Powerful tricks with modulo calculation - Boris Sokolov

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Trick #1:
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(A / B) % MOD = (A % (MOD * B)) / B

Conditions: none.

Advices: use this trick only if B can be not coprime with MOD, because new modulus - MOD B can be large. How to avoid overflow working with large modulus read at the trick #5.

Trick #2:

(A / B) % MOD = ((A % MOD) * (B $^{phi\,(90D)}$ - 1 % MOD)) % MOD ,where phi is Euler's totient function

Conditions: B and MOD are coprimes.

Proof: (A / B) % MOD = ((A % MOD) * (B⁻¹ % MOD)) % MOD from Exponentiation properties. And from Euler's theorem follows that $B^{phi (MOD)}$ % MOD = 1. Let's multiply this equation by B^{-1} and we will get that B^{-1} % MOD = $B^{phi (MOD)}$ - 1 % MOD.

Trick #3:

 $(A / B) % MOD = ((A % MOD) * (B^{NOD - 2} % MOD)) % MOD$

Conditions: B and MOD are coprimes, MOD is a prime number.

Advices: if you're sure that MOD is prime, better use this trick instead of trick #2. Remember that 10°+7 and 10°+9 are prime numbers.

Proof: if MOD is prime then phi(MOD) - MOD - 1 from properties of Fuler's totient function. As it's just a particular case of trick #2, the rest of proof is similar.

Trick #4:

AN & MOD = AN & Phi (MOD) & MOD

Conditions: A and MOD are coprimes.

Advices: use this trick only if N can't be present in any standard data type, otherwise use Fast exponentiation. Proof: from Euler's theorem follows that $A^{\text{phi} (MOD)}$ % MOD = 1. It's easy to see that A^{X} * phi (MOD) + Y then A^{N} % MOD = A^{Y} % MOD and minimal such Y = N % phi (MOD).

Trick #5:

(A * B) % MOD where MOD can't be present in int data type

```
function mulmod(A, B, MOD) {
    RES = 0;
    while (B > 0) {
        if (B is odd) {
            RES = (RES + A) %

MOD;
    }
    A = (A * 2) % MOD;
    B = B / 2;
}
return RES;
```

Conditions: 2 * MOD can be present in a standart data type.

Advices: use this trick only if (A % MOD) * (B % MOD) can't be present in any standart data type because of overflow and you don't want to use BigIntegers. But keep in mind that it works in O(logB) operations, not in O(1)as (A % MOD) * (B % MOD).

Proof: if B is even then $A \cdot B = 2 \cdot A \cdot (B/2)$, otherwise $A \cdot B = A + A \cdot (B-1)$.

F.A.Q. (in PM) - Alex Danilyuk (CF handle: Um_nik)

Q1: I'm a newbie. What should I do to become great coder?

A1: Stop doing competitive programming Solve problems

Q2: I'm doing CP for two months and I'm still not red green. What should I do?

A2: You are lazy and impatient Solve more problems

Q3: You became a red in less than two years, it is unbelievable!

A3: No, it isn't. You can do it too if you will solve [* * * ing troblems

Q4: You became a red blah-blah such a huge fan blah-blah. Oh, and what should I do to become as great as you?

A4: ... right. You already know the answer. Solve problems. Unate you.

Q5: I'm not good at DP [or something]. What can you suggest?

A5: Maybe you should try to stop asking stuppd questions and solve some problems on DP? Or read some blogs and editorials.

Q6: I can't solve a problem / understand your code. Can you help me?

(Well, it is not a bad question in general. It is a good (if you really want me to explain something not to write the solution instead of you) question, But...)

A6: Sure. Can you provide the link to the problem / code? (I'm not joking, it is a real story).

Bonus!

Q0: Hello bro/sir.

A0: Stop doing this please