

# CMSC 330, Summer 2018 Quiz 4 **Solution**

Name \_\_\_\_\_

## Instructions

- Do not start this quiz until you are told to do so.
- You have 20 minutes for this quiz.
- This is a closed book quiz. No notes or other aids are allowed.
- For partial credit, show all your work and clearly indicate your answers.

1. (4 points) Reduce the following  $\lambda$  term to normal form. If it does not have one, write “no normal form”.

$$\begin{aligned}
 & (\lambda x. \lambda y. x \ y \ y) (\lambda z. y) \ a \rightarrow_{\alpha} (\lambda x. \lambda w. x \ w \ w) (\lambda z. y) \ a \\
 & \rightarrow_{\beta} (\lambda w. (\lambda z. y) \ w \ w) \ a \\
 & \rightarrow_{\beta} (\lambda z. y) \ a \ a \\
 & \rightarrow_{\beta} y \ a
 \end{aligned}$$

2. (4 points) Here are the call-by-value substitution semantics for the  $\lambda$  calculus.

$$\begin{aligned}
 & (1) \frac{}{(\lambda x. t) \Downarrow_v (\lambda x. t)} \\
 & (2) \frac{t_1 \Downarrow_v (\lambda x. t_{12}) \quad t_2 \Downarrow_v v_2 \quad t_{12}[x \mapsto v_2] \Downarrow_v v}{(t_1 \ t_2) \Downarrow_v v}
 \end{aligned}$$

Write rules for a call-by-name  $\lambda$  calculus semantics. (Hint: You need only fill in the box for the hypotheses of rule (2).)

$$\begin{aligned}
 & (1) \frac{}{(\lambda x. t) \Downarrow_n (\lambda x. t)} \\
 & (2) \frac{t_1 \Downarrow_v (\lambda x. t_{12}) \quad t_{12}[x \mapsto t_2] \Downarrow_v v}{(t_1 \ t_2) \Downarrow_n v}
 \end{aligned}$$

3. (4 points) If possible, evaluate  $(\lambda x. \lambda y. y) ((\lambda z. z \ z) (\lambda z. z \ z))$  with  $\Downarrow_v$  (i.e. call-by-value evaluation) and prove your result with a derivation (a proof tree). If it does not evaluate explain why not.

Since  $\Downarrow_v$  is call-by-value and requires the argument to be evaluated before the substitution, this entire thing cannot evaluate as  $(\lambda z. z \ z) (\lambda z. z \ z)$  is the divergent combinator cannot be evaluated under  $\Downarrow_v$ .

4. (4 points) If possible, evaluate  $(\lambda x. \lambda y. y) ((\lambda z. z \ z) (\lambda z. z \ z))$  with  $\Downarrow_n$  (i.e. call-by-name evaluation) and prove your result with a derivation (a proof tree). If it does not evaluate explain why not.

$$\frac{\frac{}{(\lambda x. \lambda y. y) \Downarrow_n (\lambda x. \lambda y. y)}}{\quad} \quad \frac{}{(\lambda y. y)[x \mapsto (\lambda z. z \ z) (\lambda z. z \ z)] = (\lambda y. y) \Downarrow_n (\lambda y. y)}}{\quad} \\ \frac{}{(\lambda x. \lambda y. y) ((\lambda z. z \ z) (\lambda z. z \ z)) \Downarrow_n (\lambda y. y)}$$

5. (4 points) Describe the confluence property of the  $\lambda$  calculus. Are your results from (3) and (4) consistent with confluence?

Confluence states that there is a unique normal form, independent of reduction strategy. This is consistent since CBV did not yield a normal form and therefore confluence doesn't apply here.