

Domain Name System (DNS) (Contd.)

CPSC 433/533, Spring 2021
Anurag Khandelwal

Evaluating Midterm Performance

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- **Observation: 45+ scorers attend class regularly, and turn up to office hours!**
 - Correlation is not causation, but it does make you wonder...
 - Attend class (encourage your friends to do the same!)
 - Participate (ask questions, post on Piazza, show up for OH, ...)

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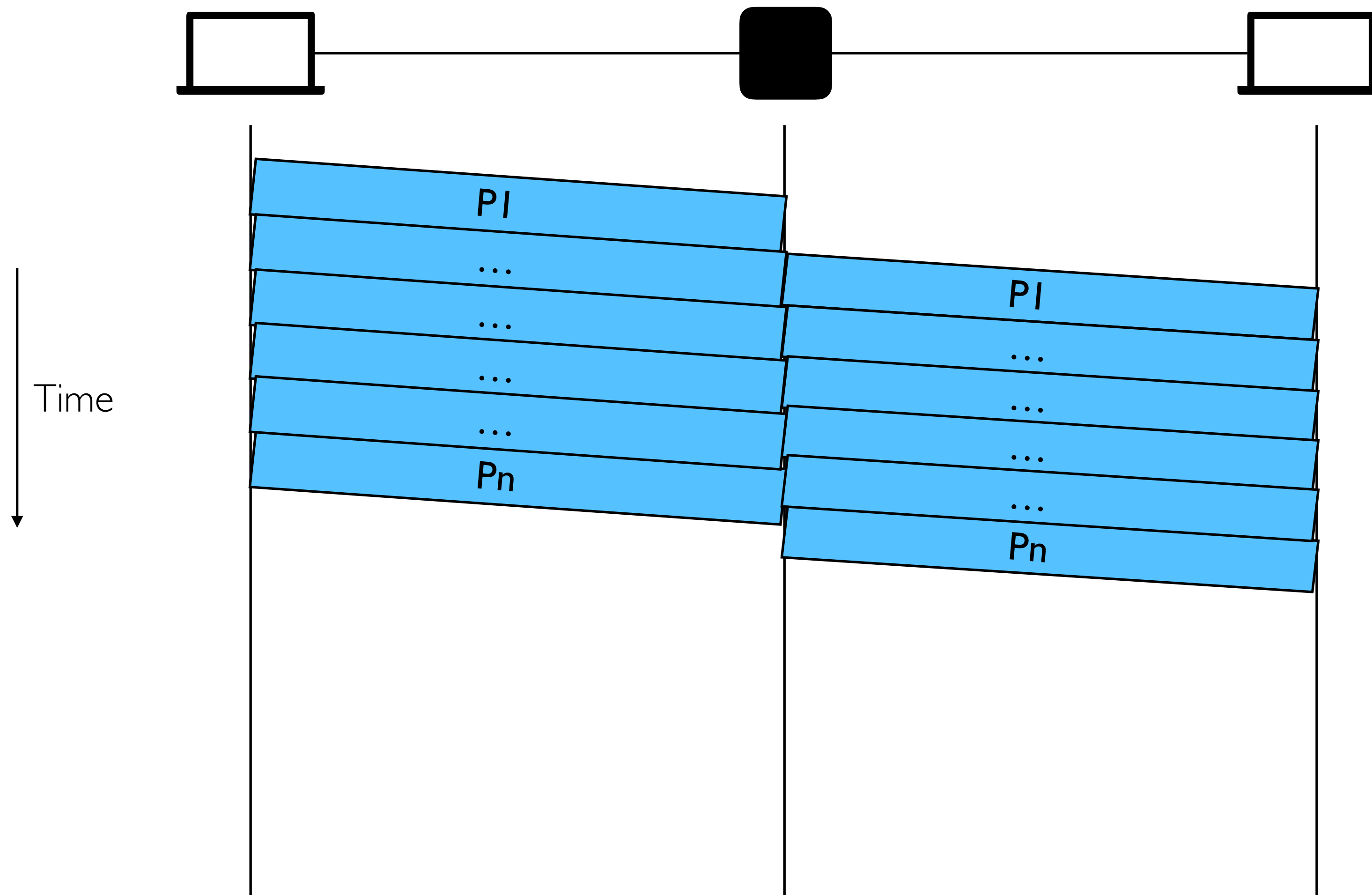
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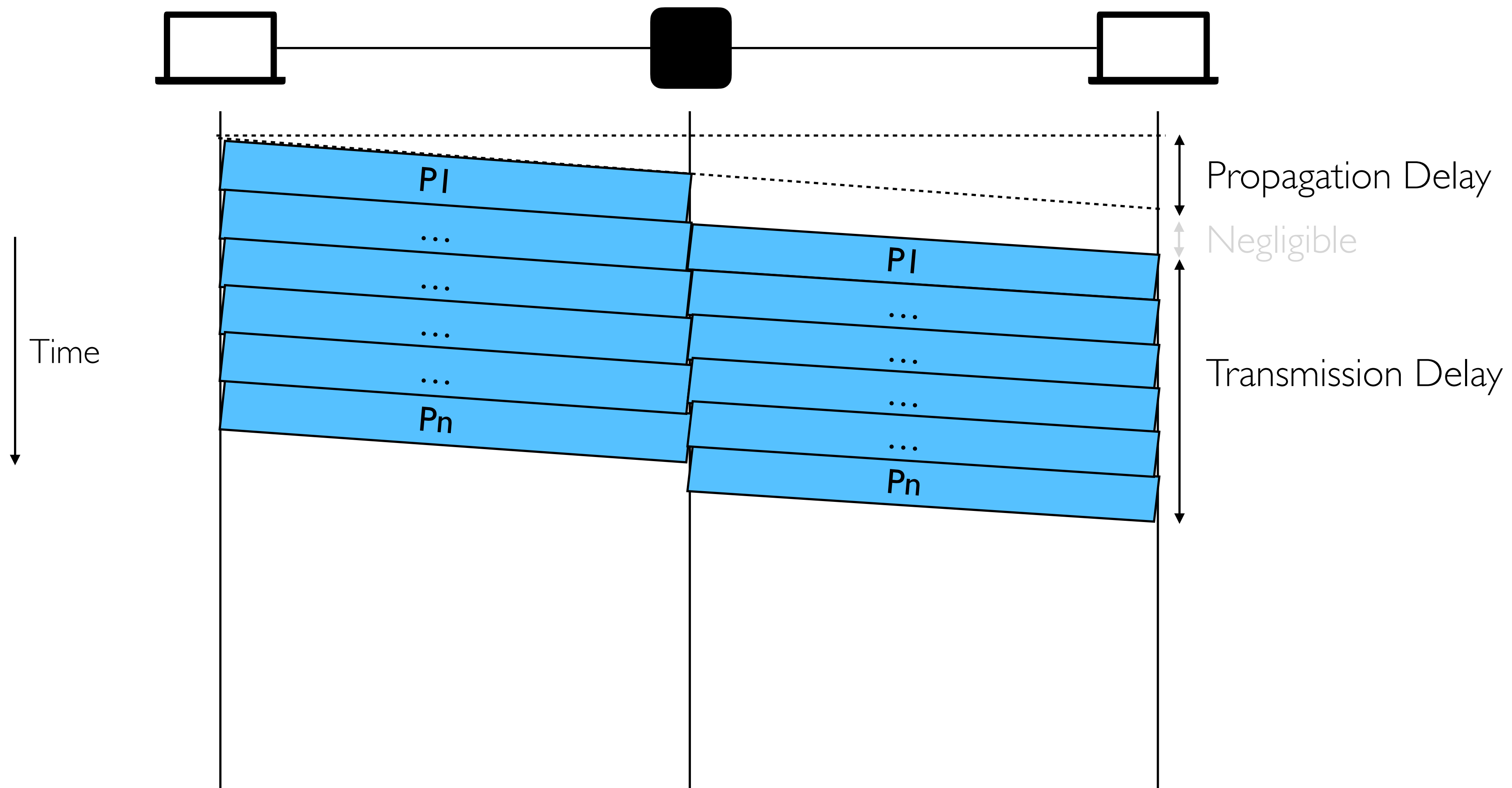
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 - We will award everyone 1 point for the question

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 - First, calibrate: the absolute scores don't matter. This was a hard exam, adjust your expectations based on class average.
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- **Please fill out the mid-semester survey!**
 - Survey link: <https://www.surveymonkey.com/r/XGYCGXD>
 - Will help us understand where you are struggling, and what we can do to fix it!
 - If you don't speak up, nothing will change...
 - ... but if you do, I promise I will do everything I can to help you learn!

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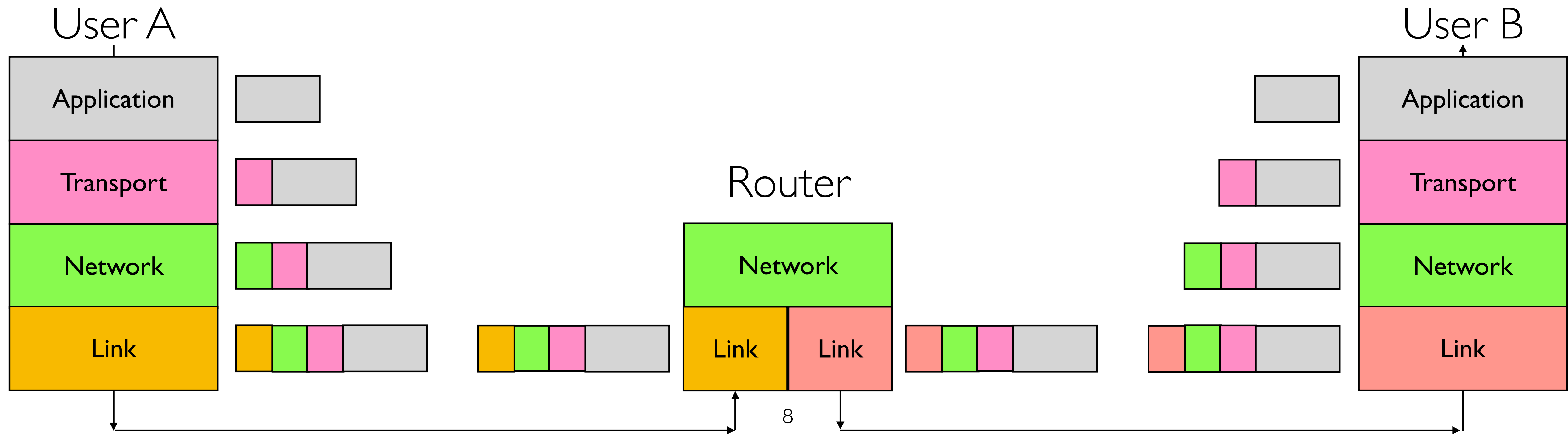
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 - 😊... will post pdfs both with & without animations

Administriva: Project 2

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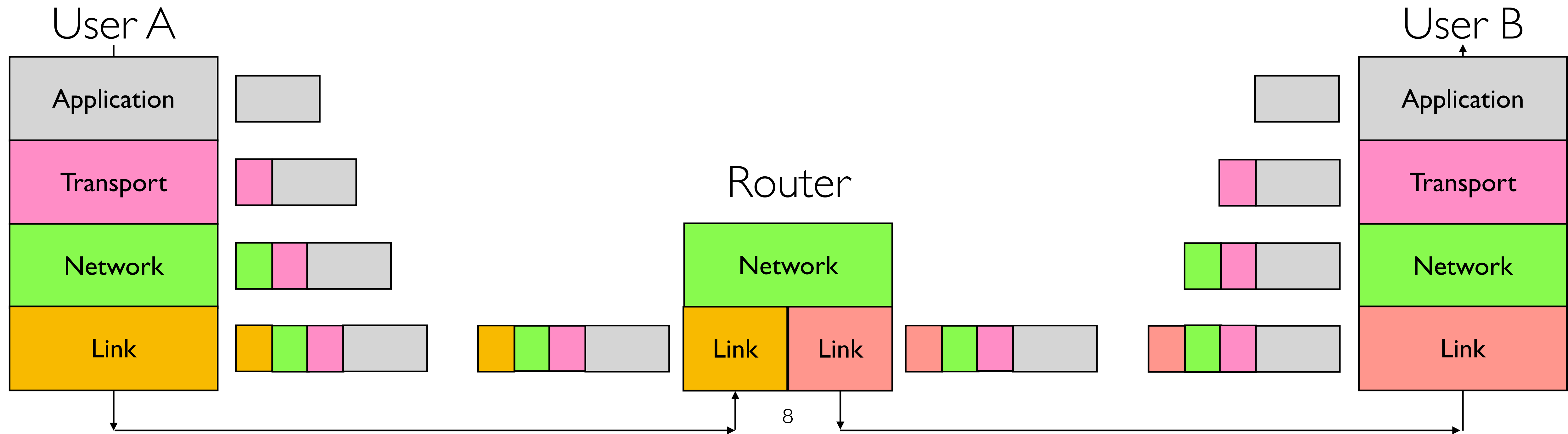
- Project 2 will be released tonight, due in 3 weeks
 - Will let Jonathan tell you more about it...

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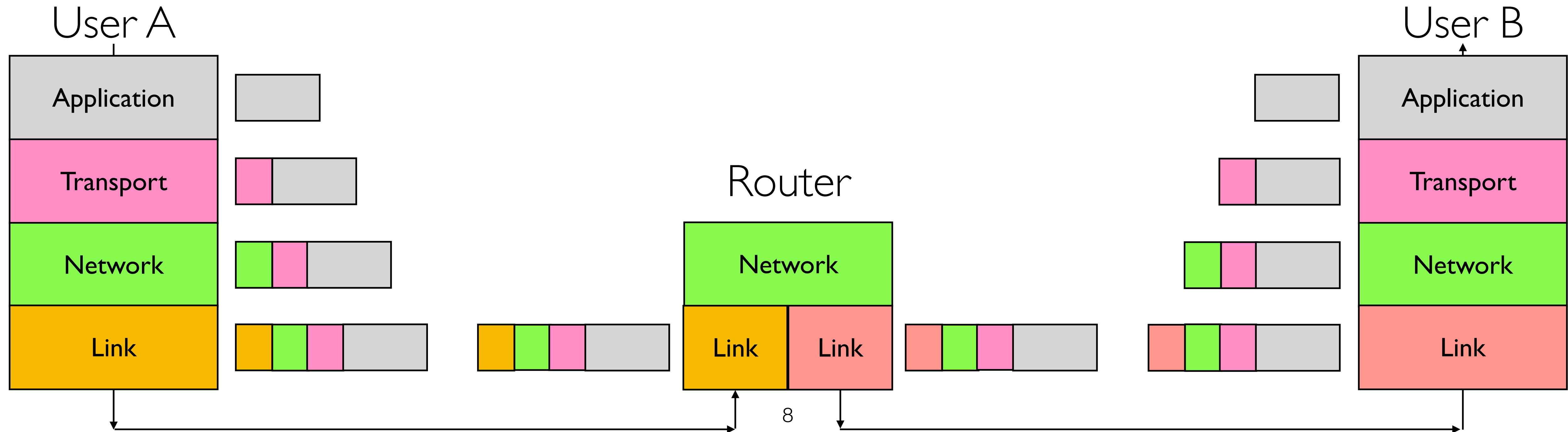
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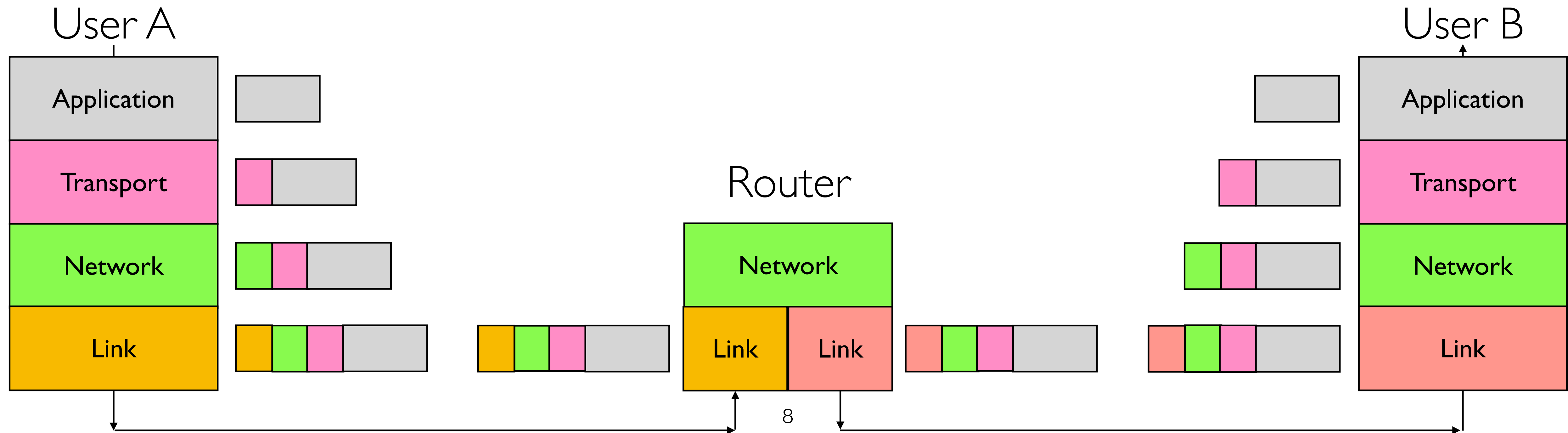
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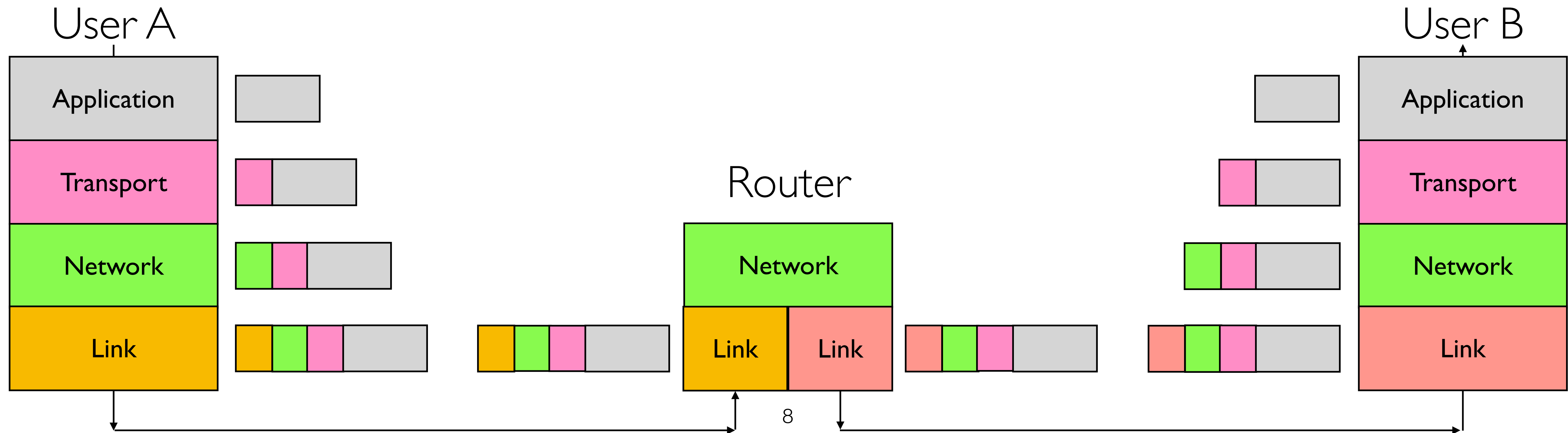
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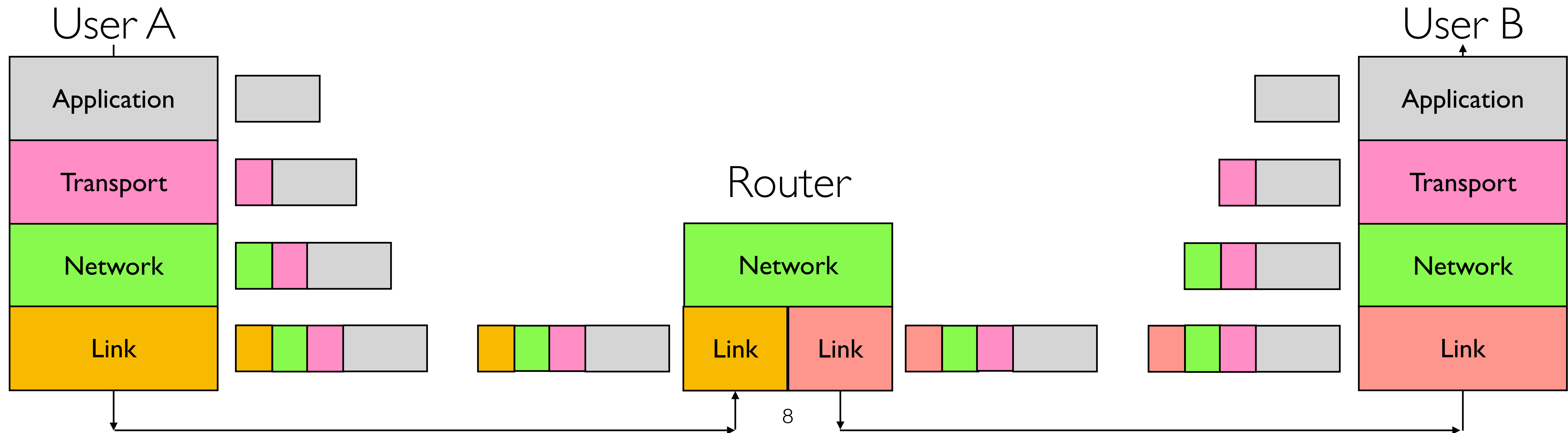
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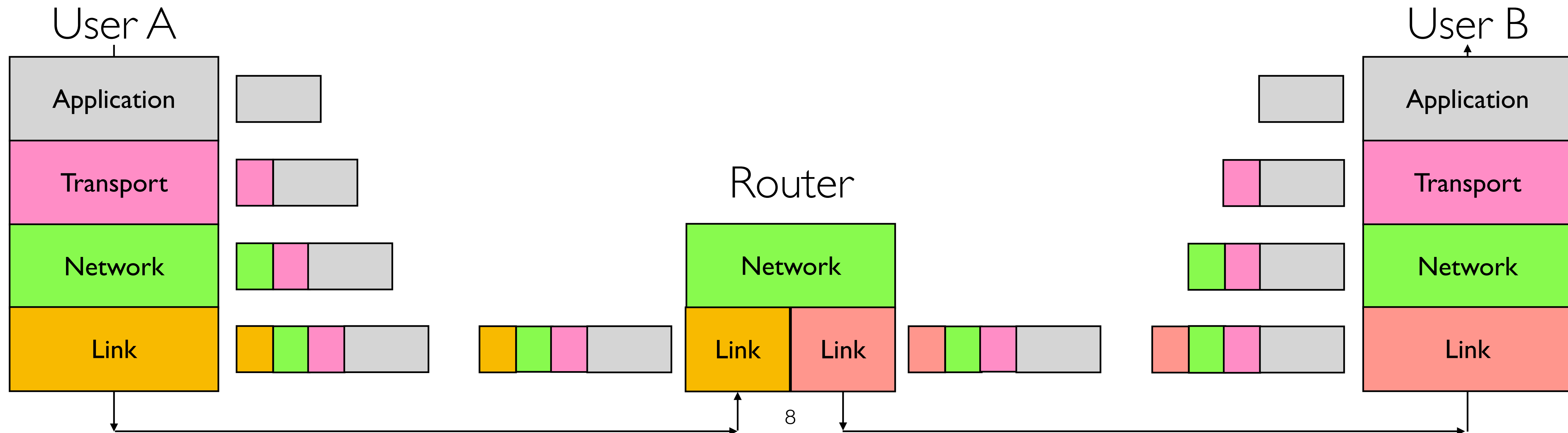


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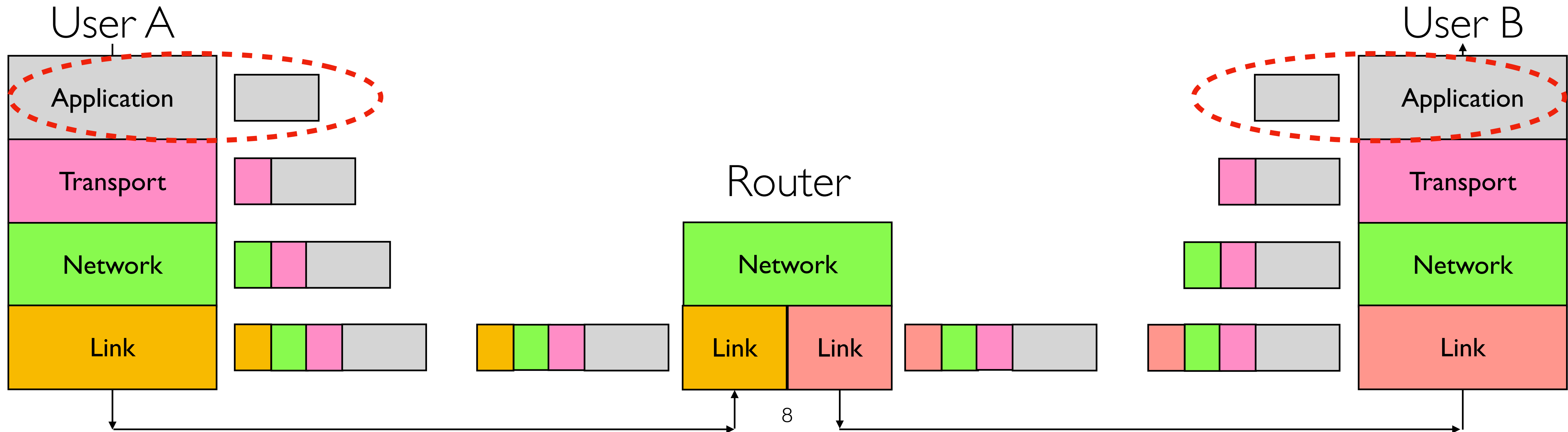
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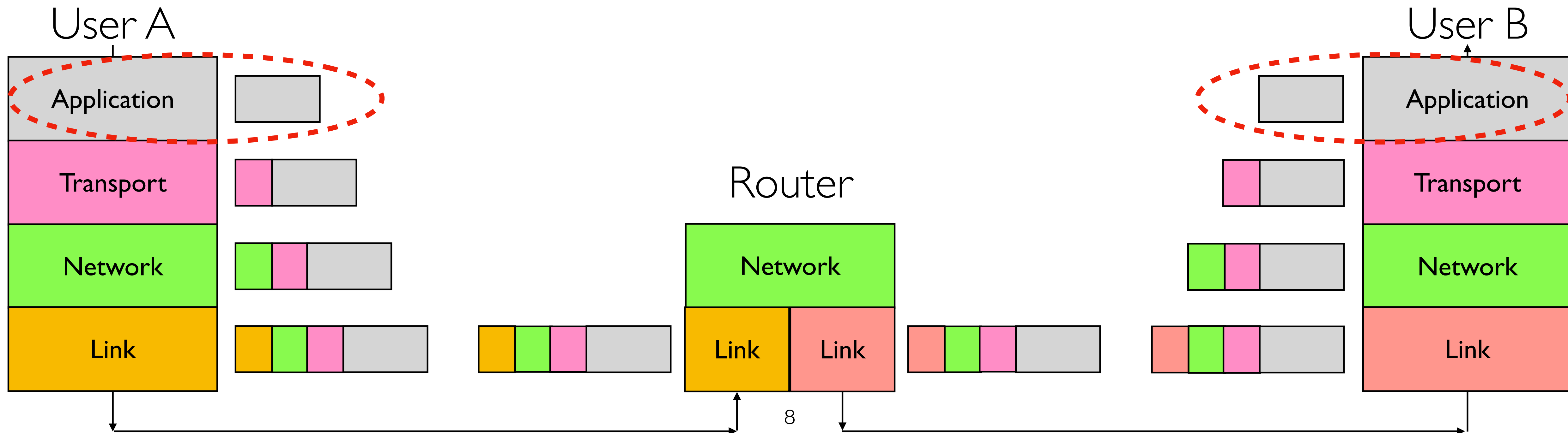
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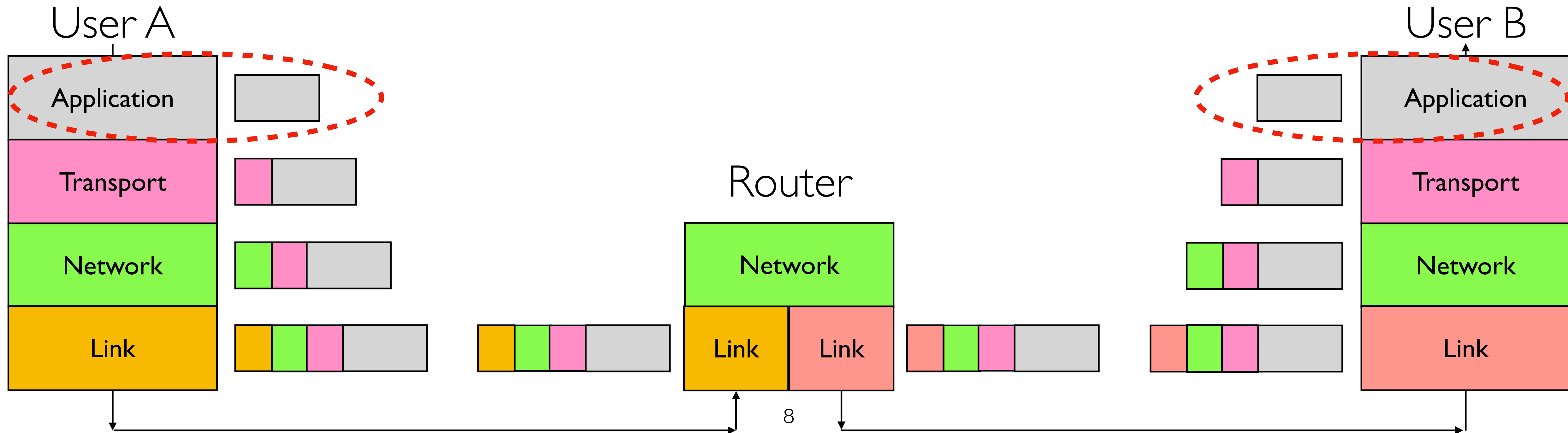
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- **Advanced Topics**, *Datacenters, SDN*



Back to DNS!

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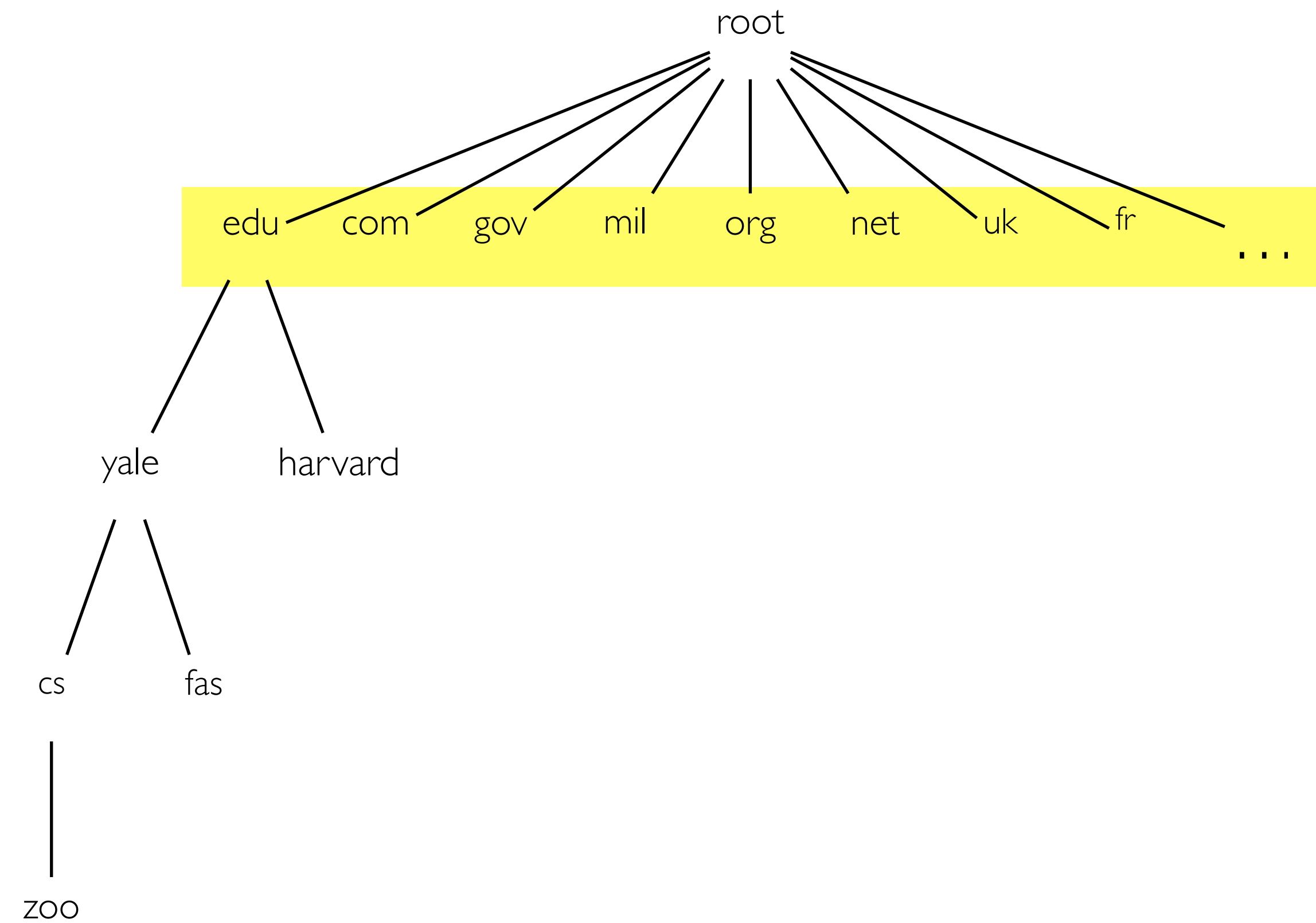
- Why bother?
- Convenience
 - Easier to remember www.google.com than 172.217.8.174
- Provides a level of indirection!
 - Decoupled names from addresses
 - Many uses beyond just naming a specific host

Recap: DNS Goals

- **Scalable**
 - *Many names*
 - *Many updates*
 - *Many users creating names*
 - *Many users looking up names*
- **Highly available**
- **Correct**
 - No naming conflicts (uniqueness)
 - Consistency → observe the latest update
- **Lookups are fast**

Recap: Scaling with Hierarchical Distribution

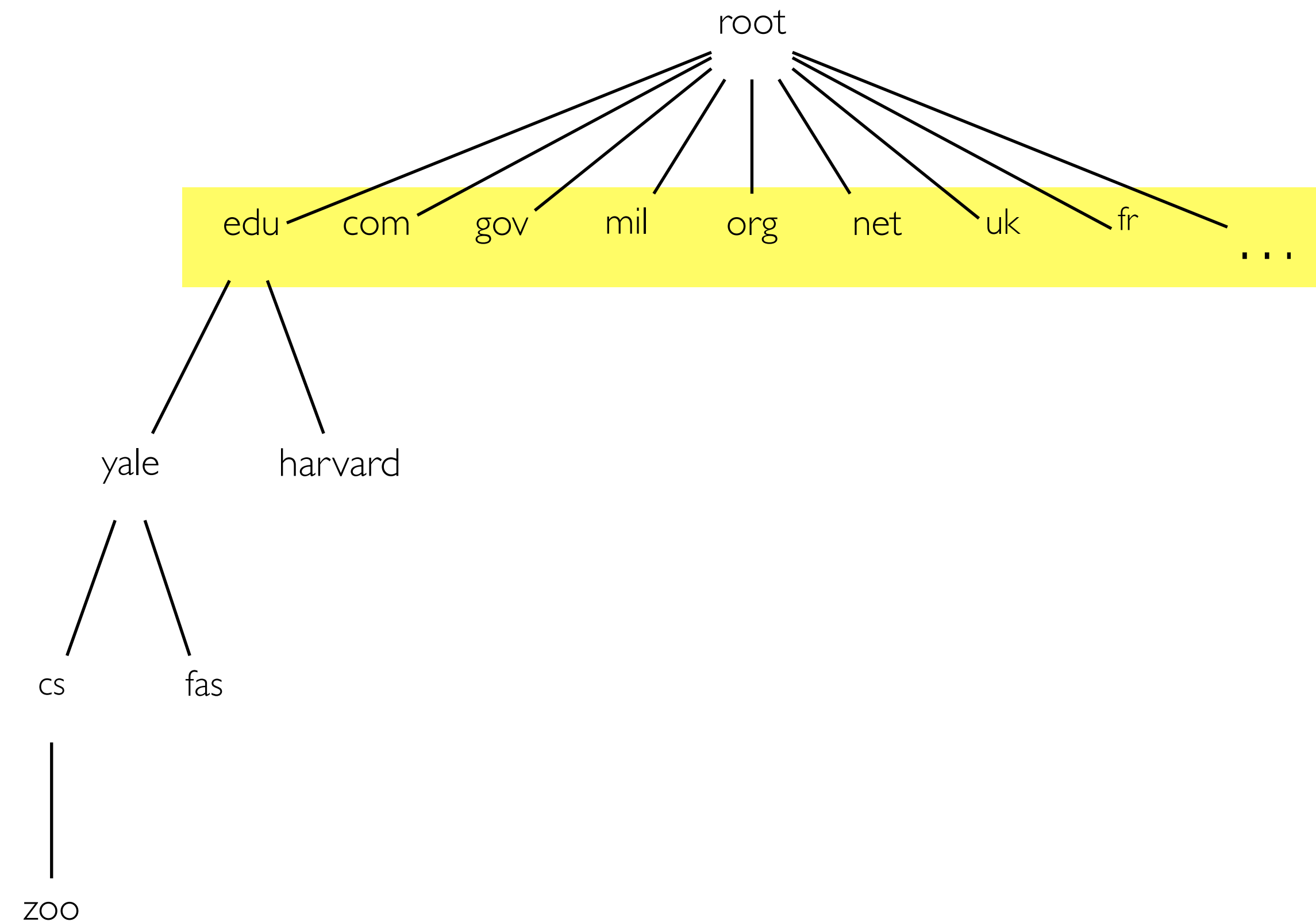
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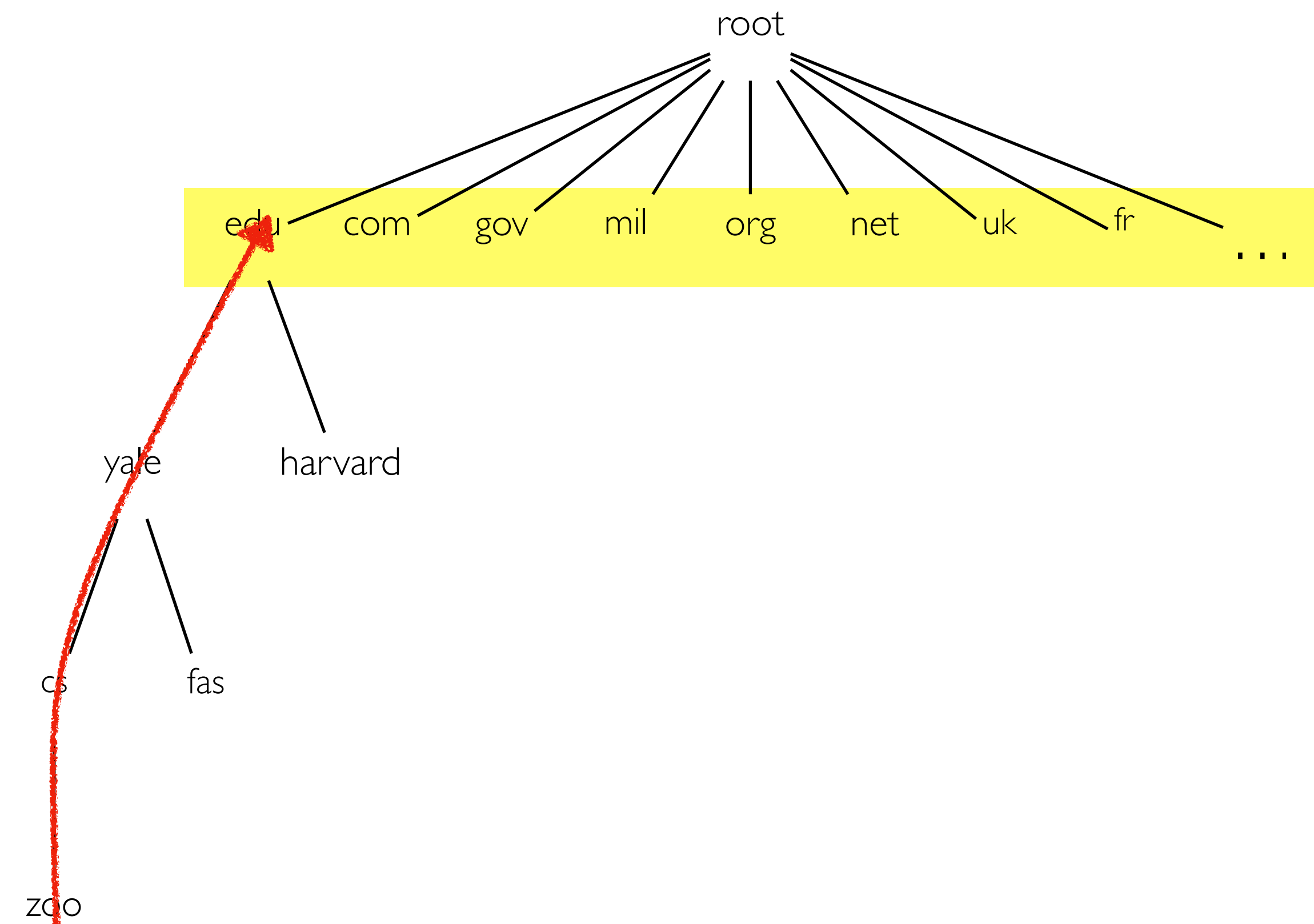
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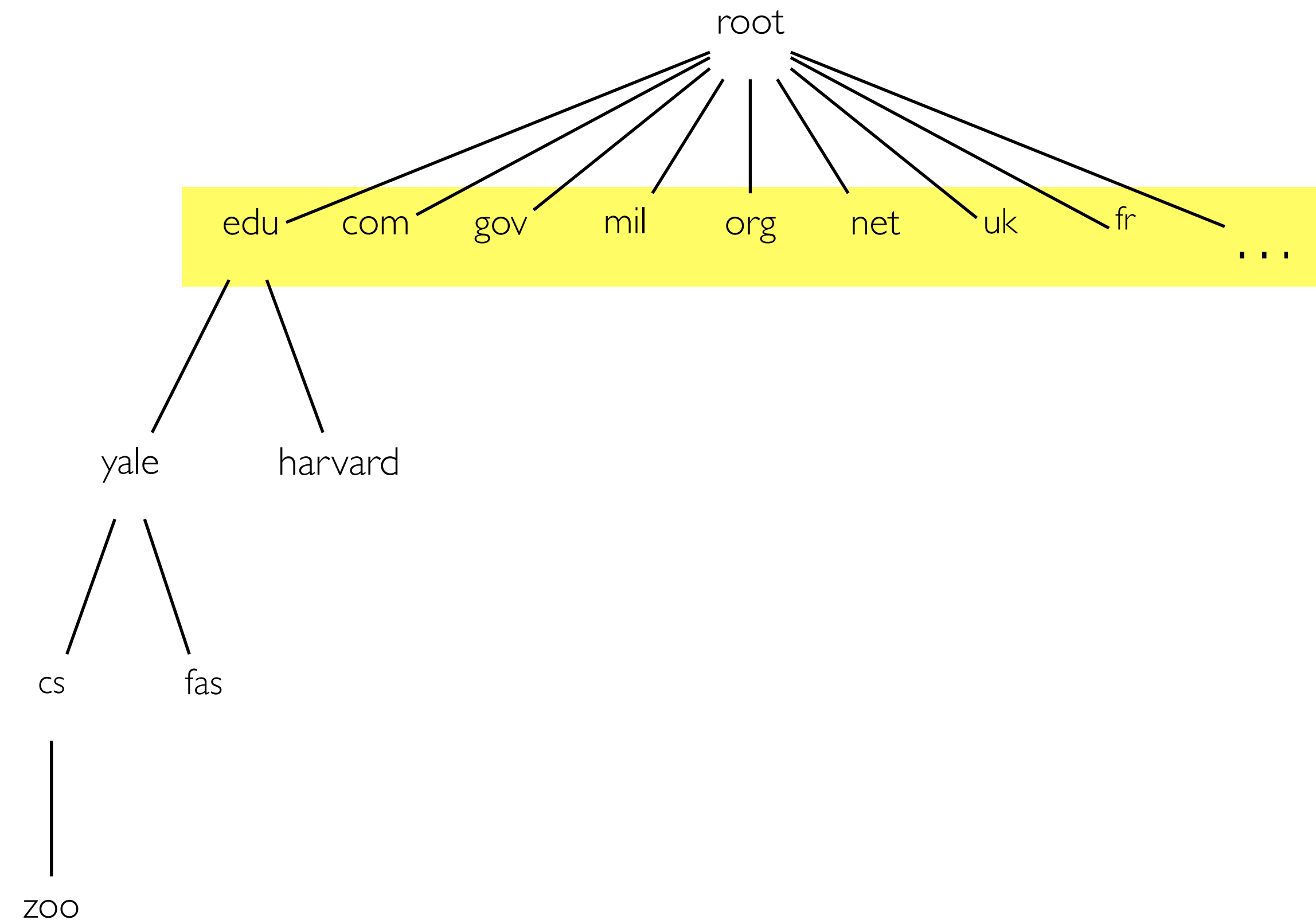
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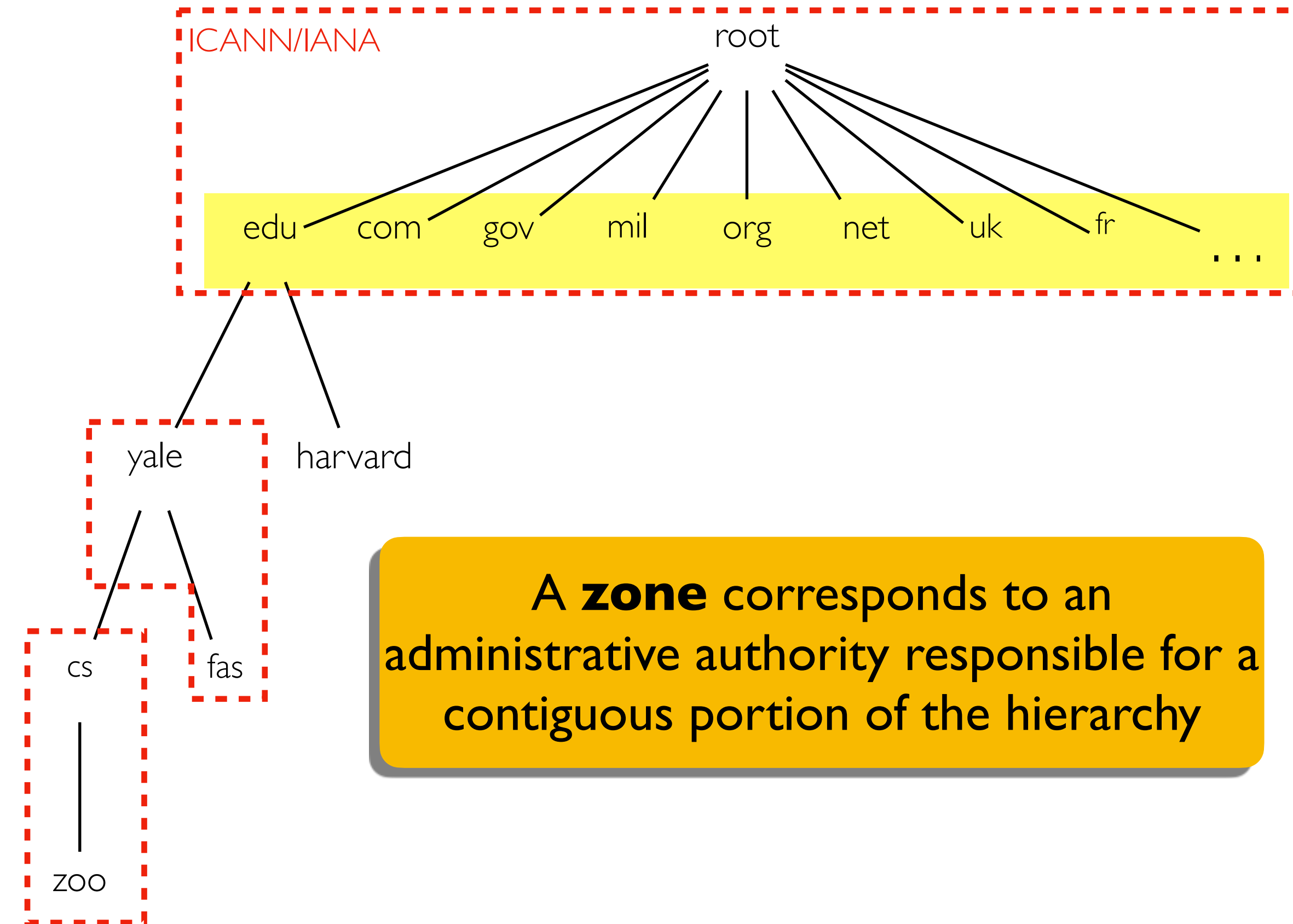
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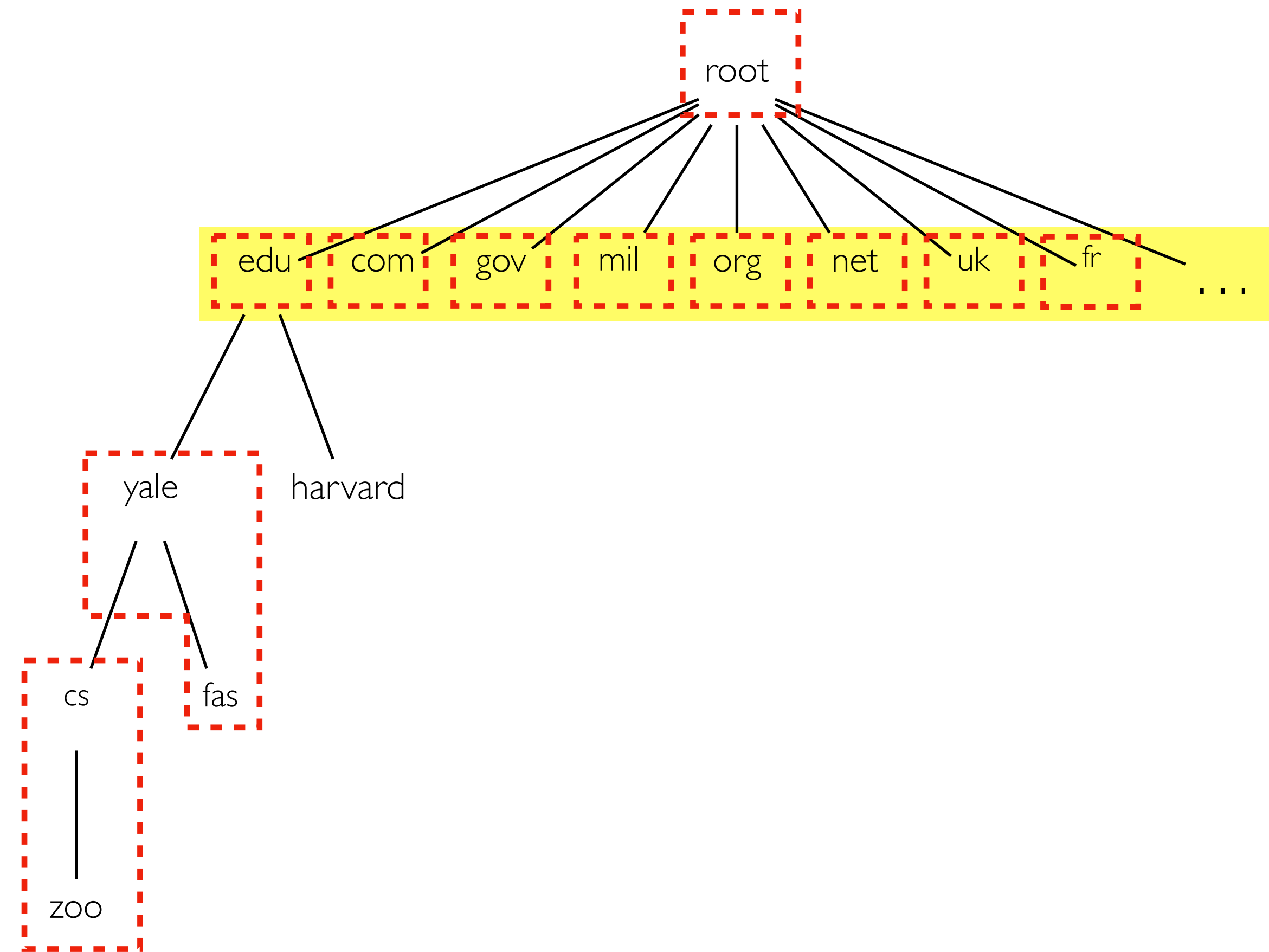
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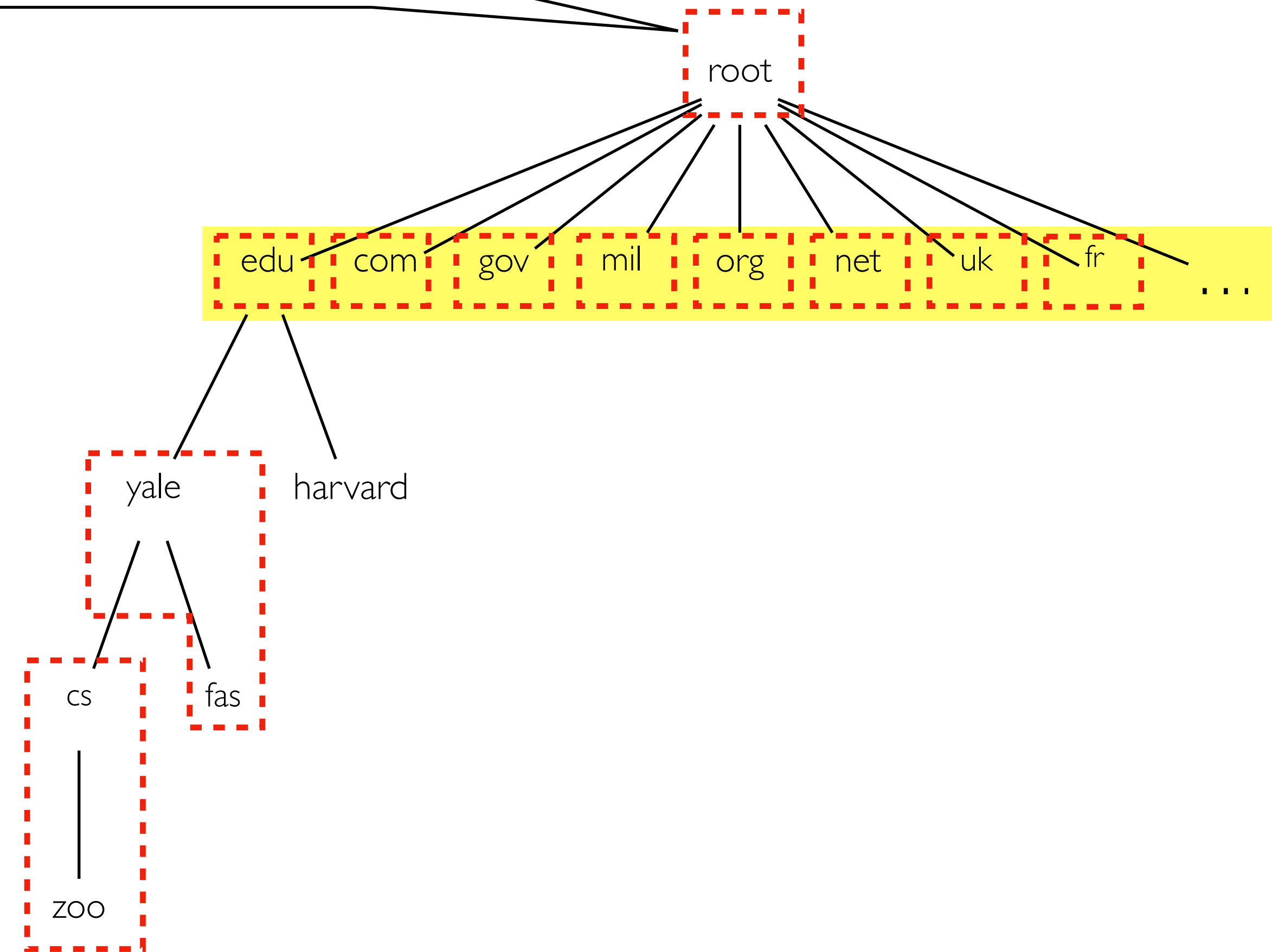


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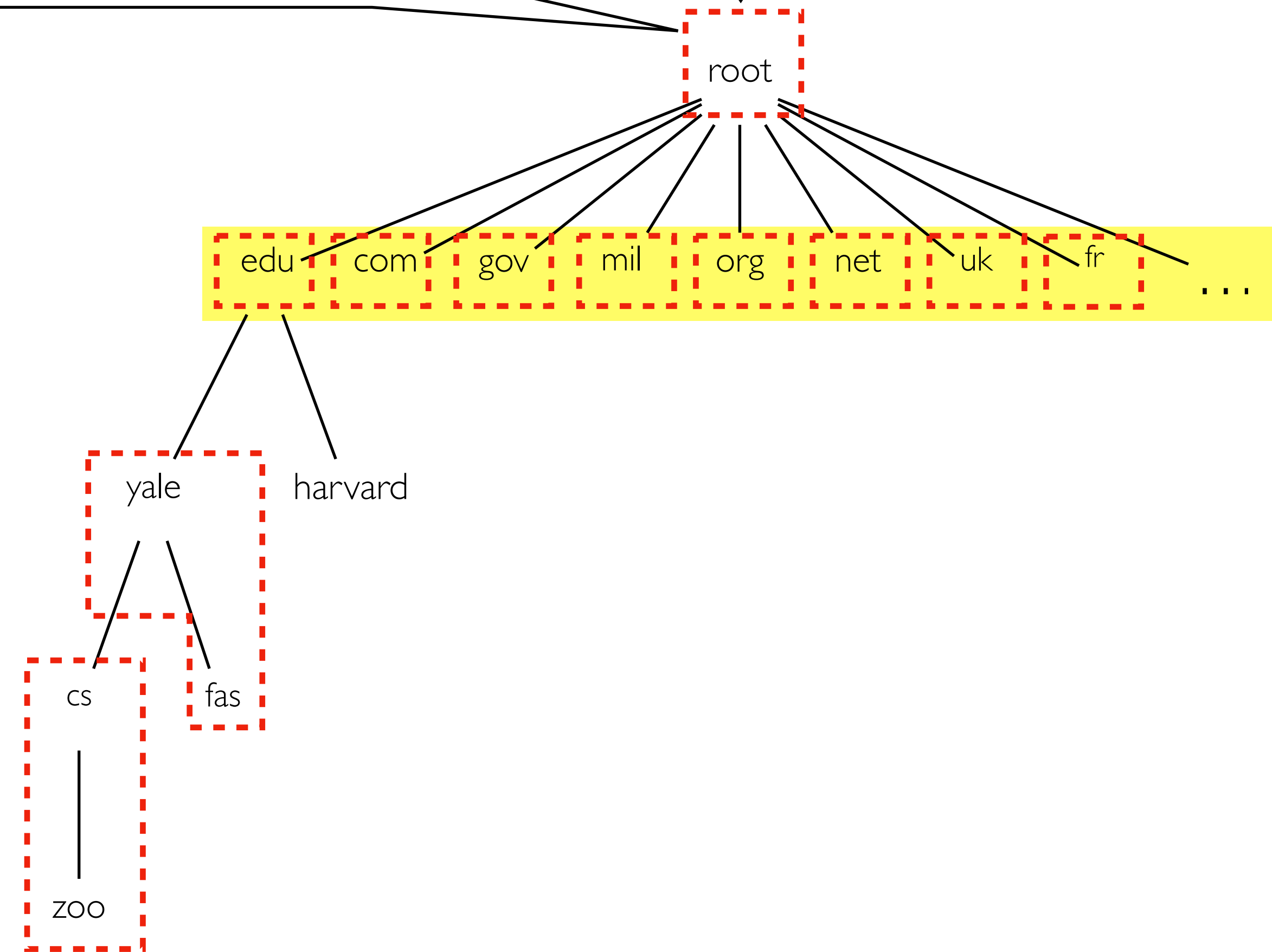
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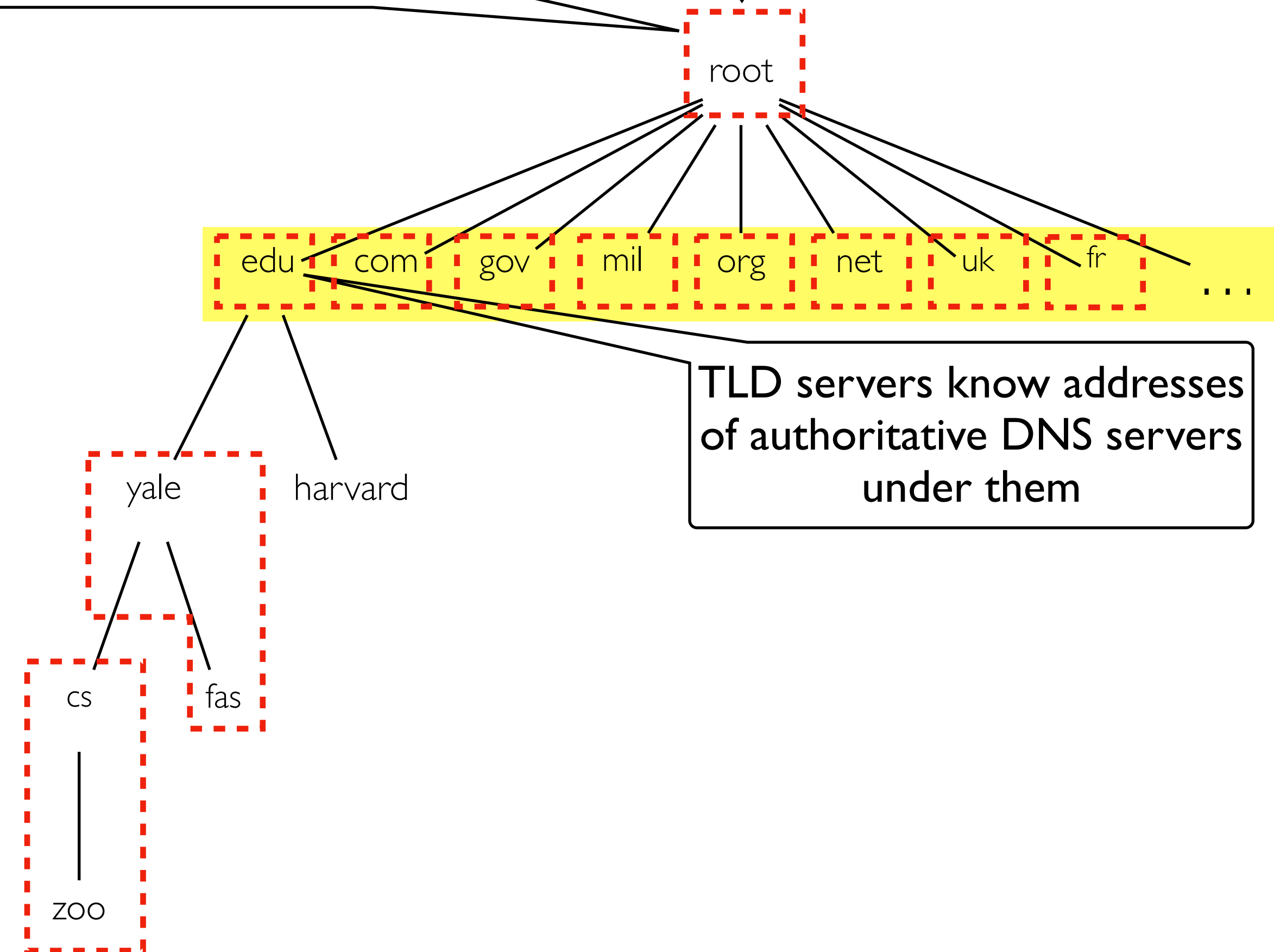
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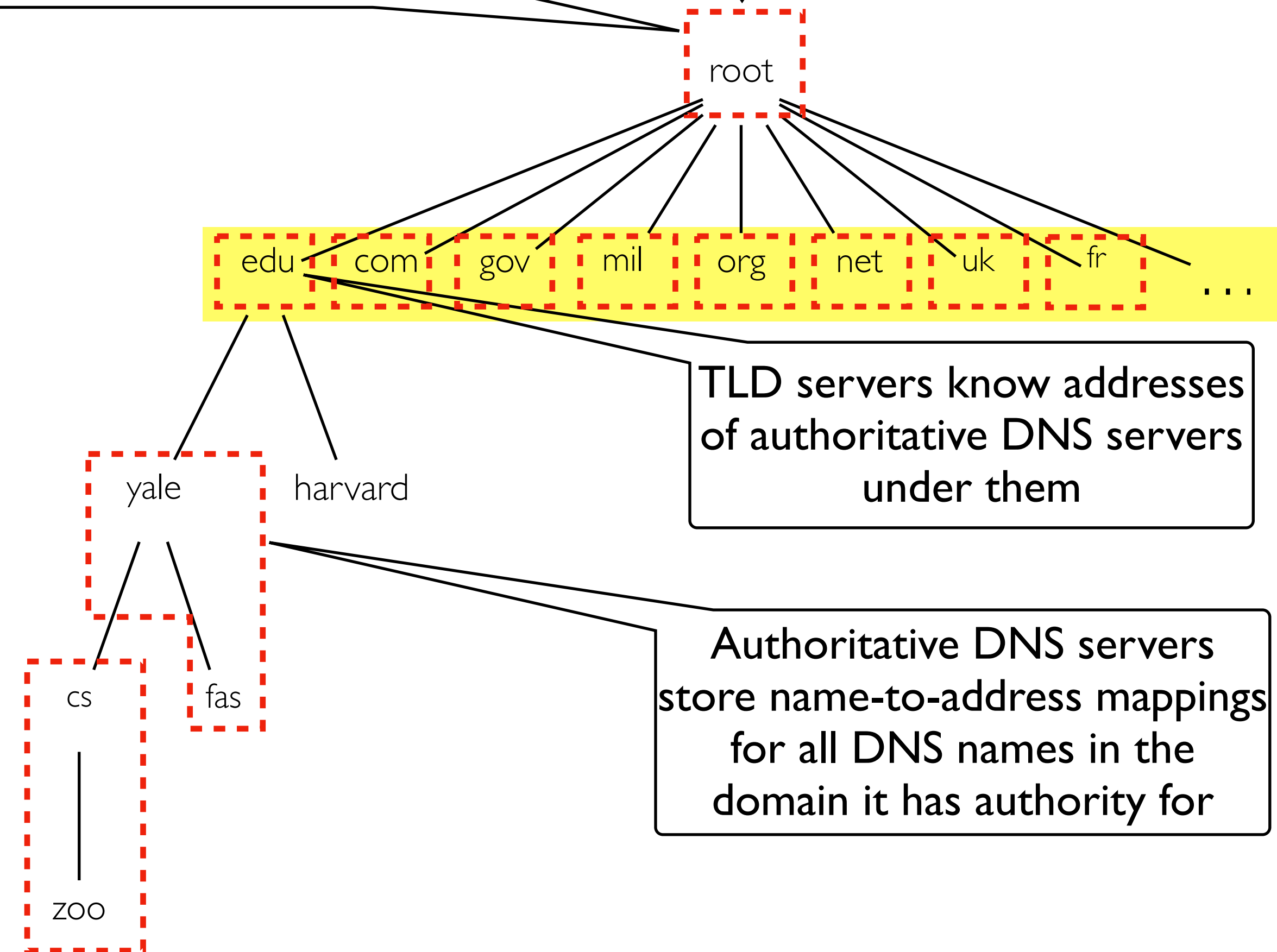
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- Name = domain
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- Type = MX: (\rightarrow Mail eXchanger)

- Name = domain in email address
- Value = name(s) of mail server(s)

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- Register foobar.com at registrar (e.g., GoDaddy)
 - Provide registrar with names and IP addresses of your authoritative name server(s)
 - Registrar inserts RR pairs into the .com TLD server
 - (foobar.com, dns1.foobar.com, NS)
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- Store resource records in your server dns1.foobar.com
 - e.g., type A records: (foobar.com, 212.44.9.130, A), (social.foobar.com, 212.44.9.131, A), etc.
 - e.g., type MX records for foobar.com

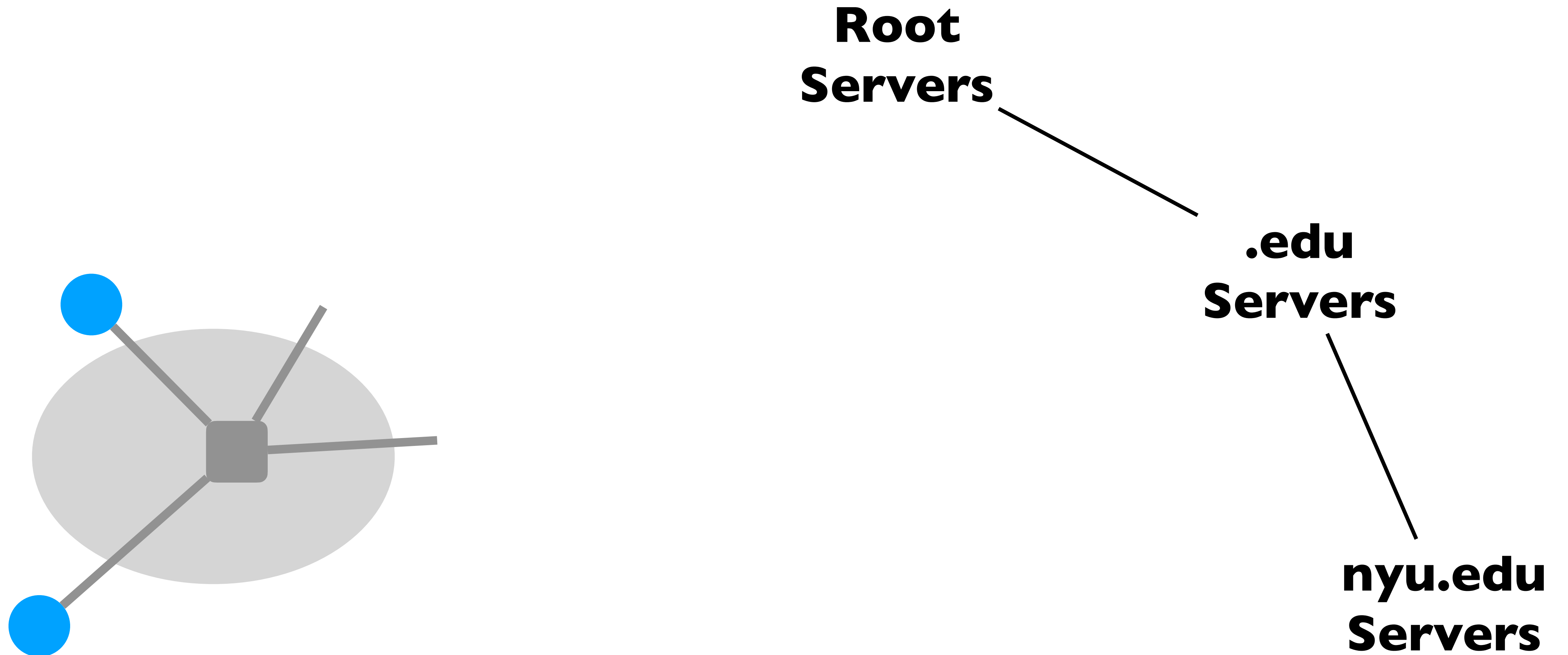
Recap: Using DNS (Client/Application view)

**Root
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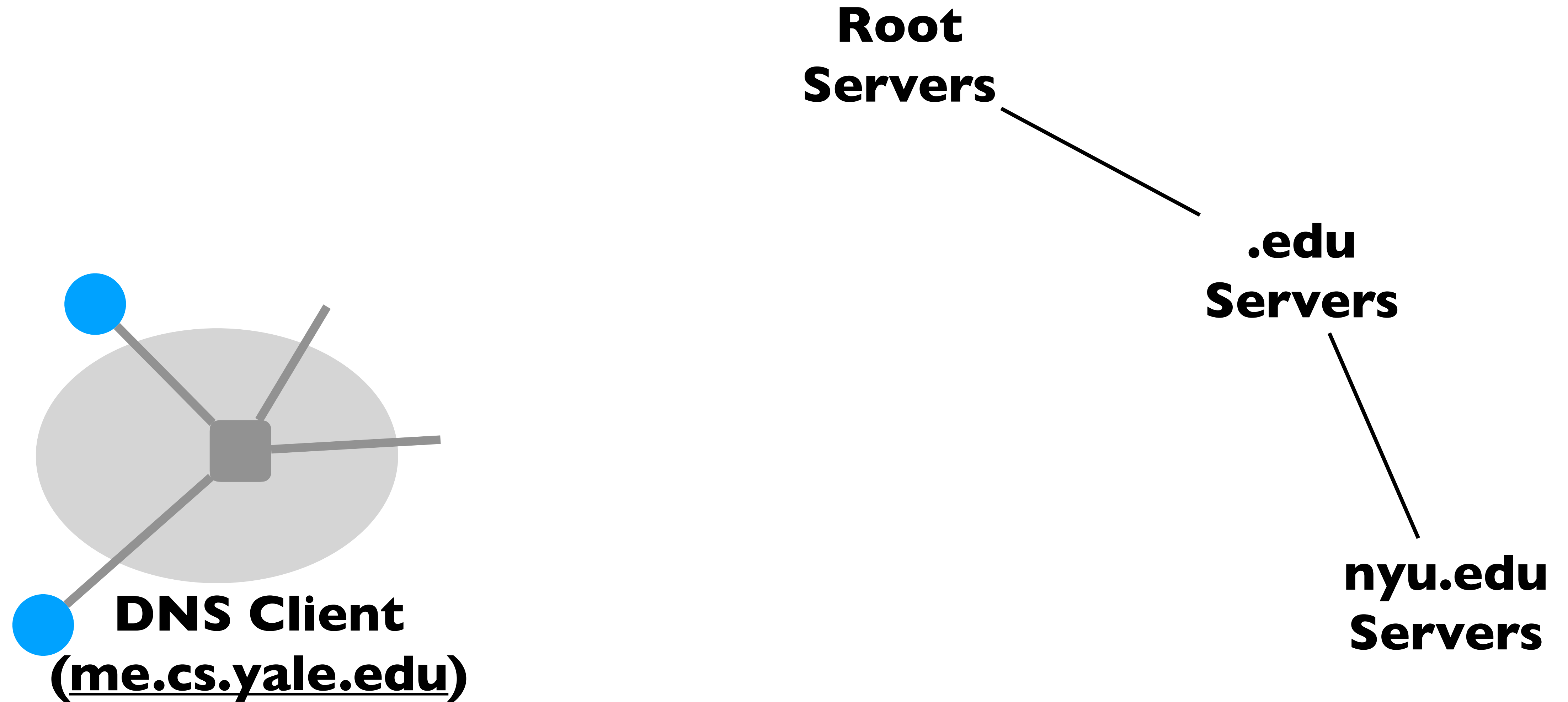
**.edu
Servers**

**nyu.edu
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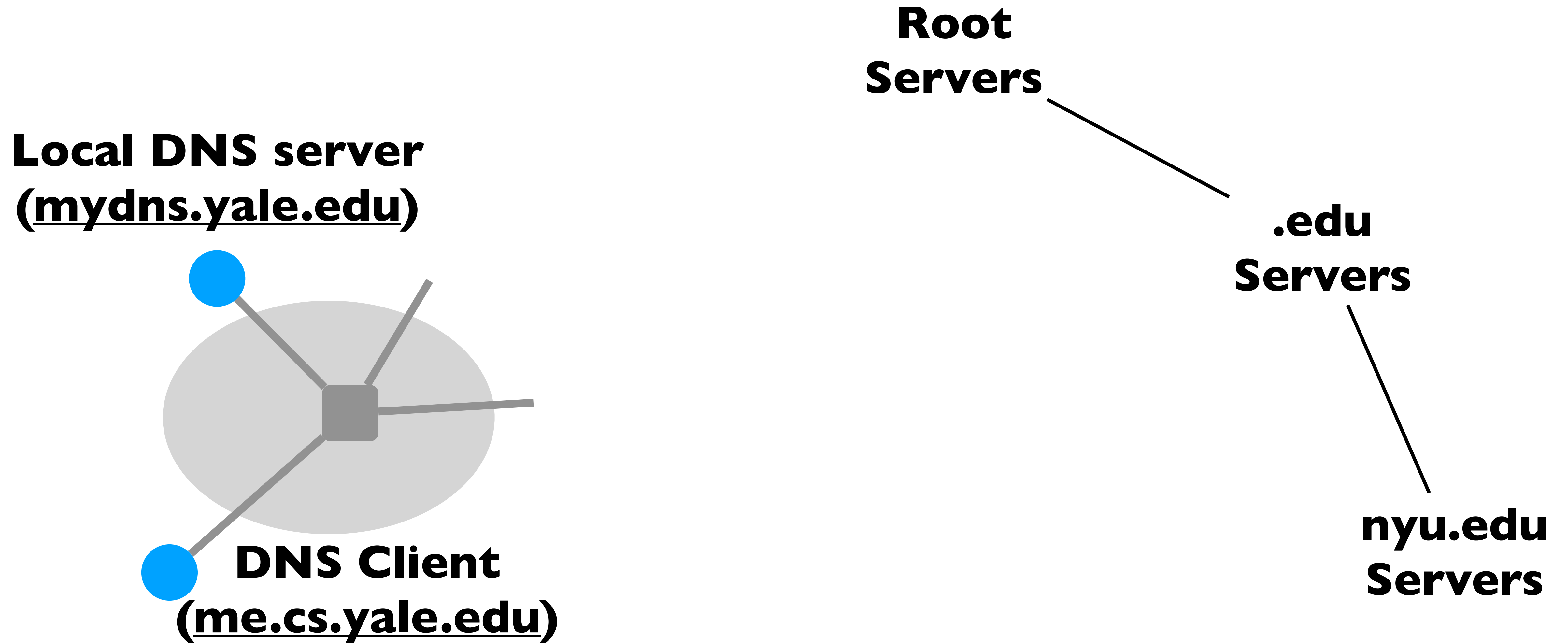
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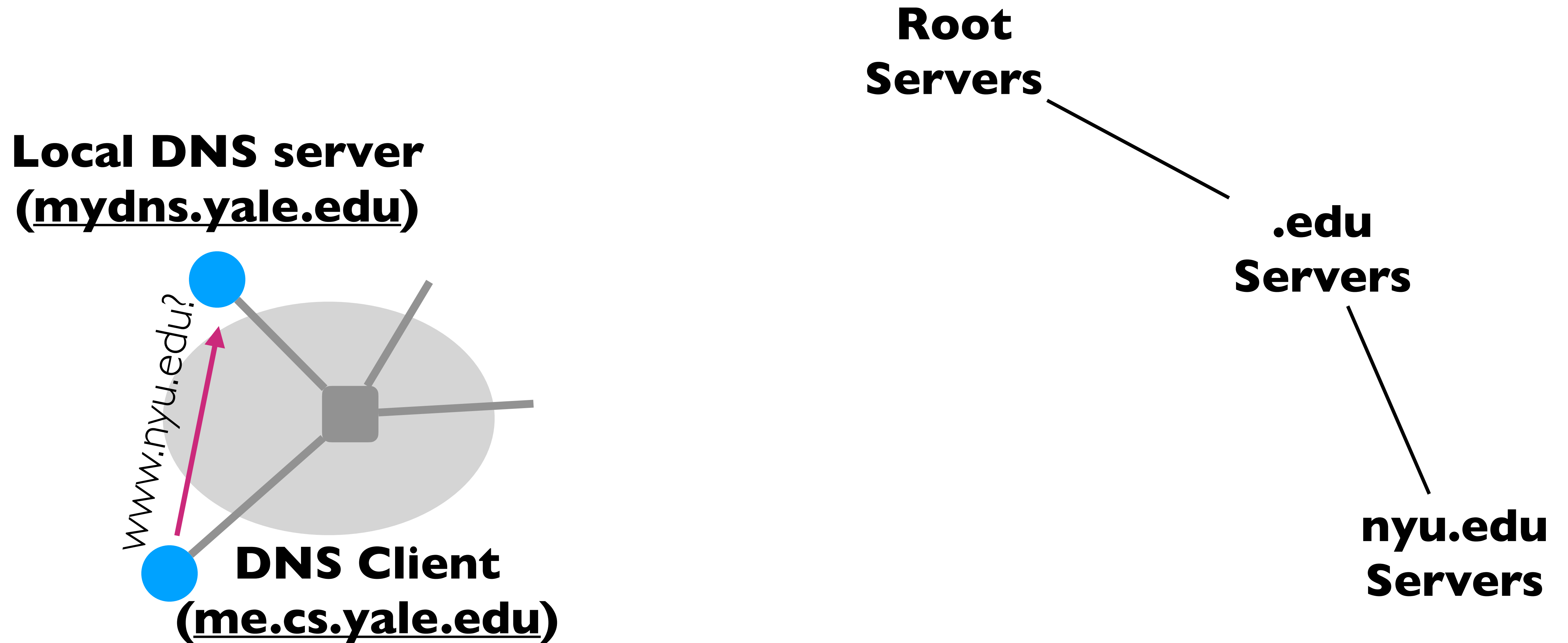
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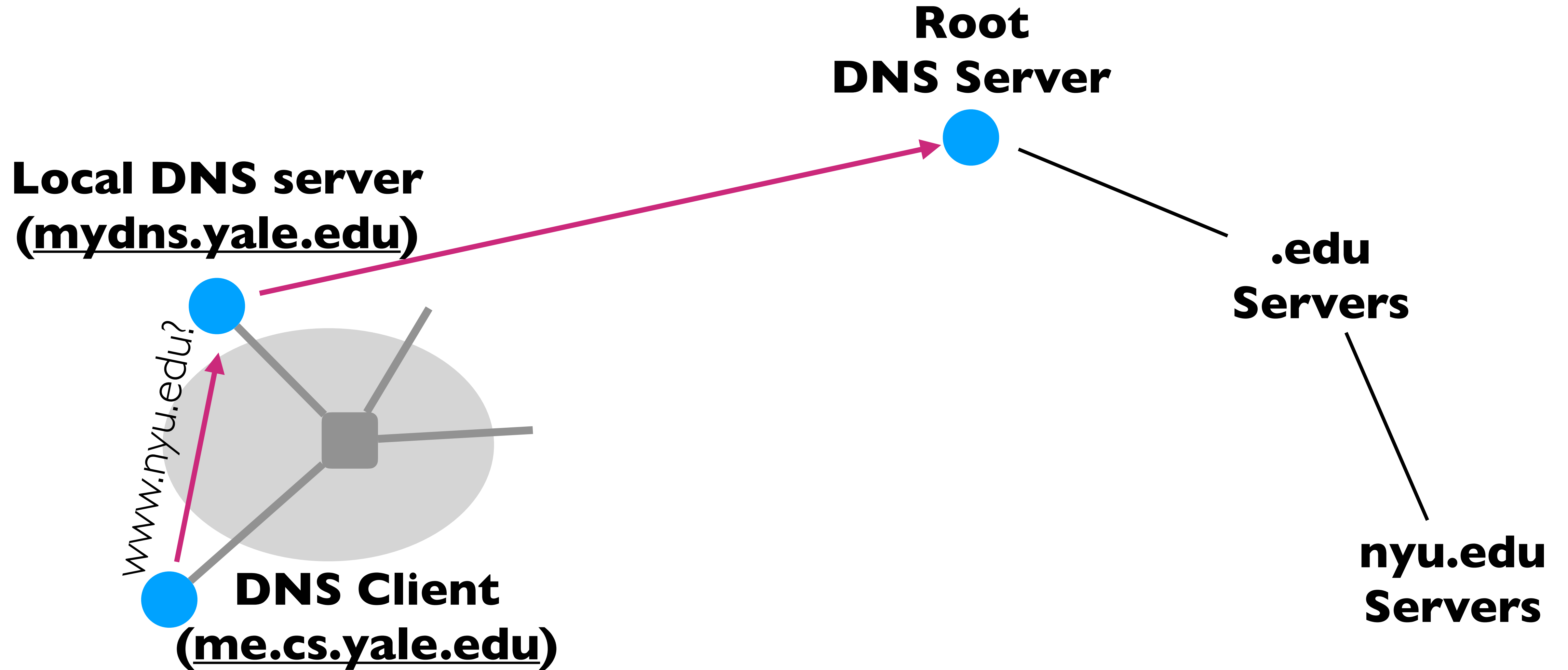


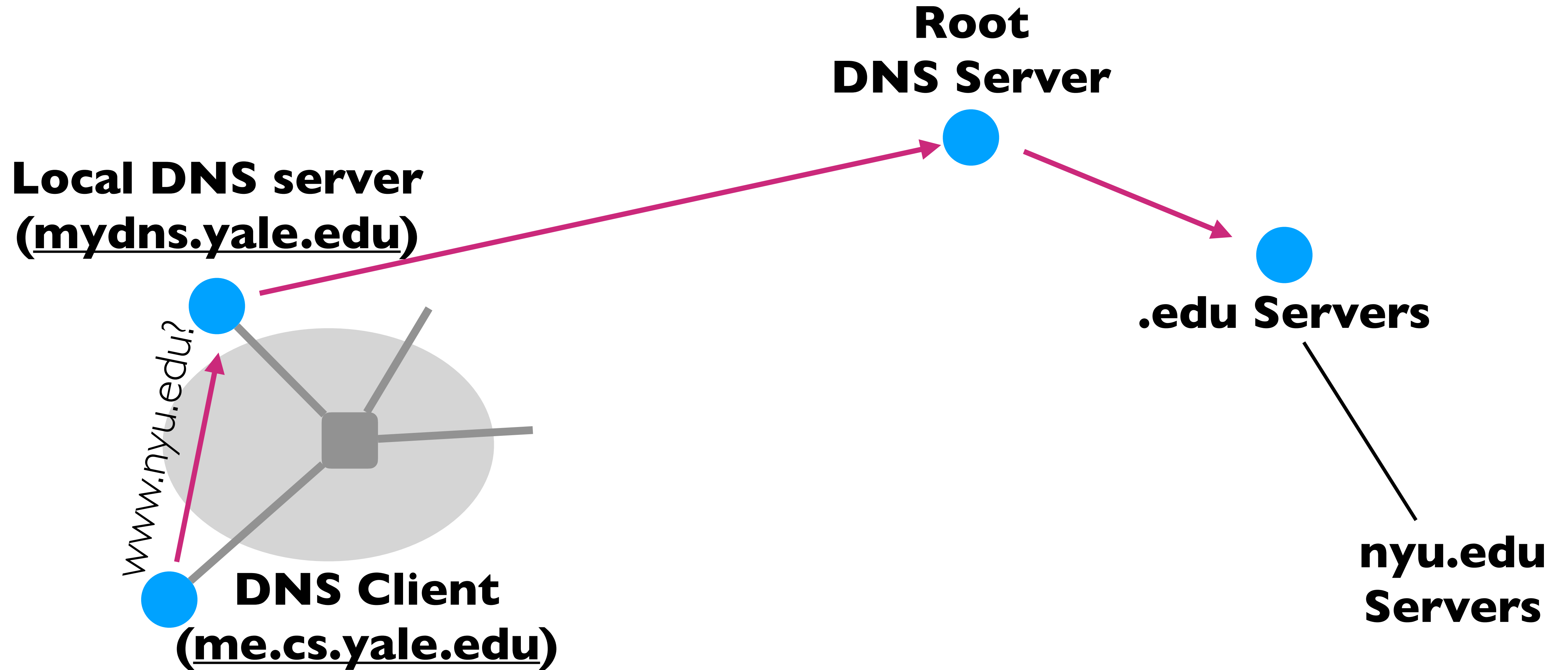
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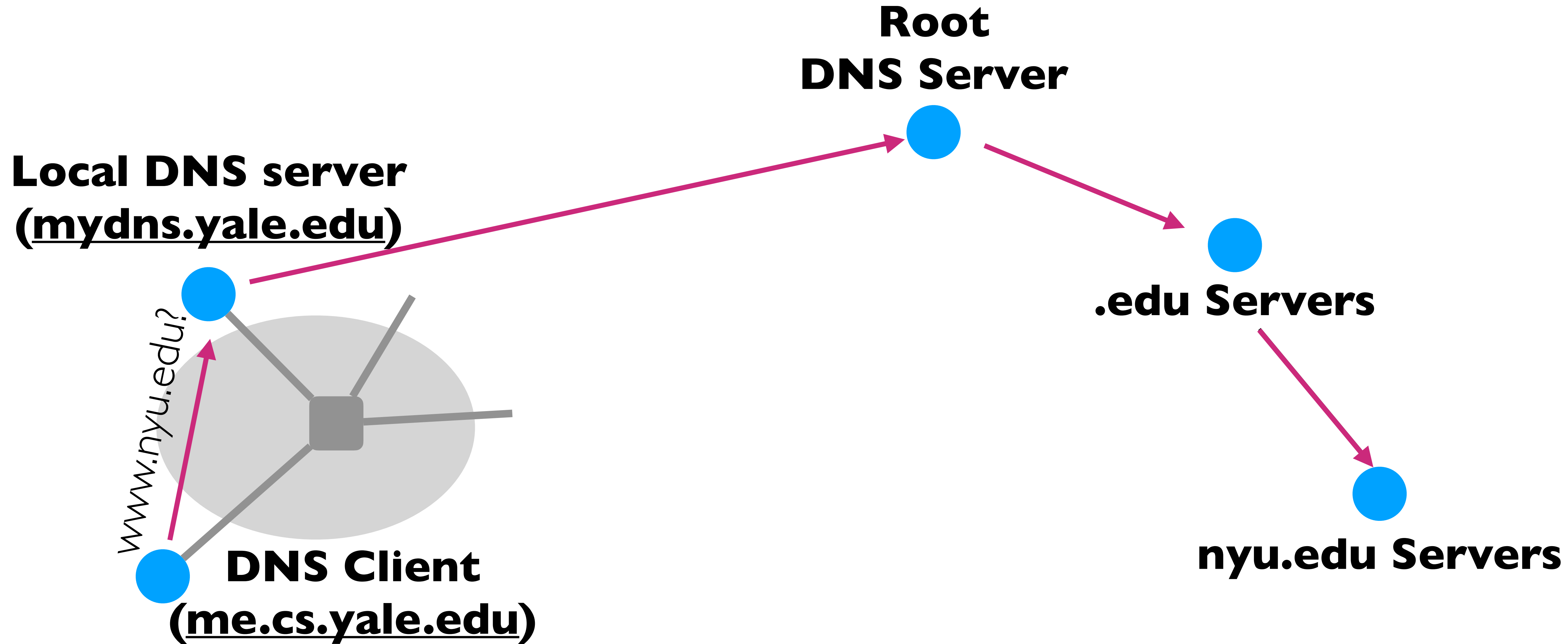


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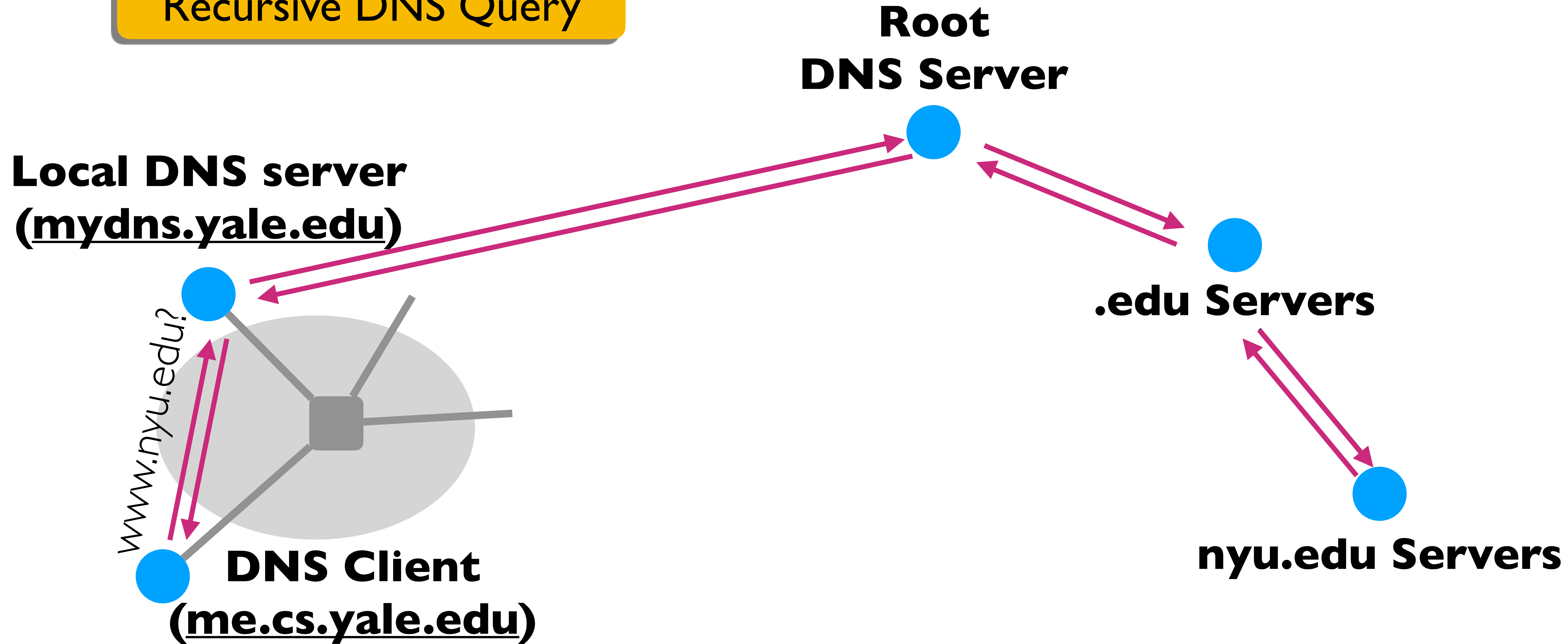




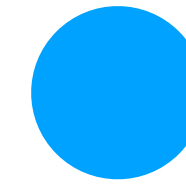




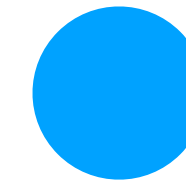
Recursive DNS Query



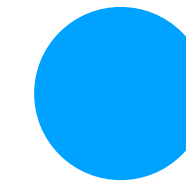
**Root
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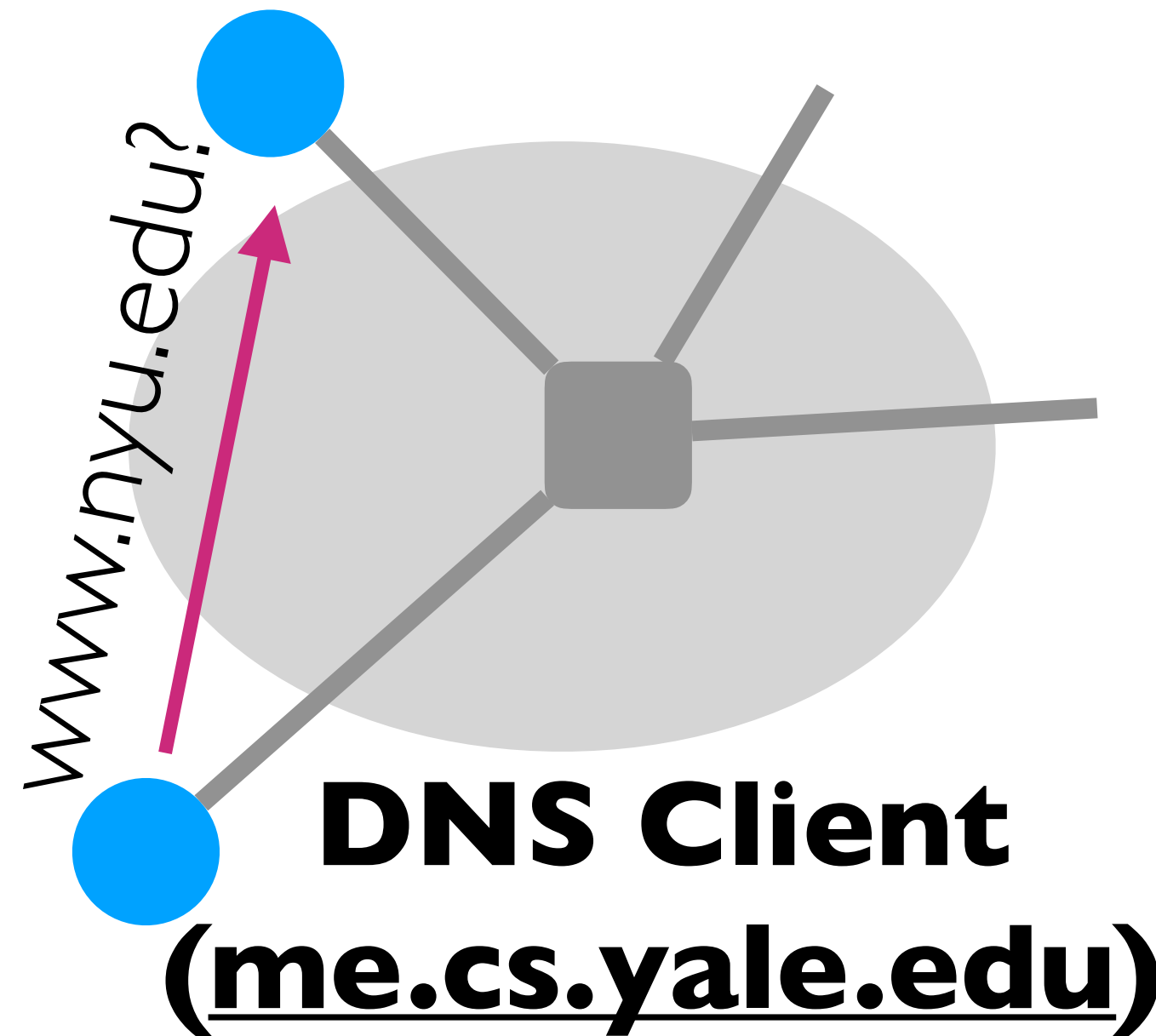
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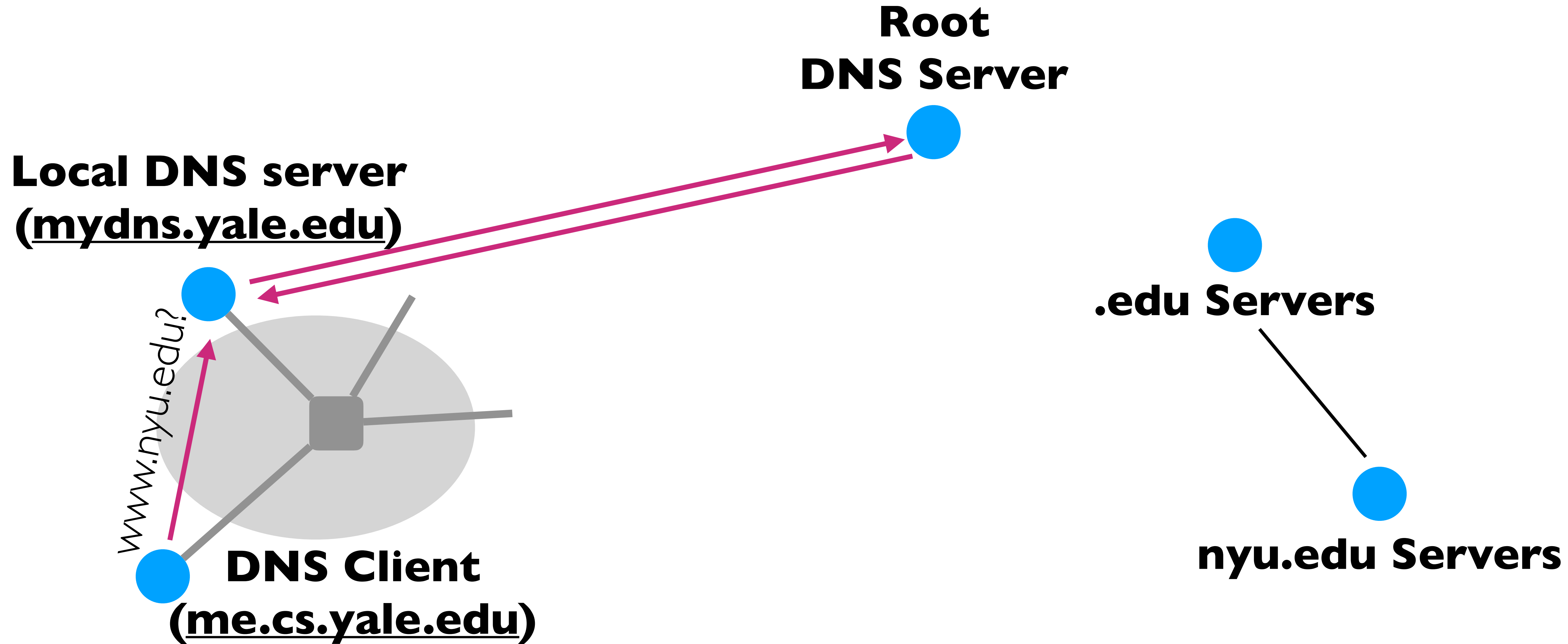


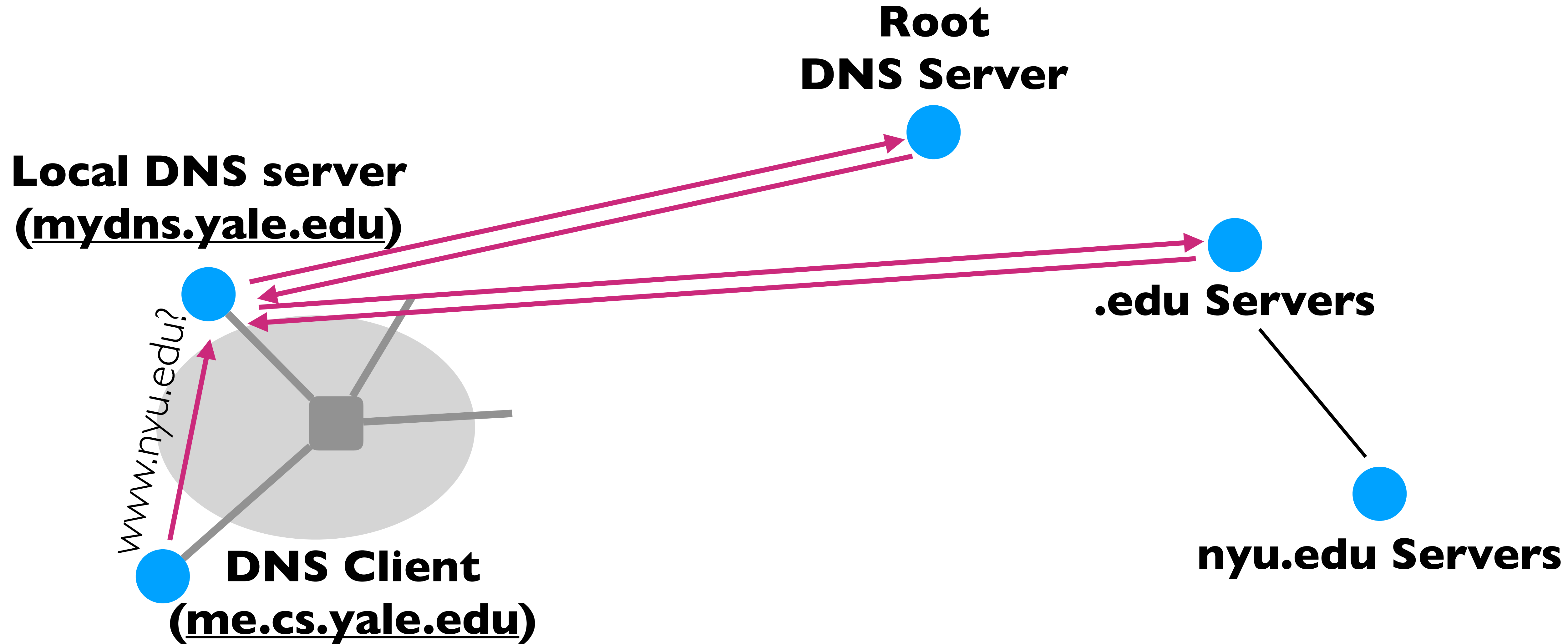
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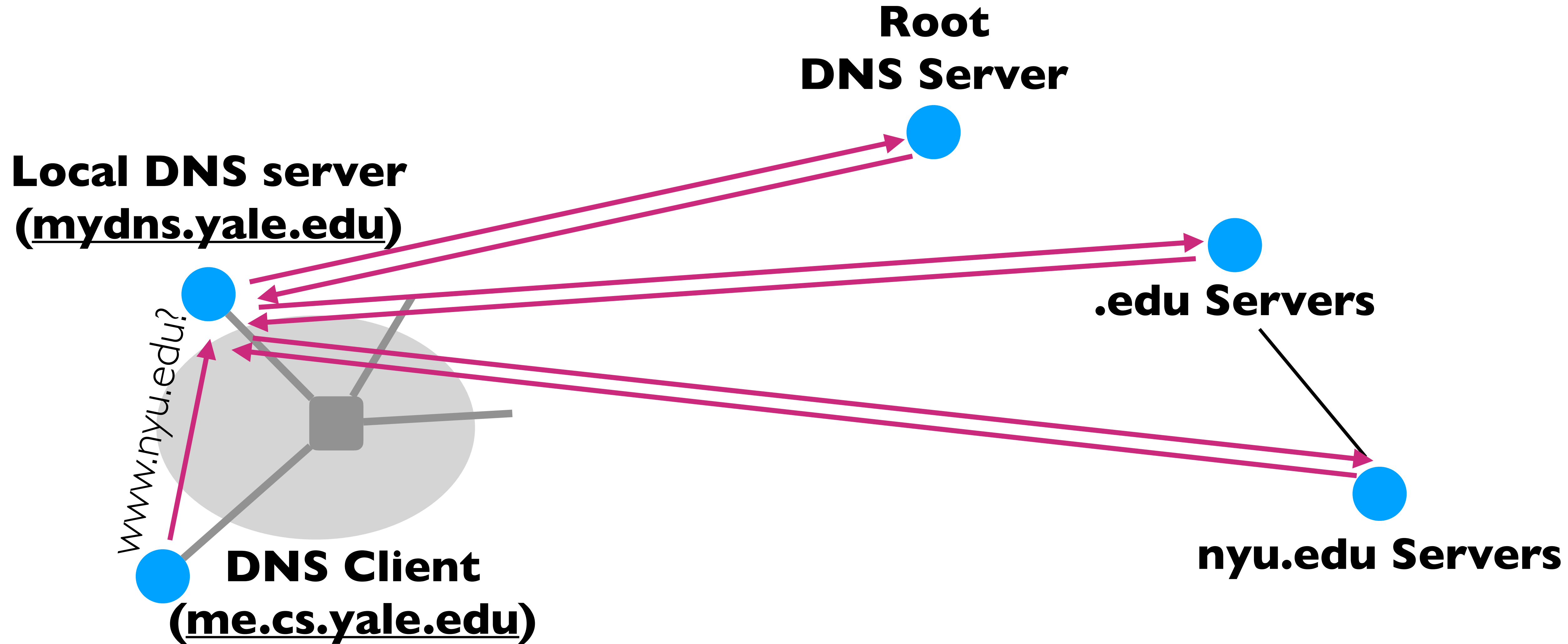


**Local DNS server
(mydns.yale.edu)**

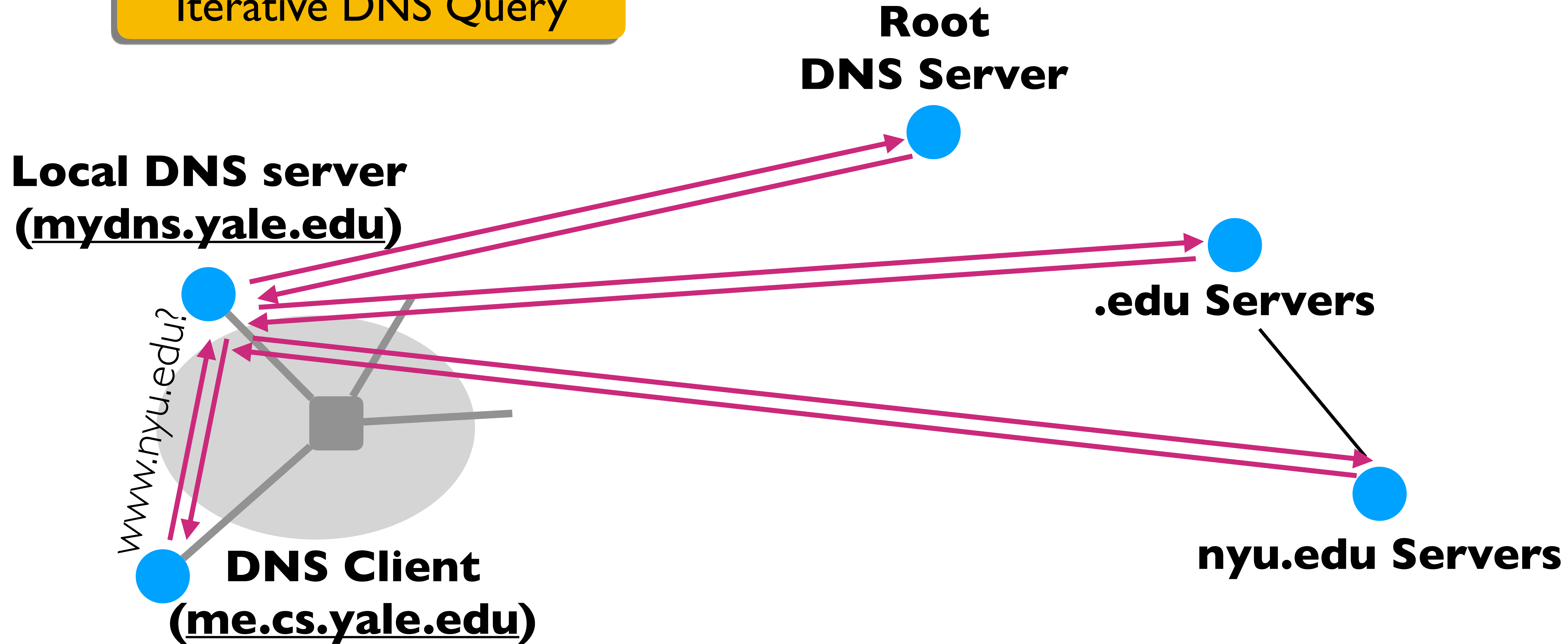








Iterative DNS Query



Recap: DNS Protocol

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- Query and Reply messages; both with the same message format
 - *See text for details*

Recap: DNS Protocol

- Query and Reply messages; both with the same message format
 - *See text for details*
- Client-Server Interaction on UDP Port 53
 - Spec. supports TCP too, but not always implemented

Questions?

Goals: How are we doing?

- Scalable
 - Many names
 - Many updates
 - Many users creating names
 - Many users looking up names
- Highly available
- Correct
 - No naming conflicts (uniqueness)
 - Consistency
- Lookups are fast

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Per-domain Availability

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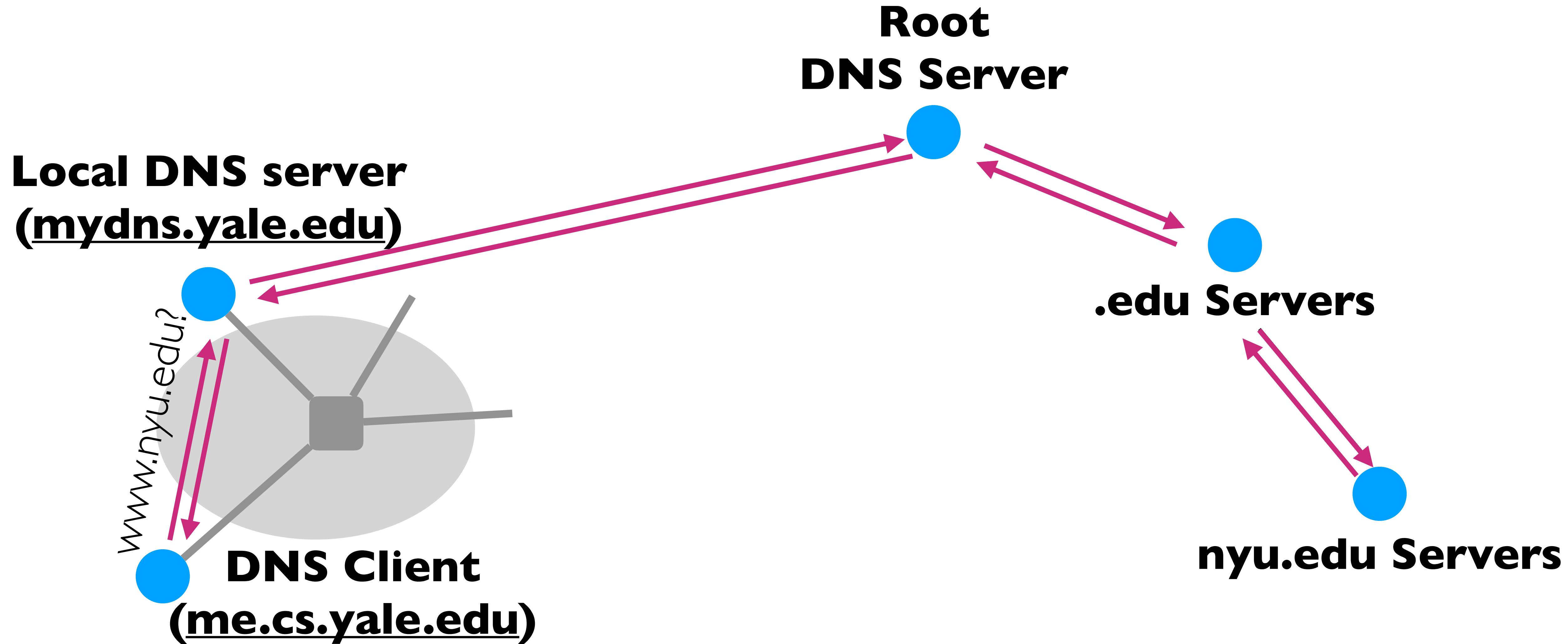
- DNS servers are **replicated**
 - Primary and secondary name servers required
 - Name service available if at least one replica is up
 - Queries can be load-balanced between replicas

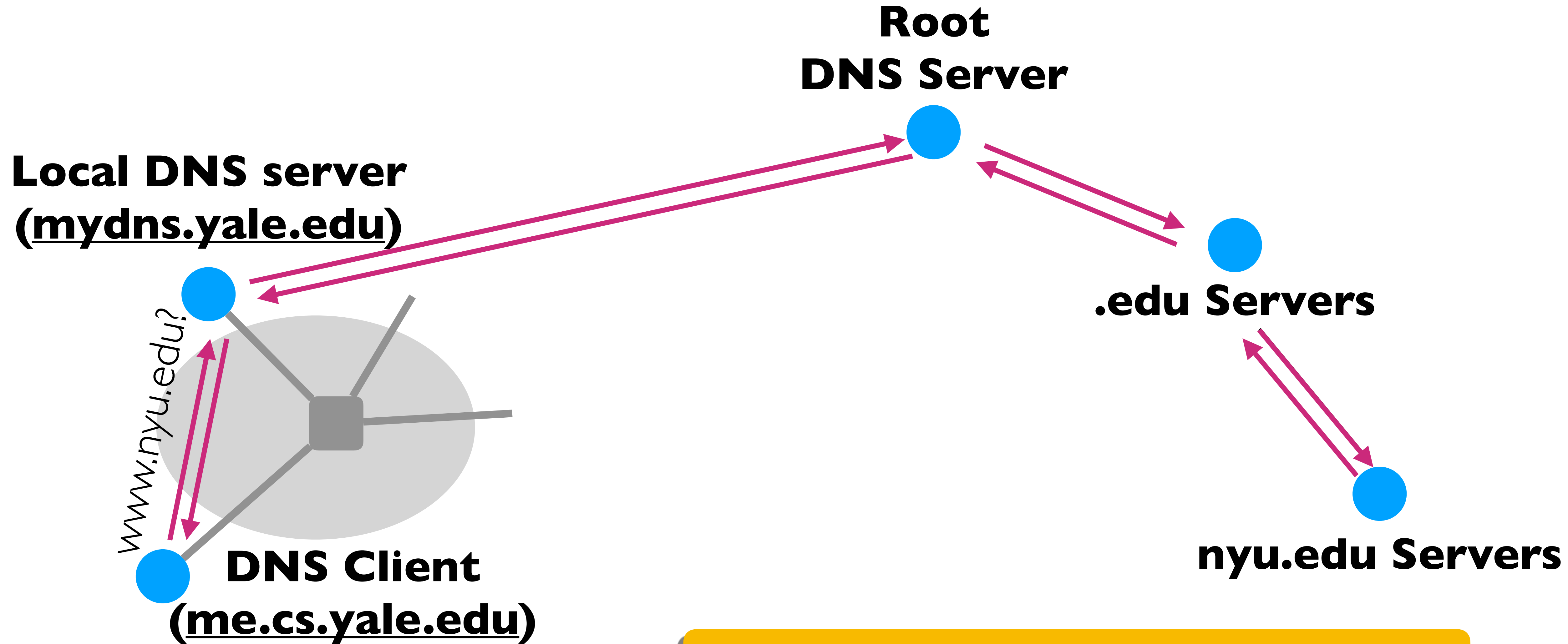
Per-domain Availability

- DNS servers are **replicated**
 - Primary and secondary name servers required
 - Name service available if at least one replica is up
 - Queries can be load-balanced between replicas
- Try alternate servers on timeout
 - *Exponential backoff* when retrying same server

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How would you speed up this process?

Caching

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- Caching of DNS responses at all levels

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- Caching of DNS responses at all levels
- Reduces load at all levels
- Reduces delay experienced by DNS client

DNS Caching

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- **How DNS caching works**
 - DNS servers cache responses to queries
 - Responses include a “time-to-live” (TTL) field
 - Server deletes cached entry after TTL expires

DNS Caching

- **How DNS caching works**
 - DNS servers cache responses to queries
 - Responses include a “time-to-live” (TTL) field
 - Server deletes cached entry after TTL expires
- **Why caching is effective**
 - The top-level servers very rarely change
 - Popular sites visited often → local DNS server often has the information cached

Negative Caching

Negative Caching

- Remember things that don't work
 - Misspellings like [www.cnn.comm](#) and [www.cnnn.com](#)
 - These can take a long time to fail the first time
 - Good to remember that they don't work
 - ... so the failure takes less time the next time around

Negative Caching

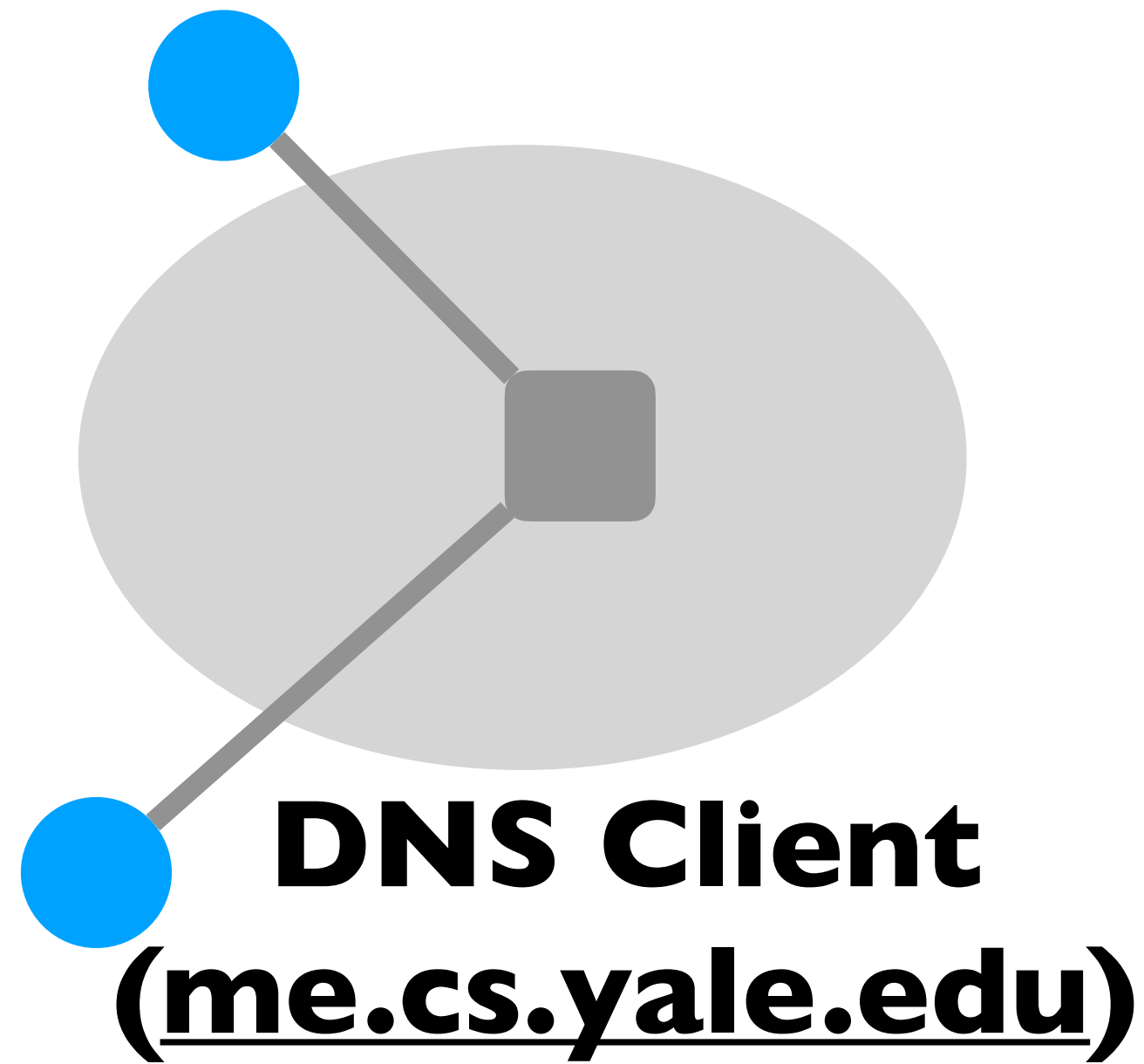
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 - ... so the failure takes less time the next time around
- Negative caching is optional

Questions?

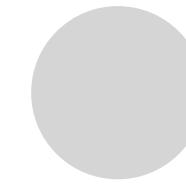
Time to put your malicious hats on...

How can one attack DNS?

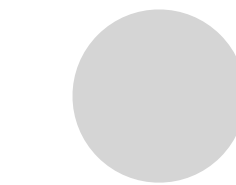
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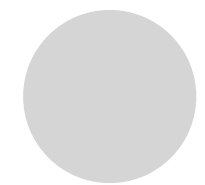
Root
DNS Server



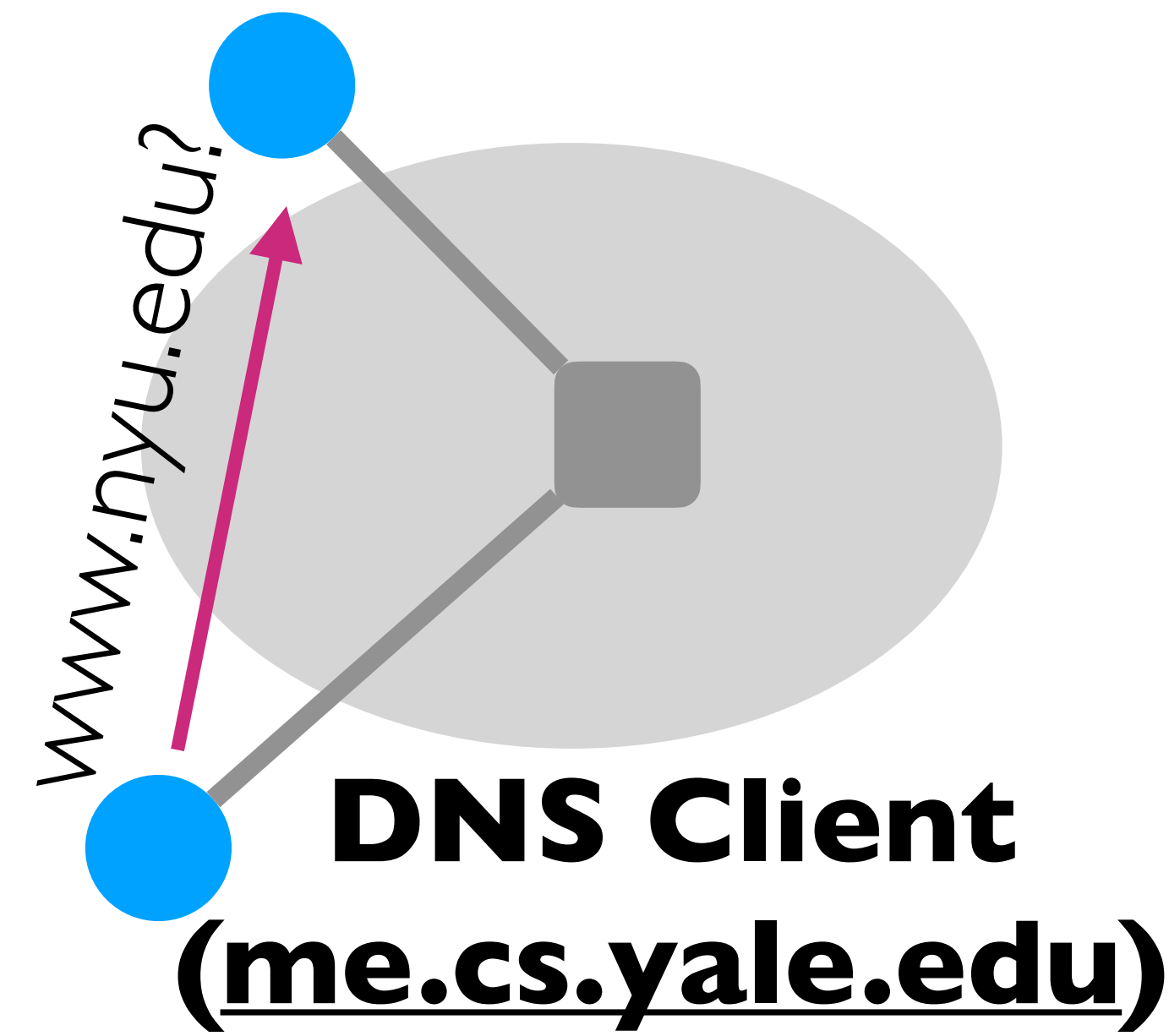
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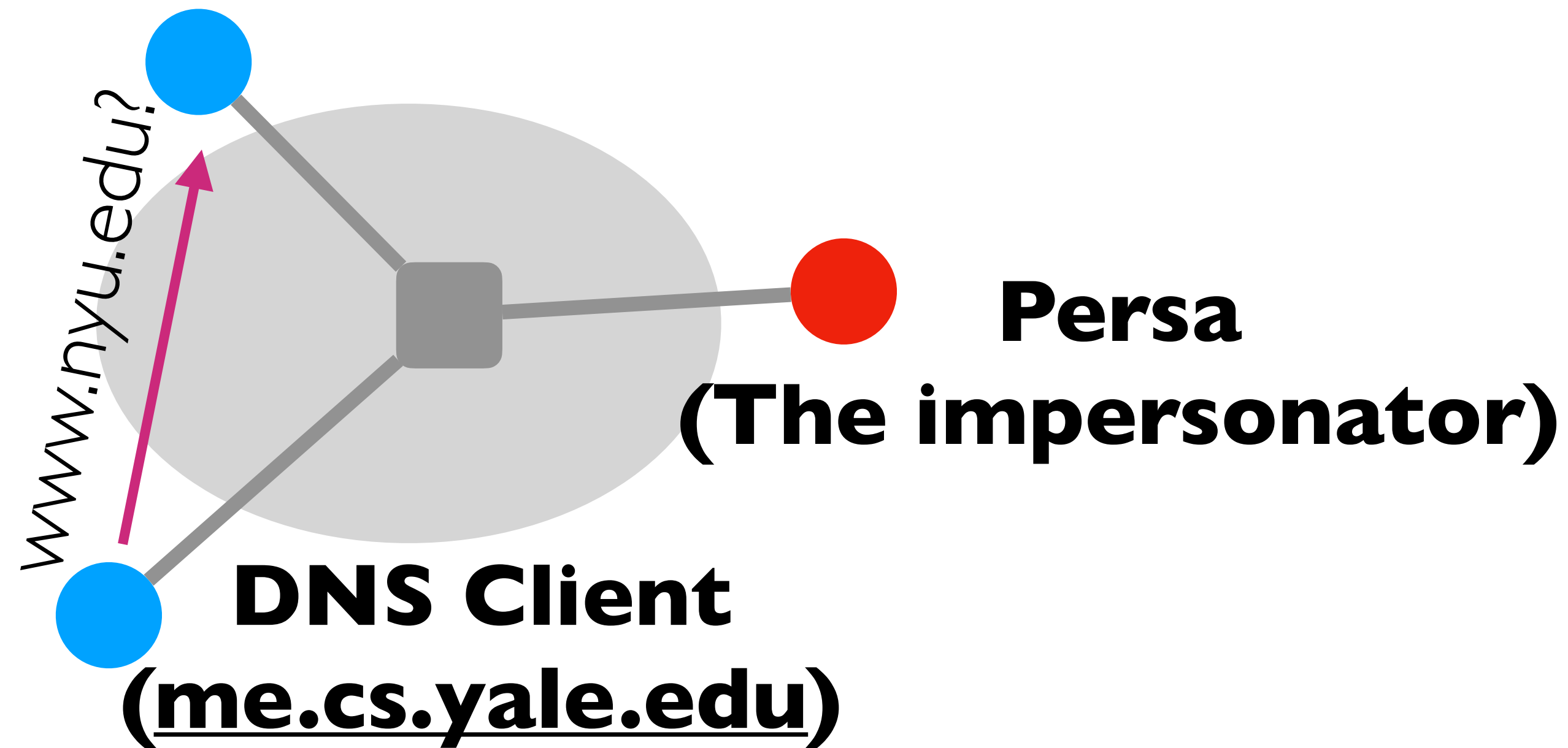
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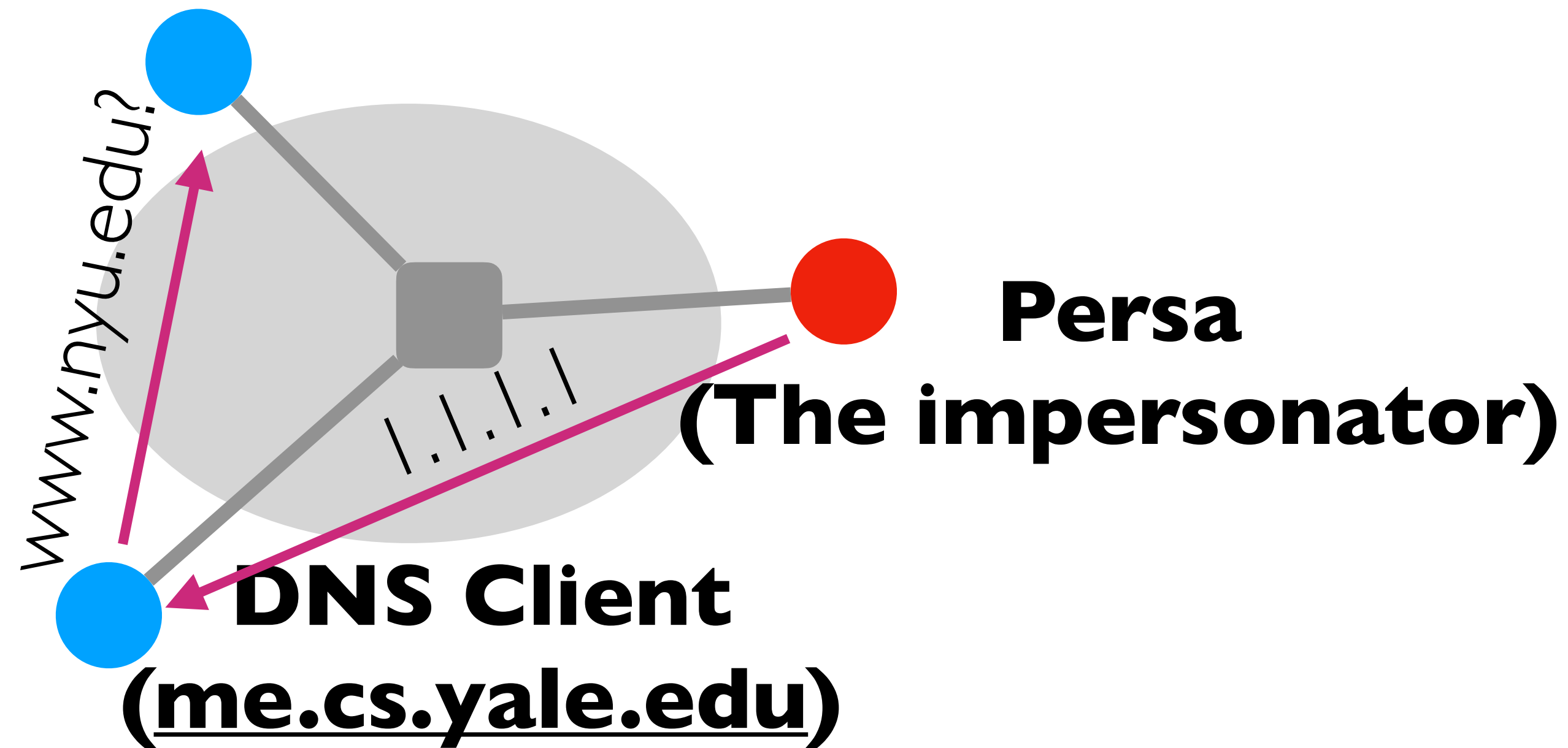
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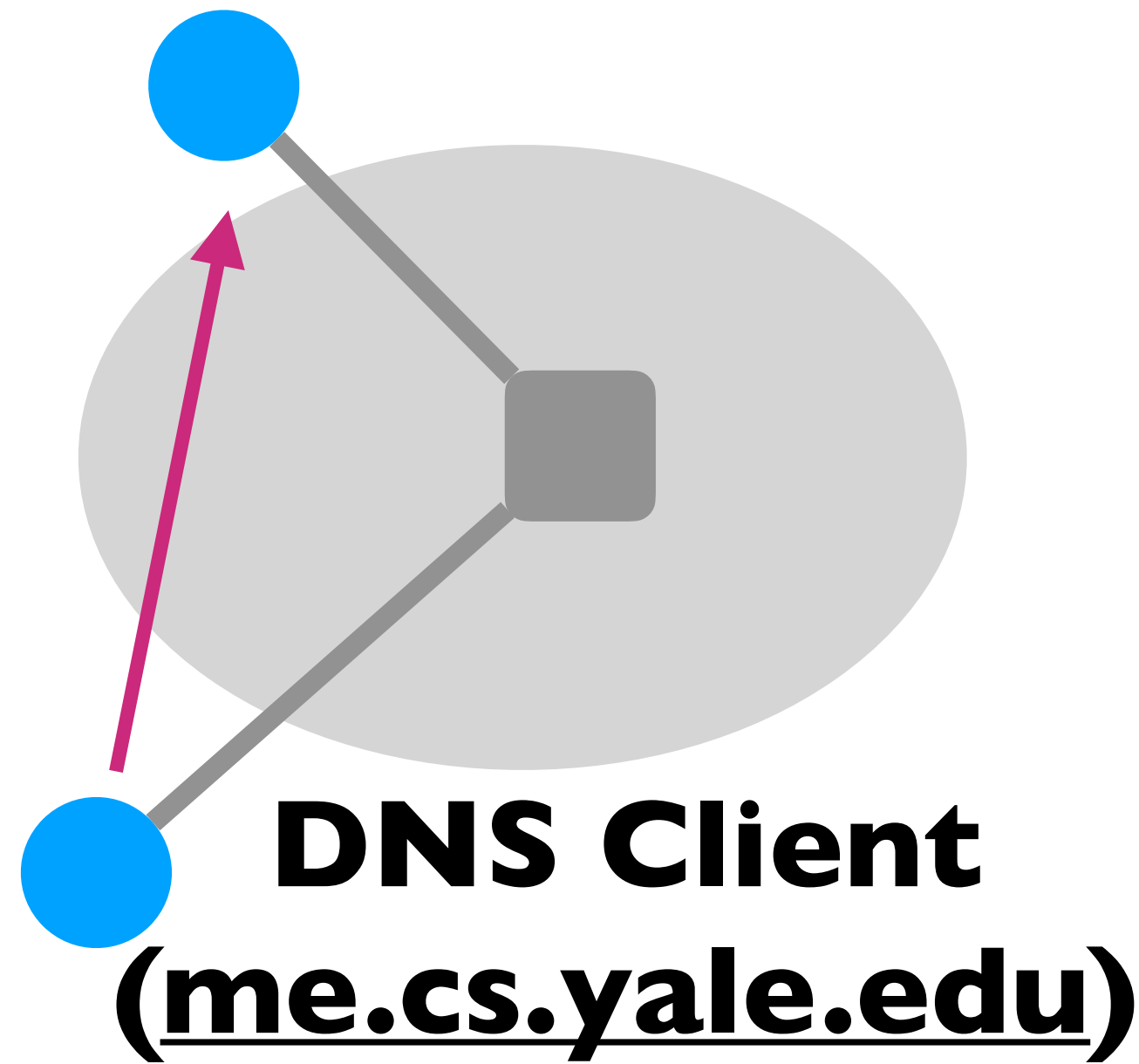
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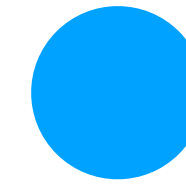
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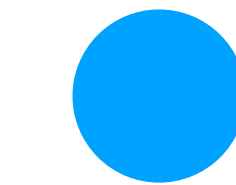
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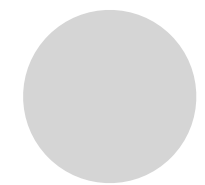
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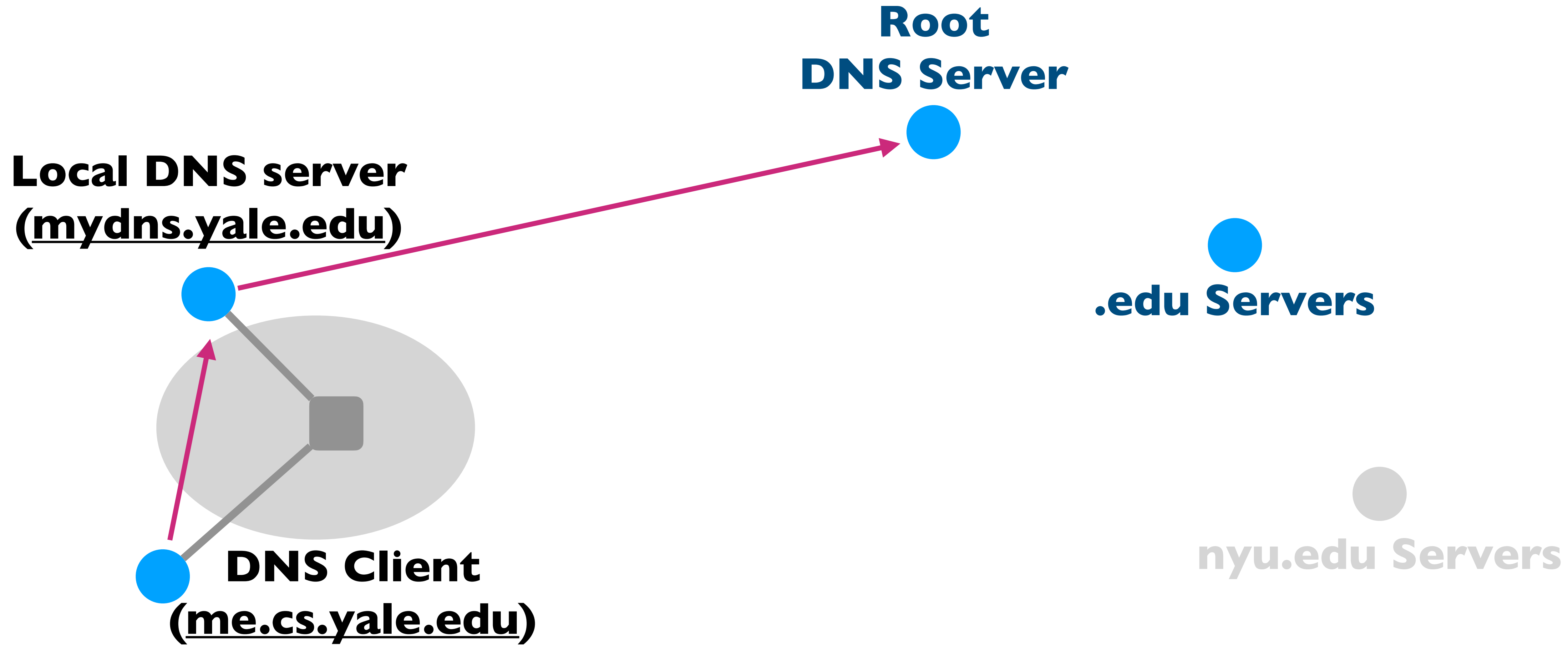


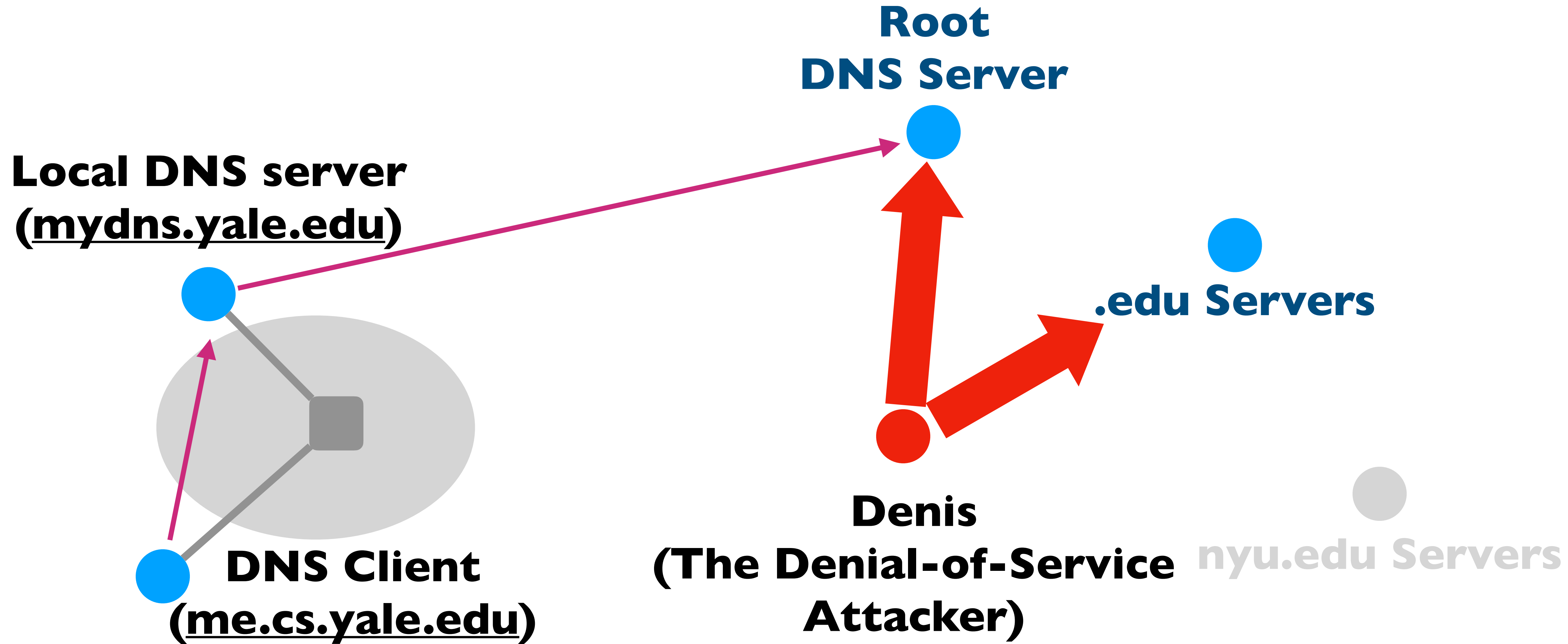
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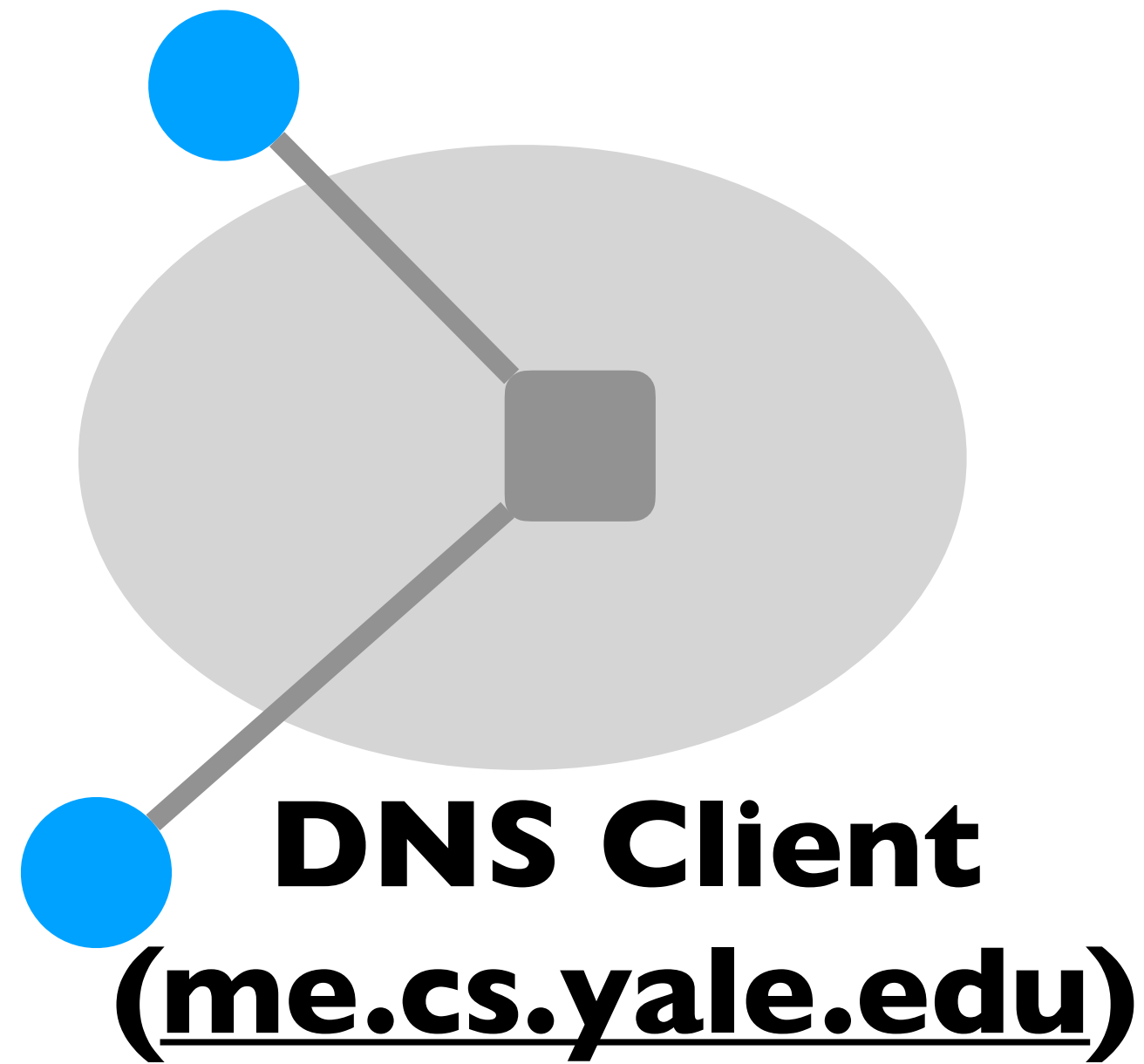




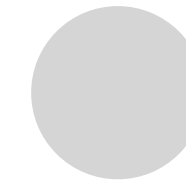
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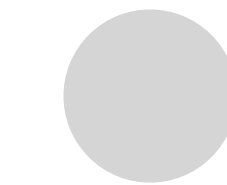
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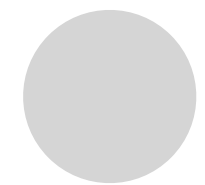
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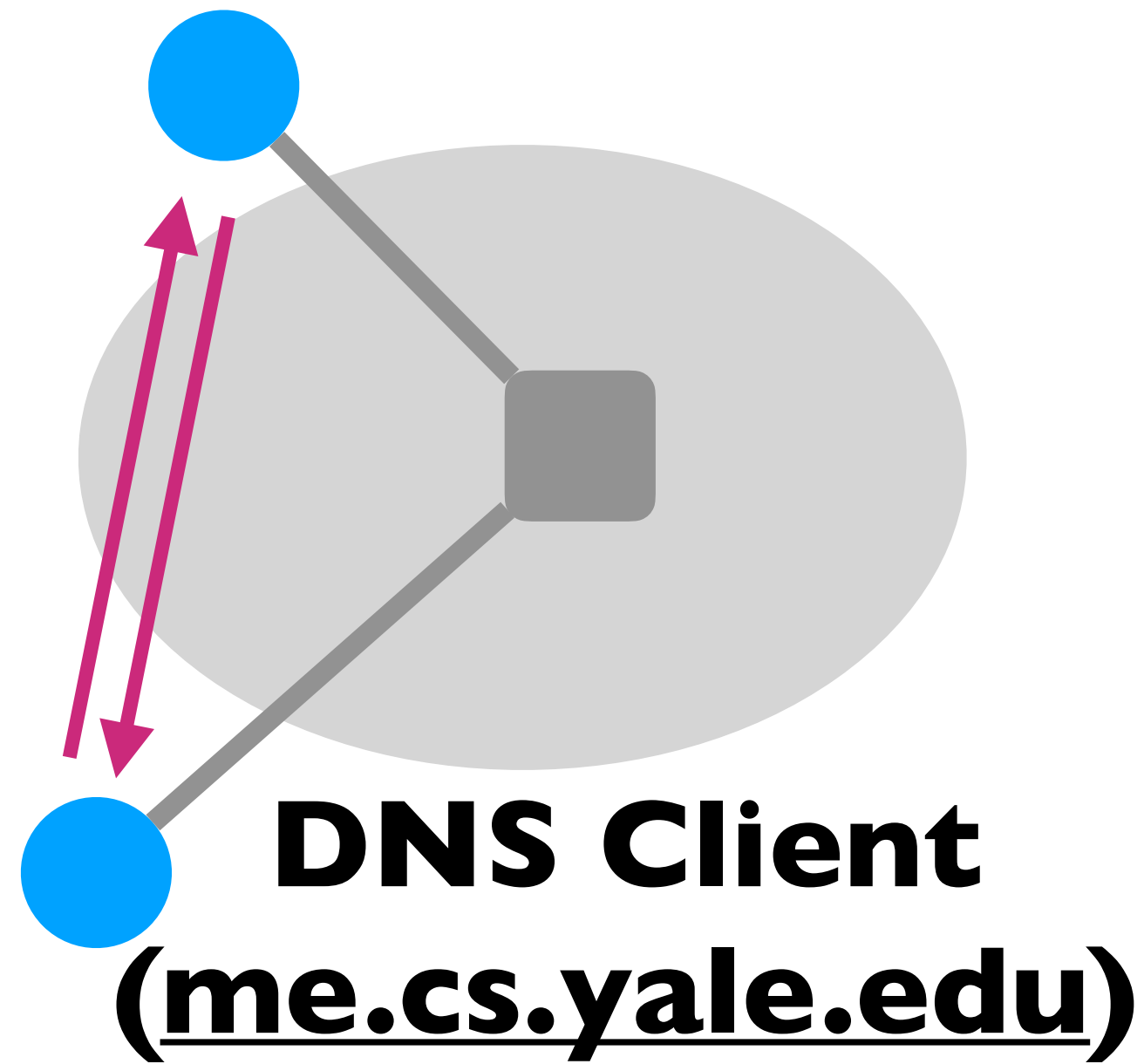
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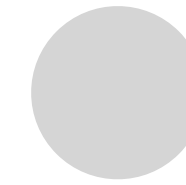
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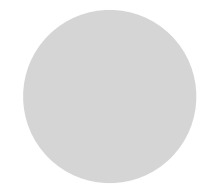
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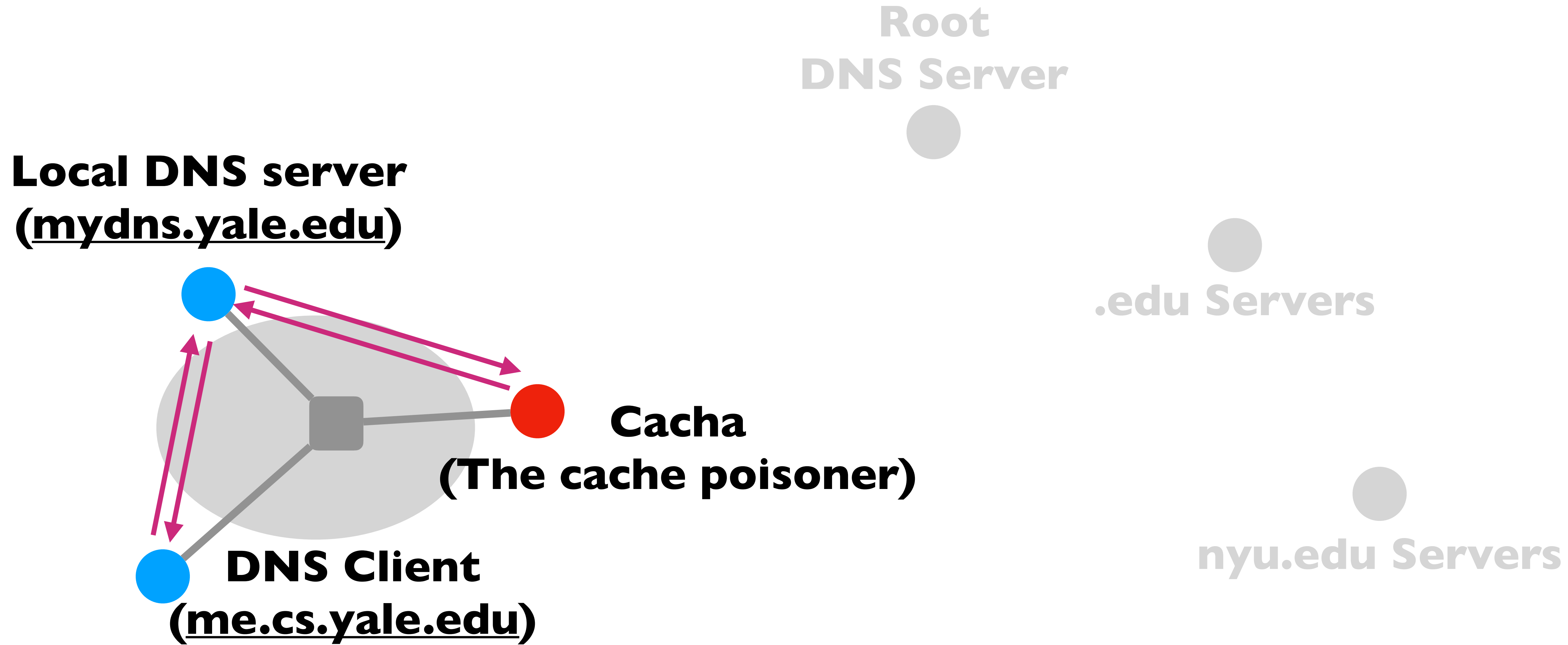


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How Can One Attack DNS?

- **Impersonate the local DNS server**
 - Give the wrong IP address to the DNS client
- **Denial-of-service the root or TLD servers**
 - Make them unavailable to the rest of the world
- **Poison the cache of a DNS server**
 - Increase the delay experienced by DNS clients

Taking Stock: Important Properties of DNS

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Administrative delegation and hierarchy results in:

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Administrative delegation and hierarchy results in:

- Easy unique naming
- “Fate sharing” for network failures
- Reasonable trust model
- Caching lends scalability, performance

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- Addresses can **change** underneath
 - Move www.cnn.com to a new IP address
 - Humans/applications are unaffected

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- Allowing “host” names to evolve into “service” names

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