



Master theorem (reprise)

$$T(n) = a T(n/b) + f(n)$$

CASE 1: $f(n) = O(n^{\log_b a - \varepsilon})$, constant $\varepsilon > 0$
 $\Rightarrow T(n) = \Theta(n^{\log_b a})$.

CASE 2: $f(n) = \Theta(n^{\log_b a} \lg^k n)$, constant $k \geq 0$
 $\Rightarrow T(n) = \Theta(n^{\log_b a} \lg^{k+1} n)$.

CASE 3: $f(n) = \Omega(n^{\log_b a + \varepsilon})$, constant $\varepsilon > 0$,
and regularity condition
 $\Rightarrow T(n) = \Theta(f(n))$.

Merge sort: $a = 2, b = 2 \Rightarrow n^{\log_b a} = n^{\log_2 2} = n$
 \Rightarrow **CASE 2** ($k = 0$) $\Rightarrow T(n) = \Theta(n \lg n)$.