

# L11.2 predicate Logic - the language

Chapter II - 1 (Chapter 1 of Part II)

## 3. Study the definitions of *sentences* and *open formulas*

### 3.1 sentences and open formulas

Occurrences of variables, free occurrence, free variables, bound variables

- (i)  $((\forall x)R(x, y))$  is a formula in which  $y$  occurs free but  $x$  does not. The formula  $((\exists y)((\forall x)R(x, y)))$  has no free variables; it is a sentence.
- (ii) A variable may have both a free and a bound occurrence in a single formula as do both  $x$  and  $y$  in  $((\forall x)R(x, y)) \vee ((\exists y)R(x, y))$ .
- (iii) If  $\varphi(x)$  is  $((\exists y)R(x, y)) \wedge ((\forall z)\neg Q(x, z))$  and  $t$  is  $f(w, u)$ , then  $\varphi(t) = \varphi(x/t)$  is  $((\exists y)R(f(w, u), y)) \wedge ((\forall z)\neg Q(f(w, u), z))$ . The term  $g(y, s(y))$  would, however, not be substitutable for  $x$  in  $\varphi(x)$ .

#### Definition 2.6:

- (i) A **subformula** of a formula  $\varphi$  is a consecutive sequence of symbols from  $\varphi$  which is itself a formula.
- (ii) An **occurrence** of a variable  $v$  in a formula  $\varphi$  is bound if there is a subformula  $\psi$  of  $\varphi$  containing that occurrence of  $v$  such that  $\psi$  begins with  $(\forall v)$  or  $(\exists v)$ . (This includes the  $v$  in  $\forall v$  or  $\exists v$  that are bound by this definition.) An occurrence of  $v$  in  $\varphi$  is **free** if it is not bound.
- (iii) A variable  $v$  is said to **occur free** in  $\varphi$  if it has at least one free occurrence there.
- (iv) A **sentence** of predicate logic is a formula with no free occurrences of any variable, i.e., one in which all occurrences of all variables are bound.
- (v) An **open formula** is a formula without quantifiers.

### 3.2 Substitutions

Recall the application of the definition of a concept to the use of the concept, we need substitution there!

1. When working backward from statement (B1),

- The main concept in (B1) is subset: name - subset; arguments -  $\Sigma, Cn(\Sigma)$ .
- Definition of subset:  
 $A$  is a **subset** of  $B$  if  $\forall x \ x \in A \implies x \in B$ .  
 $\forall x ( \in (x, A) \rightarrow \in (x, B) ) \rightarrow \text{subset}(A, B)$
- Substitution the arguments  $\Sigma, Cn(\Sigma)$  in the use of the concept subset for the variables  $A$  and  $B$  in the definition  
 $A: \Sigma$   
 $B: Cn(\Sigma)$
- Result of applying the definition to the use of the concept: the *definition instance* for the use of subset in statement (B1)  
**(A1)**  $\Sigma$  is a **subset** of  $Cn(\Sigma)$  if  $\forall x \ x \in \Sigma \implies x \in Cn(\Sigma)$ .

$$\forall x ( \in (x, \Sigma) \rightarrow \in (x, Cn(\Sigma)) ) \rightarrow \text{subset}(\Sigma, Cn(\Sigma))$$

**Definition 2.7: Substitution** (or **Instantiation**) If  $\varphi$  is a formula and  $v$  a variable, we write  $\varphi(v)$  to denote the fact that  $v$  occurs free in  $\varphi$ . If  $t$  is a term, then  $\varphi(t)$ , or  $\varphi(v/t)$ , is the result of substituting (or instantiating)  $t$  for all free occurrences of  $v$  in  $\varphi$ . We call  $\varphi(t)$  an **instance** of  $\varphi$ . If  $\varphi(t)$  contains no free variables, we call it a **ground instance** of  $\varphi$ .

**Definition 2.8:** If the term  $t$  contains an occurrence of some variable  $x$  (which is necessarily free in  $t$ ) we say that  $t$  is substitutable for the free variable  $v$  in  $\varphi(v)$  if all occurrences of  $x$  in  $t$  remain free in  $\varphi(v/t)$ .

About unique “readability” of a term or a formula

5-5-5 - syntax

--- parsing

(5-5)-5

(5-(5-5))

Study definitions. Some examination points

- Variables: bound, free,  
 $(p(x) \wedge ((\forall x) G(x)) \wedge H(x))$   
 $((\forall x) (G(x) \wedge H(x)))$   
  
 $((\exists x) (G(x) \wedge ((\exists x) H(x))))$
- Substitution
- Substitutable  
 $f(x, y)$  is not substitutable for  $x$  in formula  $((\forall y)P(x))$

Example of a formula and substitution

Sentence

- English: A is a subset of B if for any x, if x belongs to A then x belongs to B.
- Formula
  - Predicates
    - $\text{subset}(X, Y)$ : X is subset of Y
    - $\text{belongsTo}(X, Y)$ : X belongs to Y
  - Resulting formula  
 $\text{subset}(A, B) \leftarrow ($   
 $(\forall x), (\text{belongsTo}(x, A) \rightarrow \text{belongsTo}(x, B)) )$

$(A \times B)$  is a **subset** of  $(CXD)$

Do substitution:  $A / (A \times B)$  and  $B / (CXD)$  in the formula above

$\text{subset}(A \times B, CXD) \leftarrow ($   
 $(\forall x), (\text{belongsTo}(x, AXB) \rightarrow \text{belongsTo}(x, CXD)) )$

For all x, x is an even number. (i doesn't make sense to replace every x by a term)

Prove: if A is subset of B and B is subset of C, then A is a subset of C.

Unique readability of a term or a formula

- If t is a term, what is the form (the decomposition) of t?  
Is a constant, a variable,  $f(t_1, \dots, t_2)$  where f is a function symbol and  $t_i$  a term.

Counterexample in other language:  $1-1-1 \ ((1-1)-1) \ (1-(1-1))$

- If  $\alpha$  is a formula, what is the form (the decomposition) of  $\alpha$ ?  
Atomic formula, or  $(\alpha \vee \beta)$ , ...

Project

- Windows
  - Putty

Drawing for formulas (graphical form)

- Focus on intuition.

Semantics: meaning and truth - motivation