

There exist a set of teams  $T$  that provide a proof:  $M < \frac{\# \text{ edges in } T + \sum_{x \in T} P_x}{|T|}$

### Project selection:

- We are given a set of project.

- Each project  $x$  has a revenue

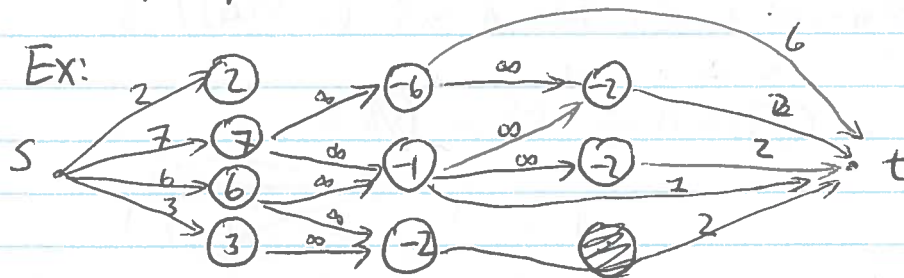
$P_x$ : If  $P_x > 0 \Rightarrow$  it provide a profit

If  $P_x < 0 \Rightarrow$  it provide a loss.

- Some project are prerequisites for other projects.

An edge from  $x \rightarrow y$  means that  $y$  is prerequisite for  $x$  (If we choose  $x$ , we have to choose  $y$ ).

Goal: Select a subset of projects that respects all the prerequisites and maximizes the total revenues.

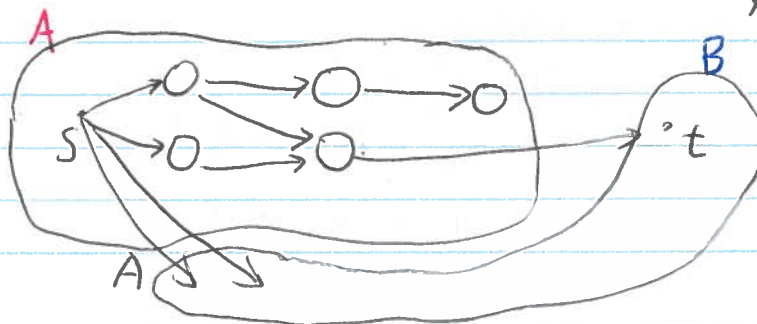


- We assign  $\infty$  capacity to all the edges. This way, if a project  $x$  is in part A of a min-cut and have the prereq  $x \rightarrow y$ , then  $y$  also has to be in A.

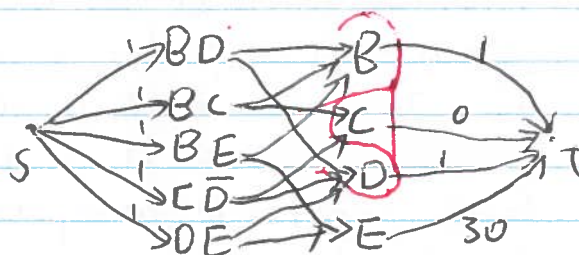
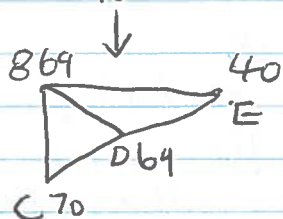
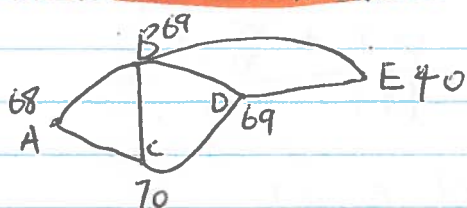
- We add sources, sink  $t$ , edges from  $S$  to  $x$  for projects with  $P_x > 0$  (capacity  $P_x$ )

- edges from  $x$  with  $P_x < 0$  to  $t$  with  $\text{Cap } |P_x|$ .

Let  $(A, B)$  be a min-cut and let  $M = \sum_{x: P_x > 0} P_x$



## Baseball Elimination



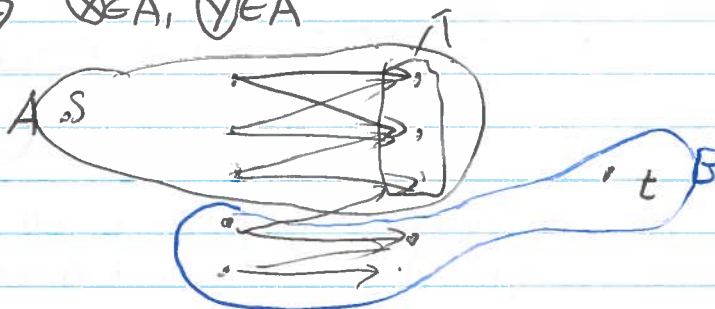
D F-F and if  $\text{max flow} = \# \text{ remaining edges} \Rightarrow$  not eliminated.

What can we say about min-cut?

Min-cut  $\neq \infty$  (Take  $A = \{s\}$   $B = V(G) \setminus \{s\}$ )

- Consider min-cut

If  $(x, y) \in A \Rightarrow x \in A, y \in A$



$$\text{Cap}(A, B) = \sum_{x \in T} (M - P_x) + \text{number of matches } xy \text{ with cut length one of } x \text{ or } y \text{ not in } T.$$

$$= \text{max-flow}$$

If  $\text{max-flow} < \# \text{ of edges} \Rightarrow$  our team is eliminated

$$\text{max flow} = \text{cap}(A, B)$$

$$= M \times |T| - \sum_{x \in T} P_x + K$$

$$\Leftrightarrow M \times |T| - \sum_{x \in T} P_x + K < \# \text{ of edges}$$

$$\Leftrightarrow M \times |T| < \# \text{ edges in } T + \sum_{x \in T} P_x$$

$$\text{Cap}(A, B) = \sum_{\substack{x \in B \\ P_x > 0}} P_x + \sum_{\substack{x \in A \\ P_x < 0}} |P_x| \quad \#$$

$$= \sum_{\substack{x \in A \\ P_x < 0}} -P_x + \sum_{\substack{x \notin A \\ P_x > 0}} P_x$$

$$= \sum_{\substack{x \in A \\ P_x < 0}} -P_x + \left( M - \sum_{\substack{x \in A \\ P_x > 0}} P_x \right)$$

$$= M - \sum_{x \in A} P_x \leftarrow \text{maximized in order to have a min-cut.}$$

We also know that the projects in A respect the prereq condition.

The total profit we can make