

GREEDY ALGORITHM

- * Greedy means greed!
- * "Greedy algorithms is the problem solving technique where we make the locally optimum choice at each stage & hope to achieve a global optimum."
 - this means, at every step we can make a choice that looks best at the moment & it will lead to an overall optimal solution.

Ex → Company HR

Round 1 : Online Test (1000)

Round 2 : Interview (100)

Round 3 : Interview HR (20)

10 Hired

or

Online Test (1000)

Interview (1000)

Interview HR (1000)

or

or

Greedy Approach

Fair Approach

* Pros

→ Simple & Easy

→ Good Enough Time Complexity.

* Cons

→ A lot of time

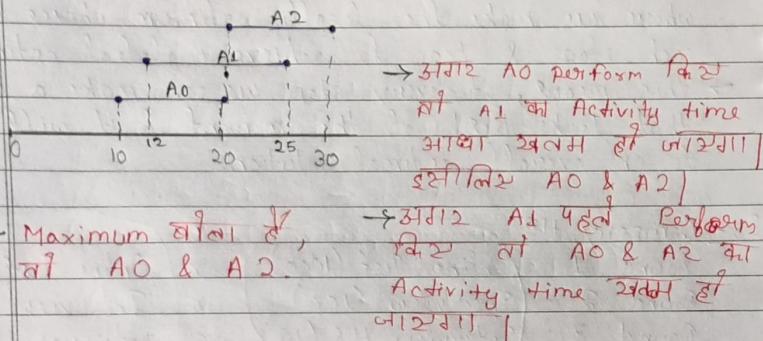
→ Many times optimum
is not achieved.

- ★ ACTIVITY SELECTION — You are given n activities with their start & end times. Select the maximum number of activities that can be performed by a single person, assuming that a person can only work on a single activity at a time.
Activities are sorted according to end time.

- * Sample Input: $[10, 12, 20] \leftarrow \text{start}$
 $[20, 25, 30] \leftarrow \text{end}$

* Sample Output: ② (AO & A2)

↓
max. 2 Activities perform की सकता है A0 & A2



* Approach

1. end time sort
 2. we have to select 1st activity i.e. A0.
 3. for ($i = 1$ to n)

non-overlapping (disjoint)

start time \geq last chosen activity end time

count ++;

Input → start = [1, 3, 0, 5, 8, 5]
→ end = [2, 4, 6, 7, 9, 9]

Output \Rightarrow count = 4
ans = A0, A1, A3, A9

II Time Complexity = $O(n)$

```

⇒ public static void main (String args [ ]) {
    int start [] = {1, 3, 0, 5, 8, 5};
    int end [] = {2, 4, 6, 7, 9, 9};
    //end time basis sorted
    int maxAct = 0;
    ArrayList < Integer > ans = new ArrayList < > ();
    //1st activity
    maxAct = 1;
    ans.add (0);
    int lastEnd = end [0];
    for (int i = 1; i < end.length; i++) {
        if (start [i] >= lastEnd) {
            //activity select
            maxAct++;
            ans.add (i);
            lastEnd = end [i];
        }
    }
}

```

```

System.out.println("max activities = " + maxAct);
for (int i = 0; i < ans.size(); i++) {
    System.out.print("A" + ans.get(i) + " ");
}
System.out.println();

```

Output: max activities = 4
A0 A1 A3 A4

If in a question, there wasn't given that activities are sorted then we have to write the sorted part also.

In Java, Comparator is an interface for sorting java objects.

// Time Complexity = $O(n \log n)$

```
public static void main (String args []) {
    int start [] = {1, 3, 0, 5, 8, 5};
    int end [] = {2, 4, 6, 7, 9, 3};
    // sorting
    int activities [][] = new int [start.length] [3];
    for (int i = 0; i < start.length; i++) {
        activities [i] [0] = i;
        activities [i] [1] = start [i];
        activities [i] [2] = end [i];
    }
}
```

// Lambda function → shorthand of any function
`Arrays.sort(activities, Comparator.comparingDouble (o -> o [0]));`

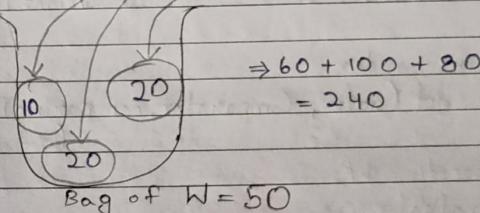
```
// end time basic sorted
int maxAct = 0;
ArrayList < Integer > ans = new ArrayList < > ();
// 1st activity
maxAct = 1;
ans.add (activities [0] [0]);
int lastEnd = activities [0] [2];
for (int i = 1; i < end.length; i++) {
    if (activities [i] [1] >= lastEnd) {
        // activity select
        maxAct++;
        ans.add (activities [i] [0]);
        lastEnd = activities [i] [2];
    }
}
```

```
System.out.println ("max activities = " + maxAct);
for (int i = 0; i < ans.length(); i++)
    System.out.print (" " + ans.get (i));
System.out.println ();
```

It is basically
a bag.

★ FRACTIONAL KNAPSACK - Given the weights & values of N items, put these items in a knapsack of capacity W to get the maximum total value in the knapsack.

* Sample: value = [60, 100, 120]
 Input: weight = [10, 20, 30]
 $W = 50$
 Output: 20 from 30.



* Approach:

* We have to think like Greedy Businessman,
 $\text{weight} \downarrow \times \text{value} \uparrow \rightarrow \text{ratio: value} / \text{weight}$

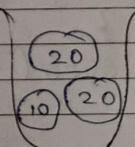
* Weight: [10, 20, 30], $W = 50$

* Value: [60, 100, 120]

$$\text{ratio: } \left[\frac{60}{10}, \frac{100}{20}, \frac{120}{30} \right] = [6; 5, 4] \quad \because \text{Capacity} \times 4 = 20 \times 4 = 80$$

$$\rightarrow \text{Capacity} = 50 \ 40 \ 20 \ 0$$

$$\rightarrow \text{value} = 60 + 100 + 80 = 240$$



```

→ public static void main (String args []) {
    int val [] = { 60, 100, 120 };
    int weight [] = { 10, 20, 30 };
    int W = 50;
    double ratio [][] = new double [val.length] [2];
    // 0th col → idx; 1st col → ratio
    for (int i=0; i<val.length; i++) {
        ratio [i][0] = i;
        ratio [i][1] = val [i] / (double) weight [i];
    }
    // ascending order
    Arrays.sort (ratio, Comparator.comparingDouble
        (o → o[1]));
    int capacity = W;
    int finalVal = 0;
    for (int i=ratio.length-1; i≥0; i--) {
        int idx = (int) ratio [i][0];
        if (capacity ≥ weight [idx]) { // include full item
            finalVal += val [idx];
            capacity -= weight [idx];
        } else {
            // include fractional item
            finalVal += (ratio [i][1] * capacity);
            capacity = 0;
            break;
        }
    }
    System.out.println ("final value = " + finalVal);
}

```

* Output: final value = 240

★ MIN ABSOLUTE DIFFERENCE PAIRS — Given two arrays A & B of equal length n. Pair each element of array A to an element in array B, such that sum S of absolute differences of all the pairs is minimum.

* Sample Input: A = [1, 2, 3] case ① |1-2| + |2-1| + |3-3| = 1+1+0=2
 ← sample output → B = [2, 1, 3] case ② |1-3| + |2-1| + |3-2| = 2+1+1=4
 Output: Ans: 0. case ③ |1-1| + |2-2| + |3-3| = 0+0+0=0

≠ Minimum ⇒ Case ③ i.e., 0.

* Greedy Approach

$a \longleftrightarrow b$
 Min. if Number close first तरीका
 absolute difference ↓

* A = [4, 1, 8, 7], B = [2, 3, 6, 5]
 A = [1, 4, 7, 8], B = [2, 3, 5, 6]

* Pair with the closest numbers means small no. with small one & big one with big one.

i.e., $|1-2| + |4-3| + |7-5| + |8-6|$

$$| -1 | + | +1 | + | 2 | + | 2 |$$

$$1 + 1 + 2 + 2$$

6

// Time Complexity, $O(n \log n)$

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```

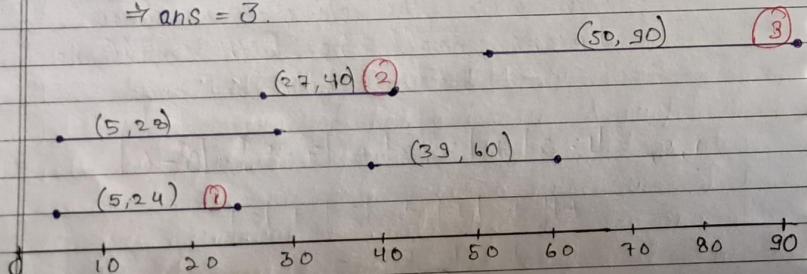
⇒ public static void main (String args[]) {
    int A [] = { 4, 1, 8, 7 };
    int B [] = { 2, 3, 6, 5 };
    Arrays.sort (A);
    Arrays.sort (B);
    int minDiff = 0;
    for (int i=0; i < A.length; i++) {
        minDiff += Math.abs (A[i] - B[i]);
    }
    System.out.println ("min absolute diff of pairs = " + minDiff);
}

```

★ MAX LENGTH CHAIN OF PAIRS - We are given n pairs of numbers. In every pair, the first number is always smaller than the second number. A pair (c, d) can come after pair (a, b) if $b < c$. Find the longest chain which can be formed from a given set of pairs.

Ex. pairs = $(5, 24)$ $(39, 60)$ $(5, 28)$ $(27, 40)$ $(50, 90)$

\Rightarrow ans = 3.



\therefore ans = 3 set of pairs.

// Time Complexity $\rightarrow O(n \log n)$

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```

⇒ public static void main (String args[]) {
    int pairs [] [] = {{5, 24}, {39, 60}, {5, 28}, {27, 40}, {50, 90}};
    Arrays.sort (pairs, Comparator.comparingDouble (o -> o[1]));
    int chainLen = 1;
    int chainEnd = pairs [0][1]; // Last selected pair end
    for (int i=1; i < pairs.length; i++) {
        if (pairs [i][0] > chainEnd) {
            chainLen++;
            chainEnd = pairs [i][1];
        }
    }
    System.out.println ("max length of chain = " + chainLen);
}

```

★ INDIAN COINS - We are given an infinite supply of denominations [1, 2, 5, 10, 20, 50, 100, 500, 2000]. Find min number of coins / notes to make change for a value V .

* Input: $V = 121$

* ans : 3 ($100+20+1$)

* Input: 590

* ans : 4 ($500+50+20+20$)

* Approach

① Sort descending [2000, 500, 100, 50, 20, 10, 5, 2, 1]

②

\Rightarrow Amount = 590 20 40 20 0

② count = 0

for (int i=0; i < n; i++)

if (coin[i] < amount)

while (coin[i] < amount)

count++
amount = amount - coin[i];

We will learn sorting of objects
in below code:

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```
⇒ static class Job {
    int deadline;
    int profit;
    int id;
    public Job (int i, int d, int p) {
        id = i;
        deadline = d;
        profit = p;
    }
}
public static void main (String args[]) {
    int jobsInfo [][] = {{4,20}, {1,10}, {1,40}, {1,30}};
    indexes → 0   1   2   3
    ArrayList <Job> jobs = new ArrayList <>();
    for (int i=0; i<jobsInfo.length; i++) {
        jobs.add (new Job (i, jobsInfo [i][0], jobsInfo [i][1]));
    }
    Collections.sort (jobs, (obj1,obj2)→obj2.profit - obj1.profit);
    ArrayList <Integer> seq = new ArrayList <>();
    int time = 0;
    for (int i=0; i<jobs.size(); i++) {
        Job curr = jobs.get (i);
        if (curr.deadline > time)
            seq.add (curr.id);
        time++;
    }
    System.out.println ("max jobs=" + seq.size ());
    for (int i=0; i<seq.size(); i++)
        System.out.print (seq.get (i) + " ");
    System.out.println ();
}
```

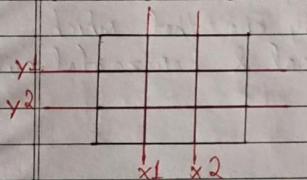
Output: max jobs=2
2 0

Min cost to cut
board into squares.

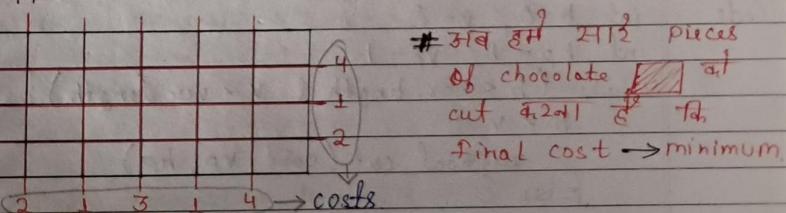
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★ CHOCOLA PROBLEM — We are given a bar of chocolate composed of $m \times n$ square pieces. One should break the chocolate into single squares. Each break of a part of the chocolate is charged a cost expressed by a positive integer. This cost does not depend on the size of the part that is being broken but only depends on the line the break goes along.

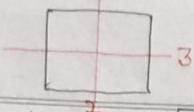
Let us denote the costs of breaking along consecutive vertical lines with x_1, x_2, \dots, x_{m-1} & along horizontal lines with y_1, y_2, \dots, y_{n-1} . Compute the minimal cost of breaking the whole chocolate into single squares.



* Approach $n=4, m=6$



- ① It's necessary to have $(n-1) + (m-1)$ cuts.
- ② after each cut → chocolate pieces ↑
 * expensive pieces की पहचान करें।
 * cheap pieces की पहचान करें।



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Case 1: Vertical cuts \rightarrow [3] [3] \rightarrow cost = $2 + 3 + 3$
 $= 8$

Case 2: Horizontal cuts \rightarrow [2] [2] \rightarrow cost = $3 + 2 + 2$
 $= 7$

- * cost of vertical cuts \rightarrow horizontal pieces * x.
- * cost of horizontal cuts \rightarrow vertical pieces * y.

$$\begin{aligned} \text{* Vertical cuts} &= [4, 3, 2, 1, 1] \\ \text{* horizontal cuts} &= [4, 2, 1] \end{aligned}$$

* Now, we will compare one by one that which cut is expensive in vertical & horizontal cuts, we cut them first.

$$\rightarrow \text{Total Cuts} = mxn = 24$$

$$\rightarrow \text{Total Costs} = 42$$

Pseudo \Rightarrow hor, ver \rightarrow Descending sort

$$\text{inf } h = 0, v = 0$$

$$\text{inf } hp = 1, vp = 1$$

while ($h < \text{hor. length}$ & $v < \text{ver. length}$) {

$n < v$ {

 vertical cut (cost, vp, hp)

 } else {

 horizontal cut

}

horizontal

Vertical

Total min cost

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\Rightarrow public static void main (String args[]) {

 inf h = 4, m = 6;

 Integer costVer[] = {2, 1, 3, 1, 4}; //m-1

 Integer costHor[] = {4, 1, 2}; //n-1

 Array8.sort (costVer, Collections.reverseOrder());

 Array8.sort (costHor, Collections.reverseOrder());

 inf h = 0, v = 0;

 inf hp = 1, vp = 1;

 inf cost = 0;

 while ($h < \text{costHor.length}$ & $v < \text{costVer.length}$) {

 if ($\text{costVer}[v] \leq \text{costHor}[h]$) { //Horizontal cut
 $\text{cost} += (\text{costHor}[h] * vp);$

 hp++;

 vp++;

 } else { //Vertical cut
 $\text{cost} += (\text{costVer}[v] * hp);$

 vp++;

 v++;

}

 while ($h < \text{costHor.length}$) {

 cost += ($\text{costHor}[h] * vp$);

 hp++;

 h++;

}

 while ($v < \text{costVer.length}$) {

 cost += ($\text{costVer}[v] * hp$);

 vp++;

 v++;

}

 System.out.println ("min cost of cuts = " + cost); //42

★ COMPARATORS IN JAVA

It is an interface that is used to order the objects of user-defined classes.

- * A comparator object is capable of comparing two objects of the same class.
- * This interface is present in the `java.util` package & contains 2 methods `compare (Object obj1, Object obj2)` & `equals (Object element)`.
- * Method of Collections class for sorting List elements is used to sort the elements of list by the given comparator.
 \Rightarrow `public void sort (List list, Comparator<Object> c)`

★ LAMBDA EXPRESSIONS IN JAVA

- * A lambda expression is a short block of code which takes in parameters & returns a value.
- * Lambda expressions are similar to methods, but they do not need a name & they can be implemented right in the body of a method.

Format: `Comparator<ClassName> comparator = Comparator<ClassName>.comparing (o → o.property);`

Example: `Comparator<Student> comparator = Comparator<Student>.comparing (o → o.age);`
`Collections.sort (students, comparator);`

Ques 1. Maximum Balanced String Partitions - We have balanced string str of size N with an equal number of L & R, the task is to find a maximum number X, such that a given string can be partitioned into X balanced substrings. A string is called to be balanced if the number of 'L's' in the string equals the number of 'R's'.

* Input - "LRRRRLLRLLRL"

* Output - 3

```
⇒ public static int BalancedPartition (String str, int n) {
    if (n == 0) {
        return 0; // TC → O(1)
    }
    int s1 = 0, l = 0;
```

```
    int ans = 0;
    for (int i = 0; i < n; i++) {
        if (str.charAt(i) == 'R') {
            l++;
        } else if (str.charAt(i) == 'L') {
            l++;
        }
        if (s1 == l) {
            ans++;
        }
    }
    return ans;
}
```

```
psvm {
    String str = "LRRRRLLRLLRL";
    int n = str.length();
    System.out.print (BalancedPartition (str, n) + "\n");
}
```

Que 2. Kth Largest Odd number in a given range -
 We have two variables L & R, indicating a range of integers from L to R inclusive, & a number K, the task is to find Kth largest odd number. If K > number of odd numbers in the range L to R then return 0.
 * Input - L = -3, R = 3, K = 1
 * Output - 3.

```
public static int kthOdd (int [] range, int k) {
    if (K <= 0) {
        return 0; // O(1)
    }
    int L = range [0];
    int R = range [1];
    if ((R < L) > 0) {
        int Count = (int) Math.ceil ((R-L+1)/2);
        if (K > Count) {
            return 0;
        } else {
            return (R-2 * K + 2);
        }
    } else {
        int Count = (R-L+1)/2;
        if (K > Count) {
            return 0;
        } else {
            return (R-2 * K + 1);
        }
    }
}

psvm {
    int [] p = {-10, 10}; int k = 8;
    System.out.println(kthOdd(p,k));
}
```

Que 3. Lexicographically smallest string of length N & sum K - We have two integers N & K. The task is to print the lexicographically smallest string of length N consisting of lower-case English alphabets such that the sum of the characters of the string equals to K where 'a' = 1, 'b' = 2, 'c' = 3 ... & 'z' = 26.

```
* Input = N = 5, K = 42           * Input = N = 3, K = 25
* Output = a amz (1+1+13+26)   * Output = aaaw (1+1+23)
⇒ public static char [] lexo_small (int n, int k) {
    char arr [] = new char [n];
    Arrays.fill (arr, 'a');
    for (int i = n-1; i >= 0; i--) {
        if (k >= i) {
            if (k >= 26) {
                arr [i] = 'z';
                k -= 26;
            } else {
                arr [i] = (char) (k+97-1);
                k -= arr [i] - 'a' + 1;
            }
        } else {
            break;
        }
        k += i;
    }
    return arr;
}

psvm {
    int n = 5, k = 42;
    char arr [] = lexo_small (n, k);
    System.out.println (new String (arr));
}
```

11 TC $\rightarrow O((N-1)c(k-1))$ - here 'c' depicts combinations i.e., $((n-1)! / ((n-k)! \times (k-1)!)$ where N is no. of elements of array & k is no. of divisions that we are having.

Ques 84 Best time to Buy & Sell Stock - Given an array prices [] of length N, representing the price of the stock on different days, the task is to find the maximum profit possible for buying & selling the stocks on different days using transactions where at most one transaction is allowed

Stock must be bought before being sold.

* Input: prices[] = {7, 6, 4, 3, 1} * Input: prices[] = {7, 1, 5, 3, 4}
* Output: 0 * Output: 5

```
⇒ public static int maxProfit (int prices[], int n) {
    int buy = prices[0], max_profit = 0;
    for (int i=1; i<n; i++) {
        if (buy > prices[i]) {
            buy = prices[i];
        } else if (prices[i] - buy > max_profit) {
            max_profit = prices[i] - buy;
        }
    }
    return max_profit;
}
```

```
public static void main (String args[]) {
    int prices[] = {7, 1, 5, 6, 4};
    int n = prices.length;
    int max_profit = maxProfit (prices, n);
    System.out.println (max_profit);
}
```

// Time Complexity - O(n)

// Space Complexity - O(1)

Ques 85 Split the given array into K sub-arrays - We have arr[] of N elements & a number K, ($1 \leq K \leq N$). Split arr[] into K subarrays (must cover all elements). The maximum subarray sum achievable out of K subarrays formed must be the minimum possible. Find that possible subarray sum.

* Input $\Rightarrow A_{8 \times 7} = \{1, 1, 2\}$, K=2 * Input $\Rightarrow A_{6 \times 7} = \{1, 2, 3, 4\}$, K=3

* Output $\Rightarrow 2$ * Output $\Rightarrow 4$

\Rightarrow public static int ans = 100000000;
public static void solve (int a[], int n, int k, int index, int sum, int maxSum) {

```
if (k == 1) {
    maxSum = Math.max (maxSum, sum);
    sum = 0;
    for (int i=index; i<n; i++) {
        sum += a[i];
    }
}
```

```
maxSum = Math.max (maxSum, sum);
ans = Math.min (ans, maxSum);
return;
```

// Space Complexity - O(n)

```
sum = 0;
for (int i=index; i<n; i++) {
    sum += a[i];
}
maxSum = Math.max (maxSum, sum);
solve (a, n, k-1, i+1, sum, maxSum);
}
```

```
public static void main (String args[]) {
    int arr[] = {1, 2, 3, 4};
    int k = 3; // K divisions
    int n = 4; // size of an Array
    solve (arr, n, k, 0, 0, 0);
    System.out.println (ans + "\n");
}
```