

Today's agenda:

Some features in Haskell adopted from LC:

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Function definition

Example 1

```
square :: Int -> Int      (1)
square x = x * x          (2)
```

Line (1) says that the function named square takes as i/p an integer and returns as o/p an integer; this is called the signature or type of the function; \rightarrow denotes maps to.

Line (2) the body of the function is written in the RHS of $=$. LHS says that the input parameter is x . The notation of function application is employed in the LHS (like $f\ x$ or $\sin\ \theta$).

In general, the structure of function definition will look like

```
<function name> :: <type of the function>
<function name> <input parameters> = <body of the function>
```

Example 2

```
plus :: Int -> Int -> Int
plus x y = x + y
```

The function named plus takes as i/p two integers and returns as o/p an integer

List

Example: $[2, 3, 7]$, $[]$, $[2]$

Empty list: $[]$ length of this list is zero

Nonempty list $[2, 3, 7]$ can be written as $2 : [3, 7]$: constructor for list
It has a head (2)—an item and a tail $[3, 7]$ —a list; we will write $x : xs$ to represent a nonempty list

How is a list constructed?

$[2, 3, 7] = 2 : [3, 7] = 2 : 3 : [7] = 2 : 3 : 7 : []$: right associative

The process works internally as per the following steps.

1. Empty list `[]`
2. `7 : [] = [7]` (element, list)
3. `3 : [7] = [3,7]` (element, list)
4. `2 : [3,7] = [2,3,7]` (element, list)

The operator `:` has type `a -> [a] -> [a]` takes as input an element of type `a` and a list of elements of type `a` and returns a list of elements of type `a`.

Recursion with lists

```
length :: [a] -> Int
length [] = 0
length (x : xs) = 1 + length xs
```

computation:

```
length [2,3,5] = length (2 : [3,5])
               = 1 + length [3,5]
               = 1 + length (2 : [5])
               = 1 + 1 + length [5]
               = 1 + 1 + length (5 : [])
               = 1 + 1 + 1 + length []
               = 1 + 1 + 1 + 0
               = 3
```

```
sum :: [Int] -> Int
sum [] = 0
sum (x : xs) = x + sum xs
```

```
sum [2, 3, 5] = sum (2 : [3,5]) = 2 + sum [3,5] = 2 + sum (3 : [5]) = 2 + 3 + sum (5 : []) = 2+3+5 + sum []
              = 2+3+5+0 = 10
```

Insertion sort

suppose we have a sorted list, say, in ascending order. we want to insert an element into this list (in the proper position)

```
insert :: Int -> [Int] -> [Int]
insert x [] = [x]
insert x (y : ys)
  | x <= y = x : (y : ys)      | denotes guard
  | x > y  = y : insert x ys
```

```
2 : [3, 6] = [2,3,6]
5 : [3,6] = 3 : insert 5 [6]
insert 5 [6] = insert 5 (6 : []) = [5,6]
```

combining we get [3,5,6]

`inSort :: [Int] -> [Int]`

`inSort [] = []`

`inSort (x : xs) = insert x (inSort xs)`

let us do insertion sort on the list 7: (5:3)

`inSort (7: [5,3])`

`insert 7 (inSort [5,3])`

`inSort [5,3] = insert 5 (inSort [3])`

`inSort (3: []) = insert 3 (inSort [])`

`inSort [] = []`

`insert 3 [] = [3]` first created

then we obtain [3,5]

finally we obtain [3,5,7]

end of lecture 24.2.

HOF

Function Composition $(f \bullet g) x = f (g x)$

Composition is associative: $(f \bullet g) \bullet h = f \bullet (g \bullet h)$ for all f, g, h

Currying

`Multiply :: Int -> Int -> Int`

\rightarrow is right associative

`Multiply 2 :: Int -> Int`

`(Multiply 2) 3 :: Int`

`Multiply 2 3 :: Int` function application is left associative

`MultiplyUC :: (Int, Int) -> Int`

`MultiplyUC (x, y) = x * y`

The type `Int -> Int -> Int` is same as `Int -> (Int -> Int)` but it is different from `(Int -> Int) -> Int`

`g :: (Int -> Int) -> Int`

`g h = (h 1) + (h 2)`

let `h = succ`, then `g h = 5`

map function

`map :: (a -> b) -> [a] -> [b]`

a, b are type variables

`map: (Int -> Int) -> [Int] -> [Int]`

instance of map

`map succ [2,4,6] = [3,5,7]`

`map sqr [2,4,6] = [4,16, 36]`

`map f [] = []`
`map f (x : xs) = f x : map f xs`

`filter iseven [2, 3, 6] = [2,6]`

`iseven n = (mod n 2 == 0)`

`filter :: (a -> Bool) -> [a] -> [a]`
`filter p [] = []`
`filter p (x : xs)`
| `p x = x : filter p xs`
| `otherwise = filter p xs`

Lambda abstraction:

`map addone [2,3,4] = [3,4,5]`
`map (\x -> x+1) [2,3,4] = [3,4,5]`

`(\x -> x+1) = \x. x + 1`

`f x y z = result` is same as `\x y z -> result` is same as `\x. \y. \z. result`

`length :: [a] -> Int`
`length [] = 0`
`length (_ : xs) = 1 + length xs` since x has no role in the RHS

concatenation operator: `++`
`[1,2,4] ++ [9,4] = [1,2,4,9,4]`

`reverse :: [a] -> [a]`
`reverse [] = []`
`reverse (x : xs) = (reverse xs) ++ [x]`

`isequal :: Int -> Int -> Bool`
`isequal x x = True`
`isequal x y = False` [what is wrong with this code?] [arguments must be distinct]

`isequal x y = (x == y)` `==` denotes test

can we check if two lists are same?

`isequalLists :: [a] -> [a] -> Bool`
`isequal ListA ListB = (ListA == ListB)`

give a recursive code that checks element wise.

End of lecture 27.2.25

List comprehension

Let `inp = [2, 4, 8]`

`[3*n | n <- inp]` will be `[6, 12, 24]` `<-` is like the set membership operator

"Perform `3*n` for all `n` in the list `inp`."

`quicksort :: [Int] -> [Int]`

`quicksort [] = []`

`quicksort (x : xs) = quicksort [y | y <- xs, y <= x] ++ [x] ++ quicksort [y | y <- xs, y > x]`

we have simply written the idea behind quicksort.

Lazy evaluation

`f x y = x + y` `let x = (9-3) y = (f 34 3)`

then `f x y = f (9-3) (f 34 3) = (9-3) + (f 34 3) = 6 + (34 + 3) = 6 + 37 = 43` [CBN]

Infinite lists

`ones = 1 : ones` will produce `[1, 1, 1,]` interrupt by ctrl-c

`addfirsttwo :: [Int] -> Int`

`addfirsttwo (x:y:zs) = x + y`

`addfirsttwo ones = addfirsttwo (1 : ones) = addfirsttwo (1 : 1 : ones) = 1+1 = 2`

in-built `[n ..]` will give for `n=2` `[2, 3, 4, 5,]`

`[n, m ..]` will give for `n=2, m=2` `[2, 4, 6, 8,]`

`[n .. m]` will give `[n, n+1, n+2, ..., m]`

will give `[]` if `m < n`

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Development of Haskell (or any FPL) from the primitive PL (i.e., LC) that in turn from the primitive recursive functions and computable functions in general.

PRF : 1920s LC : 1930s Haskell : 1990s [60-70 years]

ANN : 1950s DL : 2000 AI-ML (popularity) : 2010s [60 years]

End of lecture 27.2.25

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