

Writing Queries involving more than one Relation

CSC365
Spring 2019

Tuple Variables (Renaming Relations ρ)

- Sometimes, you want to alias the name of a relation to a short string

```
SELECT m.mid, title, year, genre
```

```
FROM Movies as m, Genre as g
```

```
WHERE m.mid=g.mid;
```

- This means that each tuple in Movies is aliased as m, each tuple in Genre is aliased as g.
- A tuple variable -> A variable for each tuple in a relation.

Interpreting Multirelation Queries as Nested Loops

```
SELECT m.title, g.genre  
FROM Movies m, Genre g, StarsIn s  
WHERE m.mid = g.mid  
AND g.genre = 'Sci-Fi'  
AND m.mid = s.mid  
AND s.sname = 'Will Smith';
```

```
For each tuple m in Movies:  
  For each tuple g in Genre:  
    For each tuple s in StarsIn:  
      If m.mid == g.mid  
        AND g.genre == 'Sci-Fi'  
        AND m.mid == s.mid  
        AND s.sname == 'Will Smith':  
          Produce m.title, g.genre
```

Converting Multirelation Queries to Relational Algebra

SELECT m.title, g.genre

FROM Movies m, Genre g, StarsIn s

WHERE m.mid = g.mid

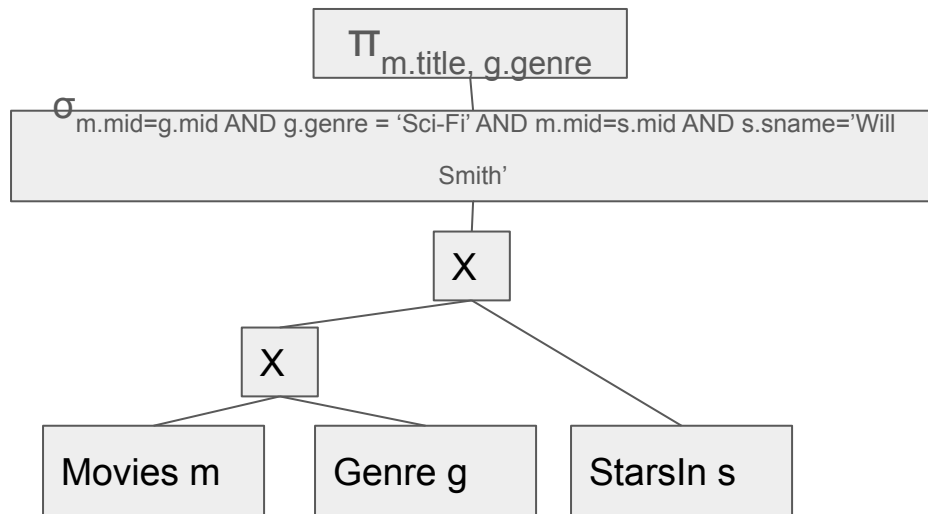
AND g.genre = 'Sci-Fi'

AND m.mid = s.mid

AND s.sname = 'Will Smith';

$\Pi_{m.title, g.genre}$

$(\sigma_{m.mid=g.mid \text{ AND } g.genre = 'Sci-Fi' \text{ AND } m.mid=s.mid \text{ AND } s.sname='Will Smith'}$
 $(\rho_m(\text{Movies}) \times \rho_g(\text{Genre}) \times \rho_s(\text{StarsIn})))$



Subqueries

A query constructed out of the results of other queries - Nesting queries

1. Subqueries can return a single constant, which can be used in **WHERE** clause.
2. Subqueries can return relations that can be used in various ways in **WHERE** clauses.
3. Subqueries can appear in **FROM** clauses, followed by a tuple variable that represents the tuples in the result of the subquery.

Subqueries that produce Scalar Values

```
SELECT MAX(imdb) FROM Movies;
```

```
SELECT * FROM Movies WHERE imdb = 9;
```

MAX(imdb)
9



```
SELECT *
```

```
FROM Movies
```

```
WHERE imdb = (SELECT MAX(imdb) FROM Movies);
```

Conditions Involving Relations

- R - A one column relation produced by a subquery.
- EXISTS R
 - TRUE if and only if R is not empty.
- S IN R
 - TRUE if and only if S is equal to one of the values in R.
- S NOT IN R
 - TRUE if and only if S is not equal to any values in R.
- S > ALL R
 - TRUE if and only if S is greater than every value in R.
- S > ANY R
 - TRUE if and only if S is greater than at least one of the values in R.

S IN / NOT IN R

- $s \text{ IN } R$ is true if and only if s is equal to one of the values in R .
 - s may be a single value or single column name, in which case R must be a single-column relation
 - If s has more than one element, the number of elements must match the number of columns in R
- $s \text{ NOT IN } R$ is true if any only if s is equal to *no value* in R

IN / NOT IN Example

#Print the team which did not win a single game.

```
SELECT t.id, t.name
```

```
FROM Team t
```

```
WHERE t.id not in (
```

```
SELECT t.id From Game g JOIN Team t ON t.id = g.home_team_id and  
g.score_home > score_away or t.id = g.away_team_id and g.score_away >  
g.score_home
```

```
);
```

IN / NOT IN Example

-- Students without a minor who are in a college other than OCOB

```
SELECT *
```

```
FROM Student
```

```
WHERE MinorCode IS NULL
```

```
AND MajorCode NOT IN (
```

```
  SELECT Code
```

```
  FROM Department
```

```
  WHERE College = 'OCOB'
```

```
)
```

How does this query change if we use JOIN rather than a subquery?

ALL / ANY

$s < \text{comparator} > \text{ALL } R$

$s < \text{comparator} > \text{ANY } R$

(comparator may be any of our six basic comparisons: $>$, $<$, $>=$, $<=$, $=$, $<>$)

$s > \text{ALL } R$ is true if and only if s is greater than every value in *unary* relation R . s

$<> \text{ALL } R$ is the same as $s \text{ NOT IN } R$

$s > \text{ANY } R$ is true if and only if s is greater than at least one value in *unary* relation

R . $s = \text{ANY } R$ is the same as $s \text{ IN } R$

NOT ALL / ANY

ALL / ANY may be negated just like any other boolean expression. For example:

NOT $s \geq \text{ALL } R$ is true if and only if s is not the maximum value in R

NOT $s > \text{ANY } R$ is true if any only if s is the minimum value in R

ALL Example

```
-- Find departments established before every department in Engineering
SELECT *
FROM Department
WHERE DateEstablished < ALL (
    SELECT DateEstablished
    FROM Department WHERE College = 'CENG'
)
```

ANY Example

-- Find the earliest-established department

SELECT *

FROM Department

WHERE NOT DateEstablished > ANY (

 SELECT DateEstablished

 FROM Department

)

(NOT) EXISTS R

- EXISTS R is true if and only if R is not empty
- NOT EXISTS R is true if and only if R is empty

(NOT) EXISTS Example

-- Find students who are the only students in their major

SELECT *

FROM Student o

WHERE MajorCode IS NOT NULL

AND NOT EXISTS (

 SELECT StudentID

 FROM Student i

 WHERE o.MajorCode = i.MajorCode AND o.StudentID <> i.StudentID

)

-- Find students who share a major with at least one other student

... (change NOT EXISTS to EXISTS) ...

EXISTS / NOT EXISTS Transformation

Subqueries in the `WHERE` clause using `IN` can be transformed into equivalent expressions using `EXISTS` & `NOT EXISTS`.

```
SELECT *  
FROM Student  
WHERE MajorCode IN (  
    SELECT Code  
    FROM Department  
    WHERE College = 'OCOB'  
)
```

```
SELECT *  
FROM Student  
WHERE EXISTS (  
    SELECT Code  
    FROM Department  
    WHERE College = 'OCOB'  
    AND Code = Student.MajorCode  
)
```

EXISTS / NOT EXISTS Transformation

List students who were enrolled on/after all departments in CENG were established.

```
SELECT *  
FROM Student  
WHERE DateEnrolled >= ALL (  
    SELECT DateEstablished  
    FROM Department  
    WHERE College = 'CENG'  
)
```

```
SELECT *  
FROM Student  
WHERE NOT EXISTS (  
    SELECT DateEstablished  
    FROM Department  
    WHERE College = 'CENG'  
    AND Student.DateEnrolled < DateEstablished  
)
```



EXISTS / NOT EXISTS Transformation

Consider relations R(A, B) and S(C)

```
SELECT C
FROM S
WHERE C IN (
    SELECT SUM(B)
    FROM R
    GROUP BY A
)
```

```
SELECT C
FROM S
WHERE EXISTS (
    SELECT SUM(B) FROM R
    GROUP BY A
    HAVING SUM(B) = S.C
)
```

Conditions Involving Tuples

We've seen how IN, ANY, and ALL work with one-column relations. It is also possible to perform tuple-based comparisons using subqueries, as long as the *degree* matches on both sides of the comparison.

-- Students who share the same major & minor as another student

SELECT *

FROM Student SO

WHERE (MajorCode, MinorCode) IN (

 SELECT MajorCode, MinorCode FROM Student SI

 WHERE SI.StudentID <> SO.StudentID

)

-- Anybody missing?

Nested Queries in the FROM clause

Since a SELECT statement returns a relational table, we can use a nested SELECT statement in the FROM clause of a query.

```
SELECT <column list>  
FROM (SELECT query) [AS] <alias>  
[WHERE <condition> ]  
[GROUP BY <attribute list>  
[HAVING <group condition>]
```

Nested Queries in the FROM clause

Two requirements to consider when using a nested `SELECT` statement in the `FROM` clause of a query:

1. Nested `SELECT` must be enclosed in **parentheses** and *must have an **alias***
2. All **computed columns** (aggregates, scalar functions, `CASE`, etc.) *must have **aliases***

Subqueries in the FROM clause

-- Print duplicate records in Movies table

```
SELECT m1.title, m1.year, m1.gross, m1.imdb
```

```
FROM Movies m1,
```

```
(SELECT title, year, count(*) as num FROM Movies GROUP BY title, year) m2
```

```
WHERE m1.title = m2.title AND m1.year = m2.year AND m2.num > 1
```

```
ORDER BY m1.title, m1.year;
```

Joins with Subquery Results

SELECT a1, a2, ..., an

FROM R

JOIN S ON jc

WHERE c;

R or S can be substituted with a relation resulting from a subquery.

#List all actors who co-starred with Tom Hanks

Select distinct s.sname

From StarsIn s

JOIN (select * from StarsIn where sname='Tom Hanks') as kb

ON kb.mid=s.mid AND kb.sname != s.sname;

Nested Queries in the FROM clause

#List all actors who co-starred with Harrison Ford

```
select distinct s.sname
```

```
From StarsIn s JOIN (
```

```
select * from StarsIn where sname='Harrison Ford'
```

```
) as hf
```

```
ON hf.mid=s.mid AND hf.sname != s.sname;
```

Nested Queries in the FROM clause

#List Movies both Harrison Ford and Sylvester Stallone appeared

```
SELECT m.*
```

```
FROM Movies m JOIN (
```

```
    select DISTINCT s.mid FROM StarsIn s JOIN (
```

```
        select * from StarsIn where sname='Harrison Ford'
```

```
    ) as hf ON hf.mid=s.mid WHERE s.sname='Sylvester Stallone'
```

```
) as hs ON m.mid=hs.mid;
```

Nesting Subqueries

Subqueries can be deeply nested, up to practical limits imposed by RDBMS implementations. Inner queries can reference outer columns, but cannot reference siblings.

```
-- Non-CENG departments with students who enrolled
-- prior to the date the department was established
SELECT *
FROM Department d1
WHERE EXISTS (
    SELECT StudentID FROM Student s
    WHERE s.MajorCode = d1.Code AND DateEnrolled < (
        SELECT DateEstablished FROM Department d2
        WHERE d2.Code = s.MajorCode
    ) AND s.MajorCode NOT IN (
        SELECT Code FROM Department d3
        WHERE d3.College = 'CENG'
    )
)
-- How would we list the corresponding students?
```

Nesting Subqueries

Subqueries can be deeply nested, up to practical limits imposed by RDBMS implementations. Inner queries can reference outer columns, but cannot reference siblings.

```
select name from (  
  select name, count(name) as countNum  
  from (select name  
        from (select * from  
              (select player_id as id  
               from Event  
               where type = "Goals scored" and  
                    count >= 3) as players  
              natural join Player) as playersAndTeam, Team  
        where playersAndTeam.team_id = Team.id) as  
        teamName  
  group by name having countNum >= 1) as NameNums;
```

Correlated Subqueries

The subquery examples seen thus far have been uncorrelated: subqueries that can be executed once for the entire outer SELECT.

It is also possible to refer to values in the outer query from within a subquery, causing the subquery to be re-evaluated multiple times during execution of the outer query.

This second form of subquery is called a **correlated subquery**

Correlated Subqueries Example

-- Check if there are movies with the same title but made earlier

```
SELECT m1.title, m1.year, m1.gross, m1.imdb
```

```
FROM Movies m1
```

```
WHERE m1.year < ANY (SELECT year FROM Movies WHERE title = m1.title)
```

```
ORDER BY m1.title, m1.year;
```

Correlated Subqueries Example

```
SELECT name, (SELECT COUNT(*) FROM StarsIn WHERE sname=name) as  
num_movies  
  
FROM Stars;
```

Can we convert this correlated
subquery to a regular join?

```
SELECT name, num_movies
```

```
FROM Stars s
```

```
JOIN (SELECT sname, COUNT(*) as num_movies FROM StarsIn group by  
sname) as n ON n.sname=name;
```

Subquery Recap

- Considerable flexibility, deep nesting allowed
- Can appear throughout SELECT statement (with some exceptions)
- Subquery variations:
 - Uncorrelated
 - Derived table in FROM clause
 - Correlated
 - No correlation permitted for subqueries that appear FROM clause
- Always consider performance (more on this later)

Set Operations in SQL

- \cup : UNION
- \cap : INTERSECT
- $-$: EXCEPT

UNION

(SELECT title, year, gross, imdb FROM Movies where imdb > 7.0)

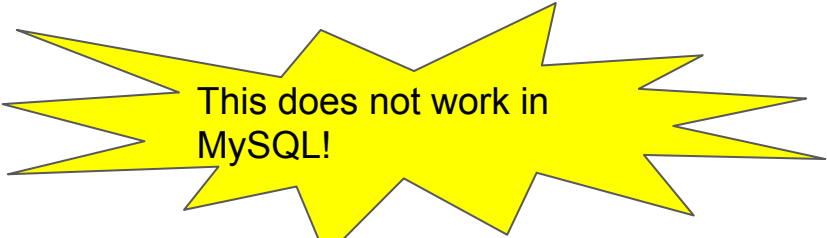
UNION

(SELECT title, year, gross, imdb FROM Movies where gross > 10000000);

INTERSECT

Unfortunately, MySQL does not support the **INTERSECT** operator.

However, you can simulate **INTERSECT** operator.



This does not work in
MySQL!

```
(SELECT title, year, gross, imdb FROM Movies where imdb > 7.0)
```

```
INTERSECT
```

```
(SELECT title, year, gross, imdb FROM Movies where gross > 10000000);
```

Simulating INTERSECT

simulate **INTERSECT** operator using DISTINCT operator and INNER JOIN

(SELECT title, year, gross, imdb FROM Movies where imdb > 7.0)

INTERSECT

(SELECT title, year, gross, imdb FROM Movies where gross > 10000000);



SELECT **DISTINCT** m1.title, m1.year, m1.gross, m1.imdb FROM Movies m1
JOIN

(SELECT title, year, gross, imdb FROM Movies where gross > 10000000) as m2

USING (title, year)

WHERE m1.imdb > 7.0;

Simulating INTERSECT

Simulate MySQL INTERSECT operator using IN operator and subquery

```
SELECT title, year, gross, imdb FROM Movies  
WHERE imdb > 7.0 AND mid IN  
(SELECT mid FROM Movies where gross > 10000000);
```

EXCEPT

Unfortunately, MySQL does not support the **EXCEPT** operator.

However, you can simulate **EXCEPT** operator using NOT IN operator, LEFT JOIN, and NOT EXISTS operator.

Simulating EXCEPT

```
SELECT title, year, gross, imdb FROM Movies
```

```
WHERE imdb > 7.0 AND mid NOT IN
```

```
(SELECT mid FROM Movies where gross > 100000000);
```

```
SELECT title, year, gross, imdb FROM Movies
```

```
WHERE imdb > 7.0 AND (title,year) NOT IN
```

```
(SELECT title,year FROM Movies where gross > 100000000);
```

Simulating EXCEPT

```
SELECT m1.title, m1.year, m1.gross, m1.imdb FROM Movies m1
```

```
LEFT JOIN
```

```
(SELECT mid FROM Movies where gross > 10000000) as m2 USING (mid)
```

```
WHERE imdb > 7.0 AND m2.mid is NULL;
```


Simulating EXCEPT

```
SELECT title, year, gross, imdb FROM Movies m
```

```
WHERE imdb > 7.0 AND NOT EXISTS
```

```
(SELECT title,year FROM Movies where mid=m.mid AND gross > 10000000);
```

Duplicates in Unions, Intersections, and Differences

- Basically, set operations produce SET of tuples, not BAG of tuples.
- This means that duplicates will be automatically eliminated.
- This is True with UNION in MySQL.
- When simulating INTERSECT and EXCEPT (MINUS) in MySQL, you have to eliminate duplicate tuples by yourself.

UNION

(SELECT title, year, gross, imdb FROM Movies where imdb > 7.0)

UNION ALL

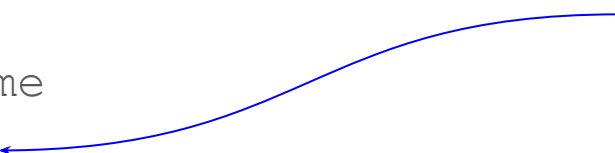
(SELECT title, year, gross, imdb FROM Movies where gross > 10000000);

Control Row Count Returned by SELECT

ANSI SQL does not define a standard way to limit the number of rows returned by a `SELECT` query.

Each RDBMS has its own syntax. In MySQL, we use `LIMIT`:

```
SELECT *  
FROM Student  
ORDER BY LastName  
LIMIT 2
```



Sorting (`ORDER BY`) is performed **before** `LIMIT`, so we see the first two students, based on A-Z ordering of last name.

MySQL LIMIT Variations

Single-argument LIMIT returns up to the specified number of rows:

`LIMIT 5` -- top 5 rows in result set

Two-argument LIMIT returns up to the requested number of rows, after a given offset (zero indexed):

`LIMIT 5,10` -- rows 6-15

Row Limiting in other RDBMSs

RDBMS	Syntax
MySQL & PostgreSQL	<code>SELECT * FROM table LIMIT 10</code>
Microsoft SQL Server & Access	<code>SELECT TOP 10 * FROM table</code>
Oracle	<code>SELECT * FROM (SELECT * FROM table) WHERE rownum <= 10</code>
IBM DB2	<code>SELECT * FROM table FETCH FIRST 10 ROWS ONLY</code>
Informix	<code>SELECT FIRST 10 * FROM table</code>

Using LIMIT in Subqueries

```
SELECT Code, Name, DateEstablished,  
       DATEDIFF((SELECT DateEstablished  
                 FROM Department  
                 ORDER BY DateEstablished DESC  
                 LIMIT 1),  
               DateEstablished) / 365 AS YearsOlderThanNewestDept  
FROM Department
```

Alternative way to write
this query, without LIMIT?