

LLVM Cookbook

Over 80 engaging recipes that will help you build a compiler frontend, optimizer, and code generator using LLVM



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Mayur Pandey Suyog Sarda



BIRMINGHAM - MUMBAI

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I would like to thank my family and friends. They made it possible for me to complete the book by taking care of my other commitments and always encouraging me.

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I would like to thank my family and friends. I would also like to thank the LLVM open-source community for always being helpful.

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I would like to thank my wife for supporting me every day with her smiles and love. I would also like to thank the entire LLVM community for all the hard work they have put into LLVM/Clang and other LLVM projects. It is amazing to see how fast LLVM evolves.

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I would like to thank the author of this book, who has done a good job. Thanks to the editors of the book at Packt Publishing, who made this book perfect and offered me the opportunity to review such a nice book.

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Table of Contents

| Preface Preface | v |
|--|----|
| Chapter 1: LLVM Design and Use | 1 |
| Introduction | 1 |
| Understanding modular design | 2 |
| Cross-compiling Clang/LLVM | 5 |
| Converting a C source code to LLVM assembly | 7 |
| Converting IR to LLVM bitcode | 9 |
| Converting LLVM bitcode to target machine assembly | 11 |
| Converting LLVM bitcode back to LLVM assembly | 13 |
| Transforming LLVM IR | 14 |
| Linking LLVM bitcode | 17 |
| Executing LLVM bitcode | 18 |
| Using the C frontend Clang | 19 |
| Using the GO frontend | 23 |
| Using DragonEgg | 24 |
| Chapter 2: Steps in Writing a Frontend | 27 |
| Introduction | 27 |
| Defining a TOY language | 28 |
| Implementing a lexer | 29 |
| Defining Abstract Syntax Tree | 32 |
| Implementing a parser | 35 |
| Parsing simple expressions | 36 |
| Parsing binary expressions | 38 |
| Invoking a driver for parsing | 41 |
| Running lexer and parser on our TOY language | 42 |
| Defining IR code generation methods for each AST class | 43 |
| Generating IR code for expressions | 45 |

| Table of Contents — | |
|--|-------------|
| Generating IR code for functions | 46 |
| Adding IR optimization support | 49 |
| Chapter 3: Extending the Frontend and Adding JIT Support | 51 |
| Introduction | 51 |
| Handling decision making paradigms – if/then/else constructs | 52 |
| Generating code for loops | 58 |
| Handling user-defined operators – binary operators | 63 |
| Handling user-defined operators – unary operators | 68 |
| Adding JIT support | 73 |
| Chapter 4: Preparing Optimizations | 77 |
| Introduction | 77 |
| Various levels of optimization | 78 |
| Writing your own LLVM pass | 79 |
| Running your own pass with the opt tool | 82 |
| Using another pass in a new pass | 83 |
| Registering a pass with pass manager | 85 |
| Writing an analysis pass | 87 |
| Writing an alias analysis pass | 90 |
| Using other analysis passes | 93 |
| Chapter 5: Implementing Optimizations | 97 |
| Introduction | 97 |
| Writing a dead code elimination pass | 98 |
| Writing an inlining transformation pass | 102 |
| Writing a pass for memory optimization | 106 |
| Combining LLVM IR | 108 |
| Transforming and optimizing loops | 110 |
| Reassociating expressions | 112 |
| Vectorizing IR | 114 |
| Other optimization passes | 121 |
| Chapter 6: Target-independent Code Generator | 125 |
| Introduction | 12 5 |
| The life of an LLVM IR instruction | 126 |
| Visualizing LLVM IR CFG using GraphViz | 12 9 |
| Describing targets using TableGen | 136 |
| Defining an instruction set | 137 |
| Adding a machine code descriptor | 138 |
| Implementing the MachineInstrBuilder class | 142 |
| Implementing the MachineBasicBlock class | 142 |
| Implementing the MachineFunction class | 144 |

| | —— Table of Contents |
|--|----------------------|
| Writing an instruction selector | 145 |
| Legalizing SelectionDAG | 151 |
| Optimizing SelectionDAG | 158 |
| Selecting instruction from the DAG | 163 |
| Scheduling instructions in SelectionDAG | 170 |
| Chapter 7: Optimizing the Machine Code | 173 |
| Introduction | 173 |
| Eliminating common subexpression from machine code | 174 |
| Analyzing live intervals | 184 |
| Allocating registers | 190 |
| Inserting the prologue-epilogue code | 196 |
| Code emission | 200 |
| Tail call optimization | 202 |
| Sibling call optimisation | 205 |
| Chapter 8: Writing an LLVM Backend | 207 |
| Introduction | 207 |
| Defining registers and registers sets | 208 |
| Defining the calling convention | 210 |
| Defining the instruction set | 211 |
| Implementing frame lowering | 212 |
| Printing an instruction | 215 |
| Selecting an instruction | 219 |
| Adding instruction encoding | 222 |
| Supporting a subtarget | 224 |
| Lowering to multiple instructions | 226 |
| Registering a target | 228 |
| Chapter 9: Using LLVM for Various Useful Projects | 241 |
| Introduction | 241 |
| Exception handling in LLVM | 242 |
| Using sanitizers | 247 |
| Writing the garbage collector with LLVM | 249 |
| Converting LLVM IR to JavaScript | 256 |
| Using the Clang Static Analyzer | 257 |
| Using bugpoint | 259 |
| Using LLDB | 262 |
| Using LLVM utility passes | 267 |
| Index | 271 |

Preface

A programmer might have come across compilers at some or the other point when programming. Simply speaking, a compiler converts a human-readable, high-level language into machine-executable code. But have you ever wondered what goes on under the hood? A compiler does lot of processing before emitting optimized machine code. Lots of complex algorithms are involved in writing a good compiler.

This book travels through all the phases of compilation: frontend processing, code optimization, code emission, and so on. And to make this journey easy, LLVM is the simplest compiler infrastructure to study. It's a modular, layered compiler infrastructure where every phase is dished out as a separate recipe. Written in object-oriented C++, LLVM gives programmers a simple interface and lots of APIs to write their own compiler.

As authors, we maintain that simple solutions frequently work better than complex solutions; throughout this book, we'll look at a variety of recipes that will help develop your skills, make you consider all the compiling options, and understand that there is more to simply compiling code than meets the eye.

We also believe that programmers who are not involved in compiler development will benefit from this book, as knowledge of compiler implementation will help them code optimally next time they write code.

We hope you will find the recipes in this book delicious, and after tasting all the recipes, you will be able to prepare your own dish of compilers. Feeling hungry? Let's jump into the recipes!

What this book covers

Chapter 1, LLVM Design and Use, introduces the modular world of LLVM infrastructure, where you learn how to download and install LLVM and Clang. In this chapter, we play with some examples to get accustomed to the workings of LLVM. We also see some examples of various frontends.

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Chapter 2, Steps in Writing a Frontend, explains the steps to write a frontend for a language. We will write a bare-metal toy compiler frontend for a basic toy language. We will also see how a frontend language can be converted into the LLVM intermediate representation (IR).

Chapter 3, Extending the Frontend and Adding JIT Support, explores the more advanced features of the toy language and the addition of JIT support to the frontend. We implement some powerful features of a language that are found in most modern programming languages.

Chapter 4, Preparing Optimizations, takes a look at the pass infrastructure of the LLVM IR. We explore various optimization levels, and the optimization techniques kicking at each level. We also see a step-by-step approach to writing our own LLVM pass.

Chapter 5, Implementing Optimizations, demonstrates how we can implement various common optimization passes on LLVM IR. We also explore some vectorization techniques that are not yet present in the LLVM open source code.

Chapter 6, Target-independent Code Generator, takes us on a journey through the abstract infrastructure of a target-independent code generator. We explore how LLVM IR is converted to Selection DAGs, which are further processed to emit target machine code.

Chapter 7, Optimizing the Machine Code, examines how Selection DAGs are optimized and how target registers are allocated to variables. This chapter also describes various optimization techniques on Selection DAGs as well as various register allocation techniques.

Chapter 8, Writing an LLVM Backend, takes us on a journey of describing a target architecture. This chapter covers how to describe registers, instruction sets, calling conventions, encoding, subtarget features, and so on.

Chapter 9, Using LLVM for Various Useful Projects, explores various other projects where LLVM IR infrastructure can be used. Remember that LLVM is not just a compiler; it is a compiler infrastructure. This chapter explores various projects that can be applied to a code snippet to get useful information from it.

What you need for this book

All you need to work through most of the examples covered in this book is a Linux machine, preferably Ubuntu. You will also need a simple text or code editor, Internet access, and a browser. We recommend installing the meld tool for comparison of two files; it works well on the Linux platform.

Who this book is for

The book is for compiler programmers who are familiar with concepts of compilers and want to indulge in understanding, exploring, and using LLVM infrastructure in a meaningful way in their work.

This book is also for programmers who are not directly involved in compiler projects but are often involved in development phases where they write thousands of lines of code. With knowledge of how compilers work, they will be able to code in an optimal way and improve performance with clean code.

Sections

In this book, you will find several headings that appear frequently (Getting ready, How to do it, How it works, There's more, and See also).

To give clear instructions on how to complete a recipe, we use these sections.

Getting ready

This section tells you what to expect in the recipe, and describes how to set up any software or any preliminary settings required for the recipe.

How to do it...

This section contains the steps required to follow the recipe.

How it works...

This section usually consists of a detailed explanation of what happened in the previous section.

There's more...

This section consists of additional information about the recipe in order to make you more knowledgeable about the recipe.

See also

This section provides helpful links to other useful information for the recipe.

Conventions

In this book, you will find a number of text styles that distinguish between different kinds of information. Here are some examples of these styles and an explanation of their meaning.

Code words in text, database table names, folder names, filenames, file extensions, pathnames, dummy URLs, user input, and Twitter handles are shown as follows: "We can include other contexts through the use of the include directive."

A block of code is set as follows:

```
primary := identifier_expr
:=numeric_expr
:=paran expr
```

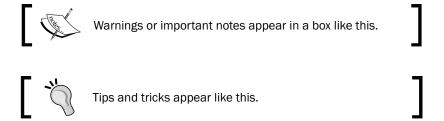
When we wish to draw your attention to a particular part of a code block, the relevant lines or items are set in bold:

```
primary := identifier_expr
:=numeric_expr
:=paran expr
```

Any command-line input or output is written as follows:

\$ cat testfile.ll

New terms and **important words** are shown in bold. Words that you see on the screen, for example, in menus or dialog boxes, appear in the text like this: "Clicking on the **Next** button moves you to the next screen."



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1

LLVM Design and Use

In this chapter, we will cover the following topics:

- Understanding modular design
- Cross-compiling Clang/LLVM
- Converting a C source code to LLVM assembly
- ▶ Converting IR to LLVM bitcode
- Converting LLVM bitcode to target machine assembly
- ► Converting LLVM bitcode back to LLVM assembly
- Transforming LLVM IR
- ▶ Linking LLVM bitcode
- Executing LLVM bitcode
- Using C frontend Clang
- Using the GO frontend
- Using DragonEgg

Introduction

In this recipe, you get to know about **LLVM**, its design, and how we can make multiple uses out of the various tools it provides. You will also look into how you can transform a simple C code to the LLVM intermediate representation and how you can transform it into various forms. You will also learn how the code is organized within the LLVM source tree and how can you use it to write a compiler on your own later.

Understanding modular design

LLVM is designed as a set of libraries unlike other compilers such as **GNU Compiler Collection** (**GCC**). In this recipe, LLVM optimizer will be used to understand this design. As LLVM optimizer's design is library-based, it allows you to order the passes to be run in a specified order. Also, this design allows you to choose which optimization passes you can run—that is, there might be a few optimizations that might not be useful to the type of system you are designing, and only a few optimizations will be specific to the system. When looking at traditional compiler optimizers, they are built as a tightly interconnected mass of code, that is difficult to break down into small parts that you can understand and use easily. In LLVM, you need not know about how the whole system works to know about a specific optimizer. You can just pick one optimizer and use it without having to worry about other components attached to it.

Before we go ahead and look into this recipe, we must also know a little about LLVM assembly language. The LLVM code is represented in three forms: in memory compiler **Intermediate Representation** (**IR**), on disk bitcode representation, and as human readable assembly. LLVM is a **Static Single Assignment** (**SSA**)-based representation that provides type safety, low level operations, flexibility, and the capability to represent all the high-level languages cleanly. This representation is used throughout all the phases of LLVM compilation strategy. The LLVM representation aims to be a universal IR by being at a low enough level that high-level ideas may be cleanly mapped to it. Also, LLVM assembly language is well formed. If you have any doubts about understanding the LLVM assembly mentioned in this recipe, refer to the link provided in the See also section at the end of this recipe.

Getting ready

We must have installed the LLVM toolchain on our host machine. Specifically, we need the opt tool.

How to do it...

We will run two different optimizations on the same code, one-by-one, and see how it modifies the code according to the optimization we choose.

1. First of all, let us write a code we can input for these optimizations. Here we will write it into a file named testfile.11:

```
$ cat testfile.11
define i32 @test1(i32 %A) {
   %B = add i32 %A, 0
   ret i32 %B
}
```

```
define internal i32 @test(i32 %X, i32 %dead) {
     ret i32 %X
   }
   define i32 @caller() {
     %A = call i32 @test(i32 123, i32 456)
     ret i32 %A
   }
2. Now, run the opt tool for one of the optimizations—that is, for combining the
   instruction:
   $ opt -S -instcombine testfile.ll -o output1.ll
3. View the output to see how instcombine has worked:
   $ cat output1.11
   ; ModuleID = 'testfile.ll'
   define i32 @test1(i32 %A) {
     ret i32 %A
   }
   define internal i32 @test(i32 %X, i32 %dead) {
     ret i32 %X
   }
   define i32 @caller() {
     %A = call i32 @test(i32 123, i32 456)
     ret i32 %A
   }
4. Run the opt command for dead argument elimination optimization:
   $ opt -S -deadargelim testfile.11 -o output2.11
```

5. View the output, to see how deadargelim has worked:

```
$ cat output2.11
; ModuleID = testfile.11'

define i32 @test1(i32 %A) {
    %B = add i32 %A, 0
    ret i32 %B
}

define internal i32 @test(i32 %X) {
    ret i32 %X
}

define i32 @caller() {
    %A = call i32 @test(i32 123)
    ret i32 %A
}
```

How it works...

In the preceding example, we can see that, for the first command, the instcombine pass is run, which combines the instructions and hence optimizes %B = add i32 %A, 0; ret i32 %B to ret i32 %A without affecting the code.

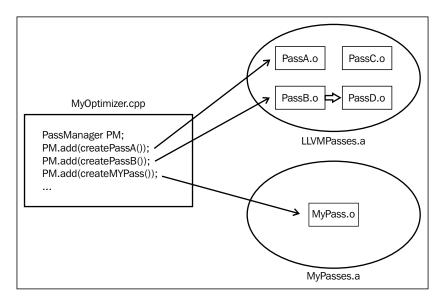
In the second case, when the deadargelim pass is run, we can see that there is no modification in the first function, but the part of code that was not modified last time gets modified with the function arguments that are not used getting eliminated.

LLVM optimizer is the tool that provided the user with all the different passes in LLVM. These passes are all written in a similar style. For each of these passes, there is a compiled object file. Object files of different passes are archived into a library. The passes within the library are not strongly connected, and it is the LLVM **PassManager** that has the information about dependencies among the passes, which it resolves when a pass is executed. The following image shows how each pass can be linked to a specific object file within a specific library. In the following figure, the **PassA** references **LLVMPasses.a** for **PassA.o**, whereas the custom pass refers to a different library **MyPasses.a** for the **MyPass.o** object file.

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There's more...

Similar to the optimizer, the LLVM code generator also makes use of its modular design, splitting the code generation problem into individual passes: instruction selection, register allocation, scheduling, code layout optimization, and assembly emission. Also, there are many built-in passes that are run by default. It is up to the user to choose which passes to run.

See also

- In the upcoming chapters, we will see how to write our own custom pass, where we can choose which of the optimization passes we want to run and in which order. Also, for a more detailed understanding, refer to http://www.aosabook.org/en/llvm.html.
- ► To understand more about LLVM assembly language, refer to http://llvm.org/docs/LangRef.html.

Cross-compiling Clang/LLVM

By cross-compiling we mean building a binary on one platform (for example, x86) that will be run on another platform (for example, ARM). The machine on which we build the binary is called the host, and the machine on which the generated binary will run is called the target. The compiler that builds code for the same platform on which it is running (the host and target platforms are the same) is called a **native assembler**, whereas the compiler that builds code for a target platform different from the host platform is called a **cross-compiler**.

In this recipe, cross-compilation of LLVM for a platform different than the host platform will be shown, so that you can use the built binaries for the required target platform. Here, cross-compiling will be shown using an example where cross-compilation from host platform x86_64 for target platform ARM will be done. The binaries thus generated can be used on a platform with ARM architecture.

Getting ready

The following packages need to be installed on your system (host platform):

- ▶ cmake
- ninja-build (from backports in Ubuntu)
- ▶ gcc-4.x-arm-linux-gnueabihf
- ▶ gcc-4.x-multilib-arm-linux-gnueabihf
- ▶ binutils-arm-linux-gnueabihf
- ▶ libgccl-armhf-cross
- ▶ libsfgcc1-armhf-cross
- ▶ libstdc++6-armhf-cross
- ▶ libstdc++6-4.x-dev-armhf-cross
- ▶ install llvm on your host platform

How to do it...

To compile for the ARM target from the host architecture, that is **X86_64** here, you need to perform the following steps:

- 1. Add the following cmake flags to the normal cmake build for LLVM:
 - -DCMAKE CROSSCOMPILING=True
 - -DCMAKE_INSTALL_PREFIX= path-where-you-want-the-toolchain(optional)
 - -DLLVM_TABLEGEN=<path-to-host-installed-llvm-toolchain-bin>/llvm-tblgen
 - -DCLANG_TABLEGEN=< path-to-host-installed-llvm-toolchain-bin >/ clang-tblgen
 - -DLLVM DEFAULT TARGET TRIPLE=arm-linux-gnueabihf
 - -DLLVM TARGET ARCH=ARM
 - -DLLVM TARGETS TO BUILD=ARM
 - -DCMAKE_CXX_FLAGS='-target armv7a-linuxgnueabihf -mcpu=cortex-a9 -I/usr/arm-linux-gnueabihf/include/
 c++/4.x.x/arm-linux-gnueabihf/ -I/usr/arm-linux-gnueabihf/
 include/ -mfloat-abi=hard -ccc-gcc-name arm-linux-gnueabihf-gcc'

- 2. If using your platform compiler, run:
 - \$ cmake -G Ninja <llvm-source-dir> <options above>

If using Clang as the cross-compiler, run:

\$ CC='clang' CXX='clang++' cmake -G Ninja <source-dir> <options
above>

If you have clang/Clang++ on the path, it should work fine.

- 3. To build LLVM, simply type:
 - \$ ninja
- 4. After the LLVM/Clang has built successfully, install it with the following command:
 - \$ ninja install

This will create a sysroot on the install-dir location if you have specified the DCMAKE_ INSTALL PREFIX options

How it works...

The cmake package builds the toolchain for the required platform by making the use of option flags passed to cmake, and the tblgen tools are used to translate the target description files into C++ code. Thus, by using it, the information about targets is obtained, for example—what instructions are available on the target, the number of registers, and so on.



If Clang is used as the cross-compiler, there is a problem in the LLVM ARM backend that produces absolute relocations on **position-independent code** (**PIC**), so as a workaround, disable PIC at the moment.

The ARM libraries will not be available on the host system. So, either download a copy of them or build them on your system.

Converting a C source code to LLVM assembly

Here we will convert a C code to intermediate representation in LLVM using the C frontend Clang.

Getting ready

Clang must be installed in the PATH.

How to do it...

1. Lets create a C code in the multiply.c file, which will look something like the following:

```
$ cat multiply.c
int mult() {
int a =5;
int b = 3;
int c = a * b;
return c;
}
```

2. Use the following command to generate LLVM IR from the C code:

```
$ clang -emit-llvm -S multiply.c -o multiply.ll
```

3. Have a look at the generated IR:

```
$ cat multiply.11
; ModuleID = 'multiply.c'
target datalayout = "e-m:e-i64:64-f80:128-n8:16:32:64-S128"
target triple = "x86_64-unknown-linux-gnu"
; Function Attrs: nounwind uwtable
define i32 @mult() #0 {
  %a = alloca i32, align 4
  %b = alloca i32, align 4
  %c = alloca i32, align 4
  store i32 5, i32* %a, align 4
  store i32 3, i32* %b, align 4
  %1 = load i32* %a, align 4
  %2 = load i32* %b, align 4
  %3 = mul nsw i32 %1, %2
  store i32 %3, i32* %c, align 4
  %4 = load i32* %c, align 4
  ret i32 %4
}
```

We can also use the ccl for generating IR:

```
$ clang -cc1 -emit-llvm testfile.c -o testfile.ll
```

How it works...

The process of C code getting converted to IR starts with the process of lexing, wherein the C code is broken into a token stream, with each token representing an Identifier, Literal, Operator, and so on. This stream of tokens is fed to the parser, which builds up an abstract syntax tree with the help of **Context free grammar** (**CFG**) for the language. Semantic analysis is done afterwards to check whether the code is semantically correct, and then we generate code to IR.

Here we use the Clang frontend to generate the IR file from C code.

See also

► In the next chapter, we will see how the lexer and parser work and how code generation is done. To understand the basics of LLVM IR, you can refer to http://llvm.org/docs/LangRef.html.

Converting IR to LLVM bitcode

In this recipe, you will learn to generate LLVM bit code from IR. The LLVM bit code file format (also known as bytecode) is actually two things: a bitstream container format and an encoding of LLVM IR into the container format.

Getting Ready

The 11vm-as tool must be installed in the PATH.

How to do it...

Do the following steps:

1. First create an IR code that will be used as input to 11vm-as:

```
$ cat test.11
define i32 @mult(i32 %a, i32 %b) #0 {
%1 = mul nsw i32 %a, %b
ret i32 %1
}
```

To convert LLVM IR in test.11 to bitcode format, you need to use the following command:

```
llvm-as test.ll -o test.bc
```

3. The output is generated in the test.bc file, which is in bit stream format; so, when we want to have a look at output in text format, we get it as shown in the following screenshot:

Since this is a bitcode file, the best way to view its content would be by using the hexdump tool. The following screenshot shows the output of hexdump:

```
$ hexdump -C test.bc
00000000 42 43 c0 de 21 0c 00 00
                                  68 00 00 00 0b 82 20 00
                                                           [BC..!...h.......
00000010
         02 00 00 00 13 00 00 00
                                  07 81 23 91 41 c8 04 49
                                                           |....#.A..I|
00000020
         06 10 32 39 92 01 84 0c
                                  25 05 08 19 1e 04 8b 62
                                                           |...29....%.....b|
00000030
         80 0c 45
                  02 42 92 0b 42
                                           14 38 08 18 4b
                                                           [..E.B..Bd.2.8..K]
                                  64 10 32
         0a 32 32 88 48 90 14 20
00000040
                                  43 46 88 a5 00 19 32 42
                                                            |.22.H.. CF....2B|
00000050
         e4 48 0e 90 91 21 c4 50
                                  41 51 81 8c e1 83 e5 8a
                                                            |.H...!.PAQ.....
00000060
         04 19 46 06 89 20 00 00
                                  0b 00 00 00 32 22 c8 08
                                                           |..F.. .....2"...|
         20 64 85 04 93 21 a4 84
                                  04 93 21 e3 84 a1 90 14
                                                           | d...!....!....
00000070
00000080
         12 4c 86 8c 0b 84 64 4c
                                  10 14 73 04 60 50 06 00
                                                           |.L....dL..s.`P...|
00000090
         94 81 80 11 00 00 00 00
                                  43 1c 01 00 00 00 00 00
                                                           .......2...3...
000000a0
         00 00 00 00 c8 c3 00 00
                                  32 00 00 00 33 08 80 1c
00000060
         c4 e1 1c 66 14 01 3d 88
                                  43 38 84 c3 8c 42 80 07
                                                           |...f..=.C8...B..
000000c0
         79
            78 07
                  73 98 71 0c e6
                                  00 Of
                                        ed 10 0e f4 80 0e
                                                            |yx.s.q.....
         33 0c 42 1e c2 c1 1d ce
                                                           |3.B....f0.=.C|
000000d0
                                  a1 1c 66 30 05 3d 88 43
000000e0
         38 84 83 1b cc 03 3d c8
                                  43 3d 8c 03 3d cc 78 8c
                                                           8.....=.C=..=.x.
000000f0
         74 70 07 7b 08 07 79 48
                                  87 70 70 07 7a 70 03 76
                                                           |tp.{..yH.pp.zp.v|
00000100
         78 87 70 20 87 19 cc 11
                                  0e ec 90 0e e1 30 0f 6e
                                                            | x.p ....0.n
                                  33 10 c4 1d de 21 1c d8
                                                            |0.....P.3....!..
00000110
         30 Of e3 f0 Oe f0 50 Oe
         21 1d c2
                                           3b d0 43 39 b4
00000120
                  61 1e 66 30 89
                                  3b bc 83
                                                           |!..a.f0.;..;.C9.|
00000130
         03
            3c bc
                  83 3c 84 03 3b
                                  cc
                                     f0
                                        14
                                           76 60 07
                                                    7b 68
                                                           |.<..<..;...v`.{h|
00000140
         07 37 68 87 72 68 07 37
                                  80 87 70 90 87 70 60 07
                                                            |.7h.rh.7..p..p
00000150
         76 28 07 76 f8 05 76 78
                                  87 77 80 87 5f 08 87 71
                                                            |v(.v..vx.w.._..q|
                                                            |..г..у..,.....
00000160
         18 87 72 98 87 79 98 81
                                  2c ee f0 0e ee e0 0e f5
00000170
         c0 0e ec 00 71 20 00 00
                                  02 00 00 00 06 40 30 d4
                                                            ....q .......@0.
00000180
         32 01 00 00 61 20 00 00
                                  07 00 00 00 13 04 81 09
                                                            2...a ......
00000190
         81 08 32 08 07 02 00 00
                                  02 00 00 00 16 10 00 26
                                                            |..2.......
000001a0
         10 04 00 00 00 00 00 00 00 00 00 00
                                                            [......
000001ac
```

How it works...

The 11vm-as is the LLVM assembler. It converts the LLVM assembly file that is the LLVM IR into LLVM bitcode. In the preceding command, it takes the test.11 file as the input and outputs, and test.bc as the bitcode file.

There's more...

To encode LLVM IR into bitcode, the concept of blocks and records is used. Blocks represent regions of bitstream, for example—a function body, symbol table, and so on. Each block has an ID specific to its content (for example, function bodies in LLVM IR are represented by ID 12). Records consist of a record code and an integer value, and they describe the entities within the file such as instructions, global variable descriptors, type descriptions, and so on.

Bitcode files for LLVM IR might be wrapped in a simple wrapper structure. This structure contains a simple header that indicates the offset and size of the embedded BC file.

See also

► To get a detailed understanding of the LLVM the bitstream file format, refer to http://llvm.org/docs/BitCodeFormat.html#abstract

Converting LLVM bitcode to target machine assembly

In this recipe, you will learn how to convert the LLVM bitcode file to target specific assembly code.

Getting ready

The LLVM static compiler 11c should be in installed from the LLVM toolchain.

How to do it...

Do the following steps:

 The bitcode file created in the previous recipe, test.bc, can be used as input to llc here. Using the following command, we can convert LLVM bitcode to assembly code:

\$ 11c test.bc -o test.s