## **Homework 5.** Predicate Logic: Syntax and Semantics

Submit your solution in PDF file (and Latex source file if you use Latex) to blackboard by **11:59pm Sun Nov 20**.

1. Consider the following sentence.

Every number greater than or equal to 4 can be written as the sum of two prime numbers.

- (a) Write *a language* (as defined in Def 2.1) such that some formulas of the language can be used to represent the sentence above.
- (b) Write a formula of your language that should reflect the meaning of the sentence above.
- (c) In terms of your language, write three example terms, three example atomic formulas, and three example formulas.
- 2. Write a formula to represent the following information. Your formula should be as close as possible to the intended meaning of these sentences.
  - (a) There is a mother to all children.
  - (b) ALL ITEMS NOT AVAILABLE AT ALL STORES.

    Note, the sentence above is a disclaimer in the weekly flyer of specials of a grocery store chain.
- 3. Given the language defined in Def 2.1, which of the following are formulas defined by Def 2.5?
  - (a) f(x,c)
  - (b) R(c, f(d, z))
  - (c)  $\forall x(P(x))$
  - (d)  $((\exists x)(((\forall y)P(z)) \rightarrow R(x,y)))$
- 4. Given

$$((\exists x)(((\forall y)P(z)) \to R(x,y))),$$

- (a) List all its subformulas.
- (b) Draw its formation tree.
- 5. Which of the following terms are substitutable for *x* in the corresponding formulas?
  - (a) f(z,y) in  $((\exists y)(P(y) \land R(x,z)))$ .
  - (b) g(f(z,y),a) in  $((\exists x)(P(x) \land R(x,y)))$ .
- 6. Connect predicate calculus to the study of this course. To focus on the substance, we need to extend (informally) predicate calculus (syntax) as follows
  - You can use sets or proposition as a parameter of a predicate.
  - You can use = as a predicate symbol in the normal way. For example, that two terms  $t_1$  and  $t_2$  are equal can be represented by  $t_1 = t_2$ .
  - You can use a variable to refer to a function or a set or a proposition. (e.g.,  $((\exists \mathcal{V})\mathcal{V}(x) = T)$  where T is constant.)

The form of formula will be expanded accordingly to the extension above. Consider the definition of *consequence* (Def 3.7 of Part I).

- (a) Represent it as a formula. You need to introduce all predicates you need in the formula, in the way we did in L11.1 for the subset example. Intuitively, a concept name and its parameters/arguments you identify in the definition is a good candidate of a predicate. You also need to introduce any constant or function symbol you need. Note the main "if" in a definition should be understood as "if and only if." Note statement " $\forall x \in A, x \in B$ " in the subset definition. We translate it to " $((\forall x)(x \in A \to x \in B))$ ."
- (b) Let your formula be  $\alpha$ . What can you logically derive (some intuition/experience from your earlier study is needed here) from the formula ( $\alpha \wedge$  "B is a consequence of  $\{A, A \to B\}$ ")? Note you should translate the English in the formula using predicate(s) you introduced. Your answer to this question has to be in the form of a formula. You need to do a variable substitution ( $\Sigma$  in English definition would be replaced by  $\{A, A \to B\}$ , and  $\sigma$  by B).
- 7. 1) Let  $A = \{1, 2\}$ . a) List all functions from A to A in the form of sets of pair. For example, one function is  $\{(1, 1), (2, 2)\}$ . If we let the function be named g. Then in the example function, g(1) is 1 and g(2) is 2. b) List all unary relations on A. Your relations must be represented as sets.
  - 2) Show that  $\forall x(p(x) \to q(f(x))) \land \forall xp(x) \land \exists x \neg q(x)$  is satisfiable. The format of your structure should follow the one in L12.1. In your structure, you must use the domain  $A = \{1,2\}$ . You must represent the function assigned to f in the set of pairs. You must represent the relations assigned to f and f in the set forms. Note you have to figure out the constants in the language on which the formula is defined.

8. Prove that  $\mathcal{A} \models \neg \exists x \varphi(x)$  if and only if  $\mathcal{A} \models \forall x \neg \varphi(x)$ . You have to follow the proof format we used earlier. Working backward again is a good idea. You have to be able to apply the definitions.