代码簿

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1. 杂项

• VIMRC

```
set hls
set is
set cb=unnamed
set ts=4
set sw=4
set si
set nu

inoremap { {}<Left>
inoremap {<CR> {<CR>}<Esc>0
inoremap {} {}
```

• Merge Sort

```
int arr[MAXLEN]; // 数组
2 int tmp[MAXLEN]; // 暂存
3 int ans = 0; // 逆序对
5 void mergeSort(int left, int right){
   if(left >= right) return;
    int mid = left + ((right - left) >> 1);
   mergeSort(left, mid);
   mergeSort(mid + 1, right);
9
   int 1 = left;
    int r = mid + 1;
11
    int idx = 0;
12
   while(1 <= mid && r <= right){</pre>
13
     tmp[idx++] = arr[1] <= arr[r] ? arr[1++]</pre>
14
      : ((ans += mid - l + 1 /* 逆序对 */ ), arr[r++]);
15
16
   while(1 <= mid){</pre>
17
18
     tmp[idx++] = arr[l++];
19
   while(r <= right){</pre>
20
21
     tmp[idx++] = arr[r++];
22
   for(int i = 0; i < idx; ++i){</pre>
23
24
      arr[left + i] = tmp[i];
25
26 }
```

• Disjoint Set

```
int djs[MAXLEN];
2 //init for(int i = 0; i < size; ++i)djs[i] = i;</pre>
4 int find(int id){
return djs[id] == id ? id : djs[id] = find(djs[id]);
6 }
7 bool uSet(int a, int b){
8 int pa = find(a), pb = find(b);
   return pa == pb ? true : (djs[pa] = pb, false);
10 }
11
12 //with dummy head
13
14 int djs[MAXLEN];
15 //init for(int i = 0; i < 2 * n; ++i) djs[i] = i; for(int i = 0; i < n; ++i) djs[i] = n + i;
16
int find(int id){
return djs[id] == id ? id : djs[id] = find(djs[id]);
```

```
19 }
bool uSet(int a, int b){
21   int pa = find(a), pb = find(b);
22   return pa == pb ? true : (djs[pa] = pb, false);
23 }
```

2. 图论

• Eular Path

```
1 int cd[MAXLEN], rd[MAXLEN]; //有向图
2 vector<pair<int, int>> adjs[MAXLEN]; //first for 点, second fir 边
3 bool visit[MAXLEN]; //边
5 int cnt = 0; //判断联通(cnt == size)
7 void dfs(int id){
      for(auto &adj : adjs){
          if(!visit[adj.second]){
              ++cnt;
10
              visit[adj.second] = true;
11
              dfs(adj.first);
12
              //一些操作
13
          }
14
15
16 }
18 bool judge(){
      //无向(边为偶数除了端点)
19
      int cnt = 0;
20
      int qidian = 0; // 起点
21
      for(int i = 0; i < n; ++i){</pre>
22
          if(adjs[i].size() & 1) ++cnt, qidian = i;
23
24
      return (cnt == 2 || cnt == 0);
26
      //有向((出度 - 入度) = 1 && (入度 - 出度) = 1 || 欧拉环)
27
      int rdc = 0, cdc = 0;
      int qidian = 0;
29
      for(int i = 0; i < n; ++i){</pre>
30
          if((cd[i] - rd[i] == 1)) ++cdc;
31
          else if((rd[i] - cd[i]) == 1) ++rdc, qidian = i;
32
          else if((rd[i] - cd[i]) != 0) return false;
33
34
35
      return ((cdc == 1 && rdc == 1) || (cdc == 0 && rdc == 0));
37 }
```

• 拓扑排序

```
vector<int> adjs[MAXLEN];
2 int rd[MAXLEN]; //入度
3 int n; //点的数量
4 vector<int> ans; //结果
6 void topSort(){
     queue<int> q;
      int cnt= 0;
      for(int i = 0; i < n; ++i) if(!rd[i]) q.push(i), ++cnt, ans.push_back(i);</pre>
9
      while(!q.empty()){
          int f = q.front(); q.pop();
          for(auto adj : adjs[f]){
12
              --rd[adj];
13
              if(!rd[adj]) q.push(adj), ++cnt, ans.push_back(adj);
14
          }
15
16
      if(cnt == n) {}//有找到
17
      else {} //没找到
18
19 }
```

• 多元最短路

```
int v[MAXLEN][MAXLEN];
//初始化成0x3f3f3f3f, 有边的地方为权重, 自己为0

for(int k = 0; k < n; ++k){
    for(int i = 0; i < n; ++i){
        for(int j = 0; j < n; ++j){
            v[i][j] = min(v[i][k] + v[k][j], v[i][j]); //最短路径
```

```
8 v[i][j] = min(v[i][j], max(v[i][k], v[k][j])); //最大的最小
9 }
10 }
11 }
```

• 单元最短路

```
1 vector<pair<int, int>> adjs[MAXLEN]; //first for 点 second for 边权
int dis[MAXLEN];
3 bool visit[MAXLEN];
5 void dij(int id){
     memset(dis, 0x3f, sizeof(dis));
      memset(visit, 0, sizeof(visit));
      dis[id] = 0;
      priority_queue<pii, vector<pii>, greater<pii>> q;
9
      //first for 目前距离 second for 点
      q.push({0, id});
12
13
      while(!q.empty()){
          pair<int, int> front = q.top(); q.pop();
14
          if(visit[front.second]) continue;
1.5
          visit[front.second] = true;
16
          for(auto adj : adjs[front.second]){
17
              if((adj.second + dis[front.second]) < dis[adj.first]){</pre>
18
                   dis[adj.first] = adj.second + dis[front.second];
19
                   if(visit[adj.first]) continue;
20
21
                   q.push({dis[adj.first], adj.first});
              }
22
          }
23
24
      }
25 }
```

• 找负环

```
vector<pair<int, int>> adjs[MAXLEN]; //first for 点 second for 权重
2 int dis[MAXLEN]; //最短路径
3 int m, n; //n 点数 m 边数
5 bool bellman(){
      memset(dis, 0x3f, 4 * n);
6
      for(int c = 0; c < n - 1; ++c){
8
        for(int i = 0; i < n; ++i){</pre>
          for(auto adj : adjs[i]){
9
            dis[adj.first] = min(dis[adj.first], dis[i] + adj.second);
11
        }
12
      }
13
      bool hasAnswer = false:
14
    for(int i = 0; i < n; ++i){</pre>
15
     for(auto adj : adjs[i]){
16
       if(dis[adj.first] > dis[i] + adj.second){
17
          hasAnswer = true;
18
          goto ans;
19
        }
20
      }
21
    }
22
23
      ans:
24
      return hasAnswer;
25 }
```

• MST

```
1 // kruskal
3 pair<int, pii> adjs[MAXLEN]; //first for 权重 //second.first for 起点 //second.second for 终点
4 int n; //边的数量
5 int m; //点的数量
6 int djs[MAXLEN];
9 int parent(int i){
      return (djs[i] == i) ? i : djs[i] = parent(djs[i]);
10
11 }
void kruskal(){
      for(int i = 0; i <= m; ++i){</pre>
14
15
          djs[i] = i;
16
17
```

```
sort(adjs, adjs + n);
18
      int cnt = m; // 计算边的数量
19
      for(int i = 0; i < n; ++i){</pre>
20
          int p1 = parent(adjs[i].second.first), p2 = parent(adj[i].second.second);
          if(p1 == p2) continue;
22
          djs[p1] = p2;
23
          //一些操作
24
          if((--cnt) == 1) break;
25
26
27 }
28
29 //Prim 普利姆算法
30
31 vector<pii> adjs[MAXLEN]; //first for 点, second for 权重
32 bool visit[MAXLEN];
33
34 int n, m;
35 void prim(){
   priority_queue<pii, vector<pii>, greater<pii>> q;
    int ans = 0; //最小距离
   q.push({0, 1});
38
    int cnt = n; //是否联通
39
    while(!q.empty()){
     pii top = q.top();q.pop();
41
      if(visit[top.second])continue;
42
     ans += top.first;
43
44
     visit[top.second] = true;
45
      --cnt;
      for(const pii &adj : adjs[top.second]){
46
47
        if(!visit[adj.first]){
48
          q.push({adj.second, adj.first});
49
     }
50
51
    if(cnt){
52
     printf("orz\n");
53
    }else{
54
     printf("%d\n", ans);
55
56
57
58 }
```

• 最大流 (Dinic)

```
struct edge
2 {
      int to;
      long long val;
5
      int rev;
      edge(int _to, long long _val, int _rev):to(_to), val(_val), rev(_rev){}
6
7 };
9 int n, m, s, t;
10
vector<edge> adjs[MAXLEN];
13
14 int depth[MAXLEN]; //深度
15
16 bool bfs(){ //分层 + 确定没有回去找
      memset(depth, 0, sizeof(depth));
17
      depth[s] = 1;
18
19
      queue<int> q;
20
      q.push(s);
21
      while(!q.empty()){
22
          int f = q.front(); q.pop();
          for(edge &adj : adjs[f]){
23
              if(!depth[adj.to] && adj.val){
24
                  depth[adj.to] = depth[f] + 1;
                  q.push(adj.to);
26
27
              }
          }
28
29
30
      return depth[t];
31 }
32
33 long long dfs(int u = s, long long in = 0x3f3f3f3f3f3f3f3f3f){ //多路增广
      if(u == t) return in;
34
      if(in == 0) return 0;
35
36 long long out = 0;
```

```
for(edge &adj : adjs[u]){
37
           if((depth[adj.to] == (depth[u] + 1)) && adj.val){
38
               long long dist = dfs(adj.to, min(in, adj.val));
39
               adj.val -= dist;
40
               adjs[adj.to][adj.rev].val += dist;
41
42
               in -= dist;
               out += dist;
43
           }
44
45
       if(!out) depth[u] = 0;
46
47
       return out;
48 }
49
50 long long ans = 0;
51
52 void dinic(){
53
       while(bfs()){
          while(long long d = dfs()){
54
               ans += d;
55
56
57
      }
58 }
```

· 费用流 (Dinic)

```
struct edge
2 {
      int to;
      int val;
      int rev;
5
6
      edge(int _to, int _val, int _rev, int _cost):to(_to), val(_val), rev(_rev), cost(_cost){} //反向边
      的cost要是负数
8 };
10 int n, m, s, t;
vector<edge> adjs[MAXLEN];
int dis[MAXLEN];
bool visit[MAXLEN];
17
19 long long zuidaliu = 0, zuixiaofeiyong = 0; //最大流and最小费用
20
21 bool spfa(int id = t){ //BellMan-Ford
memset(visit, 0, sizeof(visit));
   memset(dis, 0x3f, sizeof(dis));
23
    visit[id] = true;
   dis[id] = 0;
25
26
    deque<int> q;
    q.push_front(id);
28
29
    while(!q.empty()){
30
      int front = q.front(); q.pop_front();
31
32
      for(edge &adj : adjs[front]){
        if(adjs[adj.to][adj.rev].val > 0 && dis[adj.to] > dis[front] - adj.cost ){
33
          dis[adj.to] = dis[front] - adj.cost;
34
          if(!visit[adj.to]){
35
            visit[adj.to] = true;
36
            if(!q.empty() && dis[adj.to] < dis[q.front()]) q.push_front(adj.to); //SLF优化
37
38
            else q.push_back(adj.to);
          }
39
        }
40
41
      visit[front] = false;
42
43
    return dis[s] != 0x3f3f3f3f;
44
45 }
46
int dfs(int u = s, int flow = 0x3f3f3f3f){ //多路增广
   if(u == t){
     visit[u] = true;
49
50
      return flow;
51
   int in = flow;
52
53
    int out = 0;
visit[u] = true;
```

```
for(edge &adj : adjs[u]){
56
      if(adj.val >0 && !visit[adj.to] && dis[adj.to] == (dis[u] - adj.cost)){
57
        int dis = dfs(adj.to, min(in, adj.val));
        if(dis > 0){
59
          in -= dis;
60
         adj.val -= dis;
61
          out += dis;
62
          zuixiaofeiyong += dis * adj.cost;
63
          adjs[adj.to][adj.rev].val += dis;
64
65
66
        if(!in) break;
67
68
69
70
71
    return out;
72 }
73
74
75
76 void dinic(){
   while(spfa()){
     while(int d = dfs()){
78
        memset(visit, 0, sizeof(visit));
79
        zuidaliu += d;
80
81
   }
82
83 }
```

3. 树问题

• 树直径

```
1 //DFS两次
 3 int d[MAXN];
 vector<int> adjs[MAXN];
6 int last; //最远点, d[last] 为数的直径 不可做负边
8 int bfs(int u, int p){
       queue<pair<int, int>> q;
       q.push({u, p});
10
11
       while(!q.empty()){
           auto [u, p] = q.front(); q.pop();
for(const int adj : adjs[u]){
   if(adj == p)
13
14
15
                     continue;
                d[adj] = d[u] + 1;
16
17
                last = adj;
                q.push({adj, u});
18
            }
19
       }
21 }
22
int farPoint(int u, int p){
       last = u;
24
25
       bfs(u, p);
       d[last] = 0;
26
       bfs(last, p);
27
       return d[last];
29 }
30
31 //树形DP, 可做负边
33 int n, d = 0;
34 int d1[N], d2[N];
vector<int> adjs[N];
37 void dfs(int u, int p) {
38  d1[u] = d2[u] = 0;
    for (int v : adjs[u]) {
     if (v == p)
40
            continue;
41
42
     dfs(v, u);
int t = d1[v] + 1; //可改权重
if (t > d1[u])
d2[u] = d1[u], d1[u] = t;
```

```
46    else if (t > d2[u])
47     d2[u] = t;
48    }
49    d = max(d, d1[u] + d2[u]);
50 }
```

• LCA

```
1 // 倍增
3 int dep[MAXN];
4 int pa[MAXN][32];
5 vector<int> ch[MAXN];
8 void dfs(int u, int p){ // Pre-process
      dep[u] = dep[p] + 1;
9
      pa[u][0] = p;
10
      for(int i = 1; i < 32; ++i)</pre>
12
          pa[u][i] = pa[pa[u][i - 1]][i - 1];
13
14
      for(const int v : ch[u]){
15
          if(v == p)
16
17
               continue;
18
          dfs(v, u);
19
20 }
21
22 int lca(int a, int b){
    if(dep[a] > dep[b])
23
24
          swap(a, b);
25
      int diff = dep[b] - dep[a];
26
      int i = 0;
28
      //调整高度
29
     for(; diff; ++i, diff >>= 1)
30
          if (diff & 1)
31
32
               b = pa[b][i];
33
     if(a == b)
34
35
          return a;
36
37
      i = 0;
      for(i = 31; i >= 0; --i)
39
          if(pa[a][i] != pa[b][i])
40
41
               a = pa[a][i], b = pa[b][i];
42
43
      return pa[a][0];
44
45 }
```

4. 计算几何

• 凸包

```
1 //Andrew
4 // 点,上凸包,下凸包
6 using pdd = pair<double, double>
7 pdd ps[MAXN], u[MAXN], d[MAXN];
9 int ui = 0, di = 0;
10
inline double cross(const pdd &p1, const pdd &p2, const pdd &p3){
     return (p2.first - p1.first) * (p3.second - p1.second) - (p2.second - p1.second) * (p3.first - p1.
      first);
13 }
14
15
void andrew(){
17
    sort(ps, ps + n);
      for(int i = 0; i < n; ++i){</pre>
18
19
          while(ui >= 2 && cross(u[ui - 2], u[ui - 1], ps[i]) <= 0)</pre>
           --ui;
20
```

5. 数论

• Miller Robin(确认素数)

```
1 // allow condition 1, 1, 1, 1, .....
                       X, X, X, X, X, p - 1, 1, 1, 1, ....
p - 1, 1, 1, 1, 1, ....
2 //
3 //
4 //
                       0 (cannot check) multiply of p
5
6 bool is_prime(ll p){
     if (p < 3 || p % 2 == 0)
           return p == 2;
9
       11 u = p - 1;
       int t = 0;
10
       while(!(u & 1))
11
12
          u >>= 1, ++t;
13
14
      const static 11 ud[] = {2, 325, 9375, 28178, 450775, 9780504, 1795265022};
15
      for(const 11 a : ud){
16
17
           ll v = qpow(a, u, p);
           if (v <= 1 || v == p - 1)
18
                continue;
19
20
           for(int j = 1; j <= t; ++j){</pre>
21
               v = qmul(v, v, p);
22
                if(v == p - 1 \&\& j != t){
23
                   v = 1;
24
25
                    break;
               }
26
                if(v == 1)
27
                   return false;
29
           if (v != 1)
30
31
               return false;
       }
32
33
       return true;
34 }
```

• Pollard-Rho(确认因数)

```
// remember srand(time(NULL))
3 11 PR(11 n){
      if(n == 4)
4
          return 2;
6
      if(is_prime(n))
7
          return n;
9
      while(1){
10
          ll c = (rand() % (n - 1)) + 1;
11
           auto f = [&](const ll x){
12
               return ((111) x * x + c) % n;
13
14
           11 t = 0, r = 0, p = 1, q;
15
16
           do{
               for(int i = 0; i < 128; ++i){</pre>
17
                   t = f(t), r = f(f(r));
18
                   if(t == r \mid \mid (q = (111) p * abs(t - r) % n) == 0)
19
                       break;
20
21
                   p = q;
22
               11 d = __gcd(p, n);
23
               if (d > 1)
24
25
                  return d;
           }while(t != r);
26
27
28 }
```

• FastGCD

```
1 // 定理1: gcd(x, y) = gcd(y, x mod y) 辗转相除法
_2 // 定理2: x = a * b * c, where a, b,c <= sqrt(x) or isPrime
3 // 定理3: gcd(x, y) = gcd(a, y) * gcd(x / a, y / gcd(a, y))
6 int minp[MAXN] = {0}; //最小质因数
7 int g[1010][1010]; //1000*1000 最大公因数表
8 int fac[MAXN][3]; //三个拆分的因数
9 bitset<MAXN> vis;
10
void init(){
      fac[1][0] = fac[1][1] = fac[1][2] = 1;
12
13
       for(int i = 1; i < MAXN; ++i){</pre>
           if(!vis[i]){ // Prime
14
               fac[i][0] = fac[i][1] = fac[i][2] = i;
15
                if (111 * i * i > MAXN)
16
                    continue;
17
                for(int j = i * i; j <= MAXN; j += i){
    vis[j] = 1;</pre>
18
19
                    if(!minp[j])
20
21
                        minp[j] = i;
               }
22
           }else{
23
                int tmp = i / minp[i], a = minp[i] * fac[tmp][0], b = fac[tmp][1], c = fac[tmp][2];
24
                if(a > b)
25
26
                   swap(a, b);
27
                if(b > c)
                   swap(b, c);
28
29
                fac[i][0] = a, fac[i][1] = b, fac[i][2] = c;
           }
30
       }
31
32
      for(int i = 1; i <= 1000; ++i){</pre>
33
           g[i][0] = g[0][i] = i;
34
           for(int j = 1; j <= i; ++j)</pre>
35
                g[i][j] = g[j][i] = g[j][i % j];
36
37
38 }
39
40 itn gcd(int x, int y){
    int tmp = g[fac[x][0]][y % fac[x][0]], res = tmp;
41
42
       y /= tmp;
       tmp = g[fac[x][1]][y \% fac[x][1]], res *= tmp, y /= tmp;
tmp = fac[x][2] > 1000 ? (y \% fac[x][2] == 0 ? fac[x][2] : 1) : g[fac[x][2]][y % fac[x][2]], res
43
44
       *= tmp;
45
       return res;
46 }
```

• 欧拉函数

```
1 // 单一
2
3 int Ep(int n){
      int lim = sqrt(n);
      int res = n;
5
6
      for(int i = 2; i <= lim; ++i){</pre>
          if(n % i == 0){
8
9
              res = res / i * (i - 1);
               while (n \% i == 0)
10
                  n /= i;
11
          }
12
13
      if(n > 1)
14
          res = res / n * (n - 1);
15
      return res;
16
17 }
18
19 // 值域
21 bool vis[MAXN] = {false};
22 int primes[MAXN];
23 int phi[MAXN];
24
void build(){
    int tot = 0;
26
      phi[1] = 1;
27
```

```
29     for(int i = 2; i <= n; ++i){</pre>
          if(!vis[i])
30
               primes[tot++] = i, phi[i] = i - 1;
31
           for(int j = 0; j < tot && i * primes[j] <= n; ++j){</pre>
               vis[i * primes[j]] = true;
33
34
               if(i % primes[j])
                  phi[i * primes[j]] = phi[i] * (primes[j] - 1);
35
               elsef
36
                   phi[i * primes[j]] = primes[j] * phi[i];
37
38
                   break;
               }
39
40
          }
      }
41
42 }
```

• 扩展欧几里得

```
_1 // ax+by=gcd(a, b) if gcd(a, b) == 1
2 // ax同余c(mod b) <=> ax + by = c
3 // c 同余 a / b (mod M) = bx 同余 a (mod M)
5 11 Exgcd(11 a, 11 b, 11 &x, 11 &y){
      if(!b)
          return x = 1, y = 0, a;
      11 d = Exgcd(b, a % b, x, y);
      11 t = x;
      x = y;
y = t - (a / b) * y;
10
11
      return d;
12
13 }
15 //线性
16
17 void inv(){
    inv[1] = 1;
18
      for(int i = 2; i <= n; ++i)</pre>
19
20
       rev[i] = (ll)(p - p / i) * rev[p % i] % p;
21 }
23
24 // 随机线性
26 int arr[n + 1], ji[n + 1], rev[n + 1], inv[n + 1]; // 数组, 前缀积, 前缀积逆元
28 ji[0] = 1;
29
30 for(int i = 1; i <= n; ++i)</pre>
     ji[i] = ji[i - 1] * arr[i] % p;
31
32
33 rev[n] = (exgcd(ji[n], p), x);
34
35 for(int i = n; i >= 1; --i)
     rev[i - 1] = (ll)rev[i] * arr[i] % p;
37
38 for(int i = 1; i <= n; ++i)
inv[i] = (ll)rev[i] * ji[i - 1] % p;
```

• 中国剩余定理

```
1 // r互素
3 int a[MAXN], r[MAXN];
4 int n = 1, ans = 0;
7 for(int i = 0; i < k; ++i)</pre>
   n *= r[i];
10 for(int i = 0; i < k; ++i){</pre>
11 ll m = n / r[i];
12
   exgcd(m, r[i]);
   ans = (ans + ((((lll)a[i] * m) % n) * x % n)) % n;
13
14 }
15
16
17 // r非互素
19 ll M = r[0], ans = a[0];
21 for(int i = 1; i < k; ++i){</pre>
```

```
11 da = (((a[i] - ans) % r[i]) + r[i]) % r[i];
11 g = exgcd(M, r[i]);
24 auto tmp = (111)M * x * (da / g);
25 M = M * (r[i] / g);
26 ans = (((ans + tmp) % M) + M) % M;
27 }
```

6. 动态规划

• 区间覆盖

```
pii intv[MAXLEN];
2 int n; //区间数量
3 int end; //终点
6 int solve(){
       sort(intv, intv + n);
       int tmp = 0 /* 暂存 */ , r = 0 /* 右边更新 */;
      int sum = 0; //答案
       int i = 0;
10
       while(i < n && r < end){</pre>
1.1
          if(intv[i].first > r) break;
           \label{lem:while} \begin{picture}(i < n && intv[i].first <= r && tmp < end) \\ \end{picture}
13
               tmp = max(intv[i].second, tmp);
14
15
               ++i;
           }
16
           r = tmp;
17
           ++sum;
18
19
20
      return r >= end : sum ? 0; // 没有则回传0
21 }
```

• LIS

```
int tmp[MAXLEN]; //暂存
int arr[MAXLEN]; //待处理数组
int dp[MAXLEN]; //各个结尾的值

void LIS(){

for(int i = 0; i <= n; ++i) tmp[i] = INT_MAX;

tmp[0] = 0;

for(int i = 0; i < n; ++i){

int pos = lower_bound(tmp, tmp + n, arr[i]) - tmp; //lower_bound找 < // upper_bound找 <= tmp[pos] = arr[i], dp[i] = pos;
}

tmp[pos] = arr[i], dp[i] = pos;
}
```

7. 字串

• KMP

```
int f[MAXLEN]; // failure Function 1-index
3 void kmp(char s[], char p[]){
   int i = 0, j = -1;
    int n = strlen(s), m = strlen(p);f[0]=-1;
    while(i < m){</pre>
6
      if(j == -1 || p[i] == p[j]){
       f[++i] = ++j;
     }else{
9
10
       j = f[j];
11
    }
12
    i = 0; j = 0;
13
    while(i <= n){</pre>
14
     if(j == -1 || s[i] == p[j]){
15
        ++j, ++i;
16
        if(j == m){
17
          //一些操作
18
19
20
     }else{
21
        j = f[j];
22
23
   }
25 }
```

• 马拉车

```
1 char s[MAXLEN]; //字串 记得填充$#^(aaa-> $#a#a#a#^)
2 int p[MAXLEN]; //回文长度
4 int len = 0 //s的长度
6 void manache(){
   int i_r = 0, c = 0, r = 0; //i_r 反射点, c 中心点 r 右端点
    for(int i = 1; i < len; ++i){</pre>
     i_r = 2 * c - i;
9
      if(r > i) p[i] = min(r - i, p[i_r]);
10
      else p[i] = 0;
11
     while(s[i + 1 + p[i]] == s[i - 1 - p[i]]) ++p[i];
12
     if(i + p[i] > r) c = i, r = i + p[i];
14
   int m = 0; // 长度
15
   c = 0; //中央
   for(int i = 0; i < len - 1; ++i){</pre>
17
   if(p[i] > m) c = i, m = p[i];
}
18
19
   printf("%d\n", m);
20
21 }
```

• AC 自动机 + 拓扑优化

```
int cnt = 0;
int trie[MAXLEN][26];
3 int fail[MAXLEN];
4 int rd[MAXLEN]; //拓扑必须
5 char c;
8 void insert(char* s) { //trie插入
    int id = 0;
   for (; *s; ++s) {
  c = *s - 'a';
10
      if (!trie[id][c]) trie[id][c] = ++cnt;
12
     id = trie[id][c];
13
14
    //一些操作(在trie上产生节点)
15
16 }
17
18 void build() { //记得放在main
   queue<int> q;
19
   for (int i = 0; i < 26; ++i) {
20
     if (trie[0][i])q.push(trie[0][i]);
21
22
   while (!q.empty()) {
23
24
     int u = q.front(); q.pop();
      for (int i = 0; i < 26; ++i) {</pre>
25
       if (trie[u][i]) {
26
         int t = trie[u][i];
27
28
          fail[t] = trie[fail[u]][i];
          q.push(t);
29
          ++rd[fail[t]]; //拓扑必须
30
31
32
        else {
         trie[u][i] = trie[fail[u]][i];
        }
34
35
      }
   }
36
37 }
39 void query(char* s) {
   int id = 0;
40
    for (int i = 0; s[i]; ++i) {
41
     id = trie[id][s[i] - 'a'];
42
      //一些懒标操作(复杂度较低)
43
     for (int j = id; j; j = fail[j]) {
    //直接操作
44
45
      }
    }
47
48 }
50 void topo() {
51
   queue<int> q;
   for (int i = 1; i <= cnt; ++i)if (!rd[i])q.push(i);</pre>
52
   while (!q.empty()) {
53
int u = q.front(); q.pop();
```

8. 数据结构

• 串列

```
struct {
     int 1 = 0, r = 0;
3 } lst[MAXLEN];
5 void init() {
     memset(lst, -1, sizeof(lst));
6
      lst[1].1 = 0;
      lst[0].r = 1;
8
9 }
10
void addl(int now, int node) { //加在node的左边
12
    lst[now].r = node;
13
      lst[lst[node].1].r = now;
      lst[now].1 = lst[node].1;
14
15
     lst[node].1 = now;
16 }
17
18 void addr(int now, int node) { //加在node的右边
      lst[now].1 = node;
19
20
      lst[lst[node].r].l = now;
      lst[now].r = lst[node].r;
21
     lst[node].r = now;
22
23 }
24
void remove(int node) {
     if (!~lst[node].1)return;
      lst[lst[node].1].r = lst[node].r;
27
      lst[lst[node].r].1 = lst[node].1;
28
      lst[node].l = lst[node].r = -1;
29
30 }
```

• 稀疏表

```
1 int Log2[MAXLEN]; //预处理log2
2 int st[MAXLEN][20]; //稀疏表
3 int n, m; //n个节点 m个query
5 void pre() {
\log 2[1] = 0, \log 2[2] = 1;
    for (int i = 3; i <= MAXLEN; ++i) {</pre>
      Log2[i] = Log2[i >> 1] + 1;
9
10 }
11
void build() {
   for (int j = 1; j <= 20; ++j) {
     for (int i = 1; i + (1 << j) - 1 <= n; ++i) {
  st[i][j] = min(st[i][j - 1], st[i + (1 << (j - 1))][j - 1]);</pre>
14
15
16
    }
17
18 }
19
20 inline int query(int 1, int r) {
   int lo = Log2[r - 1 + 1];
21
    return min(st[1][lo], st[r - (1 << lo) + 1][lo]);</pre>
22
23 }
```

• Fenwick

```
#define lowbit(x) (x & -x)
int n; //总数量
3 //逆序数对由"大排到小", search(loc - 1)就行了回圈就行了

namespace fw{
int tree[MAXLEN] = {0};
void insert(int pos, int val){ //单点插入
for(;pos <= n; pos += lowbit(pos))tree[pos] += val;
```