

Lecture 5

Lexical Analysis

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January 29, 2025



• Difference between lexeme and token



- Difference between lexeme and token
- Approaches for implementing lexical analyzer



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- Difficulties while doing lexical analysis



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- Maximal Munch







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• Between m and n occurrence



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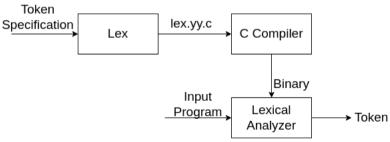
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transition rules

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```
%{
include < math.h >
int count;
%}
```

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Pattern $\{Action\}$



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Auxiliary Functions

C Code, going directly into lex.yy.c



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Important functions/variables

• yyin

Transition Rule

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Auxiliary Functions

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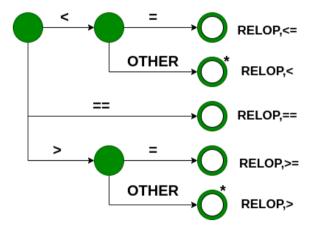
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- Regular expressions alone are not enough
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- Lexical definitions consist of regular definitions, priority rules, and maximal munch principle
- Construct an analyzer that will return \(\text{token, lexeme} \) pairs



Transition Diagram for Relational Operator





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- This can be implemented using lots of switch-cases in C programming language.



Interface to other passes

