## Theory of Automata - Homework 6

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- 1. Consider the pushdown automata  $M = (K, \Sigma, \Gamma, \Delta, s, F)$ , where
  - $K = \{s, f\},$
  - $F = \{f\},$
  - $\Sigma = \{a, b\},\$
  - $\Gamma = \{a\},$
  - $\Delta = \{((s, a, e), (s, a)), ((s, b, e), (s, a)), ((s, a, e), (f, e)), ((f, a, a), (f, e)), ((f, b, a), (f, e))\}.$
  - a). Trace all possible sequence of transitions of M on input aba
  - b). Show that aba, aa,  $abb \notin L(M)$ , but baa, bab,  $baaaa \in L(M)$
  - c). Describe L(M) in English

Sol: a) There are three possible computations of M on aba:

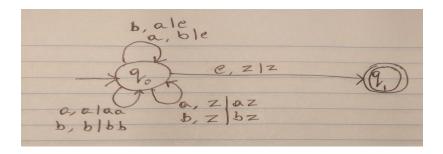
- (s, aba, e) | (s, ba, a) | (s, a, aa) | (s, e, aaa)
- (s, aba, e) | (s, ba, a) | (s, a, aa) | (f, e, aa)
- (s, aba, e) |- (f, ba, e)

None of these is an accepting configuration.

- (b) This is done by tracing the computation of M on each of the strings, as shown in (a).
- (c) L(M) is the set of strings whose middle symbol is a. In other words, L(M) =  $\{xay \in \{a, b\}^* : |x| = |y|\}$ .
  - 2. Construct a Pushdown automata that accept each of the followings:
    - a). The language  $\{w \in \{a, b\}^* : w = w^R\}$
    - b). The language  $\{w \in \{a, b\}^* : w \text{ has the same number of } a's \text{ and } b's\}$

Sol : a)

Z is the 'start stack symbol'



## b)

