

Lecture 31

Code Generation

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- For a leaf operand x, if base is b generate the instruction LD R_h , x



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 - Generate the instruction OP R_{b+k-1} , R_{b+k-2} , R_{b+k-1}





Start from the root of the tree with base b=1. Assume the number of registers are r>=2. For node N with label r or less, use same algo.

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