

Lecture 11

17.2.2025

Today's agenda:

Boolean operations:

$\text{true} = \lambda x. \lambda y. x$

$\text{false} = \lambda x. \lambda y. y$

$\text{OR} = \lambda x. \lambda y. x \text{ true } y$ [version 1]

$\text{OR} = \lambda x. \lambda y. x \ x \ y$ [version 2] does not involve any other function; thus better

$\text{AND} = \lambda x. \lambda y. x \ y \ \text{false}$ [version 1]

$\text{AND} = \lambda x. \lambda y. x \ y \ x$ [version 2] does not involve any other function; thus better

$\text{NOT} = \lambda x. x \ \text{false} \ \text{true}$

$\text{XOR} = \lambda x. \lambda y. x \ (\text{NOT } y) \ y$

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Idea: use the functions true and false appropriately.

consider OR: $x \ \text{OR} \ y$ [version 1]

$\lambda x. \lambda y. x \ ______ \ y$

When we pass the actual parameter say $x=\text{true}$, we get $\text{true} \ ______ \ y$

By the property, “ $______$ ” would be returned. Since we know the answer as true so “ $______$ ” is true.

What if $x=\text{false}$? Then the answer is whatever is the value of y . so the second argument should be y and our choice is correct.

So we get OR: $\lambda x. \lambda y. x \ \text{true} \ y$

Take a closer look. What if x is true? It is true which is x . So we have

OR: $\lambda x. \lambda y. x \ x \ y$ now the body of OR does not contain any constant (true/false) [version 2]

Now we can design AND easily

What if $x=\text{false}$? Then the answer is false irrespective of the value of y . Otherwise, it is the value of y .
[version 1]

Take a closer look. What if x is false? It is false which is x . So we have

AND: $\lambda x. \lambda y. x \ y \ x$ [version 2]

Pair: (a, b)

The function for Pair should specify the two components and the projection function.

$\text{pair} = \lambda x. \lambda y. \lambda z. (z \ x \ y)$ x,y are variables for the actual parameters e.g., a, b

$\text{fst} = \lambda p. p \ (\lambda x. \lambda y. x)$ first component of the pair [NB: *fst* is not the same as *first*]

$\text{snd} = \lambda p. p \ (\lambda x. \lambda y. y)$ second component of the pair [*snd* is not the same as *second*]

let us compute $\text{fst} (\text{pair } a \ b)$ and $\text{snd} (\text{pair } a \ b)$:

$\text{pair } a \ b = (\lambda x. \lambda y. \lambda z. (z \ x \ y)) \ a \ b$

$= (\lambda y. \lambda z. (z \ a \ y)) \ b$ “=” means beta-reduction

$= \lambda z. (z \ a \ b)$

$= N \ (\text{say})$

$\text{fst } N = (\lambda p. p \ (\lambda x. \lambda y. x)) \ N$

$= (\lambda p. p \ (\lambda x. \lambda y. x)) \ (\lambda z. (z \ a \ b))$ see the role of *p* and *z*.

$= (\lambda z. (z \ a \ b)) \ (\lambda x. \lambda y. x)$ *z* may be replaced with either *first* or *second*

$= (\lambda x. \lambda y. x) \ a \ b$ *first a b*

$= (\lambda y. a) \ b$

$= a$

$\text{snd } N = \lambda p. p \ (\lambda x. \lambda y. y) \ N$

$= \lambda p. p \ (\lambda x. \lambda y. y) \ (\lambda z. (z \ a \ b))$

$= (\lambda z. (z \ a \ b)) \ (\lambda x. \lambda y. y)$

$= (\lambda x. \lambda y. y) \ a \ b$ *second a b*

$= (\lambda y. y) \ b$

$= b$

End of lecture