



Lecture 35

Code Optimizations

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Lattice Theory

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- Greatest Lower Bound (GLB)

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- If transfer function is monotone and distributive, then MFP solution will be equal to MOP.

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Meet operator is the intersection to calculate dominators.

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 - ▶ If there is backedge $n \rightarrow d$, natural loop will be d and all the nodes that can reach n without going through d .

Thank You