

## \* Deadlock:

A deadlock can be defined as the permanent blocking of a set of processes that either compete for system resources or communicate with each other. Or deadlock involves conflicting needs for resources by two or more processes.

eg: In a multiprogramming system, suppose two processes are there & each want to print large tape files. Process A requests permission to use the printer & is granted. Process B then requests permission to use the tape drive & is also granted. Now A asks for the tape drive but the request is denied until B releases it. Instead of releasing the tape drive, B asks for the printer. At this point both processes are blocked & will remain so forever. This situation is deadlock.



## \* Conditions of Deadlock:

### a) Mutual exclusion:

The resources which require only mutually exclusive access, may give rise to a deadlock. These resource types cannot allow multiple processes to access it at the same time. For eg, a memory location, if allocated to a process, cannot be allowed to some other process. A printer cannot print the output of two processes at the same time.

### b) Hold and Wait:

There must exist a process that is holding at least one resource & is waiting to acquire additional resources that are currently being held by another process. Process currently holding resources granted earlier can request new resources.

### c) No preemption:

Resources cannot be preempted, that is a resource can only be released voluntarily by the process holding it, after the process has completed its task. Resources previously granted cannot be forcibly taken away from a process. They must be explicitly released by the process holding them.

### d) Circular wait:

There exist a set  $(P_0, P_1, \dots, P_n)$  of waiting processes such that  $P_0$  is waiting for a resource which is held by  $P_1$ ,  $P_1$  is waiting for a resource which is held by  $P_2$ ,  $P_{n-1}$  is waiting for a resource which is held by  $P_n$ , &  $P_n$  is waiting for a resource which is held by  $P_0$ . Thus there must be a circular chain of two or more processes, each of which is waiting for a resource held by the next member of the chain.



## → Resource Allocation Graph (RAG):

Deadlock can be described more precisely in terms of a directed graph called system resource allocation graph. The RAG follows the following notations & rules.

(P) → Process

[ ] R → Resource

[ • • ] → Number of instances of that resource

(P) → [ • • ]  
R → process has requested this resource but it has not been allocated. Known as request edge.

[ • • ] → (P)  
R → one instance of this resource has been allocated to the process. Known as assignment node.

