代码簿

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1. 杂项

• Merge Sort

```
1 int arr[MAXLEN]; // 数组
2 int tmp[MAXLEN]; // 暂存
3 int ans = 0; // 逆序对
5 void mergeSort(int left, int right){
    if(left >= right) return;
    int mid = left + ((right - left) >> 1);
    mergeSort(left, mid);
   mergeSort(mid + 1, right);
9
   int 1 = left;
10
    int r = mid + 1;
11
    int idx = 0;
12
    while(1 <= mid && r <= right){</pre>
13
     tmp[idx++] = arr[1] <= arr[r] ? arr[1++]</pre>
14
      : ((ans += mid - l + 1 /* 逆序对 */ ), arr[r++]);
15
16
   while(1 <= mid){</pre>
17
18
     tmp[idx++] = arr[1++];
19
   while(r <= right){</pre>
20
21
     tmp[idx++] = arr[r++];
22
23
   for(int i = 0; i < idx; ++i){</pre>
     arr[left + i] = tmp[i];
25
26 }
```

• Disjoint Set

```
int djs[MAXLEN];
//init for(int i = 0; i < size; ++i)djs[i] = i;</pre>
4 int find(int id){
   return djs[id] == id ? id : djs[id] = find(djs[id]);
5
7 bool uSet(int a, int b){
   int pa = find(a), pb = find(b);
    return pa == pb ? true : (djs[pa] = pb, false);
10 }
12 //with dummy head
13
14 int djs[MAXLEN];
15 //init for(int i = 0; i < 2 * n; ++i) djs[i] = i; for(int i = 0; i < n; ++i) djs[i] = n + i;
int find(int id){
  return djs[id] == id ? id : djs[id] = find(djs[id]);
18
19 }
20 bool uSet(int a, int b){
int pa = find(a), pb = find(b);
return pa == pb ? true : (djs[pa] = pb, false);
23 }
```

2. 图论

• Eular Path

```
1 int cd[MAXLEN], rd[MAXLEN]; //有向图
2 vector<pair<int, int>> adjs[MAXLEN]; //first for 点, second fir 边
3 bool visit[MAXLEN]; //边
```

```
int cnt = 0; //判断联通(cnt == size)
7 void dfs(int id){
     for(auto &adj : adjs){
8
9
          if(!visit[adj.second]){
              ++cnt:
10
              visit[adj.second] = true;
11
12
              dfs(adj.first);
              //一些操作
13
          }
14
15
      }
16 }
17
18 bool judge(){
     //无向(边为偶数除了端点)
19
20
      int cnt = 0;
      int qidian = 0; // 起点
21
      for(int i = 0; i < n; ++i){
22
          if(adjs[i].size() & 1) ++cnt, qidian = i;
24
      return (cnt == 2 || cnt == 0);
25
26
      //有向((出度 - 入度) = 1 && (入度 - 出度) = 1 || 欧拉环)
int rdc = 0, cdc = 0;
27
28
      int qidian = 0;
29
      for(int i = 0; i < n; ++i){</pre>
30
          if((cd[i] - rd[i] == 1)) ++cdc;
31
          else if((rd[i] - cd[i]) == 1) ++rdc, qidian = i;
32
          else if((rd[i] - cd[i]) != 0) return false;
33
35
36
      return ((cdc == 1 && rdc == 1) || (cdc == 0 && rdc == 0));
37 }
```

• 拓扑排序

```
vector<int> adjs[MAXLEN];
2 int rd[MAXLEN]; //入度
3 int n; //点的数量
4 vector<int> ans; //结果
6 void topSort(){
      queue<int> q;
      int cnt= 0;
      for(int i = 0; i < n; ++i) if(!rd[i]) q.push(i), ++cnt, ans.push_back(i);</pre>
9
10
      while(!q.empty()){
          int f = q.front(); q.pop();
11
          for(auto adj : adjs[f]){
12
13
               --rd[adj];
              if(!rd[adj]) q.push(adj), ++cnt, ans.push_back(adj);
14
15
          }
16
      if(cnt == n) {}//有找到
17
      else {} //没找到
19 }
```

• 多元最短路

```
int v[MAXLEN][MAXLEN];

//初始化成0x3f3f3f3f, 有边的地方为权重, 自己为0

for(int k = 0; k < n; ++k){
    for(int i = 0; i < n; ++i){
        for(int j = 0; j < n; ++j){
            v[i][j] = min(v[i][k] + v[k][j]); //最短路径
            v[i][j] = min(v[i][j], max(v[i][k], v[k][j])); //最大的最小
        }

}

}
```

• 单元最短路

```
rector<pair<int, int>> adjs[MAXLEN]; //first for 点 second for 並权
int dis[MAXLEN];
bool visit[MAXLEN];

void dij(int id){
memset(dis, 0x3f, sizeof(dis));
memset(visit, 0, sizeof(visit));
```

```
dis[id] = 0;
8
      priority_queue<pii, vector<pii>, greater<pii>> q;
9
10
      //first for 目前距离 second for 点
11
      q.push({0, id});
12
13
      while(!q.empty()){
         pair<int, int> front = q.top(); q.pop();
14
          if(visit[front.second]) continue;
1.5
16
          visit[front.second] = true;
17
          for(auto adj : adjs[front.second]){
               if((adj.second + dis[front.second]) < dis[adj.first]){</pre>
18
19
                   dis[adj.first] = adj.second + dis[front.second];
                   if(visit[adj.first]) continue;
20
21
                   q.push({dis[adj.first], adj.first});
               }
22
          }
23
      }
24
25 }
```

• 找负环

```
1 vector<pair<int, int>> adjs[MAXLEN]; //first for 点 second for 权重
2 int dis[MAXLEN]; //最短路径
3 int m, n; //n 点数 m 边数
5 bool bellman(){
6
      memset(dis, 0x3f, 4 * n);
      for(int c = 0; c < n - 1; ++c){
        for(int i = 0; i < n; ++i){</pre>
          for(auto adj : adjs[i]){
9
            dis[adj.first] = min(dis[adj.first], dis[i] + adj.second);
11
12
        }
13
      bool hasAnswer = false;
14
15
    for(int i = 0; i < n; ++i){</pre>
16
     for(auto adj : adjs[i]){
       if(dis[adj.first] > dis[i] + adj.second){
17
          hasAnswer = true;
          goto ans;
19
        }
20
      }
21
    }
22
23
      ans:
      return hasAnswer;
25 }
```

• MST

```
1 // kruskal
3 pair<int, pii> adjs[MAXLEN]; //first for 权重 //second.first for 起点 //second.second for 终点
4 int n; //边的数量
5 int m; //点的数量
6 int djs[MAXLEN];
9 int parent(int i){
      return (djs[i] == i) ? i : djs[i] = parent(djs[i]);
10
11 }
void kruskal(){
     for(int i = 0; i <= m; ++i){</pre>
14
          djs[i] = i;
15
16
17
      sort(adjs, adjs + n);
18
      int cnt = m; //计算边的数量
19
      for(int i = 0; i < n; ++i){</pre>
20
         int p1 = parent(adjs[i].second.first), p2 = parent(adj[i].second.second);
21
22
          if(p1 == p2) continue;
          djs[p1] = p2;
23
          //一些操作
24
          if((--cnt) == 1) break;
25
26
27 }
29 //Prim 普利姆算法
31 vector<pii> adjs[MAXLEN]; //first for 点, second for 权重
```

```
32 bool visit[MAXLEN];
33
34 int n, m;
35 void prim(){
   priority_queue<pii, vector<pii>, greater<pii>> q;
36
    int ans = 0; //最小距离
37
    q.push({0, 1});
38
    int cnt = n; //是否联通
39
40
    while(!q.empty()){
     pii top = q.top();q.pop();
41
      if(visit[top.second])continue;
42
43
      ans += top.first;
      visit[top.second] = true;
44
45
      --cnt;
      for(const pii &adj : adjs[top.second]){
46
        if(!visit[adj.first]){
47
48
          q.push({adj.second, adj.first});
49
      }
50
    }
51
    if(cnt){
52
      printf("orz\n");
53
    }else{
54
      printf("%d\n", ans);
55
56
57
58 }
```

• 最大流 (Dinic)

```
1 struct edge
2 {
3
       int to;
4
       long long val;
       int rev:
5
       edge(int _to, long long _val, int _rev):to(_to), val(_val), rev(_rev){}
7 };
9 int n, m, s, t;
10
11
vector<edge> adjs[MAXLEN];
13
int depth[MAXLEN]; //深度
15
16 bool bfs(){ //分层 + 确定没有回去找
17
       memset(depth, 0, sizeof(depth));
       depth[s] = 1;
18
       queue<int> q;
19
20
       q.push(s);
       while(!q.empty()){
21
22
           int f = q.front(); q.pop();
           for(edge &adj : adjs[f]){
23
               if(!depth[adj.to] && adj.val){
24
25
                    depth[adj.to] = depth[f] + 1;
26
                    q.push(adj.to);
               }
27
28
           }
29
30
       return depth[t];
31 }
32
33 long long dfs(int u = s, long long in = 0x3f3f3f3f3f3f3f3f3f){ //多路增广
      if(u == t) return in;
34
       if(in == 0) return 0;
35
36
       long long out = 0;
       for(edge &adj : adjs[u]){
37
           if((depth[adj.to] == (depth[u] + 1)) && adj.val){
   long long dist = dfs(adj.to, min(in, adj.val));
38
               adj.val -= dist;
40
41
               adjs[adj.to][adj.rev].val += dist;
               in -= dist;
42
               out += dist;
43
           }
44
45
       if(!out) depth[u] = 0;
46
47
       return out;
48 }
50 long long ans = 0;
```

```
51
52  void dinic(){
53   while(bfs()){
54       while(long long d = dfs()){
55            ans += d;
56       }
57    }
58 }
```

• 费用流 (Dinic)

```
struct edge
2 {
      int to;
3
       int val;
      int rev;
5
6
      int cost;
      edge(int _to, int _val, int _rev, int _cost):to(_to), val(_val), rev(_rev), cost(_cost){} //反向边
       的cost要是负数
8 };
10 int n, m, s, t;
vector<edge> adjs[MAXLEN];
13
14 int dis[MAXLEN];
15 bool visit[MAXLEN];
16
17
18
19 long long zuidaliu = 0, zuixiaofeiyong = 0; //最大流and最小费用
20
bool spfa(int id = t){ //BellMan-Ford
   memset(visit, 0, sizeof(visit));
   memset(dis, 0x3f, sizeof(dis));
23
24
   visit[id] = true;
25
    dis[id] = 0;
   deque<int> q;
26
    q.push_front(id);
28
    while(!q.empty()){
29
      int front = q.front(); q.pop_front();
30
31
      for(edge &adj : adjs[front]){
32
        if(adjs[adj.to][adj.rev].val > 0 && dis[adj.to] > dis[front] - adj.cost ){
33
          dis[adj.to] = dis[front] - adj.cost;
34
35
           if(!visit[adj.to]){
            visit[adj.to] = true;
36
            if(!q.empty() && dis[adj.to] < dis[q.front()]) q.push_front(adj.to); //SLF优化
37
38
             else q.push_back(adj.to);
          }
39
        }
40
41
      visit[front] = false;
42
43
44
    return dis[s] != 0x3f3f3f3f;
45 }
46
47 int dfs(int u = s, int flow = 0x3f3f3f3f){ //多路增广
48
    if(u == t){
     visit[u] = true;
      return flow;
50
   }
51
    int in = flow;
52
    int out = 0;
visit[u] = true;
53
54
    for(edge &adj : adjs[u]){
55
56
      if(adj.val >0 && !visit[adj.to] && dis[adj.to] == (dis[u] - adj.cost)){
        int dis = dfs(adj.to, min(in, adj.val));
58
59
        if(dis > 0){
          in -= dis;
60
          adj.val -= dis;
61
62
          out += dis;
          zuixiaofeiyong += dis * adj.cost;
63
64
          adjs[adj.to][adj.rev].val += dis;
        if(!in) break;
66
67
68
```

```
69 }
70
71
    return out;
72 }
73
74
75
76 void dinic(){
   while(spfa()){
77
     while(int d = dfs()){
78
        memset(visit, 0, sizeof(visit));
79
80
        zuidaliu += d;
81
   }
82
83 }
```

3. 动态规划

• 区间覆盖

```
pii intv[MAXLEN];
2 int n; //区间数量
3 int end; //终点
6 int solve(){
     sort(intv, intv + n);
      int tmp = 0 /* 暂存 */ , r = 0 /* 右边更新 */;
      int sum = 0; //答案
9
      int i = 0;
11
      while(i < n && r < end){</pre>
          if(intv[i].first > r) break;
12
          while(intv[i].first <= r && tmp < end){</pre>
13
              tmp = max(intv[i].second, tmp);
              ++i;
15
16
         }
         r = tmp;
17
         ++sum;
18
19
      return r >= end : sum ? 0; // 没有则回传0
20
21 }
```

• LIS

```
int tmp[MAXLEN]; //暂存
int arr[MAXLEN]; //待处理数组
int dp[MAXLEN]; //各个结尾的值

void LIS(){

for(int i = 0; i <= n; ++i) tmp[i] = INT_MAX;

tmp[0] = 0;

for(int i = 0; i < n; ++i){
    int pos = lower_bound(tmp, tmp + n, arr[i]) - tmp; //lower_bound找 < // upper_bound找 <= tmp[pos] = arr[i], dp[i] = pos;
}

}
```

4. 字串

• KMP

```
int f[MAXLEN]; // failure Function 1-index
3 void kmp(char s[], char p[]){
   int i = 0, j = -1;
int n = strlen(s), m = strlen(p);f[0]=-1;
    while(i < m){</pre>
      if(j == -1 || p[i] == p[j]){
        f[++i] = ++j;
     }else{
        j = f[j];
10
      }
11
12
    i = 0; j = 0;
13
14
    while(i <= n){</pre>
     if(j == -1 || s[i] == p[j]){
15
   ++j, ++i;
16
```

马拉车

```
1 char s[MAXLEN]; //字串 记得填充$#^(aaa-> $#a#a#a#^)
2 int p[MAXLEN]; //回文长度
4 int len = 0 //s的长度
6 void manache(){
   int i_r = 0, c = 0, r = 0; //i_r 反射点, c 中心点 r 右端点
    for(int i = 1; i < len; ++i){</pre>
      i_r = 2 * c - i;
9
      if(r > i) p[i] = min(r - i, p[i_r]);
      else p[i] = 0;
11
     while (s[i + 1 + p[i]] == s[i - 1 - p[i]]) ++p[i];
12
      if(i + p[i] > r) c = i, r = i + p[i];
13
   }
14
   int m = 0; // 长度
15
   c = 0; // 中央
16
   for(int i = 0; i < len - 1; ++i){</pre>
17
     if(p[i] > m) c = i, m = p[i];
18
19
   printf("%d\n", m);
20
21 }
```

• AC 自动机 + 拓扑优化

```
1 int cnt = 0;
int trie[MAXLEN][26];
3 int fail[MAXLEN];
4 int rd[MAXLEN]; //拓扑必须
5 char c:
8 void insert(char* s) { //trie插入
   int id = 0;
    for (; *s; ++s) {
  c = *s - 'a';
10
11
     if (!trie[id][c]) trie[id][c] = ++cnt;
     id = trie[id][c];
13
   }
14
   //一些操作(在trie上产生节点)
15
16 }
18 void build() { //记得放在main
19
   queue<int> q;
    for (int i = 0; i < 26; ++i) {
     if (trie[0][i])q.push(trie[0][i]);
21
22
23
    while (!q.empty()) {
     int u = q.front(); q.pop();
24
      for (int i = 0; i < 26; ++i) {</pre>
       if (trie[u][i]) {
26
         int t = trie[u][i];
27
         fail[t] = trie[fail[u]][i];
         q.push(t);
29
          ++rd[fail[t]]; //拓扑必须
30
31
       else {
32
33
         trie[u][i] = trie[fail[u]][i];
34
      }
35
    }
36
37 }
38
39 void query(char* s) {
40 int id = 0;
    for (int i = 0; s[i]; ++i) {
41
   id = trie[id][s[i] - 'a'];
42
     //一些懒标操作(复杂度较低)
43
for (int j = id; j; j = fail[j]) {
```

```
45 //直接操作
46 }
46
47 }
48 }
49
50 void topo() {
queue < int > q;
   for (int i = 1; i <= cnt; ++i)if (!rd[i])q.push(i);</pre>
52
   while (!q.empty()) {
53
     int u = q.front(); q.pop();
54
     //一些操作(更新当前值)
55
56
     int f = fail[u];
     --rd[f];
57
    //一些操作(懒标回推)
58
     if (!rd[f])q.push(f);
59
60
61 }
```

5. 数据结构

• 串列

```
struct {
     int 1 = 0, r = 0;
3 } lst[MAXLEN];
5 void init() {
      memset(lst, -1, sizeof(lst));
6
      lst[1].1 = 0;
      lst[0].r = 1;
9 }
10
void addl(int now, int node) { //加在node的左边
      lst[now].r = node;
12
13
      lst[lst[node].l].r = now;
      lst[now].1 = lst[node].1;
14
      lst[node].1 = now;
15
16 }
17
18 void addr(int now, int node) { //加在node的右边
      lst[now].1 = node;
20
      lst[lst[node].r].l = now;
21
      lst[now].r = lst[node].r;
      lst[node].r = now;
22
23 }
24
void remove(int node) {
    if (!~lst[node].1)return;
26
      lst[lst[node].1].r = lst[node].r;
      lst[lst[node].r].l = lst[node].l;
28
29
      lst[node].l = lst[node].r = -1;
30 }
```

• 稀疏表

```
int Log2[MAXLEN]; //预处理log2
2 int st[MAXLEN][20]; //稀疏表
3 int n, m; //n个节点 m个query
5 void pre() {
   Log2[1] = 0, Log2[2] = 1;
    for (int i = 3; i <= MAXLEN; ++i) {
  Log2[i] = Log2[i >> 1] + 1;
9
10 }
11
void build() {
   for (int j = 1; j <= 20; ++j) {
13
     for (int i = 1; i + (1 << j) - 1 <= n; ++i) {
        st[i][j] = min(st[i][j-1], st[i+(1 << (j-1))][j-1]);
15
16
   }
17
18 }
inline int query(int 1, int r) {
int lo = Log2[r - 1 + 1];
return min(st[1][lo], st[r - (1 << lo) + 1][lo]);</pre>
23 }
```

Fenwick

```
#define lowbit(x) (x & -x)
int n; //总数量

namespace fw{
    int tree[MAXLEN] = {0};
    void insert(int pos, int val){ //单点插入
        for(;pos <= n; pos += lowbit(pos))tree[pos] += val;
    }
    int search(int pos){ //单点前缀和查询
    int ans = 0;
    for(;pos; pos -= lowbit(pos)){
        ans += tree[pos];
    }
    return ans;
}

int query(int l, int r){ //区间查询
    return search(r) - search(l - 1);
}

}
```