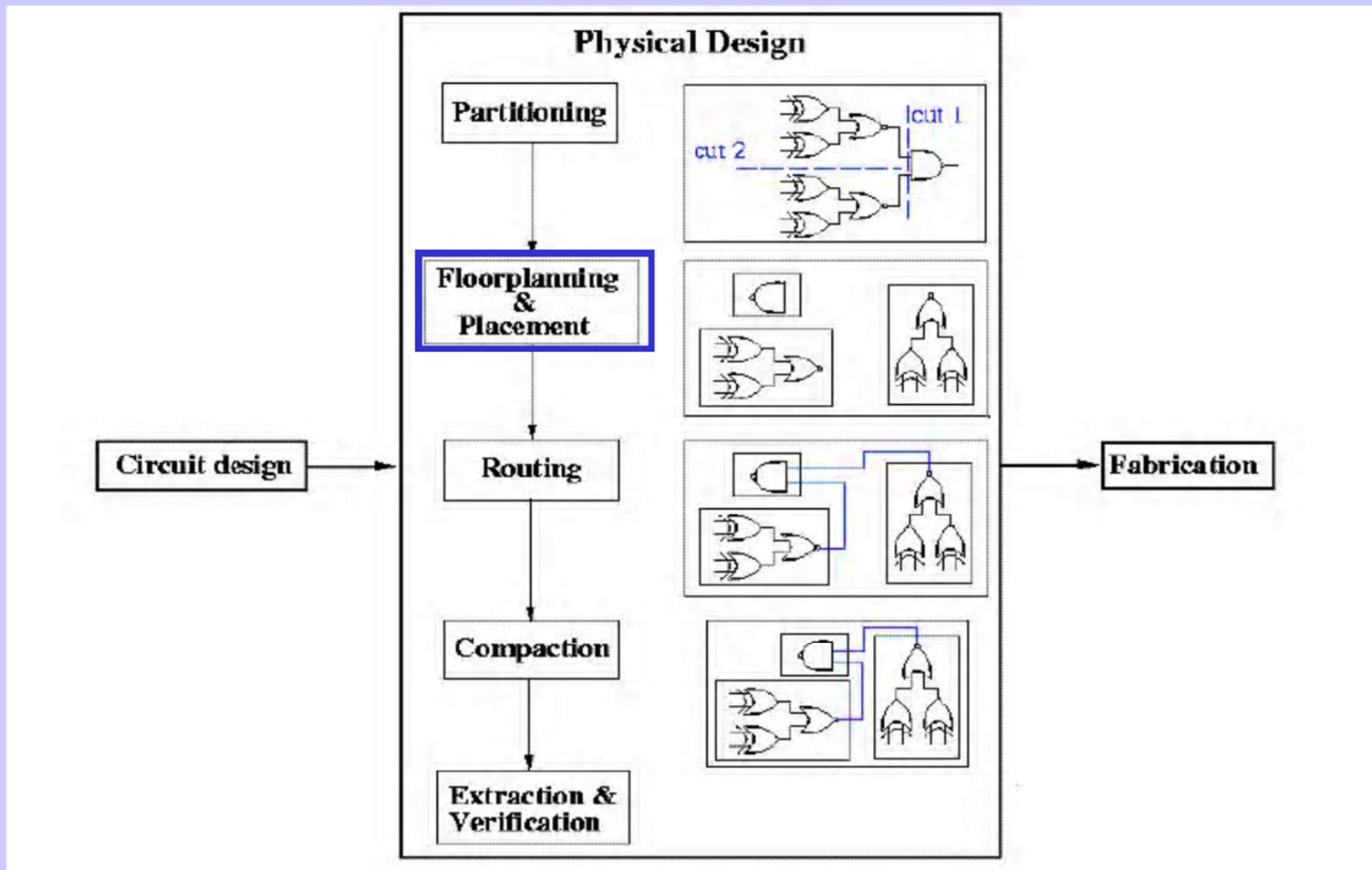


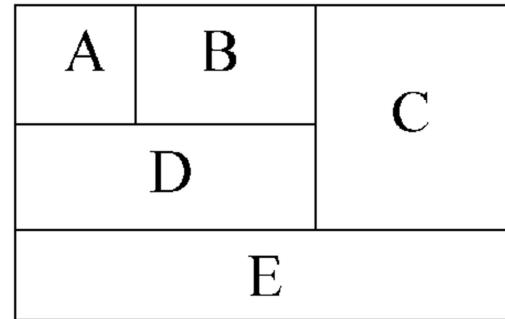
# Floorplanning-1

# Physical Design Flow



# Floorplans

- A *floorplan* is the placement of (flexible) blocks with fixed area but unknown dimensions
- Blocks are usually assumed to be rectangular in shape
- Some work done with L-shaped blocks
- Usually a lower and an upper bound on the Aspect ratio are given



A Floorplan with 5 rectangular blocks

A floorplan is usually represented by a rectangular dissection

# Floorplanning Problem

Objective :

Finding a suitable topology of the blocks and their dimensions such that

- Area minimized
- Critical net delay minimized

# Floorplanning algorithm

A good Floorplanning algorithm should

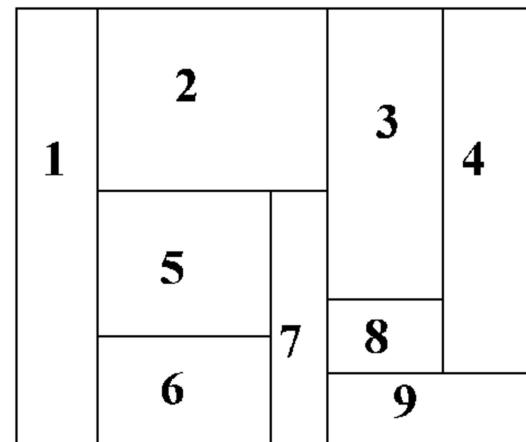
- minimize total chip area
- make the subsequent routing phase easy
- improve the performance, by, for example, reducing signal delays

It is difficult to

- consider all these goals simultaneously, as they mutually conflict
- consider algorithmically several objectives simultaneously

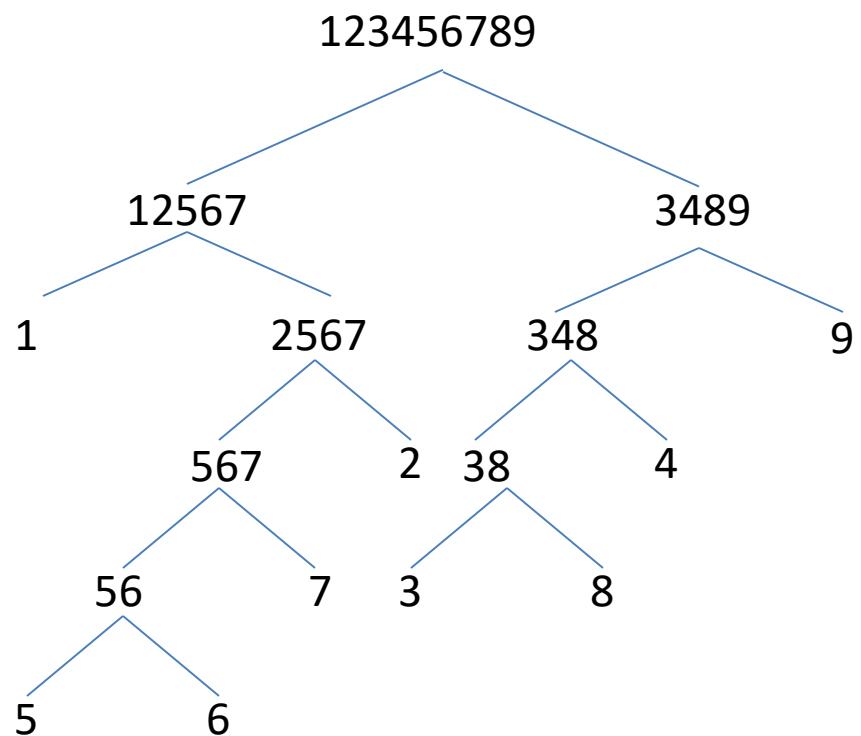
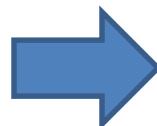
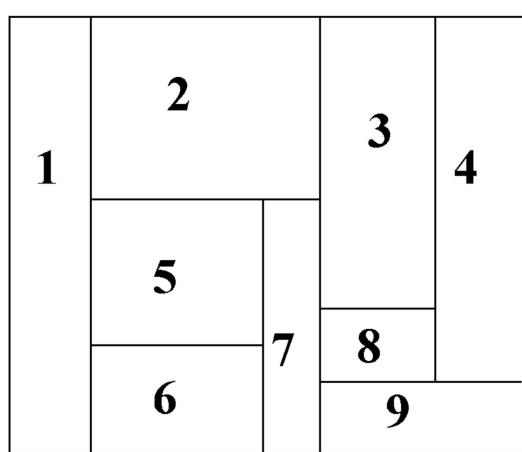
## Slicable Floorplans or Rectangular Duals

A floorplan is *Slicable* if it is obtained by recursively bipartitioning a rectangle into two sliceable floorplans either by a horizontal or a vertical line



Slicable Floorplan

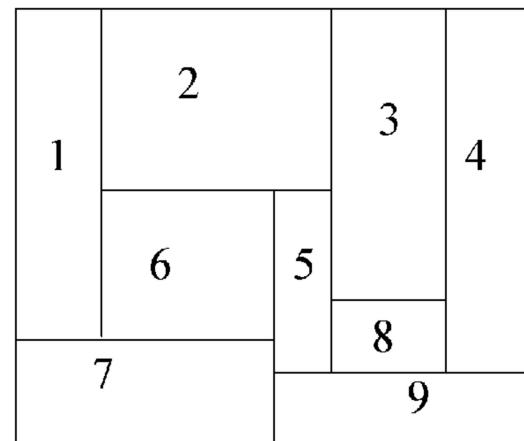
## Sliceable Floorplans



Easy to represent with binary tree  
- facilitates efficient algorithm design  
Sizing and routing algorithms execute in polynomial time

# Non-Slicable Floorplans

A floorplan that is not *Slicable*

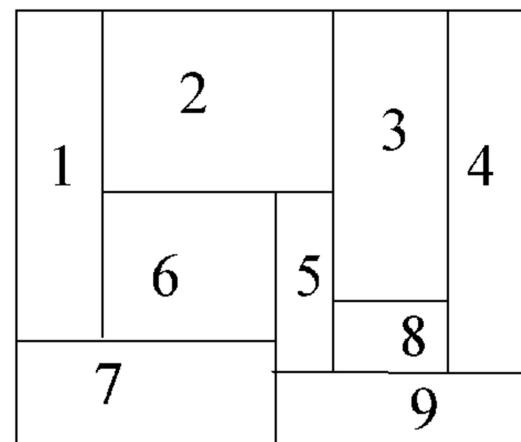
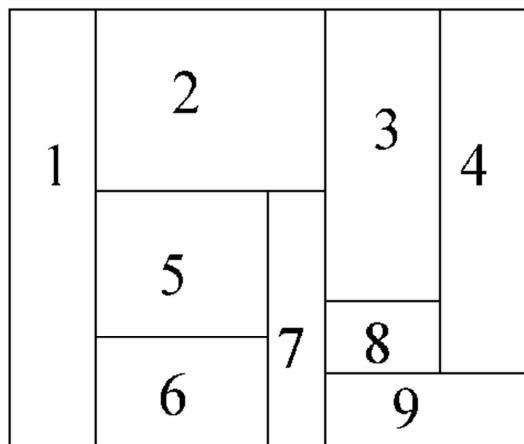


NonSlicable Floorplan

Floorplan sizing problem for nonslicable floorplan is NP complete

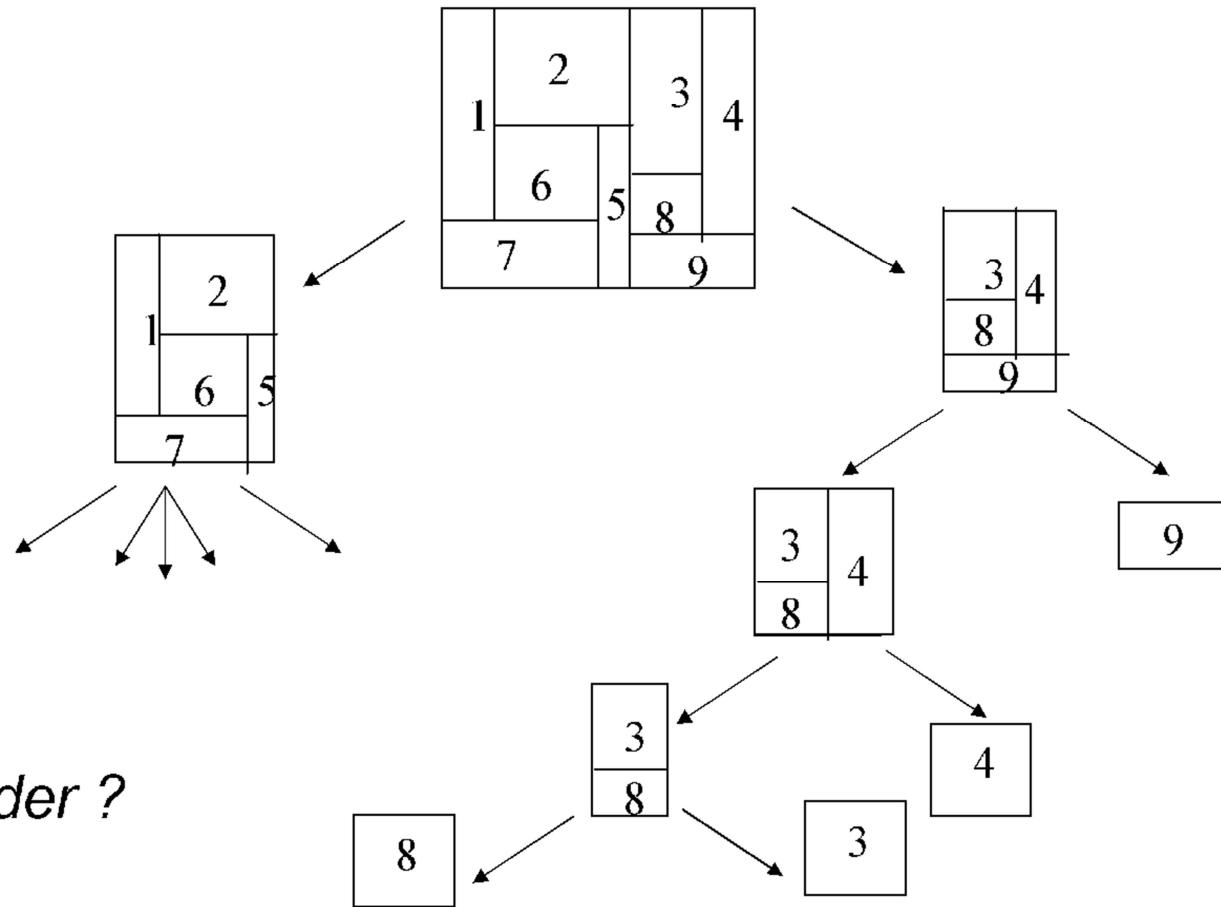
## Floorplans

A *floorplan* is said to be hierarchical order of  $k$ , if it can be obtained by recursively partitioning a rectangle into at most  $k$  parts



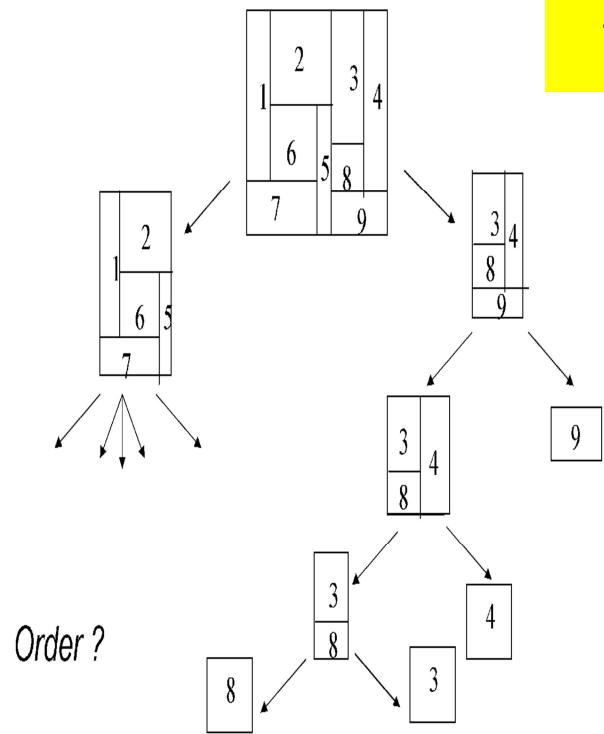
What are the hierarchical orders of the above floorplans ?

# Hierarchical Floorplan Trees



Hierarchically defined floorplan is a floorplan  
that can be described by a floorplan tree

## Hierarchical Floorplan Trees



The leaves of the tree correspond to modules

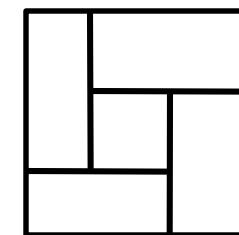
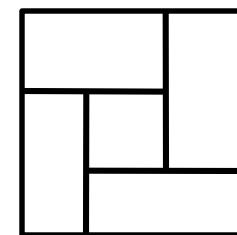
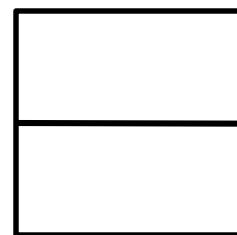
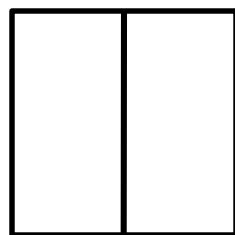
An internal node defines how its child floorplans are combined to form a partial floorplan

Various restrictions may be imposed on internal nodes

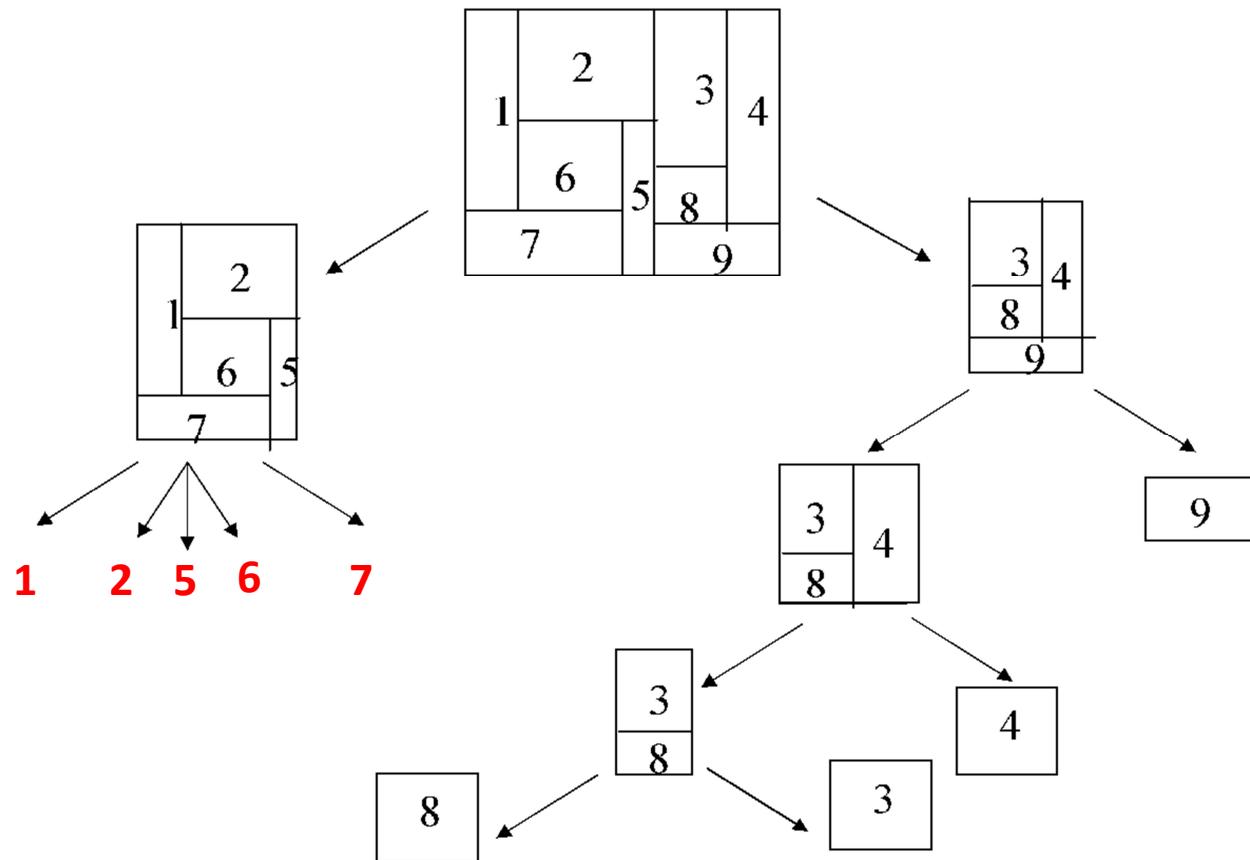
- (a) The number of child floorplans cannot exceed a certain limit
- (b) The shape of the floorplan corresponding to internal node is limited, for example, to rectangles or other simple shapes

# Hierarchical Floorplan

Four types of Internal nodes



# Hierarchical Floorplan Trees



Floorplan sizing problem for nonsliceable floorplan is NP complete

# Floorplanning by Rectangular Dualization

## *Input :*

- A set of rectangular blocks
- A set of realizations, i.e., (width, height) pairs, for each block
- Adjacency graph for the blocks

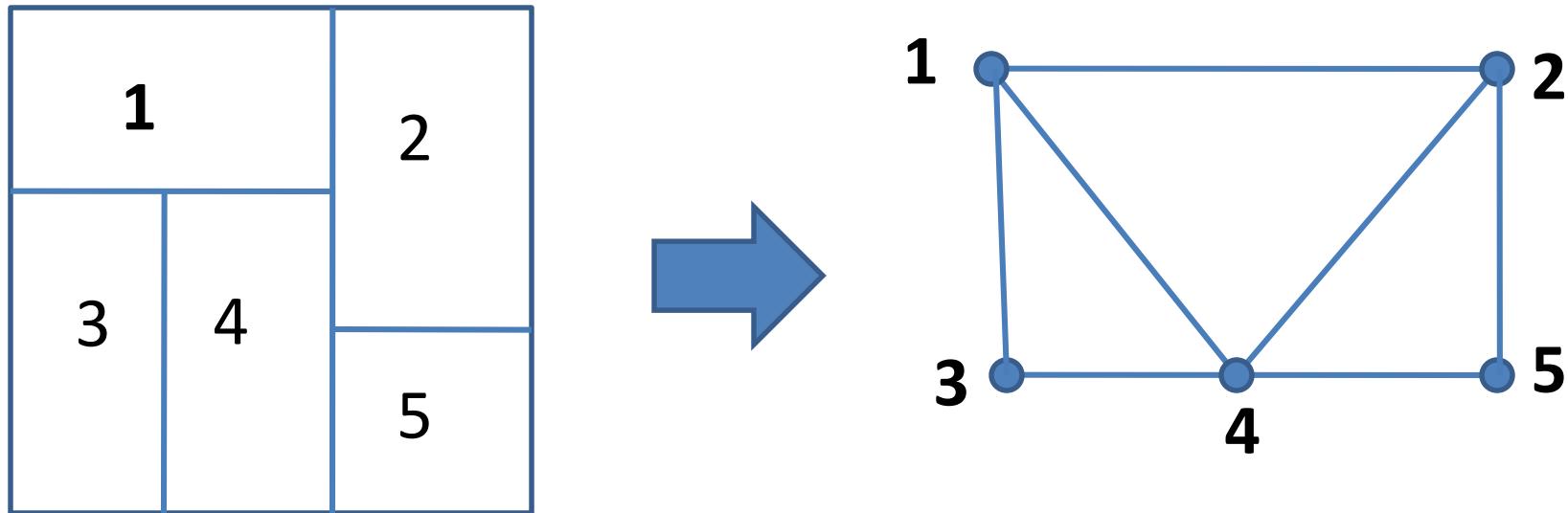
## *Requirements :*

- Topology generation --- Location of each block within a rectangular envelope such that no two blocks overlap
- Sizing --- An appropriate size, i.e., width and height, of each block

## *Objectives :*

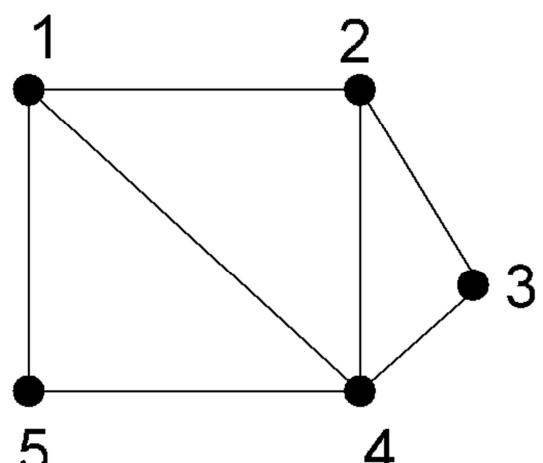
- Minimize area of the rectangular envelope
- Reduce net-length for critical nets

# Rectangular Dual Graph



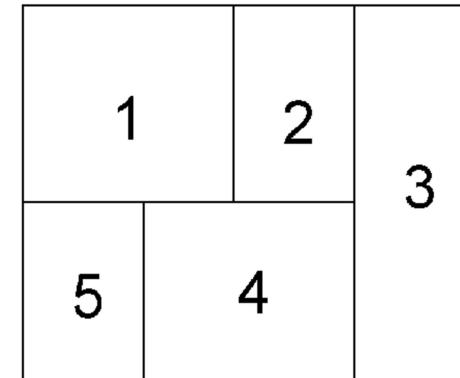
A rectangular dual graph of a rectangular floorplan is a plane graph  $G=(V,E)$ , where  $V$  is the set of modules and  $(M_i, M_j) \in E$  if and only if modules  $M_i$  and  $M_j$  are adjacent in the floorplan.

## Rectangular Dualization - An Example

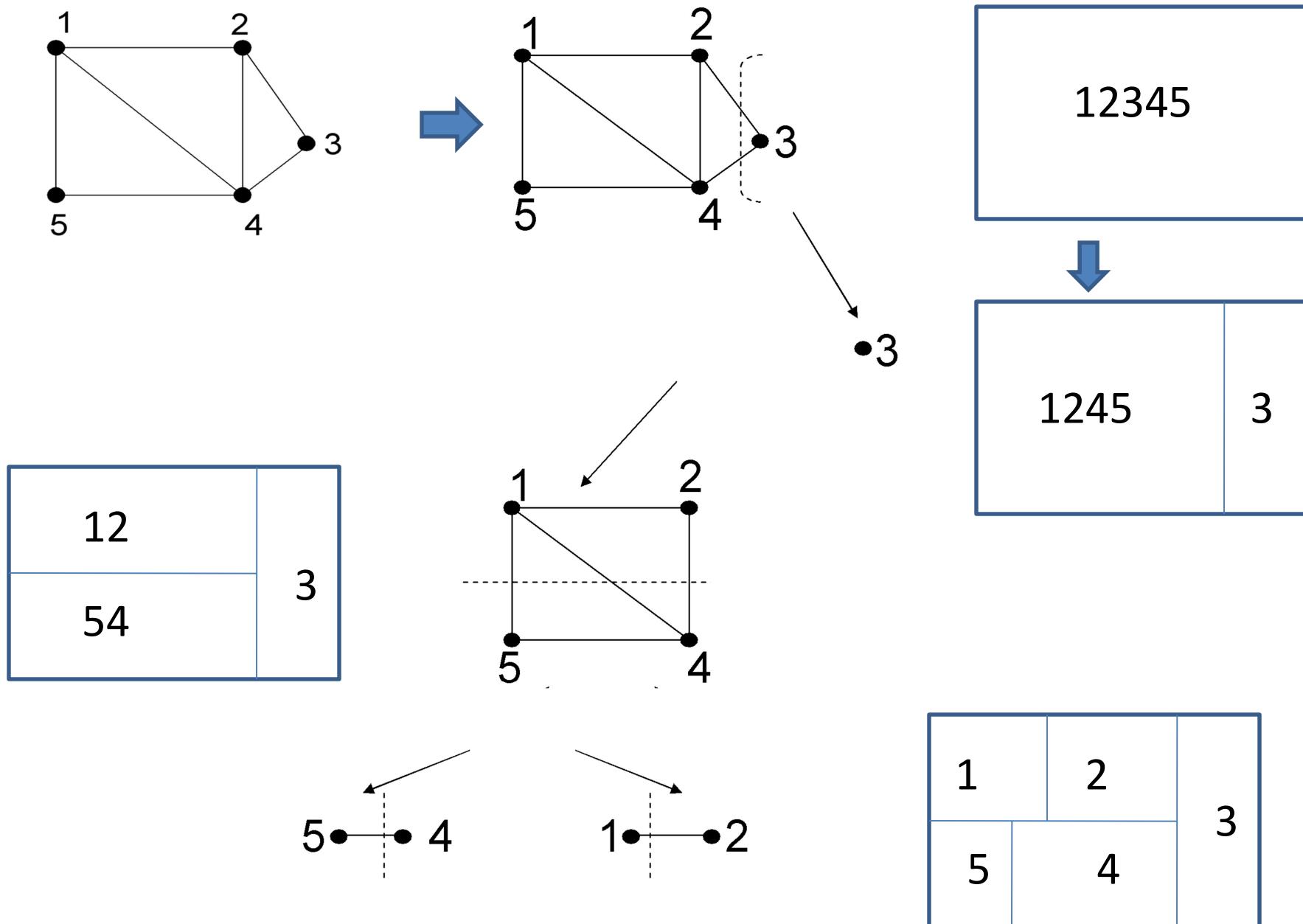


Adjacency Graph

Geometric Dual of Adjacency Graph

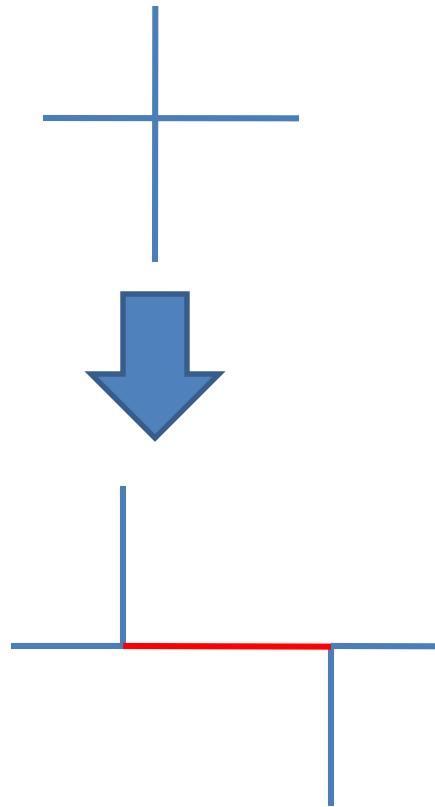


## Topology Generation : A simple example



# Rectangular Dual

Assumed that a rectangular floorplan contains no cross junctions



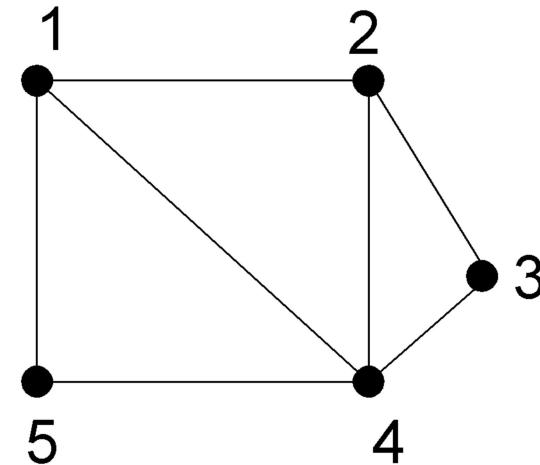
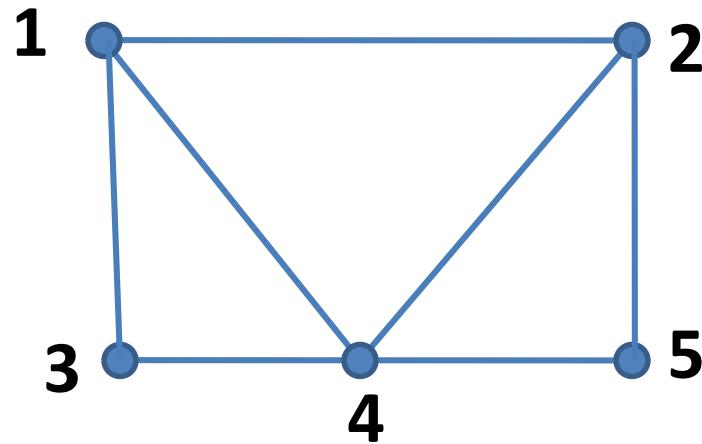
A cross junction is replaced by two T-junctions with another short edge

A graph is said to be planar if there exists some geometric representation of  $G$  which can be drawn on a plane such that no two of its edges intersect

Dual graph of a rectangular floorplan is a planar triangulated graph (PTG)

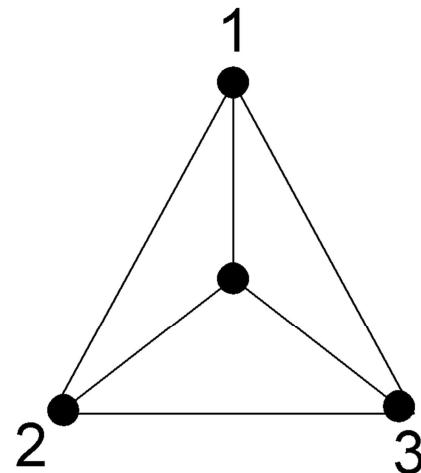
# Rectangular Dualization: Planar Triangulated Graph

Dual graph of a rectangular floorplan is a planar triangulated graph (PTG)



- Every face, except the exterior, is a triangle
- All internal vertices have degree  $\geq 4$
- all cycles that are not faces and the exterior face have length  $\geq 4$

# Rectangular Dualization: Planar Triangulated Graph



No corresponding rectangular floorplan exist

Only

Forbidden pattern in rectangular dual graph

Every dual graph of a rectangular floorplan is a planar triangulated graph (PTG)

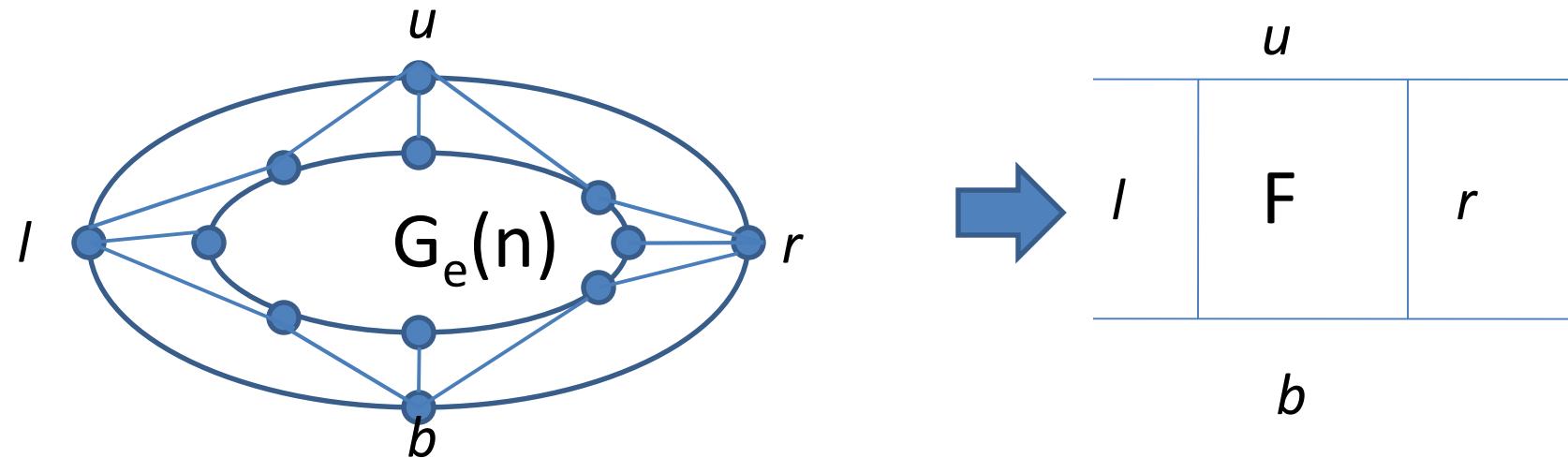
**BUT not every PTG corresponds to a rectangular floorplan**

**A rectangular floorplan exists iff the dual of the floorplan contains no complex triangle**

A rectangular graph is an adjacency graph

that yields a rectangular dual

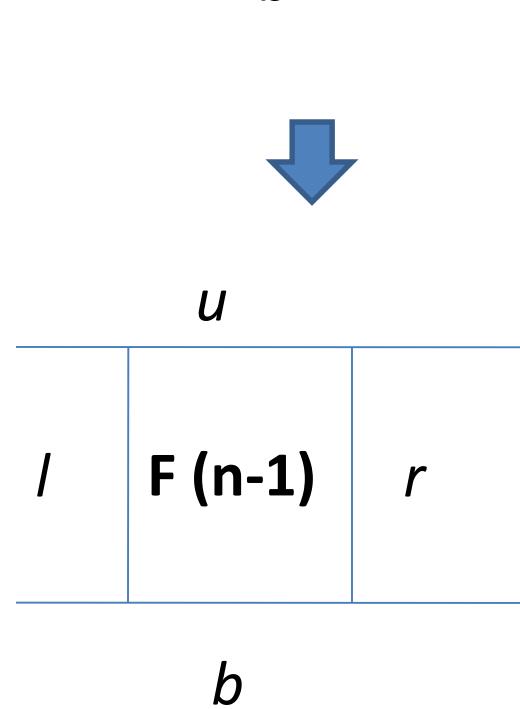
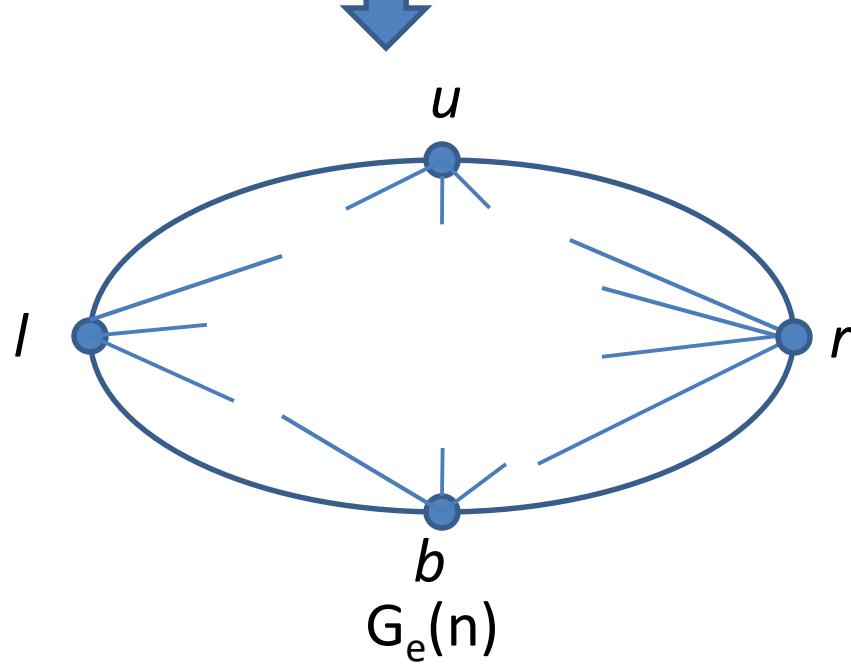
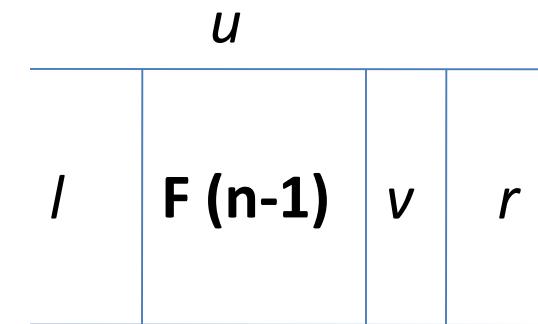
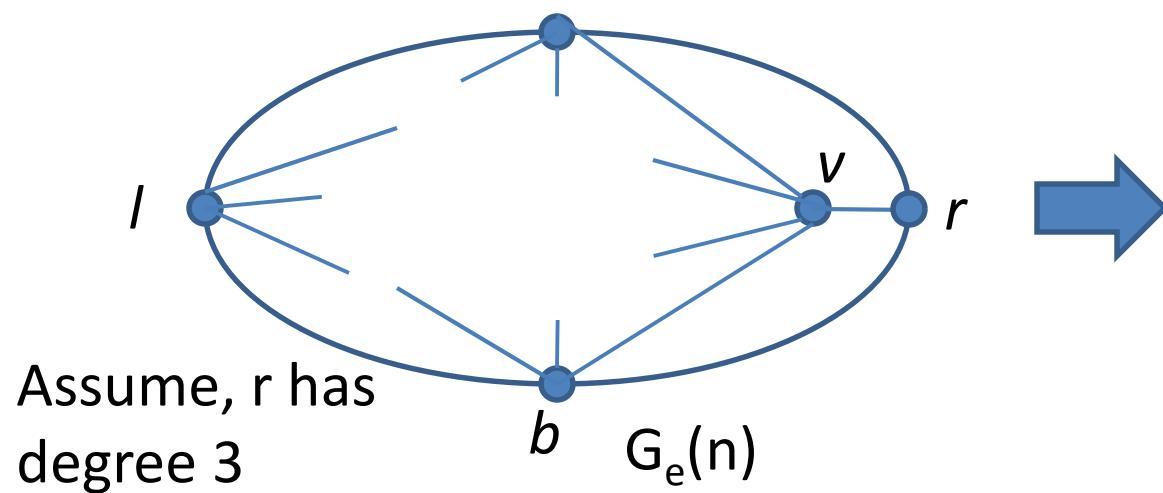
## Extended Dual



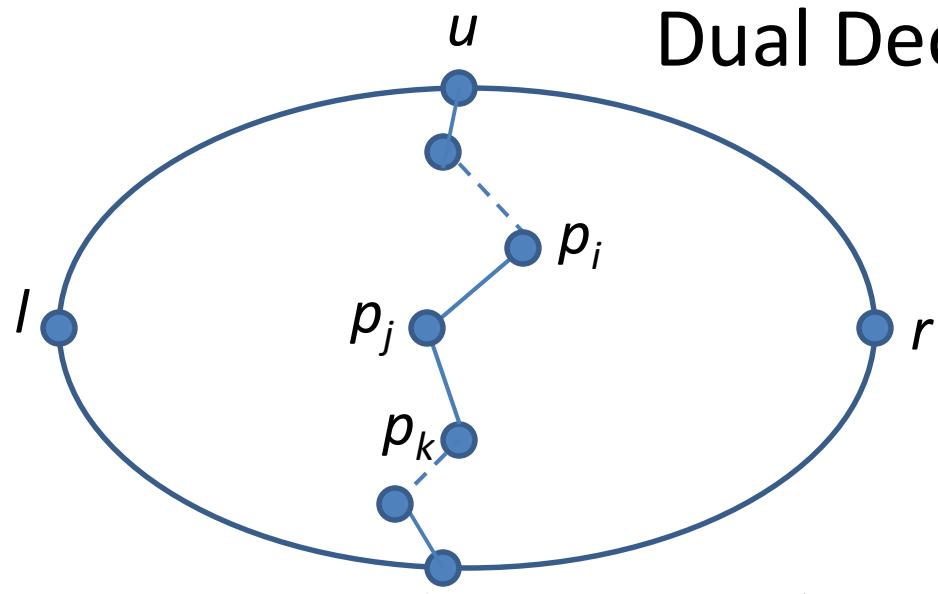
A rectangular floorplan exists if and only if the extended dual of the floorplan contains no complex triangles

Hypothesis: It is possible to generate a floorplan for any dual graph  $G_e(k)$  for  $k < n$

# Extended Dual



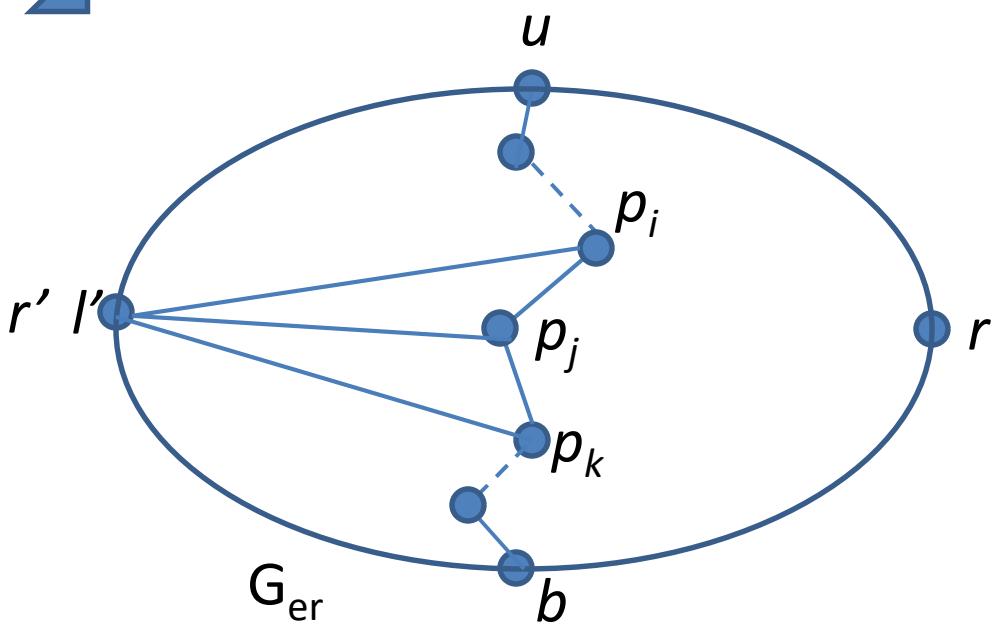
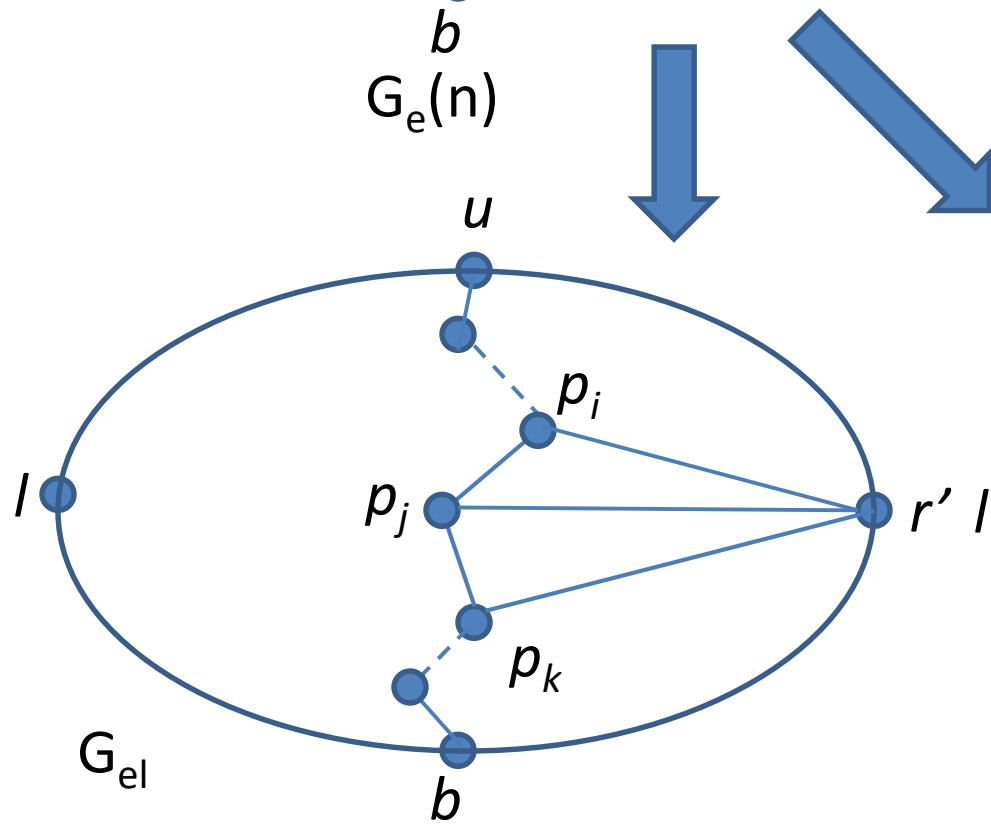
# Dual Decomposition



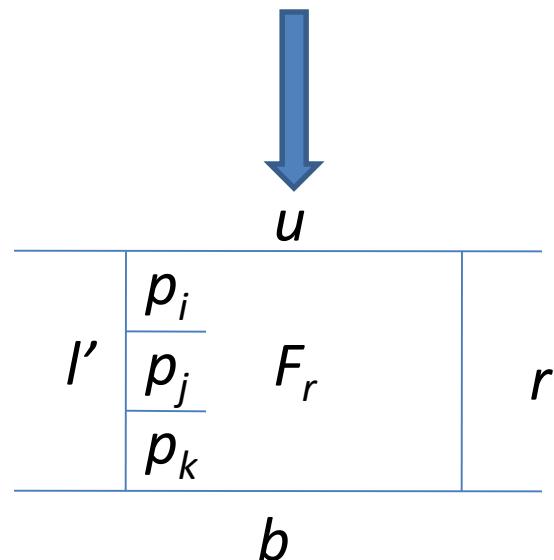
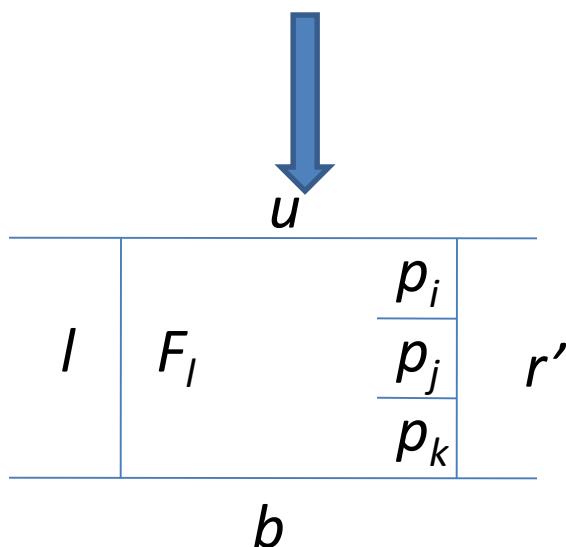
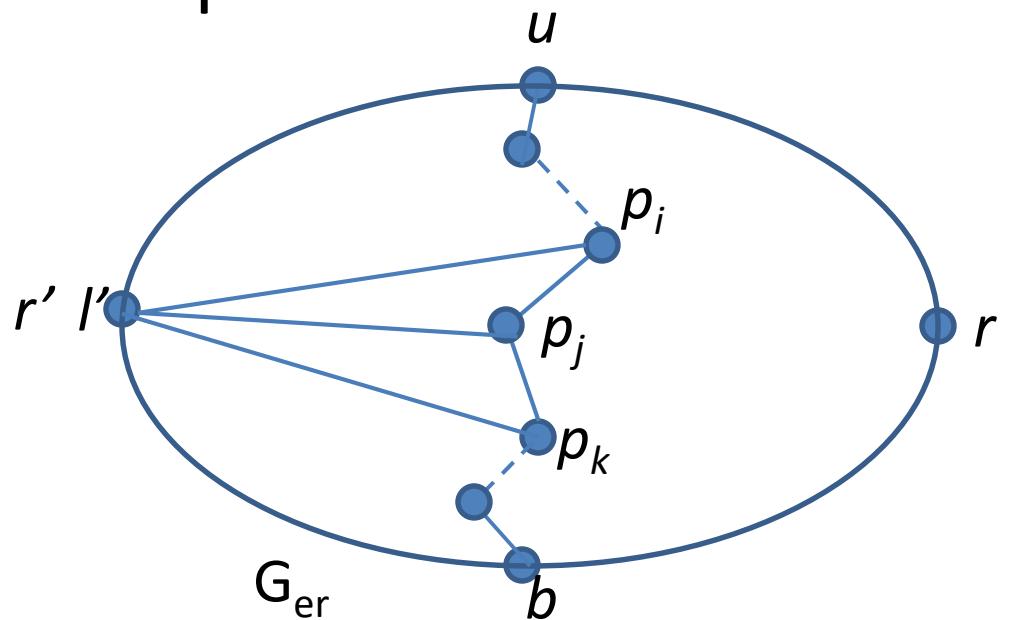
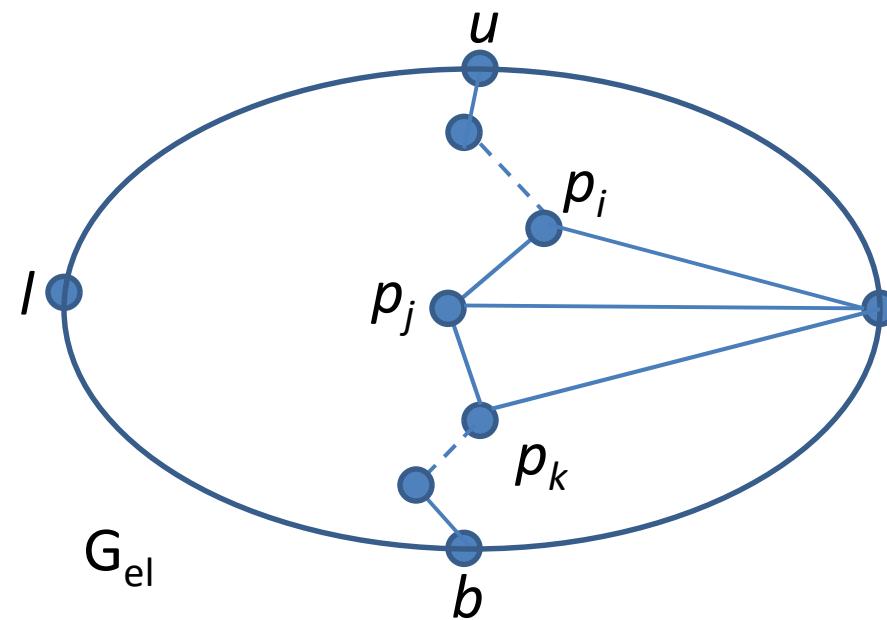
A path  $P_V = \{u=p_1, p_2, \dots, p_m=b\}$

- (i)  $p_2, \dots, p_{m-1} \notin \{r, u, l, b\}$
- (ii)  $(p_i, p_j) \notin G_e(n)$  for all distinct  $i, j$
- (iii)  $(p_i, l) \notin G_e(n)$  for some  $i$
- (iv)  $(p_j, r) \notin G_e(n)$  for some  $j$

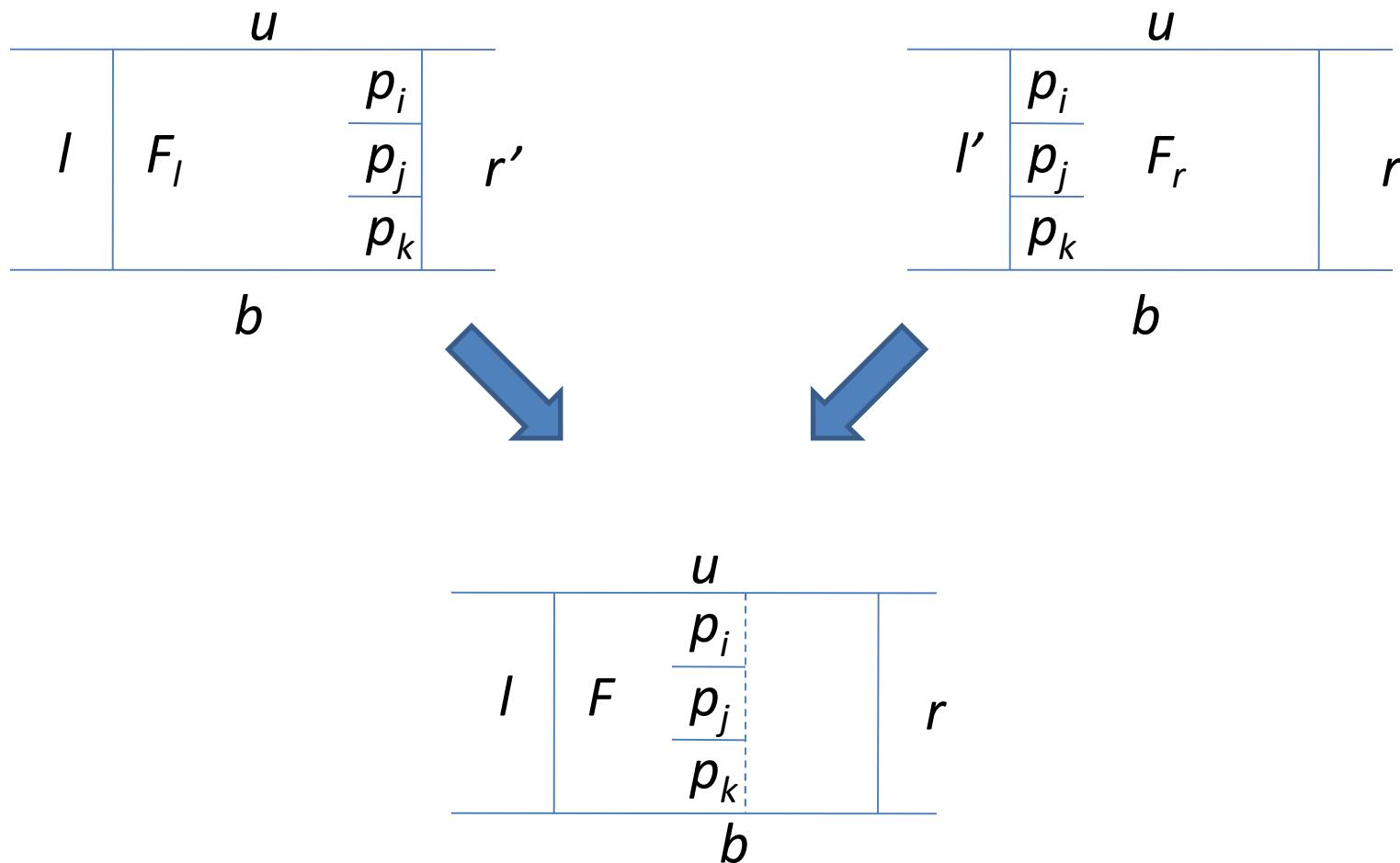
**Splitting path**



# Dual Decomposition



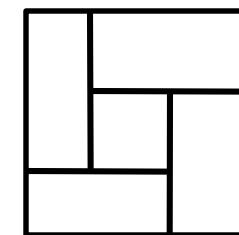
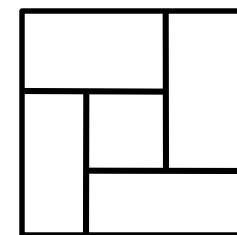
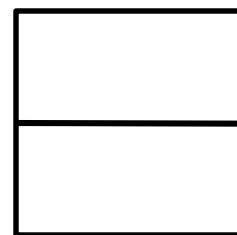
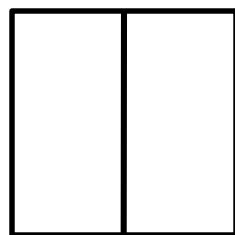
## Merging Floorplans



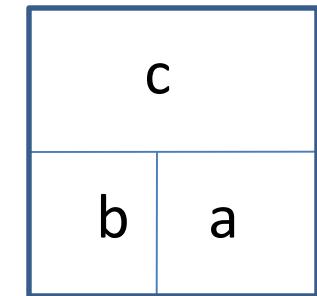
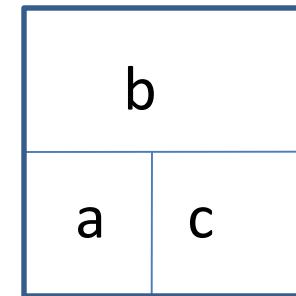
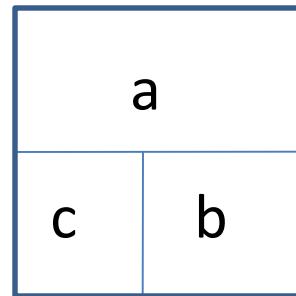
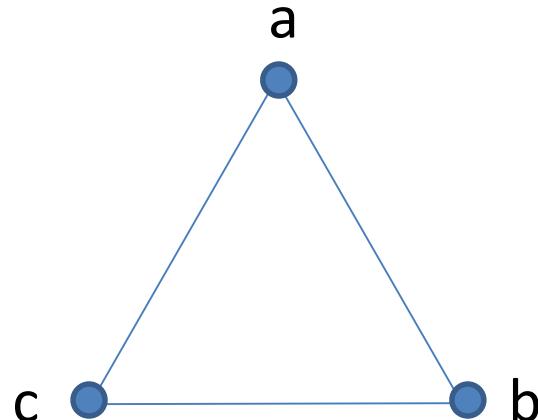
It is possible to generate a floorplan for any extended dual graph provided it contains no complex triangle

# Hierarchical Floorplan

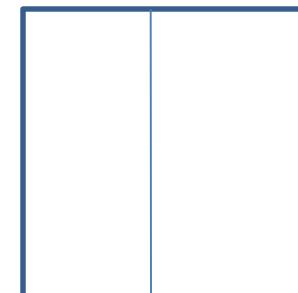
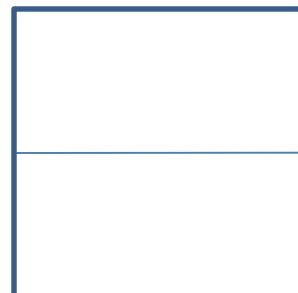
Four types of Internal nodes



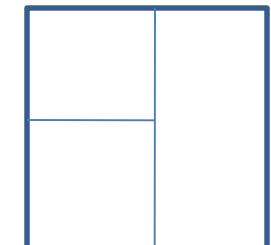
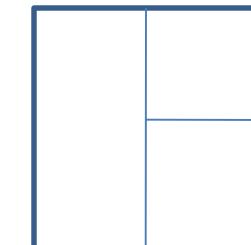
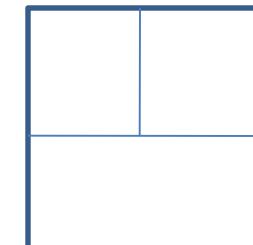
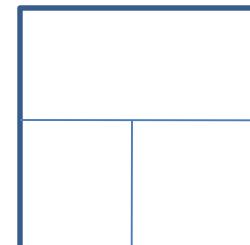
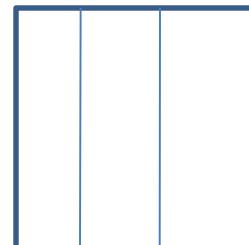
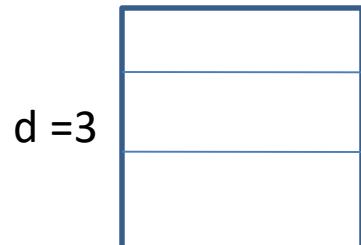
# Hierarchichal Approach



The number of possible floorplans increases exponentially with number of modules



$d = 2$



$d = 3$