Rabin-Karp for string matching

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History and application

Michael O. Rabin and Richard M. Karp (1987)

Example application: Detecting plagiarism

String matching

 $T \leftarrow Text$

P ← Pattern

Problem: $P \subseteq T$?

Naive algorithm

$$i, j \leftarrow 0, 0$$

- 1. Compare T[i] with P[j]
- 2. If match: increase i and j

2.1. If
$$j = len(P)$$
: $P \subseteq T$

- 3. If not match: increase i, $j \leftarrow 0$
- 4. If i = len(T): $P \notin T \times$
- 5. Repeat step 1

Rabin-Karp

What makes Rabin Karp special? Rolling hash!

- 1. Hash P
- 2. Hash |P| first characters of T
- 3. After this, we can find the next hash values in T in constant time by applying Horner's method!

Importance of hash function

We want to limit the amount of collisions

The hash function should map every unique input to a unique value

Easier said than done in practice

Demonstration

Assume we have -

a rolling hash function: hash(string) \rightarrow Int

a pattern: "ABA"

a text: "AABABCABACABBA"

 $hash("AAB") \rightarrow 17$

 $hash("ABA") \rightarrow 33$







A A B A B C A B A C A B B A

 $hash("ABA") \rightarrow 33$

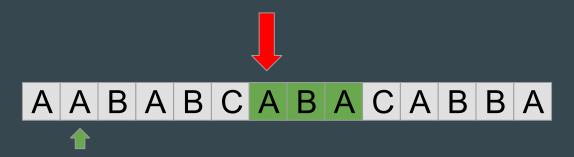


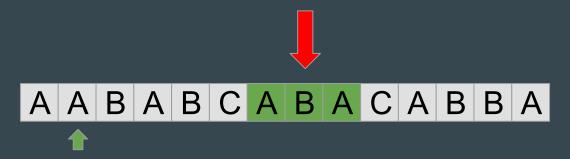
 $hash("ABC") \rightarrow 28$

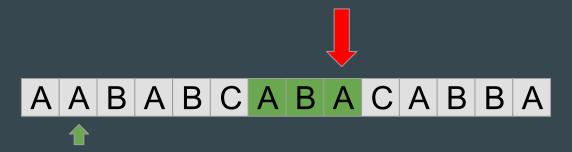
 $hash("BCA") \rightarrow 13$

A A B A B C A B A C A B B A

hash("CAB") \rightarrow 17









A A B A B C A B A C A B B A



 $hash("BAC") \rightarrow 51$



hash("ACA") \rightarrow 43

A A B A B C A B A C A B B A



hash("CAB") \rightarrow 17



hash("ABB") \rightarrow 33





hash("ABB") \rightarrow 33





hash("ABB") \rightarrow 33









hash("BBA") \rightarrow 47

Complexity and advantages/disadvantages

Average/best case: O(n)

Worst case: O(mn) = O(m) + O(m) * O(n)

Hash collisions and false positives!

Thanks for watching!

Sources:

- https://en.wikipedia.org/wiki/Rabin%E2%80%93Karp_algorithm
- https://www.uio.no/studier/emner/matnat/ifi/IN3130/h21/slides/forelesning-3---string-search-pk.pdf