

Theory of Automata – Homework 5

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1. Consider the grammar $G = (V, \Sigma, R, S)$, where

$$V = \{a, b, S, A\},$$

$$\Sigma = \{a, b\},$$

$$R = \{S \rightarrow AA,$$

$$A \rightarrow AAA, .$$

$$A \rightarrow a,$$

$$A \rightarrow bA,$$

$$A \rightarrow Ab\}.$$

- a). Give a string of $L(G)$ that can be produced by applying the rules at most 4 times
b). Same string can be derived in different ways, e.g., $S \Rightarrow AA \Rightarrow aA \Rightarrow aa$, $S \Rightarrow AA \Rightarrow Aa \Rightarrow aa$. Give at least 2 distinct derivations for the string *babbab*
c). For any $m, n, p > 0$, describe a derivation in G of the string $b^m ab^n ab^p$

Sol: (a) Length no more than 4:

$$S \Rightarrow AA \Rightarrow aA \Rightarrow aa$$

$$S \Rightarrow AA \Rightarrow aA \Rightarrow abA \Rightarrow aba$$

$$S \Rightarrow AA \Rightarrow aA \Rightarrow aAb \Rightarrow aab$$

$$S \Rightarrow AA \Rightarrow bAA \Rightarrow baA \Rightarrow baa$$

$$S \Rightarrow AA \Rightarrow bAA \Rightarrow bAa \Rightarrow baa$$

$$S \Rightarrow AA \Rightarrow AbA \Rightarrow abA \Rightarrow aba$$

$$S \Rightarrow AA \Rightarrow AbA \Rightarrow Aba \Rightarrow aba$$

$$S \Rightarrow AA \Rightarrow Aa \Rightarrow aa$$

$$S \Rightarrow AA \Rightarrow Aa \Rightarrow bAa \Rightarrow baa$$

$$S \Rightarrow AA \Rightarrow Aa \Rightarrow Aba \Rightarrow aba$$

$$S \Rightarrow AA \Rightarrow AbA \Rightarrow abA \Rightarrow aba$$

$$S \Rightarrow AA \Rightarrow AbA \Rightarrow Aba \Rightarrow aba$$

$$S \Rightarrow AA \Rightarrow AAb \Rightarrow aAb \Rightarrow aab$$

$$S \Rightarrow AA \Rightarrow AAb \Rightarrow Aab \Rightarrow aab$$

Applying the rules in different orders. strings that can be generated are: aa, aab, aba, baa.

(b) Note that $A \Rightarrow bA \Rightarrow bAb \Rightarrow bab$, and also that $A \Rightarrow Ab \Rightarrow bAb \Rightarrow bab$. (8 distinct derivations):

$$S \Rightarrow AA \Rightarrow AbA \Rightarrow AbAb \Rightarrow Abab \Rightarrow^* babbab$$

$$S \Rightarrow AA \Rightarrow AAb \Rightarrow AbAb \Rightarrow Abab \Rightarrow^* babbab$$

$$S \Rightarrow AA \Rightarrow bAA \Rightarrow bAbA \Rightarrow babA \Rightarrow^* babbab$$

$$S \Rightarrow AA \Rightarrow AbA \Rightarrow bAbA \Rightarrow babA \Rightarrow^* babbab$$

Where each of these four has 2 ways to reach babbab in the last steps.

(c) Producing a sequence in terms of m, n, p that will produce the string $b^m ab^n ab^p$.

$$S \Rightarrow /* \text{ by rule } S \rightarrow AA */$$

$$AA \Rightarrow^* /* \text{ by m applications of rule } A \rightarrow bA */$$

$b^m A A \Rightarrow$ /* by rule $A \rightarrow a$ */
 $b^m a A \Rightarrow$ /* by n applications of rule $A \rightarrow bA$ */
 $b^m ab^n A \Rightarrow$ /* by p applications of rule $A \rightarrow Ab$ */
 $b^m ab^n Ab^p \Rightarrow$ /* by rule $A \rightarrow a$ */
 $b^m ab^n ab^p$
 Clearly this produces $b^m ab^n ab^p$ for each m, n, p .

2. Consider the grammar $G = (V, \Sigma, R, S)$, where

$$V = \{a, b, S, A\},$$

$$\Sigma = \{a, b\},$$

$$R = \{S \rightarrow aAa,$$

$$S \rightarrow bAb,$$

$$S \rightarrow e,$$

$$A \rightarrow SS\}.$$

Give a derivation of the string $baabbb$ in G .

Sol: Derivation of $baabbb$

$$\begin{aligned}
 S &\Rightarrow bAb \\
 &\Rightarrow bSSb \\
 &\Rightarrow baAaSb \\
 &\Rightarrow baSSaSb \\
 &\Rightarrow baSaSb \\
 &\Rightarrow baaSb \\
 &\Rightarrow baabAbb \\
 &\Rightarrow baabSSbb \\
 &\Rightarrow baabSbb \\
 &\Rightarrow baabbb
 \end{aligned}$$

3. Show that the following languages are context-free by exhibiting context-free grammars generating each.
- $\{a^m b^n : m \geq n \geq 0\}$
 - $\{a^m b^n c^{2m+n} : m, n \geq 0\}$

Sol: **(a)** : $S \rightarrow aAB / AB$
 $A \rightarrow aA / a$
 $B \rightarrow bB / b / e$

(b) : $S \rightarrow ABC / e$
 $A \rightarrow aA / a / e$
 $B \rightarrow bB / b / e$
 $C \rightarrow cC / c / e$