

CS61B – Algorithms

Anmol Parande

Spring 2019 - Professor Josh Hug

1 Tree Traversals

1.1 Preorder Traversal

In a preorder traversal a node is visited before any of its children are visited.

```
public void preorder(TreeNode node) {  
    if (node == null) {  
        return;  
    }  
    visit(node);  
    for (TreeNode child : node.children) {  
        preorder(child);  
    }  
}
```

1.2 In Order Traversal

For a Binary Tree, an in order traversal is simply the elements when read from left to right.

```
public void inorder(BNode node) {  
    if (node == null) {  
        return;  
    }  
    inorder(node.left);  
    visit(node);  
    inorder(node.right);  
}
```

1.3 Post Order Traversal

In a postorder traversal, a node is visited after its children are visited.

```

public void postorder(TreeNode node) {
    if (node == null) {
        return;
    }
    for (TreeNode child : node.children) {
        postorder(child);
    }
    visit(child);
}

```

1.4 Level Order Traversal

In a level order traversal, the nodes of each level of the tree are visited in order. It is done through a Breadth-First Search of the tree (see below)

2 Breath-First Search

The idea of breadth first search is to visit all children before visiting the children's children. Its runtime is $O(|V| + |E|)$ where V is the set of vertices and E is the set of edges.

```

repeat until queue is empty:
    remove vertex v from queue
    for each unmarked vertex n adjacent to v
        mark n
        keep track of some statistic
        add n to end the queue

```

3 Depth-First Search

The idea of depth-first search is to visit all the children of a child before visiting the nodes other children. The preorder is the order in which DFS calls are made. The postorder is the order of DFS returns. It's runtime is $O(|V| + |E|)$.

3.1 Pre-Order

```

add start node to the stack
while stack is not empty:
    v = stack.pop
    mark(v)
    for vertex n adjacent to v:
        mark(n)
        stack.push(n)

```

3.2 PostOrder

```
DFS(Vertex v) {
    mark v visited
    set color of v to gray
    for each successor v' of v {
        if v' not yet visited {
            DFS(v')
        }
    }
    doAction(v)
}
```

4 Dijkstra's Algorithm

Dijkstra's algorithm helps build the shortest paths tree starting from a particular source node.

```
add all vertices to priority queue with infinity priority
set the source's priority to 0
while priority queue is not empty
    remove the smallest vertex v
    mark v
    relax edges from v
```

The relaxation procedure is as follows

```
if w is visited
    return
w = weight(parent) + edge.weight
if w < v.priority
    priority = w
```

Dijkstra's algorithm runs in $O(|E|\log|V|)$

5 A^* Search

A^* search is like Dijkstra's in that it finds the shortest path between two nodes. However, it is different because the priority of each vertex is the distance to that vertex plus a heuristic value

A^* requires a goal vertex to run. Its runtime is $O(|E|\log|V|)$

6 MSTs

A Minimum Spanning Tree (MST) is a tree connecting all vertices in a graph which has the minimum total weight (the sum of its edges is minimized).

6.1 Prim's Algorithm

```
add source vertex to PQ
set all other vertex priorities to infinity
for edge out of V:
    if destination is unmarked:
        relax(V) using the distance to the tree
```

Notice that this is basically Dijkstra's algorithm but we are using distance to tree rather than distance to source. The runtime is $O(|E|\log|V|)$.

6.2 Kruskal's Algorithm

```
Make a priority queue of all edges
Create a Disjoint Set of all the vertices
while !pq.empty() && mst.size() < V - 1:
    e = pq.pop
    if source not connected to target in DisjointSet
        connect(source, target)
    add e to MST
```

The runtime is also $O(|E|\log|V|)$.

7 Sorting

In sorting, an **Inversion** is a pair of elements which is out of order. The goal of sorting is to reduce the number of inversions to 0.

A sort is **stable** if the ordering of equal elements in the sorted array is the same as their ordering in the original array.

All comparison sorts are bounded below by $\Omega(N\log N)$, but sorts which do not rely on comparisons can do better.

7.1 Selection Sort

Selection sort maintains a sorted portion of the array by searching through the array for the minimum element and swapping it to the front.

```
def selectionSort(arr):
    for (i=0; i < arr.length; i++):
        minIndex = i
        for (j=i+1; j < arr.leng; j++):
            if (arr[j] < arr[i]):
                minIndex = j;
        temp = arr[i];
        arr[i] = arr[minIndex];
        arr[minIndex] = temp;
```

Time Complexity: $\Theta(N^2)$

Space Complexity: $\Theta(1)$

Stable?: Yes

Identifying Characteristics: In intermediate steps, the beginning portion of the array will always be in sorted order.

7.2 Heap Sort

The essential idea of heap sort is to put the entire array into a max-heap, and then build up the sorted array backwards by popping of the top element.

If we do this naively (i.e by copying the original array into a new heap and then building a new array, the space complexity is $\Theta(N)$). However, it can be done in place using array heapification.

Best Case Time Complexity: $\Theta(N)$ (All elements are equal)

Worst Case Time Complexity: $\Theta(N \log N)$

Space Complexity: $\Theta(1)$

Stable?: No

Identifying Characteristics: The maximum values are sent to the back of the array during intermediate steps.

7.2.1 Bottom-Up Heapification

Bottom-Up Heapification turns an array into a Max-Heap in place. Simply sink the nodes in reverse level order and you will get a Max Heap. Once this is done, heap sort can proceed mostly as usual.

7.3 Merge Sort

Mergesort is a recursive sort which divides the array in two, sorts both halves, and then merges the results.

```
def mergeSort(arr):
    if arr.length == 1:
        return arr
    mid = arr.length // 2
    leftHalf = mergeSort(arr[:mid])
    rightHalf = mergeSort(arr[mid:])
    i, j, k = 0, 0, 0
    sortedArr = new T[arr.length]
    while (i < leftHalf.length && j < rightHalf.length):
        if (leftHalf[i] <= rightHalf[j]):
            sortedArr[k] = leftHalf[i]
            i++
        else:
            sortedArr[k] = rightHalf[j]
            j++
    while i < leftHalf.length:
        sortedArr[k] = leftHalf[i]
        i++
    while j < rightHalf.length:
        sortedArr[k] = rightHalf[j]
        j++
    return sortedArr
```

```

        k++
    }

    while (i < leftHalf.length):
        sortedArr[k] = leftHalf[i]
        i++
        k++

    while (j < rightHalf.length):
        sortedArr[k] = rightHalf[j]
        j++
        k++

    return sortedArr

```

Time Complexity: $\Theta(N \log N)$

Space Complexity: $\Theta(N)$

Stable?: Yes

Identifying Characteristics: Elements from different halves of the array do not cross during intermediate steps of the algorithm.

7.4 Insertion Sort

Insertion sort maintains a sorted portion in the beginning of the array. Out of place elements are swapped backwards until they arrive in their correct position. Insertion sort performs very well on arrays with few inversions and with few elements.

```

def insertionSort(arr):
    n = arr.length;
    for (int i = 1; i < n; i++) {
        key = arr[i]
        j = i - 1
        while (j >= 0 && arr[j] > key):
            arr[j + 1] = arr[j]
            j--
        arr[j + 1] = key;
    }

```

Best Case Time Complexity: $\Theta(N)$

Worst Case Time Complexity: $\Theta(N^2)$

Space Complexity: $\Theta(1)$

Stable?: Yes

Identifying Characteristics: A distinct sorted portion exists at the beginning of the array.

7.5 QuickSort

QuickSort works based on the idea of partitioning. To partition an array based on a particular pivot element, reorder the array so everything left of the pivot is

smaller and everything to the right is larger or equal. Essentially, if we partition the array, and then QuickSort both halves, we will arrive at a sorted array.

```
def quicksort(arr):
    pivotIndex = choosePivot(arr)
    arr = partition(arr, pivotIndex)
    leftHalf = quicksort(arr[:pivotIndex])
    rightHalf = quicksort(arr[pivotIndex+1:])
    return leftHalf + rightHalf
```

7.5.1 3 Scan Partitioning

```
pivot = arr[0]
partitionedArr = new array
copy elements < pivot into partitionedArr
copy elements = pivot into partitionedArr
copy elements > pivot into partitionedArr
return partitionedArr
```

7.5.2 Hoare Partitioning

```
left, right = 0, arr.length - 1
while (left < right):
    while arr[left] < pivot:
        left++
    while arr[right] < pivot:
        right--
    swap(arr[left], arr[right])
    left++
    right--
```

Best Case Time Complexity: $\Theta(N \log N)$

Worst Case Time Complexity: $\Theta(N^2)$

Space Complexity: $\Theta(1)$

Stable?: No

Identifying Characteristics: During intermediary steps, everything left of where the first element in the unsorted array ends up is less than the pivot and everything to the right is larger or equal.

Note: The properties of quicksort depends on the partitioning scheme, pivot strategy, and whether or not randomization was used.

7.6 Counting Sort

Counting sort is a non-comparison sort which requires a finitely sized alphabet. It merely counts the occurrence of each item and then copies them into the proper locations in the sorted array.

```
def countingSort(arr):
```

```

max = maxElement(arr)
int n= arr.length
int B[]=new int [n]
int C[]=new int [max+1]

for (int i=0; i <=max; i++)
    C[i] = 0

for (int j=0; j<n; j++):
    C[arr[j]] = C[arr[j]] + 1

for (int i=1; i<max+1; i++):
    C[i] = C[i] + C[i-1]

for (int j=n-1; j>=0; j--):
    B[C[arr[j]]-1] = arr[j]
    C[arr[j]] = C[arr[j]] - 1

```

If N is the number of elements in the array and R is the size of the alphabet

Time Complexity: $\Theta(N + R)$

Space Complexity: $\Theta(N + R)$

Stable?: Yes

Note: As long as $N \gg R$ then counting sort will be very fast.

7.7 LSD Radix Sort

Least Significant Digit Radix sort uses counting sort to sort elements comprised of a sequence of a finite alphabet (like Strings or Integers). The idea is to apply counting sort for each digit starting with the least significant digit and working leftwards.

```

for (i=0; i < longestElement.length; i++):
    counting sort based on ith digit

```

If N is the number of elements, W is the length of the longest word, and R is the size of the alphabet. **Time Complexity:** $\Theta(WN + WR)$

Space Complexity: $\Theta(N + R)$

Stable?: Yes

7.8 MSD Radix Sort

Most Significant Radix Sort is like LSD Radix Sort except it starts with the most significant digit and works its way rightwards. However, if we proceed the same way as we did LSD radix sort, our sort will not be accurate. For MSD radix sort, it is necessary to only apply counting sort to each subproblem after we sort each subsequent digit.


```

subproblems = [arr]
for (i=longestElement.length-1; i >= 0; i--):
    for each problem in subproblems:
        counting sort problem on ith digit
        add new subproblem to array
combine sorted subproblems to get sorted array

```

Best Case Time Complexity: $\Theta(N + R)$

Worst Case Time Complexity: $\Theta(WN + WR)$

Space Complexity: $\Theta(N + WR)$

Stable?: Yes

Notes: The best case time complexity comes about when each element is distinguished only by the most significant digit.

7.9 Comparing Sorts

| Sort | Space Complexity | Best Case Time Complexity | Worst Case Time Complexity | Stable |
|-----------|------------------|---------------------------|----------------------------|--------|
| Insertion | $\Theta(1)$ | $\Theta(N)$ | $\Theta(N^2)$ | Yes |
| Selection | $\Theta(1)$ | $\Theta(N^2)$ | $\Theta(N^2)$ | No |
| Heap | $\Theta(1)$ | $\Theta(N)$ | $\Theta(N \log N)$ | No |
| Merge | $\Theta(N)$ | $\Theta(N \log N)$ | $\Theta(N \log N)$ | Yes |
| Quick | $\Theta(\log N)$ | $\Theta(N \log N)$ | $\Theta(N^2)$ | No |
| Counting | $\Theta(N + R)$ | $\Theta(N + R)$ | $\Theta(N + R)$ | Yes |
| LSD | $\Theta(N + R)$ | $\Theta(WN + WR)$ | $\Theta(WN + WR)$ | Yes |
| MSD | $\Theta(N + R)$ | $\Theta(N + R)$ | $\Theta(WN + WR)$ | Yes |

8 Topological Sort

The problem of a topological sort is given a directed acyclic graph (DAG), find an ordering of vertices that satisfies the directionality of the edges in the graph. I.e, if the edges of a graph indicate that a task must come after a specific task, then find a valid ordering of tasks.

```

procedure:
    while not all visited:
        DFS(v) for arbitrary v
        if not all visited:
            repeat on unmarked vertex
        record DFS post order
    reverse DFS post-order

```

Time Complexity: $O(V + E)$ **Space Complexity:** $\Theta(V)$