

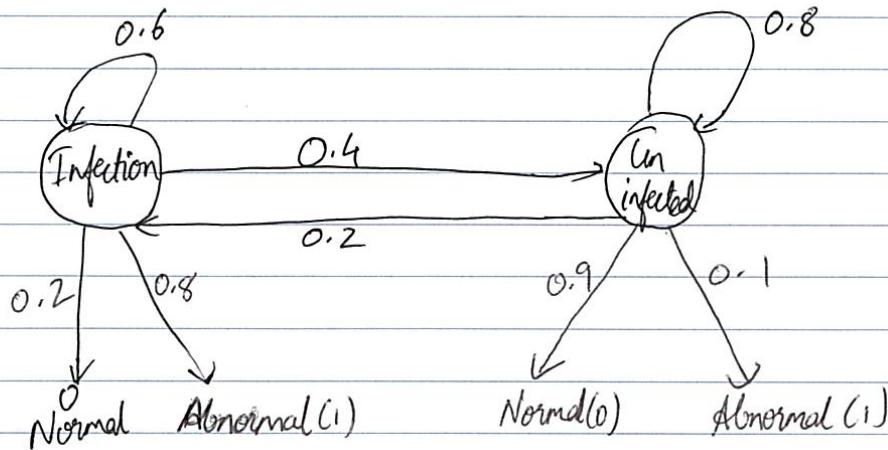
CSC591- HW2 (Group: ALL_12)

Anusha (amanur)

Kartik (kcshah)

Shailaja (smallic)

1



a) $N = \# \text{ of hidden states. Here } N=2.$

$M = \# \text{ of possible observations of all states. Here } M=2.$

There are 3 probabilities :

π_i : initial probability. There are N states so N parameters are required.

a_{ij} : Transitional probability. We can transition from any N state to any of N states. So, $N \times N$ i.e N^2 parameters are required.

$b_i(k)$: Emission or observation probability - Each of N state can emit M observations. So, $N \times M$ i.e NM parameters are required.

So, total number of parameters required are
 $N + N^2 + NM \Rightarrow 2 + 2^2 + 2 \times 2 \Rightarrow \underline{\underline{10}}$

P)

b)

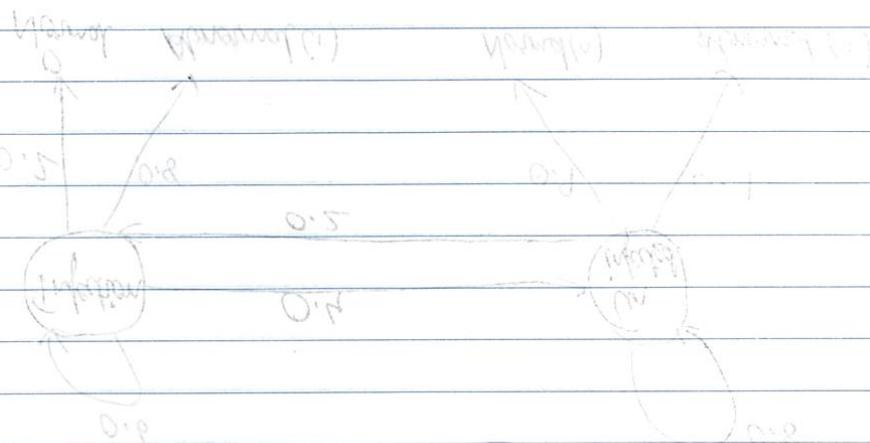
	State (i)	$\pi(i)$
Infected (1)		0.75
Uninfected (2)		0.25

Transition probability: $a_{ij} \Rightarrow$

	from 10	
	infected (1)	uninfected (2)
infected (1)	0.6	0.4
uninfected (2)	0.2	0.8

Observation or emission probability: $b_{ik}(k) = P(O_t=k | q_t=s_i)$

state observation	Normal (0)	Abnormal (1)
infected (1)	0.2	0.8
uninfected (2)	0.9	0.1



(2)

c) Observation : $\{ \begin{matrix} O_1 \\ O_2 \\ O_3 \end{matrix} \} = \begin{cases} 0 & \text{normal} \\ 1 & \text{abnormal} \end{cases}$

Forward algo : initial : $\alpha_1(i) = b_i(O_1) \cdot \pi_i$

at a time : $\alpha_{t+1}(i) = b_i(O_{t+1}) \cdot \sum_j a_{ji} \cdot \alpha_t(j)$

so,

$$\alpha_1(1) = b_1(O_1) \cdot \pi_1 \Rightarrow b_1(\text{normal}) \cdot \pi_1 \Rightarrow 0.2 \times 0.75 \Rightarrow 0.15$$

$$\alpha_1(2) = b_2(O_1) \cdot \pi_2 \Rightarrow 0.9 \times 0.25 \Rightarrow 0.225$$

$$\alpha_2(1) = b_1(O_2) \cdot \sum_{\substack{i=1, j=1 \\ i=2, j=2}} [a_{11} \cdot \alpha_1(1) + a_{21} \cdot \alpha_1(2)] \Rightarrow 0.8 \times [0.6 \times 0.15 + 0.2 \times 0.225] \Rightarrow 0.108$$

$$\alpha_2(2) = b_2(O_2) \cdot \sum_{\substack{i=1, j=1 \\ i=2, j=2}} [a_{12} \cdot \alpha_1(1) + a_{22} \cdot \alpha_1(2)] \Rightarrow 0.1 \times [0.4 \times 0.15 + 0.8 \times 0.225] \Rightarrow 0.024$$

$$\alpha_3(1) = b_1(O_3) \cdot \sum_{\substack{i=1, j=1 \\ i=2, j=2}} [a_{11} \cdot \alpha_2(1) + a_{21} \cdot \alpha_2(2)] \Rightarrow 0.8 \times [0.6 \times 0.108 + 0.2 \times 0.024] \Rightarrow 0.05568$$

$$\alpha_3(2) = b_2(O_3) \cdot \sum_{\substack{i=1, j=1 \\ i=2, j=2}} [a_{12} \cdot \alpha_2(1) + a_{22} \cdot \alpha_2(2)] \Rightarrow 0.1 \times [0.4 \times 0.108 + 0.8 \times 0.024] \Rightarrow 0.00624$$

$$P(\text{observed sequence}) = 0.05568 + 0.00624$$

$$(P(O|N) = \sum_{i=1}^3 \alpha_i(i)) = \underline{\underline{0.06192}}$$

$$b^T(3) = 1$$

$$b^T(1) = 1$$

$$P(O|N) = \min \{ b^T(1), b^T(2), b^T(3) \}$$

$$P(O|N) = \max \{ b^T(1), b^T(2), b^T(3) \}$$

d) Observation: $\{0, 1, 1\}$

Backward algo: $\sum_j \alpha_{ij} b_j(o_{t+1}) \cdot \beta_{t+1}(j)$

β_0 ,

$$\beta_3(1) = 1$$

$$\beta_3(2) = 1$$

$$\begin{aligned}\beta_2(1) &= [a_{11} \cdot b_1(o_3) \cdot \beta_3(1) + a_{12} \cdot b_2(o_3) \cdot \beta_3(2)] \\ &= [0.6 \times 0.8 \times 1 + 0.4 \times 0.1 \times 1] \Rightarrow 0.52\end{aligned}$$

$$\begin{aligned}\beta_2(2) &= [a_{21} \cdot b_1(o_3) \cdot \beta_3(1) + a_{22} \cdot b_2(o_3) \cdot \beta_3(2)] \\ &= [0.2 \times 0.8 \times 1 + 0.8 \times 0.1 \times 1] \Rightarrow 0.24\end{aligned}$$

$$\begin{aligned}\beta_1(1) &= [a_{11} \cdot b_1(o_2) \cdot \beta_2(1) + a_{12} \cdot b_2(o_2) \cdot \beta_2(2)] \\ &= [0.6 \times 0.8 \times 0.52 + 0.4 \times 0.1 \times 0.24] \Rightarrow 0.2592\end{aligned}$$

$$\begin{aligned}\beta_1(2) &= [a_{21} \cdot b_1(o_2) \cdot \beta_2(1) + a_{22} \cdot b_2(o_2) \cdot \beta_2(2)] \\ &= [0.2 \times 0.8 \times 0.52 + 0.8 \times 0.1 \times 0.24] \Rightarrow 0.1024\end{aligned}$$

$$P(\text{Observation}) = P(O|A) = \sum_{i=1}^N \pi_i \beta_0(i)$$

$$= \sum [\pi_1 \cdot \beta_1(1) + \pi_2 \cdot \beta_1(2)]$$

$$= 0.15 \times 0.2592 + 0.225 \times 0.1024$$

$$= \underline{\underline{0.06192}}$$

The forward & backward algorithm gives the same probability.

(5)

e) forward Backward $(\alpha_t(i), \beta_t(i))$

$$\alpha_1(1) \cdot \beta_1(1) = 0.15 \times 0.2592 \Rightarrow 0.03888$$

$$\alpha_1(2) \cdot \beta_1(2) = 0.225 \times 0.1024 \Rightarrow 0.02304$$

$$\alpha_2(1) \cdot \beta_2(1) = 0.108 \times 0.52 \Rightarrow 0.05616$$

$$\alpha_2(2) \cdot \beta_2(2) = 0.024 \times 0.24 \Rightarrow 0.00576$$

$$\alpha_3(1) \cdot \beta_3(1) = 0.05568 \times 1 \Rightarrow 0.05568$$

$$\alpha_3(2) \cdot \beta_3(2) = 0.00624 \times 1 \Rightarrow 0.00624$$

Since for every place $\alpha_t(1) \cdot \beta_t(1) > \alpha_t(2) \cdot \beta_t(2)$
 we can say the most likely sequence
 is infected, infected, infected.

$$\gamma(s) = \max\{\alpha_1(1) \cdot \beta_1(s), \alpha_1(2) \cdot \beta_1(s)\}$$

$$\gamma(s) = \max\{0.15 \times 0.2592, 0.225 \times 0.1024\} \times 0.8 \Rightarrow 0.035$$

$$\gamma(s) = [\alpha_1(1) \cdot \beta_1(s) \quad \alpha_1(2) \cdot \beta_1(s)] \cdot \gamma(s)$$

$$\gamma(s) = 0.15 \cdot \gamma(s) \Rightarrow 0.15 \times 0.035 \Rightarrow 0.00525$$

$$\gamma(s) = 0.15 \cdot \gamma(s) \Rightarrow 0.15 \times 0.035 \Rightarrow 0.00525$$

Answer: (0, 1, 1)

$$\alpha_t(i) = \max\{\alpha_{t-1}(1) \cdot \beta_{t-1}(1) \cdot \gamma(s), \alpha_{t-1}(2) \cdot \beta_{t-1}(2) \cdot \gamma(s)\}$$

f) Viterbi algorithm : $S_t(i) = \pi_i \cdot b_i(O_1)$
At time $t+1$: $S_{t+1}(j) = \max_{1 \leq i \leq N} [S_t(i) \cdot a_{ij}] \cdot b_j(O_{t+1})$

Observation: (0, 1, 1)

$$S_1(1) = \pi_1 \cdot b_1(O_1) \Rightarrow 0.75 \times 0.2 \Rightarrow 0.15$$

$$S_1(2) = \pi_2 \cdot b_2(O_1) \Rightarrow 0.25 \times 0.9 \Rightarrow 0.225$$

$$\begin{aligned} S_2(1) &= \max_{\substack{(t=1, j=1) \\ (i=1, 2)}} [S_1(1) \cdot a_{11}, S_1(2) \cdot a_{21}] \cdot b_1(O_2) \\ &= \max[0.15 \times 0.6, 0.225 \times 0.2] \times 0.8 \Rightarrow 0.072 \end{aligned}$$

$$\begin{aligned} S_2(2) &= \max_{\substack{(t=1, j=2) \\ (i=1, 2)}} [S_1(1) \cdot a_{12}, S_1(2) \cdot a_{22}] \cdot b_2(O_2) \\ &= \max[0.15 \times 0.4, 0.225 \times 0.8] \times 0.1 \Rightarrow 0.018 \end{aligned}$$

$$\begin{aligned} S_3(1) &= \max_{\substack{(t=2, j=1) \\ (i=1, 2)}} [S_2(1) \cdot a_{11}, S_2(2) \cdot a_{21}] \cdot b_{11}(O_3) \\ &= \max[0.072 \times 0.6, 0.018 \times 0.2] \times 0.8 \Rightarrow 0.03456 \end{aligned}$$

$$\begin{aligned} S_3(2) &= \max_{\substack{(t=2, j=2) \\ (i=1, 2)}} [S_2(1) \cdot a_{12}, S_2(2) \cdot a_{22}] \cdot b_{12}(O_3) \\ &= \max[0.072 \times 0.4, 0.018 \times 0.8] \times 0.1 \Rightarrow 0.00288 \end{aligned}$$

Since $S_1(2) > S_1(1)$ uninfected

This is because the value came from 2.

$S_2(1) > S_2(2)$ infected

This is because the highest value came from 1.

$S_3(1) > S_3(2)$ infected

This is because the highest value came from 1.

The most probable sequence will be
uninfected, infected, infected.

2 || HMM

- a) The output sequence of length 2 that cannot be generated by the given HMM is 5, 4.

This is because if S_1 gets 5 ~~then~~ first then there is no way of getting 4 since the S_1 has probability of $P(X=4) \Rightarrow 0$. Also, it cannot go to S_2 since there is no transition probability.

Now, for S_2 getting 5 first probability is $P(x=5) \neq 0$.
 So it cannot have 5 in first place.

So, this sequence cannot be obtained.

$$(b) \underset{1^{\text{st}} \text{ toss}}{P(S_1 | x=1)} = \frac{P(S_1 \cap x=1)}{P(x=1)} = \frac{0.01 \times 0.5}{0.33 \times 0.5} = \frac{0.005}{0.33 \times 0.5} = 0.0303$$

$$P(S_2 | x=1) = \frac{P(S_2 \cap x=1)}{P(x=1)} = \frac{0.32 \times 0.5}{0.33 \times 0.5} = \frac{0.16}{0.33 \times 0.5} = 0.9696$$

$$P(x=1) = P(x=1 | S_1) P(S_1) + P(x=1 | S_2) P(S_2)$$

$$= 0.01 \times 0.5 + 0.32 \times 0.5$$

$$= 0.33 \times 0.5$$

\therefore After 1st toss most likely by player 2.

2nd toss —

$$P(x=1) = P(x=1 | S_1) P(S_1) + P(x=1 | S_2) P(S_2)$$

$$= 0.01 \times \left[\frac{0.5 \times 1}{S_1 \underset{1^{\text{st}} \text{ toss}}{S_1} \underset{2^{\text{nd}} \text{ toss}}{S_2}} + \frac{0.5 \times 0.01}{S_2 \underset{1^{\text{st}} \text{ toss}}{S_2} \underset{2^{\text{nd}} \text{ toss}}{S_1}} \right] + 0.32 \left[\frac{0.5 \times 0}{S_1 \underset{S_1}{S_2}} + \frac{0.5 \times 0.99}{S_2 \underset{S_2}{S_2}} \right] = 0.15505$$

$$P(S_1 | x=1) = \frac{P(S_1 \cap x=1)}{P(x=1)} = \frac{0.00505}{0.15505} = 0.032$$

$$P(S_2 | x=1) = \frac{P(S_2 \cap x=1)}{P(x=1)} = \frac{0.15}{0.15505} = 0.9674$$

2nd toss most likely by player 2

But if we look at the probability of player 2's decreases while player 1's increases. And if the observation increases till $20,60^{2020}$, player 2's probability of getting $x=1$ decreases while player 1's increases.

Hence, without knowing if the values converge or not, it is most likely that S_1 is the hidden state.

(9)

c) Viterbi : $S_t(i) = \pi_i \cdot b_i(O_i)$
 $S_{t+1}(j) = \max_{1 \leq i \leq N} [S_t(i) \cdot a_{ij}] \cdot b_j(O_{t+1})$

Observation : (1,1)

$S_1(1) = \pi_1 \cdot b_1(O_1) \Rightarrow 0.5 \times 0.01 \Rightarrow 0.005$

$S_1(2) = \pi_2 \cdot b_2(O_1) \Rightarrow 0.5 \times 0.32 \Rightarrow 0.16$

$S_2(1) = \max_{\substack{(t=1, j=1) \\ (i=1, 2)}} [S_1(i) \cdot a_{11}, S_1(2) \cdot a_{21}] \cdot b_1(O_2)$
 $= \max[0.005 \times 1, 0.16 \times 0.01] \times 0.01$
 $\Rightarrow 0.00005$

$S_2(2) = \max_{\substack{(t=1, j=2) \\ (i=1, 2)}} [S_1(i) \cdot a_{22}, S_1(1) \cdot a_{12}] \cdot b_2(O_2)$
 $= 0.16 \times 0.99 \times 0.32 \Rightarrow 0.0506$

Since $S_1(2) > S_1(1)$ we select S_2 . Since the high value came from S_2 .

& $S_2(2) > S_2(1)$ we select S_2 since again the value came from S_2 .

The most likely sequence will be S_2, S_2 if this is the observation given.

Q) Applying Viterbi algorithm to the following sequence of observations {1, 1, 3}

Q) Applying Viterbi algorithm to the following sequence of observations {1, 1, 3}

d) Output sequence $\{1, 1, 3\}$

→ As we found in "c" for $\{1, 1\}$ observation the states were S_2, S_2 hereby, in this as well it will be S_2, S_2 for first two states.

→ However, for 3 we have shown our calculations & the state which we get is below.

Using Viterbi

$$\begin{aligned} \delta_3(1) &= \max [S_2(1) \cdot a_{11}, S_2(2) \cdot a_{21}] \cdot b_1(O_3) \\ (\text{t=2, j=1}) &= \max [0.00005 \times 1, 0.0506 \times 0.01] \times 0.3 \\ &= 0.00015 \end{aligned}$$

$$\begin{aligned} \delta_3(2) &= \max [S_2(2) \cdot a_{22}, 0] \cdot b_2(O_3) \\ (\text{t=2, j=2}) &= 0.0506 \times 0.99 \times 0.14 \\ &= 0.00701 \end{aligned}$$

Since $\delta_3(2) > \delta_3(1)$ & the highest value came from S_2 we see S_2 as output state.

So, for the given output sequence $\{1, 1, 3\}$ we see S_2, S_2, S_2 .

$$2^1(3) = 0.2 \cdot 0.14 = 0.028$$

$$2^1(1) = 0.2 \cdot 0.01 = 0.002$$

$$\text{Minimum } (0.002, 0)$$

$$2^1(1) = \text{Minimum}(0.002, 0) = 0$$

$$\therefore \text{Answer} : 2^1(1) = 0 \cdot 2 \cdot 0.01 = 0$$

Q3 (a) optimal policy -

s1: Right

s2: right

s3: right

s4: Right

s5: left

(b) Using the Bellman equation -

$$V_{ss} = 2 + 0.8(1 * V_{sy}) \quad (1)$$

$$V_{sy} \text{ right} - V_{sy} = 20 + 0.8(1 * V_{ss}) \quad (2).$$

Solving (1) & (2) -

$$\begin{aligned} 0.8 V_{sy} &= V_{ss} - 2 \\ V_{sy} &= \frac{0.8 V_{ss} + 20}{0.8} \\ 0 &= 0.36 V_{ss} - 18 \end{aligned} \Rightarrow V_{ss}^* = \frac{18}{0.36} = 50.$$

$$\boxed{V_{ss}^* = 50}$$

$$(l) \alpha = 0.5$$

Start from s_3^0

$$Q(s_3, L) = (1-\alpha) Q(s_3, L) + \frac{\alpha}{0.5} \left[r(s_3, L) + \frac{Q(s_2)}{2} \right]$$

$$= 1$$

$$Q(s_2, L) = (1-\alpha) Q(s_2, L) + 0.5 \left[2 + \frac{Q(s_1)}{0} \right]$$

$$= 1$$

$$Q(s_1, R) = (1-\alpha) Q(s_1, R) + 0.5 \left[1 + 0.8 \max \left[Q(s_2, R), \underline{Q(s_2, L)} \right] \right]$$

$$= 0.5 [1 + 0.8 \times 1] = 0.9$$

$$Q(s_2, L) = (1-\alpha) Q(s_2, L) + 0.5 \left[2 + 0.8 \times 0.9 \right]$$

$$= 0.5 \times 1 + 0.5 (2 + 0.72)$$

! gets updated .

	L	R
s_1		0.9
s_2	1	
s_3	1	
s_4		
s_5		

<u>10 pairs</u>	
(s_3, L)	(s_2, L)
(s_3, L)	(s_1, R)
(s_2, L)	(s_1, R)
(s_1, R)	(s_2, L)
(s_2, L)	

	V_{S1}	V_{S2}	V_{S3}	V_{S4}	V_{S5}
t	20	30	20	30	10
$t+1$	25	21.2	23.6	36	26

$$Q(S_1, R) = 1 + 0.8(1 \times 30) = 25 \quad R$$

$$Q(S_2, R) = 1 + 0.8(0.7 \times 20 + 0.3 \times 30) = 19.4 \quad \left. \begin{array}{l} \\ 21.2 \end{array} \right\}$$

$$Q(S_2, L) = 1 + 0.8(0.4 \times 30 + 0.6 \times 20) = 21.2$$

$$Q(S_3, R) = 2 + 0.8(0.8 \times 20 + 0.2 \times 30) = 19.6 \quad \left. \begin{array}{l} \\ 23.6 \end{array} \right\}$$

$$Q(S_3, L) = 2 + 0.8(0.7 \times 30 + 0.3 \times 20) = 23.6$$

$$Q(S_4, R) = 2 + 0.8(0.5 \times 30 + 0.5 \times 10) = 36 \quad \left. \begin{array}{l} \\ 36 \end{array} \right\}$$

$$Q(S_4, L) = 3 + 0.8(0.6 \times 20 + 0.4 \times 30) = 22.2$$

$$Q(S_4, L) = 3 + 0.8(1 \times 30) = 26 - L$$

$$Q(S_5, L) = 2 + 0.8(1 \times 30) = 26 - L$$

(e) (1) Based on the value function at time t —

The values of Q for each state is from part (d)

$$Q(S_1, R) = 25 \quad - \quad S_1: \text{Right}$$

$$Q(S_2, R) = 19.4 \quad \left. \begin{array}{l} \\ \end{array} \right\} 21.2 \quad - \quad S_2: \text{Left}$$

$$Q(S_2, L) = 21.2$$

$$Q(S_3, R) = 19.6 \quad \left. \begin{array}{l} \\ \end{array} \right\} 23.6 \quad - \quad S_3: \text{Left}$$

$$Q(S_3, L) = 23.6$$

$$Q(S_4, R) = 36 \quad \left. \begin{array}{l} \\ \end{array} \right\} 36 \quad - \quad S_4: \text{Right}$$

$$Q(S_4, L) = 22.2 \quad - \quad S_5: \text{Left}$$

$$Q(S_5, L) = 26$$

\Rightarrow The optimal policy π —

$(S_1: \text{Right})$, $(S_2: \text{Left})$, $(S_3: \text{Left})$, $(S_4: \text{Right})$,
 $(S_5: \text{Left})$.

(2) Based on value function at time $t+1$ —

$$Q(S_1, R) = 17.96 \quad - \quad \text{Right}$$

$$Q(S_2, R) = 19.3 \quad \left. \begin{array}{l} \\ \end{array} \right\} 20.78 \quad - \quad \text{Left}$$

$$Q(S_2, L) = 20.78 \quad Q(S_2, R) = 1 + 0.8(0.3 * 21.2 + 0.7 * 23.6) = 19.3$$

$$Q(S_3, R) = 22.86 \quad \left. \begin{array}{l} \\ \end{array} \right\} 22.86 \quad - \quad \text{Right}$$

$$Q(S_3, L) = 19.5 \quad \text{Similarly done for rest states.}$$

$$Q(S_4, R) = 44.8 \quad \left. \begin{array}{l} \\ \end{array} \right\} 44.8 \quad - \quad \text{Right}$$

$$Q(S_4, L) = 25.8 \quad - \quad \text{Left}$$

$$Q(S_5, R) = 30.8 \quad - \quad \text{All Q values are calculated at t+1}$$

\Rightarrow The optimal policy is —

(S1: Right), (S2: Left), (S3: Right), (S4: Right), (S5: Left)

(3) Yes, the two policies are different.