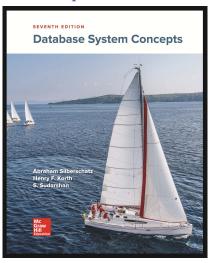
Database System Concepts, 7th Edition Chapter 4: Intermediate SQL

Silberschatz, Korth and Sudarshan

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Database System Concepts



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Joined Relations

- ▶ Join operations take two relations and return as a result another relation.
- ▶ A join operation is a Cartesian product which requires that tuples in the two relations match (under some condition). It also specifies the attributes that are present in the result of the join.
- ► The join operations are typically used as subquery expressions in the FROM clause.
- ► Three types of joins:
 - ► Natural join
 - ► Inner join
 - ▶ Outer join

Natural Join in SQL

- ▶ Natural join matches tuples with the same values for all common attributes, and retains only one copy of each common column.
- ▶ For all students in the university who have taken some course, find their names and the course ID of all courses they took.

```
SELECT
name, course_id
FROM
students, takes
WHERE
student.ID = takes.ID;
```

Natural Join in SQL

▶ Same query in SQL with "natural join" construct.

```
1 SELECT
2 name, course_id
3 FROM
4 student NATURAL JOIN takes;
```

Natural Join in SQL (Cont.)

► The FROM clause can have multiple relations combined using natural join:

```
SELECT A_1, A_2, \ldots, A_n
FROM r_1
NATURAL JOIN r_2
NATURAL JOIN \ldots
NATURAL JOIN r_n
WHERE P;
```

Student Relation

ID	name	dept_name	tot_cred
00128	Zhang	Comp. Sci.	102
12345	Shankar	Comp. Sci.	32
19991	Brandt	History	80
23121	Chavez	Finance	110
44553	Peltier	Physics	56
45678	Levy	Physics	46
54321	Williams	Comp. Sci.	54
55739	Sanchez	Music	38
70557	Snow	Physics	0
76543	Brown	Comp. Sci.	58
76653	Aoi	Elec. Eng.	60
98765	Bourikas	Elec. Eng.	98
98988	Tanaka	Biology	120

Takes Relation

ID	course_id	sec_id	semester	year	grade
00128	CS-101	1	Fall	2017	A
00128	CS-347	1	Fall	2017	A-
12345	CS-101	1	Fall	2017	C
12345	CS-190	2	Spring	2017	A
12345	CS-315	1	Spring	2018	A
12345	CS-347	1	Fall	2017	A
19991	HIS-351	1	Spring	2018	В
23121	FIN-201	1	Spring	2018	C+
44553	PHY-101	1	Fall	2017	B-
45678	CS-101	1	Fall	2017	F
45678	CS-101	1	Spring	2018	B+
45678	CS-319	1	Spring	2018	В
54321	CS-101	1	Fall	2017	A-
54321	CS-190	2	Spring	2017	B+
55739	MU-199	1	Spring	2018	A-
76543	CS-101	1	Fall	2017	A
76543	CS-319	2	Spring	2018	A
76653	EE-181	1	Spring	2017	C
98765	CS-101	1	Fall	2017	C-
98765	CS-315	1	Spring	2018	В
98988	BIO-101	1	Summer	2017	A
98988	BIO-301	1	Summer	2018	null

Student Natural Join Takes

ID	name	dept_name	tot_cred	course_id	sec_id	semester	year	grade
00128	Zhang	Comp. Sci.	102	CS-101	1	Fall	2017	A
00128	Zhang	Comp. Sci.	102	CS-347	1	Fall	2017	A-
12345	Shankar	Comp. Sci.	32	CS-101	1	Fall	2017	C
12345	Shankar	Comp. Sci.	32	CS-190	2	Spring	2017	A
12345	Shankar	Comp. Sci.	32	CS-315	1	Spring	2018	A
12345	Shankar	Comp. Sci.	32	CS-347	1	Fall	2017	A
19991	Brandt	History	80	HIS-351	1	Spring	2018	В
23121	Chavez	Finance	110	FIN-201	1	Spring	2018	C+
44553	Peltier	Physics	56	PHY-101	1	Fall	2017	B-
45678	Levy	Physics	46	CS-101	1	Fall	2017	F
45678	Levy	Physics	46	CS-101	1	Spring	2018	B+
45678	Levy	Physics	46	CS-319	1	Spring	2018	В
54321	Williams	Comp. Sci.	54	CS-101	1	Fall	2017	A-
54321	Williams	Comp. Sci.	54	CS-190	2	Spring	2017	B+
55739	Sanchez	Music	38	MU-199	1	Spring	2018	A-
76543	Brown	Comp. Sci.	58	CS-101	1	Fall	2017	A
76543	Brown	Comp. Sci.	58	CS-319	2	Spring	2018	A
76653	Aoi	Elec. Eng.	60	EE-181	1	Spring	2017	C
98765	Bourikas	Elec. Eng.	98	CS-101	1	Fall	2017	C-
98765	Bourikas	Elec. Eng.	98	CS-315	1	Spring	2018	В
98988	Tanaka	Biology	120	BIO-101	1	Summer	2017	A
98988	Tanaka	Biology	120	BIO-301	1	Summer	2018	null

Beware of unrelated attributes with same name which get equated incorrectly

Example:

List the names of students along with the titles of courses that they have taken.

```
1 SELECT
2 name, title
3 FROM
4 student
5 natural join
6 takes
7 natural join
8 course;
```

Beware of unrelated attributes with same name which get equated incorrectly

Example:

List the names of students along with the titles of courses that they have taken.

```
SELECT
name, title
FROM
student
natural join
takes
natural join
course;
```

Incorrect! double check dept_name attribute

Example:

List the names of students along with the titles of courses that they have taken.

```
1 SELECT
2 name, title
3 FROM
4 student
5 NATURAL JOIN
6 takes
7 NATURAL JOIN
8 course;
```

▶ This query omits all (student name, course title) pairs where the student takes a course in a department other than the student's own department.

Example:

List the names of students along with the titles of courses that they have taken.

```
1 SELECT
2 name, title
3 FROM
4 student
5 NATURAL JOIN
6 takes, course
7 WHERE
8 takes.course_id = course.course_id;
```

▶ The correct version (above), correctly outputs such pairs.

Outer Join

- ► An extension of the join operation that avoids loss of information.
- ► Computes the JOIN and then adds tuples form one relation that does not match tuples in the other relation to the result of the join.
- ▶ Uses null values.
- ► Three forms of outer join:
 - left outer join
 - right outer join
 - ▶ full outer join

Outer Join Examples

▶ Relation *course*:

$course_id$	title	$\mathbf{dept}_{-}\mathbf{name}$	credits
BIO-301	Genetics	Biology	4
CS-190	Game Design	Comp. Sci.	4
CS-315	Robotics	Comp. Sci.	3

▶ Relation *prereq*:

$course_id$	$prereq_id$
BIO-301	BIO-101
CS-190	CS-101
CS-347	CS-101

- ▶ Observe that
 - course information is missing for CS-347
 - ▶ prereq information is missing for CS-315

Left Outer Join

► course NATURAL LEFT OUTER JOIN prereq

course_id	title	$dept_name$	credits	$prereq_id$
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null

▶ In relational algebra: course ⋈ prereq

Right Outer Join

► course NATURAL RIGHT OUTER JOIN prereq

$course_id$	title	$dept_name$	credits	$\mathbf{prereq_id}$
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-347	null	null	null	CS-101

▶ In relational algebra: course ⋈ prereq

Full Outer Join

► course NATURAL FULL OUTER JOIN prereq

$course_id$	title	$dept_name$	credits	$\operatorname{prereq_id}$
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null
CS-347	null	null	null	CS-101

► In relational algebra: course ⋈ prereq

Joined Types and Conditions

- ▶ **Join operations** take two relations and return as a result another relation.
- ► These additional operations are typically used as subquery expressions in the FROM clause
- ▶ Join condition defines which tuples in the two relations match, and what attributes are present in the result of the join.
- ▶ Join type defines how tuples in each relation that do not match any tuple in the other relation (based on the join condition) are treated.

Join types
inner join
left outer join
right outer join
full outer join

Join conditions
natural
on < predicate>
using $(A_1, A_2,, A_n)$

Joined Relations – Examples

► course NATURAL RIGHT OUTER JOIN prereq

course_id	title	$dept_name$	credits	prereq_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-347	null	null	null	CS-101

► course FULL OUTER JOIN prereq USING (course_id)

course_id	title	dept_name	credits	prereq_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null
CS-347	null	null	null	CS-101

Joined Relations – Examples

course INNER JOIN prereq ON course.course_id =
prereq.course_id

course_id	title	dept_name	credits	prereq_id	course_id
BIO-301	Genetics	Biology	4	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190

- ▶ What is the difference between the above, and a natural join?
- course LEFT OUTER JOIN prereq ON course.course_id = prereq.course_id

course_id	title	dept_name	credits	prereq_id	course_id
BIO-301	Genetics	Biology	4	BIO-101	BIO-301
CS-190	Game Design	Comp. Sci.	4	CS-101	CS-190
CS-315	Robotics	Comp. Sci.	3	null	null

Joined Relations – Examples

► course NATURAL RIGHT OUTER JOIN prereq

course_id	title	$dept_name$	credits	prereq_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-347	null	null	null	CS-101

► course FULL OUTER JOIN prereq USING (course_id)

course_id	title	dept_name	credits	prereq_id
BIO-301	Genetics	Biology	4	BIO-101
CS-190	Game Design	Comp. Sci.	4	CS-101
CS-315	Robotics	Comp. Sci.	3	null
CS-347	null	null	null	CS-101

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Views

- ▶ In some cases, it is not desirable for all users to see the entire logical model (that is, all the actual relations stored in the database.)
- ► Consider a person who needs to know an instructors name and department, but not the salary. This person should see a relation described, in SQL, by

```
SELECT

ID, name, dept_name

FROM

instructor;
```

- ► A **view** provides a mechanism to hide certain data from the view of certain users.
- Any relation that is not of the conceptual model but is made visible to a user as a "virtual relation" is called a view.

View Definition

► A view is defined using the create view statement which has the form:

CREATE VIEW v **AS** $< query_expression > ;$

where $< query_expression >$ is any legal SQL expression. The view name is represented by v.

- ▶ Once a view is defined, the view name can be used to refer to the virtual relation that the view generates.
- ▶ View definition is not the same as creating a new relation by evaluating the query expression
 - ▶ Rather, a view definition causes the saving of an expression; the expression is substituted into queries using the view.

View Definition and Use

▶ A view of instructors without their salary.

```
CREATE VIEW faculty AS
SELECT ID, name, dept_name
FROM instructor
```

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View Definition and Use

▶ A view of instructors without their salary.

```
CREATE VIEW faculty AS
SELECT ID, name, dept_name
FROM instructor
```

► Find all instructors in the Biology department

```
SELECT name
FROM faculty
WHERE dept_name = 'Biology'
```

View Definition and Use

▶ A view of instructors without their salary.

```
CREATE VIEW faculty AS
SELECT ID, name, dept_name
FROM instructor
```

► Find all instructors in the Biology department

```
SELECT name
FROM faculty
WHERE dept_name = 'Biology'
```

Create a view of department salary totals

```
CREATE VIEW departments_total_salary(dept_name, total_salary) AS

SELECT dept_name, SUM(salary)

FROM instructor

GROUP BY dept_name;
```

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Views Defined Using Other Views

- ▶ One view may be used in the expression defining another view.
- A view relation v_1 is said to depend directly on a view relation v_2 if v_2 is used in the expression defining v_1 .
- A view relation v_1 is said to depend on view relation v_2 if either v_1 depends directly to v_2 or there is a path of dependencies from v_1 to v_2 .
- ightharpoonup A view relation v is said to be recursive if it depends on itself.

Views Defined Using Other Views

```
CREATE VIEW
             physics_fall_2017 AS
         SELECT
             course.course_id, sec_id, building, room_number
        FROM
5
             course, section
6
         WHERE
             course.course_id = section.course_id AND
             course.dept_name = 'Physics' AND
9
             section.semester = 'Fall' AND
10
11
             section.year = '2017';
```

Views Defined Using Other Views

```
create View
physics_fall_2017_watson AS

SELECT
course_id, room_number

FROM
physics_fall_2017

WHERE
building= 'Watson';
```

View Expansion

Expand the view:

```
CREATE VIEW physics_fall_2017_watson AS

SELECT course_id, room_number

FROM physics_fall_2017

WHERE building= 'Watson';
```

To:

```
CREATE VIEW physics_fall_2017_watson AS
1
        SELECT course_id, room_number
2
3
        FROM ( SELECT course.course_id, sect_id, building, room_number
4
                 FROM course, section
5
                 WHERE course.course id = section.course id
                 AND course.dept_name = 'Physics'
6
                 AND section.semester = 'Fall'
8
                 AND section.year = '2017' )
9
        WHERE building= 'Watson';
```

View Expansion (Cont.)

- ▶ A way to define the meaning of views defined in terms of other views.
- Let view v_1 be defined by an expression e_1 that may itself contain uses of view relations.
- ▶ View expansion of an expression repeats the following replacement step:

REPEAT

Find any view relation v_i in e_i Replace the view relation v_i by the expression defining v_i UNTIL no more view relations are present in e_i

► As long as the view definitions are not recursive, this loop will terminate.

Materialized Views

- Certain database systems allow view relations to be physically stored.
 - ▶ Physical copy created when the view is defined.
 - ► Such views are called **Materialized view**.
- ▶ If relations used in the query are updated, the materialized view result becomes out of date.
 - ▶ Need to **maintain** the view, by updating the view whenever the underlying relations are updated.

Update of a View

▶ Add a new tuple to faculty view which we defined earlier.

```
INSERT INTO faculty
VALUES ('30765', 'Green', 'Music');
```

- ► This insertion must be represented by the insertion into the instructor relation.
 - Must have a value for salary.
- ► Two approaches:
 - ▶ Reject the insert.
 - ► Insert the tuple: ('30765', 'Green', 'Music', null) into the *instructor* relation.

Some Updates Cannot be Translated Uniquely

```
CREATE VIEW instructor_info AS

SELECT

ID, name, building

FROM

instructor, department

WHERE

instructor.dept_name= department.dept_name;
```

```
INSERT INTO instructor_info
VALUES ('69987', 'White', 'Taylor');
```

► Issues:

- ▶ Which department, if multiple departments in Taylor?
- ▶ What if no department is in Taylor?

And Some Not at All

```
CREATE VIEW history_instructors AS

SELECT *

FROM instructor

WHERE dept_name= 'History';
```

What happens if we insert: ('25566', 'Brown', 'Biology', 100000) into history_instructors?

View Updates in SQL

- Most SQL implementations allow updates only on simple views:
 - ▶ The FROM clause has only one database relation.
 - ▶ The SELECT clause contains only attribute names of the relation, and does not have any expressions, aggregates, or distinct specification.
 - ► Any attribute not listed in the SELECT clause can be set to null.
 - ▶ The query does not have a GROUP BY or HAVING clause.

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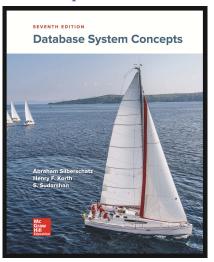
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