

[CSN212] Assignment 2  
Divide and Conquer  
[Maximum Marks 100]

## 1 Rank in two sorted arrays [10 Marks]

Given two sorted arrays of size  $m$  and  $n$ . *Design* and *analyze* a divide and conquer algorithm to compute the  $k^{th}$  smallest number in the union of two arrays.

## 2 Smallest Triangle [15 Marks]

Given a set of  $n$  points on a 2D plane, *design* and *analyze* an algorithm to find three points making a triangle of the smallest size (perimeter).

## 3 Maximum Sum Subarray and Maximum Difference [15 Marks]

Consider the two problems described in Section 4.1 CLRS textbook.

1. Given an array  $A$  of  $n$  positive integers, the difference of the ordered pair  $(i, j)$  where  $i \leq j$  is  $A[j] - A[i]$ .
2. Given an array  $A$  of  $n$  integers, the sum of its sub-array is  $A[i, j] = A[i] + A[i + 1] + \dots + A[j]$ .

The textbook algorithm uses a reduction amongst the problems and gives  $O(n \log n)$  solutions for them.

1. *Design* and *analyze* divide and conquer algorithms (without reduction) to solve them in  $O(n)$  time.
2. *Design* and *analyze* linear sweep algorithms (without reduction) to solve them in  $O(n)$  time.

**Hint:** Scan the Array from left to right and maintain the necessary data structures.

## 4 Extended FFT [20 Marks]

In the FFT algorithm covered in the class, we divide the problems of computing  $A$  into two subproblems as of computing  $A_{odd}$  and  $A_{even}$ . This requires us to use roots of unity to ensure  $|X^2| = |X|/2$ .

1. Design a divide step to divide each problem into 4 subproblems. What would be the samples used? Compute the complexity of the algorithm.
2. Repeat the above when we divide it into 3 sub-problems.
3. Repeat the above for general case of  $k$  sub-problems.

## 5 $k$ th quantiles, multiary Quicksort [20 Marks]

In Quicksort a *pivot* element partitions an array of numbers to those smaller than the pivot and those larger than the pivot (assume for simplicity that all numbers are unique). Consider a multiary version of Quicksort, where  $k - 1$  pivots are chosen to partition the numbers to  $k$  classes analogously to the standard Quicksort. Show that such multiary partitioning can be produced in  $O(n \log k)$  time so that the all classes are of the same size (to within 1).

## 6 FFT Applications [20 Marks]

1. Recall the pattern matching algorithm from class using FFT. Modify the algorithm to report whether a pattern  $A$  and any substring of the string  $B$  matches  $k$ -cyclically for a given  $k \in [0, 25]$ , where  $k$ -cyclic matching of two strings  $P$  and  $Q$  over characters  $[a, z]$  of length  $n$  is defined as:

*For each  $i \in [0, k]$ , there are at least  $\lfloor \frac{n}{k} \rfloor$  matches between characters of  $P$  and  $i^{\text{th}}$  cyclic shift of character in  $Q$  along  $[a, z]$ , eg. third cyclic shift of  $a$  and  $y$  are  $d$  and  $b$  respectively.*

Also, analyze the complexity of the algorithm.

2. Recall the pairwise sum problem from the class. Given two arrays of numbers  $A[n]$  and  $B[n]$  having numbers in the range  $[0, m]$ . Use FFT once to report the number of elements of  $A$  and  $B$  that can be used to generate the sum  $k$  when a pair is made using one element from  $A$  and one element from  $B$ .

## 7 Textbook Problems (Solve don't submit)

Solve the following problems from CLRS Textbook (3rd Edition)

4-5, 4-6, 9-2, 30.1-7, 30-3, 30-6