

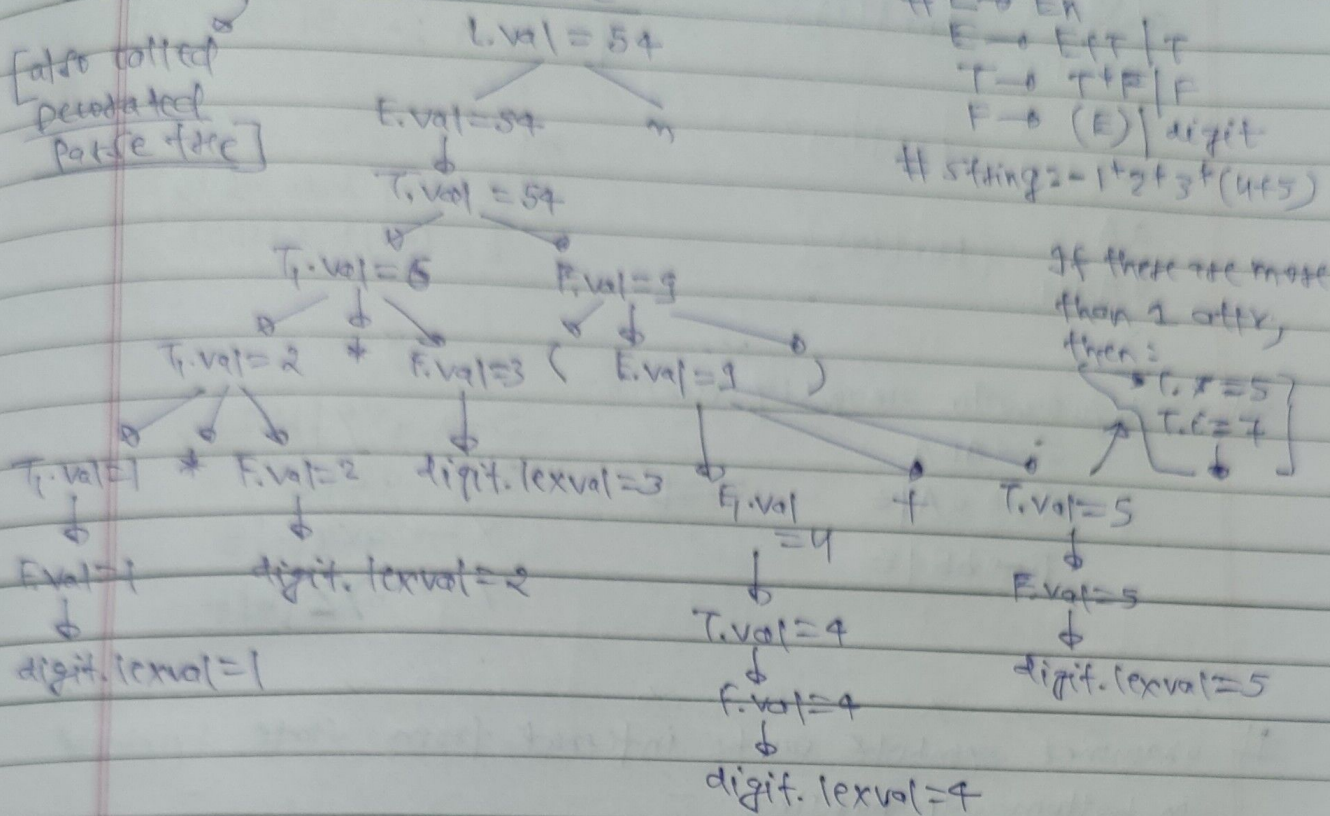
MTE Practice

1.

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Annotated parse tree :-

[also called
annotated
parse tree]



Bottom up parsing:-

$X \rightarrow a A M_1 \mid C M_2$

$A \rightarrow b B M_3 \mid a M_4$

$B \rightarrow x C M_5 \mid C M_2$

$M_1 \rightarrow \epsilon \{ \text{print}(3) \}$

$M_2 \rightarrow \epsilon \{ \text{print}(5) \}$

$M_3 \rightarrow \epsilon \{ \text{print}(2) \}$

$M_4 \rightarrow \epsilon \{ \text{print}(4) \}$

$M_5 \rightarrow \epsilon \{ \text{print}(1) \}$

ababccc

abab M_2 cc {print(5)}

abab~~X~~cc

abab $X C M_5 c$ {print(1)}

ababBC

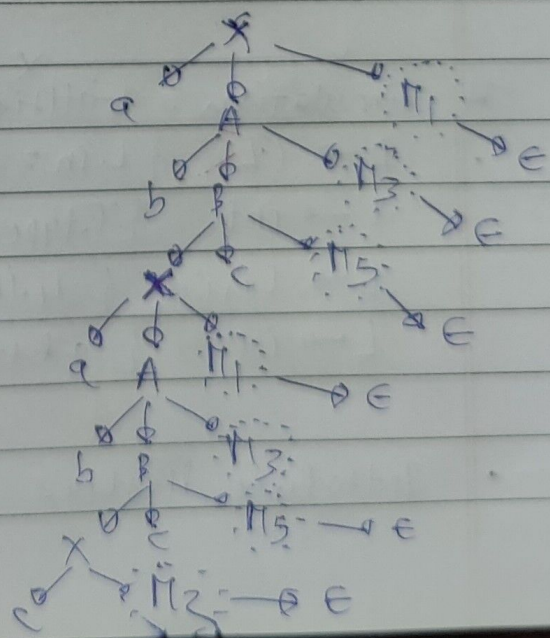
abab $B M_3 c$ {print(2)}

abaAe

aba $A M_1 c$ {print(3)}

Introduce Marker Symbols and make highest derivation of string to be parsed.

Bottom-up will be in reverse



abX < {print(1)}

abX < M5

abB

abBM3 {print(2)}

aA

aAM1 {print(5)}

Output :- 5123123

Ans

string reduced to start symbol

Ambiguous to unambiguous :-

S → ABA

A → aA | ε

B → bB | ε

S → aXY | bYZ | ε

Z → aZ | ε

X → aXY | a | ε

Y → bYZ | b | ε

Grammar symbols can be inferred from state symbols in bottom-up parsing :-

- state stack = 0, 2, 4, 7 {At some point of time, the config. of state stack}
- From state 0 to 2, we have 'T' as the transition
- From 2 to 4, (*)
- From 4 to 7, (F)

Then state 2 will correspond to (T)

4 will be (*)

7 will be (F)

Simulation: L-affiliated + LR-parsing (bottom-up) :-

D → TL {L.in = T.type}

T → int {T.type = int}

L → L, id {L.in = L.in; addtype(id.entry, L.in);}

L → id {addtype(id.entry, L.in);}

- Introduce Markers and form SPT :-

$D \rightarrow TM, L$

$T \rightarrow \text{int}$

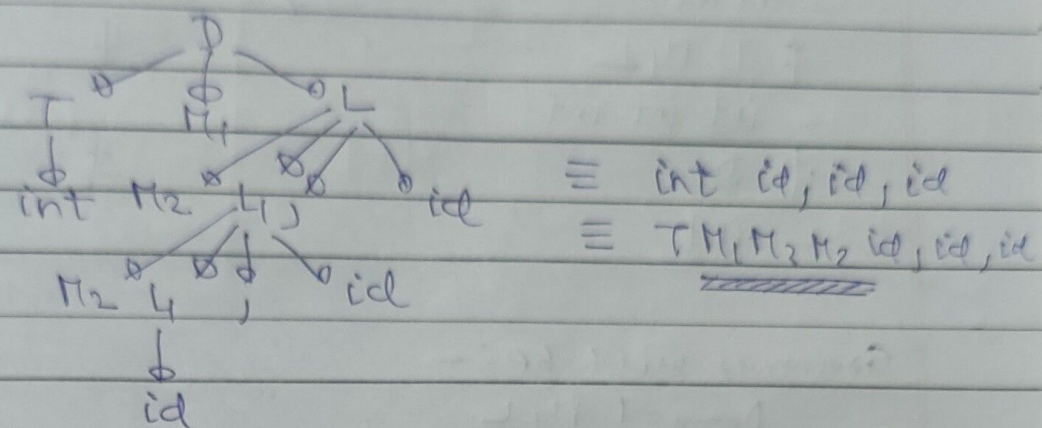
$L \rightarrow M_2 L, id \{ \text{addtype}(id.entry, L.in); \}$

$L \rightarrow id \{ \text{addtype}(id.entry, L.in); \}$

$M_1 \rightarrow \epsilon \{ M_1.i = T.type, M_1.s = M_1.t \}$

$M_2 \rightarrow \epsilon \{ M_2.i = L.inh, M_2.s = M_2.t \}$

For simulation, do rightmost derivation of the string :-



stack	input
\$	int id, id, id \$
\$ int	id, id, id \$
\$ T	id, id, id \$
\$ T M ₁	id, id, id \$
\$ T M ₁ M ₂	id, id, id \$
\$ T M ₁ M ₂ M ₂	id, id, id \$
\$ T M ₁ M ₂ M ₂ id	, id, id \$
\$ T M ₁ M ₂ M ₂ id id.entry	

[undefined] →

[attribute or value stack]

[undefined stack]

M₁ will take value from top-t always

(1)

$\$ T M_1 M_2 L$ $id, id \$$ $\{ addtype(id.entry, L.in) \}$
 $- int int int int int$
 $\$ T M_1 M_2 M_2 L, id$ $id \$$ [from "top-1"]
 $- int int int int int - id.entry$

 $\$ T M_1 M_2 L$ $, id \$$ reduce it
 $- int int int int$ $\{ id.entry at$
 $\$ T M_1 M_2 L, id$ $\$$ top and
 $- int int int int - id.entry$ $L.in will be$

 $\$ T M_1 L$ $\$$ at top-4
 $- int int int$

 $\$ D$ $\$$ \Rightarrow [string reduced
 $- int$ to start symbol]

Grammar will be:-

$D \rightarrow T M_1 L$
 $T \rightarrow int \{ addtype(value[top], entry, value[top-1].in) \}$
 ~~$L \rightarrow M_2 L, id \{ addtype((value[top], entry), value[top-4].in) \}$~~
 $L \rightarrow M_2 L, id \{ addtype((value[top], entry), value[top-4].in) \}$
 $L \rightarrow id \{ addtype(value[top].entry, value[top-1].in) \}$
 $M_1 \rightarrow \epsilon \{ value[++top].in = value[top-1].type; value[top].syn = value[top].inh \}$
 $M_2 \rightarrow \epsilon \{ value[++top].in = value[top-1].syn; value[top].~~inh~~syn = value[top].inh \}$

[use synthesized / here]

X X X X

L-attributed + LL-Grammar (Top-down Parsing): -

- "action records" will have pointer to actions
 - these actions will copy the value in the bottom of the stack and execute some print statements or side effects above the symbol.
- "Synthesized Records"
 - will contain synthesized attribute.
 - there will be parent with the symbol itself. (just below it)

$A \rightarrow BC$

		action(B)
		B
	\Rightarrow	action(C)
		B
A		syn(B,C)
syn(A)		syn(A)

When 'A' gets popped off, $inh(A)$ will also get popped out and hence the value of $A.i$ will have copied to all the locations where it is required.
 \hookrightarrow action(B),
 action(C)

$\overline{S} \rightarrow$ while ({ $L1 = new()$
 $L2 = new()$
 $c.false = S.next$
 $c.true = L2$
 print (label $L1$) }
 <) { $S.next = L1$;
 print (label $L2$); }
 \overline{S}

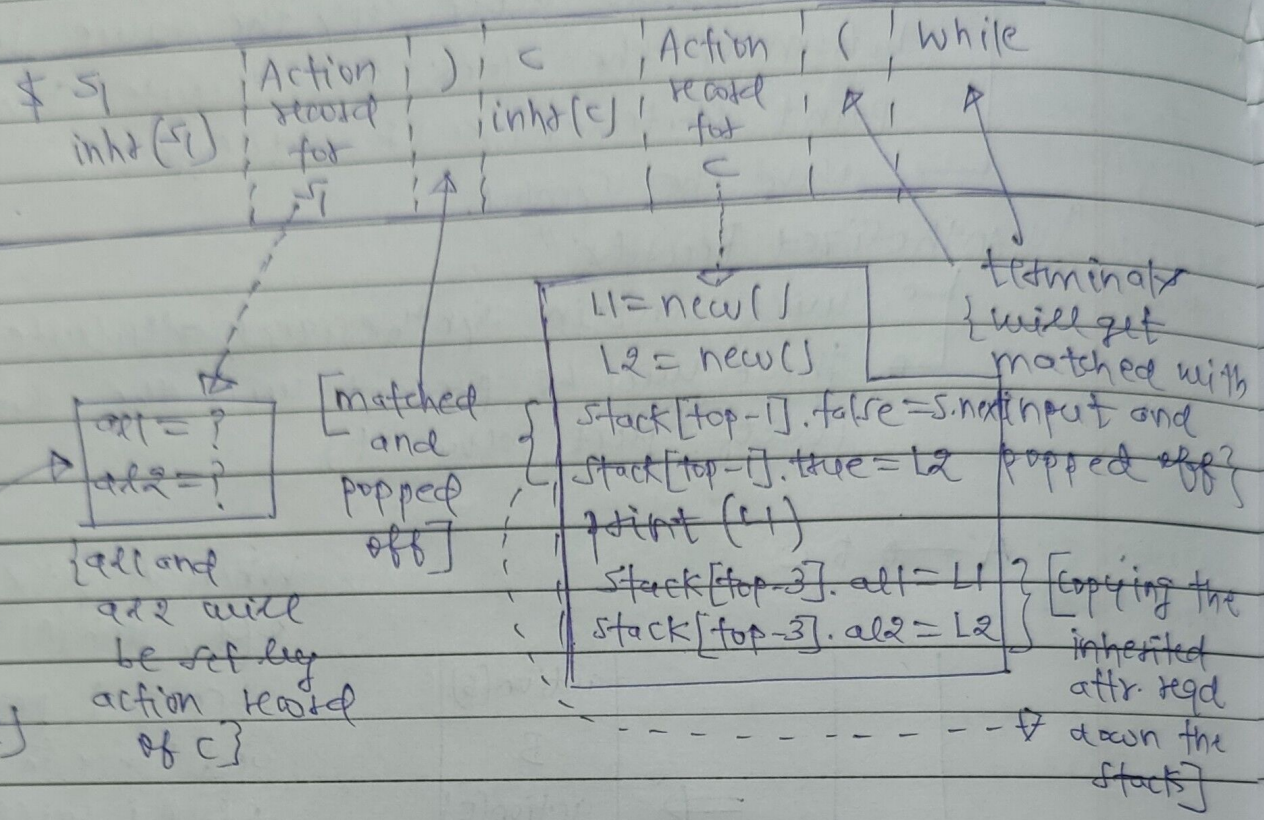
\hookrightarrow { $L1$ and $L2$ are labels for conditional jumps. }

[is called inherited record for S]

stack

\$ S
- s.next

symbol stack
attribute stack



s.next will be set earlier by the action record of S before popping off.

- # S-attribute + LR grammar → Bottom-up parsing
 - # L-attribute + LL grammar → top-down parsing
 - # L-attribute + LR grammar → No parsing
- ↳ We can't do syntax direction translation along with parsing here.

	LR	LL
S-attr.	B	B/T
L-attr.	X	B/T