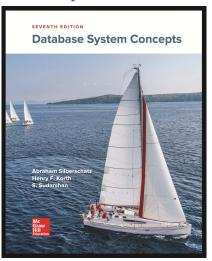
Database System Concepts, 7th Edition Chapter 3: Introduction to SQL

Silberschatz, Korth and Sudarshan

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Database System Concepts



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Plan

Overview of The SQL Query Language

SQL Data Definition

Basic Query Structure of SQL Queries

Additional Basic Operations

Set Operations

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Aggregate Functions

Nested Subqueries

Modification of the Database

History

- ▶ IBM Sequel language developed as part of System R project at IBM San Jose Research Laboratory.
- ► Renamed Structured Query Language (SQL).
- ► ANSI and ISO standard SQL:
 - ► SQL-86
 - ► SQL-89
 - ► SQL-92
 - ► SQL:1999¹
 - ► SQL:2003
- Commercial systems offer most, if not all, SQL-92 features, plus varying feature sets from later standards and special proprietary features.
 - ▶ Not all examples here may work on your particular system.

¹language name became Y2K compliant!.

SQL Parts

- ▶ DML provides the ability to query information from the database and to insert tuples into, delete tuples from, and modify tuples in the database.
- ► Integrity the DDL includes commands for specifying integrity constraints.
- ➤ View definition the DDL includes commands for defining views.

SQL Parts

- ► Transaction control includes commands for specifying the beginning and ending of transactions.
- ► Embedded SQL and dynamic SQL define how SQL statements can be embedded within general-purpose programing languages.
- ► Authorization includes commands for specifying access rights to relations and views.

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Data Definition Language

The SQL data-definition language (DDL) allows the specification of information about relations, including:

- ▶ The schema for each relation.
- ▶ The type of values associated with each attribute.
- ► The integrity constraints.
- ▶ The set of indices to be maintained for each relation.
- ▶ Security and authorization information for each relation.
- ▶ The physical storage structure of each relation on disk.

Domain Types in SQL

- ► char(n). Fixed length character string, with user-specified length n.
- ▶ varchar(n). Variable length character string, with user-specified maximum length n.
- ▶ int. Integer (a finite subset of integers that is machine-dependent).
- ▶ smallint. Small integer (a machine-dependent subset of integer domain type).

Domain Types in SQL

- ▶ numeric(p, d). Fixed point number, with user-specified precision of p digits, with d digits to the right of decimal point. (ex., numeric(3,1), allows 44.5 to be stores exactly, but not 444.5 or 0.32).
- ▶ real, double precision. Floating point and double-precision floating point numbers, with machine-dependent precision.
- ▶ float(n). Floating point number, with user-specified precision of at least n digits.
- ▶ More are covered in Chapter 4.

Create Table Construct

► An SQL relation is defiend using the create table command:

```
CREATE TABLE r (
A_1D_1, A_2D_2, \ldots, A_nD_n,
(integrity\_constraints_1),
(integrity\_constraints_2),
\vdots
(integrity\_constraints_k)
);
```

- ightharpoonup r is the name of the relation.
- \triangleright each A_i is an attribute name in the schema of relation r.
- $ightharpoonup D_i$ is the data type of values in the domain of attribute A_i .

Create Table Construct

```
CREATE TABLE instructor (

ID char(5),

name varchar(20),

dept_name varchar(20),

salary numeric(8,2)

);
```

Integrity Constraints in Create Table

- ► Types of integrity constraints:
 - ightharpoonup primary $\ker(A_{j_1},\ldots,A_{j_m})$
 - ightharpoonup foreing $\ker(A_{k_1},\ldots,A_{k_n})$
 - ▶ not null
- ▶ SQL prevents any update to the database that violates an integrity constraint.

Integrity Constraints in Create Table

```
CREATE TABLE instructor (

D char(5),

name varchar(20),

dept_name varchar(20),

salary numeric(8,2),

primary key (ID),

foreign key (dept_name) references department

);
```

Few more Relation Definitions

```
CREATE TABLE student (

ID varchar(5),

name varchar(20) not null,

dept_name varchar(20),

tot_cred numeric(3,0),

primary key (ID),

foreign key (dept_name) references department

);
```

Few more Relation Definitions

```
CREATE TABLE takes (
            TD
                       varchar(5),
            course_id varchar(8),
3
            sec_id varchar(8),
            semester varchar(6),
            year numeric(4,0),
6
            grade varchar(2),
            primary key (ID, course_id, sec_id, semester, year),
            foreign key (ID) references student,
            foreign key (course_id, sec_id, semester, year)
10
                references section
11
        );
12
```

And more still

```
CREATE TABLE course (

course_id varchar(8),

title varchar(50),

dept_name varchar(20),

credits numeric(2,0),

primary key (course_id),

foreign key (dept_name) references department

);
```

Updates to Tables

► Insert:

```
INSERT INTO instructor
   VALUES ('10211', 'Smith', 'Biology', 66000);
```

▶ Delete:

```
DELETE FROM studentes;
```

▶ Drop a table:

```
DROP TABLE r;
```

Updates to Tables

- ► Alter:
 - \triangleright alter table r add A D:
 - where A is the name of the attribute to be added to relation r and D is the domain of A.
 - ▶ all existing tuples in the relation are assigned *null* as the value for the new attribute.
 - ightharpoonup alter table r drop A:
 - \triangleright where A is the name of an attribute of relation r.
 - dropping of attributes not supported by many databases.

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Basic Query Structure

► A typical SQL query has the form:

SELECT

$$A_1, A_2, \ldots, A_n$$

FROM

$$r_1, r_2, \ldots, r_m$$

WHERE

F

- \triangleright A_i represents an attribute.
- $ightharpoonup r_i$ represents a relation.
- \triangleright P is a predicate.
- ▶ The result of an SQL is a relation.

The SELECT Clause

- ► The SELECT clause lists the attributes desired in the result of a query.
 - ▶ It corresponds to the projection (Π) operation of the relational algebra.

Example:

"Find the names of all instructors."

SELECT name FROM instructor;

The SELECT Clause

Note

SQL names are case sensitive (i.e., you may use upper- or lower-case letters.)

- ightharpoonup E.g., Name \equiv NAME \equiv name.
- ▶ Some people use upper case wherever we use bold font.

- ➤ SQL allows duplicates in relations as well as in query results.
- ➤ To force the elimination of duplicates, insert the keyword DISTINCT after the SELECT clause.

Example:

"Find the department names of all instructors, and remove duplicates."

SELECT DISTINCT dept_name FROM instructor;

► The keyword ALL specifies that duplicates should not be remove.

SELECT ALL dept_name FROM instructor;

► An asterisk (*) in the SELECT clause denotes "all attributes".

```
SELECT * FROM instructor;
```

► An attribute can be a literal with FROM clause.

```
SELECT 'A' FROM instructor;
```

Result is a table with one column and N rows (number of tuples in the *instructor* table), each row with value "A".

▶ An attribute can be a literal with no FROM clause.

```
SELECT '42';
```

- ▶ Results is a table with one column and a single row with value "42".
- ► Can give the column a name using:

```
SELECT '42' AS foo;
```

- ▶ The SELECT clause can contain arithmetic expression involving the operation, +, −, *, and /, and operating on constants or attributes of tuples.
 - ► The query:

```
SELECT ID, name, salary/12 FROM instructor;
```

would return a relation that is the same as the *instructor* relation, except that the value of the attribute salary is divided by 12.

► Can rename ''salary/12'' using the AS clause

```
SELECT ID, name, salary/12 AS monthly_salary FROM instructor;
```

The WHERE Clause

- ► The WHERE clause specifies conditions that the result must satisfy
 - \blacktriangleright Corresponds to the selection (σ) predicate of the relational algebra.

Example:

"Find all instructors in Comp. Sci. department."

```
SELECT name
FROM instructor
WHERE dept_name = 'Comp. Sci.';
```

The WHERE Clause

- ➤ SQL allows the use of the logical connectives **and**, **or**, and **not**.
- ► The operands of the logical connectives can be expressions involving the comparison operators <, <=, >, >=, =, and <>>.
- ► Comparisons can be applied to results of arithmetic expressions.

Example:

"Find all instructors in Comp. Sci. with salary > 80000."

```
SELECT name

FROM instructor

WHERE dept_name = 'Comp. Sci.' AND salary > 80000;
```

The FROM Clause

- ► The FROM clause lists the relations involved in the query.
 - ightharpoonup Corresponds to the Cartesian product (\times) operation of the relational algebra.

Example:

"Find the Cartesian producto $instructor \times teaches$."

```
SELECT *
FROM instructor, teaches;
```

The FROM Clause

- generates every possible instructor teaches pair, with all attributes from both relations.
- ▶ for common attributes (e.g., ID), the attributes in the resulting table are renamed using the relation name (e.g., instructor.ID)
- ▶ Cartesian product is not very useful directly, but it is useful combined with the WHERE clause condition (σ operation in relational algebra).

"Find the names of all instructors who have taught some course and the course id."

"Find the names of all instructors who have taught some course and the course id."

```
1 SELECT
2 name, course_id
3 FROM
4 instructor, teaches
5 WHERE
6 instructor.ID = teaches.ID;
```

"Find the names of all instructors in the Art department who have taught some course and the course_id."

"Find the names of all instructors in the Art department who have taught some course and the course_id."

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The Rename Operation

► The SQL allows renaming relations and attributes using the AS clause:

old_name AS new_name

"Find the names of all instructors who have a higher salary than some instructor in Comp. Sci. department."

```
SELECT DISTINCT
T.name
FROM
instructor AS T, instructor AS S
WHERE
T.salary > S.salary AND
S.dept_name = 'Comp. Sci.';
```

The Rename Operation

keyword AS is optional and may be omitted.

 $instructor AS \equiv instructor T$

Self Join Example

ightharpoonup Relation employee-supervisor:

person	supervisor
Bob	Alice
Mary	Susan
Alice	David
David	Mary

- **Exercises:**
 - ► "Find the supervisor of Bob"
 - ► "Find the supervisor of the supervisor of Bob"
 - "Could you find ALL the supervisor (direct and indirect) of Bob?"

String Operations

- ➤ SQL includes a string-matching operator for comparisons on character strings. The operator LIKE uses patterns that are described using two special characters:
 - ▶ percent (%) The % character matches any substring.
 - ▶ underscore(_) The _ character any character.

Example:

"Find the names of all instructors whose name includes the substring 'dar'."

String Operations

```
SELECT
name
FROM
instructor
WHERE
name LIKE '%dar%';
```

► Match the string '100%'

```
LIKE '100\%';
```

in that above we use backslash $(\)$ as the escape character.

String Operations (Cont.)

- ▶ Patterns are case sensitive.
- ▶ Pattern matching examples:
 - ► 'Intro%' matches any string beginning with "Intro".
 - ▶ '%Comp%' matches any string containing "Comp" as a substring.
 - '___' matches any string of exactly three characters.
 - '___%' matches any string of at least three characters.
- ▶ SQL supports a variety of string operations such as:
 - concatenation (using "||").
 - converting from upper to lower case (and viceversa).
 - finding string length, extracting substrings, etc.

Ordering the Display of Tuples

Example:

"List in alphabetic order the names of all instructors."

```
1 SELECT DISTINCT
2 name
3 FROM
4 instructor
5 ORDER BY
6 name;
```

Ordering the Display of Tuples

▶ We may specify DESC for descending order of ASC for ascending order, for each attribute; ASC is the default.

ORDER BY name DESC

► Can sort in multiple attributes.

ORDER BY dept_name, name

Where Clause Predicates

► SQL includes a BETWEEN comparison operator.

Example:

"Find the names of all instructors with salary between \$90000 and \$100000 (that is, \geq \$90000 and \leq \$100000).

```
1 SELECT
2 name
3 FROM
4 instructor
5 WHERE
6 salary BETWEEN 90000 AND 100000;
```

Where Clause Predicates (Cont.)

► Tuple comparison:

```
1     SELECT
2     name, course_id
3     FROM
4     instructor, teaches
5     WHERE
6     (instructor.ID, dept_name) = (teaches.ID = 'Biology');
```

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Set Operations

"Find courses that ran in Fall 2017 or in Spring 2018."

```
(SELECT course_id FROM section WHERE sem = 'Fall' AND year = 2017)
UNION
(SELECT course_id FROM section WHERE sem = 'Spring' AND year = 2018)
```

"Find courses that ran in Fall 2017 and in Spring 2018."

```
(SELECT course_id FROM section WHERE sem = 'Fall' AND year = 2017)
INTERSECT
(SELECT course_id FROM section WHERE sem = 'Spring' AND year = 2018)
```

Set Operations (Cont.)

"Find courses that ran in Fall 2017 but not in Spring 2018."

```
1 (SELECT course_id FROM section WHERE sem = 'Fall' AND year = 2017)
2 EXCEPT
3 (SELECT course_id FROM section WHERE sem = 'Spring' AND year = 2018)
```

- ► Set operations UNION, INTERSECT, and EXCEPT automatically eliminate duplicates.
- ► To retain all duplicate use UNION ALL, INTERSECT ALL, and EXCEPT ALL.

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Null Values

- ► It is possible for tuples to have a null value, denoted by null, for some of their attributes.
- null signifies an unknown value or that a value does not exist.
- ► The results of any arithmetic expression involving null is null.

Example

5 + null returns null.

Null Values (Cont.)

▶ The predicate IS NULL can be used to check for null values.

Example

"Find all instructors whose salary is null."

SELECT name FROM instructor WHERE salary IS NULL

▶ The predicate IS NOT NULL succeeds if the value on which it is applied is not null.

Null Values (Cont.)

- ➤ SQL treats as **unknown** the result of any comparison involving a null value (other than predicates IS NULL and IS NOT NULL).
 - Example: 5 < null or null <> null or null = null

Null Values (Cont.)

▶ The predicate in a WHERE clause can involve Boolean operations (AND, OR, NOT); thus the definitions of the Boolean operations need to be extended to deal with the value unknown.

AND:

```
 \begin{array}{l} (\texttt{true AND } \ \mathbf{unknown}) = \mathbf{unknown} \\ (\texttt{false AND } \ \mathbf{unknown}) = \mathtt{false} \\ (\mathbf{unknown } \ \mathtt{AND } \ \mathbf{unknown}) = \mathbf{unknown} \end{array}
```

DR:

```
(\mathbf{unknown} \ \mathtt{OR} \ \mathsf{true}) = \mathsf{true}
(\mathbf{unknown} \ \mathtt{OR} \ \mathsf{false}) = \mathbf{unknown}
(\mathbf{unknown} \ \mathtt{OR} \ \mathbf{unknown}) = \mathbf{unknown}
```

► Result of WHERE clause predicate is treated as false if it evaluates to unknown.

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Aggregate Functions

- ► These functions operate on the multiset of values of a column of a relation, and return a value:
 - ► AVG: average value.
 - ► MIN: minimum value.
 - ► MAX: maximum value.
 - SUM: sum of values.
 - COUNT: number of values.

Aggregate Functions Examples

"Find the average salary of instructors in the Computer Science department."

```
SELECT
AVG(salary)
FROM
instructor
WHERE
dept_name = 'Comp. Sci.';
```

Aggregate Functions Examples

"Find the total number of instructors who teach a course in the Spring 2018 semester."

```
SELECT
COUNT(DISTINCT ID)
FROM
teaches
WHERE
semester = 'Spring' AND year = 2018
```

Aggregate Functions Examples

"Find the number of tuples in the course relation."

```
SELECT
COUNT(*)
FROM
course;
```

Aggregate Functions – Group By

"Find the average salary of instructors in each department."

```
SELECT dept_name, AVG(salary) AS avg_salary
FROM instructor
GROUP BY dept_name;
```

ID	name	dept_name	salary
76766	Crick	Biology	72000
45565	Katz	Comp. Sci.	75000
10101	Srinivasan	Comp. Sci.	65000
83821	Brandt	Comp. Sci.	92000
98345	Kim	Elec. Eng.	80000
12121	Wu	Finance	90000
76543	Singh	Finance	80000
32343	El Said	History	60000
58583	Califieri	History	62000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
22222	Einstein	Physics	95000

dept_name	avg_salary
Biology	72000
Comp. Sci.	77333
Elec. Eng.	80000
Finance	85000
History	61000
Music	40000
Physics	91000

Aggregation (Cont.)

► Attributes in SELECT clause outside of aggregate functions must appear in GROUP BY list.

```
/* erroneous query */
SELECT
dept_name, ID, AVG(salary)
FROM
instructor
GROUP BY
dept_name;
```

Aggregate Functions – Having Clause

"Find the names and average salaries of all departments whose average salary is greater than 42000".

```
SELECT dept_name, AVG(salary) as avg_salary
FROM instructor
GROUP BY dept_name
HAVING AVG(salary) > 42000;
```

Note:

Predicates in the HAVING clause are applied after the formation of groups whereas predicates in the WHERE clause are applied before forming groups.

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Nested Subqueries

- ► SQL provides a mechanism for the nesting of subqueries. A **subquery** is SELECT-FROM-WHERE expression that is nested within another query.
- ▶ The nesting can be done in the following SQL query:

SELECT

$$A_1, A_2, \ldots, A_n$$

FROM

$$r_1, r_2, \ldots, r_m$$

WHERE

P

as follows:

Nested Subqueries

SELECT

$$A_1, A_2, \ldots, A_n$$

FROM

$$r_1, r_2, \ldots, r_m$$

WHERE

P

as follows:

- \triangleright From clause: r_i can be replaced by any valid subquery.
- Where clause: P can be replaced with an expression of the form:

B is an attribute and < operation > will be defined later.

 \triangleright Select clause: A_i can be replaced be a subquery that generates a single value.

Set Membership

Find courses offered in Fall 2017 and in Spring 2018.

```
SELECT DISTINCT
 1
              course_id
         FROM
3
             section
         WHERE
              semester = 'Fall' AND year = 2017 AND
              course_id IN (SELECT
                               course_id
                          FROM
9
                               section
10
                          WHERE
11
                               semester = 'Spring' AND year = 2018
12
                          );
13
```

Set Membership (Cont.)

Find courses offered in Fall 2017 but not in Spring 2018.

```
SELECT DISTINCT
 1
             course_id
         FROM
3
             section
         WHERE
             semester = 'Fall' AND year = 2017 AND
             course_id NOT IN (SELECT
                               course_id
                          FROM
9
                               section
10
                          WHERE
11
                               semester = 'Spring' AND year = 2018
12
                          );
13
```

Set Membership (Cont.)

Name all instructors whose name is neither "Mozart" nor "Einstein".

```
SELECT DISTINCT
name
FROM
instructor
WHERE
name NOT IN ('Mozart', 'Einstein');
```

Set Membership (Cont.)

Find the total number of (distinct) students who have taken course sections taught by the instructor with ID 10101.

```
SELECT
1
              COUNT(DISTINCT ID)
         FROM
              takes
          WHERE
              (course_id, sec_id, semester, year) IN (
                  SELECT
                       course_id, sec_id, semester, year
                  FROM
9
10
                       teaches
11
                  WHERE
                       teaches.TD = 10101
12
13
              );
```

Note

Above query can be written in a much simpler manner. The formulation above is simply to illustrate SQL features.

Set Comparison – "some" Clause

Find names of instructors with salary greater than that of some (at least one) instructor in the Biology department.

```
SELECT DISTINCT
T.name
FROM
instructor AS T, instructor AS S
WHERE
T.salary > S.salary AND
S.dept_name = 'Biology';
```

Set Comparison – "some" Clause (Cont.)

Same query using > SOME clause.

```
SELECT DISTINCT
              name
         FR.OM
3
              instructor
         WHERE
              salary > SOME (
                  SELECT
                       salary
                  FROM
9
10
                       instructor
                  WHERE
11
                       dept_name = 'Biology'
12
              );
13
```

Definition of "some" Clause

 $F < comp > \mathtt{SOME} \ r \Leftrightarrow \exists t \in r \ \mathrm{such \ that} \ (F < comp > t) \ \mathrm{where} < comp > \mathrm{can \ be:} \ <, \leq, >, \geq, =, \neq.$

$$\left(5 < \text{SOME} \boxed{ 5 \atop 5 \atop 6} \right) = true.$$

read: 5 < some tuple in the relation.

$$\begin{pmatrix} 5 < \text{ SOME } \boxed{0 \\ 5} \end{pmatrix} = false.$$

$$\begin{pmatrix} 5 = \text{ SOME } \boxed{0 \\ 5} \end{pmatrix} = true.$$

$$\begin{pmatrix} 5 \neq \text{ SOME } \boxed{0 \\ 5} \end{pmatrix} = true.$$

Note (= SOME) \equiv IN. However, (\neq SOME) $\not\equiv$ NOT IN.

Set Comparison – "all" Clause

Find the names of all instructors whose salary is greater than the salary of all instructors in the Biology department.

```
SELECT
              name
          FROM
3
               instructor
          WHERE
5
               salary > ALL (
                   SELECT
                        salary
                   FR.OM
9
10
                        instructor
                   WHERE.
11
12
                        dept_name = 'Biology'
               );
13
```

Definition of "all" Clause

$$F < comp > \mathtt{ALL} \ r \Leftrightarrow \exists t \in r (F < comp > t).$$

$$\left(5 < \text{ALL} \begin{bmatrix} 0 \\ 5 \\ 6 \end{bmatrix} \right) = false.$$

$$\left(5 < \text{ALL} \left[\begin{array}{c} 6 \\ \hline 10 \end{array} \right] \right) = true.$$

$$\left(5 = \text{ALL} \left| \frac{4}{5} \right| \right) = false.$$

$$\left(5 \neq \text{ ALL} \left| \frac{4}{6} \right| \right) = true.$$

Note $(\neq ALL) \equiv NOT$ IN. However, $(= ALL) \not\equiv IN$.

Test for Empty Relations

- ► The EXISTS construct returns the value true if the argument subquery is nonempty.
- ► EXISTS

$$r \Leftrightarrow r \neq \emptyset$$

NOT EXISTS

$$r \Leftrightarrow r = \emptyset$$

Use of "exists" Clause

➤ Yet another way of specifying the query "Find all courses taught in both the Fall 2017 semester and in the Spring 2018 semester".

```
SELECT
 1
                  course id
2
3
              FR.OM
                  section AS S
              WHERE
5
                  semester = 'Fall' AND year = 2017 AND
6
                  EXISTS (
                      SELECT *
                      FROM section AS T
9
                       WHERE semester = 'Spring' AND
10
                           vear= 2018 AND
11
                           S.course_id = T.course_id
12
                  );
13
```

Use of "exists" Clause

```
SELECT
                   course_id
 3
               FROM
                   section AS S
               WHERE
                   semester = 'Fall' AND year = 2017 AND
                   EXISTS (
 8
                       SELECT *
                       FROM section AS T
10
                       WHERE semester = 'Spring' AND
11
                           year= 2018 AND
12
                           S.course_id = T.course_id
13
                   );
```

- ► Correlation name: variable S in the outer query.
- ► Correlated subquery: the inner query,

Use of "not exists" Clause

► Find all students who have taken all courses offered in the Biology department.

```
SELECT DISTINCT
 1
                  S.ID, S.name
2
              FR.OM
3
                  student AS S
 4
              WHERE
5
                  NOT EXISTS (
6
                       ( SELECT course_id
7
                           FROM course
                           WHERE dept_name = 'Biology' )
9
                       EXCEPT
10
11
                       ( SELECT T.course_id
12
                           FROM takes AS T
13
                           WHERE S.ID = T.ID)
14
                  );
```

Use of "not exists" Clause

```
1
               SELECT DISTINCT
 2
                   S.ID. S.name
3
               FROM
                   student AS S
 5
               WHERE
6
                   NOT EXISTS (
 7
                        ( SELECT course id
                            FROM course
9
                            WHERE dept_name = 'Biology' )
10
                        EXCEPT
11
                        ( SELECT T.course_id
12
                            FROM takes AS T
13
                            WHERE S.ID = T.ID )
14
                   );
```

- First nested query lists all courses offered in Biology.
- Second nested query lists all courses a particular student took.
- ▶ Note: That $X Y = \emptyset \Leftrightarrow X \subseteq Y$.
- ▶ Note: Cannot write this query using = ALL and its variants.

Test for Absence of Duplicate Tuples

- ► The UNIQUE construct tests whether a subquery has any duplicate tuples in its result.
- ► The UNIQUE construct evaluates to 'true' if a given subquery contains no duplicates.

Test for Absence of Duplicate Tuples (Cont.)

▶ Find all courses that were offered at most once in 2017.

```
SELECT
 1
2
                   T.course id
              FROM
3
                   course AS T
4
              WHERE
5
                   UNIQUE (
6
                       SELECT
7
                            R.course id
8
                       FR.OM
9
                            section AS R
10
                       WHERE
11
                            T.course_id = R.course_id AND
12
                            R.year = 2017
13
                   );
14
```

Subqueries in the From Clause

► SQL allows a subquery expression to be used in the FROM clause.

Example:

Find the average instructors' salaries of those departments where the average salary is greater than \$42,000.

```
select
dept_name, avg_salary
FROM (Select dept_name, avg (salary) as avg_salary
FROM instructor
GROUP BY dept_name )
WHERE
avg_salary > 42000;
```

Subqueries in the From Clause (Cont.)

- Nother that we do not need to use the HAVING clause.
- ► Another way to write the previous query:

```
dept_name, avg_salary

from (SELECT dept_name, avg (salary)

From instructor

GROUP BY dept_name

AS dept_avg (dept_name, avg_salary)

WHERE

avg_salary > 42000;
```

WITH Clause

▶ The WITH clause provides a way of defining a temporary relation whose definition is available only to the query in which the WITH clause occurs.

Example:

Find all departments with the maximum budget.

```
WITH max_budget (value) AS

( SELECT MAX(budget)

FROM department )

SELECT

department.name

FROM

department, max_budget

WHERE

department.budget = max_budget.value;
```

Complex Queries using With Clause

Find all departments where the total salary is greater than the average of the total salary at all departments.

```
WITH dept_total(dept_name, value) AS
             ( SELECT dept_name, sum(salary)
               FROM instructor
               GROUP BY dept_name ),
             dept_total_avg(value) AS
5
             ( SELECT AVG(value)
                FROM dept_total )
             SELECT
9
                  dept_name
10
             FR.OM
                  dept_total, dept_total_avg
11
12
             WHERE
                  dept_total.value > dept_total_avg.value;
13
```

Scalar Subquery

Scalar subquery is one which is used where a single value is expected.

Example:

List all departments along with the number of instructors in each department.

```
SELECT

dept_name,

(SELECT count(*)

FROM instructor

WHERE department.dept_name = instructor.dept_name)

AS num_instructors

FROM

department;
```

► Runtime error if subquery returns more than one result tuple.

Plan

Overview of The SQL Query Language

SQL Data Definition

Basic Query Structure of SQL Queries

Additional Basic Operations

Set Operations

Null Values

Aggregate Functions

Nested Subqueries

Modification of the Database

Modification of the Database

- ▶ Deletion of tuples from a given relation.
- ▶ Insertion of new tuples into a given relation.
- ▶ Updating of values in some tuples in a given relation.

Deletion

Delete all instructors.

DELETE FROM instructor;

Deletion

Delete all instructors.

```
DELETE FROM instructor;
```

Delete all instructors from the Finance department.

```
DELETE FROM instructor
WHERE dept_name= 'Finance';
```

Deletion

Delete all instructors.

```
DELETE FROM instructor;
```

Delete all instructors from the Finance department.

```
DELETE FROM instructor
WHERE dept_name= 'Finance';
```

Delete all tuples in the instructor relation for those instructors associated with a department located in the Watson building.

```
DELETE FROM instructor

WHERE dept_name IN ( SELECT dept_name

FROM department

WHERE building = 'Watson' );
```

Deletion (Cont.)

▶ Delete all instructors whose salary is less than the average salary of instructors.

```
DELETE FROM instructor
WHERE salary < ( SELECT AVG(salary)
FROM instructor );
```

- ▶ Problem: as we delete tuples from instructor, the average salary changes
- ► Solution used in SQL:
 - 1. First, compute AVG(salary) and find all tuples to delete.
 - 2. Next, delete all tuples found above (without recomputing average or retesting the tuples).

Insertion

Add a new tuple to course.

```
INSERT INTO course
VALUES('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

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```
INSERT INTO course
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```

or equivalently.

```
INSERT INTO course (course_id, title, dept_name, credits)
VALUES('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

Insertion

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```
INSERT INTO course
VALUES('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

or equivalently.

```
INSERT INTO course (course_id, title, dept_name, credits)
VALUES('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

Add a new tuple to student with credits set to null.

```
INSERT INTO student
VALUES('3003', 'Green', 'Finance', null);
```

Insertion (Cont.)

Make each student in the Music department who has earned more than 144 credit hours an instructor in the Music department with a salary of \$18,000.

```
INSERT INTO instructor

SELECT ID, name, dept_name, 18000

FROM student

WHERE dept_name = 'Music' AND total_cred > 144;
```

Insertion (Cont.)

The SELECT FROM WHERE statement is evaluated fully before any of its results are inserted into the relation. Otherwise queries like:

```
INSERT INTO table1 SELECT * FROM table1;
```

would cause problem.

Updates

Give a 5% salary raise to all instructors.

```
UPDATE instructor
SET salary = salary * 1.05;
```

Updates

Give a 5% salary raise to all instructors.

```
UPDATE instructor
SET salary = salary * 1.05;
```

Give a 5% salary raise to those instructors who earn less than 70000.

```
UPDATE instructor
SET salary = salary * 1.05
WHERE salary < 70000;
```

Updates

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```
UPDATE instructor
SET salary = salary * 1.05
WHERE salary < 70000;</pre>
```

Give a 5% salary raise to instructors whose salary is less than average.

```
UPDATE instructor

SET salary = salary * 1.05

WHERE salary < ( SELECT AVG(salary)

FROM instructor );
```

Updates (Cont.)

Increase salaries of instructors whose salary is over \$100,000 by 3%, and all others by a 5%.

▶ Write two update statements:

```
1     UPDATE instructor
2     SET salary = salary * 1.03
3     WHERE salary > 100000;
4     UPDATE instructor
5     SET salary = salary * 1.05
6     WHERE salary <= 100000;</pre>
```

- ► The order is important.
- ► Can be done better using the case statement (next slide)...

Case Statement for Conditional Updates

Same query as before but with case statement

```
UPDATE instructor

SET salary = CASE

WHEN salary <= 100000 THEN salary * 1.05

ELSE salary * 1.03

END
```

Updates with Scalar Subqueries

Recompute and update tot_creds value for all students.

```
1    UPDATE student S
2    SET tot_cred = ( SELECT SUM(credits)
3     FROM takes, course
4    WHERE takes.course_id = course.course_id AND
5     S.ID = takes.ID AND
6     takes.grade <> 'F' AND
7    takes.grade IS NOT NULL );
```

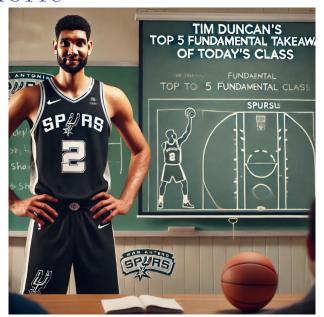
Updates with Scalar Subqueries (Cont.)

- ➤ Sets tot_creds to null for students who have not taken any course.
- ► Instead of SUM(credits), use:

```
CASE
WHEN SUM(credits) IS NOT NULL THEN SUM(credits)
ELSE 0
END
```

End of Chapter 3.

TDT5FTOTTC



5 Handling Nulls and Special Cases: SQL provides special handling for null values (unknown or missing data) through specific predicates (IS NULL, IS NOT NULL).

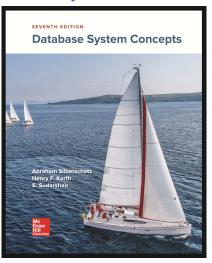
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- 1 **SQL Basics and Structure:** SQL is a standardized language used for querying and managing relational databases using commands like SELECT, FROM, and WHERE.

Database System Concepts



Content has been extracted from Database System Concepts, Seventh Edition, by Silberschatz, Korth and Sudarshan. Mc Graw Hill Education. 2019.

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