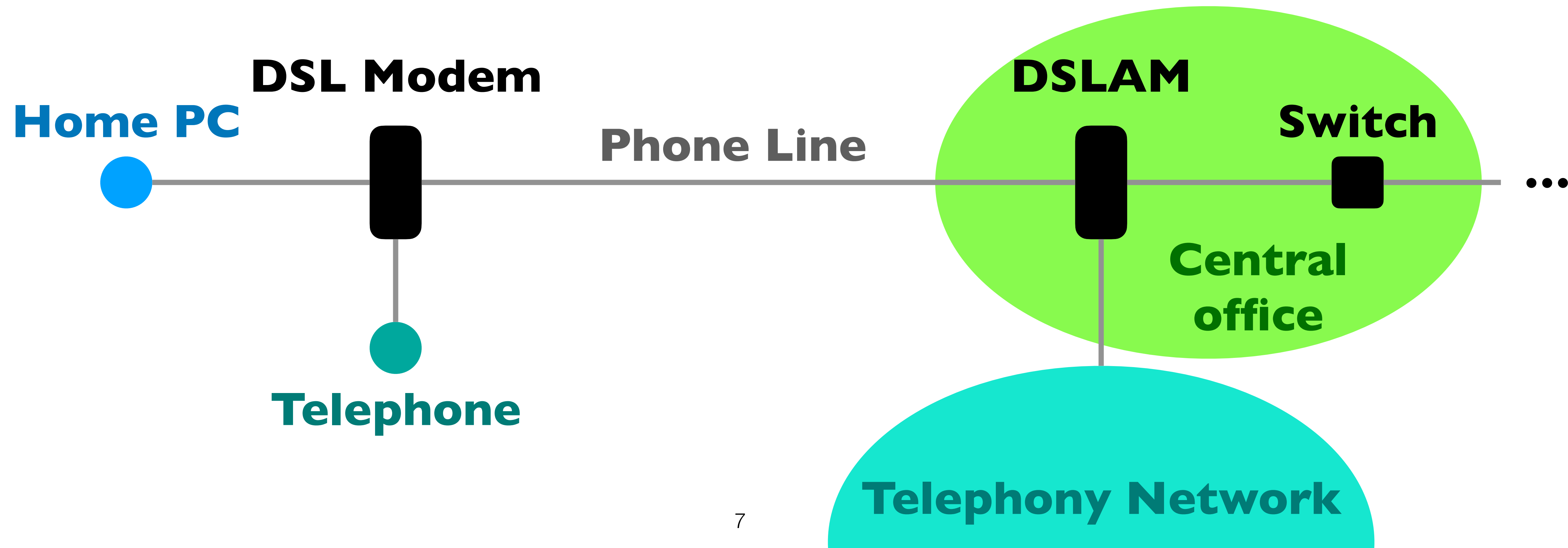
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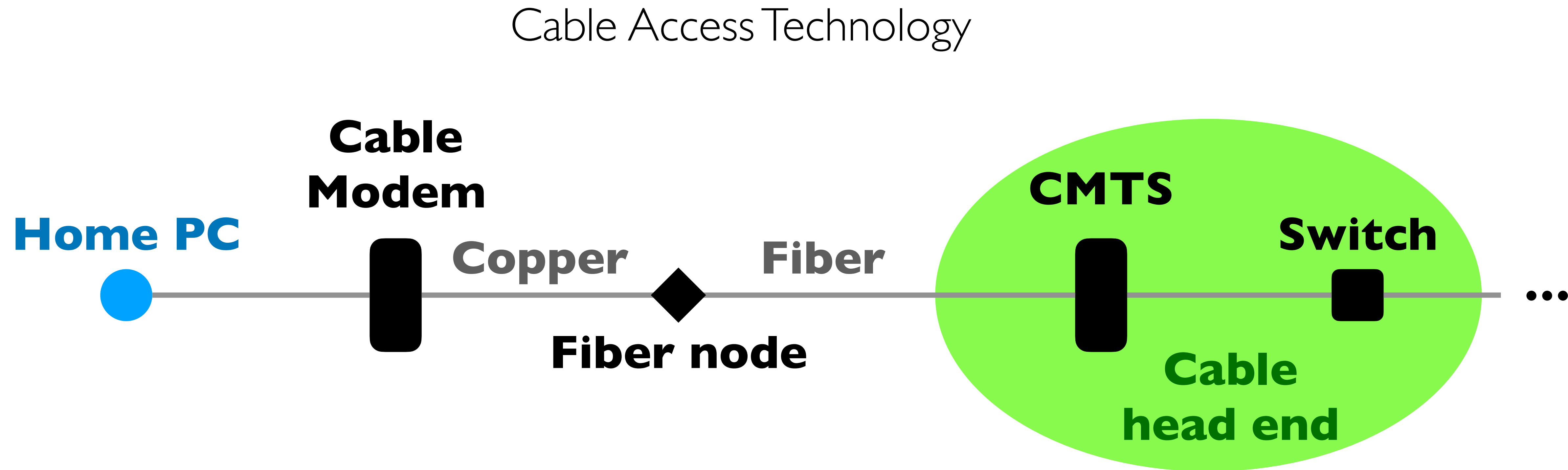


Edge Networks: The Last Mile

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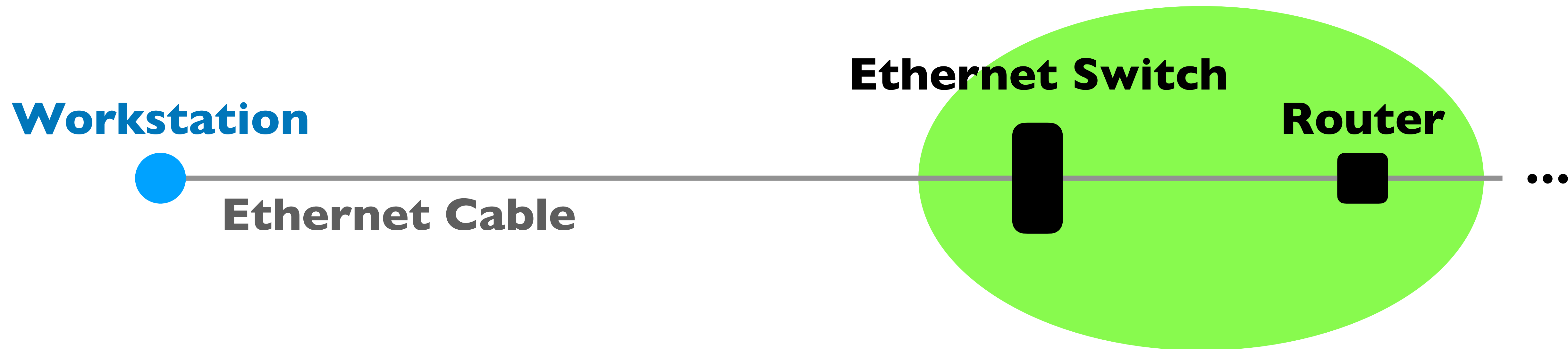


Edge Networks: The Last Mile



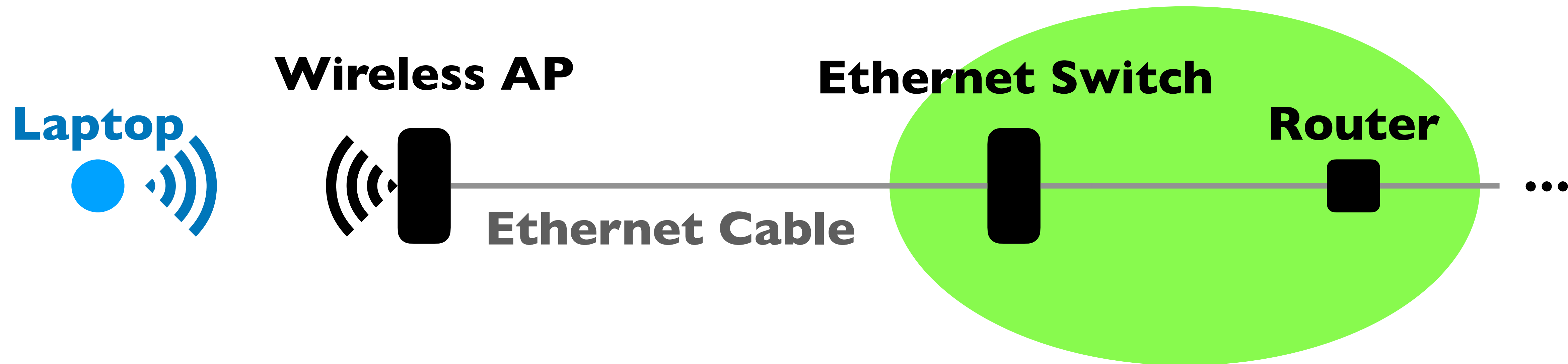
Edge Networks: The Last Mile

Ethernet Access Technology

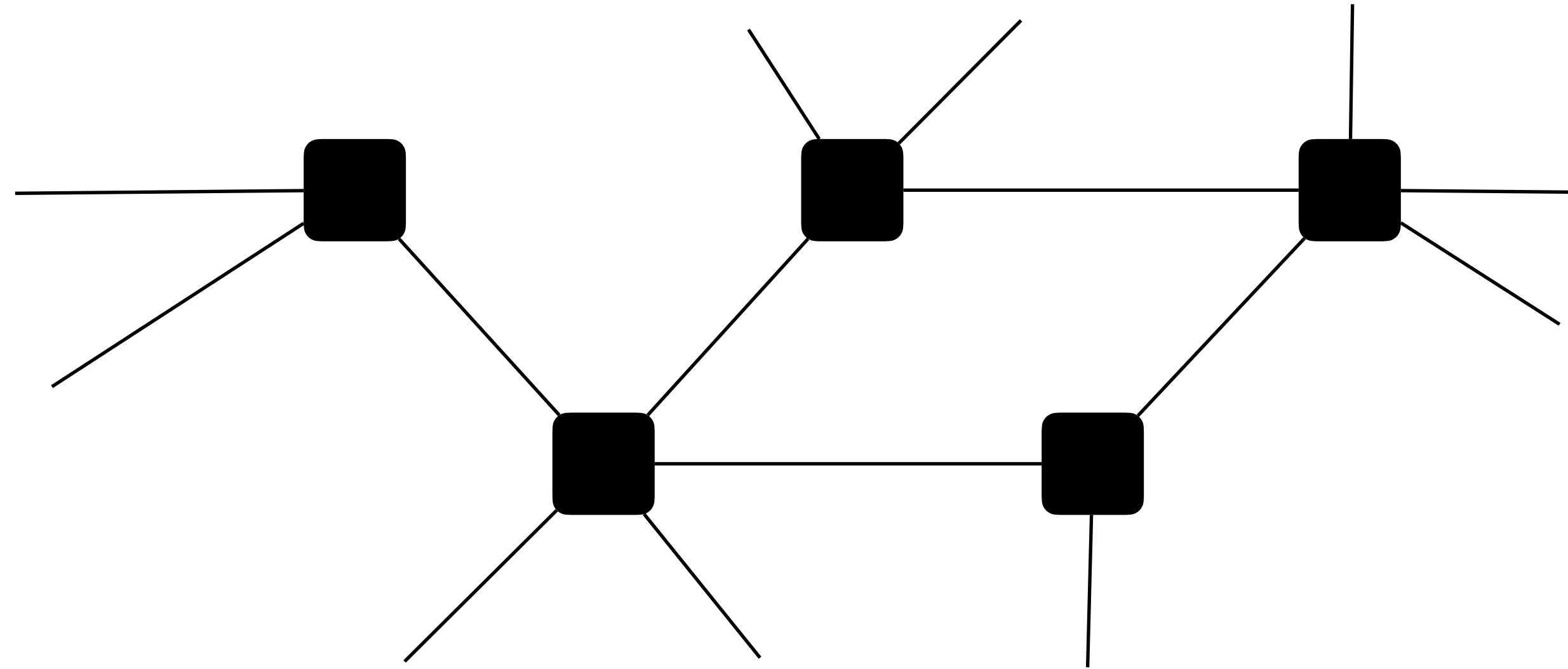


Edge Networks: The Last Mile

Ethernet Access Technology

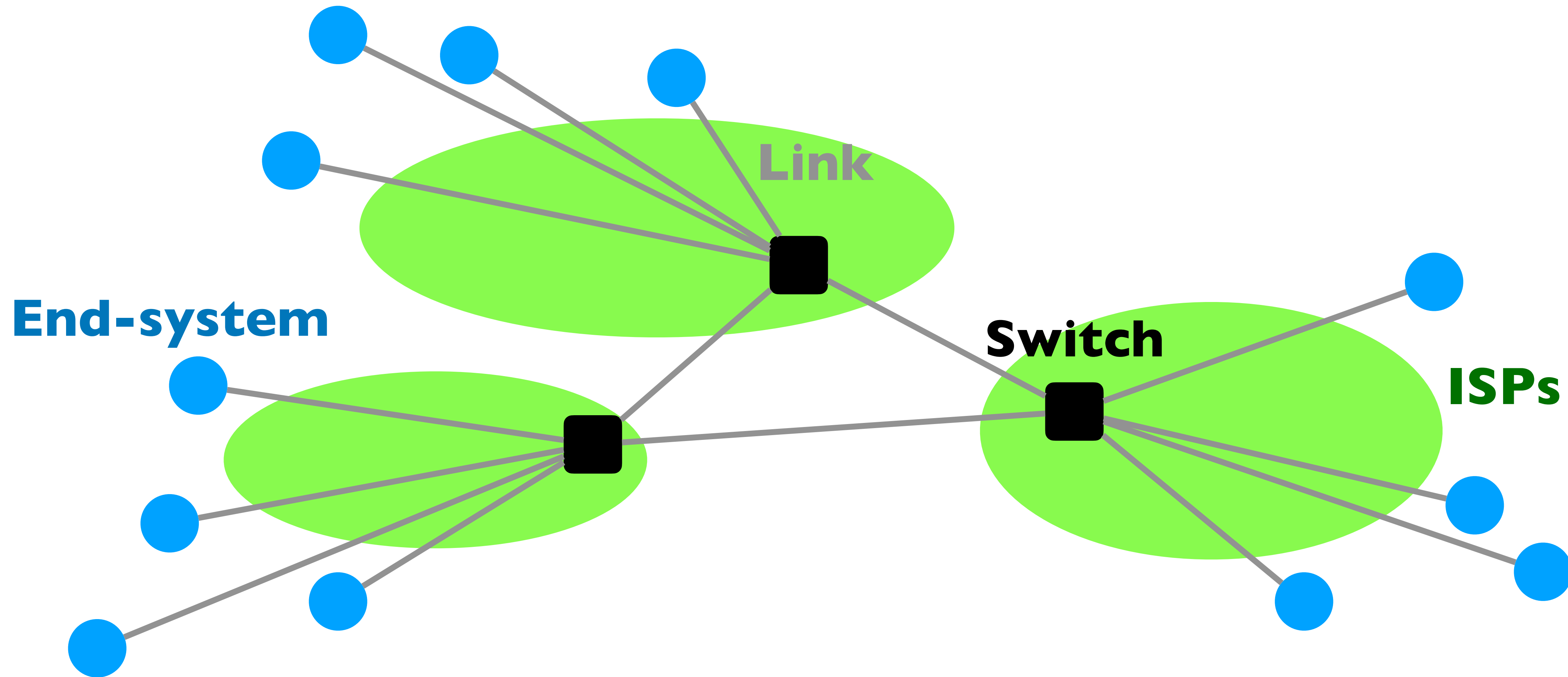


Core Network

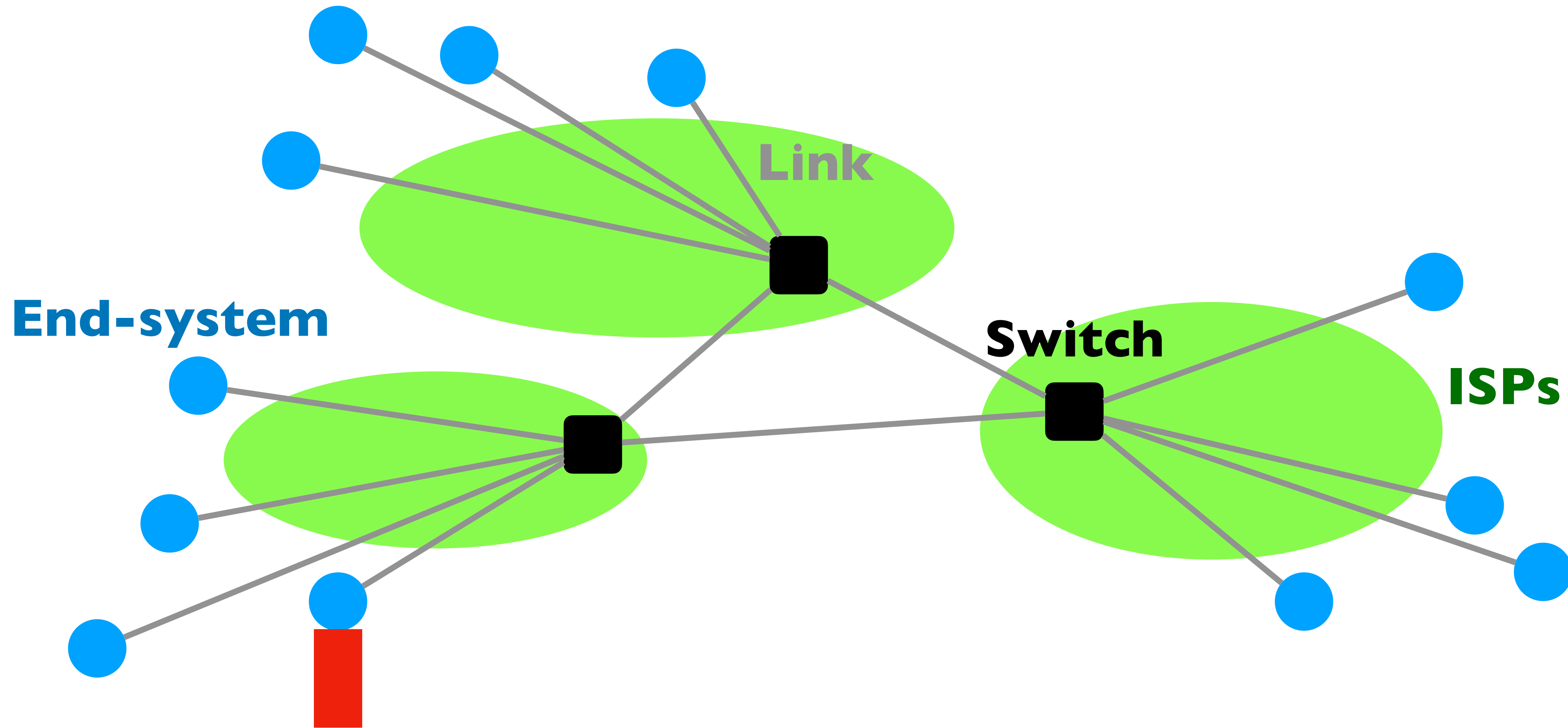


Switched networks enable efficient scaling!

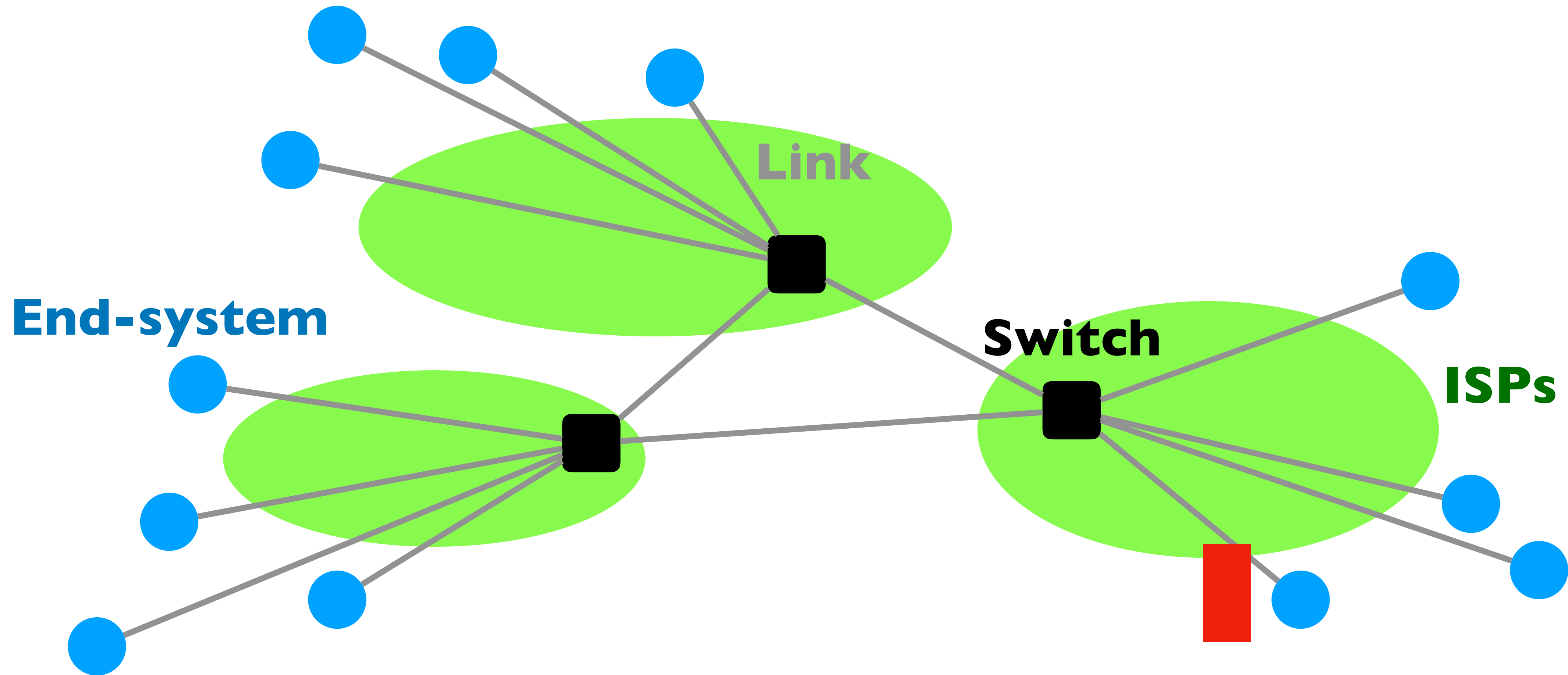
The How



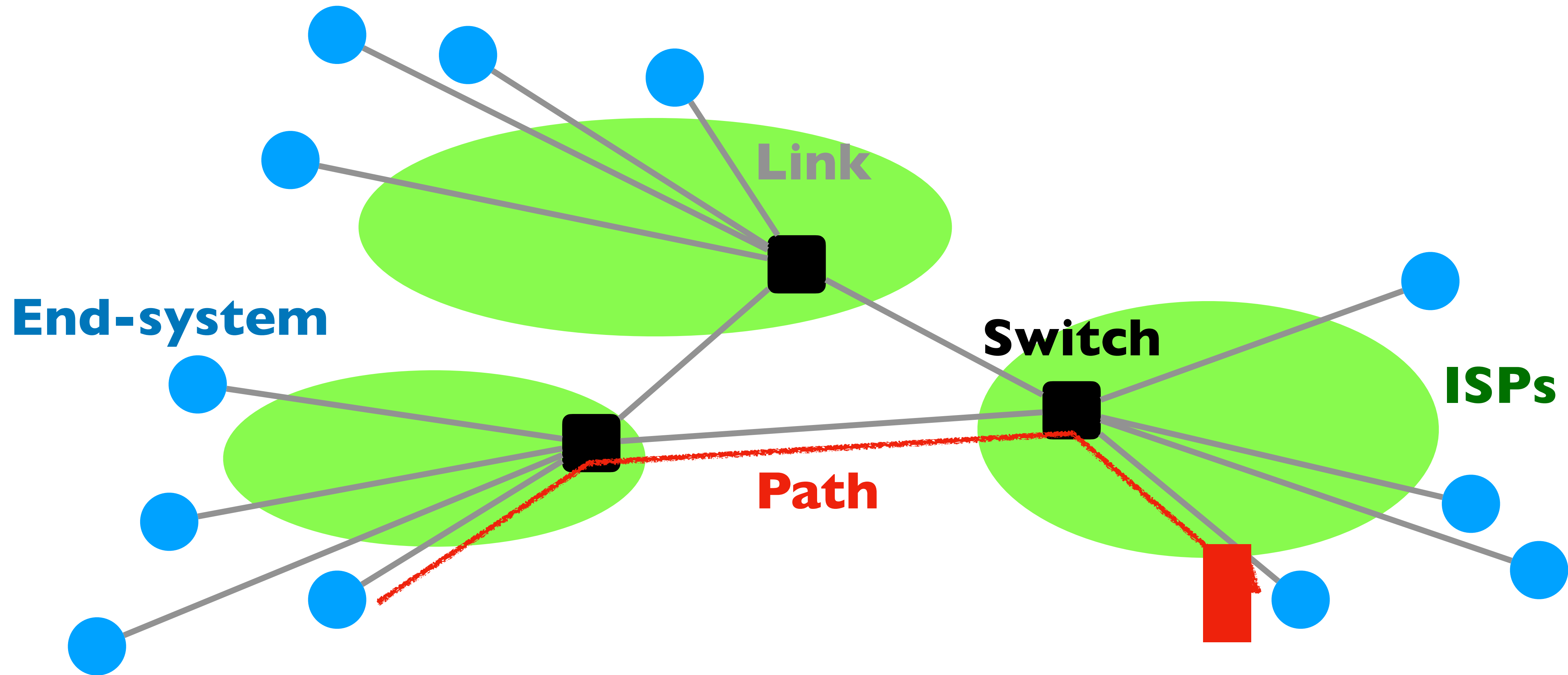
The How



The How



The How



Digging Deeper into the How

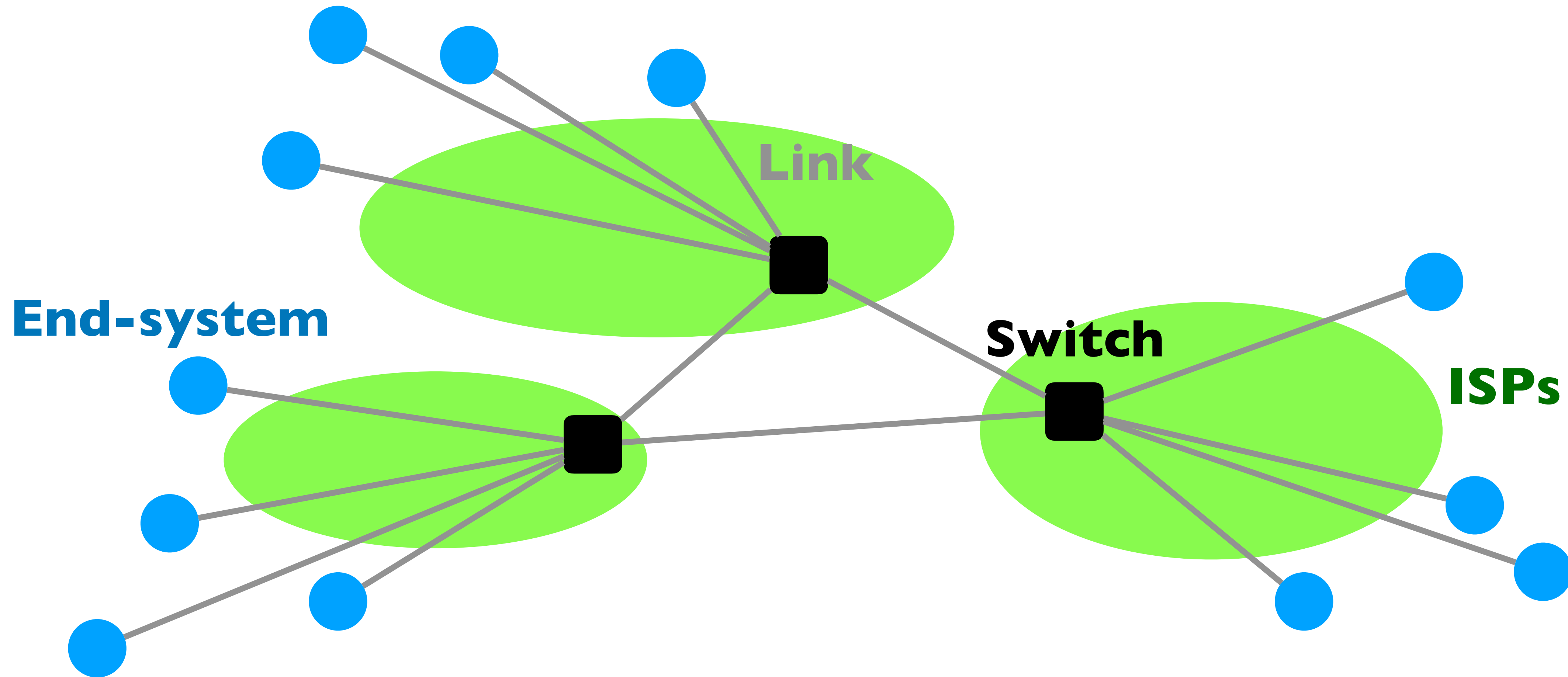
- **Locating the destination:** Naming, Addressing
- **Finding a path to the destination:** Routing
- **Sending data to the destination:** Forwarding
- **Handling cases where things go horribly wrong:** Distributed routing, congestion control
- We will spend (almost) the entire course to understand these concepts!

Questions?

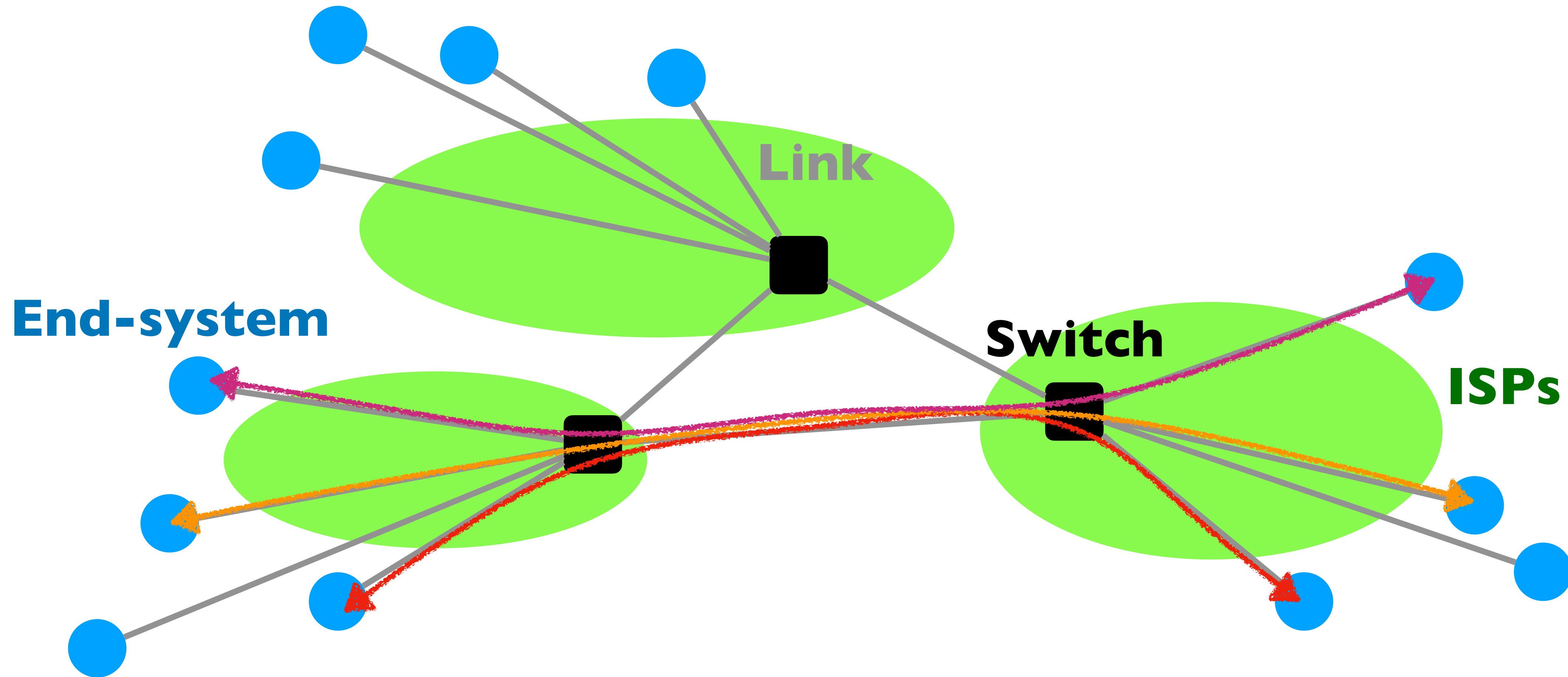
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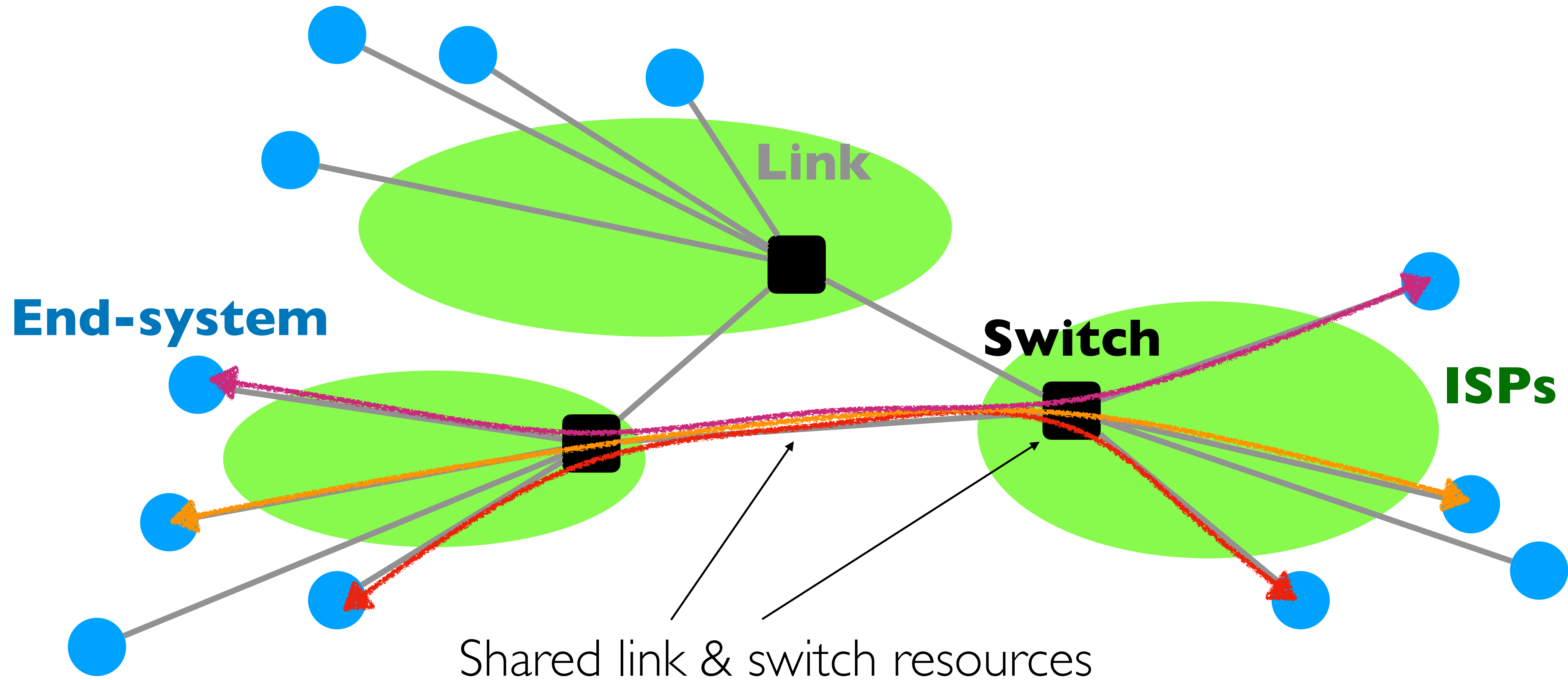
Sharing Network Resources



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Sharing Network Resources

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 - Switch processing resources, link bandwidths
 - **Link:** “Pipe” abstraction
 - **Bandwidth:** Number of bits sent per unit time
 - **Network traffic:** “flows”



How can we share?

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- Suggestions?

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 - 100MB in the 1st second, 10MB in the second, etc.

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- First approach: **Reservations**
 - Reserve bandwidth needed in advance
 - Must reserve at **every** link or switch in path
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- **Peak** bandwidth?
 - An app may generate data at varying rates
 - 100MB in the 1st second, 10MB in the second, etc.
 - Must reserve for peak (e.g., 100MB/s)

Reservation: Circuit Switching

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- One way of implementing reservations in networks

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Grandma

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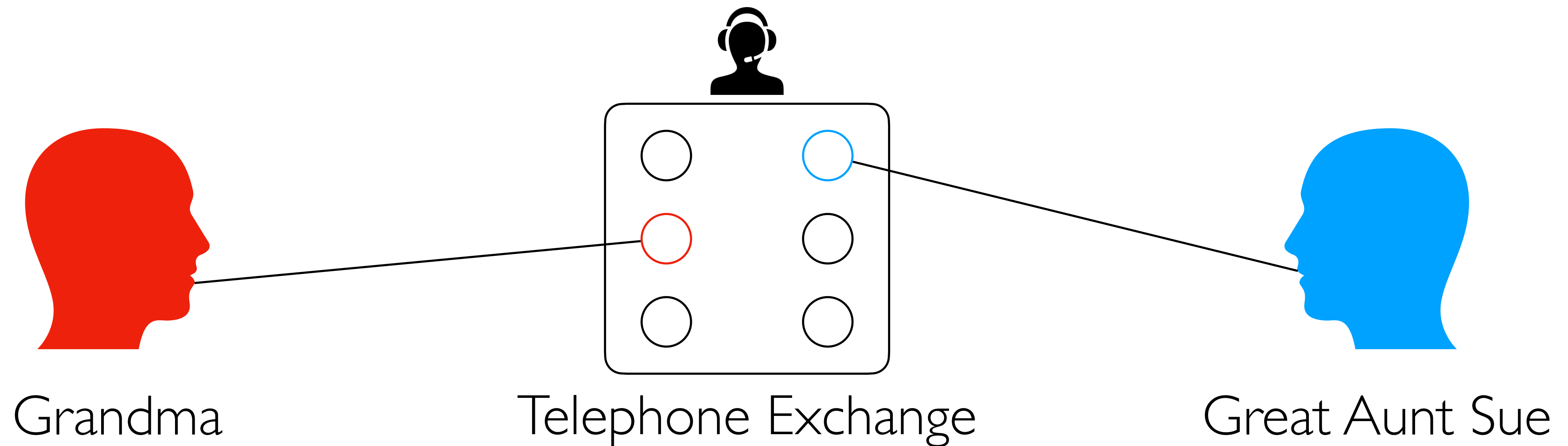
Grandma



Great Aunt Sue

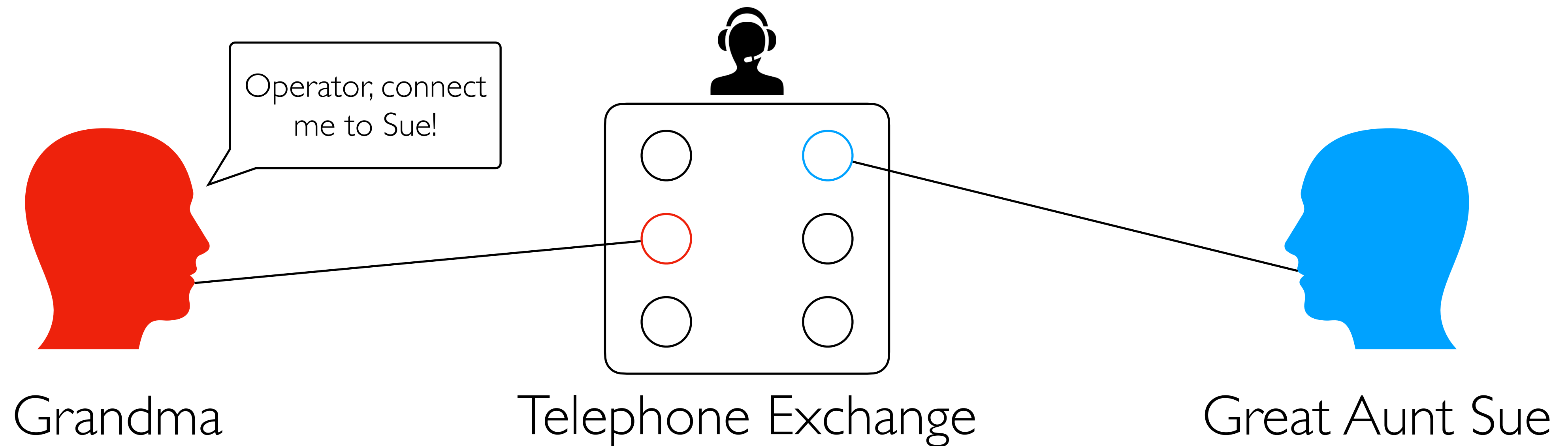
Reservation: Circuit Switching

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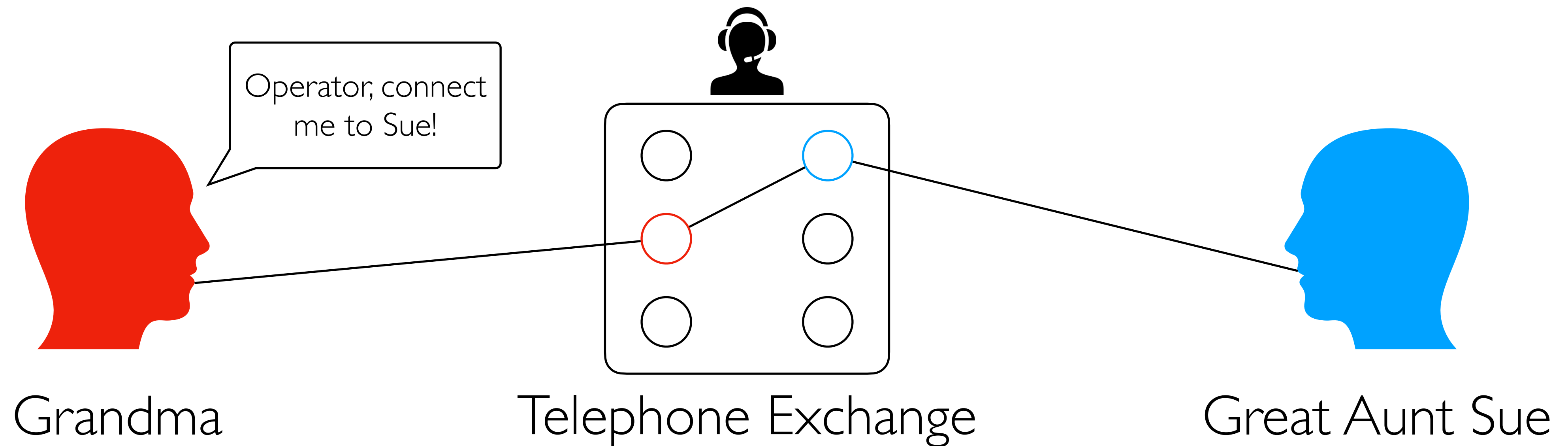
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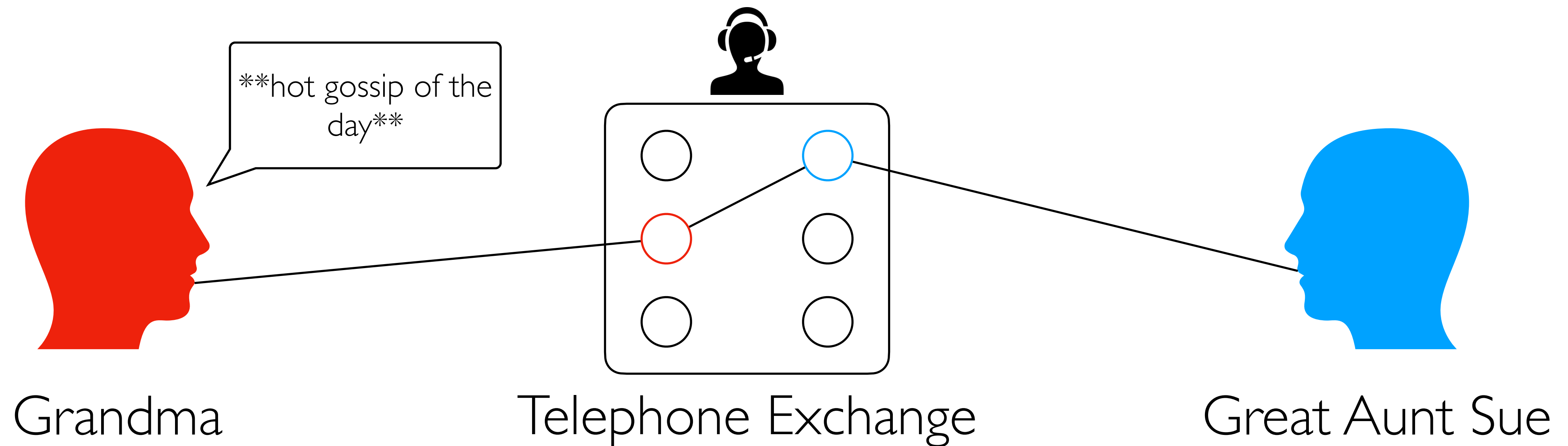
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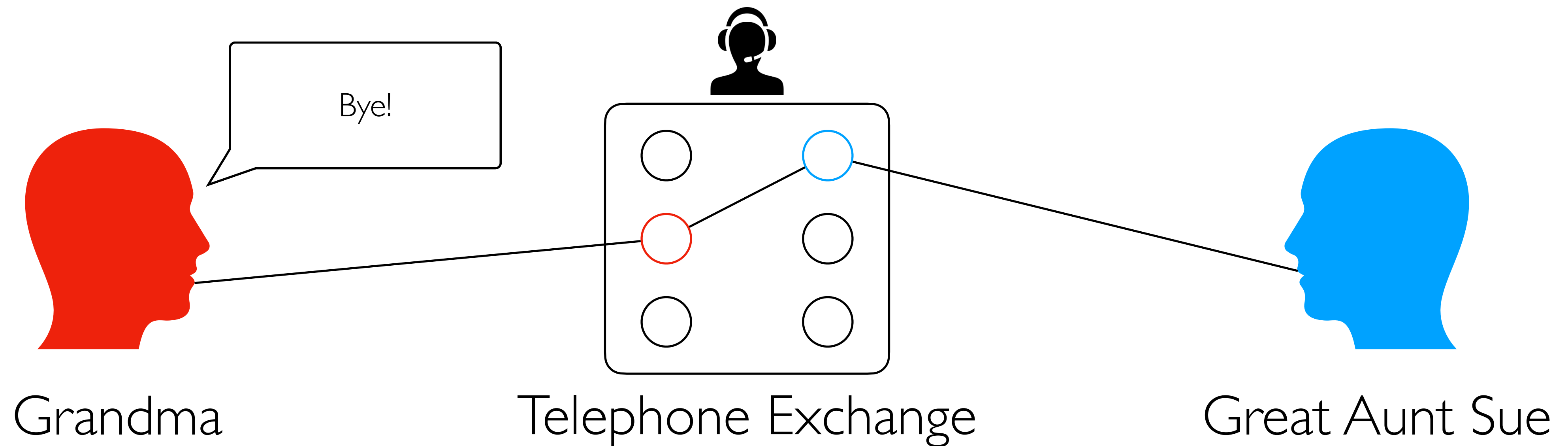
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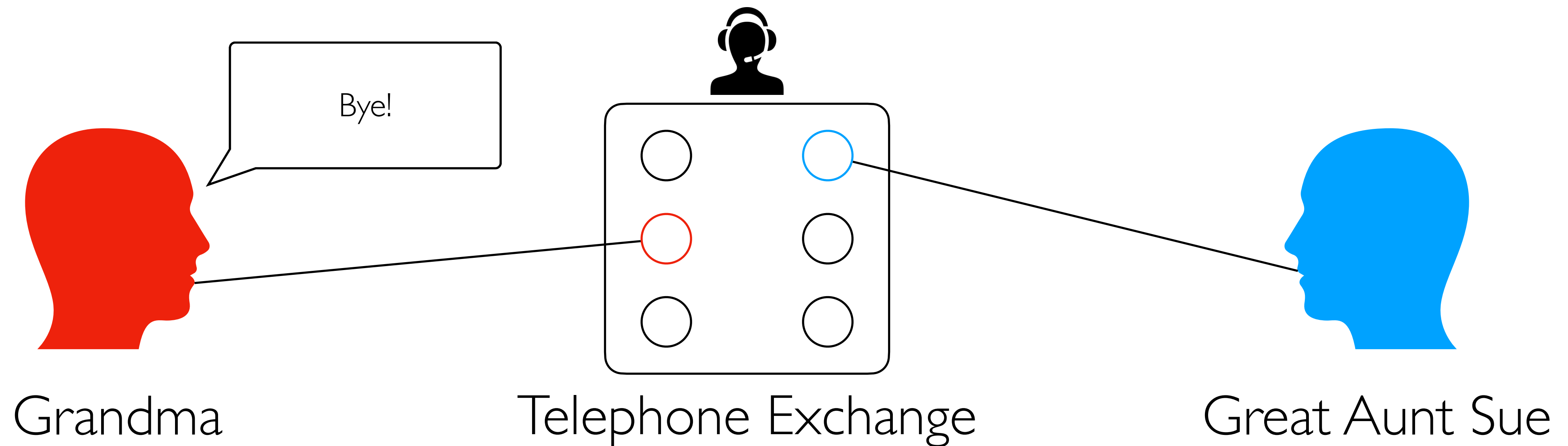
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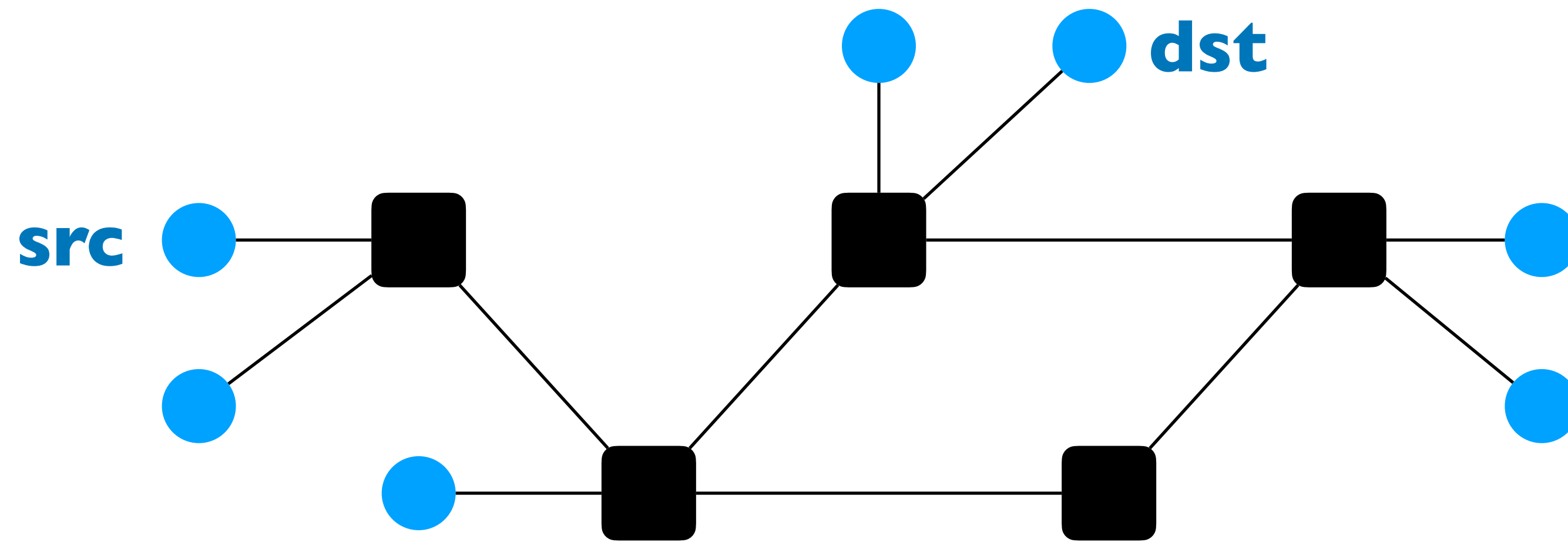


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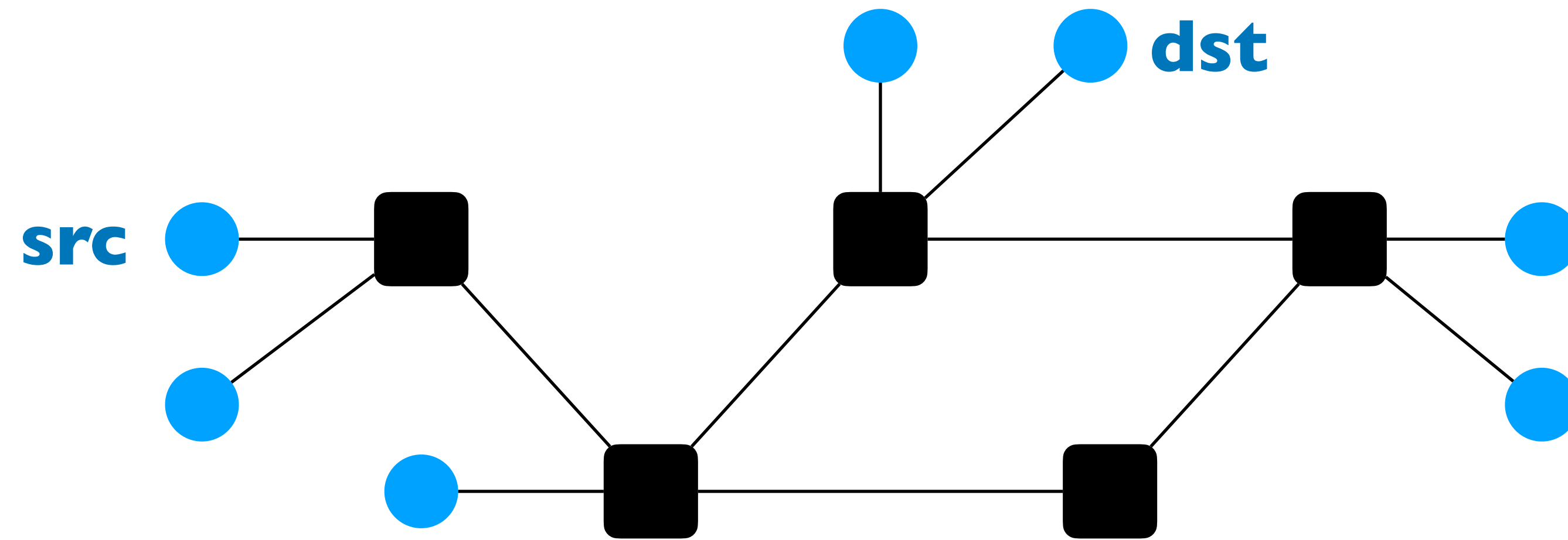
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Circuit Switching

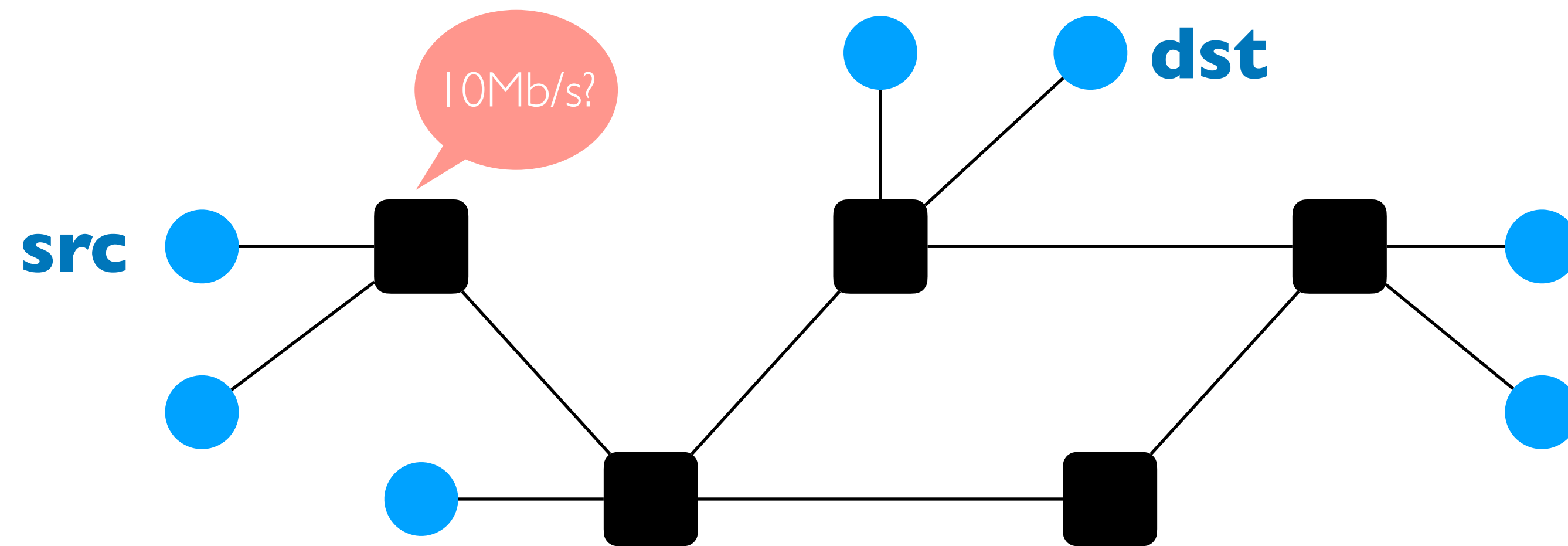


Circuit Switching



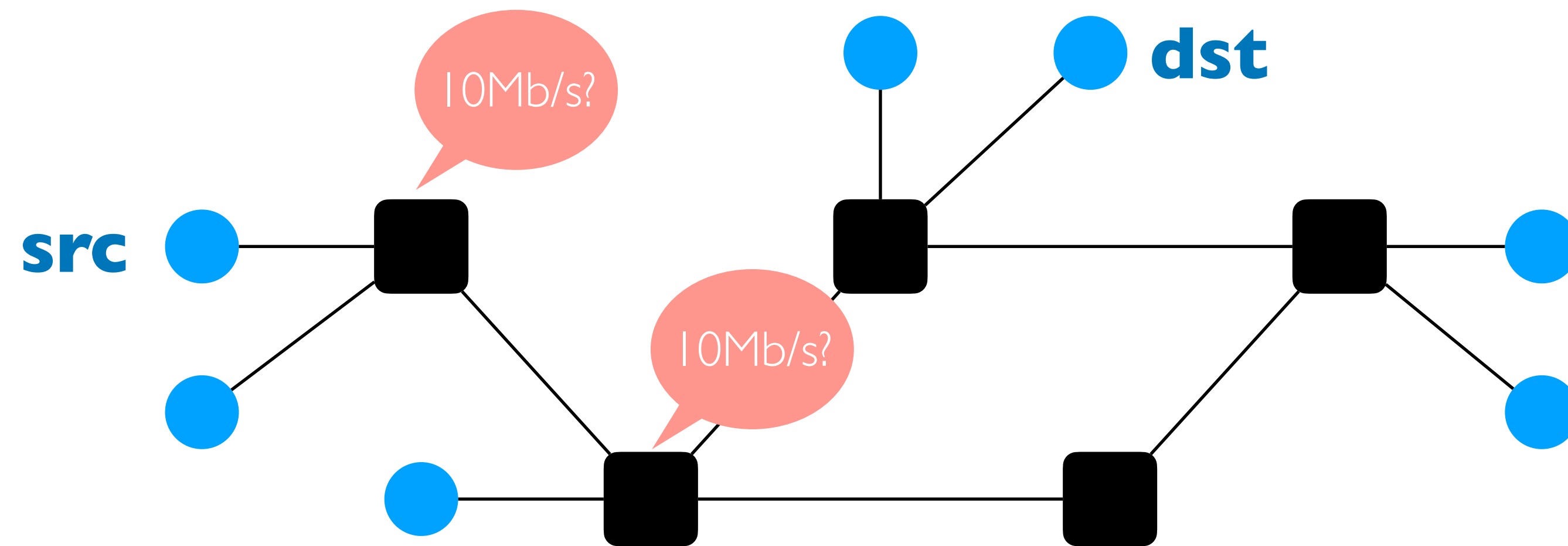
1. The **src** sends a reservation request to **dst**

Circuit Switching



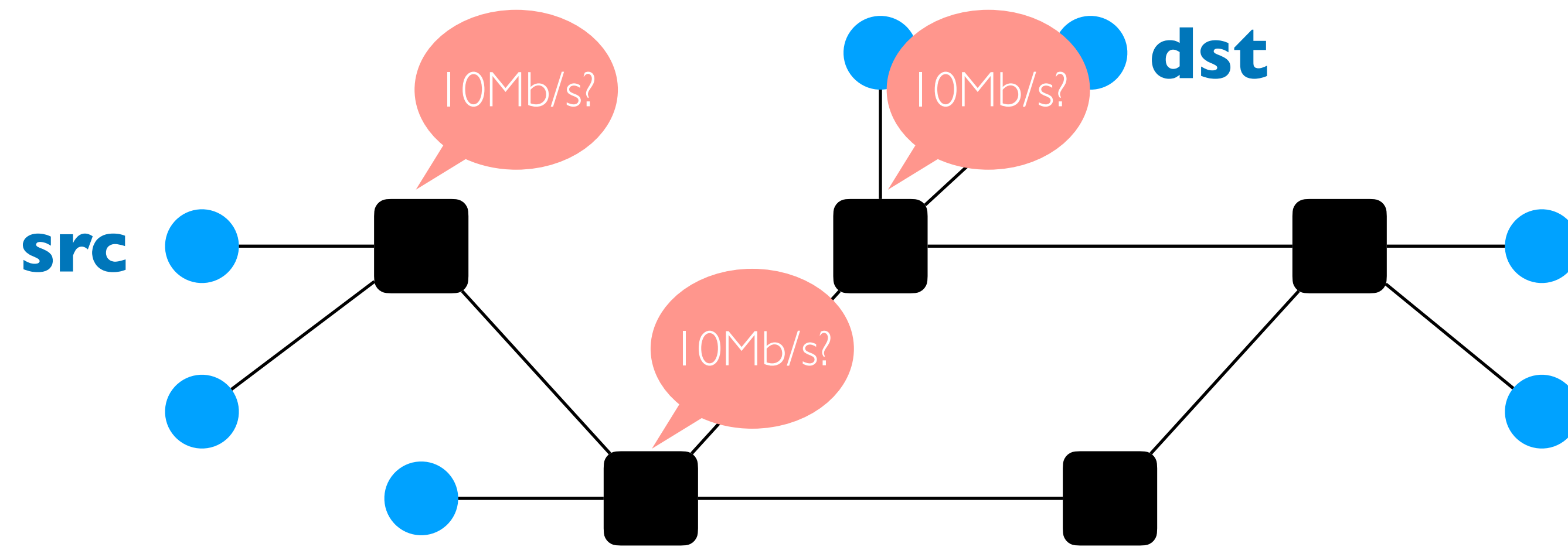
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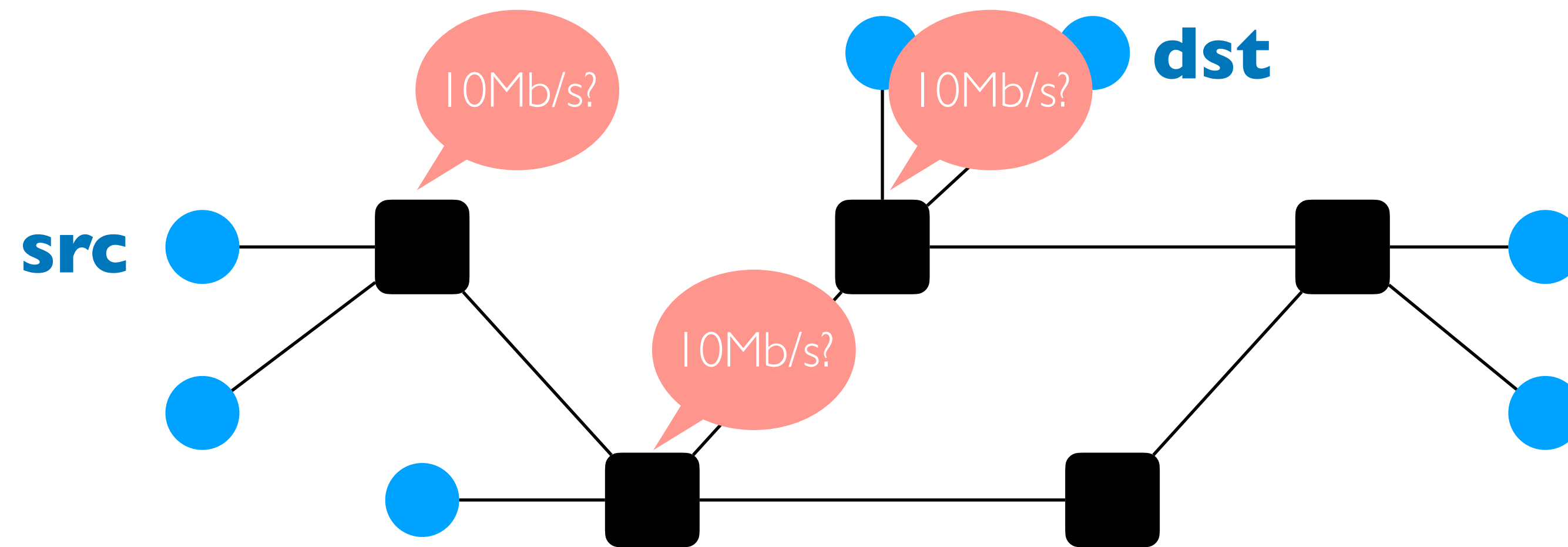
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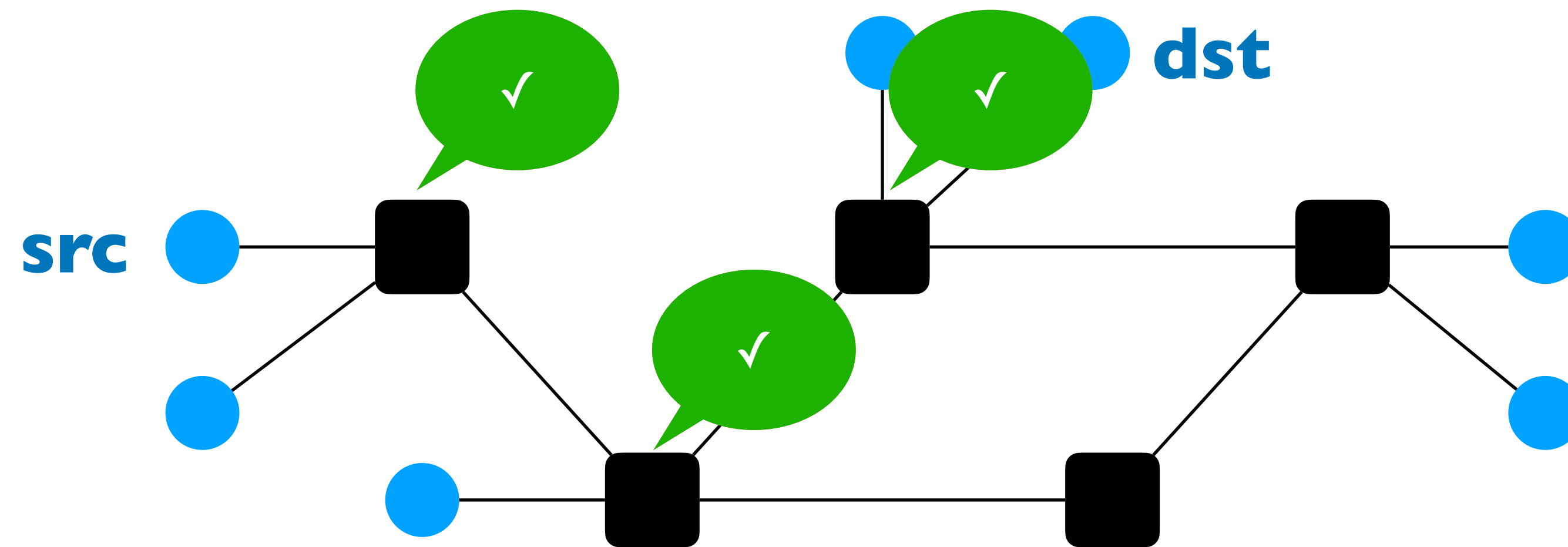
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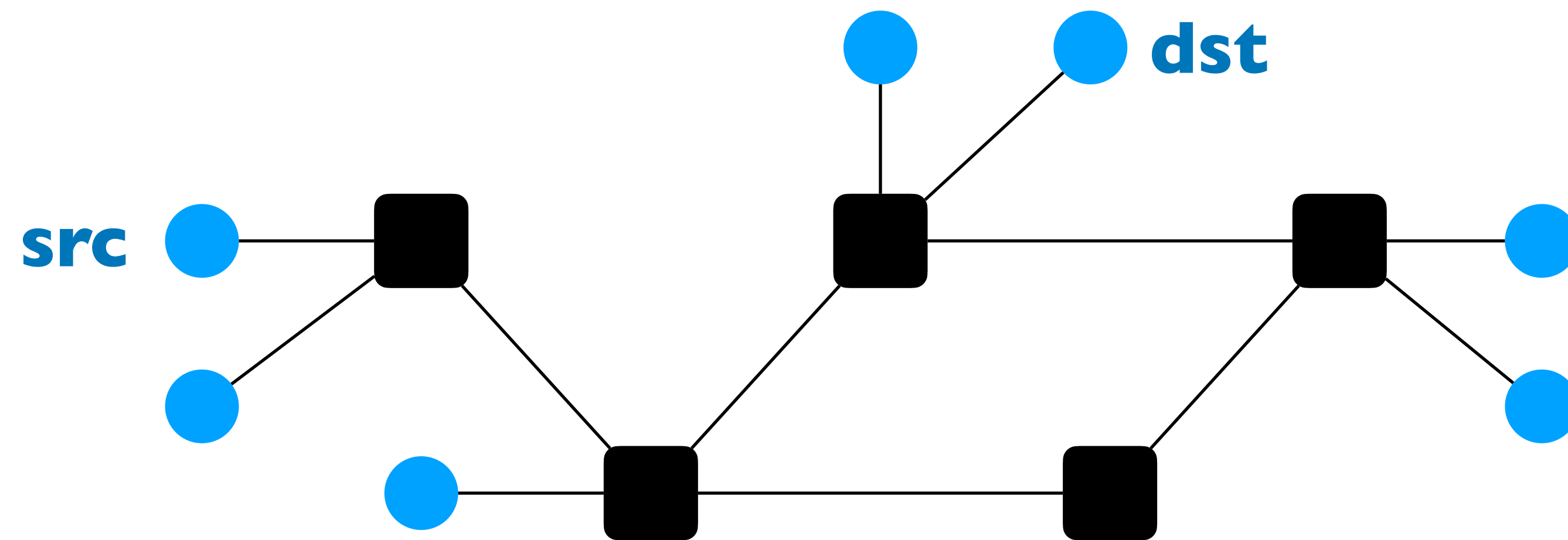
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2. Switches "establish a circuit"

Circuit Switching



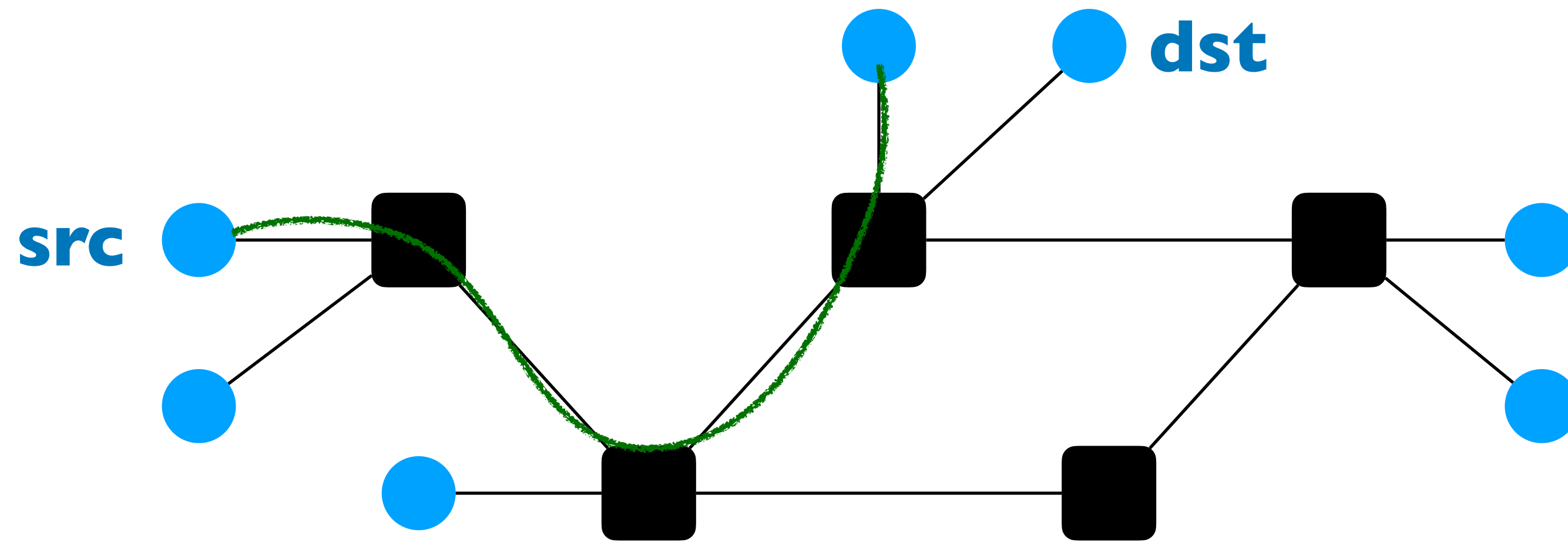
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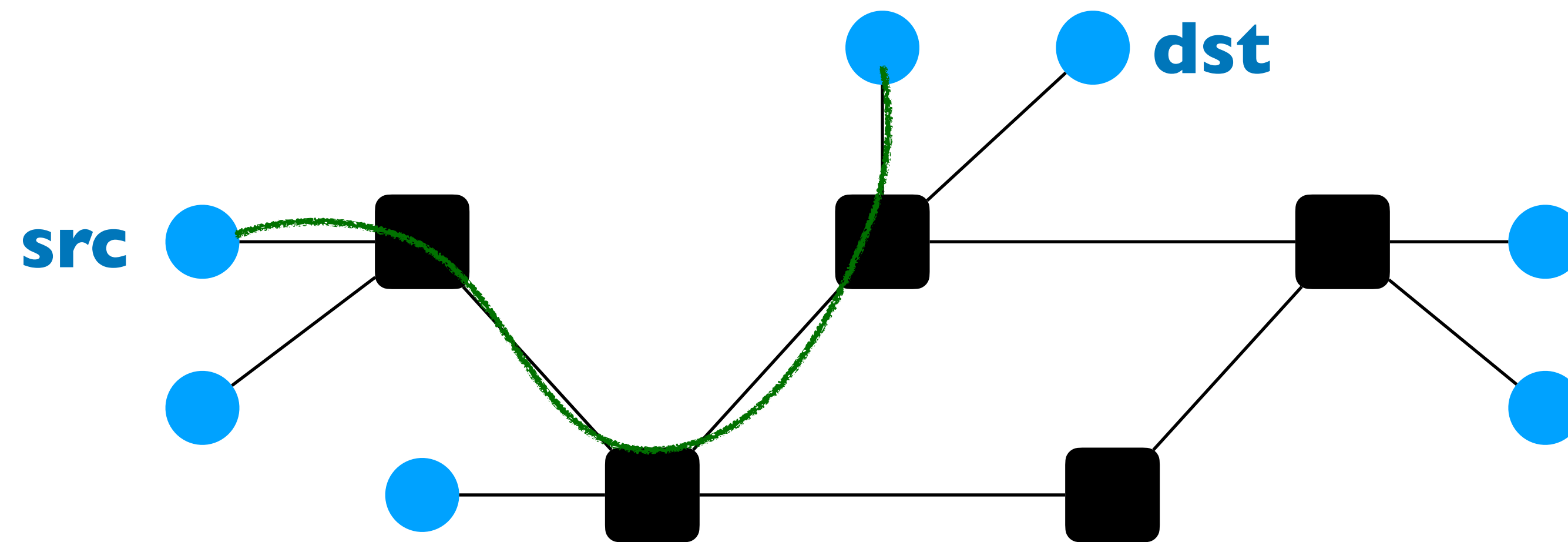
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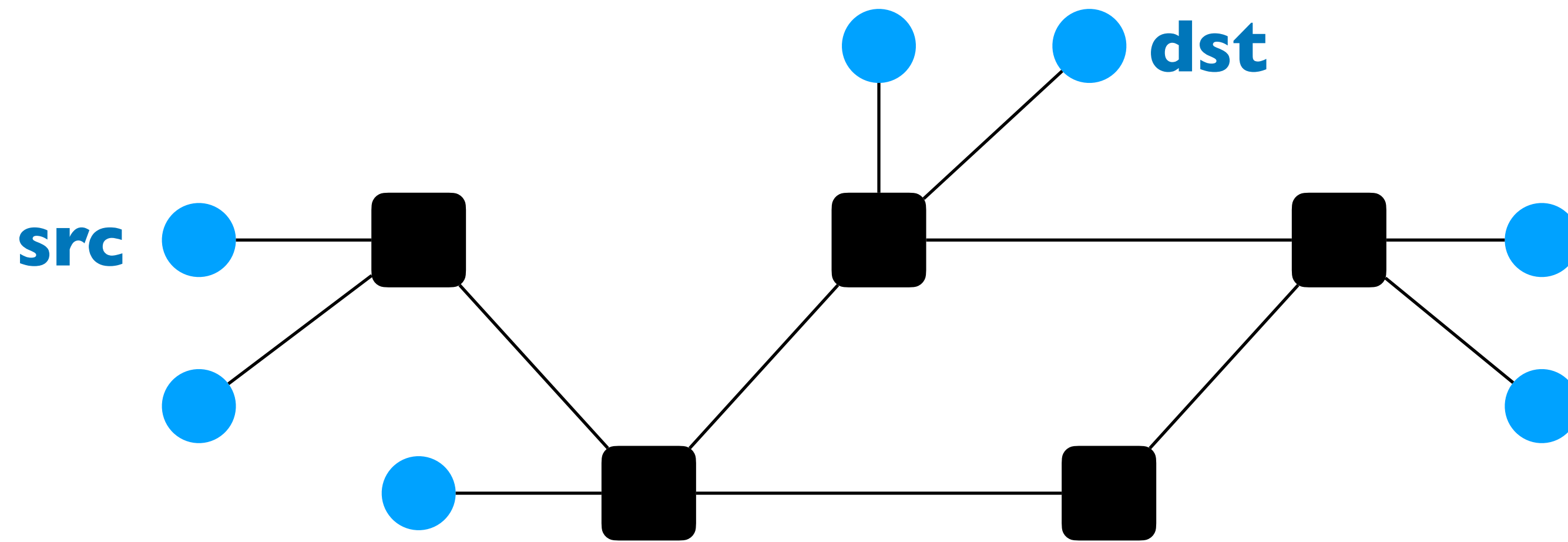
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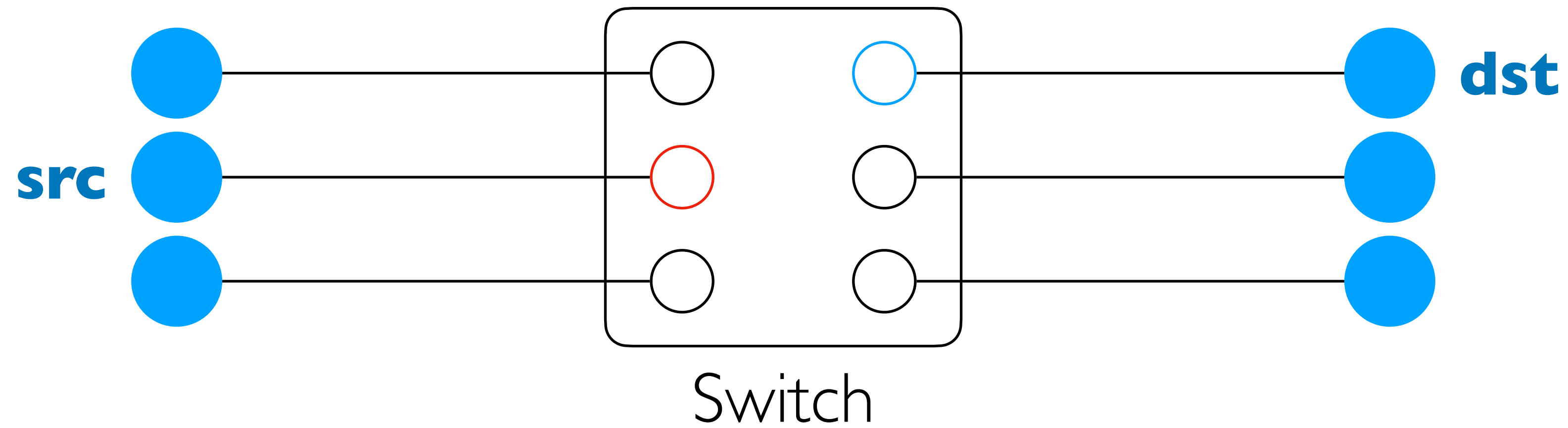
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Circuit Switching

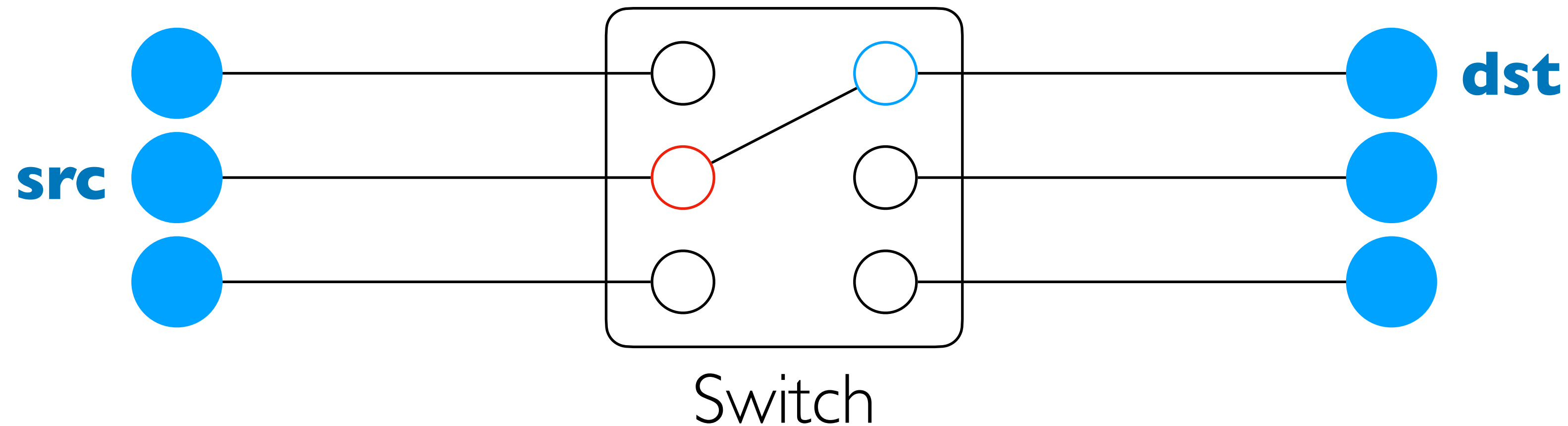


1. The **src** sends a reservation request to **dst**
2. Switches “establish a circuit”
3. The **src** starts sending data
4. The **src** sends a “teardown circuit” message

Reservation: Circuit Switching

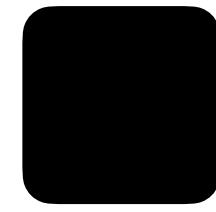
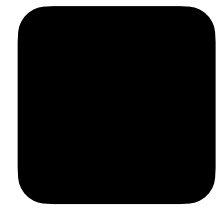


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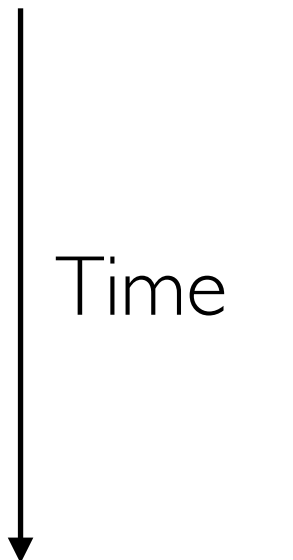
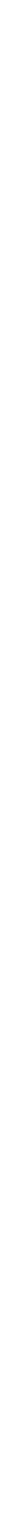
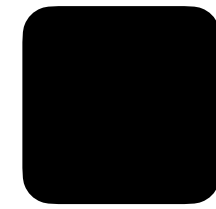
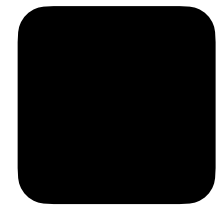


Reservation establishes a “circuit” within a switch

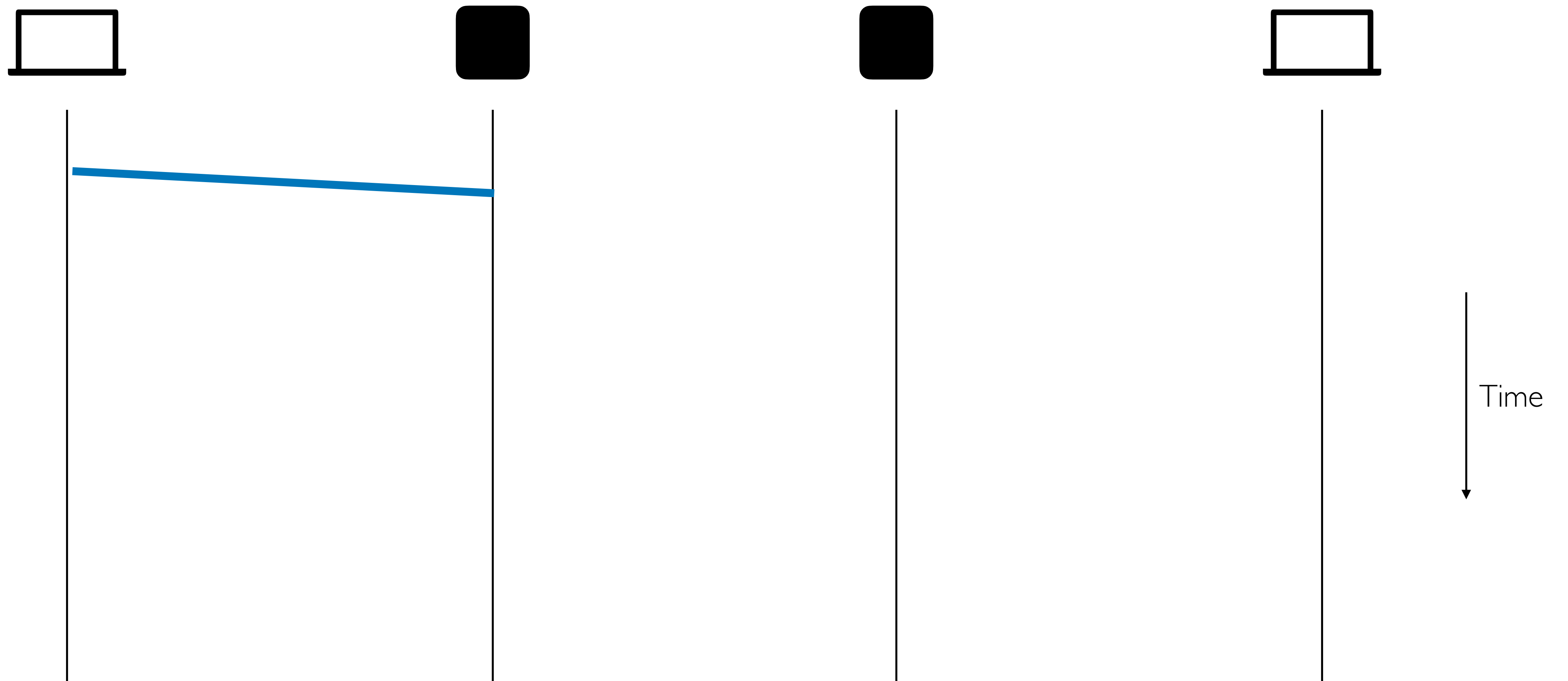
Timing in Circuit Switching



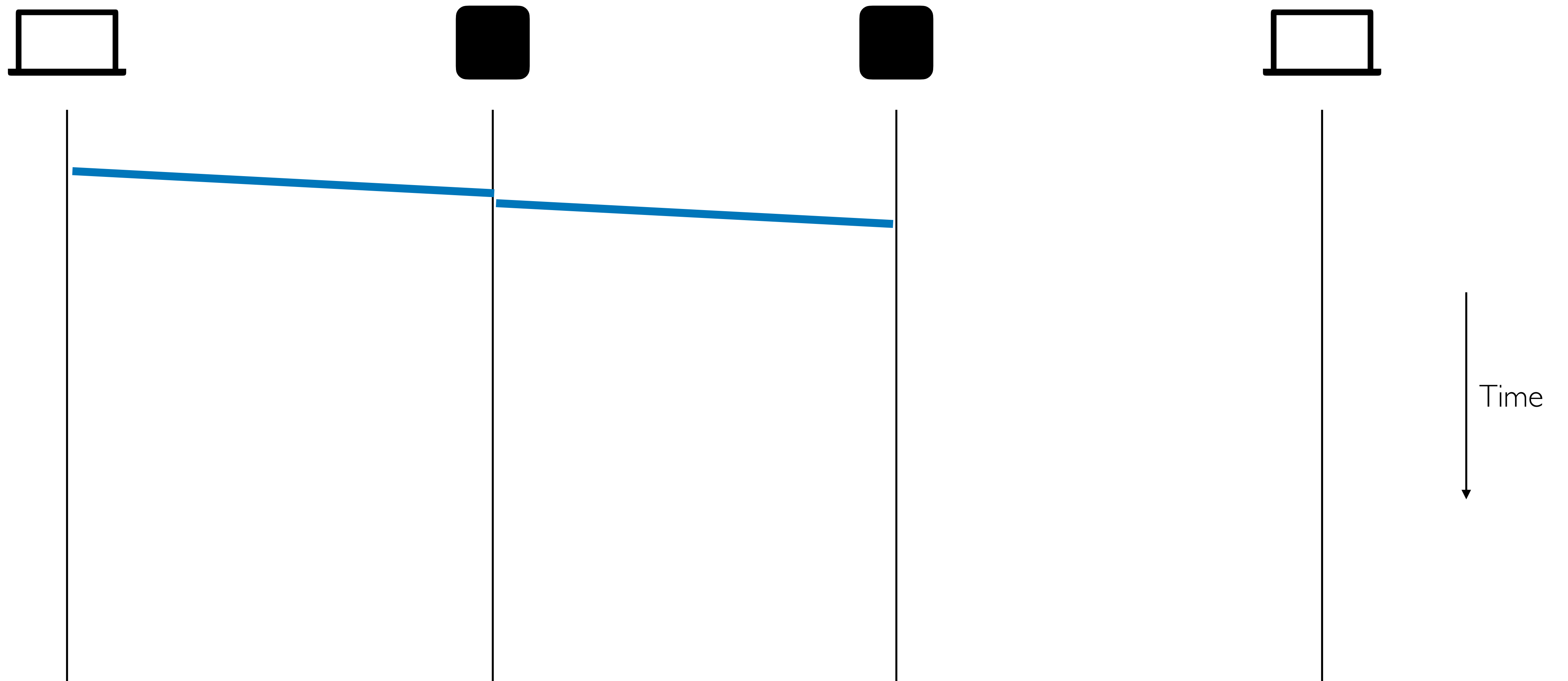
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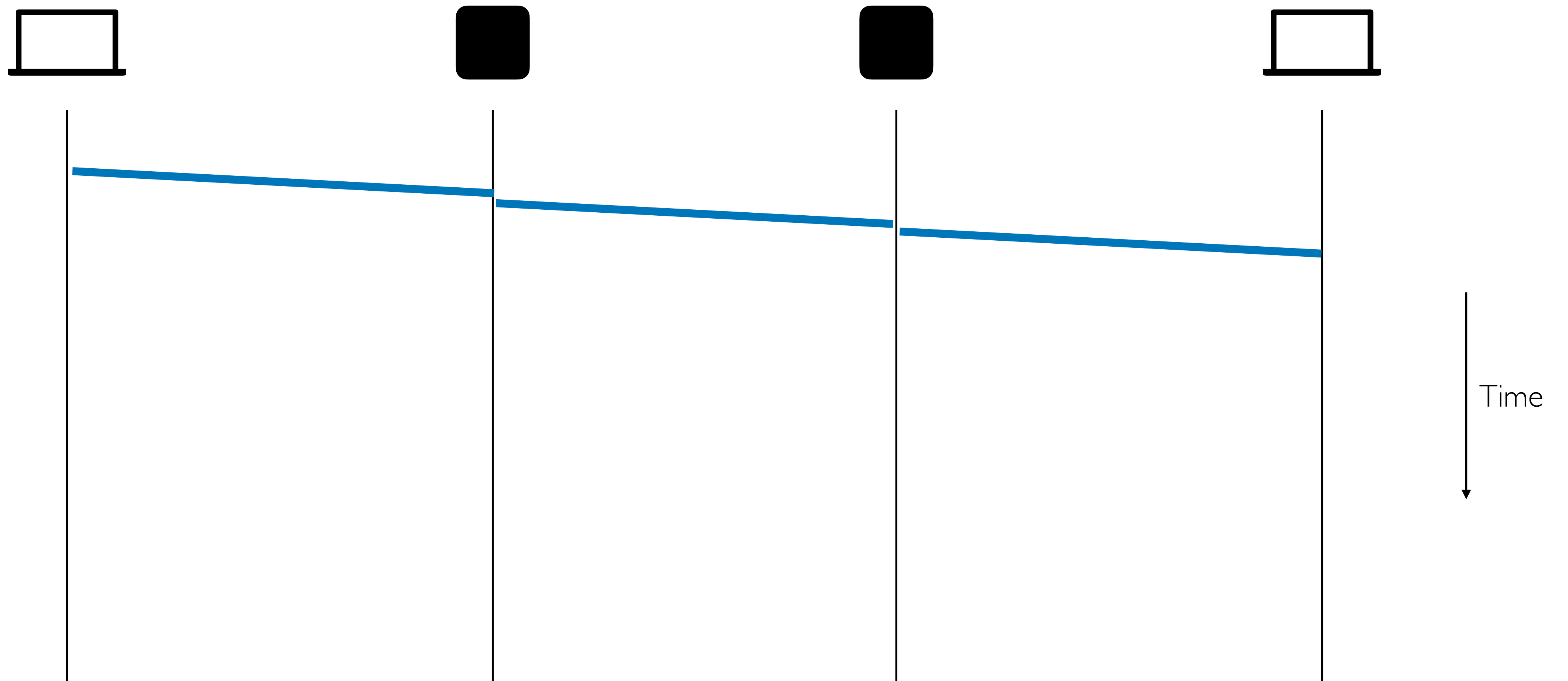
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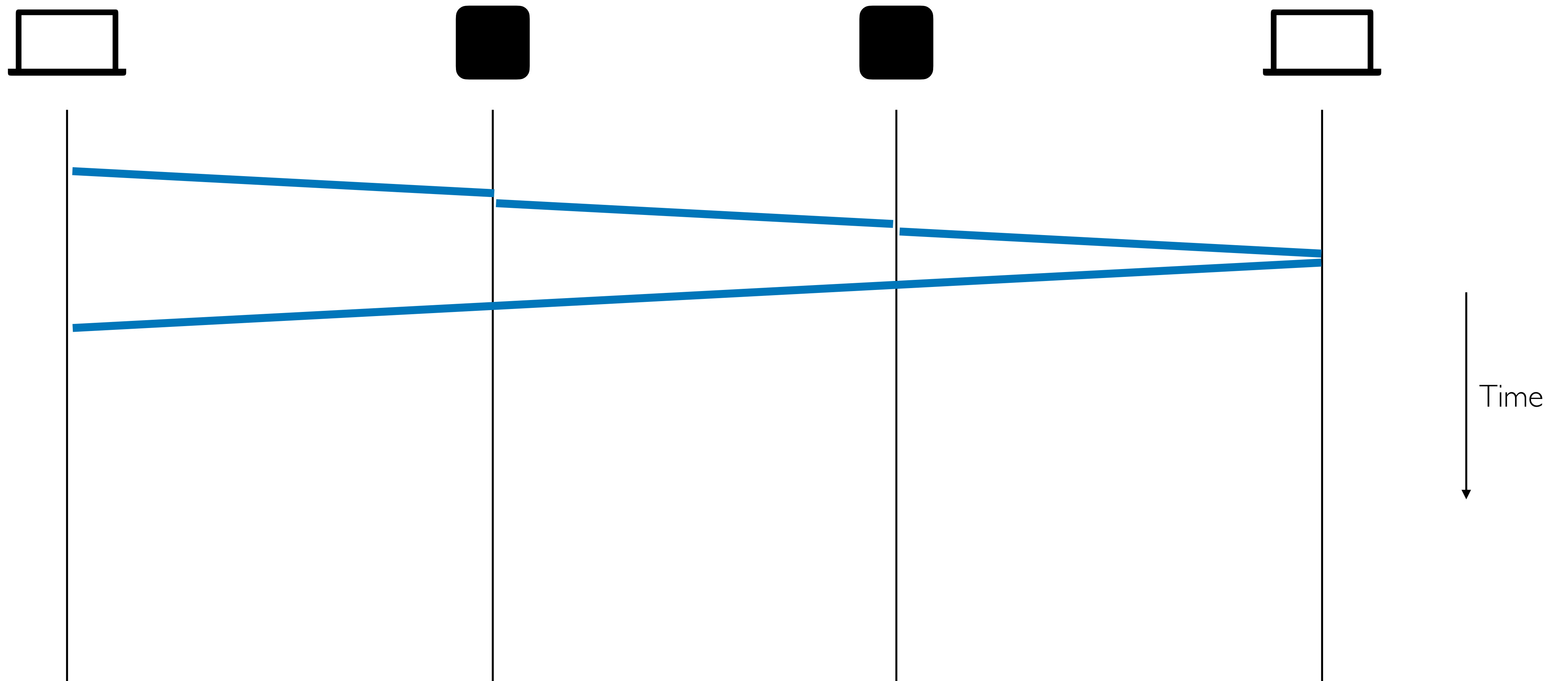
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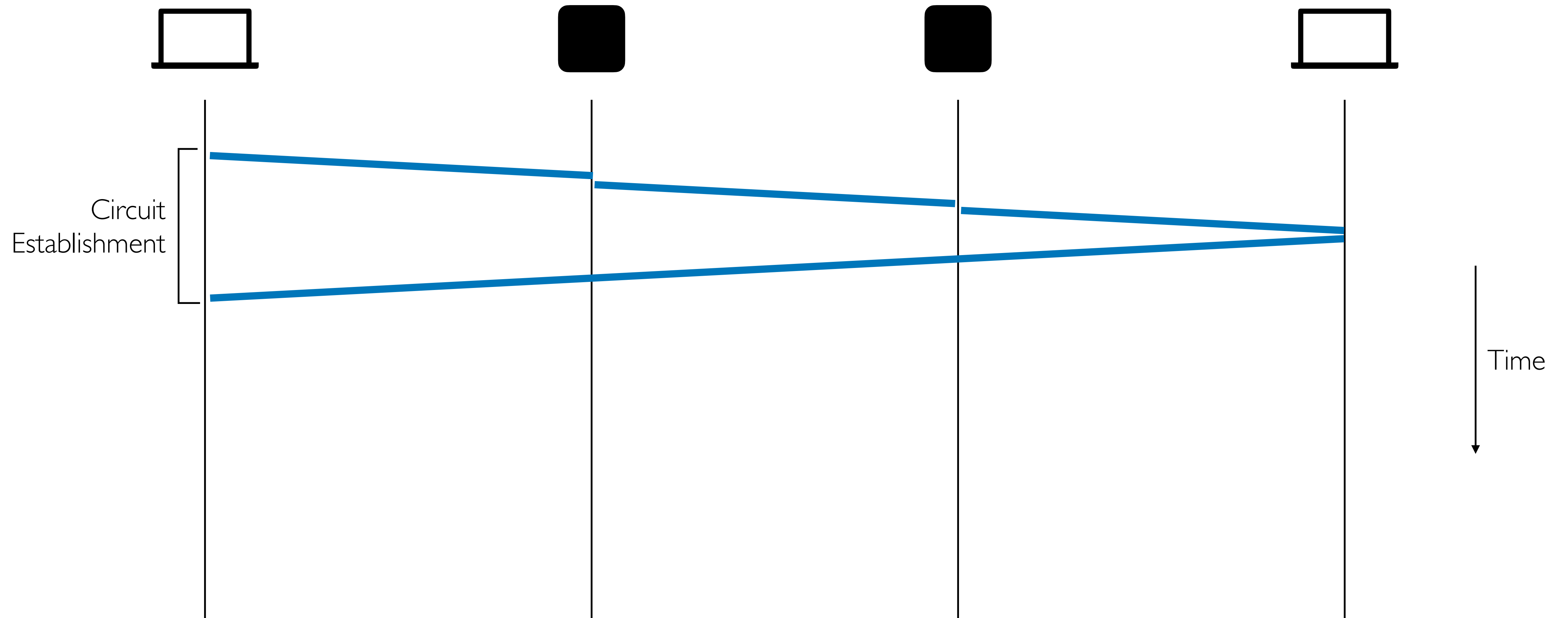
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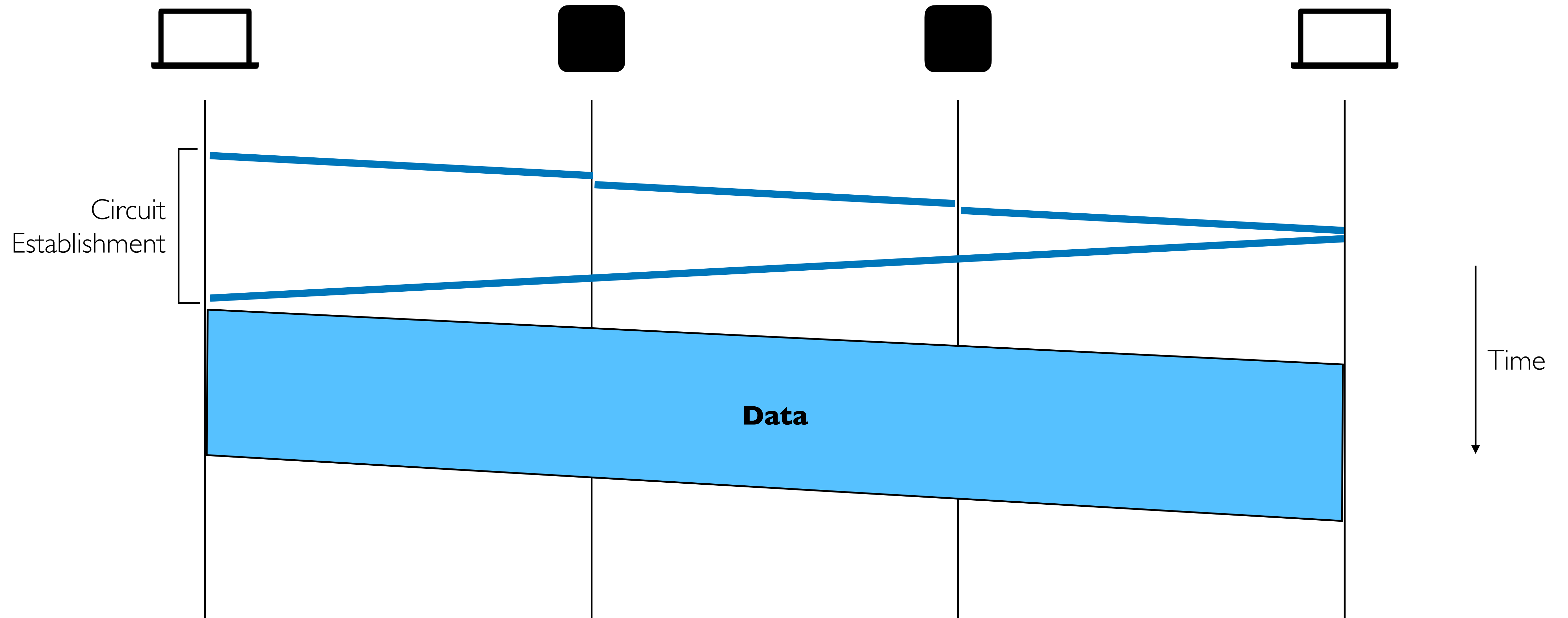
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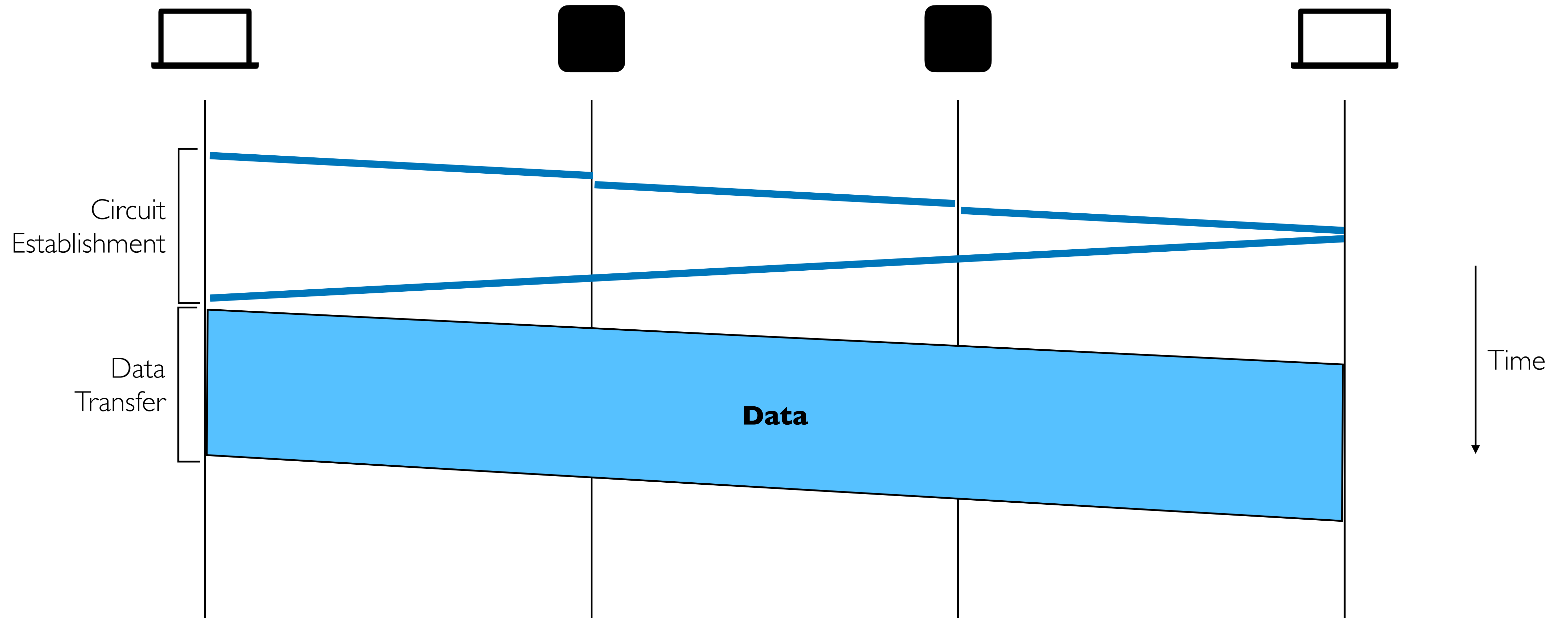
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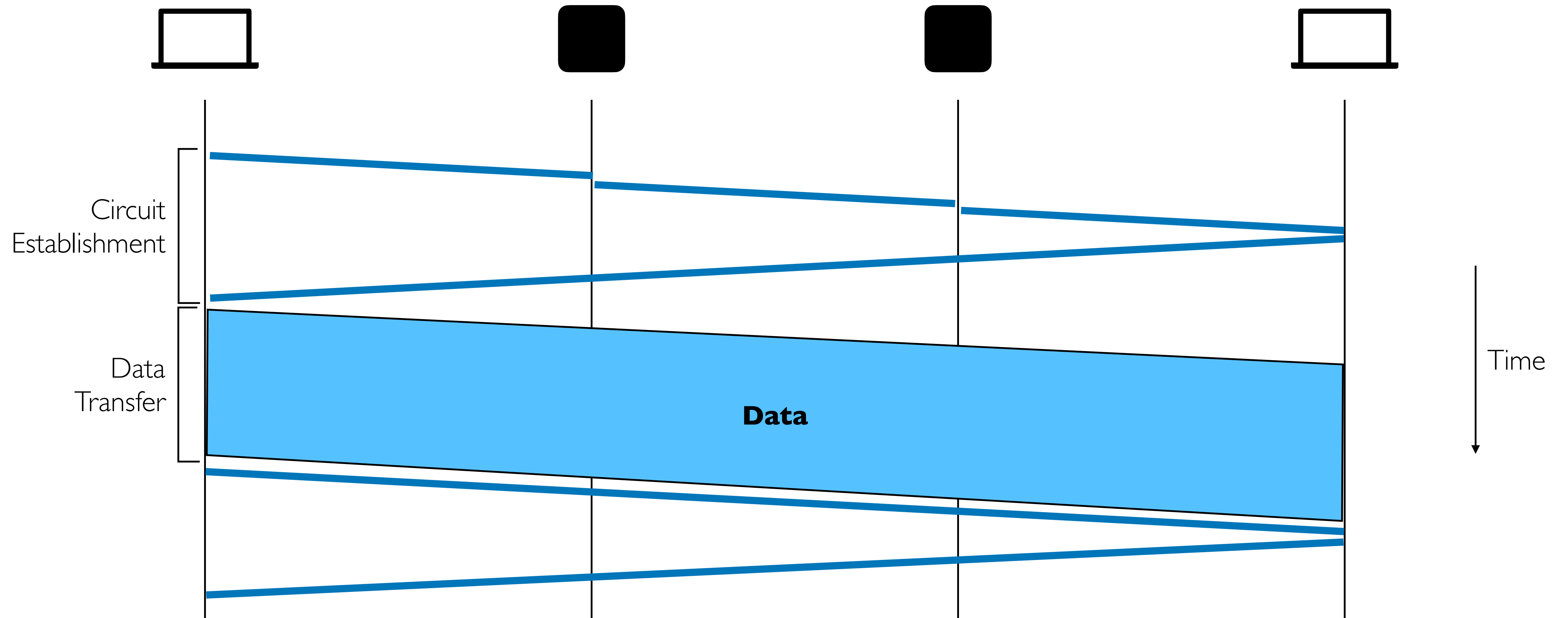
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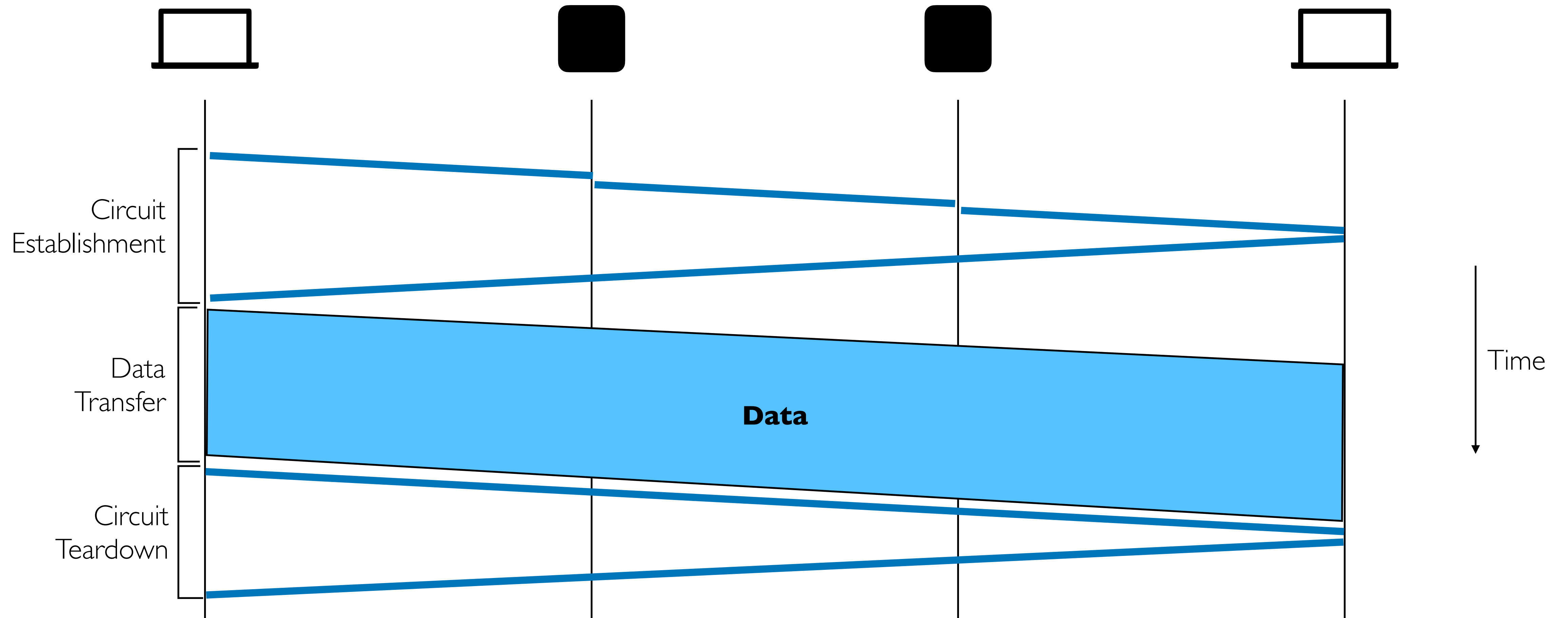
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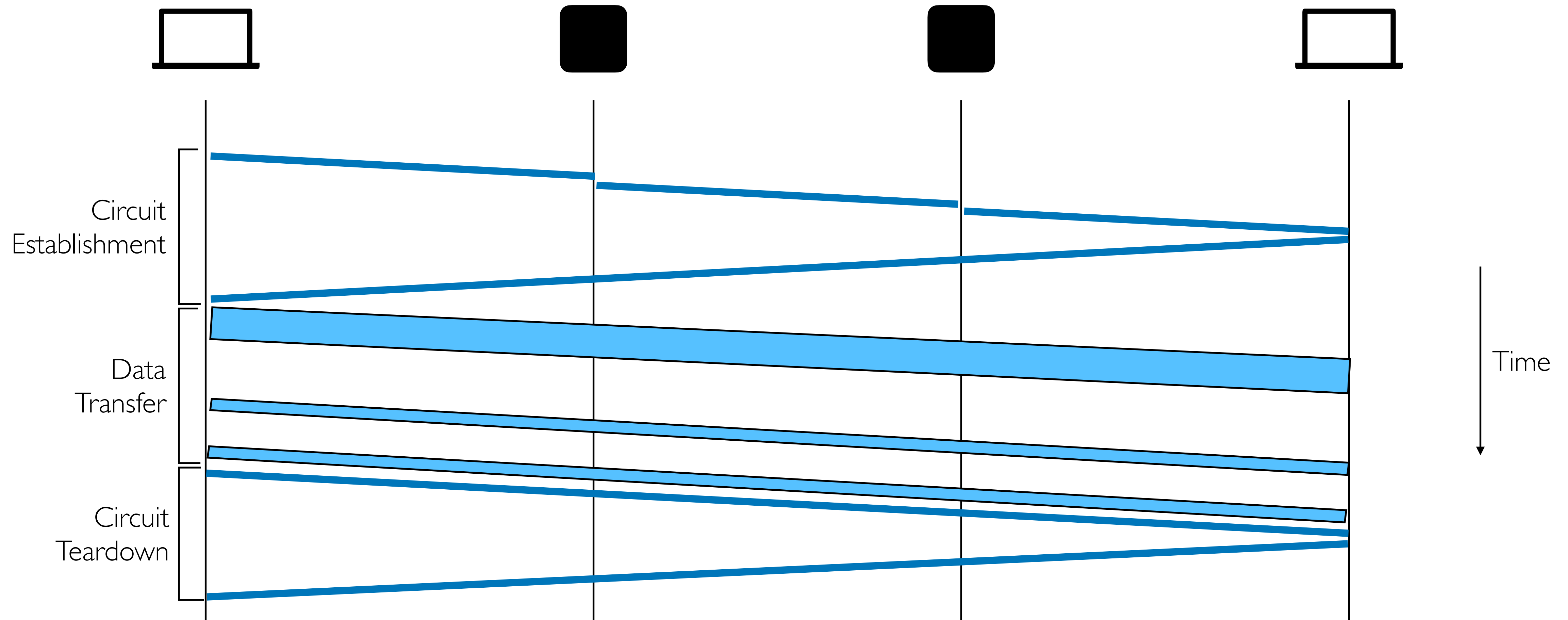
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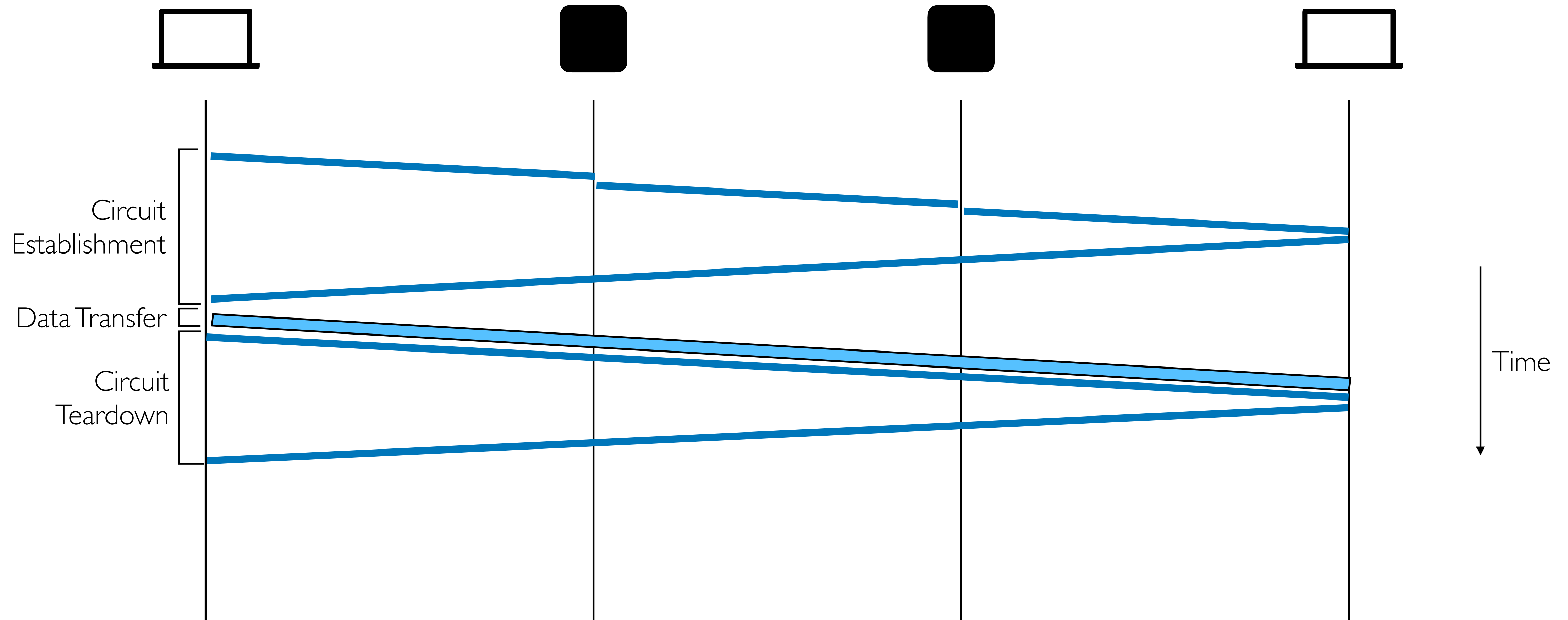
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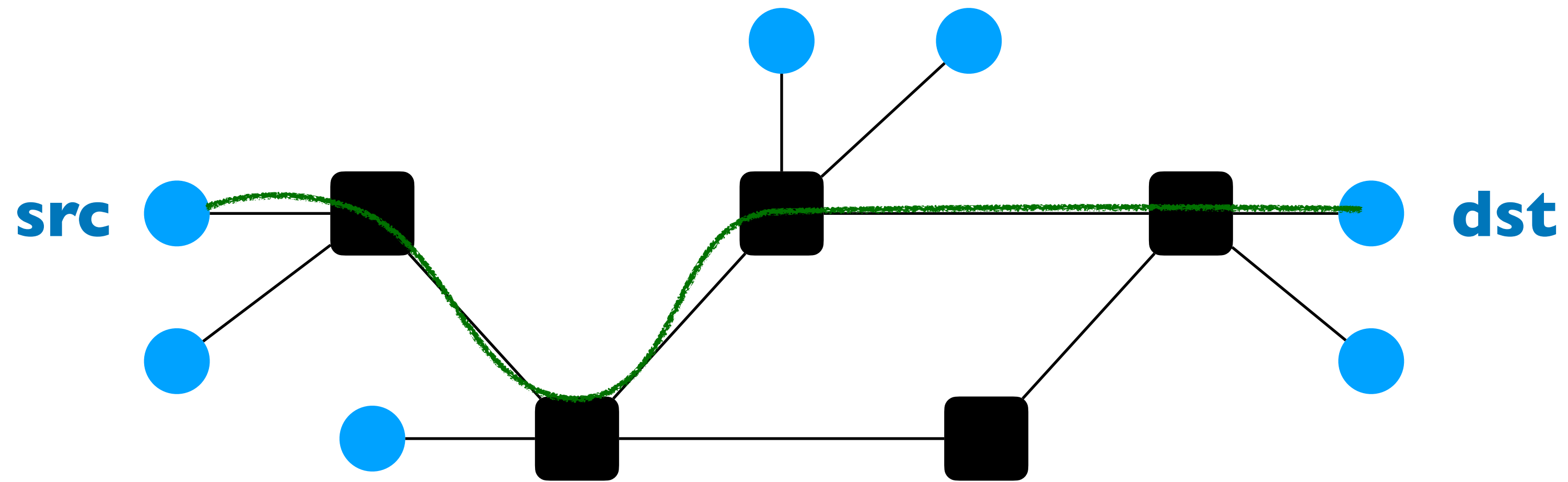
What if there are sporadic periods of data transfer?



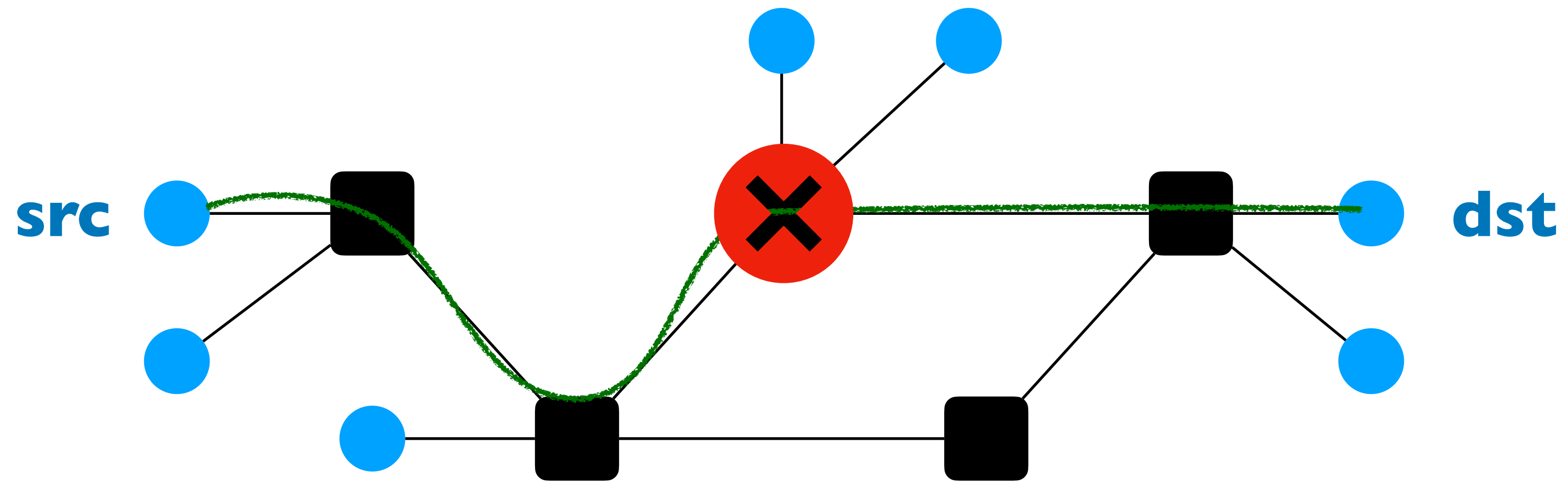
What if there is very little data to transfer?



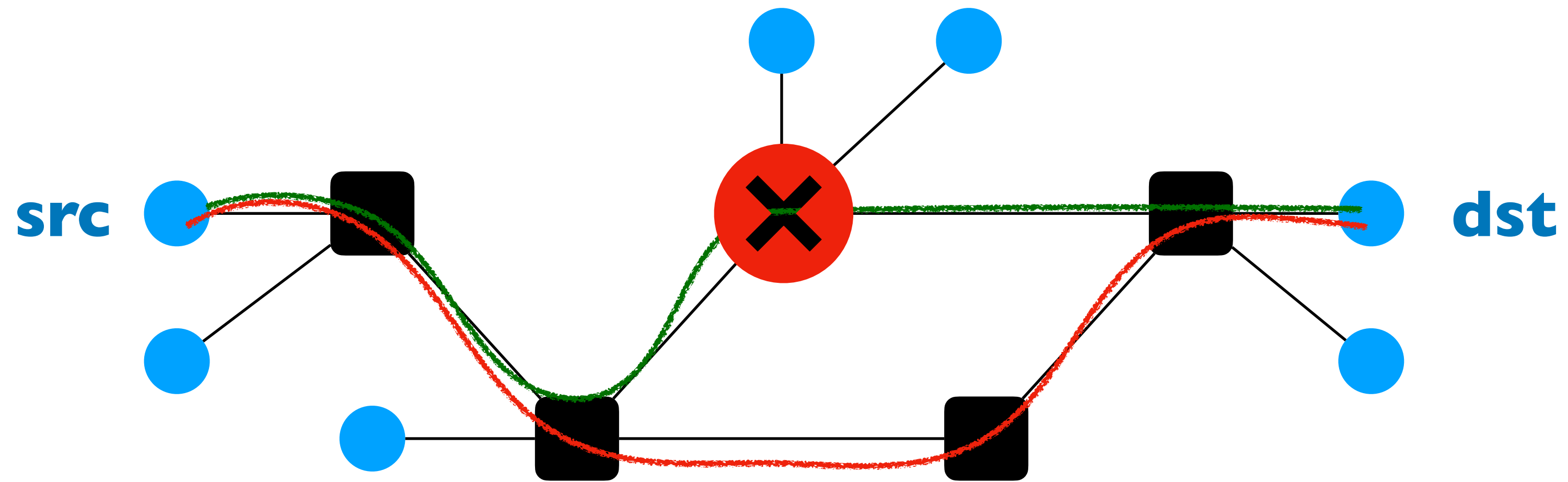
What if switches fail?



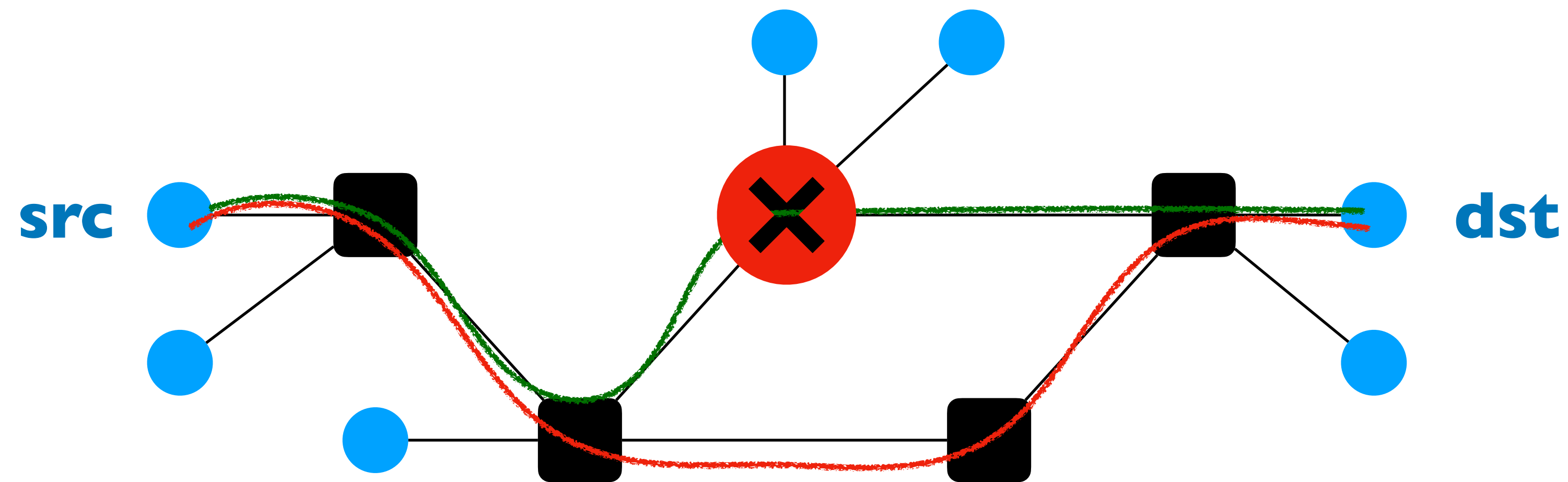
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What if switches fail?



Circuit switching doesn't ***route around trouble***

Circuit Switching

Circuit Switching

- **What's good?**
 - *Predictable performance*
 - *Simple/fast switching (once circuit is established)*

Circuit Switching

- **What's good?**
 - *Predictable performance*
 - *Simple/fast switching (once circuit is established)*
- **What's not-so-good?**
 - *Complexity of circuit setup/teardown*
 - *Inefficient when traffic is bursty*
 - *Circuit setup adds delay*
 - *Switch fails -> its circuits fail*

Questions?

How else can we share?

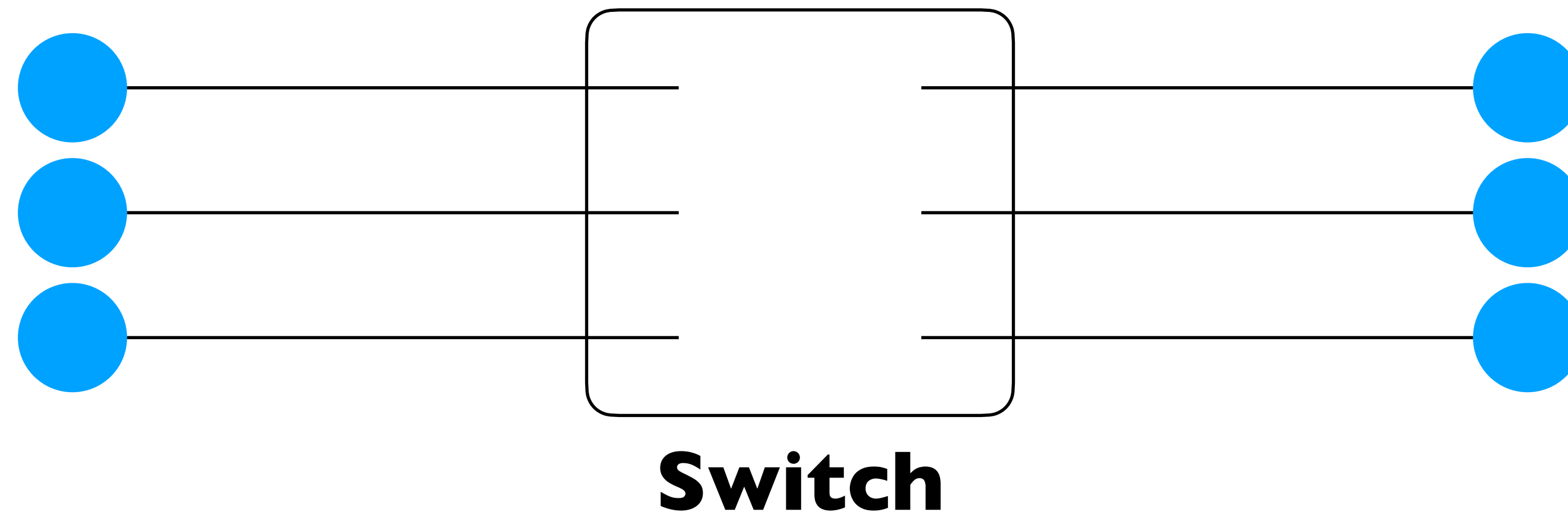
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- **Second approach:** *On demand* (aka “best-effort”)

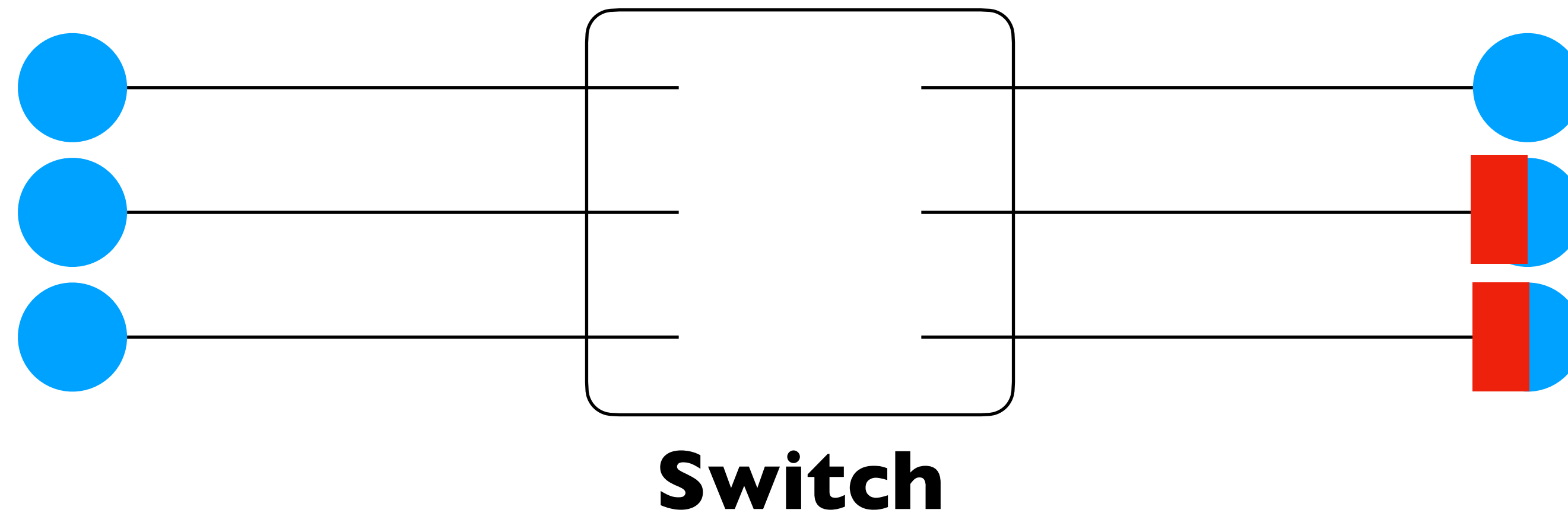
How else can we share?

- **Second approach:** *On demand* (aka “best-effort”)
- **Mechanism:**
 - Break down data into packets
 - Send packets when you have them
 - Hope for the best...

Best-effort: Packet Switching

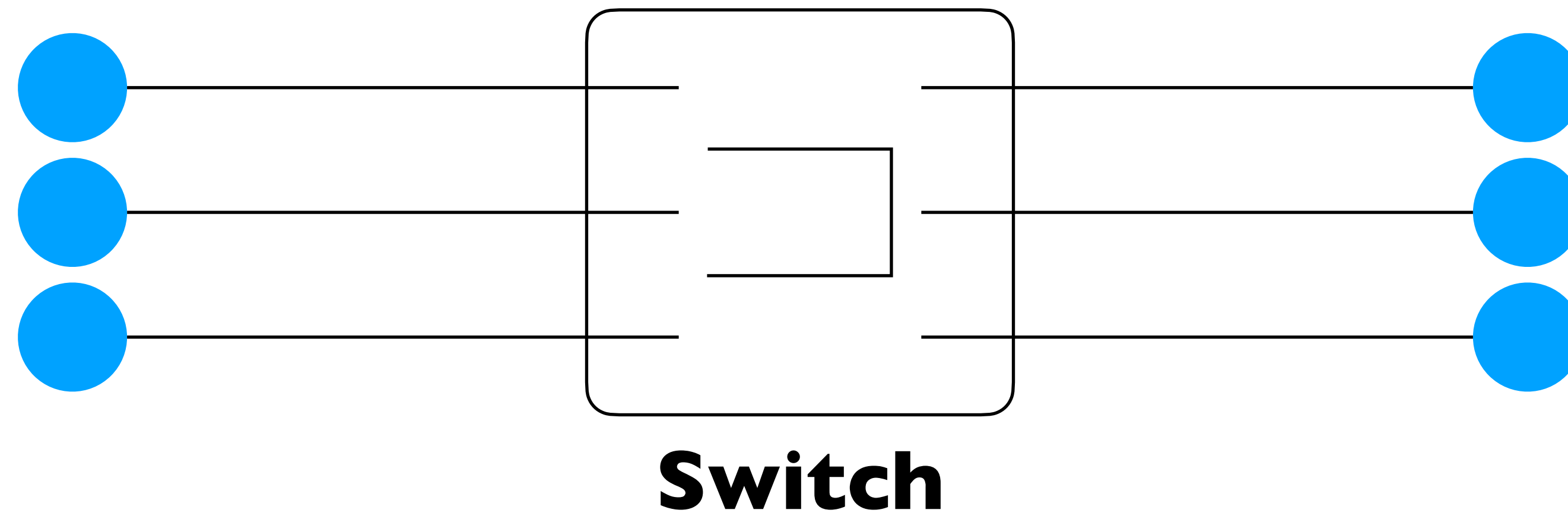


Best-effort: Packet Switching



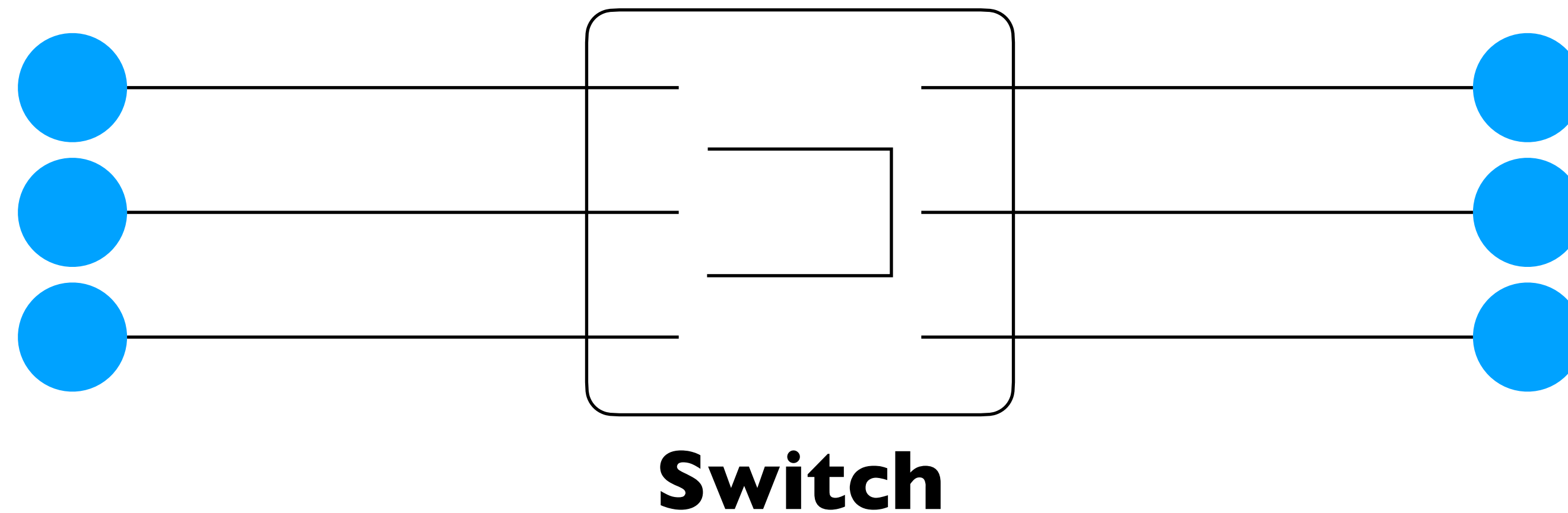
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Each packet treated independently

Best-effort: Packet Switching



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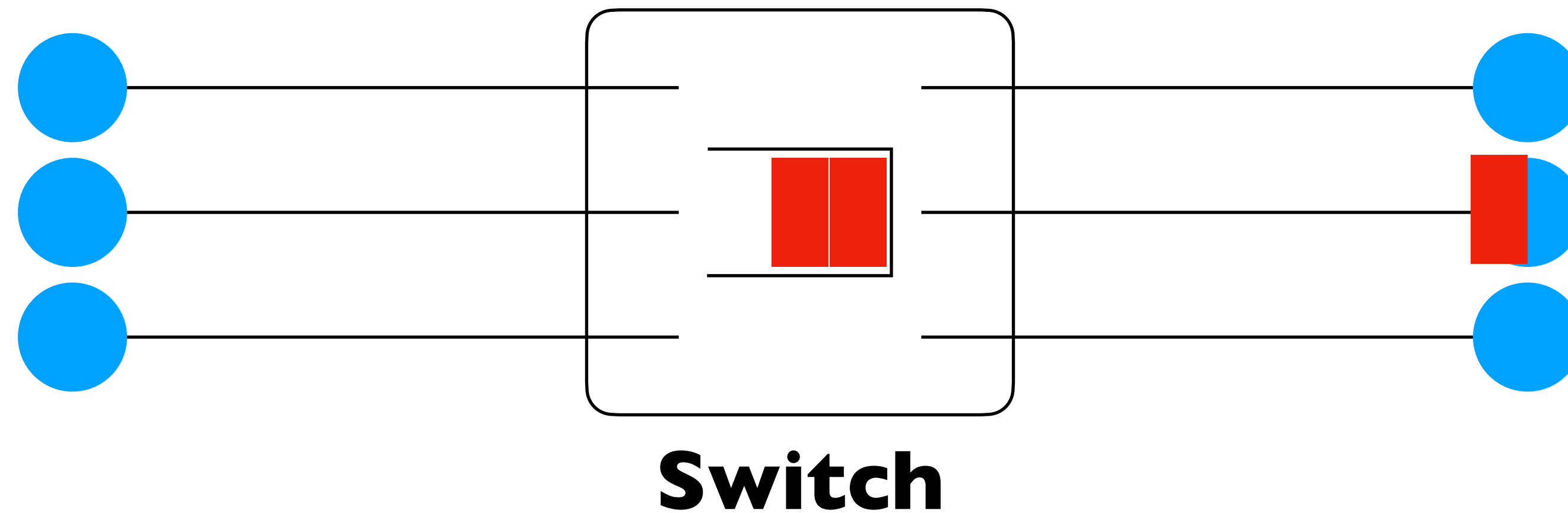


Each packet contains destination

Each packet treated independently

With buffers to absorb transient overloads

Best-effort: Packet Switching



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 - *Simpler to implement*
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Packet Switching

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- *Efficient use of network resources*
- *Simpler to implement*
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- **What's not-so-good?**

- *Unpredictable performance*
- *Requires buffer management and congestion control*

Circuits vs Packets

Circuits vs Packets

- **When are circuits better?**
 - *When you need dedicated performance (reserved bandwidth)*
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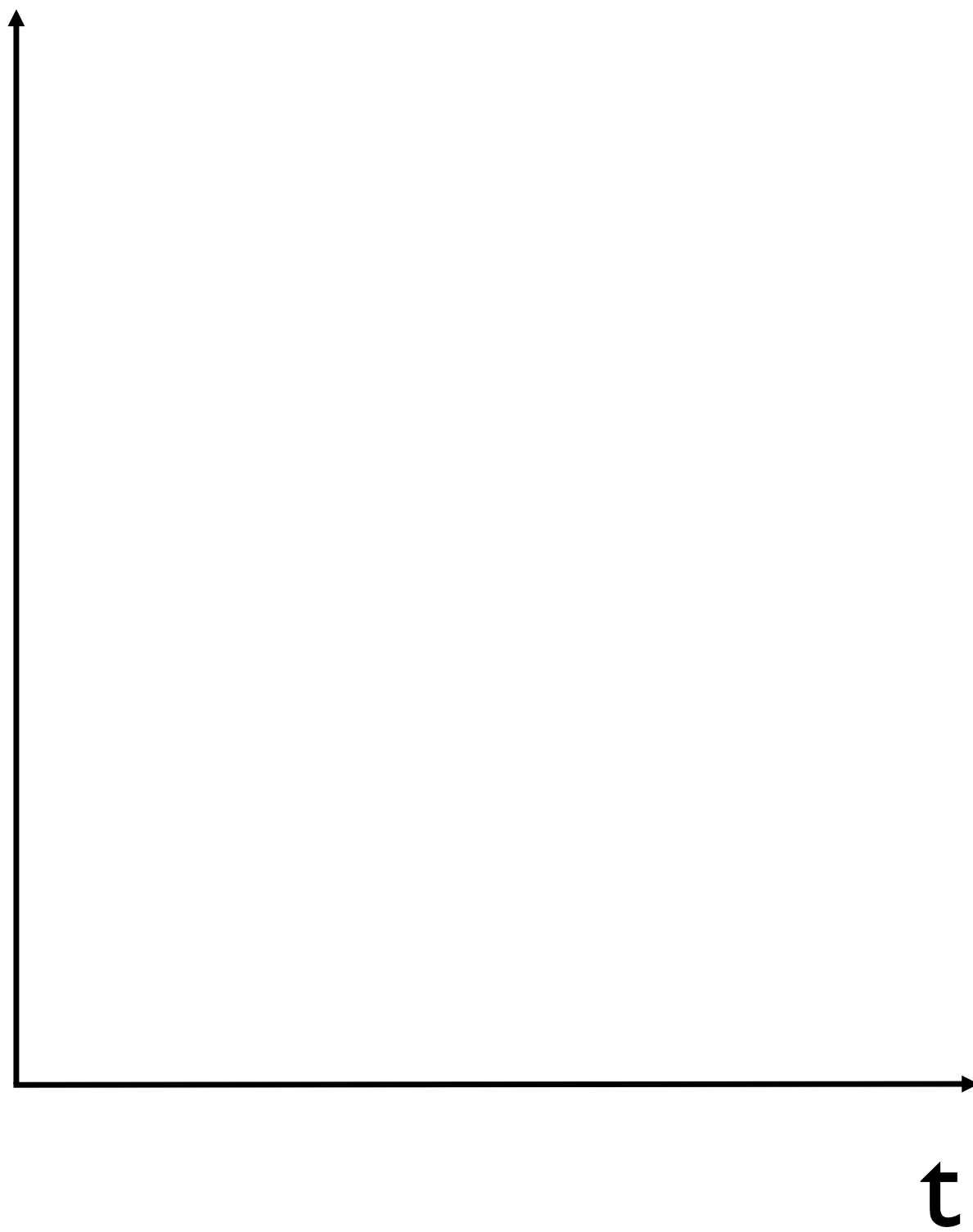
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Statistical Multiplexing: Intuition

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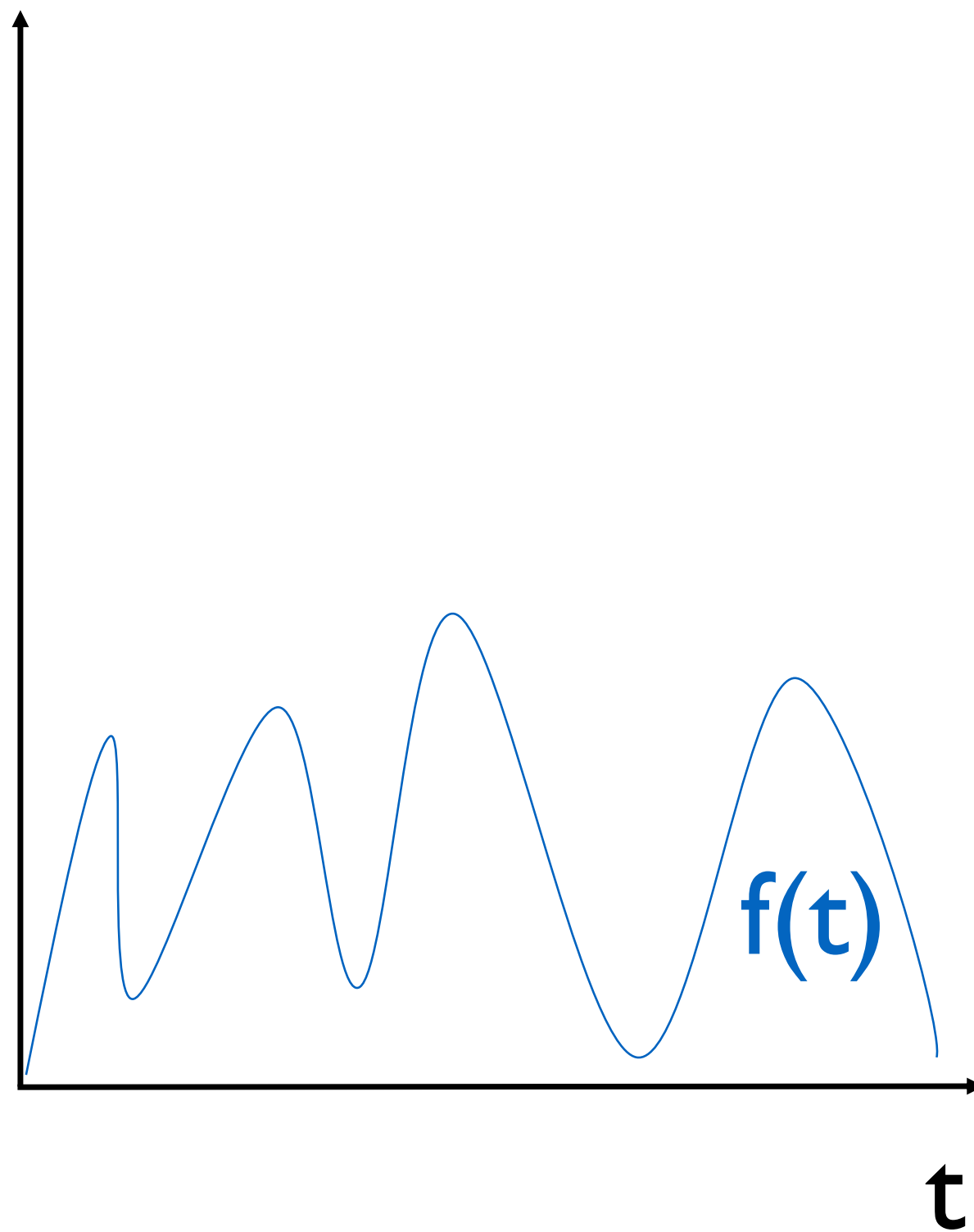
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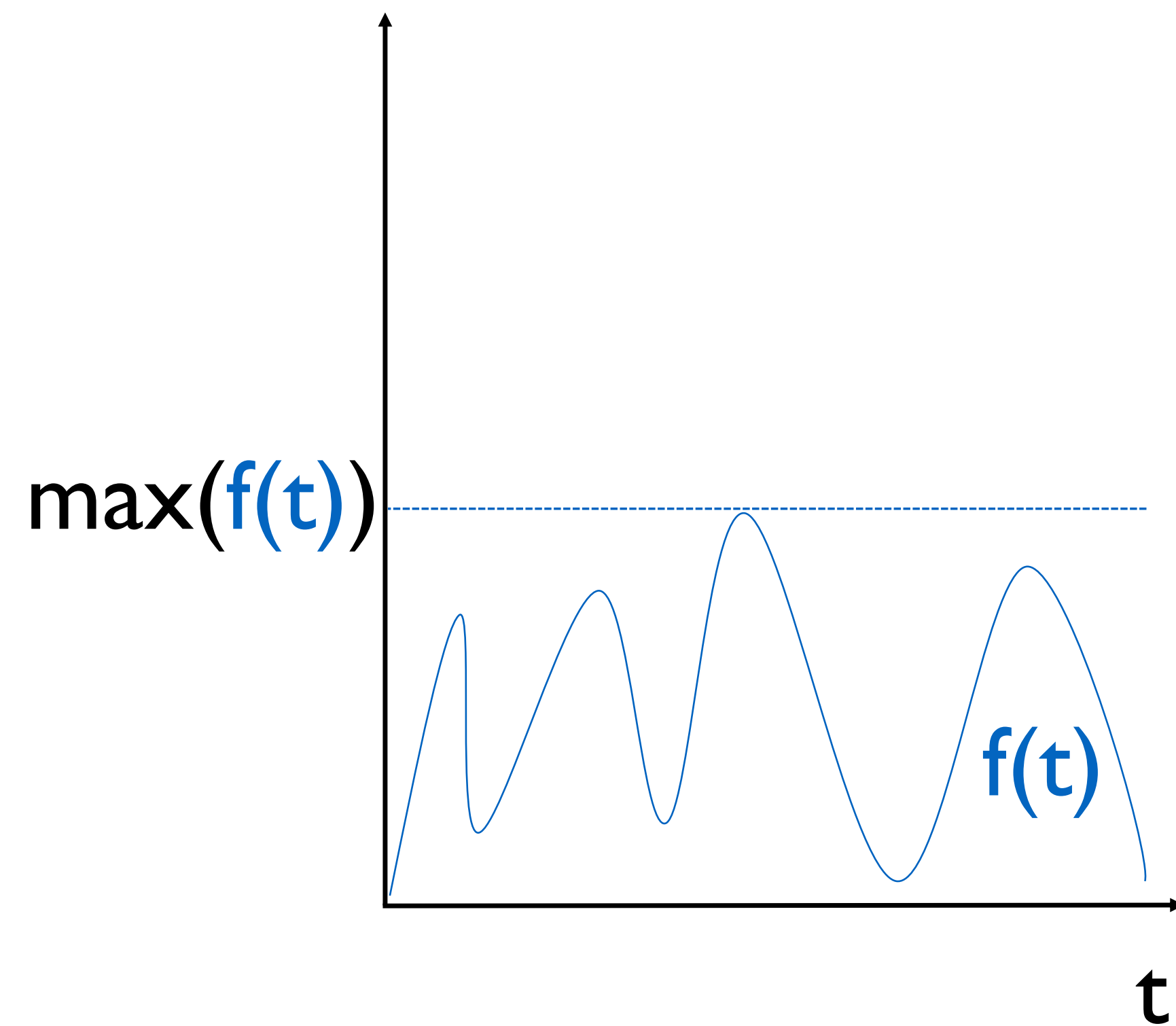
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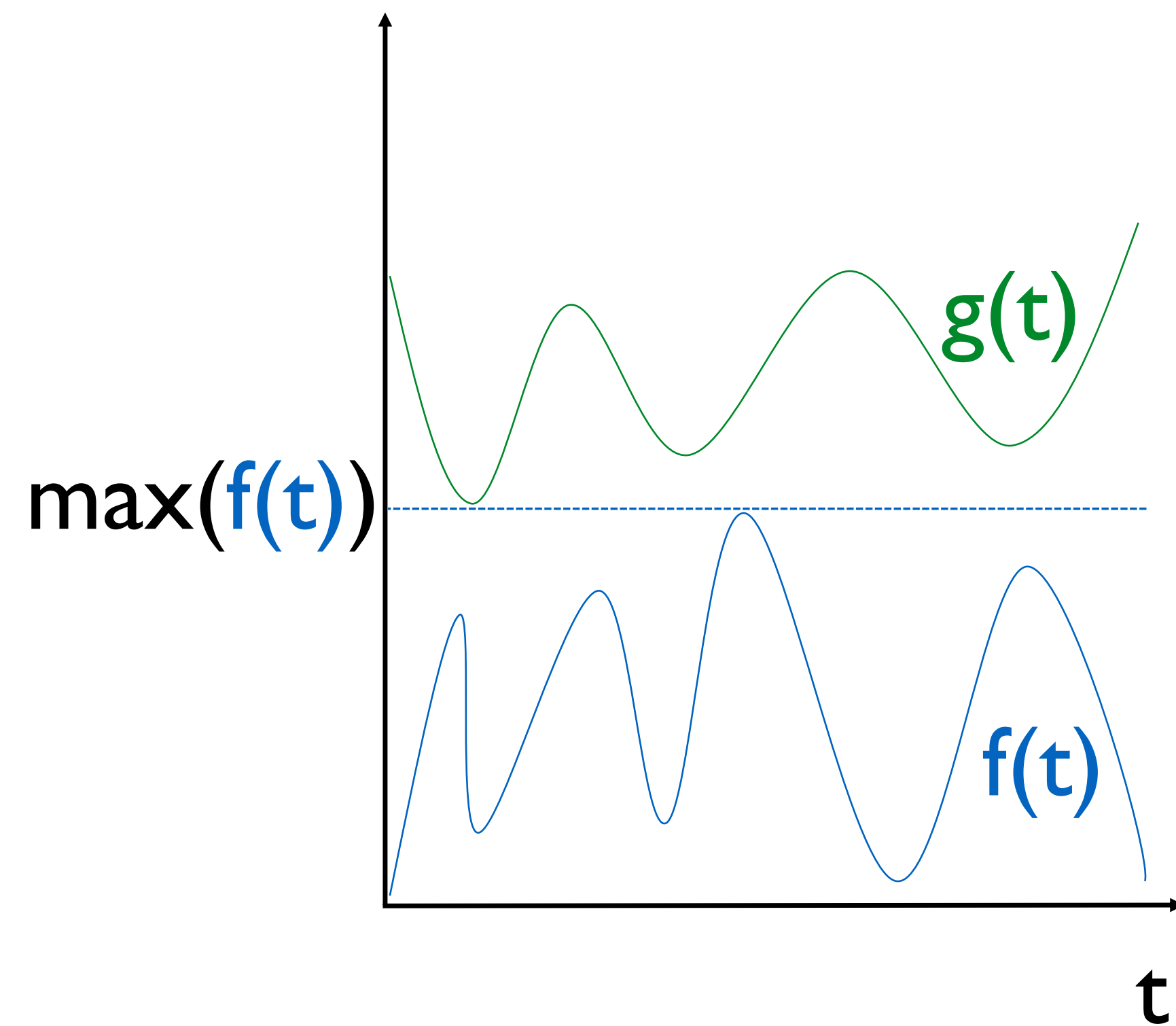
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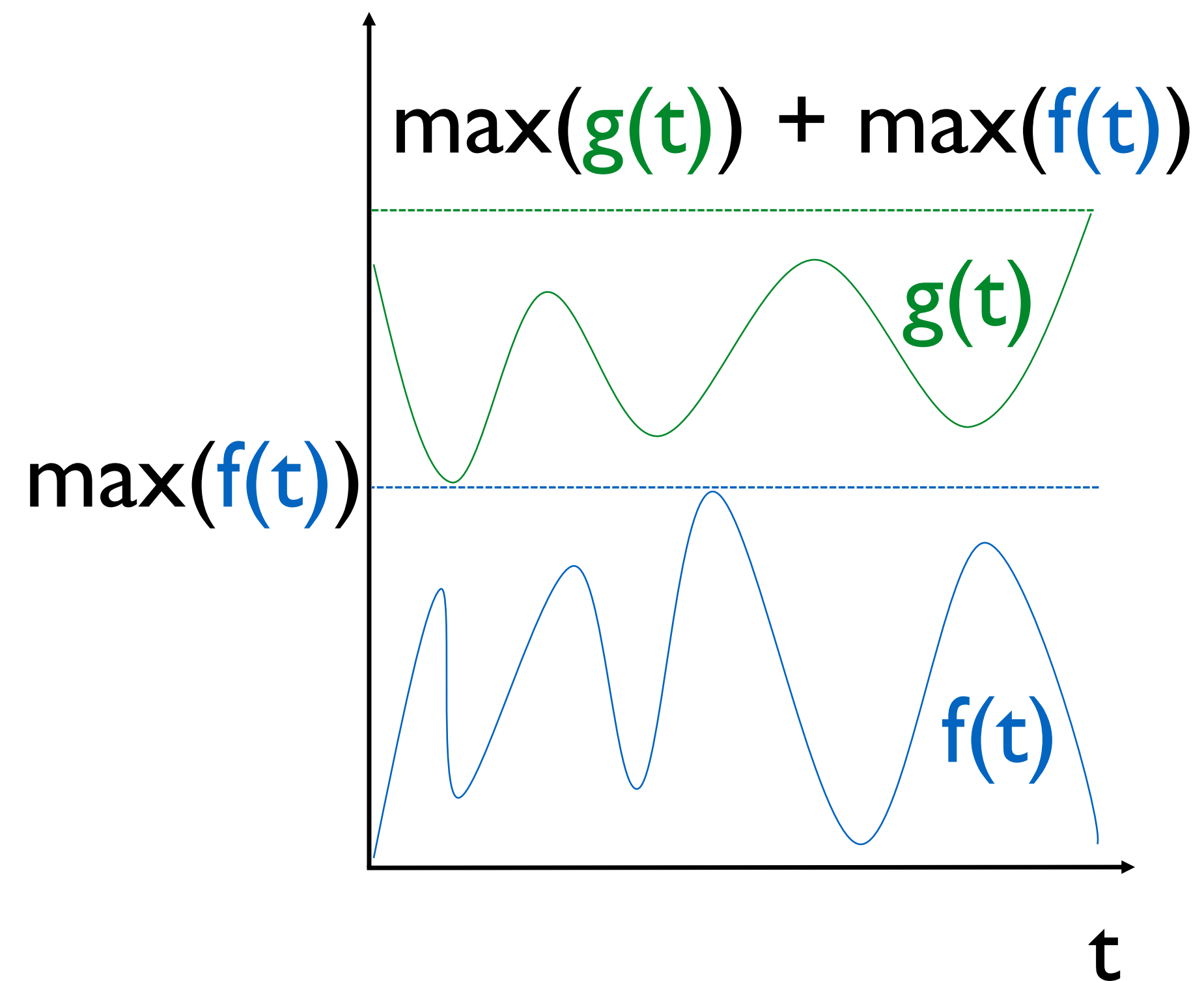
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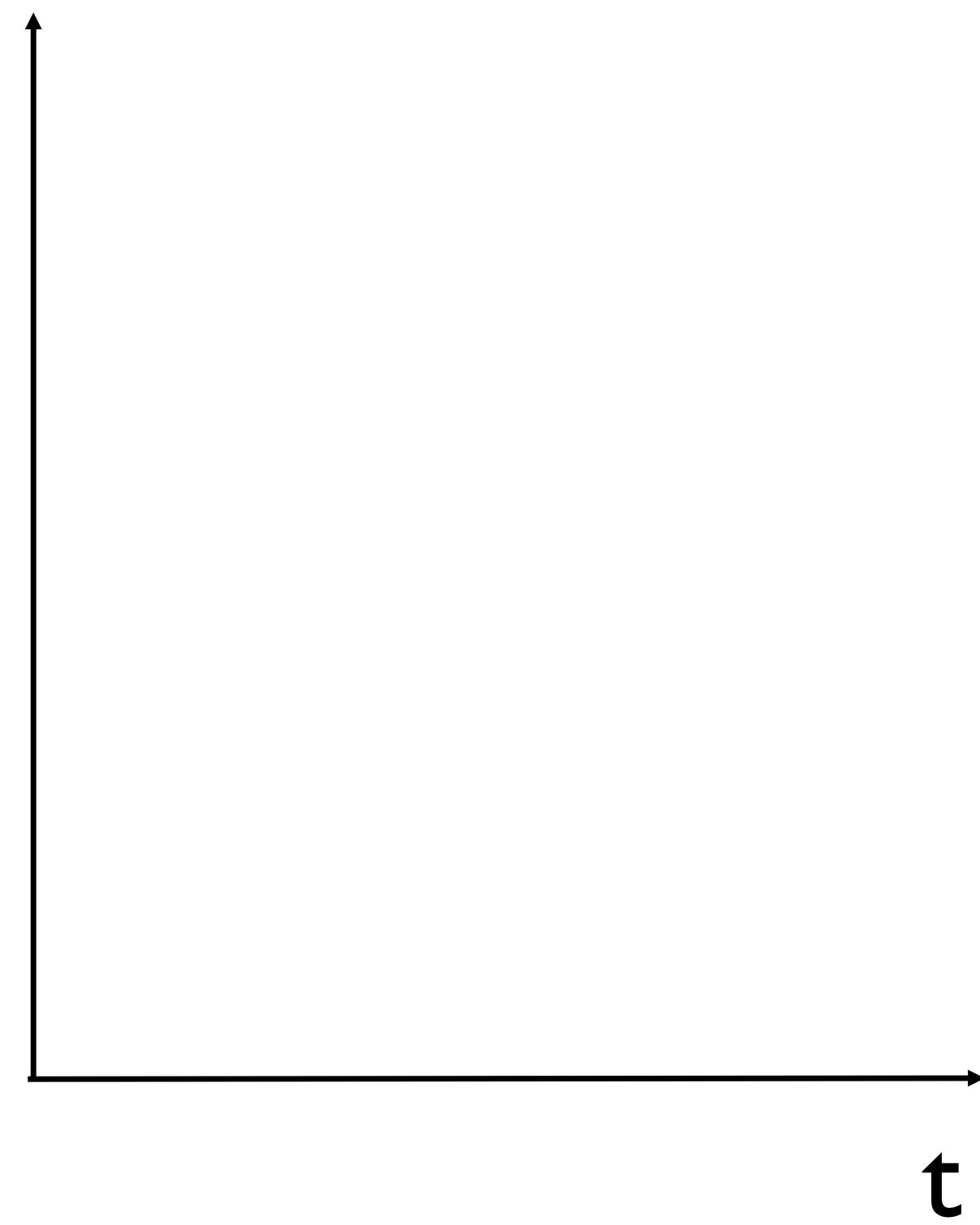
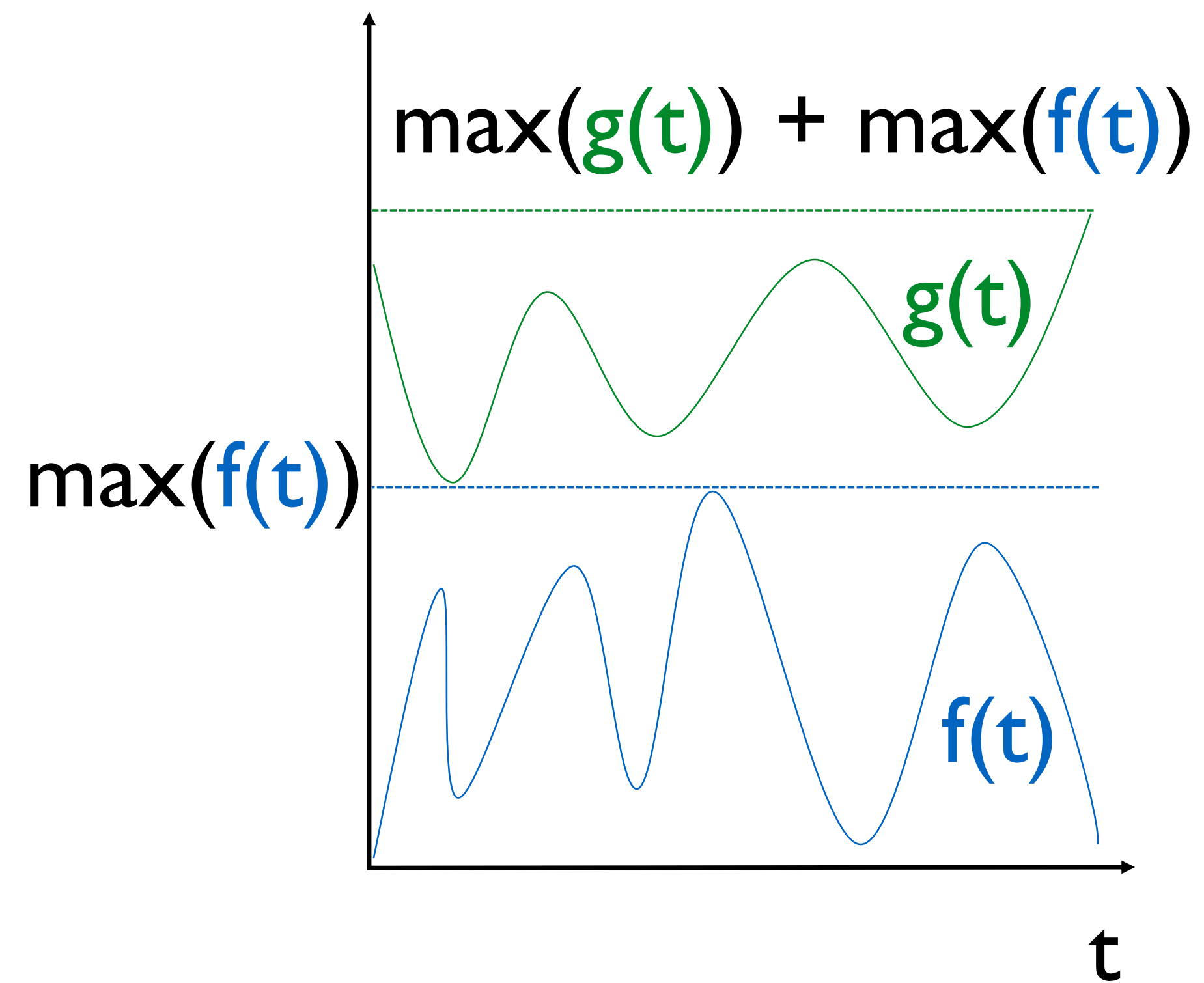
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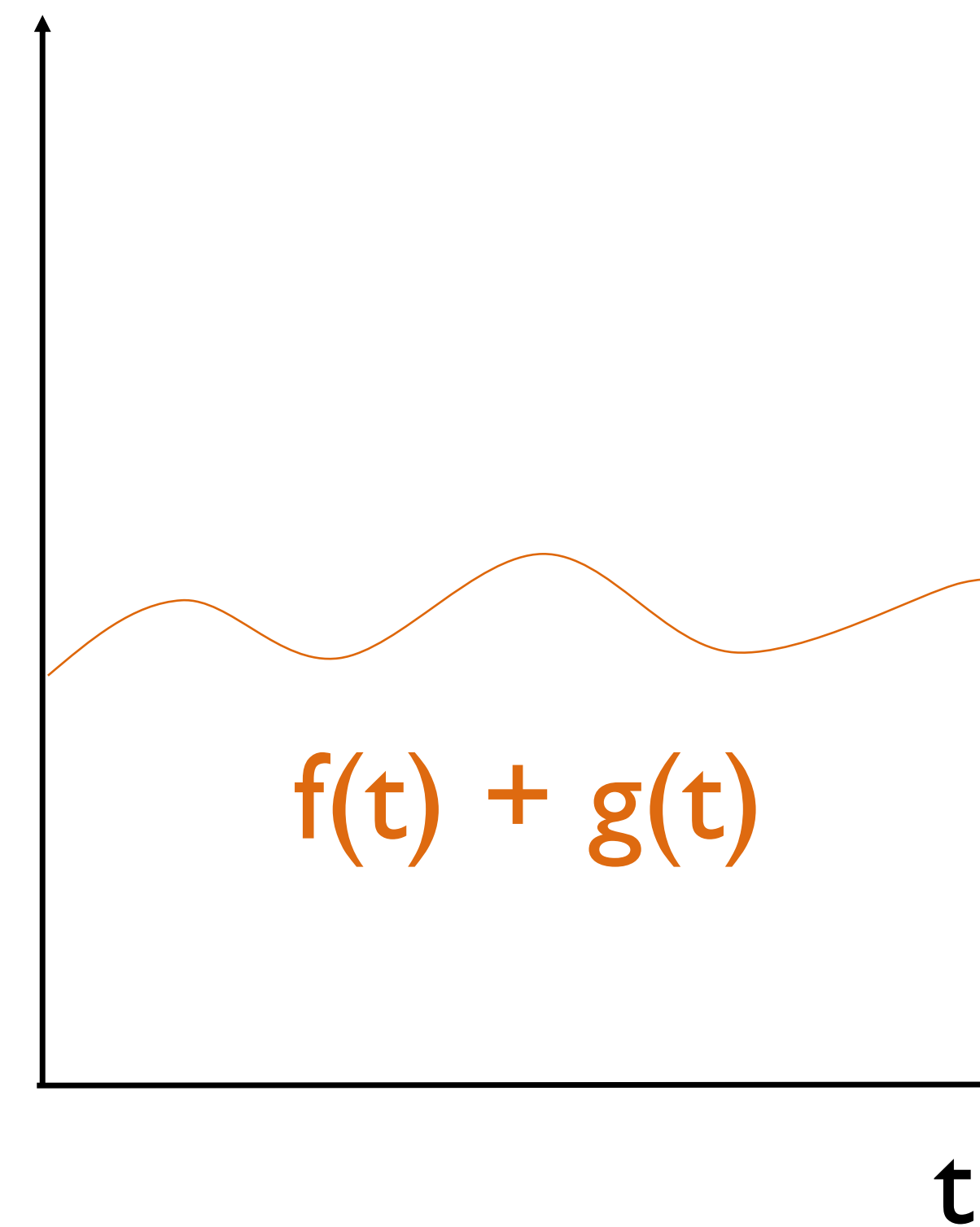
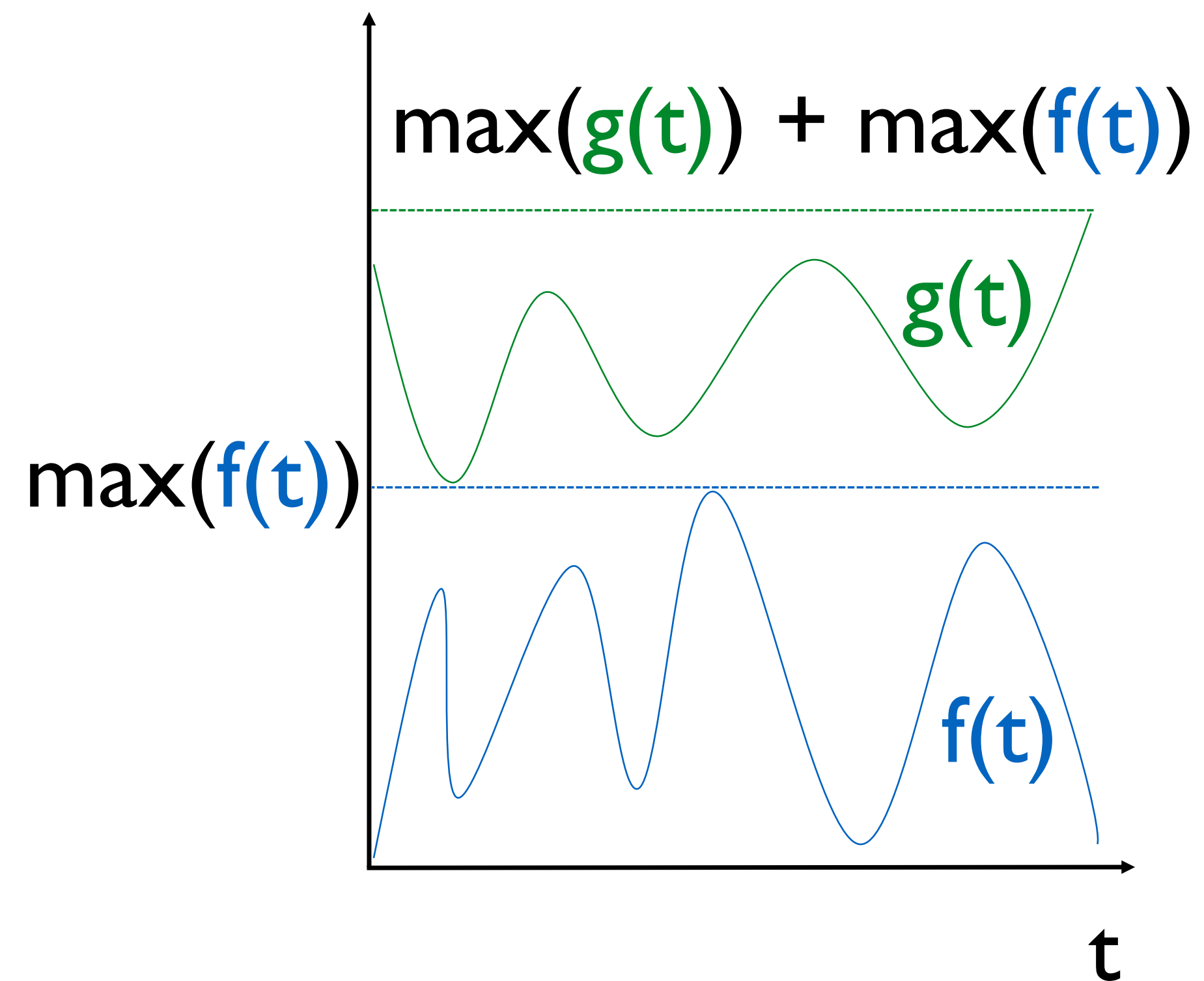
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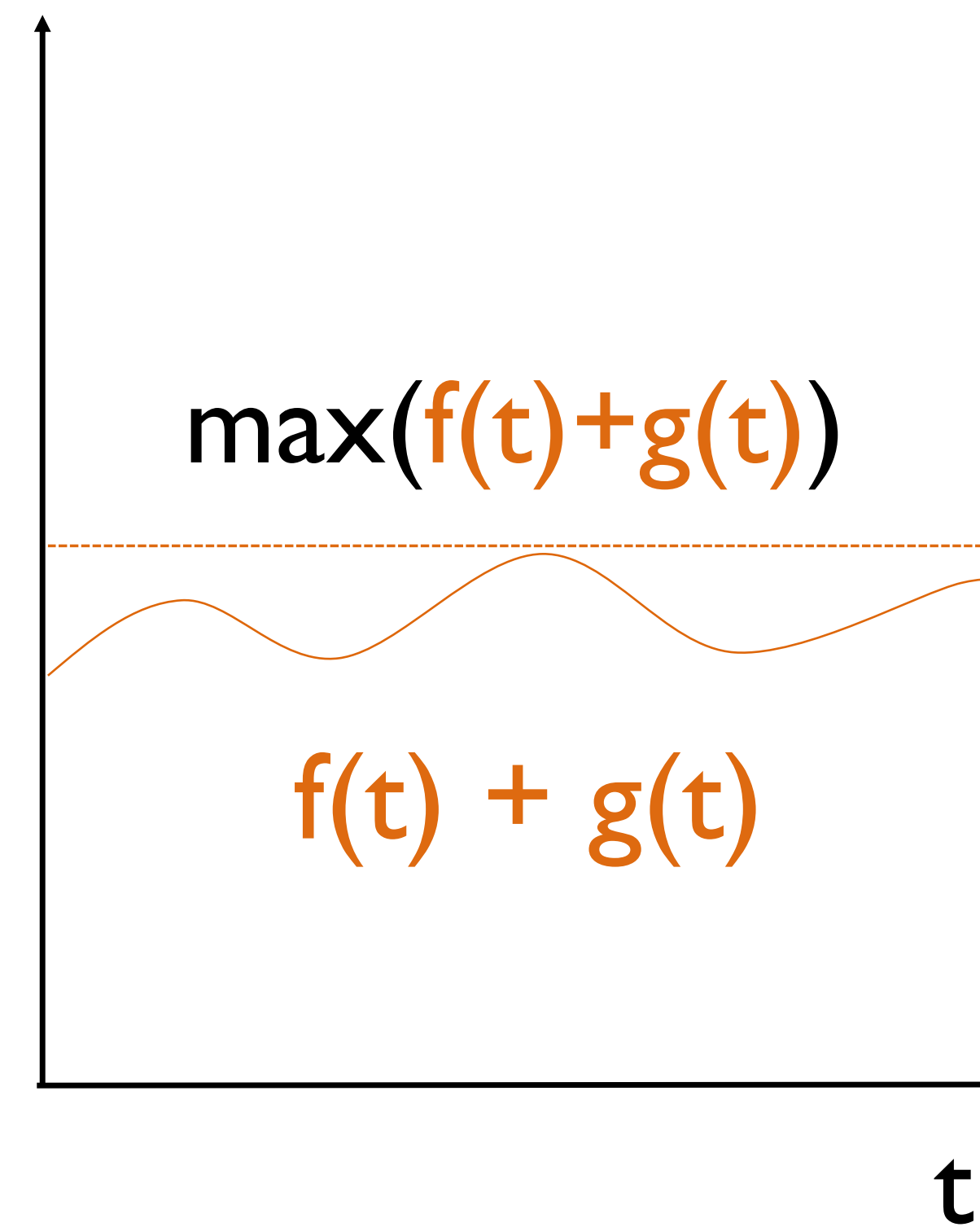
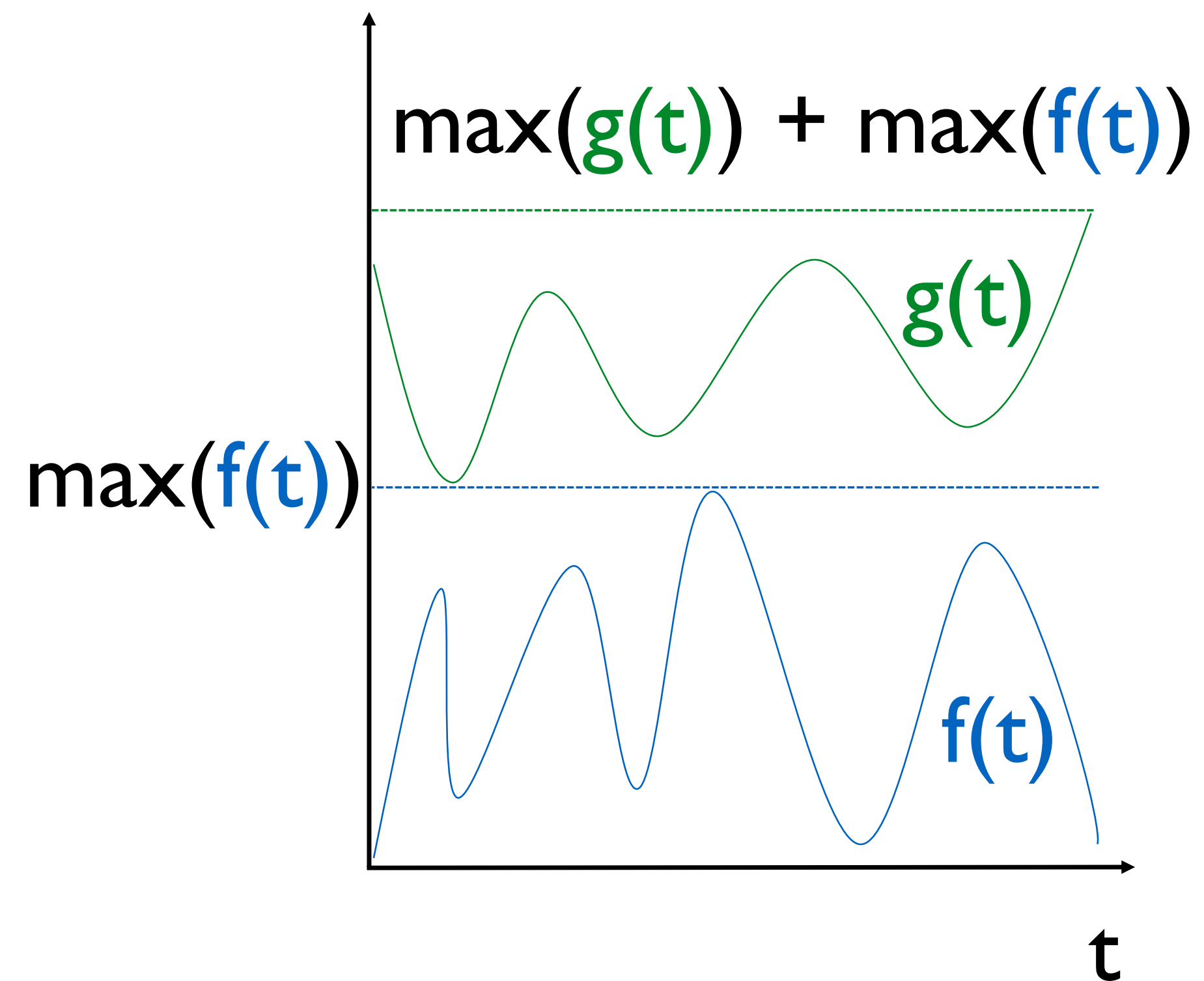
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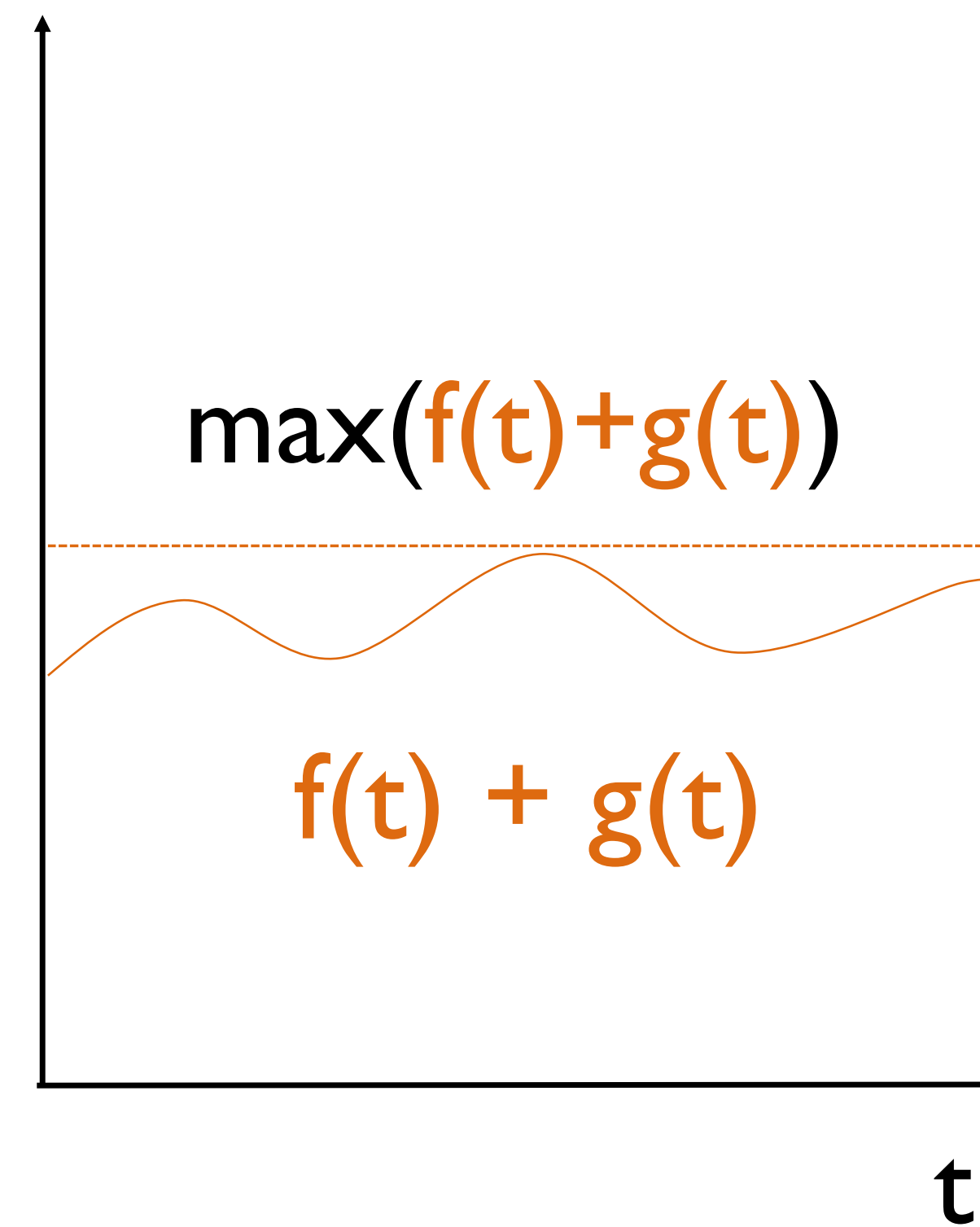
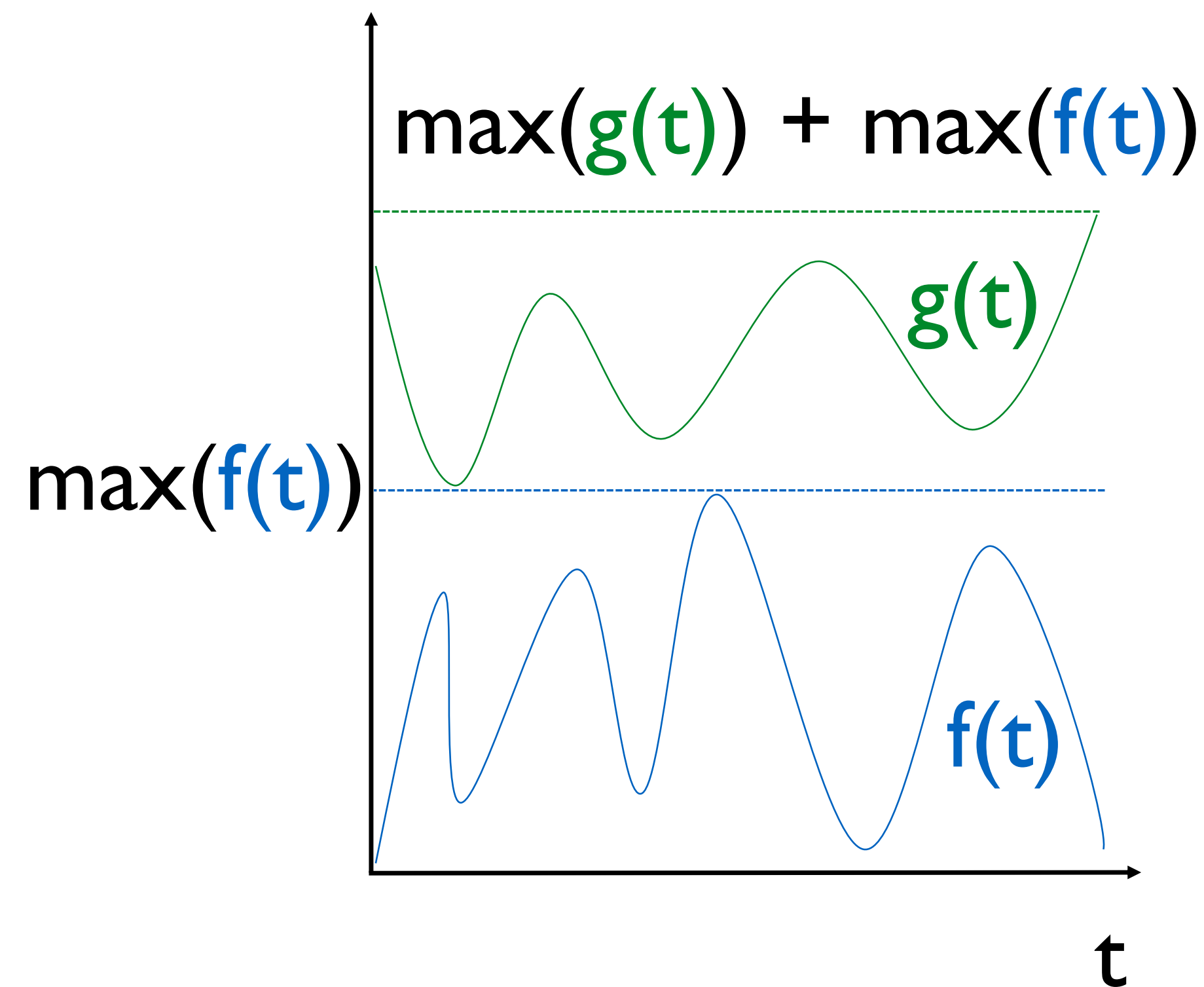
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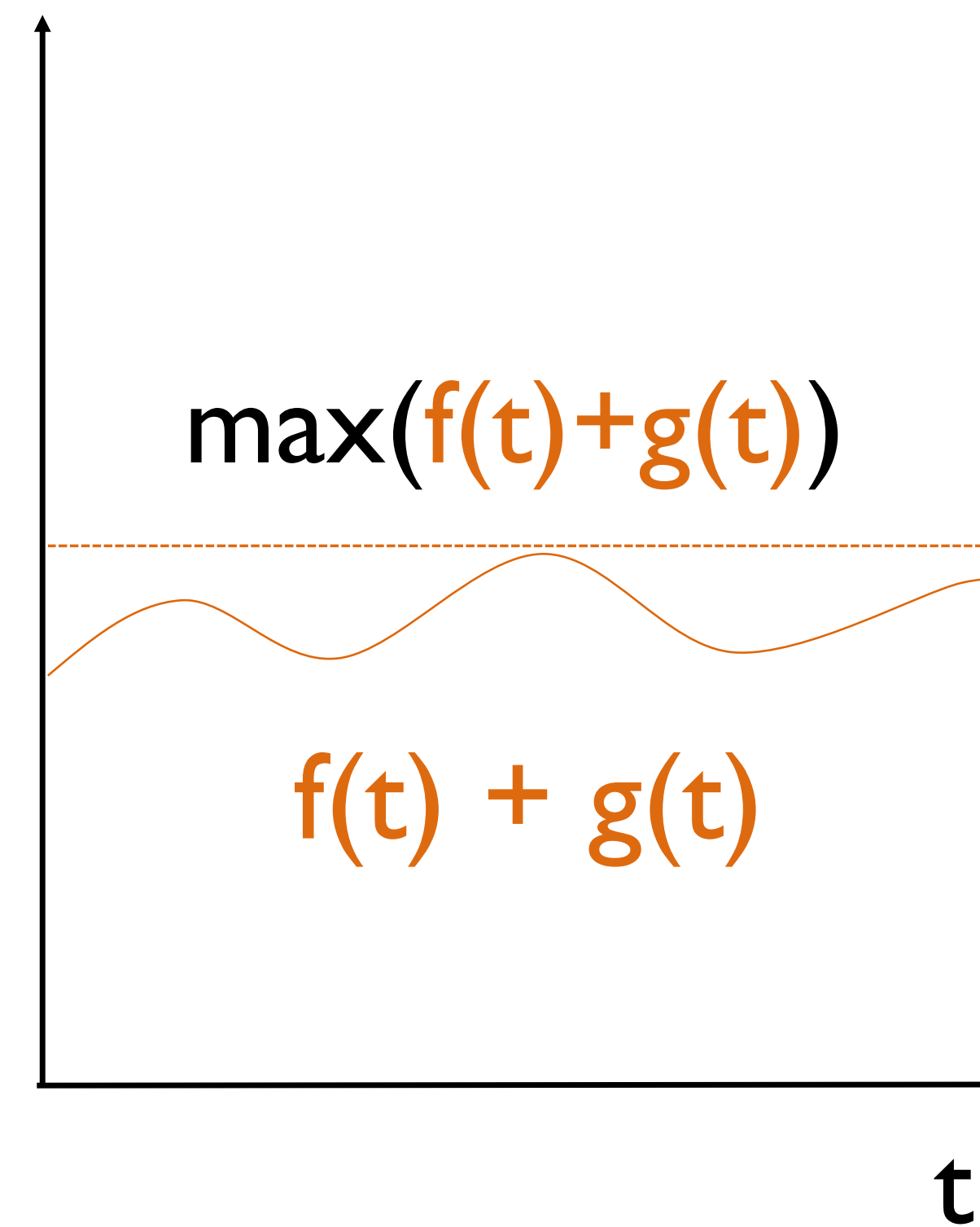
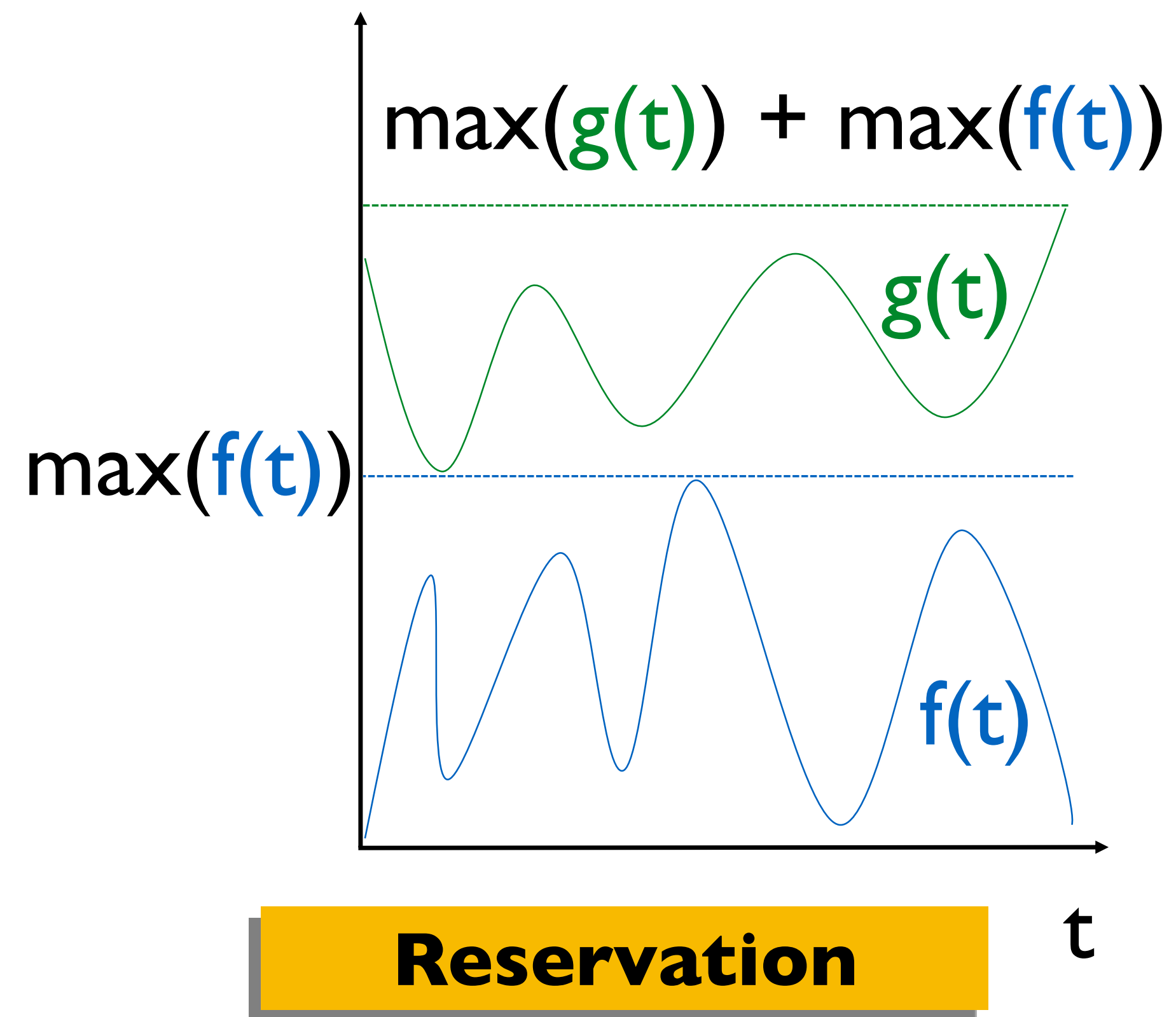
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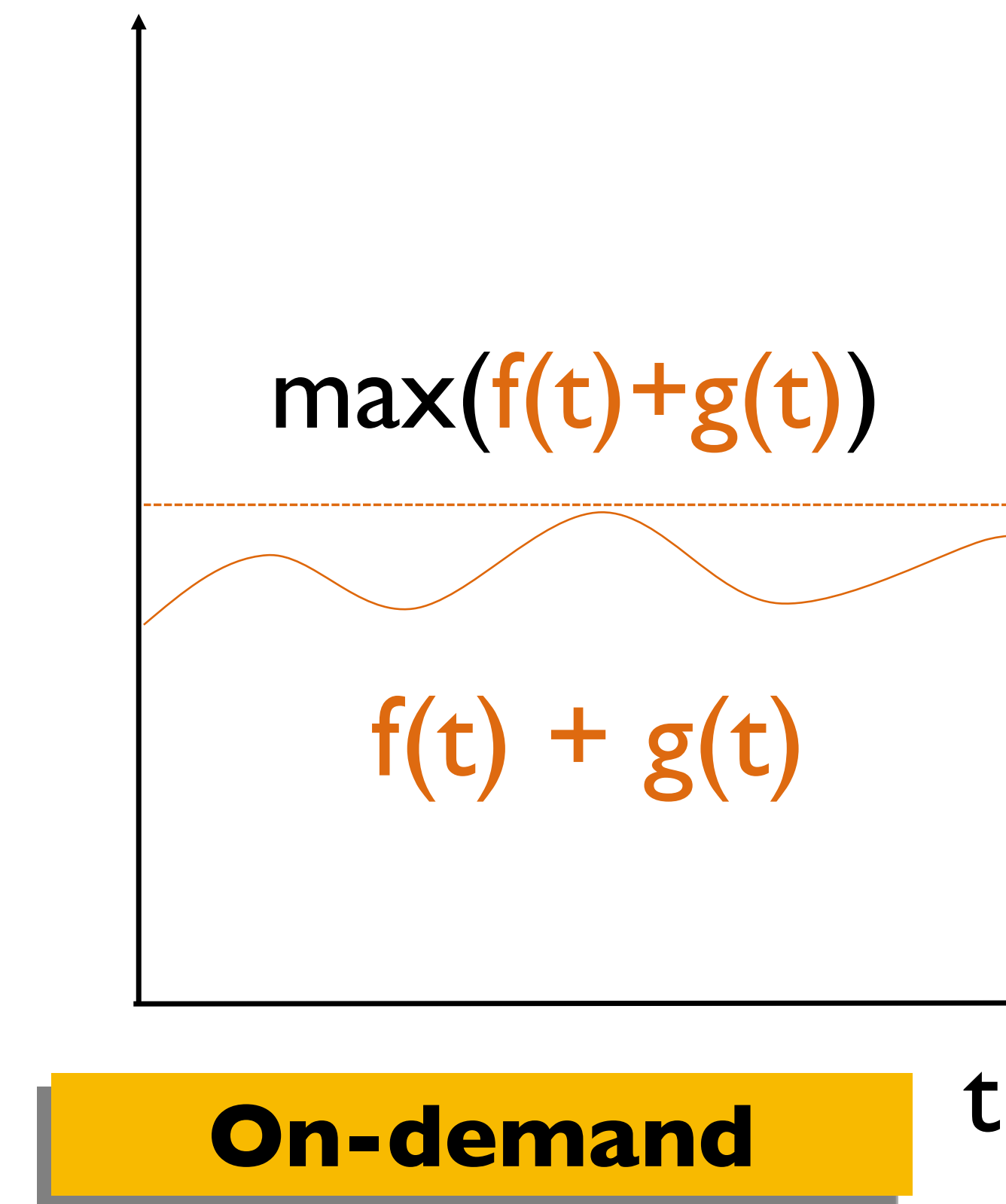
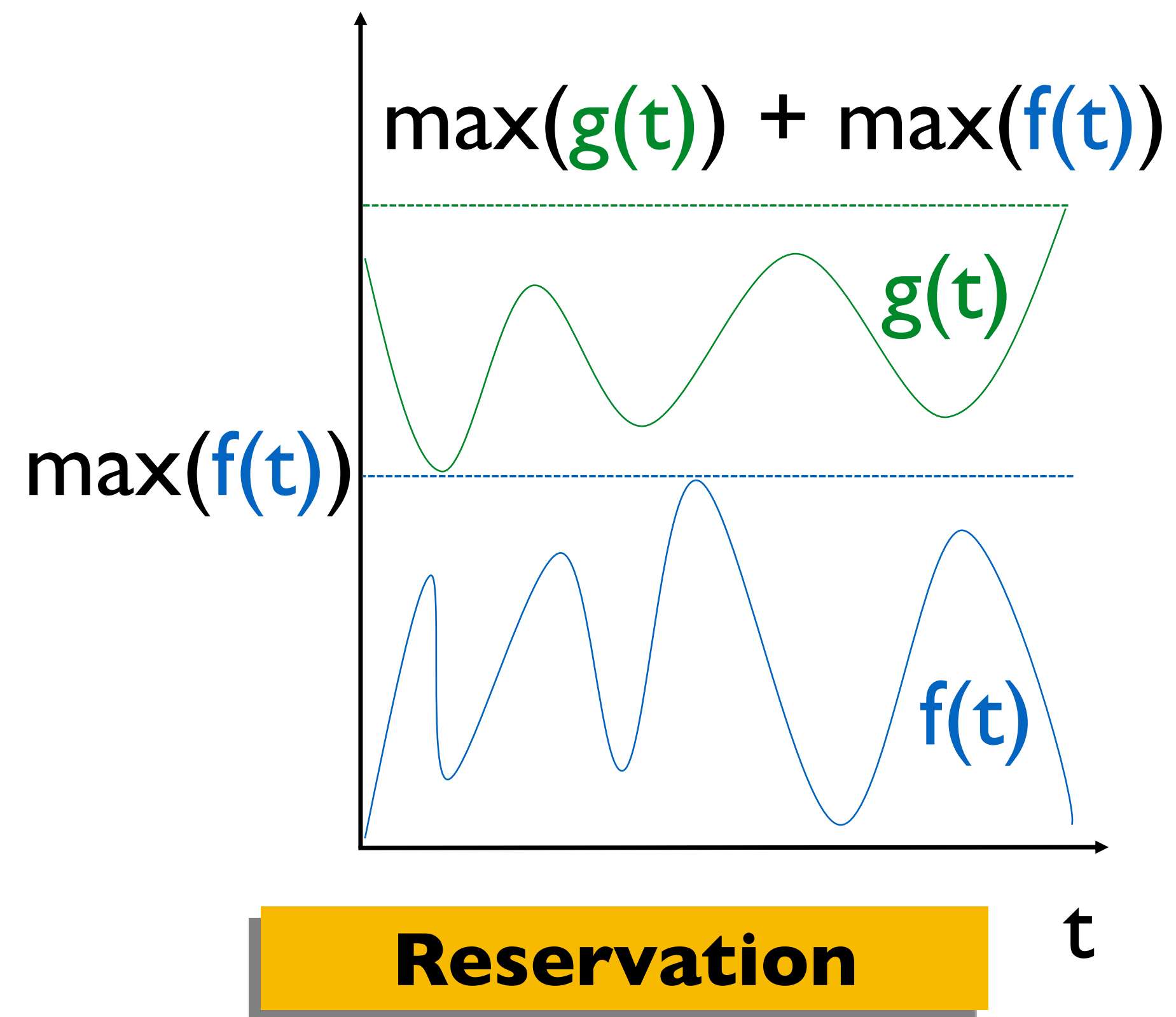
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Questions?

Today's Agenda

- More details on What & How of the Internet
- Sharing a network
- Evaluating a network

Performance Metrics

Performance Metrics

- **Delay:** How long does it take to send a packet from its source to destination?

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Performance Metrics

- **Delay:** How long does it take to send a packet from its source to destination?
- **Loss:** What fraction of the packets sent to the destination are dropped?
- **Throughput:** At what rate is the destination receiving data from the source?

Delay

Delay

- Consists of four components

Delay

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 - **Transmission Delay**

Delay

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 - **Propagation Delay**

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Due to link properties

- Queueing Delay

- Processing Delay

Due to traffic matrix and switch internals

A Network Link

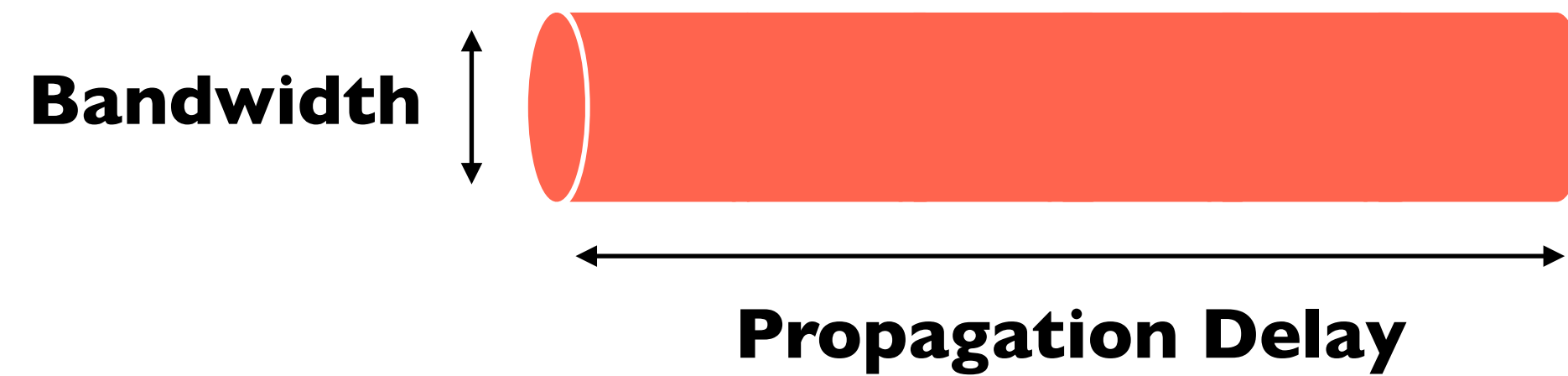


A Network Link



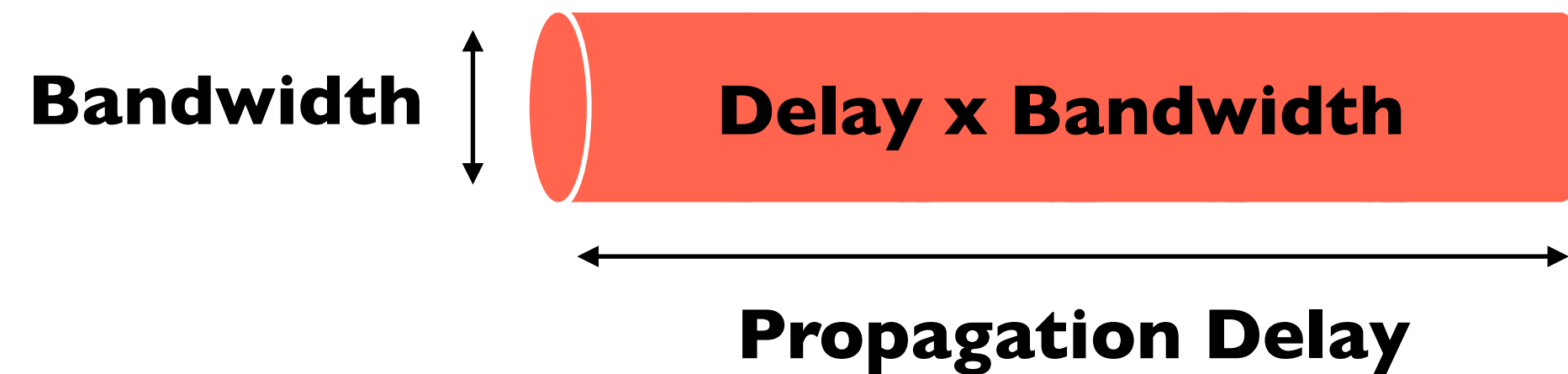
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- **Bandwidth-Delay Product (BDP)**
 - Number of bits “in flight” at any time
 - $BDP = \text{Bandwidth} \times \text{Propagation delay}$

Examples

Examples

- **Same city over a slow link:**
 - Bandwidth: $\sim 100\text{Mbps}$
 - Propagation delay: $\sim 0.1\text{ms}$
 - BDP: 10,000 bits (1.25KB)

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- **Cross country over a fast link:**
 - Bandwidth: $\sim 10\text{Gbps}$
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 - BDP: 10^8 bits (12.5MB)

Transmission Delay

Transmission Delay

- How long does it take to put all the bits of a packet into a link?

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- $\text{Packet size} / \text{Transmission rate of link}$

Transmission Delay

- How long does it take to put all the bits of a packet into a link?
- Packet size / Transmission rate of link
- **Example:**
 - Packet size: 1000 bits
 - Rate: 100 Mbits/s
 - Delay: $1000 \text{ bits} / (100 * 10^6 \text{ Mbits/s}) = 10^{-5} \text{ s}$

Propagation Delay

Propagation Delay

- How long does it take to one bit from one end of the link to the other?

Propagation Delay

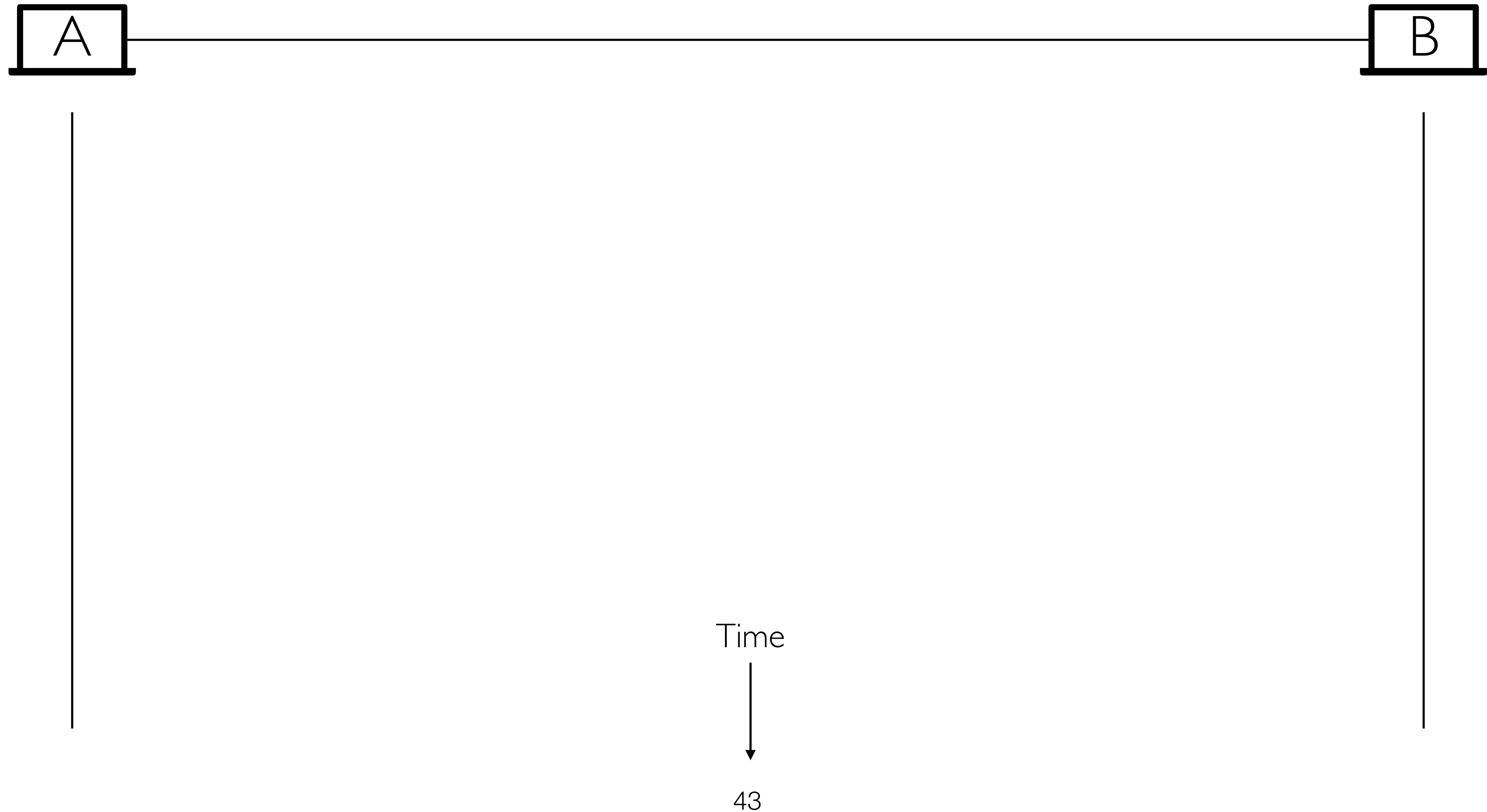
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Propagation Delay

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- Link length / Propagation speed of link
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- **Example:**
 - **Length:** 30 km = 30,000 m
 - **Delay:** $30,000 \text{ m} / (3 * 10^8 \text{ m/s}) = 10^{-4} \text{ s}$

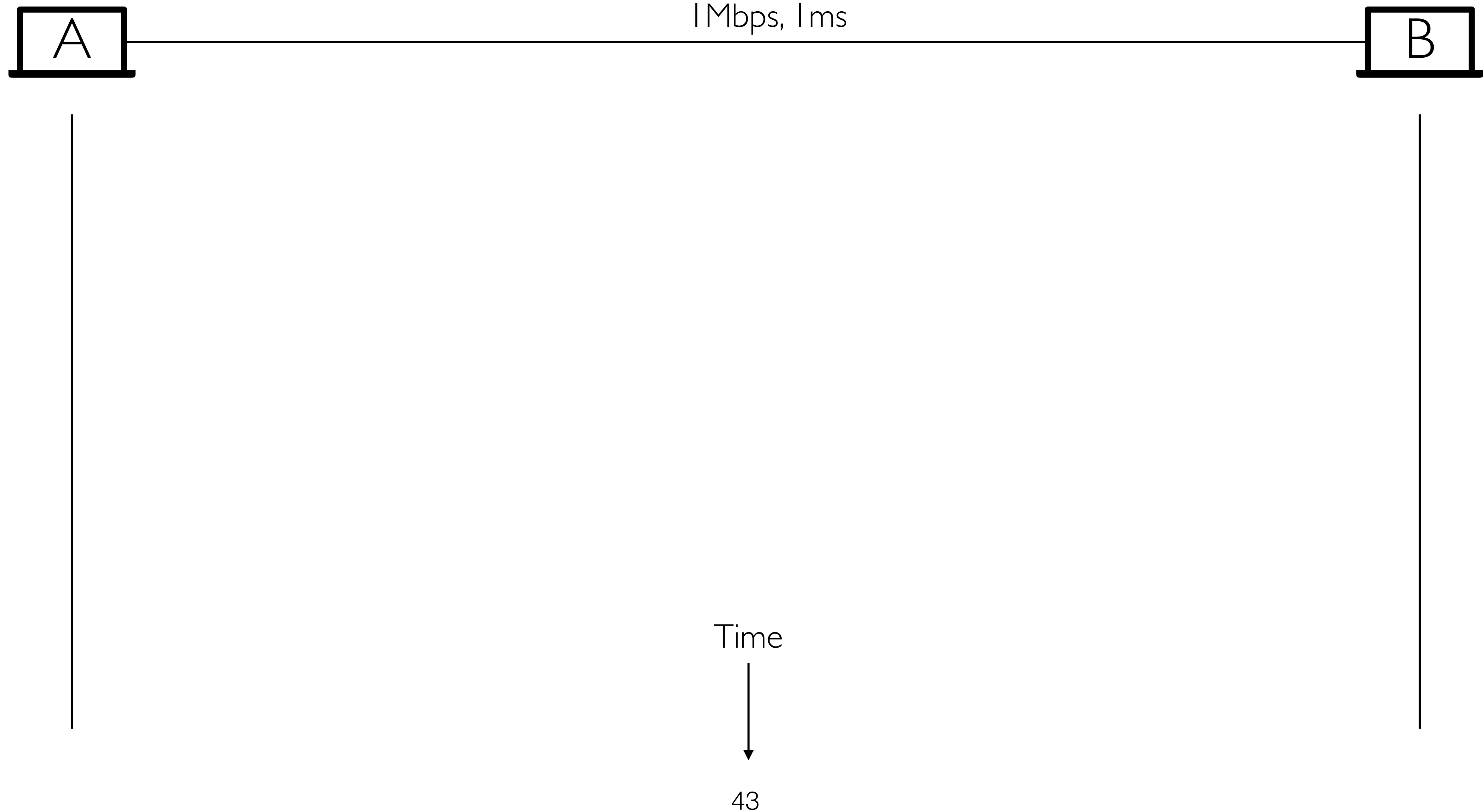
Packet Delay

Sending 100B packets from A to B?



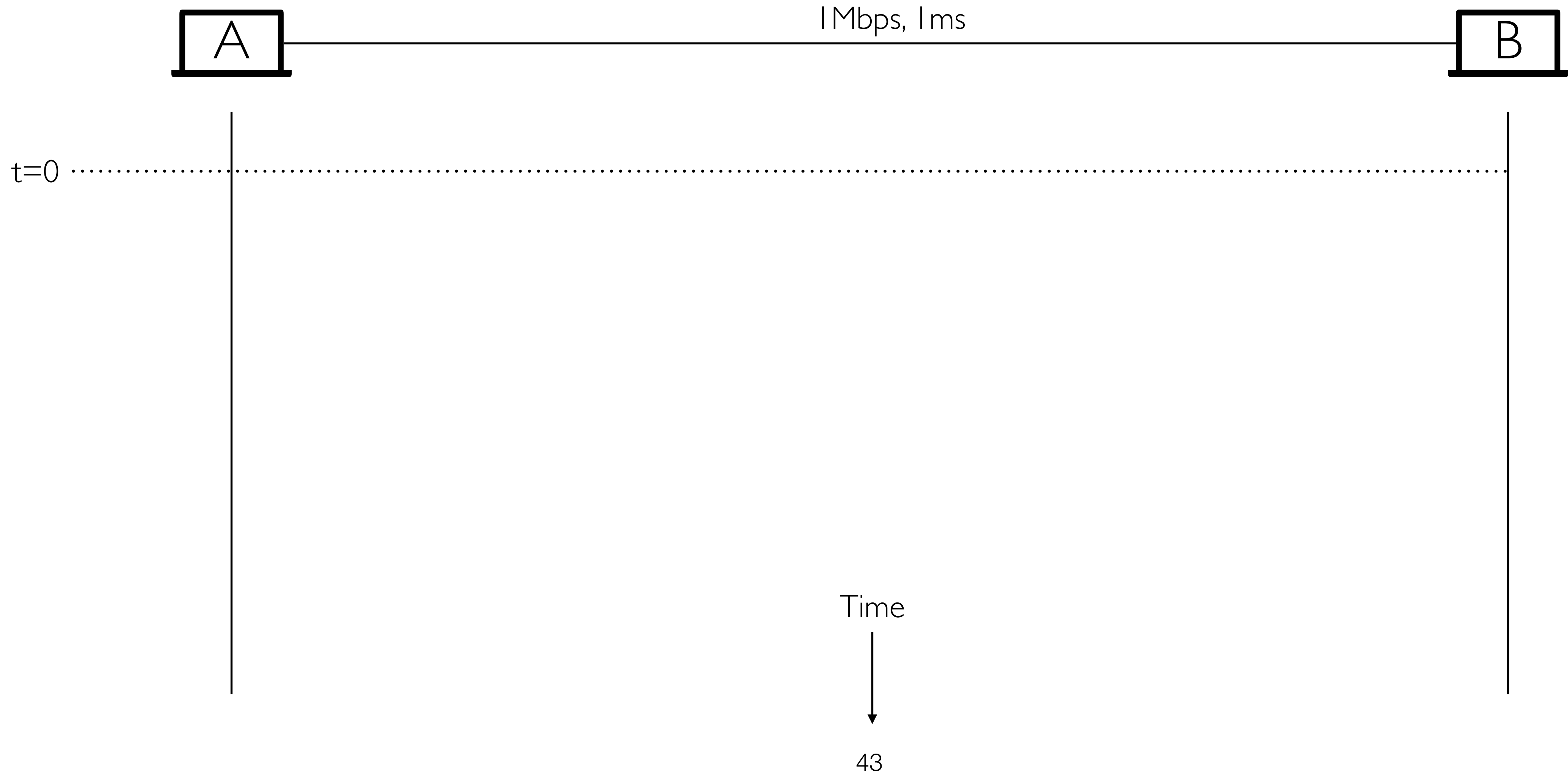
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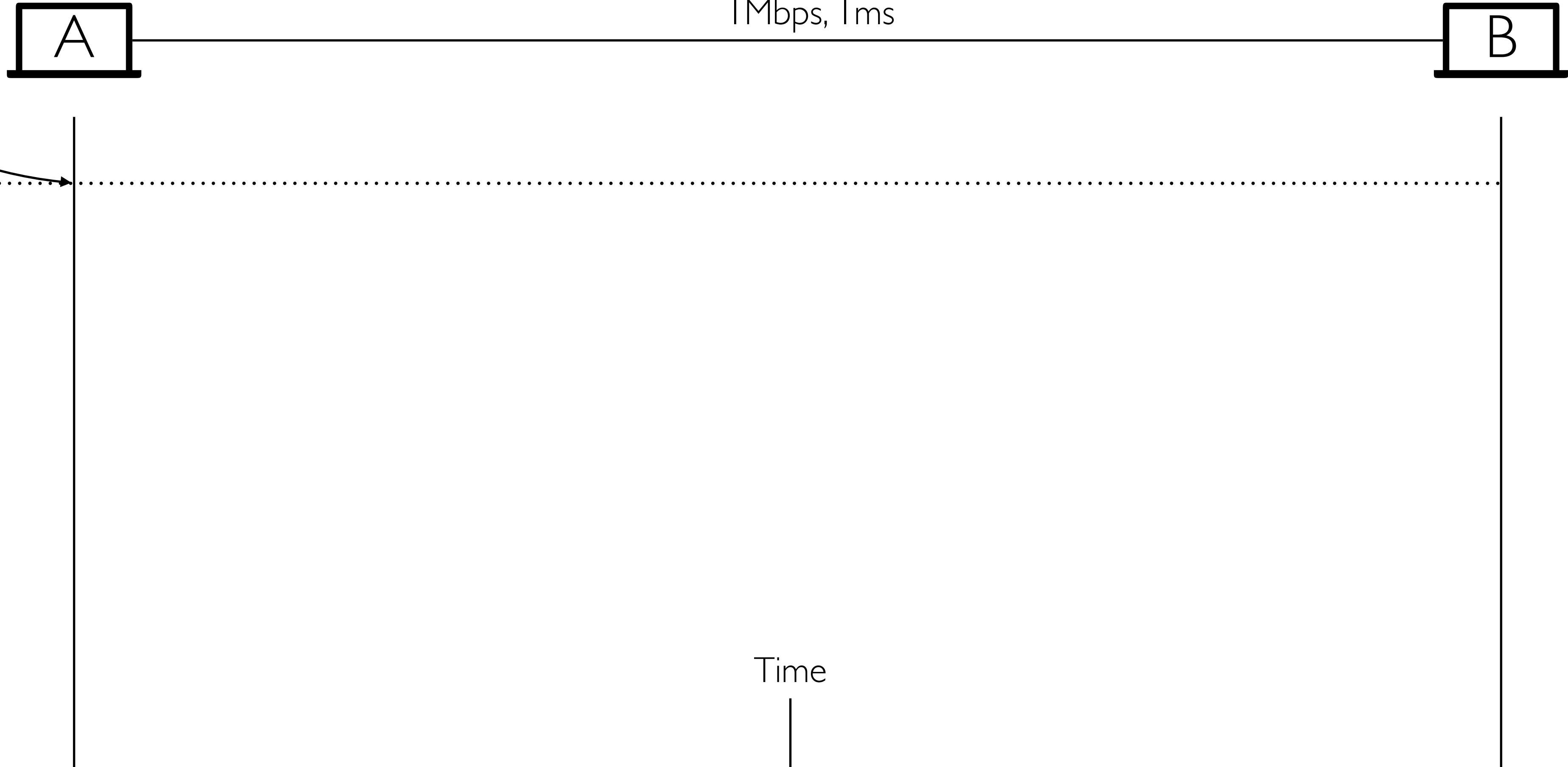
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1Mbps, 1ms

Time to transmit
one bit = $1/10^6$ s

t=0

Time



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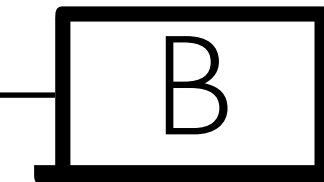
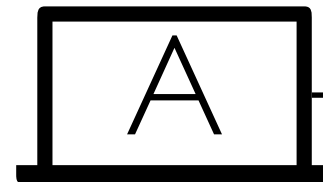
Time when that bit
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 $= 1/10^6 + 1/10^3$ s

Time

Packet Delay

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Time to transmit
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Packet Delay

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Data

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Data

Time when last bit
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 $= 1.8$ ms

Time

Packet Delay

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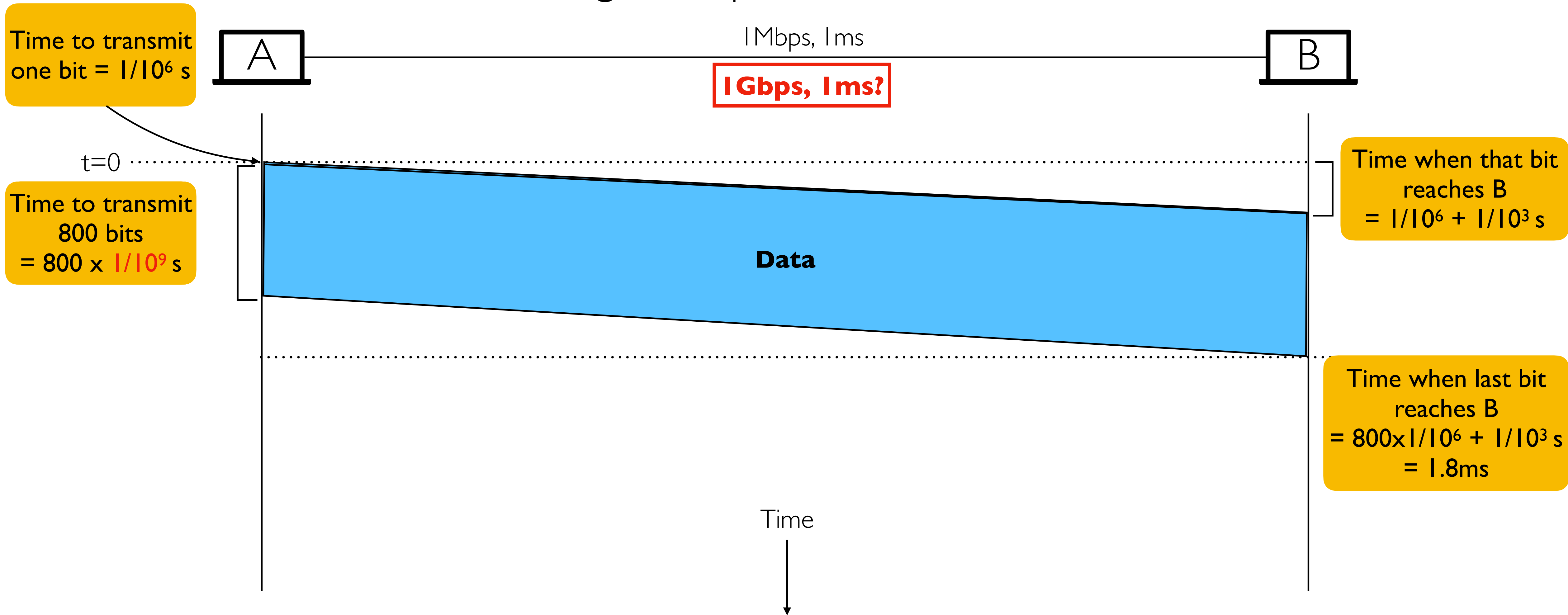
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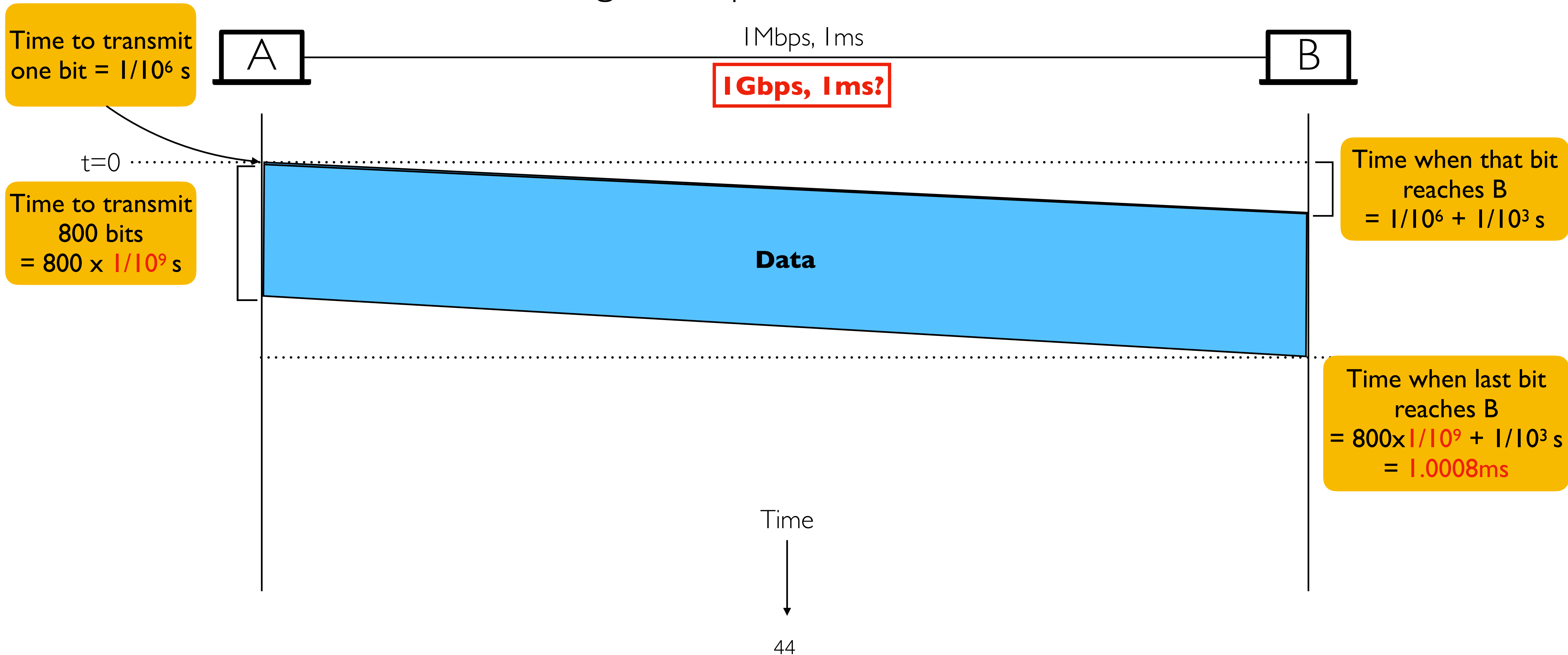
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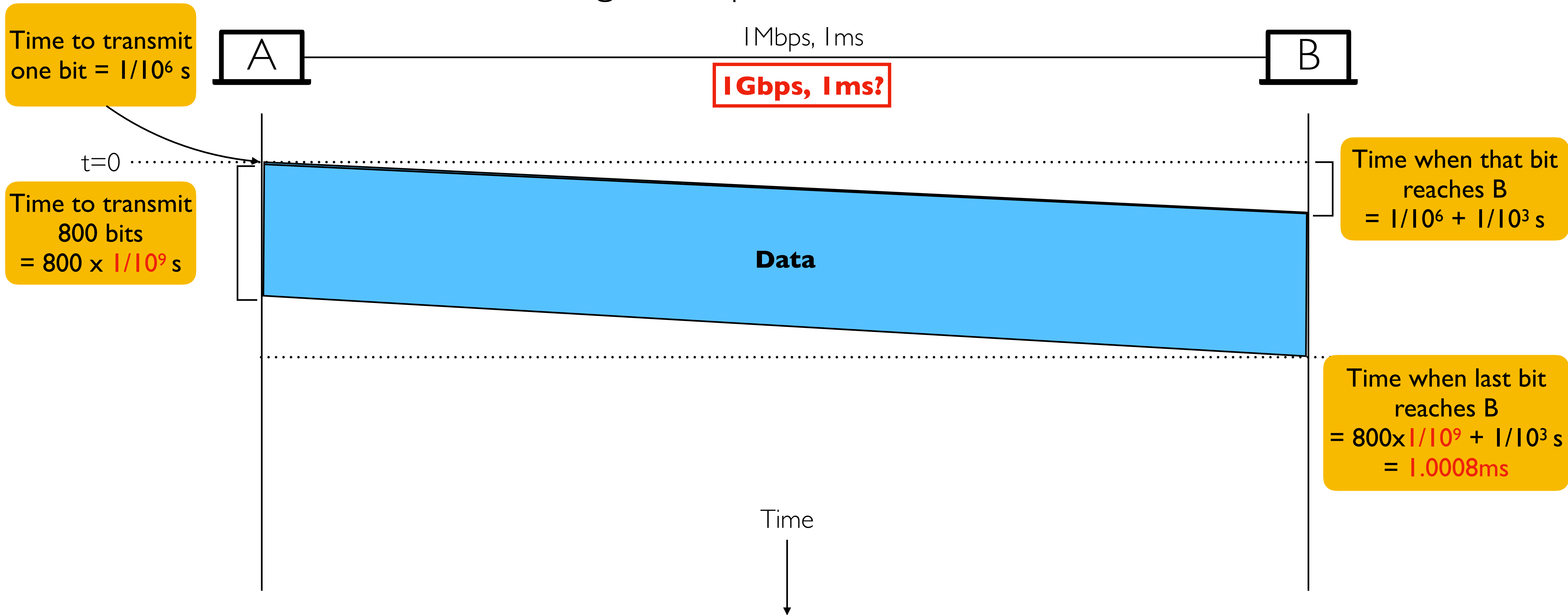
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1 GB file in 100B packets

Packet Delay

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1 Mbps, 1 ms

1 Gbps, 1 ms?

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Time to transmit
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= $800 \times 1/10^9$ s

Time when that bit
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Time when last bit
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= 1.0008 ms

Time

1 GB file in 100B packets

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Time to transmit
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 $= 10^7 \times 800 \times$
 $1/10^9$ s

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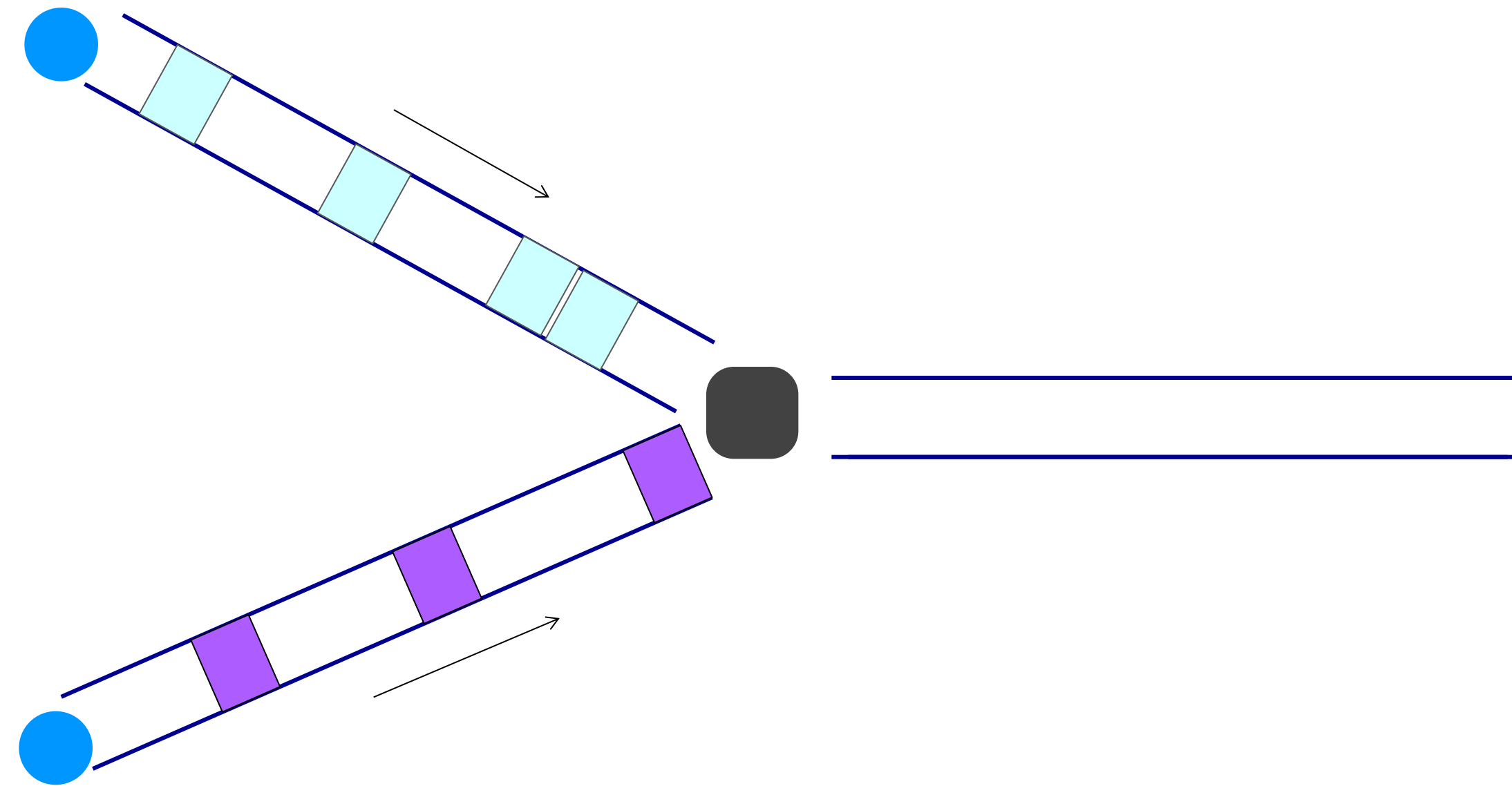
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Time

Questions?

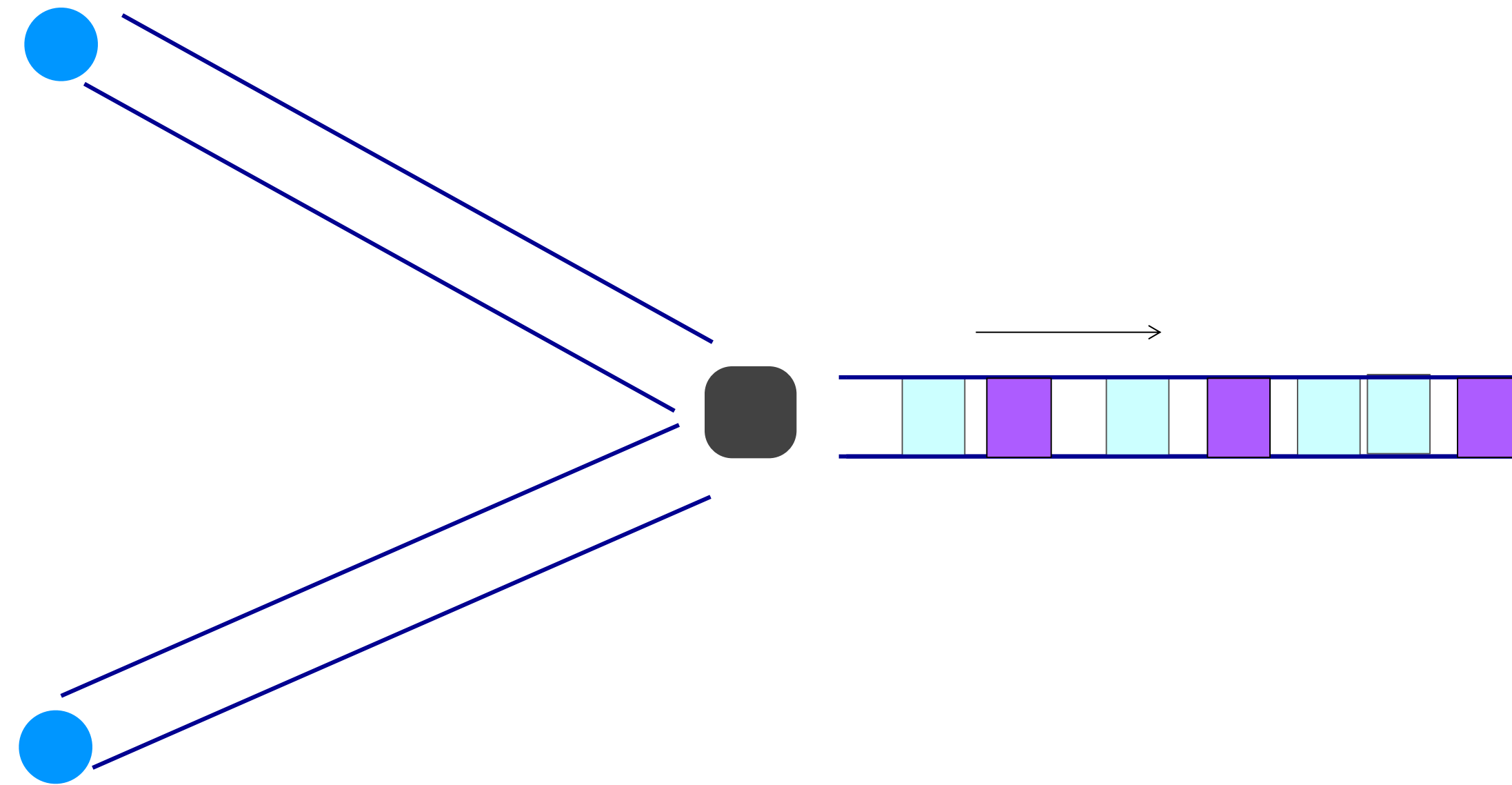
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- How long does a packet have to sit in a buffer before it is processed



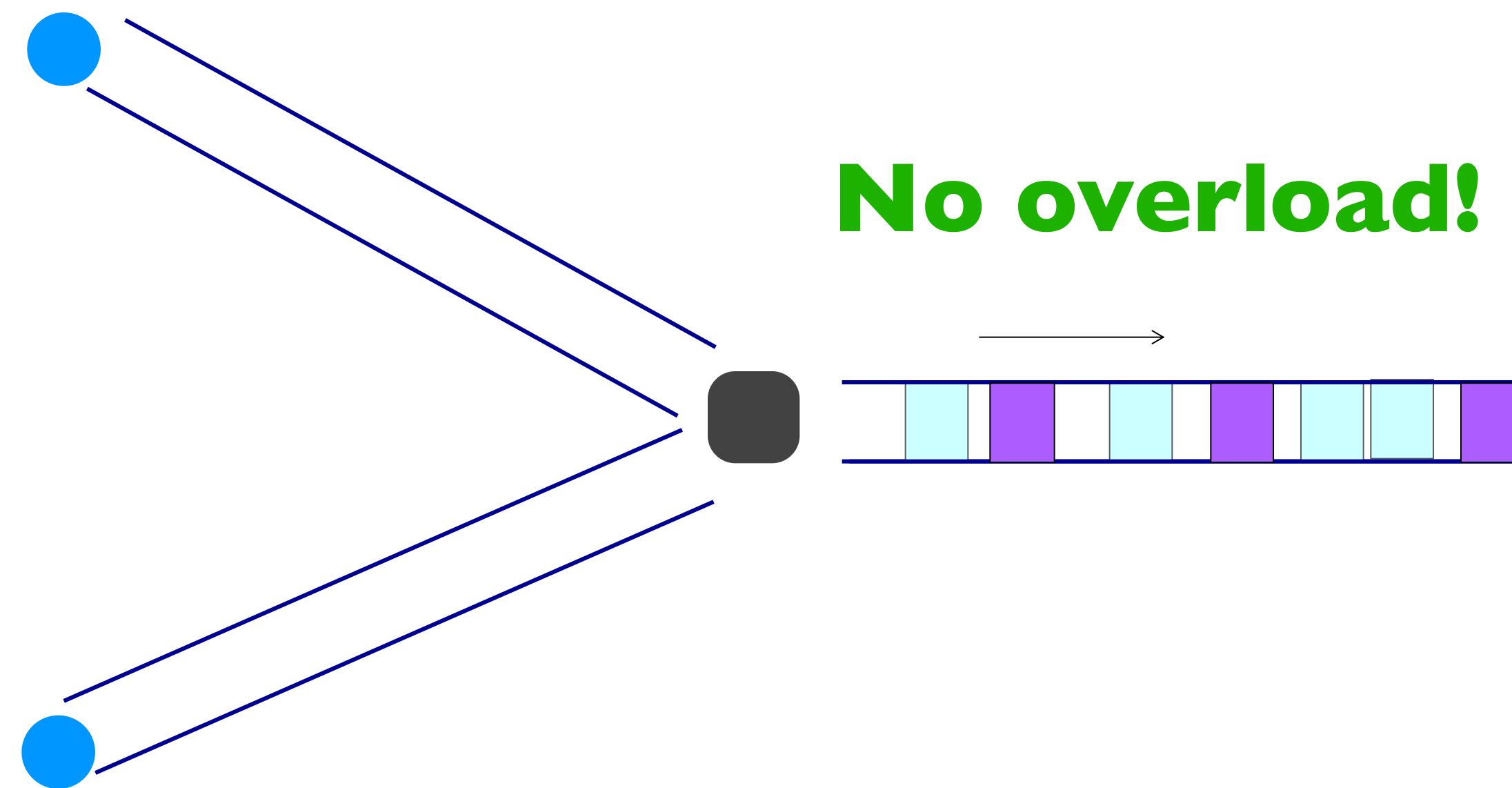
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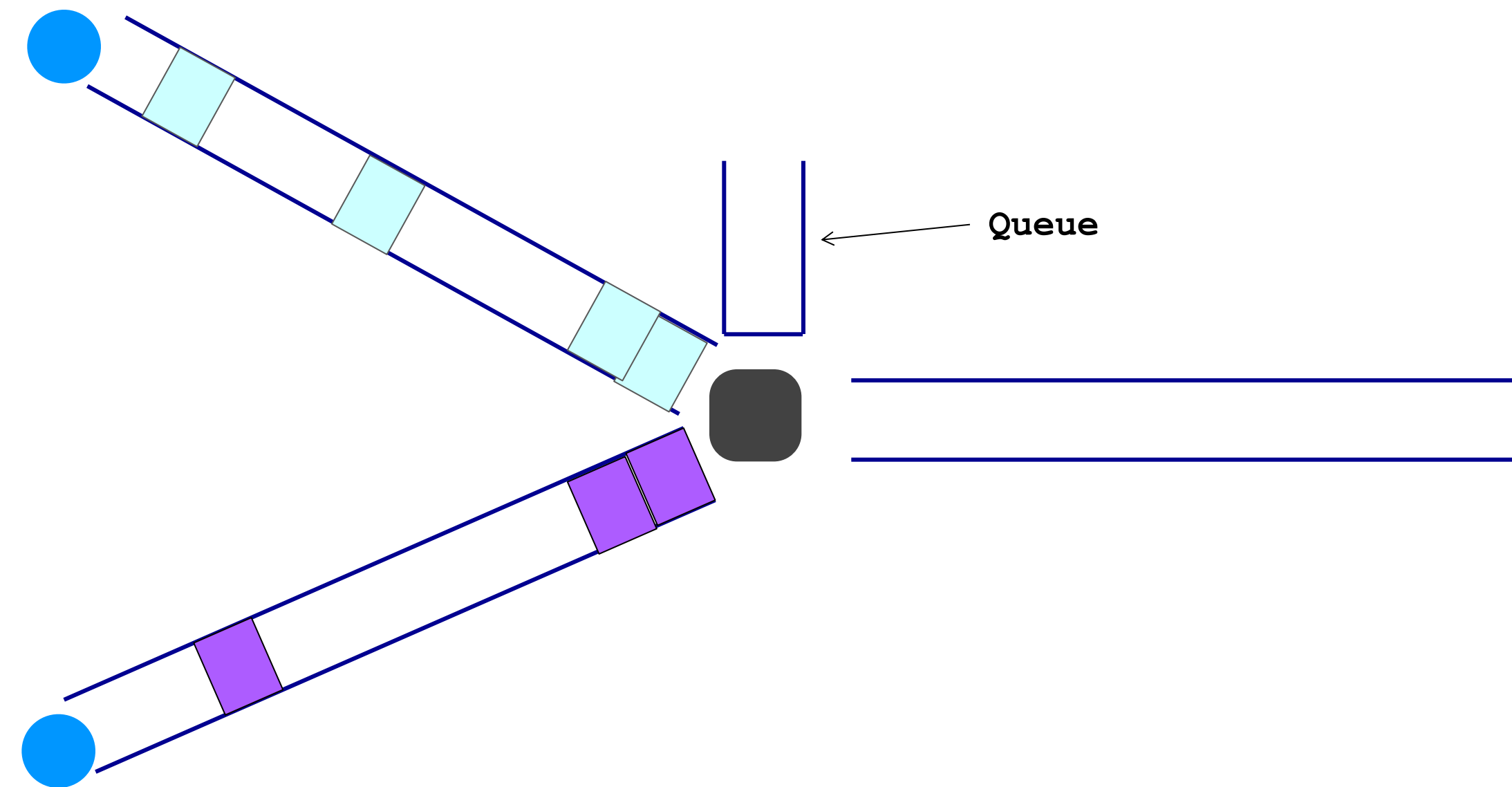
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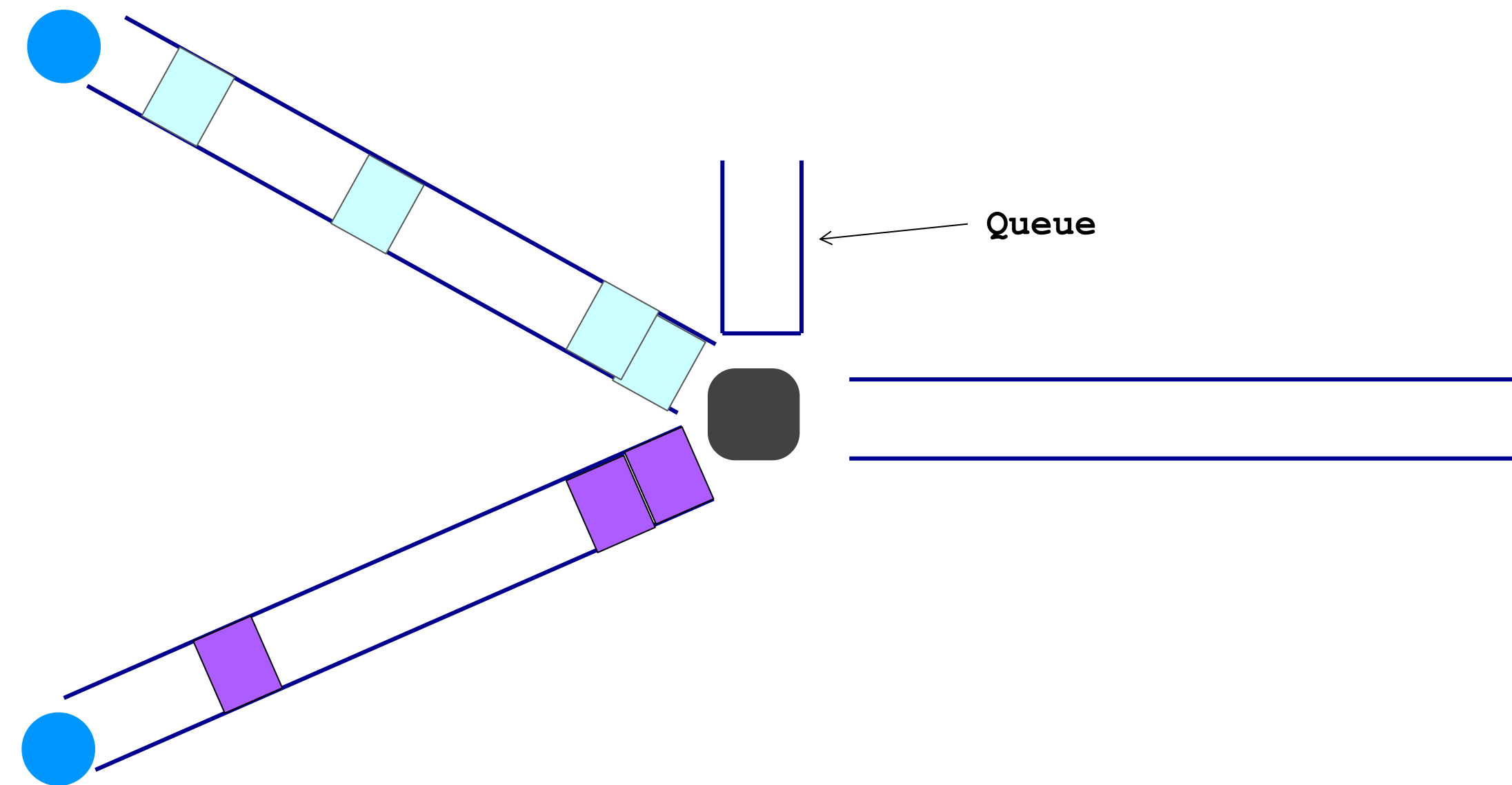
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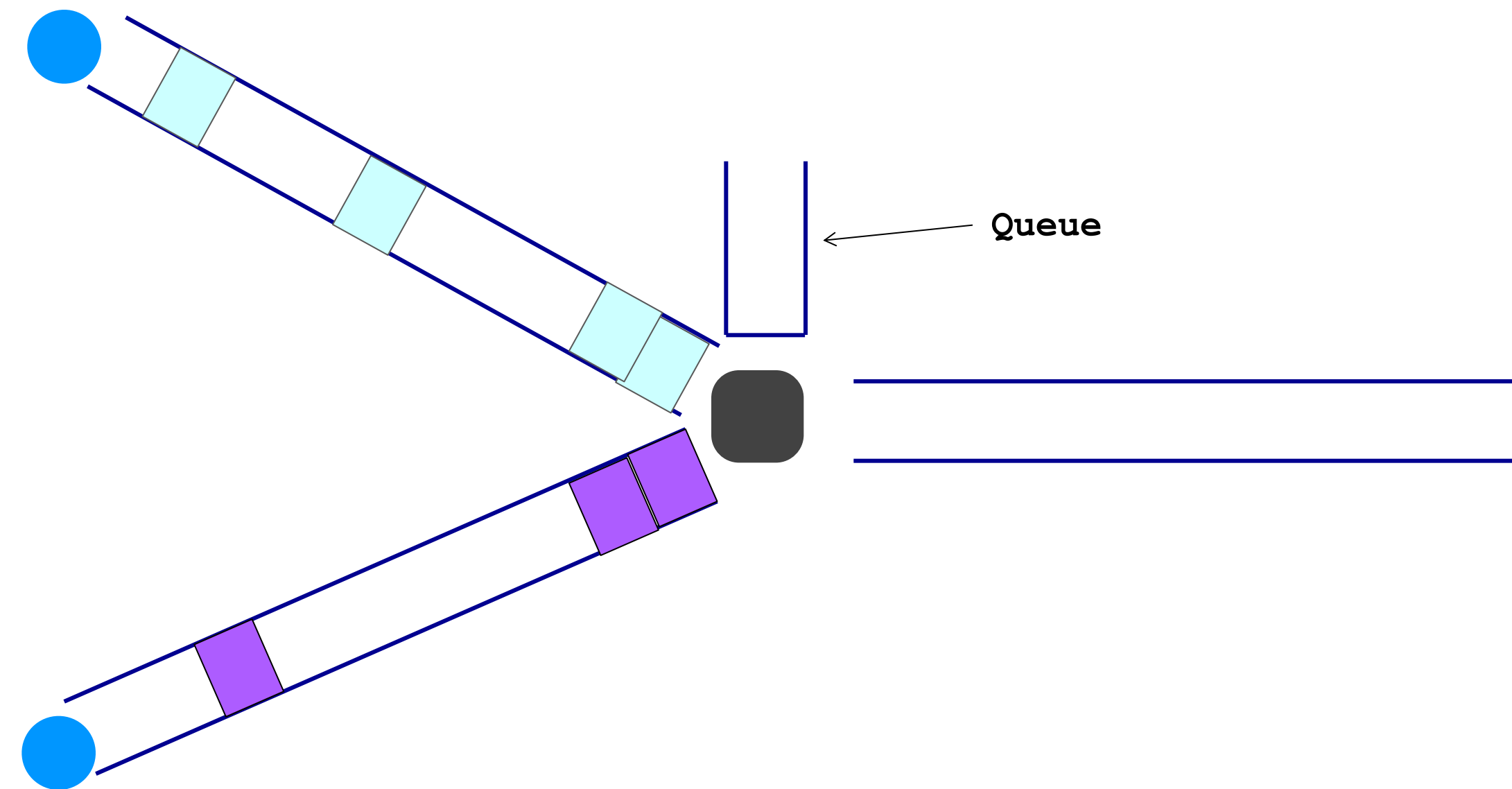
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Transient overload!

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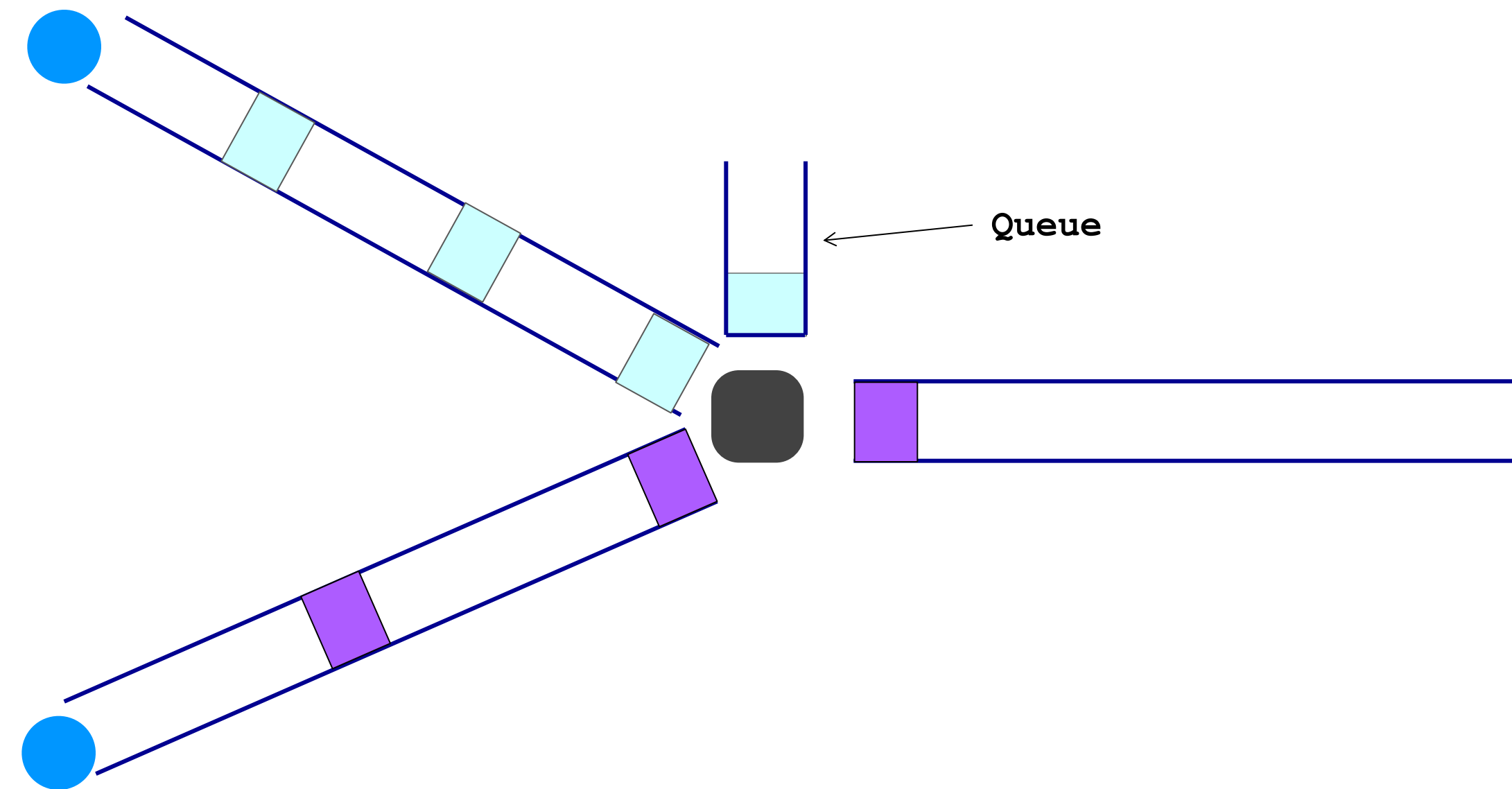


Transient overload!

Not a rare event.

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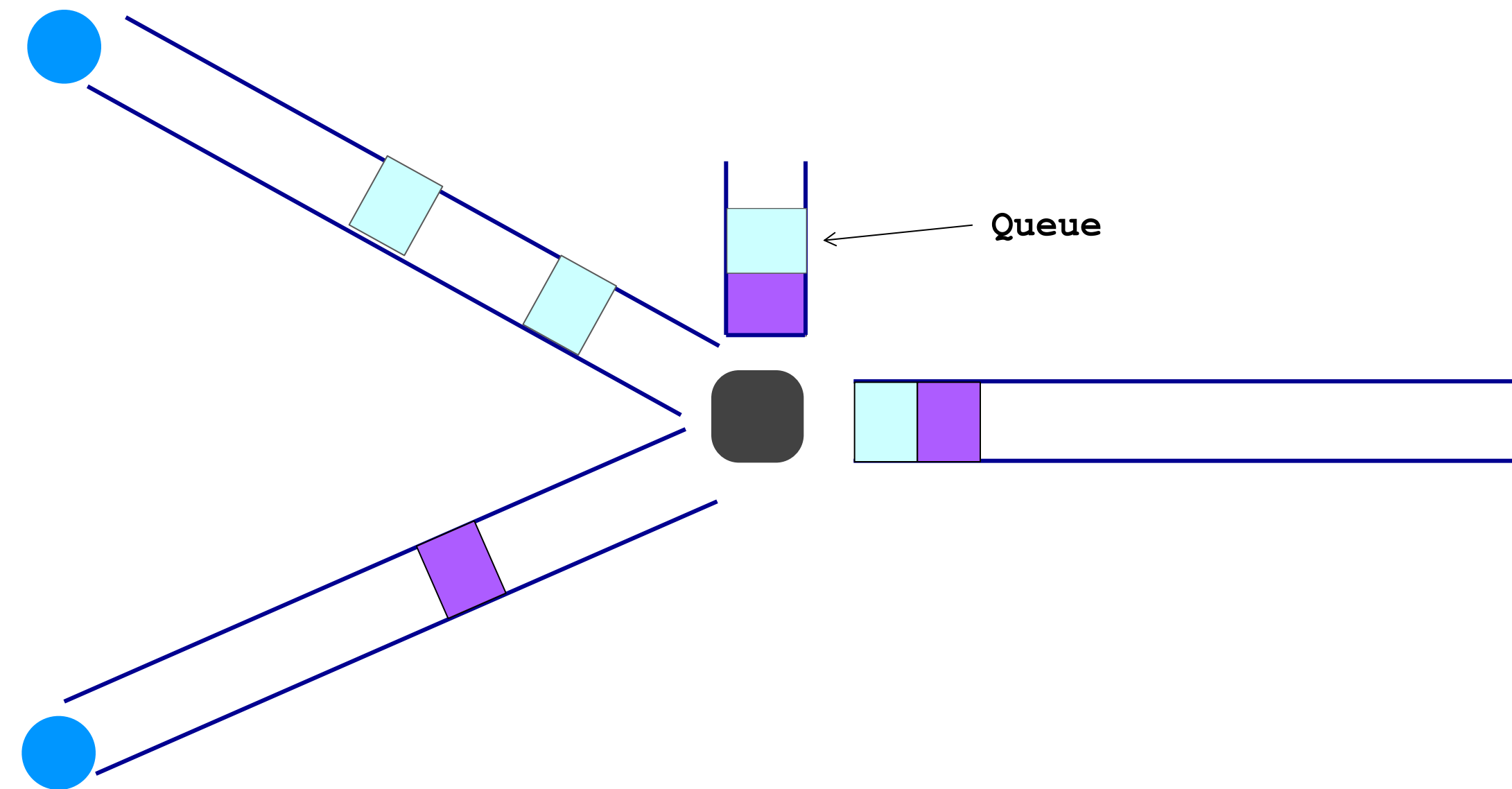


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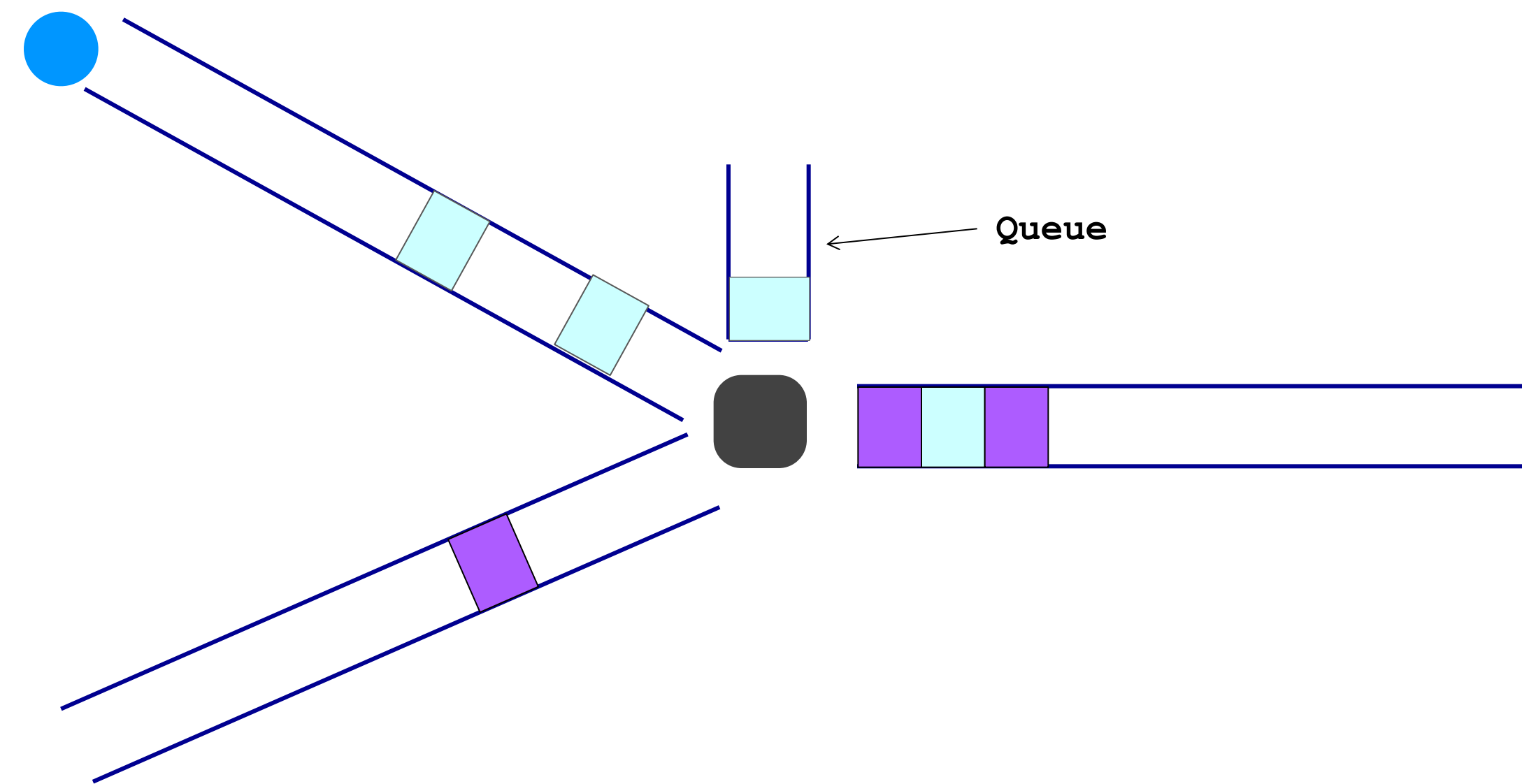


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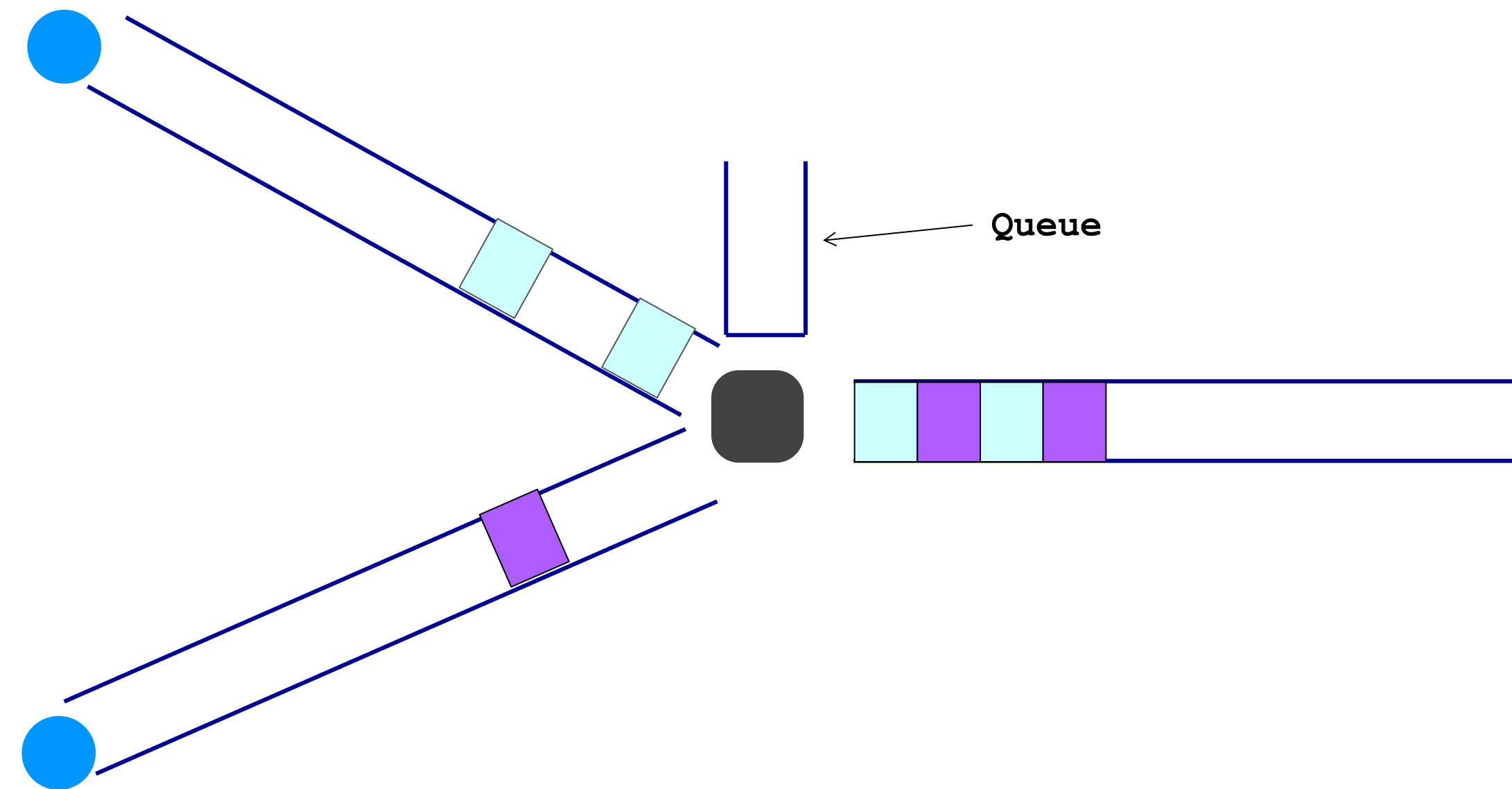


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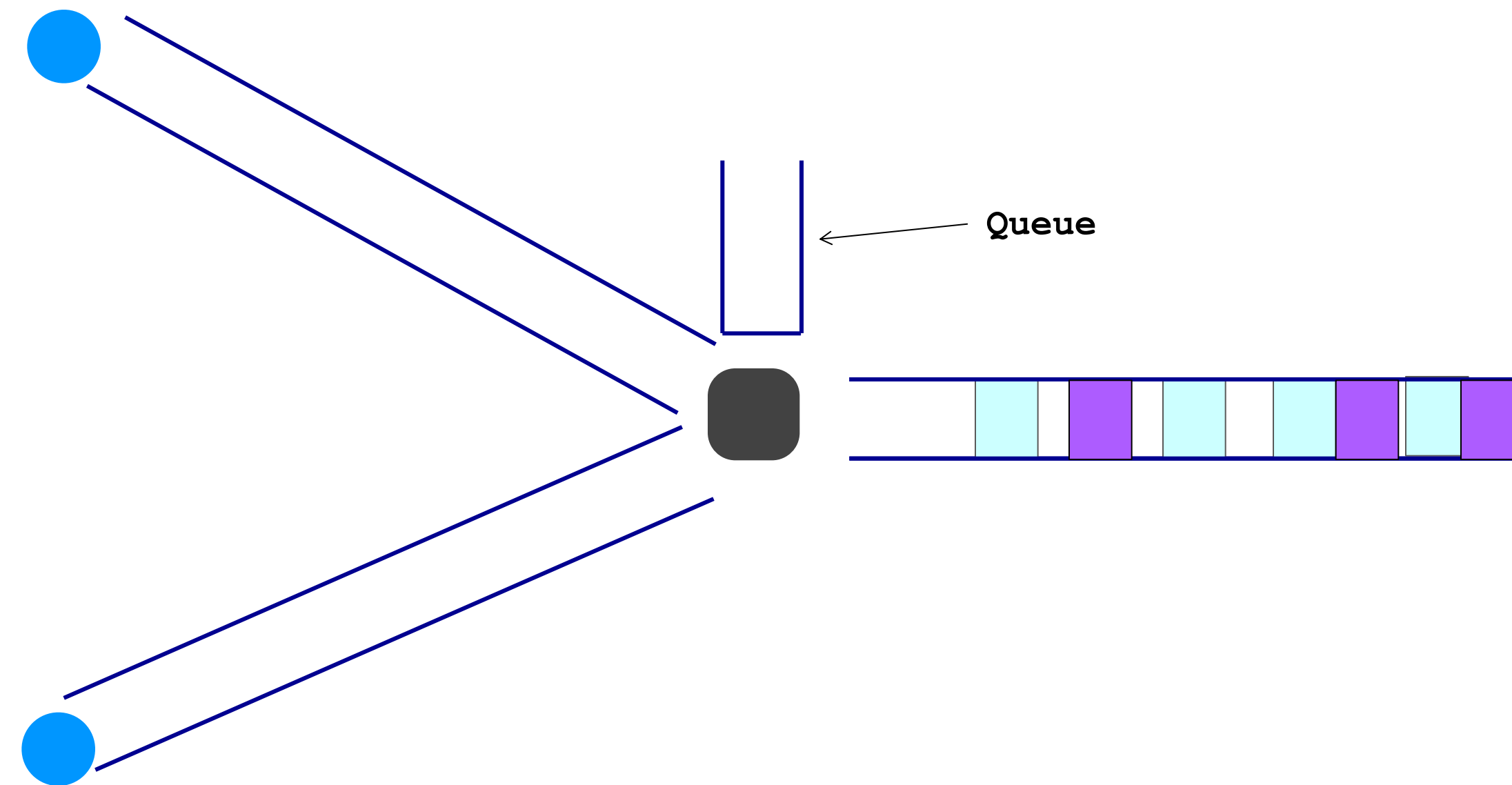


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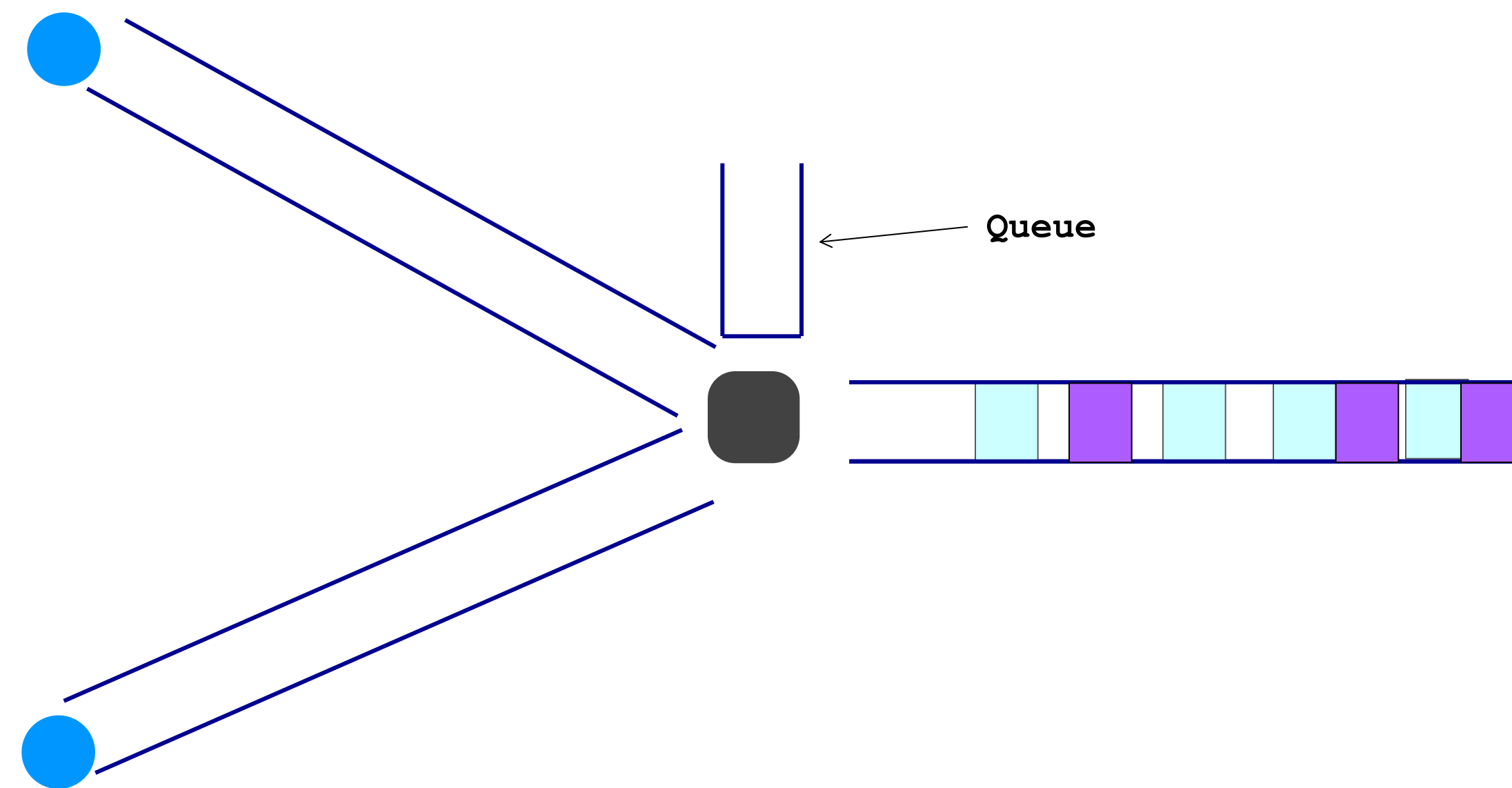


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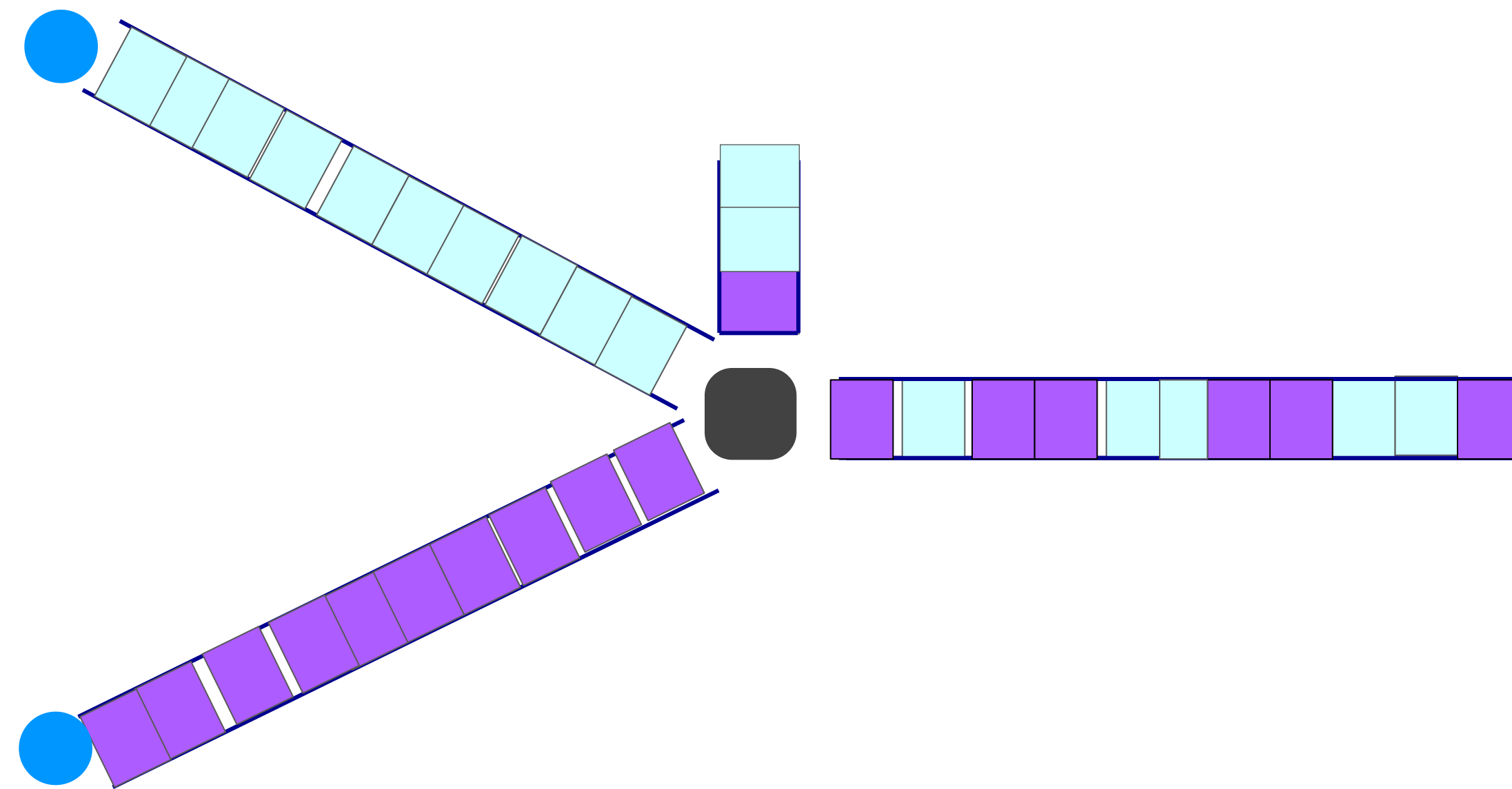
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Queues absorb transient bursts but introduce queueing delay

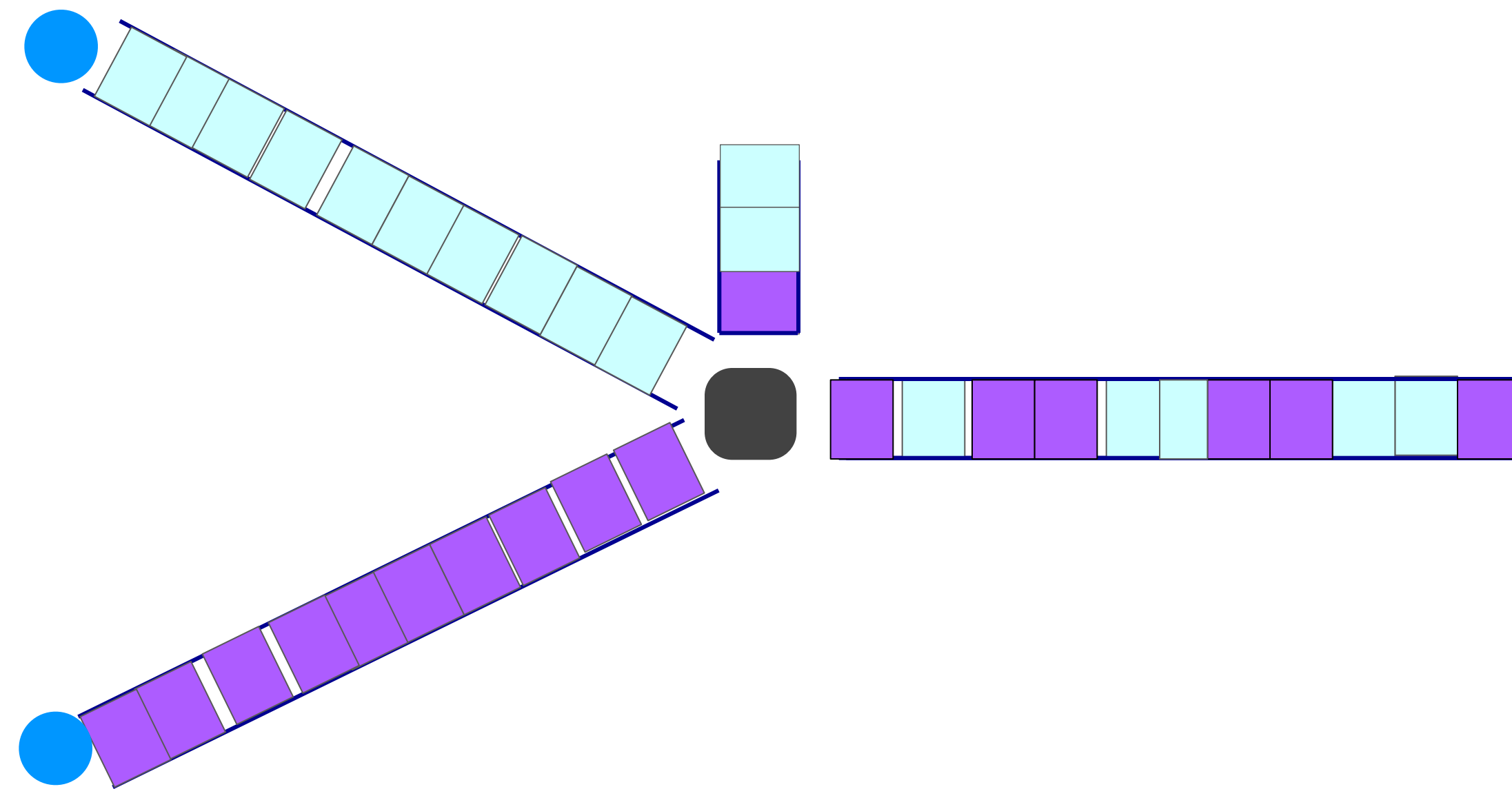
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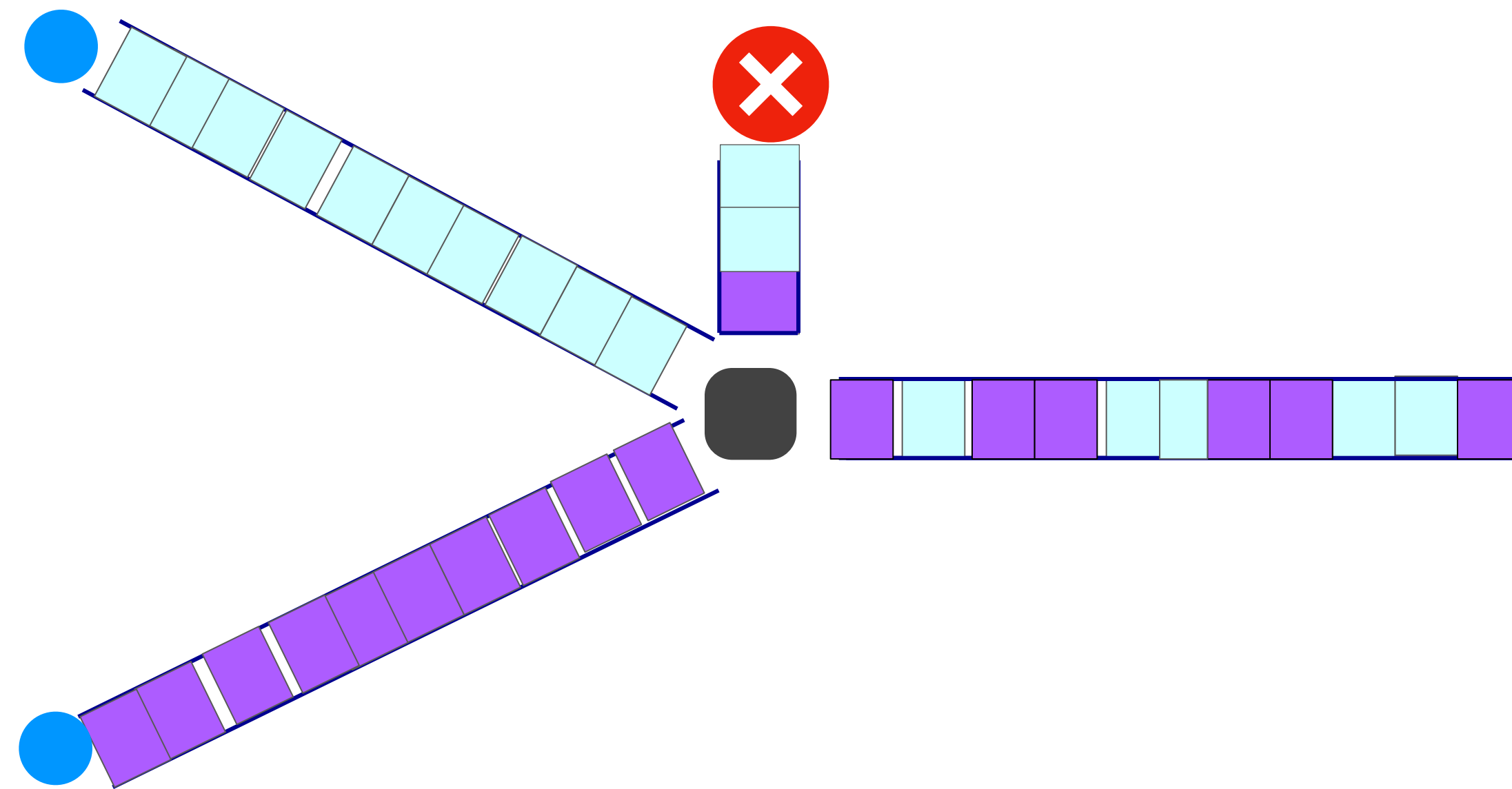
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What about persistent overload?

Queueing Delay: “Pipe” View

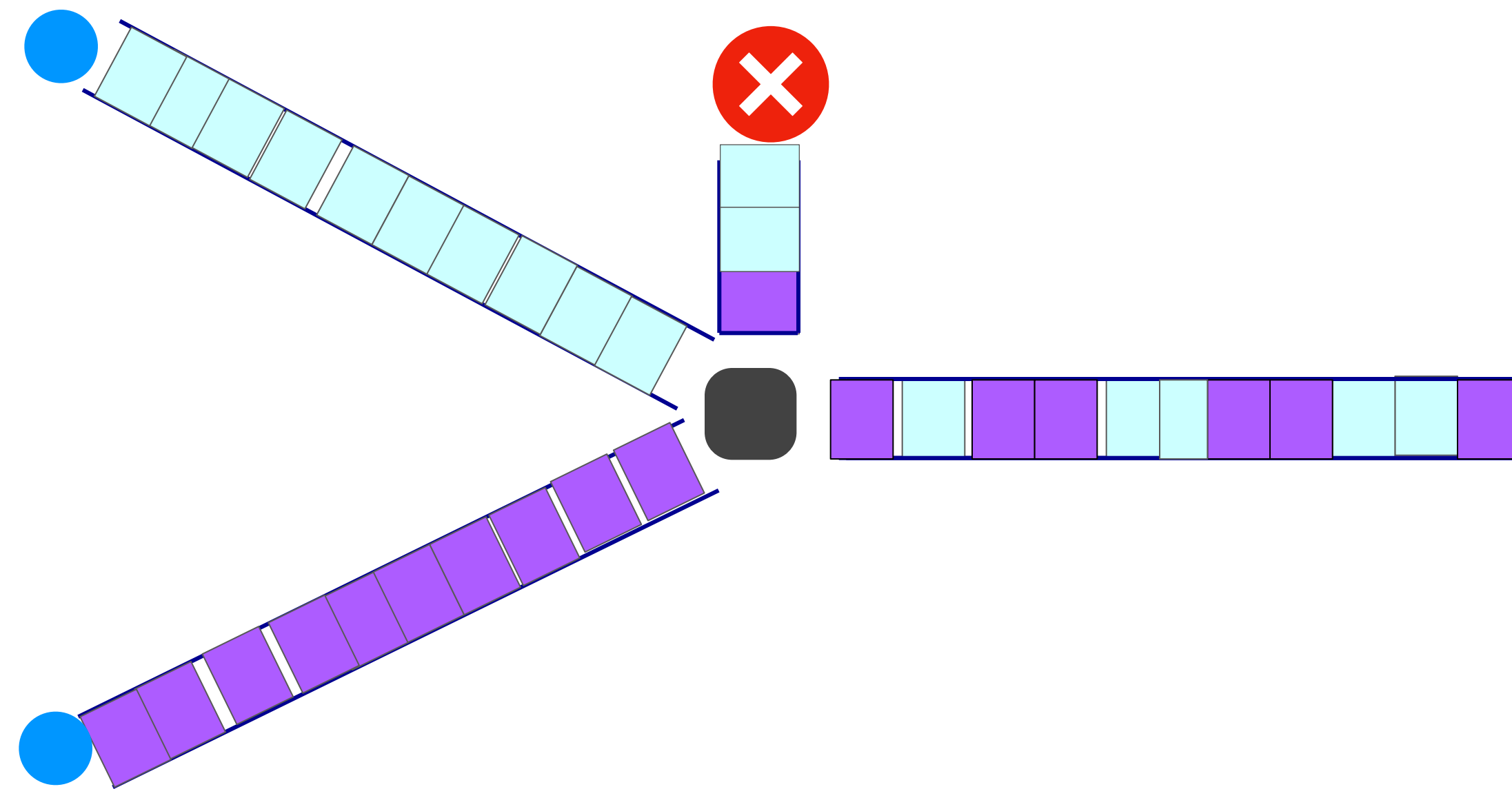
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Queueing Delay: “Pipe” View

- How long does a packet have to sit in a buffer before it is processed



What about persistent overload?

Leads to packet loss

Queueing Delay

Queueing Delay

- **If arrival rate $>$ departure rate**
 - Approaches infinity (assuming an infinite buffer)
 - In practice, finite buffer = loss

Queueing Delay

- **If arrival rate $>$ departure rate**
 - Approaches infinity (assuming an infinite buffer)
 - In practice, finite buffer = loss
- **If arrival rate $<$ departure rate**
 - Depends on burst rate

Queueing Delay

- How long does a packet have to sit in a buffer before it is processed

Queueing Delay

- How long does a packet have to sit in a buffer before it is processed
- Depends on traffic pattern

Queueing Delay

- How long does a packet have to sit in a buffer before it is processed
- Depends on traffic pattern
- Characterized by statistical measures
 - Average queueing delay
 - Average arrival rate
 - Average departure rate

Basic Queueing Theory Terminology

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- **Arrival process:** how packets arrive
 - Characterized by average arrival rate **λ**

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Basic Queueing Theory Terminology

- **Arrival process:** how packets arrive
 - Characterized by average arrival rate **A**
- **W:** average time packets wait in the queue
 - W for “waiting time”
- **L:** average number of packets waiting in the queue
 - L for “length of the queue”

Little's Law (1969)

A: avg. packet arrival rate (/s)

L: avg. # of packets waiting in queue

W: avg. time packets wait in queue

Little's Law (1969)

A: avg. packet arrival rate (/s)

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$$L = A \times W$$

Little's Law (1969)

A: avg. packet arrival rate (/s)

L: avg. # of packets waiting in queue

W: avg. time packets wait in queue

$$L = A \times W$$

or,

$$W = L / A$$

Processing Delay

Processing Delay

- How long does a switch take to process this packet?
 - Typically assume this is negligible

Delay

- Consists of four components

- Transmission Delay

- Propagation Delay

Due to link properties

- Queueing Delay

- Processing Delay

Due to traffic matrix and switch internals

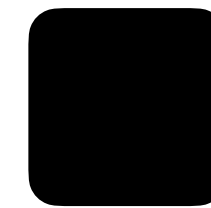
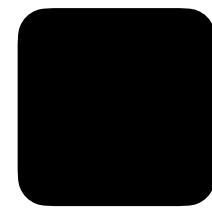
End-to-end Delay

T Transmission Delay

Q Queuing Delay

P Propagation Delay

Pr Processing Delay



Time

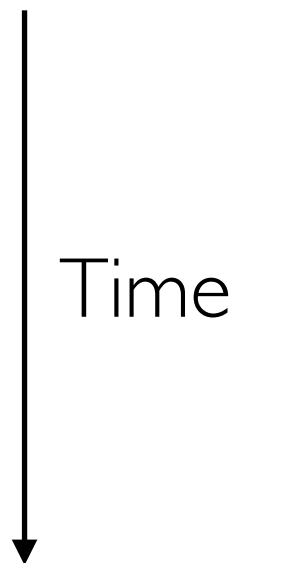
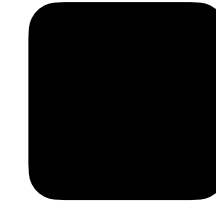
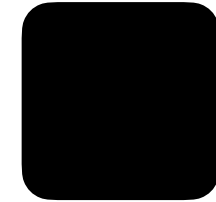
End-to-end Delay

T Transmission Delay

Q Queuing Delay

P Propagation Delay

Pr Processing Delay



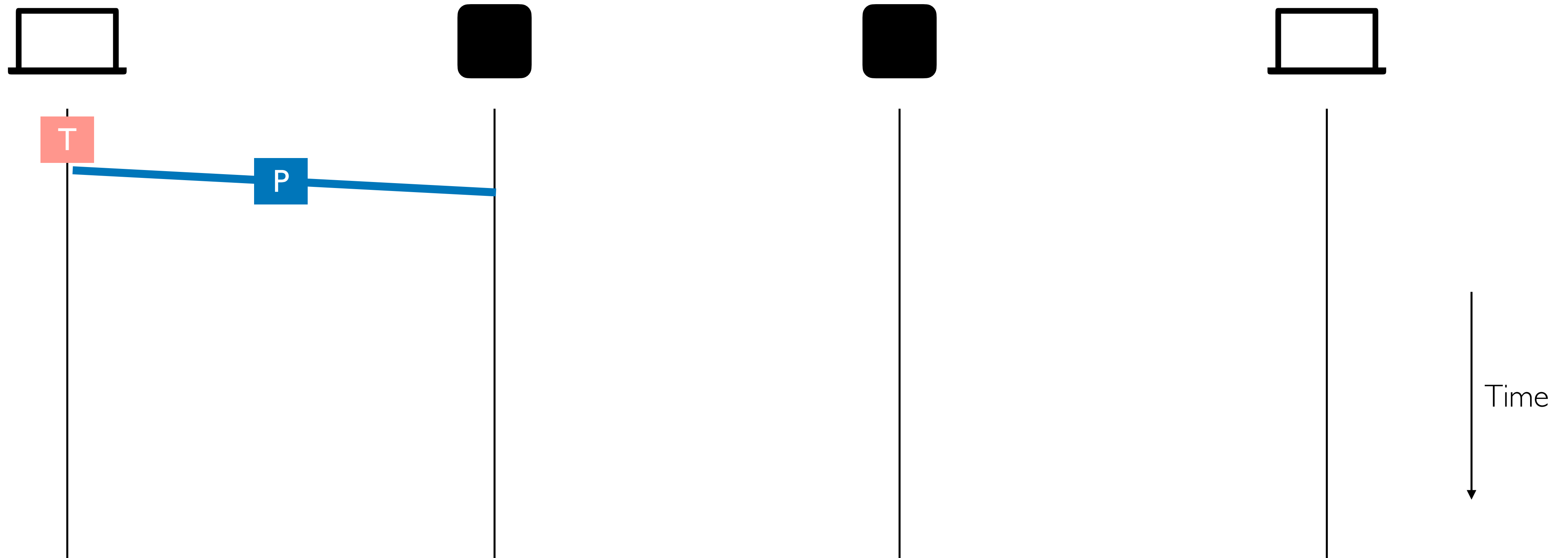
End-to-end Delay

T Transmission Delay

Q Queuing Delay

P Propagation Delay

Pr Processing Delay



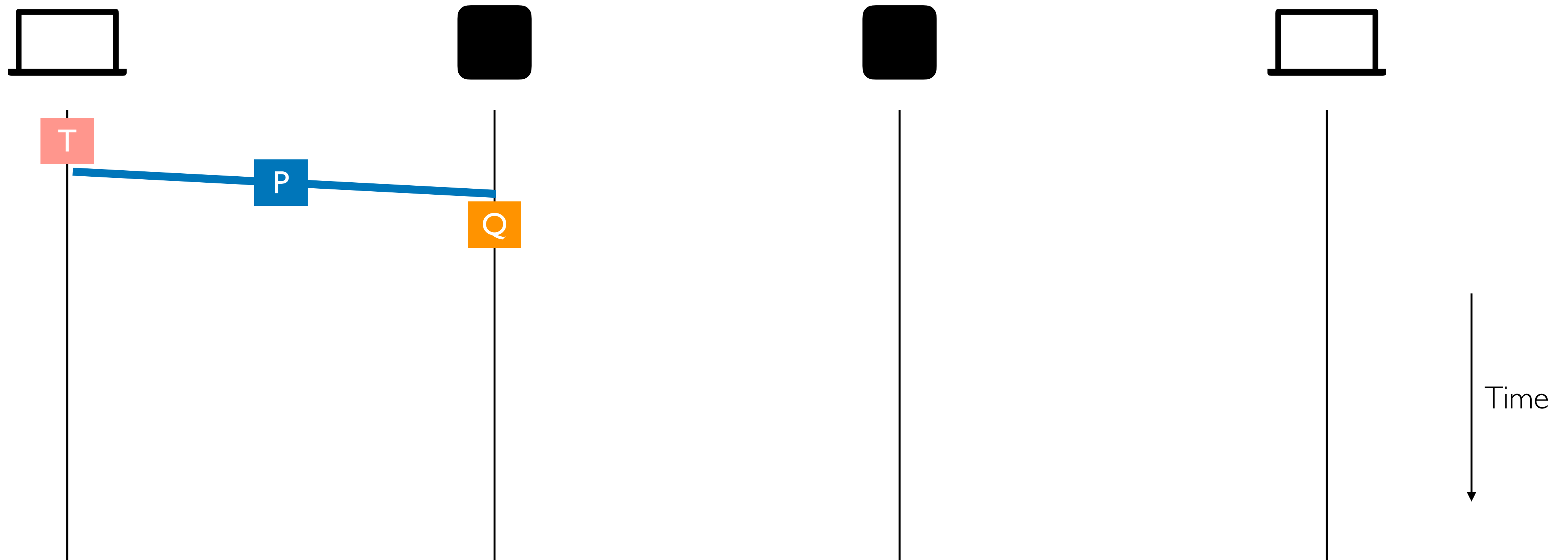
End-to-end Delay

T Transmission Delay

Q Queuing Delay

P Propagation Delay

Pr Processing Delay



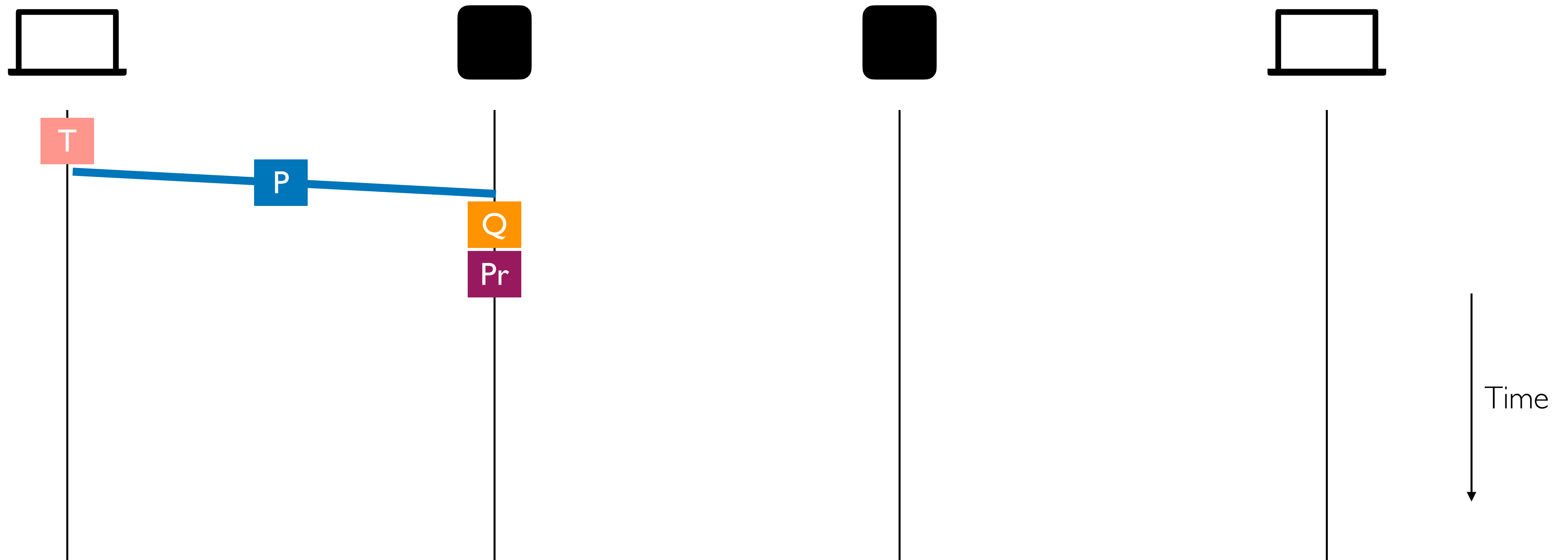
End-to-end Delay

T Transmission Delay

Q Queuing Delay

P Propagation Delay

Pr Processing Delay



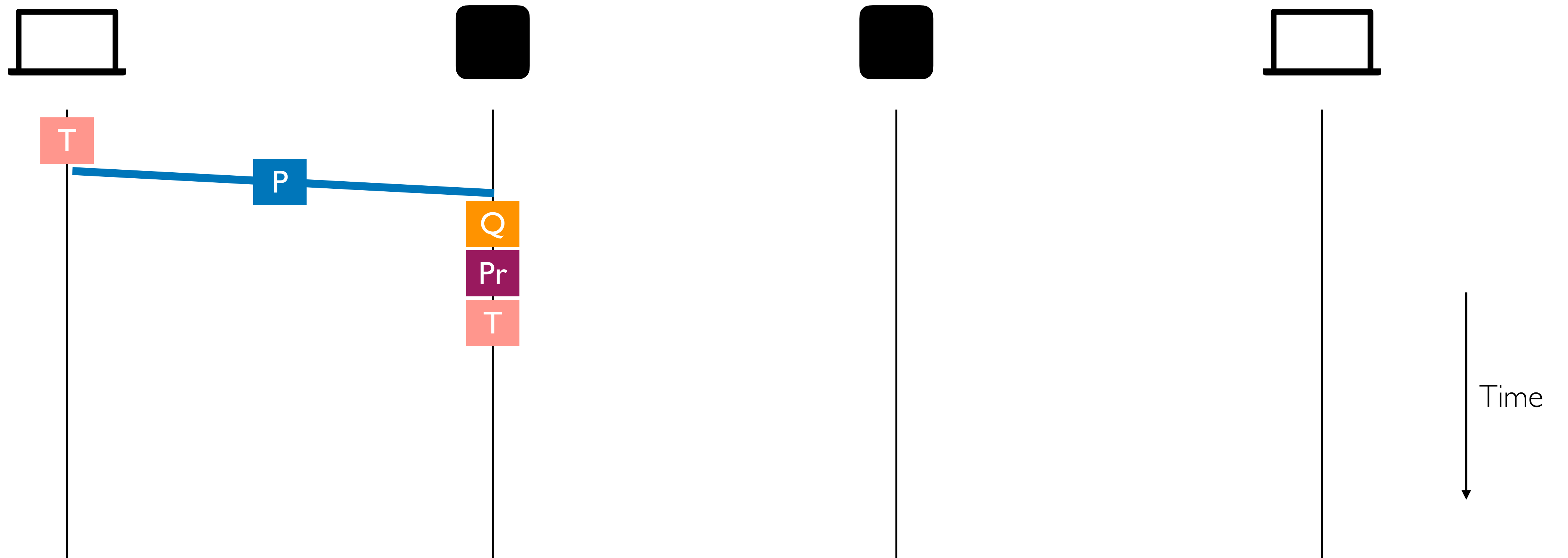
End-to-end Delay

T Transmission Delay

Q Queuing Delay

P Propagation Delay

Pr Processing Delay



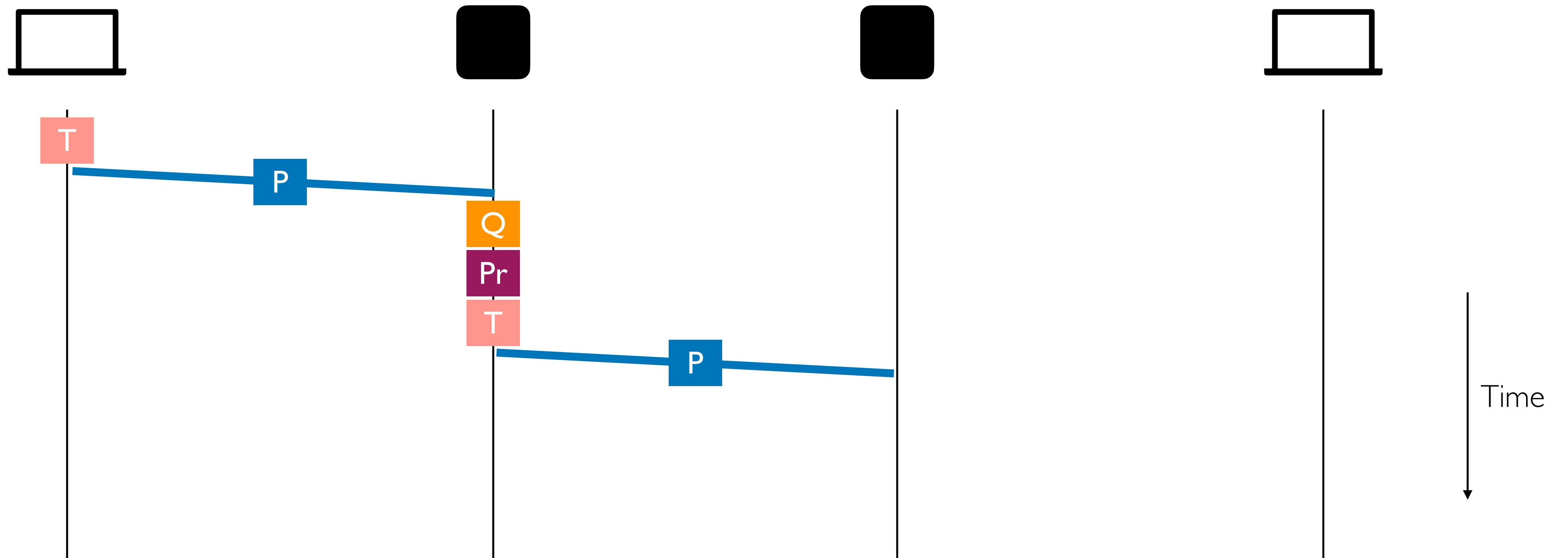
End-to-end Delay

T Transmission Delay

Q Queuing Delay

P Propagation Delay

Pr Processing Delay



End-to-end Delay



Transmission Delay



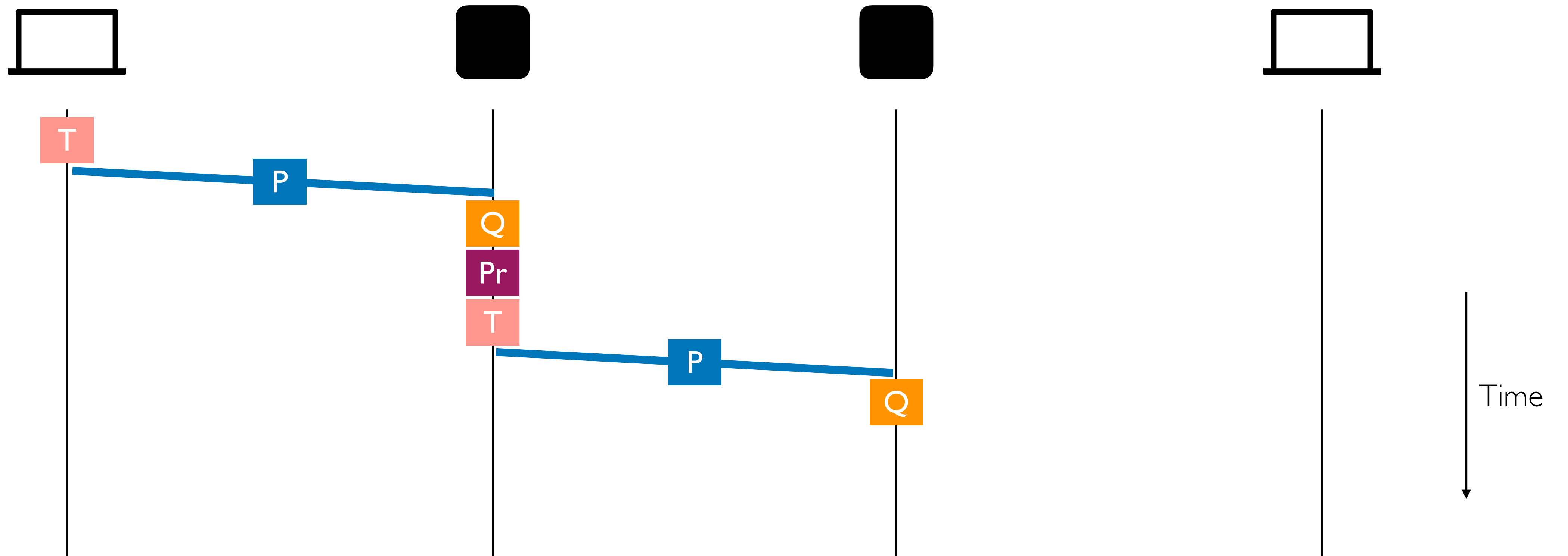
Queuing Delay



Propagation Delay



Processing Delay



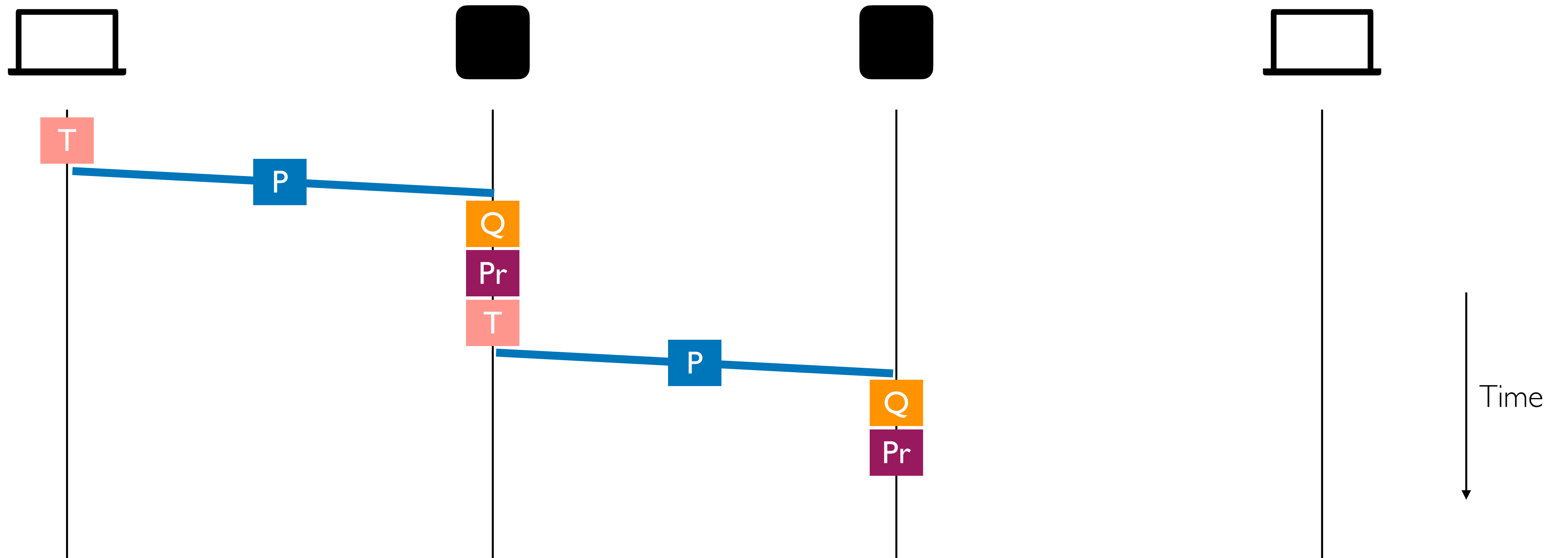
End-to-end Delay

T Transmission Delay

Q Queuing Delay

P Propagation Delay

Pr Processing Delay



End-to-end Delay



Transmission Delay



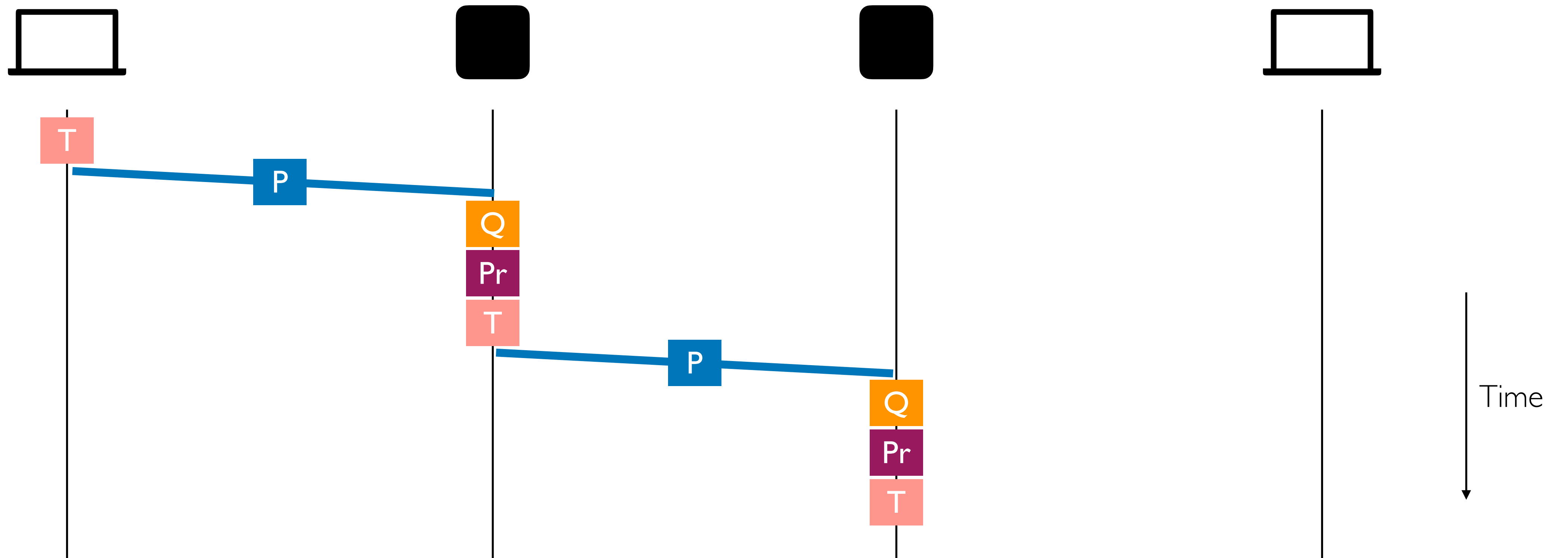
Queuing Delay



Propagation Delay



Processing Delay



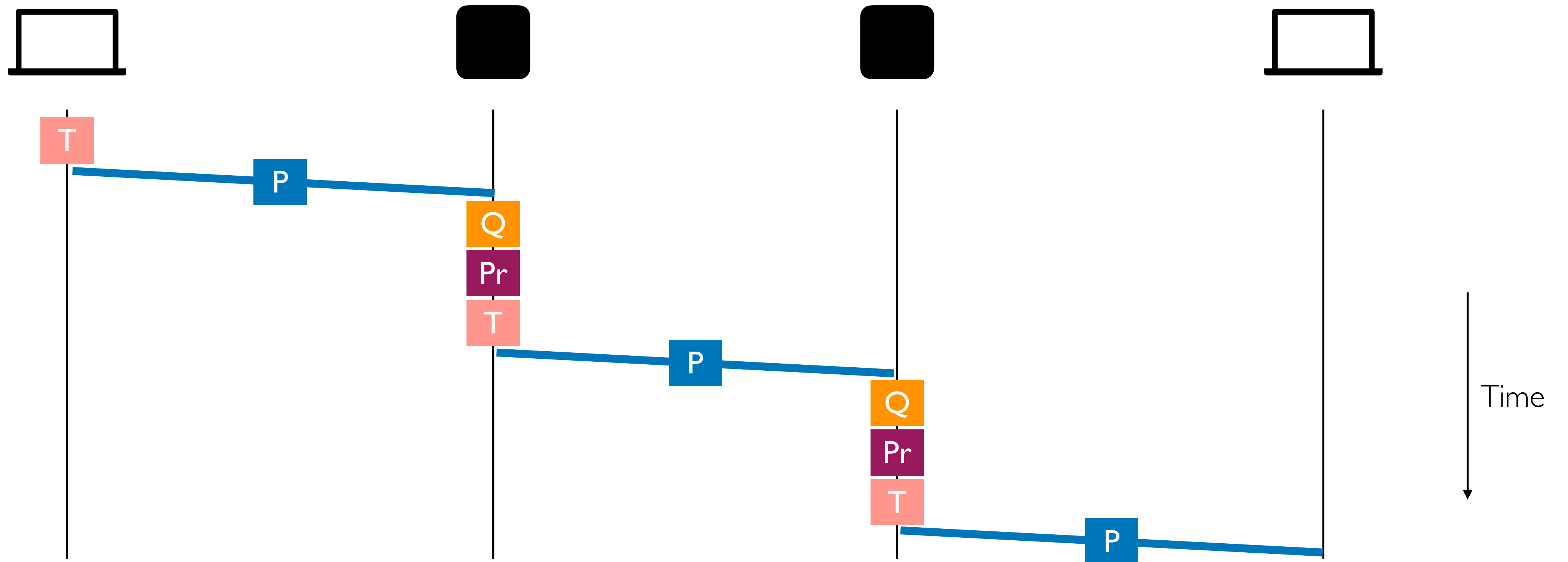
End-to-end Delay

T Transmission Delay

Q Queuing Delay

P Propagation Delay

Pr Processing Delay



Questions?

Loss

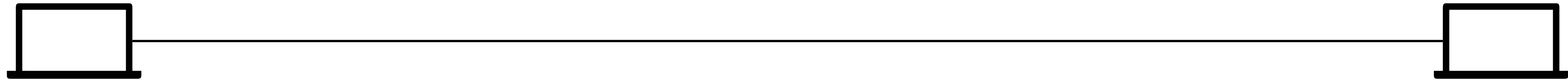
Loss

- What fraction of the packets sent to a destination are dropped?

Throughput

- At what rate is the destination receiving data from the source?
 - **Data size / transfer time**

Throughput

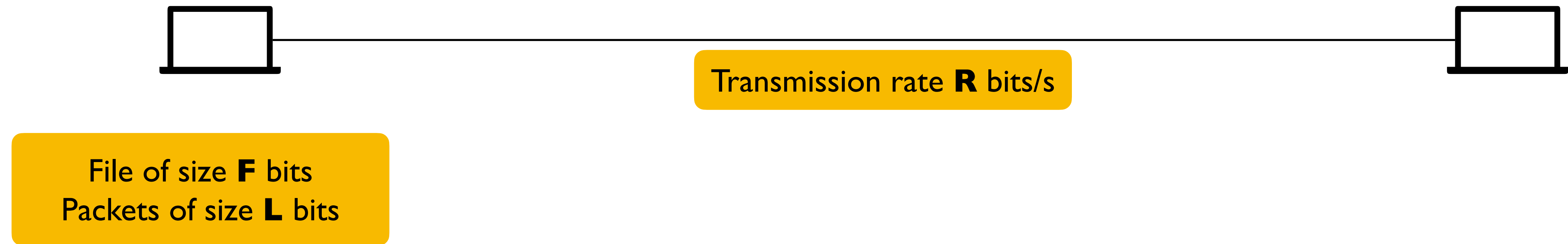


Throughput

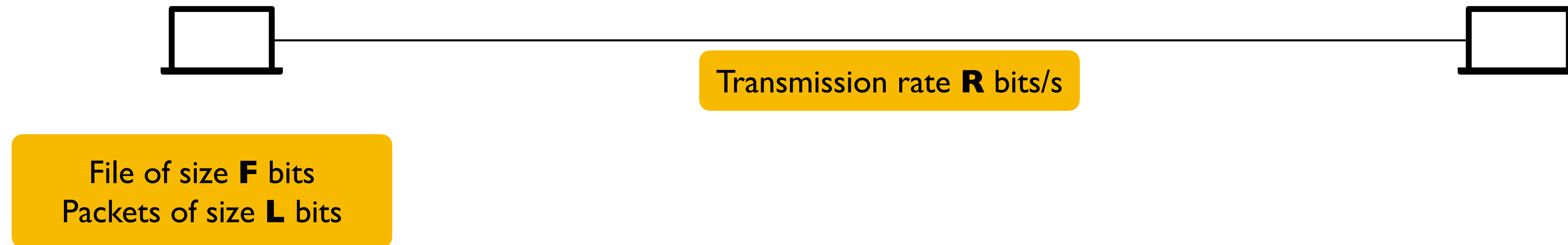


File of size **F** bits
Packets of size **L** bits

Throughput

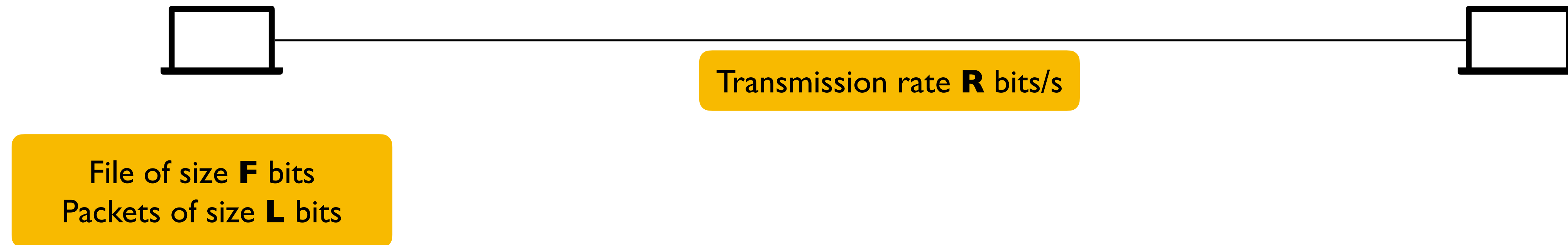


Throughput



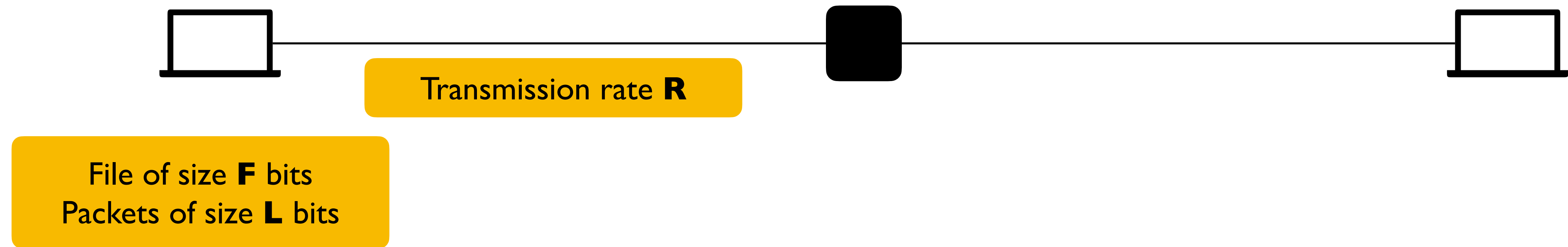
- Transfer time = **F/R** + propagation delay

Throughput

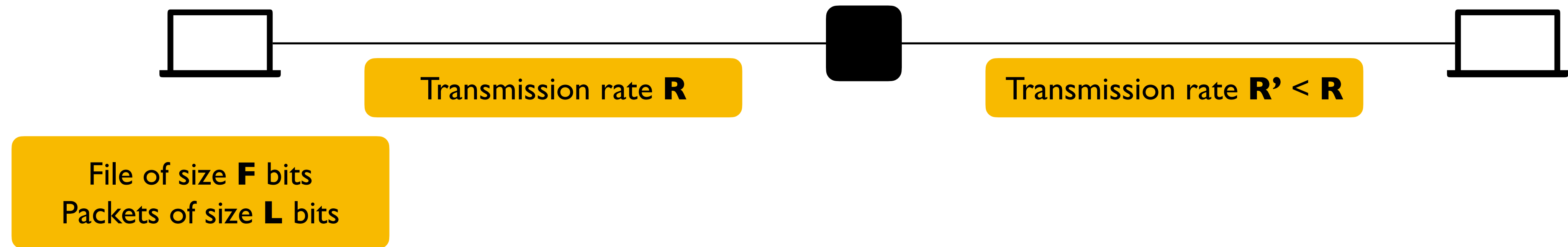


- Transfer time = $\mathbf{F/R}$ + propagation delay
- Average throughput = \mathbf{R}

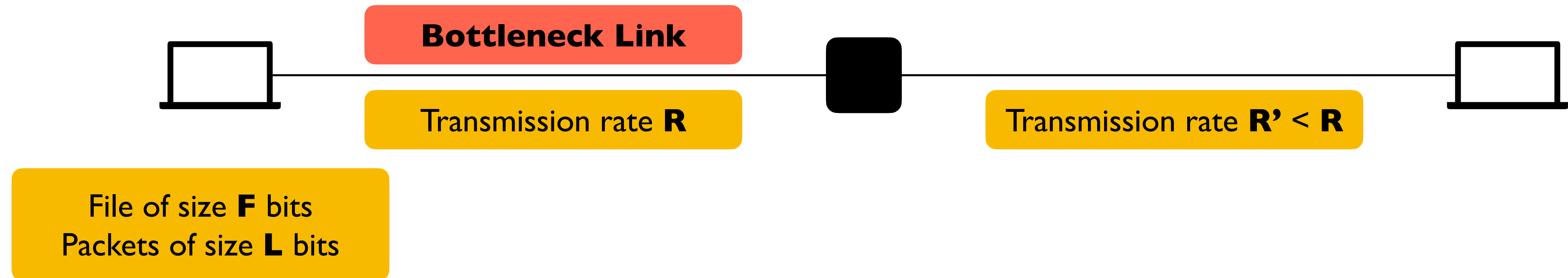
Throughput



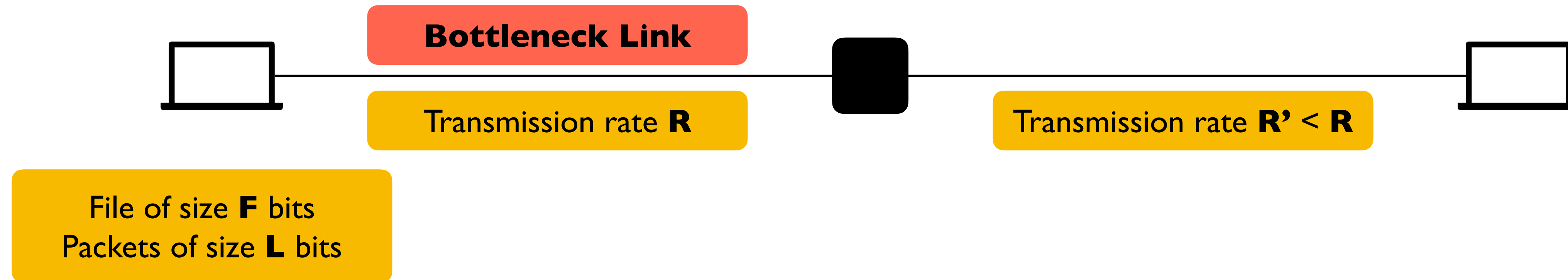
Throughput



Throughput



Throughput

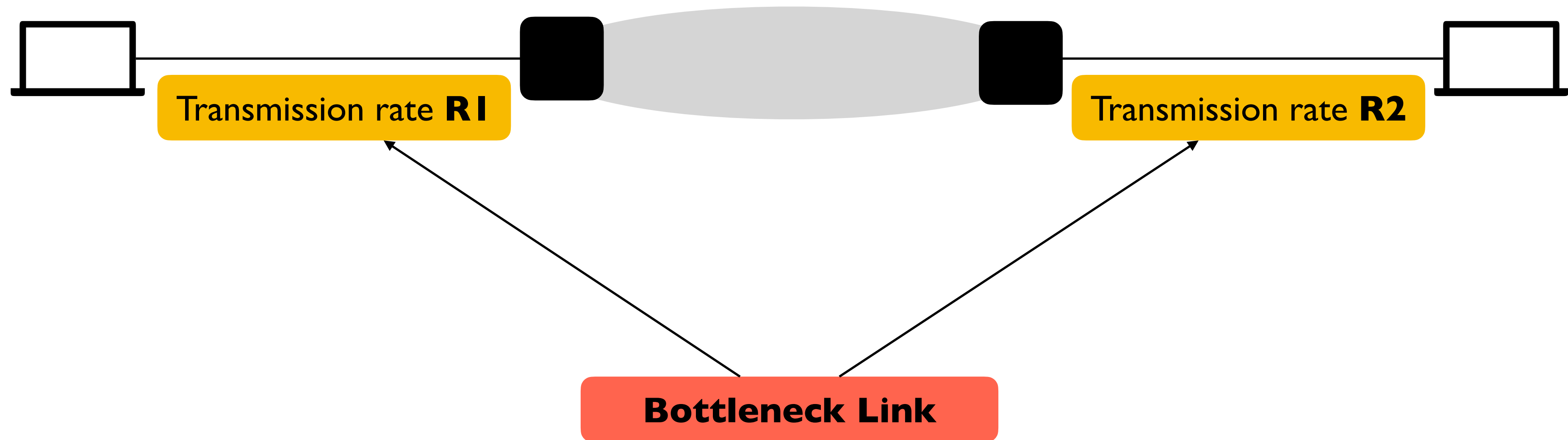


- Average throughput = $\min \{\mathbf{R}, \mathbf{R}'\} = \mathbf{R}$

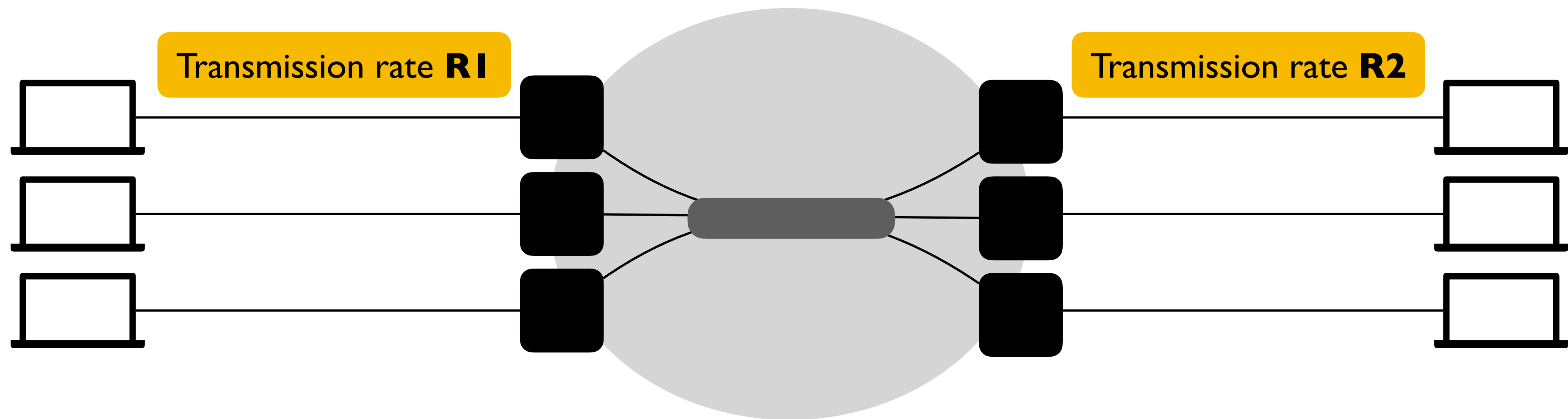
Throughput



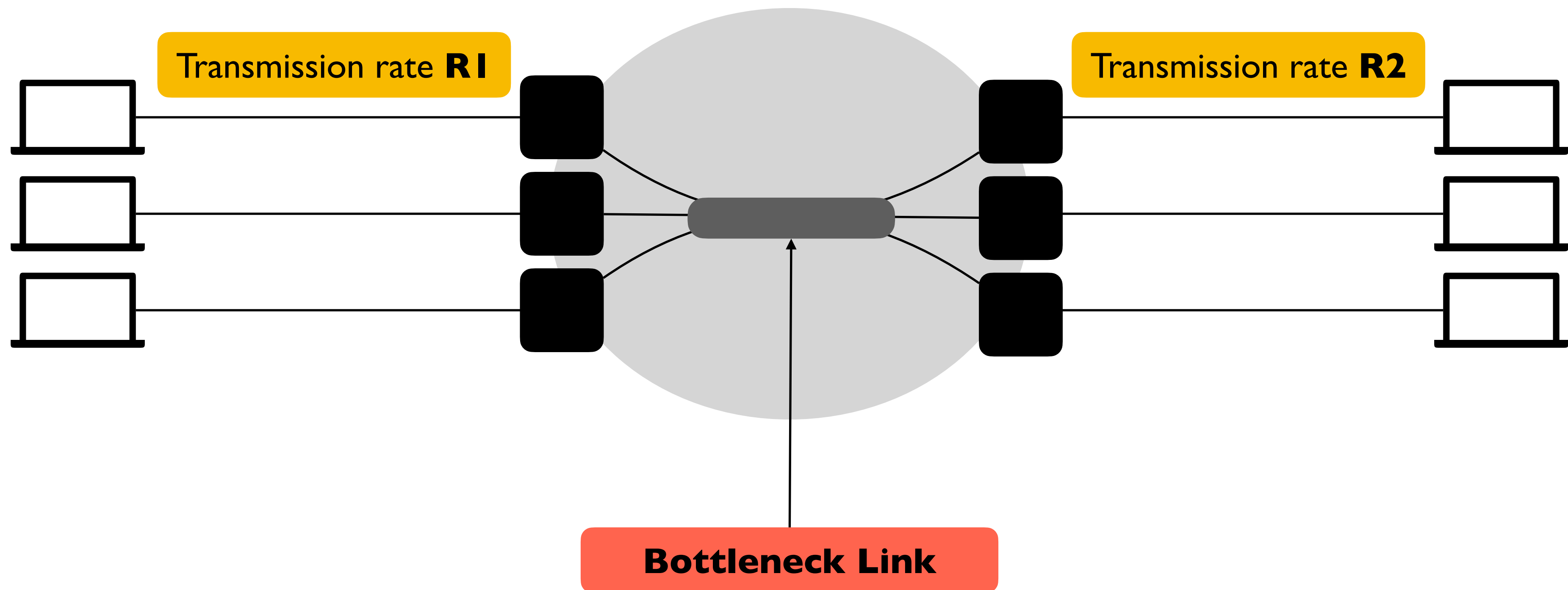
Throughput



Throughput



Throughput



Questions?

Taking Stock

Taking Stock

- What & how of the Internet
 - A slightly detailed look at physical infrastructure

Taking Stock

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- Sharing networks
 - Reserve or on-demand?

Taking Stock

- **What & how of the Internet**
 - A slightly detailed look at physical infrastructure
- **Sharing networks**
 - Reserve or on-demand?
- **Evaluating networks**
 - Delay (transmission, propagation, queueing, processing), loss and throughput