
Algorithm 1: Find candidate disks with plane sweeping technique

Input: $\mathcal{T}[t_i]$: positions in timestamp t_i , sorted by x-axis values, ϵ : flock diameter, μ : minimum size of flock

Output: \mathcal{C} : candidate disks for timestamp t_i , \mathcal{B} : active boxes in timestamp t_i

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1  $\mathcal{C} \leftarrow \emptyset, \mathcal{B} \leftarrow \emptyset$ 
2 foreach  $p_r \in \mathcal{T}[t_i]$  do // analyze elements in increasing x-values
3    $\mathcal{P} \leftarrow \emptyset$  // list of elements of current box defined by  $p_r$ 
4   foreach  $p_s \in \mathcal{T}[t_i] : |p_s.x - p_r.x| \leq \epsilon$  do // test only elements inside  $2\epsilon$  x-band
5     if  $|p_s.y - p_r.y| \leq \epsilon$  then // check if  $p_s$  is inside  $2\epsilon$  y-band
6        $\mathcal{P} \leftarrow \mathcal{P} \cup p_s$  // add element  $p_s$  to box
7   foreach  $p \in \mathcal{P} : p.x \geq p_r.x$  do (// elements inside right half of box
8     if  $\text{dist}(p_r, p) \leq \epsilon$  then // calculate pair distance
9       let  $\{c_1, c_2\}$  be disks defined by  $\{p_r, p\}$  and radius  $\epsilon/2$ 
10      foreach  $c \in \{c_1, c_2\}$  do
11        if  $|c \cap \mathcal{P}| \geq \mu$  then // check the number of entries in disk
12           $\mathcal{C} \leftarrow \mathcal{C} \cup c$  // add  $c$  to candidate disks
13           $\mathcal{B} \leftarrow \mathcal{B} \cup \text{box}(p_r)$  // add box to active boxes
14 return  $\mathcal{C}, \mathcal{B}$ 
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