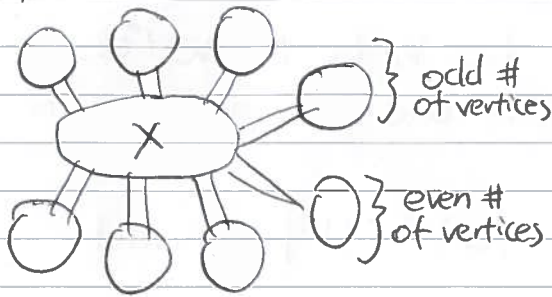


**Thm:** Every bridgeless cubic graph has a perfect matching (Peterson's thm)

**Pf:**

**Thm (Tutte's theorem):**  $G$  has a perfect matching  $\Leftrightarrow \forall S \subseteq V(G)$ , the graph induced by  $V(G) \setminus S$  has a number of components with an odd # of vertices less than or equal to the number of vertices in  $S$ .

**Pf:**



If  $G$  has a perfect matching, then  
 $| \text{odd components of } G \setminus X | \leq |X| \quad \forall X \subseteq V(G)$

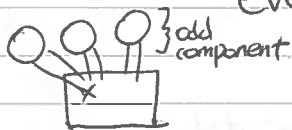
Now we show that if  $| \text{odd comp}(G \setminus X) | \leq |X| \quad \forall X \subseteq V(G)$ , then  $G$  has a perfect matching.  
 We apply induction on  $|V(G)|$

**Claim 1:**  $|V(G)|$  is even, as  $|V(G)| = |X| + \text{odd}(G \setminus X)$ , take  $|X| = 0$  gives  $|V(G)| \equiv 0 \pmod{2}$

We say  $X$  is **critical** if  $\text{odd}(G \setminus X) = |X|$ . By Claim 1, if  $\text{odd}(G \setminus X) \geq |X| - 1$ , then  $X$  is critical.

**Claim :** Since  $\emptyset$  is a critical set,  $G$  has at least 1 critical set.

Let  $X$  be the maximal critical set,  $G \setminus X$  has no even components. Otherwise, let  $V \in V(C)$  and  $X' = X \cup \{V\}$

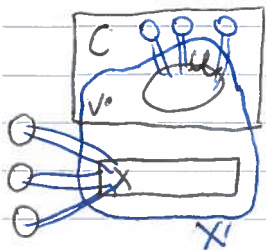


$\text{odd}(G \setminus X') \geq \text{odd}(G \setminus X) = |X| = |X'| - 1$ .  
 So  $X'$  is critical.  $\nRightarrow$

even comp

**Claim :** If  $C$  is an odd component of  $G \setminus X$  and  $V \in V(C)$ , then  $C \setminus V$  has a perfect matching.

We can show  $\forall$  subset  $U$  of  $C \setminus V$ ,  $\text{odd}(C \setminus V \setminus U) \leq |U|$  for all  $U \subseteq V(C \setminus V)$ . then apply induction.

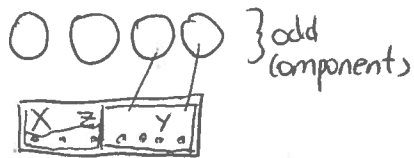


Let  $X' = X \cup \{V\} \cup U$

$$\begin{aligned} \text{odd}(G \setminus X') &= \text{odd}(G \setminus X) - 1 + \text{odd}(C \setminus V \setminus U) \\ &= |X| - 1 + \text{odd}(C \setminus V \setminus U) \end{aligned}$$

$\leq |X'| - 2$  as  $X'$  is not critical.

$$= (|X| + 1 + |U|) - 2 = |X| + |U| - 1 \Rightarrow \text{odd}(C \setminus V \setminus U) \leq |U|$$



We want a matching covering  $X$  and hitting every odd component of  $G \setminus X$  exactly once.

Suppose it doesn't exist, by Hall's theorem, there exists  $Y \subseteq X$  such that  $\# \text{ odd components adjacent to } Y < |Y|$ .

Let  $Z = X - Y$ , then  $\text{odd}(G|Z) > |Z|$   $\nexists$  So  $\square$

Pf (Peterson's theorem):

We will prove that  $\forall U \subseteq V(G)$ ,  $\text{odd}(G-U) \leq |U|$ , then the theorem will follow from Tutte's theorem.

Let  $G_i$  be an odd component of  $G \setminus U$ . Let  $V_i$  be the vertices in  $G_i$  and  $m_i$  be the set of edges with one end in  $U$  and the other in  $V_i$ .

then  $\sum_{v \in V_i} \deg(v) = 2|E_i| + m_i$ ,  $\sum_{v \in V_i} \deg(v) = 3|V_i|$ , is odd

hence  $m_i$  is odd.

As  $G$  is bridgeless,  $m_i \geq 3$ .

Let  $m$  be the number of edges with one end in  $U$  and the other in  $G \setminus U$ . then All odd component contribute to at least 3 vertices in  $m$ . Hence, the number of component is at most  $m/3$ .

In worse case, all edges with one end in  $U$  contribute to  $m$

So  $m \leq 3|U| \Rightarrow \text{odd}(U) \leq m/3 \leq \frac{3|U|}{3} = |U|$

So by Tutte thm,  $G$  has a perfect matching  $\square$

Def:  $m(G)$  is used to denote the number of perfect matching in a graph  $G$ .

Theorem  $\square$  For any  $\alpha > 0$  there exists a constant  $\beta > 0$  such that every  $n$ -vertex cubic bridgeless graph has at least  $\alpha n^\beta$  perfect matchings.