

Lecture 35

Code Optimizations

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• Way to represent a Partial Order set



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- Greatest Lower Bound (GLB)



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If transfer function is monotone and distributive, then MFP solution will be equal to MOP.



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Meet operator is the intersection to calculate dominators.



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- Natural Loops



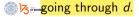
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 - ▶ There must be a back edge that enters loop header
 - ▶ If there is backedge $n \rightarrow d$, natural loop will be d and all the nodes that can reach n without



Thank You

