Monads: why should you care?

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Outline

- 1 abstract, me, what you will learn
- 2 non-monadic evaluator
- equivalent monadic evaluator
- add unbound variable error handling
- **5** handle dynamic type errors
- 6 hide the environment
- add state (i.e., simulated mutability)
- add logging
- add IO
- final version
- resources

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Monads: why should you care?

Most monad explanations

- use analogies like "container" or "context"
- show how monads are implemented
- how they relate to category theory

This presentation shows

- the usefulness of monads—what they can do for you
- (not how they work)

Goal

- provide a glimpse of benefits of monads
- so you can use them in your work
- motivate you to understand how they work on your own

Source of examples: Martin Grabmuller :

• http://www.grabmueller.de/martin/www/pub/Transformers.en.html

Harold Carr, software architect/developer

- WebLogic InfiniBand Communication at Oracle
- SOAP Web Services Technology at Oracle/Sun
- CORBA/IIOP, RMI-IIOP Technology at Sun
- Visual LISP Technology at Autodesk
- Distributed C++ at University of Utah/HP
- (Concurrent) Utah Scheme, Utah Common Lisp, Portable Standard Lisp at University of Utah/HP
- Logic Simulation at Patel Systems/Cirrus Logic

what you will learn

initial intuition

- a Monad is used to pipe the output of one function into another
- a monadic "pipe" has code that executes "behind the scenes"
 - rather than embedded in main lines of program

monads and side-effects

- a Monad has nothing to do with "real" side-effects
 - e.g., reading/writing a file
- a monadic type is often used to simulate side-effects in a purely functional way
 - aka "effectful"
- the IO monad does "real" side-effects

combining monads

• use monad "transformers" to combine two or more monads

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expression evaluator: types

```
type Name = String
                               -- variable names
3
   data Exp = Lit Integer -- expressions
4
               | Var Name
5
               | Plus Exp Exp
6
               | Abs Name Exp
               | App Exp Exp
8
              deriving (Eq, Show)
9
10
   data Value = IntVal Integer -- values
11
              | FunVal Env Name Exp
12
              deriving (Eq. Show)
13
14
   type Env = Map.Map Name Value -- var names to vals
```

non-monadic evaluator : types/code

```
eval0
                           :: Env -> Exp -> Value
3
    evalO env (Lit i) = IntVal i
5
   evalO env (Var n) = fromJust (Map.lookup n env)
8
    eval0 env (Plus e1 e2) = let IntVal i1 = eval0 env e1
9
                                 IntVal i2 = eval0 env e2
10
                             in IntVal (i1 + i2)
11
12
    evalO env (Abs n e) = FunVal env n e
13
14
    eval0 env (App e1 e2) = let v1 = eval0 env e1
15
                                 v2 = eval0 env e2
16
                             in case v1 of
17
                                  FunVal env' n body ->
18
                                    eval0 (Map.insert n v2 env') body
```

non-monadic evaluator: in action

```
1 - 12 + (\x -> x) (4 + 2)
2 exampleExp
   => Plus (Lit 12)
            (App (Abs "x" (Var "x"))
4
5
            (Plus (Lit 4) (Lit 2)))
6
   eval0 Map.empty exampleExp
8
   => IntVal 18
9
10
    evalO Map.empty (Plus (Lit 2) (Abs "x" (Lit 1)))
11
    => IntVal *** Exception: m.hs:59:31-55: Irrefutable pattern failed for
12
13
   eval0 Map.empty (Var "x")
    => *** Exception: Maybe.fromJust: Nothing
14
```

"fixing" unbound variable handling using Either

```
eval0e
                           :: Env -> Exp -> Either String Value
2
3
   evalOe env (Lit i) = Right $ IntVal i
4
5
   evalOe env (Var n) = case Map.lookup n env of
6
                               Nothing -> Left $ "unbound: " ++ n
                                Just v -> Right v
8
   evalOe env (Plus e1 e2) = let Right (IntVal i1) = evalOe env e1
10
                                  Right (IntVal i2) = eval0e env e2
11
                              in Right $ IntVal (i1 + i2)
12
13
   evalOe env (Abs n e) = Right $ FunVal env n e
14
15
   eval0e env (App e1 e2) = let Right v1 = eval0e env e1
16
                                  Right v2 = eval0e env e2
17
                              in case v1 of
18
                                  FunVal env' n body ->
19
                                      eval0e (Map.insert n v2 env')
20
                                             body
```

unbound Either: in action

```
1 (eval0e Map.empty (Var "x"))
2 => (Left "unbound: x")
```

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monadic evaluator: types

```
1 type Eval1 alpha = Identity alpha
2
3 runEval1 :: Eval1 alpha -> alpha
4 runEval1 ev = runIdentity ev
5
6 eval1 :: Env -> Exp -> Eval1 Value
```

```
1 eval1 env (Lit i) = return $ IntVal i
2
   eval1 env (Var n) = return $ fromJust (Map.lookup n env)
4
   eval1 env (Plus e1 e2) = do IntVal i1 <- eval1 env e1
6
                               IntVal i2 <- eval1 env e2
                               return $ IntVal (i1 + i2)
8
   eval1 env (Abs n e) = return $ FunVal env n e
10
11
   eval1 env (App e1 e2) = do v1 <- eval1 env e1
12
                               v2 <- eval1 env e2
13
                               case v1 of
14
                                 FunVal env' n body ->
15
                                  eval1 (Map.insert n v2 env') body
```

non/monadic evaluator : code side-by-side

```
return $ IntVal i
1 Lit: IntVal i
   Var : fromJust (Map.lookup n env)
                                         return $ fromJust (Map.lookup n
4
   Plus: let IntVal i1 = eval0 env e1
                                         do IntVal i1 <- eval1 env e1
6
             IntVal i2 = eval0 env e2
                                            IntVal i2 <- eval1 env e2
7
         in IntVal (i1 + i2)
                                            return $ IntVal (i1 + i2)
8
   Abs: FunVal env n e
                                         return $ FunVal env n e
10
11
   App : let v1 = eval0 env e1
                                       do v1 <- eval1 env e1
12
             v2 = eval0 env e2
                                            v2 <- eval1 env e2
13
         in case v1 of
                                           case v1 of
14
             FunVal env' n body -> FunVal env' n body ->
15
              eval0 (Map.insert n v2 env') eval1 (Map.insert n v2 env'
16
                    body
                                                    body
```

monadic "bind" evaluator

```
eval1 env (Lit i) = return $ IntVal i
3
   eval1 env (Var n) = return $ fromJust (Map.lookup n env)
4
5
    eval1 env (Plus e1 e2) = eval1 env e1 >>= \dummy ->
6
                             case dummy of
                               IntVal i1 -> eval1 env e2 >>= \dummy ->
8
                                 case dummy of
9
                                   IntVal i2 -> return $ IntVal (i1 + i2)
10
                                             -> fail "pattern match failure
                               _ -> fail "pattern match failure"
11
12
13
    eval1 env (Abs n e) = return $ FunVal env n e
14
15
    eval1 env (App e1 e2) = eval1 env e1 >>= \v1 ->
16
                             eval1 env e2 >>= \v2 ->
17
                             case v1 of
18
                               FunVal env' n body ->
19
                                 eval1 (Map.insert n v2 env') body
```

monadic evaluator : in action

```
1 runEval1 (eval1 Map.empty exampleExp)
2 => IntVal 18
3
4 runEval1 (eval1 Map.empty (Var "x"))
5 => *** Exception: Maybe.fromJust: Nothing
```

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unbound variable error handling: types/code

```
type Eval2 alpha = ErrorT String Identity alpha
2
   runEval2
                   :: Eval2 alpha -> Either String alpha
   runEval2 ev
                    = runIdentity (runErrorT ev)
5
6
   eval2a
                   :: Env -> Exp -> Eval2 Value
8
9
   eval2a env (Var n)
                         = case (Map.lookup n env) of
10
                            Nothing -> fail $ "unbound: " ++ n
11
                            Just v -> return v
```

unbound variable error handling: in action

runEval2 (eval2a Map.empty exampleExp)

```
2 => Right (IntVal 18)
3
4 runEval2 (eval2a Map.empty (Var "no-way"))
5 => Left "unbound: no-way"
6
7 -- type error, but not apparent in error message
8 runEval2 (eval2a Map.empty (Plus (Lit 12) (Abs "x" (Var "x"))))
9 => Left "Pattern match failure in do expression at transformers.hs:138:3
```

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handle dynamic type errors: code

no change in types

```
eval2b env (Plus e1 e2) = do e1' <- eval2b env e1
                                  e2' <- eval2b env e2
3
                                  case (e1', e2') of
4
                                   (IntVal i1, IntVal i2)
5
                                     -> return $ IntVal (i1 + i2)
6
                                   _ -> throwError "dyn type err: Plus"
7
8
    eval2b env (App e1 e2) = do v1 <- eval2b env e1
9
                                  v2 <- eval2b env e2
10
                                  case v1 of
11
                                   FunVal env' n body
12
                                     -> eval2b (Map.insert n v2 env') body
13
                                   _ -> throwError "dyn type err: App"
```

handle dynamic type errors: in action

```
1 runEval2 (eval2b Map.empty (Plus (Lit 12) (Abs "x" (Var "x"))))
2 => Left "dyn type err: Plus"
```

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hide the environment: types

```
Env only
```

- extended in App
- used in Var and Abs

```
1 type Eval3 alpha = ReaderT Env (ErrorT String Identity) alpha
2
3 runEval3 :: Env -> Eval3 alpha -> Either String alpha
```

3 runEval3 :: Env -> Eval3 alpha -> Either String alpha 4 runEval3 env ev = runIdentity (runErrorT (runReaderT ev env))

5

6 eval3 :: Exp -> Eval3 Value

hide the environment: code

```
eval3 (Var n) = do env <- ask
2
                            case Map.lookup n env of
3
                             Nothing -> throwError ("unbound: " ++ n)
4
                             Just val -> return val
5
6
    eval3 (Abs n e) = do env \leftarrow ask
                           return $ FunVal env n e
8
    eval3 (App e1 e2) = do v1 <- eval3 e1
10
                            v2 <- eval3 e2
11
                            case v1 of
12
                            FunVal env' n body
13
                               -> local (const (Map.insert n v2 env'))
                                        (eval3 body)
14
15
                             _ -> throwError "dyn type err: App"
```

hide the environment : in action

```
1 runEval3 Map.empty (eval3 exampleExp)
```

2 => Right (IntVal 18)

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```
type Eval4 alpha =
2
        ReaderT Env (ErrorT String (StateT Integer Identity)) alpha
3
   runEval4 :: Env
5
                -> Integer
6
                -> Eval4 alpha
7
                -> (Either String alpha, Integer)
   runEval4 env st. ev =
9
             runIdentity (runStateT (runErrorT (runReaderT ev env)) st)
10
11
   eval4 :: Exp -> Eval4 Value
```

add profiling to interpreter: in action

```
1 runEval4 Map.empty 0 (eval4 exampleExp)
```

2 => (Right (IntVal 18),8) -- 8 reduction steps

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add logging: types

```
type Eval5 alpha =
2
         ReaderT Env
3
                 (ErrorT String (WriterT [String]
4
                                           (StateT Integer Identity)))
5
                 alpha
6
7
   runEval5 :: Env
8
                -> Integer
9
                -> Eval5 alpha
10
                -> ((Either String alpha, [String]), Integer)
11
   runEval5 env st ev =
12
             runIdentity (runStateT (runWriterT (runErrorT
13
                                                      (runReaderT ev env)))
14
                                     st)
15
16
   eval5
             :: Exp -> Eval5 Value
```

add logging: code

add logging: in action

```
1 runEval5 Map.empty 0 (eval5 exampleExp)
2 => ((Right (IntVal 18),["x"]),8)
```

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```
type Eval6 alpha =
         ReaderT Env
3
                 (ErrorT String (WriterT [String] (StateT Integer IO)))
4
                 alpha
5
6
   runEval6 :: Env
                -> Integer
8
                -> Eval6 alpha
9
                -> IO ((Either String alpha, [String]), Integer)
10
    runEval6 env st ev =
             runStateT (runWriterT (runErrorT (runReaderT ev env))) st
11
12
13
    eval6
             :: Exp -> Eval6 Value
```

add IO: code

add IO: in action

```
1 runEval6 Map.empty 0 (eval6 exampleExp)
2 12
3 4
4 2
5 => IO ((Right (IntVal 18),["x"]),8)
```

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final version, page 1

```
eval6 (Lit i) = do tick -- profiling
                          liftIO $ print i -- print each int
3
                          return $ IntVal i
4
5
   eval6 (Var n) = do tick
6
                          tell [n] -- log var
                          env <- ask -- consult env
8
                          case Map.lookup n env of
9
                           Nothing -> throwError ("unbound: " ++ n)
10
                           Just val -> return val
11
12
   eval6 (Plus e1 e2) = do tick
13
                          e1' <- eval6 e1
14
                          e2' <- eval6 e2
15
                          case (e1', e2') of
16
                            (IntVal i1, IntVal i2)
17
                             -> return $ IntVal (i1 + i2)
18
                           _ -> throwError "dyn type err: Plus"
```

final version, page 2

```
eval6 (Abs n e) = do tick
                            env <- ask
3
                            return $ FunVal env n e
4
5
    eval6 (App e1 e2) = do tick
6
                            v1 <- eval6 e1
                            v2 <- eval6 e2
8
                            case v1 of
9
                             FunVal env' n body
10
                               -> local (const (Map.insert n v2 env'))
                                         (eval6 body)
11
12
                             _ -> throwError "dyn type err: App"
```

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- interactive version of this presentation at FPComplete :
 - https://www.fpcomplete.com/user/haroldcarr/ example-of-why-to-use-monads-what-they-can-do
- Maybe, Either, [], IO monads (and more to come):
 - http://haroldcarr.com/posts/2014-02-19-monad-series.html
- Dan Piponi's :
 - http://blog.sigfpe.com/2006/08/you-could-have-invented-monads-and.html
- illustrated :
 - http://adit.io/posts/2013-04-17-functors,_applicatives,_and_monads_in_ pictures.html
- Brent Yorgey :
 - http://www.haskell.org/haskellwiki/Typeclassopedia
- deep dive with Mike Vanier :
 - http://mvanier.livejournal.com/3917.html
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