
Algorithm 1: Find candidate disks with plane sweeping technique

Input: $\mathcal{T}[t_i]$: positions in timestamp t_i , sorted by x-axis values, ϵ : flock diameter, μ : minimum size of flock

Output: \mathcal{C} : candidate disks for timestamp t_i , \mathcal{B} : active boxes in timestamp t_i

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1  $\mathcal{C} \leftarrow \emptyset, \mathcal{B} \leftarrow \emptyset$ 
2 foreach  $p_r \in \mathcal{T}[t_i]$  do // analyze elements in increasing x-values
3    $\mathcal{P} \leftarrow \emptyset$  // list of elements of current box defined by  $p_r$ 
4   foreach  $p_s \in \mathcal{T}[t_i] : |p_s.x - p_r.x| \leq \epsilon$  do // test only elements inside  $2\epsilon$  x-band
5     if  $|p_s.y - p_r.y| \leq \epsilon$  then // check if  $p_s$  is inside  $2\epsilon$  y-band
6        $\mathcal{P} \leftarrow \mathcal{P} \cup p_s$  // add element  $p_s$  to box
7   foreach  $p \in \mathcal{P} : p.x \geq p_r.x$  do (// elements inside right half of box
8     if  $\text{dist}(p_r, p) \leq \epsilon$  then // calculate pair distance
9       let  $\{c_1, c_2\}$  be disks defined by  $\{p_r, p\}$  and radius  $\epsilon/2$ 
10      foreach  $c \in \{c_1, c_2\}$  do
11        if  $|c \cap \mathcal{P}| \geq \mu$  then // check the number of entries in disk
12           $\mathcal{C} \leftarrow \mathcal{C} \cup c$  // add  $c$  to candidate disks
13         $\mathcal{B} \leftarrow \mathcal{B} \cup \text{box}(p_r)$  // add box to active boxes
14 return  $\mathcal{C}, \mathcal{B}$ 
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Algorithm 2: Filter out disks which are subsets

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1 Algorithm FilterCandidates( $\mathcal{B}$ )
   Input:  $\mathcal{B}$ : active boxes of timestamp  $t_i$ , sorted by x-axis values
   Output:  $\mathcal{C}$ : final set of disks for timestamp  $t_i$ 
2    $\mathcal{C} \leftarrow \emptyset$ 
3   for  $j \leftarrow 0$  to  $j \leq |\mathcal{B}|$  do
4       for  $k \leftarrow j + 1$  to  $k \leq |\mathcal{B}|$  do
5           if IntersectsWith( $\mathcal{B}[j]$ ,  $\mathcal{B}[k]$ ) then
6               foreach  $c \in \mathcal{B}[j].\text{disks}$  do
7                    $\mathcal{C} \leftarrow \text{InsertDisk}(\mathcal{C}, c)$ 
8               else // No intersection.
9                   break
10 Procedure InsertDisk( $\mathcal{C}, c$ )
   Input:  $\mathcal{C}$ : set of disks,  $c$ : new disk
11   foreach  $d \in \mathcal{C}$  do
12       if  $c.\text{sign} \wedge d.\text{sign} = c.\text{sign} \ \&\& \ \text{dist}(c, d) \leq \epsilon$  then //  $c$  can be a subset of  $d$ 
13           if  $d \cap c = c$  then // Remove chance of false-positive
14               return  $\mathcal{C}$  // No need to insert  $c$ 
15       else if  $c.\text{sign} \wedge d.\text{sign} = d.\text{sign}$  then //  $d$  can be a subset of  $c$ 
16           if  $c \cap d = d$  then // Remove chance of false-positive
17                $\mathcal{C} \leftarrow \mathcal{C} \setminus d$  // Remove  $d$ 
18   return  $\mathcal{C} \cup c$ 
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