

# LR parsing (Canonical LR) $\equiv$ CLR

DATE: / /

# Grammar:-

$$S' \rightarrow S$$

$$S \rightarrow AA$$

$$A \rightarrow aA \mid b$$

# { Grammar already augmented with new start symbol  $S'$  }

# Collection computation:-

- Compute collection  $C = \{I_0, I_1, \dots, I_n\}$  where  $I_i = \text{set of LR(1) items}$ .

- $I_0 = \text{closure}([S' \rightarrow \cdot S, \$])$

[e.g.]  $= \{ [S' \rightarrow \cdot S, \$], [S \rightarrow \cdot AA, \$], [A \rightarrow \cdot aA, a/b], [A \rightarrow \cdot b, a/b] \}$

Algorithm for computing closure

for computing closure

# [lookahead of  $S$  must be  $\$$  for reduction by oth rule]

$\text{first}(A\$)$

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- $I_1 = \text{goto}(I_0, S)$

~~impracticable~~  
 $= \{ [S' \rightarrow \cdot S, \$] \}$

For the first LR(1) item, just copy the lookahead from the set, it is deriving [here, that set is  $I_0$ ].

- $I_2 = \text{goto}(I_0, A)$

$$= \{ [S \rightarrow \cdot AA, \$], [A \rightarrow \cdot aA, \$], [A \rightarrow \cdot b, \$] \}$$

these two are

computed by closure operation

on  $[S \rightarrow \cdot AA, \$]$  item.



$$\begin{aligned}
 \cdot \quad I_3 &= \text{goto}(I_0, a) \\
 &= \text{closure}([A \rightarrow a.A, a|b]) \\
 &= \text{closure}([A \rightarrow a.A, a]) \cup \text{closure}([A \rightarrow a.A, b]) \\
 &= \{ [A \rightarrow a.A, a|b], \\
 &\quad [A \rightarrow .aA, a|b], \\
 &\quad [A \rightarrow .b, a|b] \}
 \end{aligned}$$

$$\begin{aligned}
 \cdot \quad I_4 &= \text{goto}(I_0, b) \\
 &= \text{closure}([A \rightarrow b., a|b]) \\
 &= \{ [A \rightarrow b., a|b] \}
 \end{aligned}$$

[Note  $I_3$  and  $I_6$  are diff. in look-ahead]

$$\begin{aligned}
 \cdot \quad I_5 &= \text{goto}(I_2, A) \\
 &= \text{closure}([S \rightarrow AA., \$]) \\
 &= \{ [S \rightarrow AA., \$] \}
 \end{aligned}$$

$$\begin{aligned}
 \cdot \quad \text{goto}(I_2, a) &= \text{closure}([A \rightarrow a.A, \$]) \\
 &= \{ [A \rightarrow a.A, \$], \\
 &\quad [A \rightarrow .aA, \$], \\
 &\quad [A \rightarrow .b, \$] \} = I_6
 \end{aligned}$$

$$\begin{aligned}
 \cdot \quad \text{goto}(I_2, b) &= \text{closure}([A \rightarrow b., \$]) \\
 &= \{ [A \rightarrow b., \$] \} = I_7
 \end{aligned}$$

$$\begin{aligned}
 \cdot \quad \text{goto}(I_3, A) &= \text{closure}([A \rightarrow aA., a|b]) \\
 &= \{ [A \rightarrow aA., a|b] \} = I_8
 \end{aligned}$$



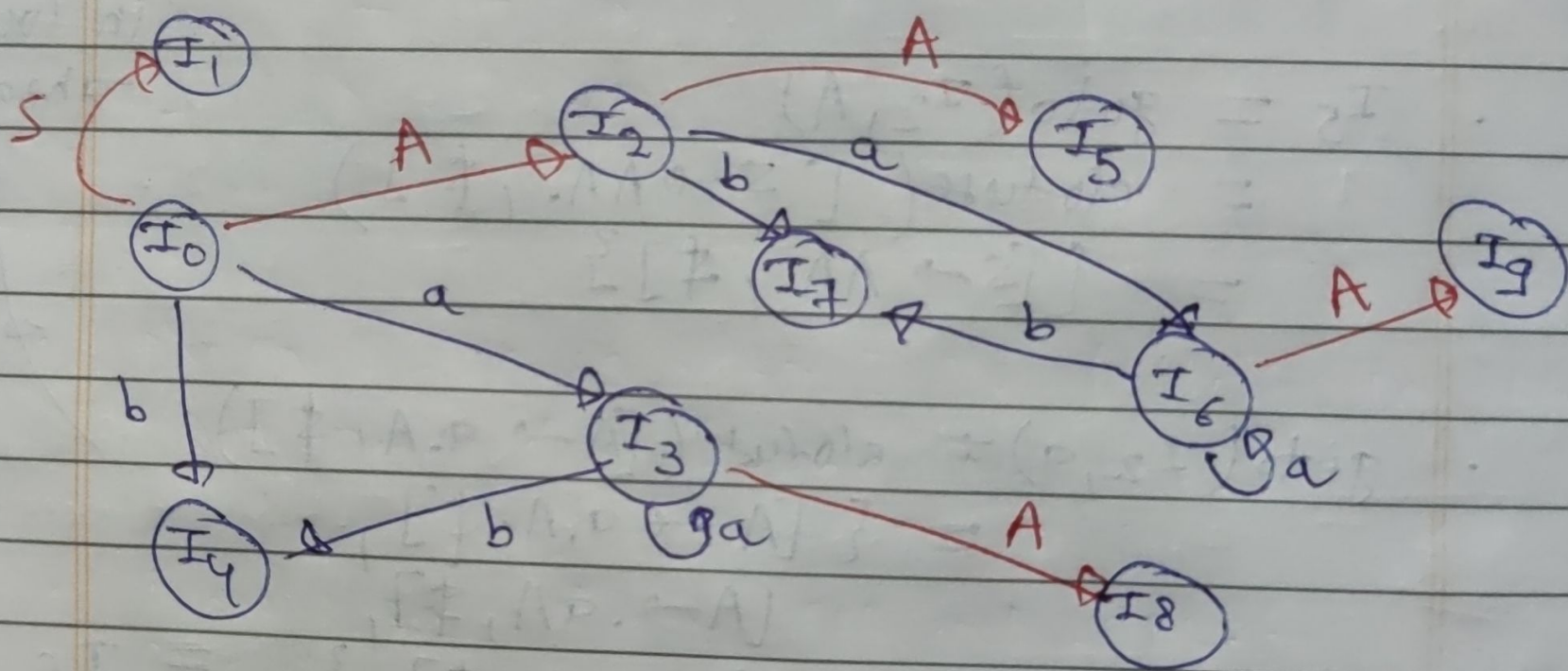
$$\cdot \text{goto}(I_3, a) = I_3 ; \text{goto}(I_3, b) = I_4$$

$$\cdot \text{goto}(I_6, a) = I_6 ; \text{goto}(I_6, b) = I_7$$

$$\cdot \text{goto}(I_6, A) = \text{close}([A \rightarrow aA., \$]) \\ = \{[A \rightarrow aA., \$]\} = I_9.$$

These are 10 states in the automaton of CLR parser.

# automata:-



# Note :- For reduce operation, in action table, now go to every  $I_i$  and look for  $[A \rightarrow a., a]$ , and

take cost of

$[s., \$]$

##  $\{ \text{action}[i, a] = \text{reduce by } (A \rightarrow a.) \}$



⇒ No conflict; No multiple entries.

Hence, grammar is CLR grammar or LR(1) grammar.

# Goto table:-

	S	A
0	1	2
1		
2		5
3		8
4		
5		
6		9
7		
8		
9		

# action table:-

	a	b	\$
0	S3	S4	
1			accept
2	S6	S7	
3	S3	S4	
4	R3	R3	
5			R1
6	S6	S7	
7			R3
8	R2	R2	
9			R2

(A-a)