

Notes on Mathematics

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1 Calculus

1.1 Differentiation and Integration

Lemma 1.1 (Simple Calculations).

1. For $1 = xx^{-1}$ the product rule yields $0 = x^{-1} + x(x^{-1})'$. Hence

$$\frac{d}{dx}x^{-1} = -\frac{1}{x^2}$$

2. Similarly $x = \sqrt{x^2}$ and $1 = 2\sqrt{x}\sqrt{x}'$ and so

$$\frac{d}{dx}\sqrt{x} = \frac{1}{2\sqrt{x}}$$

3. It is

$$\frac{d}{dx}x^n = nx^{n-1}$$

since via induction the product rule yields

$$\frac{d}{dx}x^n = \frac{d}{dx}xx^{n-1} = x^{n-1} + \frac{d}{dx}x^{n-1} = x^{n-1} + (n-1)x^{n-1} = nx^{n-1}$$

4. Again, applying the product rule gives

$$\left(\frac{1}{g}\right)' = \left(\frac{1}{x} \circ g\right)' = -\frac{g'}{g^2}$$

and the quotient rule

$$\left(\frac{f}{g}\right)' = \frac{f'}{g} + f\left(\frac{1}{g}\right)' = \frac{f'}{g} - \frac{fg'}{g^2} = \frac{gf' - fg'}{g^2}$$

5. Also $x = f \circ f^{-1}$ and $1 = (f^{-1})'f' \circ f^{-1}$. Thus

$$(f^{-1})' = \frac{1}{f' \circ f^{-1}}$$

where defined. Especially for $x \neq 0$

$$\log'(x) = \frac{1}{\exp'(\log(x))} = \frac{1}{x}$$

6. $(1-q)(1+q+q^2+\dots+q^n) = 1-q+q-q^2+q^2-q^3+\dots+q^{n+1}$ gives

$$\sum_{k=0}^n q^k = \frac{1-q^{n+1}}{1-q} \text{ and } \sum_{k=m}^n q^k = \frac{q^m - q^{n+1}}{1-q}$$

Lemma 1.2 (Exponential Function).

1. It is

$$\exp(x + y) = \exp(x) \exp(y)$$

Hence

$$\begin{aligned}\exp(0) &= 1 \\ \exp(-x) &= \exp(x)^{-1} \\ \exp(nx) &= \exp(x)^n\end{aligned}$$

2. For the derivative

$$\exp'(x) = \sum_{k=0}^{\infty} \frac{1}{k!} (x^k)' = \sum_{k=0}^{\infty} \frac{1}{k!} k x^{k-1} = \sum_{k=1}^{\infty} \frac{1}{(k-1)!} x^{k-1} = \exp(x)$$

Lemma 1.3 (Sinus and Cosinus).

1. Sinus and Cosinus power series

$$\begin{aligned}\cos(x) &= \sum_{k=0}^{\infty} \frac{(-1)^k}{2k!} x^{2k} \\ \sin(x) &= \sum_{k=0}^{\infty} \frac{(-1)^k}{(2k+1)!} x^{2k+1}\end{aligned}$$

2. Symmetry

$$\begin{aligned}\cos(-x) &= \sum_{k=0}^{\infty} \frac{(-1)^k}{2k!} (-x)^{2k} = \cos(x) \\ \sin(x) &= \sum_{k=0}^{\infty} \frac{(-1)^k}{(2k+1)!} (-x)^{2k+1} = -\sin(x)\end{aligned}$$

3. Derivatives

$$\begin{aligned}\cos'(x) &= \sum_{k=1}^{\infty} \frac{(-1)^k}{(2k-1)!} x^{2k-1} = \sum_{k=0}^{\infty} \frac{(-1)^{k+1}}{(2k+1)!} x^{2k+1} = -\sin(x) \\ \sin'(x) &= \sum_{k=0}^{\infty} \frac{(-1)^k}{2k!} x^{2k} = \cos(x)\end{aligned}$$

Theorem 1.4 (Fermat Stationary Point). Let $\Omega \subseteq \mathbb{R}$ be open and $f \in C^1(\Omega)$. If $x^* \in \Omega$ is local extremum then $f'(x^*) = 0$.

Proof. Assume x^* is the minimum of f in Ω and let $f(x^*) > 0$. Since $f \in C^1(\Omega)$ there exist $\varepsilon, \delta > 0$, so that for $|h| \leq \varepsilon$

$$\frac{f(x^* + h) - f(x^*)}{h} > \delta$$

Pick a negative $h \in [-\varepsilon, 0)$. Then

$$f(x^* + h) < f(x^*) + \delta h < f(x^*)$$

and x^* cannot be the minimum. Analog for maximum with a positive h , then apply to $-f$. \square

Theorem 1.5 (Rolle). *Let $f \in C[a, b]$ with $f(a) = f(b)$. If f is differentiable in (a, b) then there exists a $\xi \in (a, b)$ with $f'(\xi) = 0$.*

Proof. Assume f is not constant. Since $[a, b]$ is compact there exists either a global minimum or maximum $\xi \in (a, b)$ and Theorem 1.4 can be applied. \square

Theorem 1.6 (Mean Value). *Let $f \in C[a, b]$ be differentiable in (a, b) . Then there exists a $\xi \in (a, b)$ with*

$$f'(\xi) = \frac{f(b) - f(a)}{b - a}$$

Proof. Apply Theorem 1.5 to

$$g(x) = f(x) - \frac{f(b) - f(a)}{b - a}(x - a)$$

\square

Remark 1.7.

1. More generally choose any $\varphi \in C^1[a, b]$ with $\varphi(a) = 0$ and $\varphi(b) = f(b) - f(a)$. Set $g(x) = f(x) - \varphi(x)$ to see there is a $\xi \in (a, b)$ with $f'(\xi) = \varphi'(\xi)$.
2. Let f be differentiable in (a, b) with $f' = 0$. For $x, y \in (a, b)$

$$0 = f'(\xi) = \frac{f(y) - f(x)}{y - x}$$

and f is a constant.

3. Another useful generalization: let $\Omega \subseteq \mathbb{R}^n$ be open and $f \in C^1(\Omega)$. For $x, y \in \Omega$ define $\varphi(t) = f(tx + (1 - t)y)$ and apply the chain rule for differentiation

$$\varphi'(\xi) = \nabla f(\xi x + (1 - \xi)y)^T(x - y) = f(x) - f(y)$$

4. The Cauchy Schwarz inequality then yields

$$\|f(x) - f(y)\| \leq \|\nabla f(\xi x + (1 - \xi)y)\| \|x - y\|$$

Theorem 1.8 (Differentiation Theorem). *Let $f \in C[a, b]$ and define*

$$F(x) = \int_a^x f(t) dt$$

Then $F \in C^1[a, b]$ with $F'(x) = f(x)$ for $x \in [a, b]$.

Proof. Applying the Mean Value Theorem of Integration gives

$$F(x + h) - F(x) = \int_x^{x+h} f(t) dt = f(\xi)h$$

for some $\xi \in (x, x + h)$. \square

Theorem 1.9 (Fundamental Theorem of Calculus). *Let $F \in C^1[a, b]$ with $F' = f$. Then*

$$F(b) - F(a) = \int_a^b f(t) dt$$

Lemma 1.10 (Integration by Substitution). *Let $I \subseteq \mathbb{R}$ be an interval and $f \in C(I)$. For $\varphi \in C([a, b], I)$ it follows*

$$\int_{\varphi(a)}^{\varphi(b)} f(x) dx = \int_a^b f(\varphi(t))\varphi'(t) dt$$

Proof. Let $F \in C^1(I)$ with $F' = f$. Then the chain rule for differentiation yields

$$\begin{aligned} \int_{\varphi(a)}^{\varphi(b)} f(x) dx &= F(\varphi(b)) - F(\varphi(a)) \\ &= F \circ \varphi(b) - F \circ \varphi(a) \\ &= \int_a^b (F \circ \varphi)'(t) dt \\ &= \int_a^b f(\varphi(t))\varphi'(t) dt \end{aligned}$$

□

Examples 1.11.

1. For $\varphi(x) = x^2 + 1$ it is $\varphi(0) = 1$ and $\varphi(2) = 5$. Thus

$$\int_0^2 x \cos(x^2 + 1) dx = \frac{1}{2} \int_0^2 2x \cos(x^2 + 1) dx = \frac{1}{2} \int_1^5 \cos(t) dt = \frac{1}{2}(\sin(5) - \sin(1))$$

2. Consider $\varphi(x) = \sin(x)$ where $\varphi(0) = 0$ and $\varphi(\pi/2) = 1$. Since $\cos(t) = \sqrt{1 - \sin^2(t)}$ it follows

$$\int_0^1 \sqrt{1 - x^2} dx = \int_{\cos(0)}^{\cos(\pi/2)} \sqrt{1 - x^2} dx = \int_0^{\pi/2} \sqrt{1 - \sin^2(t)} \cos(t) dt = \int_0^{\pi/2} \cos^2(t) dt$$

3. Let $f \in C[a, b]$ and $\varphi(x) = a + t(b - a)$. Then

$$\int_a^b f(x) dx = (b - a) \int_0^1 f(a + t(b - a)) dt$$

4. Let $f(x) = x^n$ and $\varphi(x) = t^m$. As expected

$$\int_0^1 x^n dx = \int_0^1 t^{nm} m t^{m-1} dt = m \int_0^1 t^{m(n+1)-1} dt = \left[\frac{m}{m(n+1)} t^{m(n+1)} \right]_0^1 = \frac{1}{n+1}$$

1.2 Directional Derivative and Gradients

Lemma 1.12 (Directional Derivative). *Let $\Omega \subseteq \mathbb{R}^n$ be open and $f \in C^1(\Omega)$. Then*

$$\frac{\partial f}{\partial d}(x) = \nabla f(x)^T d$$

for any $d \in \mathbb{R}^n$.

Proof. Let $\varphi(t) = f(x + td)$. Then $\varphi \in C^1[-\varepsilon, \varepsilon]$ for some $\varepsilon > 0$ and the chain rule yields

$$\varphi'(t) = \nabla f(x + td)^T d$$

Hence

$$\varphi'(0) = \lim_{t \rightarrow 0} \frac{\varphi(x + td) - \varphi(0)}{t} = \lim_{t \rightarrow 0} \frac{f(x + td) - f(x)}{t} = \nabla f(x)^T d$$

□

Remarks 1.13.

1. Note that by definition the directional derivative is invariant under multiplication with any $\lambda \neq 0$.
2. A similar proposition holds under the weaker assumption that d is a only feasible direction for f in x
3. For $d = \nabla f(x) / \|\nabla f(x)\|$ it follows that

$$\frac{\partial f}{\partial d}(x) = \|\nabla f(x)\| > 0$$

and for any other $d \in \mathbb{R}^n$ with $\|d\| = 1$ the Cauchy Schwarz inequality yields

$$\left| \frac{\partial f}{\partial d}(x) \right| = |\nabla f(x)^T d| \leq \|\nabla f(x)\| \|d\| = \|\nabla f(x)\|$$

Hence $\nabla f(x)$ is the direction of the greatest ascent and respectively, $-\nabla f(x)$ is the direction of the greatest descent.

Theorem 1.14 (First Order Necessary Condition). *Let $\Omega \subseteq \mathbb{R}^n$ be open and $f \in C^1(\Omega)$. If $x^* \in \Omega$ is a local minimizer then $\nabla f(x^*) = 0$.*

Proof. Let $h \in \mathbb{R}^n$ and $\delta > 0$ so that $x^* + th \in \Omega$ for all $t \in (-\delta, \delta)$. Then 0 is local minimizer for $\varphi(t) = f(x^* + th)$ and

$$\varphi'(0) = \nabla f(x^*)^T h = 0$$

Now let $h = \nabla f(x^*)$.

□

Theorem 1.15 (Banach Fixed-Point Theorem). *Let X be a Banach space and $f \in C(X, X)$ a contraction*

$$\|f(x) - f(y)\| \leq q \|x - y\| \text{ for all } x, y \in X$$

for some $0 < q < 1$. Then there exists a unique fix point $x^* \in X$ with

$$f(x^*) = x^*$$

Furthermore for any $x_0 \in X$ the sequence defined by

$$x_{n+1} = f(x_n)$$

converges against x^* .

Proof. Since $\|x_{n+1} - x_n\| = \|f(x_n) - f(x_{n-1})\| \leq q\|x_n - x_{n-1}\|$ it follows, that

$$\|x_{n+1} - x_n\| \leq q^n \|x_1 - x_0\|$$

Furthermore

$$\|x_n - x_m\| \leq \sum_{k=m}^n q^k \|x_1 - x_0\| = \frac{q^m - q^{n+1}}{1 - q} \|x_1 - x_0\|$$

and (x_n) is a Cauchy sequence. For its limit x^* we have

$$x^* = \lim_{n \rightarrow \infty} x_{n+1} = \lim_{n \rightarrow \infty} f(x_n) = f(x^*)$$

For any other $y^* \in X$ with $f(y^*) = y^*$ it follows, that

$$\|x^* - y^*\| = \|f(x^*) - f(y^*)\| \leq q\|x^* - y^*\|$$

and therefore $x^* = y^*$.

□

2 Nonlinear Optimization

2.1 Minimization without Constraints

Lemma 2.1 (Gradient Inequality). *Let $M \subseteq \mathbb{R}^n$ be a convex set and $f \in C^1(M)$. Then f is convex if and only if*

$$f(x) \geq f(y) + \nabla f(y)^T(x - y)$$

for all $x, y \in M$.

Proof. Let f be convex and $x, y \in M$. For $0 \leq \lambda \leq 1$ we have

$$f(\lambda x + (1 - \lambda)y) \leq \lambda f(x) + (1 - \lambda)f(y) = \lambda f(x) - \lambda f(y) + f(y)$$

and

$$f(x) - f(y) \geq \frac{f(\lambda x + (1 - \lambda)y) - f(y)}{\lambda} = \frac{f(y + \lambda(x - y)) - f(y)}{\lambda}$$

For $d = x - y$ and $\lambda \rightarrow 0$ the term on the right converges to the direction derivative of f in d

$$\frac{\partial f}{\partial d}(y) = \nabla f(y)^T d = \nabla f(y)^T(x - y)$$

Now let $x, y \in M$ and $0 \leq \lambda \leq 1$. For $z = \lambda x + (1 - \lambda)y \in M$ it follows that

$$\begin{aligned} \lambda f(x) &\geq \lambda f(z) + \lambda \nabla f(z)^T(x - z) \\ (1 - \lambda)f(y) &\geq (1 - \lambda)f(z) + (1 - \lambda)\nabla f(z)^T(y - z) \end{aligned}$$

Adding the two inequalities gives

$$\begin{aligned} \lambda f(x) + (1 - \lambda)f(y) &\geq f(z) + \nabla f(z)^T(\lambda x - \lambda z + (1 - \lambda)y - (1 - \lambda)z) \\ &= f(z) + \nabla f(z)^T(\lambda x + (1 - \lambda)y - z) \\ &= f(z) \end{aligned}$$

□

Exercise 2.2 (Facility Locations). *The facilities are located at:*

$$(3, 0), (0, -3), (1, 4)$$

Proof. Let

$$\begin{aligned} f(x) &= (x - 3)^2 + y^2 + x^2 + (y + 3)^2 + (x - 1)^2 + (y - 4)^2 \\ &= x^2 - 6x + 9 + y^2 + x^2 + y^2 + 6y + 9 + x^2 - 2x + 1 + y^2 - 8y + 16 \\ &= 3x^2 + 3y^2 - 8x - 2y + 35 \end{aligned}$$

Then

$$\nabla f(x, y) = (6x - 8, 6y - 2) \text{ and } \nabla^2 f(x, y) = \begin{pmatrix} 6 & 0 \\ 0 & 6 \end{pmatrix} > 0$$

Hence $(4/3, 1/3)$ is the global minimum.

□

Exercise 2.3 (Convex Functions). *The sum of convex functions is convex.*

Proof. Let $x, y \in M$. Since $\alpha_i > 0$ we have

$$\begin{aligned} f(\lambda x + (1 - \lambda)y) &= \sum_{i=1}^m \alpha_i f_i(\lambda x + (1 - \lambda)y) \\ &\leq \sum_{i=1}^m \alpha_i \lambda f_i(x) + \sum_{i=1}^m \alpha_i (1 - \lambda) f_i(y) = \lambda f(x) + (1 - \lambda) f(y) \end{aligned}$$

Let $f(x) = x^2$. Then $-f$ is not convex, e.g. $x = 1, y = -1$ and $\lambda = 0.5$.

Exercise 2.4 (Solution of Quadratic Inequality). *Let*

$$f(x) = x^T A x + b^T x + c$$

Proof. The product rule gives

$$\nabla f(x) = x^T A + A x + b = (A^T + A)x + b = 2A x + b$$

Thus $\nabla^2 f(x) = 2A > 0$ and f is convex. Hence the level set Γ_{-c} is convex. Since the intersection of convex sets is convex $\Gamma_{-c} \cap \{x \in \mathbb{R}^n : g^T x + h = 0\}$ is convex, too.

Exercise 2.5 (Line Search on Compact Convex Sets). *Let $S \subset \mathbb{R}^n$ be compact and convex. Furthermore let $f \in C^1(S)$ be convex, $x \in S$ and $d \in \mathbb{R}^n$ a descent direction of f in x with $\nabla f(x)^T d < 0$.*

Proof. If $x + \lambda^* d$ is an optimal solution then $\nabla f(x + \lambda^* d)^T d = 0$ according to Theorem 1.14. Let $\nabla f(x + \lambda^* d)^T d = 0$. Then Lemma 2.1 gives

$$f(x + \lambda d) \geq f(x + \lambda^* d) + (\lambda - \lambda^*) \nabla f(x + \lambda^* d)^T d = f(x + \lambda^* d)$$

and $x + \lambda^* d$ is an optimal solution.

Exercise 2.6 (Steepest Descent). *Let*

$$f(x) = \frac{1}{2} x^T A x + b^T x + c$$

where A is symmetrical and positive definite.

Proof. Since $\nabla f(x) = A x + b$ and $\nabla^2 f(x) = A > 0$ it follows $x^* = -A^{-1}b$. Let v be eigenvector with $Av = \mu v$. For $x_0 = x^* + \theta v$ we have

$$\nabla f(x_0) = A x^* + \mu \theta v + b = \mu \theta v$$

and for $\lambda \geq 0$

$$\arg \min \{f(x_0 - \lambda \nabla f(x_0))\} = \arg \min \{f(x^* + \theta v - \lambda \mu \theta v)\} = \mu^{-1}$$

Thus

$$x_1 = x_0 - \mu^{-1} \nabla f(x_0) = x^* + \theta v - \mu^{-1} \mu \theta v = x^*$$

and $\nabla f(x_1) = 0$. Hence the algorithm stops after the first iteration. Now let

$$x_0 = x^* + \sum_{i=0}^m \theta_i v_i$$

for orthogonal eigenvectors with $Av_i = \mu_i$ and $m \leq n$. Then

$$\nabla f(x_0) = Ax^* + \sum_{i=0}^m \mu_i \theta_i v_i + b = \sum_{i=0}^m \mu_i \theta_i v_i$$

and

$$x_1 = x_0 - \lambda \sum_{i=0}^m \mu_i \theta_i v_i = x^* + \sum_{i=0}^m \theta_i v_i - \lambda \sum_{i=0}^m \mu_i \theta_i v_i = x^* + \sum_{i=0}^m (1 - \lambda \mu_i) \theta_i v_i$$

Since x^* is the minimum we have $\nabla f(x_1) = 0$ iff $\lambda = \mu_i^{-1}$ for all $0 \leq i \leq m$. \square

2.2 One Dimensional Minimization and Direct Search

Definition 2.7 (Unimodal Function). *A function $f : [a, b] \rightarrow \mathbb{R}$ is called unimodal if there exists a $\xi \in [a, b]$, so that f is strictly decreasing in $[a, \xi]$ and strictly increasing in $[\xi, b]$.*

In fact ξ is the unique minimum of f in $[a, b]$. According to the definition, for $a \leq x < y \leq b$ we have

$$f(x) > f(y) \text{ for } x, y \in [a, \xi] \text{ and } f(x) < f(y) \text{ for } x, y \in [\xi, b]$$

Thus

$$\xi \in [a, y] \text{ if } f(x) < f(y) \text{ and } \xi \in [x, b] \text{ if } f(x) \geq f(y)$$

Consider now a symmetrical partitioning of the interval $[0, 1]$ where two consecutive partitionings hold the same ratio respectively:

$$\sigma = 1 - \tau \text{ and } \frac{1}{\tau} = \frac{\tau}{\sigma}$$

Then $1 - \tau = \tau^2$ and solving the quadratic equation $\tau^2 + \tau = 1$ yields

$$\tau = \frac{\sqrt{5} - 1}{2} \approx 0.61803$$

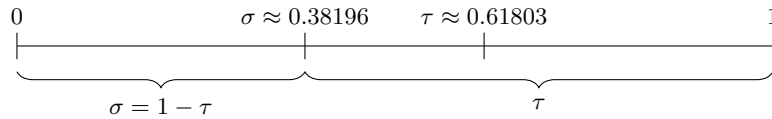


Figure 1: Golden Section

Let now $[a_0, b_0] = [a, b]$ and define

$$[a_{k+1}, b_{k+1}] = \begin{cases} [a_k, y_k] & \text{if } f(x_k) < f(y_k) \\ [x_k, b_k] & \text{if } f(x_k) \geq f(y_k) \end{cases}$$

where

$$\begin{aligned}x_k &= b_k - \tau(b_k - a_k) \\ y_k &= a_k + \tau(b_k - a_k)\end{aligned}$$

It follows that $[a_k, b_k] \supset [a_{k+1}, b_{k+1}]$ is a decreasing series of intervals with

$$(b_{k+1} - a_{k+1}) = \tau(b_k - a_k)$$

where the interval converges to ξ . This leads to the following algorithm:

Algorithm 2.8 (Golden Section Search).

```
"""Basic implementation of the golden section search, this easily can be
improved by storing and resuing the results of the previous iteration
"""
```

```
import math

def golden_section_search(f, I, eps=0.00001):
    t = 0.5 * (math.sqrt(5) - 1)
    a, b = I
    while abs(b - a) > eps:
        x, y = b - t * (b - a), a + t * (b - a)
        if f(x) > f(y):
            a = x
        else:
            b = y
    return (a + b) / 2

if __name__ == '__main__':
    p, q, I = 0, 0, (-10, 10)
    p, q, I = -4, 1, (-10, 10)
    f = lambda x: (x + p) ** 2 + q
    x0 = golden_section_search(f, I)
    print(f'arg min f on {I}: {x0}')
```

Algorithm 2.9 (Steepest Descent).

Let $f \in C^1(\mathbb{R}^n)$ and $x_0 \in \mathbb{R}^n$. For $0 < \alpha \leq \beta < 1$ and $\gamma < 1$ let

Exercise 2.10 (Surprising Convergence). *Example for $f \in C^2(\mathbb{R})$ with a sequence of strict local minima converging to a strict local maximum.*

Proof. Let $f \in C[a, b]$ and $\xi \in (a, b)$ so that f is strictly increasing in $(a, \xi]$ and strictly decreasing in $[\xi, b)$. Define

$$g(x) = \int_{\xi-x}^{\xi+x} f(t) dt$$

□

3 The Road to Reality

3.1 Hyperbolic Geometry

The ratio between the area A and A' of two similar shapes is given by

$$A' = k^2 A$$

Theorem 3.1 (Pythagoras).

$$a^2 + b^2 = c^2$$

Proof. Let A, B and C be the areas of the three triangles respectively. All triangles are similar, hence

$$B = \frac{b^2}{a^2} A \text{ and } C = \frac{c^2}{b^2} B$$

Since $A + B = C$ it follows that

$$a^2 + b^2 = \frac{b^2 A}{B} + b^2 = \frac{b^2(A + B)}{B} = \frac{b^2 C}{B} = c^2$$

□

Lemma 3.2 (label and Projective Representation). *The mapping from conformal and projective representation of any point is given by the radial expansion of the following factor*

$$\frac{2R}{R^2 + r^2}$$

Proof. For any point the distance from the origin with regard to the two representations is given by

$$\log \frac{R+r}{R-r} = \frac{1}{2} \log \frac{R+r'}{R-r'} = \log \frac{(R+r')^2}{(R-r')^2}$$

This gives

$$(R-r)^2(R+r') = (R+r)^2(R-r') \text{ and } -4R^2r + 2R^2r' + 2r^2r' = 0$$

Hence

$$r' = \frac{2R^2}{R^2 + r^2} r$$

□

3.2 Complex Numbers

Lemma 3.3 (Basic Formulas).

1. *It is*

$$(a + ib)(c + id) = (ac - bd) + i(ad + bc)$$

2. *Thus*

$$(a + ib)^2 = (a^2 - b^2) + i2ab$$

and

$$(a + ib)(a - ib) = a^2 + iab - iab - i^2b^2 = a^2 + b^2$$

3. *Hence*

$$\frac{a + ib}{c + id} = \frac{(a + ib)(c - id)}{c^2 + d^2} = \frac{ac + bd}{c^2 + d^2} + i\frac{bc - ad}{c^2 + d^2}$$

4. *For*

$$z = \sqrt{\frac{1}{2}(a + \sqrt{a^2 + b^2})} + i\sqrt{\frac{1}{2}(-a + \sqrt{a^2 + b^2})}$$

it follows

$$z^2 = \frac{1}{2}(a + \sqrt{a^2 + b^2}) - \frac{1}{2}(-a + \sqrt{a^2 + b^2}) + i2\sqrt{\frac{1}{4}(\sqrt{a^2 + b^2})^2} - a^2 = a + ib$$

Lemma 3.4 (Binomial Theorem).

1. *For the binomial coefficient Pascal's identity holds*

$$\binom{n}{k-1} + \binom{n}{k} = \binom{n+1}{k}$$

2. *The following equation states the binomial identity*

$$(a + b)^n = \sum_{k=0}^n \binom{n}{k} a^k b^{n-k} = \sum_{k=0}^n \binom{n}{k} a^{n-k} b^k$$

3. *For $a = 1$ follows*

$$(1 + x)^n = \sum_{k=0}^n \binom{n}{k} x^k$$

Proof. It is

$$\binom{n}{k} + \binom{n}{k-1} = \frac{n!}{k!(n-k)!} + \frac{n!}{(k-1)!(n-k+1)!} = \frac{n!(n+1-k) + n!k!}{k!(n+1-k)!} = \binom{n+1}{k}$$

Furthermore by using induction

$$\begin{aligned}
(a+b)^{n+1} &= \sum_{k=0}^n \binom{n}{k} a^{k+1} b^{n-k} + \sum_{k=0}^n \binom{n}{k} a^k b^{n+1-k} \\
&= \sum_{k=1}^{n+1} \binom{n}{k-1} a^k b^{n+1-k} + \sum_{k=0}^n \binom{n}{k} a^k b^{n+1-k} \\
&= \sum_{k=0}^{n+1} \binom{n+1}{k} a^k b^{n+1-k}
\end{aligned}$$

□

3.3 Exponential Function and Logarithms

Exercise 3.5 (Exponential Function). *The Cauchy product yields*

$$\sum_{n=0}^{\infty} a_n \sum_{n=0}^{\infty} b_n = \sum_{n=0}^{\infty} \sum_{k=0}^n a_k b_{n-k}$$

if at least one of the series is absolutely convergent. Hence

$$\begin{aligned}
\sum_{n=0}^{\infty} \frac{1}{n!} z^n \sum_{n=0}^{\infty} \frac{1}{n!} w^n &= \sum_{n=0}^{\infty} \sum_{k=0}^n \frac{1}{k!} z^k \frac{1}{(n-k)!} w^{n-k} \\
&= \sum_{n=0}^{\infty} \frac{1}{n!} \sum_{k=0}^n \binom{n}{k} z^k w^{n-k} \\
&= \sum_{n=0}^{\infty} \frac{1}{n!} (z+w)^n
\end{aligned}$$

Let $t \in \mathbb{R}$. Then

$$\begin{aligned}
e^{it} &= \sum_{k=0}^{\infty} \frac{1}{k!} (it)^k \\
&= \sum_{k=0}^{\infty} \frac{1}{2k!} (it)^{2k} + \sum_{k=0}^{\infty} \frac{1}{(2k+1)!} (it)^{2k+1} \\
&= \sum_{k=0}^{\infty} \frac{(-1)^k}{2k!} t^{2k} + i \sum_{k=0}^{\infty} \frac{(-1)^k}{(2k+1)!} t^{2k+1} \\
&= \cos t + i \sin t
\end{aligned}$$

More generally for $z = \log r + it$

$$e^z = e^{\log r + it} = r e^{it} = r(\cos t + i \sin t)$$

For $r = 1$ and $t = 2\pi$ this yields

$$e^{2\pi i} = \cos 2\pi + i \sin 2\pi = 1$$

and for $t = 2\pi$ we get

Lemma 3.6 (Euler Equation).

$$e^{\pi i} + 1 = 0$$

Exercise 3.7.

1. If $e^z = w$ then $z + \pi i$ is a logarithm to $-w$: $e^{z+\pi i} = e^z e^{\pi i} = -e^z = -w$.

2. Since $e^{i(s+t)} = e^{is} e^{it}$ it follows

$$\begin{aligned}\cos(s+t) + i \sin(s+t) &= (\cos s + i \sin s)(\cos t + i \sin t) \\ &= \cos s \cos t - \sin s \sin t + i(\cos s \sin t + \sin s \cos t)\end{aligned}$$

Hence

$$\begin{aligned}\cos(s+t) &= \cos s \cos t - \sin s \sin t \\ \sin(s+t) &= \cos s \sin t + \sin s \cos t\end{aligned}$$

3. It is $e^{3it} = (e^{it})^3$ and thus

$$\cos 3t + i \sin 3t = (\cos t + i \sin t)^3 = \cos^3 t - 3 \cos t \sin^2 t + i(\cos^2 t \sin t - \sin^3 t)$$

4. Fun facts

$$e^{1-4\pi^2} = e^{1+(2i\pi)^2} = e e^{2\pi i} e^{2\pi i} = e$$

and $i = e^{i\pi/2}$ gives

$$i^i = e^{i \log i} = e^{i i \pi/2} = e^{-\pi/2} \in \mathbb{R}$$

3.4 Complex Analysis

Definition 3.8 (holomorphic Function). Let $\Omega \subseteq \mathbb{C}$ be open. A function $f : \Omega \rightarrow \mathbb{C}$ is called differentiable at $z \in \Omega$ if the limit

$$f'(z) = \lim_{h \rightarrow 0} \frac{f(z+h) - f(z)}{h}$$

exists. f is called holomorphic on Ω if f is complex differentiable at all points of Ω and $f' : \Omega \rightarrow \mathbb{C}$ is called the derivative of f .

Remarks 3.9.

1. f is differentiable at $z_0 \in \Omega$ iff the limit

$$f'(z_0) = \lim_{z \rightarrow z_0} \frac{f(z) - f(z_0)}{z - z_0}$$

exists

2. If f is differentiable at $z_0 \in \Omega$ and $\varepsilon > 0$ then there exists a small enough environment of z_0 , so that

$$|f(z) - f(z_0) - f'(z_0)(z - z_0)| < \varepsilon |z - z_0|$$

Theorem 3.10 (Cauchy Riemann Equations). *Let $f = u + iv$ be holomorphic. Then f satisfies the Cauchy Riemann equations*

$$\begin{aligned}\frac{\partial u}{\partial x} &= \frac{\partial v}{\partial y} \\ \frac{\partial u}{\partial y} &= -\frac{\partial v}{\partial x}\end{aligned}$$

Proof. For $h \in \mathbb{R}$ follows

$$\lim_{h \rightarrow 0} \frac{f(z+h) - f(z)}{h} = \frac{\partial u}{\partial x}(z) + i \frac{\partial v}{\partial x}(z)$$

and

$$\lim_{h \rightarrow 0} \frac{f(z+ih) - f(z)}{ih} = \frac{\partial u}{i\partial y}(z) + \frac{\partial v}{\partial y}(z) = \frac{\partial v}{\partial y}(z) - i \frac{\partial u}{\partial y}(z)$$

□

Examples 3.11.

1. Let $f(z) = z^3$. Then $u(x, y) + iv(x, y) = x^3 - 3xy^2 + i(3x^2y - y^3)$ and as expected

$$\begin{aligned}\frac{\partial u}{\partial x}(x, y) &= x^3 - 3y^2 & \text{and} & & \frac{\partial u}{\partial y}(x, y) &= -6xy \\ \frac{\partial v}{\partial x}(x, y) &= 6xy & \text{and} & & \frac{\partial v}{\partial y}(x, y) &= x^3 - 3y^2\end{aligned}$$

Lemma 3.12. *Let $D \subseteq \mathbb{C}$ be connected. For arbitrary $z, w \in D$ there exists a polygonal path from z to w .*

Proof. For any path from z to w the image is compact, which can be used to define a finite subcover of disks. Use the center points to define the polygonal path. □

Lemma 3.13. *Let $\gamma : [a, b] \rightarrow \mathbb{C}$ a smooth path, $\psi : [c, d] \rightarrow [a, b]$ a smooth and increasing bijection and f continuous.*

$$\int_{\gamma} f(z) dz = \int_{\gamma \circ \psi} f(z) dz$$

Proof. It is

$$\begin{aligned}\int_{\gamma \circ \psi} f(z) dz &= \int_c^d f(\gamma \circ \psi(t))(\gamma \circ \psi)'(t) dt \\ &= \int_{\psi(a)}^{\psi(b)} f(\gamma(\psi(t)))\gamma'(\psi(t))\psi'(t) dt \\ &= \int_a^b f(\gamma(s))\gamma'(s) ds = \int_{\gamma} f(z) dz\end{aligned}$$

□

Lemma 3.14. For a smooth path $\gamma : [a, b] \rightarrow \mathbb{C}$ define $-\gamma(t) = a + b - t$. Then

$$\int_{-\gamma} f(z) dz = - \int_{\gamma} f(z) dz$$

Proof. Using integration by substitution

$$\int_{-\gamma} f(z) dz = - \int_a^b f(\gamma(a + b - t)) \gamma'(a + b - t) dt = \int_b^a f(\gamma(s)) \gamma'(s) ds = - \int_{\gamma} f(z) dz$$

□

In order to use the results from real calculus recall the fact, that for every $z \in \mathbb{C}$ there exists a $t \in [0, 2\pi]$, so that $z = |z|e^{it}$ and hence $|z| = ze^{-it}$.

Lemma 3.15. Let $f \in C[a, b]$. Then

$$\left| \int_a^b f(x) dx \right| \leq \int_a^b |f(x)| dx$$

Proof. Using the estimation for integrals from real calculus

$$\left| \int_a^b f(x) dx \right| = e^{-it} \int_a^b f(x) dx \leq \int_a^b |e^{-it} f(x)| dx = \int_a^b |f(x)| dx$$

□

Let $\gamma : [a, b] \rightarrow \mathbb{C}$ be a smooth path and $a = t_0 < t_1 < \dots < t_n = b$ a partitioning of $[a, b]$. Then

$$\sum_{k=1}^n |\gamma(t_k) - \gamma(t_{k-1})| = \sum_{k=1}^n \left| \frac{\gamma(t_k) - \gamma(t_{k-1})}{t_k - t_{k-1}} \right| (t_k - t_{k-1}) = \sum_{k=1}^n |\gamma'(\xi_k)| (t_k - t_{k-1})$$

yields a reasonable approximation of the length of the path. Hence

Definition 3.16. For a smooth path $\gamma : [a, b] \rightarrow \mathbb{C}$

$$L(\gamma) = \int_a^b |\gamma'(t)| dt$$

is called the length of γ .

Lemma 3.17 (Estimation Lemma). Let $\gamma : [a, b] \rightarrow \mathbb{C}$ be a smooth path. Then

$$\left| \int_{\gamma} f(z) dz \right| \leq L(\gamma) \max_{\gamma[a, b]} f$$

Proof. Using the definition above

$$\left| \int_{\gamma} f(z) dz \right| = \left| \int_a^b f(\gamma(t)) \gamma'(t) dt \right| \leq \int_a^b |f(\gamma(t)) \gamma'(t)| dt \leq \max_{\gamma[a,b]} |f| \int_a^b |\gamma'(t)| dt$$

□

Examples 3.18.

1. Let $\gamma(t) = t + it$. Then

$$\int_{\gamma} z^2 dz = \int_0^1 (t + it)^2 (1 + i) dt = (1 + i) \int_0^1 2it^2 dt = [(-2 + 2i)t^3]_0^1 = -\frac{2}{3} + i\frac{2}{3}$$

2. For $\gamma(t) = t^2 + it$

$$\begin{aligned} \int_{\gamma} z^2 dz &= \int_0^1 (t^2 + it)^2 (2 + i) dt = \int_0^1 (2t^5 - 4t^3) + i(5t^4 - t^2) dt \\ &= \left[\frac{1}{3}t^6 - t^4 \right]_0^1 + i \left[t^5 - \frac{1}{3}t^3 \right]_0^1 = -\frac{2}{3} + i\frac{2}{3} \end{aligned}$$

3. And $\gamma(t) = i + e^{it}$

$$\begin{aligned} \int_{\gamma} z^2 dz &= \int_{3/2\pi}^{2\pi} (i + e^{it})^2 ie^{it} dt = \int_{3/2\pi}^{2\pi} (-1 + 2ie^{it} + e^{2it}) ie^{it} dt \\ &= \int_{3/2\pi}^{2\pi} -ie^{it} - 2e^{2it} + ie^{3it} dt = \left[-e^{it} + ie^{2it} + \frac{1}{3}e^{3it} \right]_{3/2\pi}^{2\pi} \\ &= \left(-1 + i + \frac{1}{3} \right) - \left(i - i + \frac{1}{3}i \right) = -\frac{2}{3} + i\frac{2}{3} \end{aligned}$$

4. Let $\gamma(t) = e^{it}$ and $k \neq -1$. Then

$$\int_{\gamma} z^k dz = \int_0^{2\pi} e^{ikt} ie^{it} dt = \int_0^{2\pi} ie^{i(k+1)t} dt = 0$$

Theorem 3.19. Let $D \subseteq \mathbb{C}$ be a connected domain and $f \in C(D)$. Then the following assertions are equivalent

1. f has an antiderivative
2. For every closed path γ

$$\int_{\gamma} f(z) dz = 0$$

Proof. Let $F' = f$. Since γ is closed

$$\int_{\gamma} f(z) dz = \int_a^b f(\gamma(t)) \gamma'(t) dt = \int_a^b (F \circ \gamma)'(t) dt = F(\gamma(b)) - F(\gamma(a)) = 0$$

Now fix some arbitrary $a \in D$. For $z \in D$ let γ_z be a path from a to z and define

$$F(z) = \int_{\gamma_z} f(\zeta) d\zeta$$

This is well defined since the integral of f vanishes over each closed path. Moreover, since $\gamma_{z+h} + [z+h, z] - \gamma_z$ defines a closed path

$$F(z+h) - F(z) = \int_{\gamma_{z+h}} f(z) dz - \int_{\gamma_z} f(z) dz = \int_{[z, z+h]} f(z) dz = h \int_0^1 f(z+th) dt$$

Here the latter integral is continuous at 0 with respect to h

$$\left| \int_0^1 f(z+th) - f(z) dt \right| \leq \int_0^1 |f(z+th) - f(z)| dt \leq \max_{t \in [0,1]} |f(z+th) - f(z)|$$

□

Corollary 3.20. *The second assertion can be weakened to*

$$\int_{\partial \Delta} f(z) dz = 0$$

for every triangle $\Delta \subset D$, where e.g. D is convex or star shaped. Here the antiderivative can directly be defined as

$$F(z) = \int_{[a,z]} f(\zeta) d\zeta$$

similar to the real calculus approach. Note, that under this conditions f always has a local antiderivative.

Examples 3.21.

1. Let $z_0 \in \mathbb{C}$ and $\gamma(t) = z_0 + e^{it}$ for $t \in [0, 2\pi]$. Then

$$\int_{\gamma} \frac{1}{z - z_0} dz = \int_0^{2\pi} \frac{ie^{it}}{z_0 + e^{it} - z_0} dt = \int_0^{2\pi} i dt = 2\pi i$$

and thus $1/(z - z_0)$ has no antiderivative on $\mathbb{C} \setminus \{z_0\}$

2. Let $z_0 \in \mathbb{C}$ and $z \in D = D_r(z_0)$. Applying [Theorem 3.19](#). to ∂D and a small enough circle around z gives

$$\int_{\partial D} \frac{1}{\zeta - z} d\zeta = \int_{\partial D} \frac{1}{\zeta - z_0} d\zeta = 2\pi i$$

Theorem 3.22 (Goursat). *Let $\Omega \subseteq \mathbb{C}$ be open and f holomorphic on Ω . Then*

$$\int_{\partial\Delta} f(z) dz = 0$$

for every triangle $\Delta \subset \Omega$.

Proof. Choose a sequence of triangles $\Delta \supset \Delta_0 \supset \Delta_1 \cdots \supset \Delta_k$ as depicted. Since all the triangles are compact with a vanishing diameter there exists a unique $z_0 \in \Omega$ with $\bigcap \Delta_k = \{z_0\}$. Thus

$$\left| \int_{\partial\Delta} f(z) dz \right| \leq 4^k \left| \int_{\partial\Delta_k} f(z) dz \right| = 4^k \left| \int_{\partial\Delta_k} f(z) - f(z_0) - f'(z_0)(z - z_0) dz \right|$$

Furthermore $L(\partial\Delta) = 2^{-k} L(\partial\Delta_k)$ and

$$|z - z_0| < L(\partial\Delta_k) = 2^{-k} L(\partial\Delta)$$

for any $z \in \Delta_k$. Since f is holomorphic at z_0 for any given $\varepsilon > 0$ there exists a sufficiently large enough k , so that

$$\begin{aligned} \left| \int_{\partial\Delta} f(z) dz \right| &\leq 4^k L(\partial\Delta_k) \max_{z \in \Delta_k} |f(z) - f(z_0) - f'(z_0)(z - z_0)| \\ &\leq 4^k L(\partial\Delta_k) \varepsilon \max_{z \in \Delta_k} |z - z_0| \\ &\leq L(\partial\Delta)^2 \varepsilon \end{aligned}$$

□

Corollary 3.23.

1. *A holomorphic function always has a local antiderivative*
2. *A holomorphic function on a star shaped domain has a global antiderivative and*

$$\int_{\gamma} f(z) dz = 0$$

for any closed path

3. *The prerequisites of Goursat theorem can be weakened to continuous and holomorphic with the exception of a finite number of points: adequate partitioning of the original triangle*

Theorem 3.24 (Cauchy's Integral Formula). *Let $\Omega \subseteq \mathbb{C}$ be open and f holomorphic on Ω . Further let $D \subset \Omega$ be a disc. Then*

$$f(z) = \frac{1}{2\pi i} \int_{\partial D} \frac{f(\zeta)}{\zeta - z} dz$$

for $z \in D$.

Proof. For $z \in D$ define

$$h(\zeta) = \frac{f(\zeta) - f(z)}{\zeta - z}$$

for $\zeta \neq z$ and $f'(z)$ for $\zeta = z$. Then h is holomorphic on $D \setminus \{z\}$ and continuous at z

$$0 = \int_{\partial D} h(\zeta) d\zeta = \int_{\partial D} \frac{f(\zeta)}{\zeta - z} d\zeta - f(z) \int_{\partial D} \frac{1}{\zeta - z} d\zeta = \int_{\partial D} \frac{f(\zeta)}{\zeta - z} d\zeta - 2\pi i f(z)$$

□

4 Neural Networks

4.1 The Perceptron

Definition 4.1 (Binary Classifiers). *Let $X \subset \mathbb{R}^n$ be the union of two finite disjoint sets $X = M \cup N$.*

1. *A binary classification problem is the task to find a mapping $f : X \rightarrow \{0, 1\}$ with*

$$f(x) = \begin{cases} 1 & \text{for } x \in M \\ 0 & \text{for } x \in N \end{cases}$$

f then is called a binary classifier for X

2. *X is called separable if there exists a weight vector $w \in \mathbb{R}^n$ and a bias $b \in \mathbb{R}$ so that*

$$\begin{aligned} wx + b &> 0 & \text{for } x \in M \\ wx + b &< 0 & \text{for } x \in N \end{aligned}$$

3. *The weight w and the bias b are called solution to the classification problem. They implicitly define a binary classifier via*

$$f(x) = \begin{cases} 1 & \text{if } wx + b > 0 \\ 0 & \text{if } wx + b < 0 \end{cases}$$

Examples 4.2.

1. *Let $X = \{0, 1\} \times \{0, 1\}$ and consider the and operator $f(1, 1) = 1$ and $f(x, y) = 0$ elsewhere. Then $w = (3, 3)$ and $b = -5$ yield a solution to the classification problem $M = f^{-1}(1)$ and $N = f^{-1}(0)$*
2. *Again let $X = \{0, 1\} \times \{0, 1\}$ and $f(1, 0) = f(0, 1) = 1$ and $f(0, 0) = f(1, 1) = 0$, the xor operator. Thus for any weight (w_1, w_2) and any bias b*

$$\begin{aligned} w_1 + b &> 0 \\ w_2 + b &> 0 \end{aligned}$$

$$\begin{aligned} w_1 + w_2 + b &\leq 0 \\ b &\leq 0 \end{aligned}$$

Adding two equations respectively shows that there is no solution

3. *The bias can be integrated into the weight vector via $w' = (w, b) \in \mathbb{R}^{n+1}$ and $x' = (x, 1) \in \mathbb{R}^{n+1}$. Separability then reduces to*

$$w'x' > 0$$

Geometrical Interpretation

The idea for the perceptron most likely has its origin in a simple geometrical observation. Recall that for $x, y \in \mathbb{R}^n$ the dot product can be expressed as

$$xy = \|x\|\|y\|\cos \alpha$$

where α is the angle between the two vectors. Hence the product is positive if the angle is less than 90° degrees and negative if the angle is between 90° and 180° degrees

$$xy > 0 \quad \text{for } 0 \leq \alpha < \pi/2$$

$$xy < 0 \quad \text{for } \pi/2 < \alpha \leq \pi$$

Note, that the sign does not depend on the vector lengths, but solely on the angle.

For any two vectors it is easy enough to find a weight that satisfies $wx > 0$ and $wy > 0$. Generally $w = x + y$ is a good guess, but not always correct as shown below in [Figure 2](#).

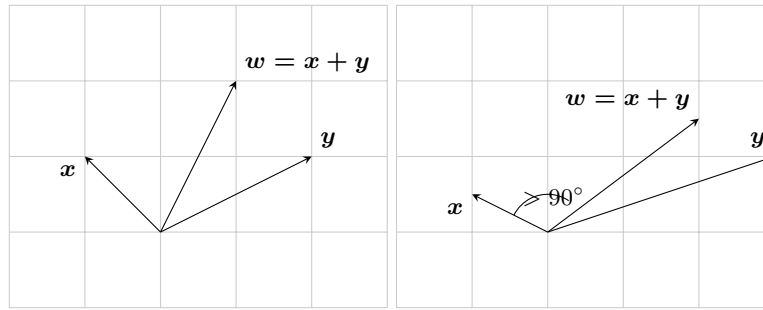


Figure 2: Dot Product and Angle

But, the more similar the lengths of the two vectors are the more likely $x + y$ works. Actually the threshold is tbd. An iterative approach is to increase $w = x + y$ in the direction of vector with the angle greater than 90° degrees.

$$w_{n+1} = \begin{cases} w_n + x & \text{if } w_n x < 0 \\ w_n + y & \text{if } w_n y < 0 \end{cases}$$

Examples 4.3.

1. Let $x = (-1, 1)$ and $y = (6, 1)$. Then

$$w_0 = (5, 2) \quad w_0 x = -3 \quad w_0 y = 32$$

$$w_1 = (4, 3) \quad w_1 x = -1 \quad w_1 y = 27$$

$$w_2 = (3, 4) \quad w_2 x = 1 \quad w_2 y = 22$$

2. Let $x = (4, -6)$ and $y = (-10, 5)$. Then

$$\begin{array}{lll} w_0 = (-6, -1) & w_0x = -18 & w_0y = 55 \\ w_1 = (-2, -7) & w_1x = 34 & w_1y = -15 \\ w_2 = (-12, -2) & w_2x = -36 & w_2y = 110 \\ w_3 = (-8, -8) & w_3x = 16 & w_3y = 40 \end{array}$$

Algorithm 4.4 (Weight).

```

from random import randint

def add(v, w):
    return tuple(x + y for x, y in zip(v, w))

def dotproduct(v, w):
    return sum(x * y for x, y in zip(v, w))

def weight(x, y):
    w = add(x, y)
    while True:
        if dotproduct(w, x) < 0:
            w = add(w, x)
        elif dotproduct(w, y) < 0:
            w = add(w, y)
        else:
            return w

def test_weight(x, y):
    print(f'x = {x}, y = {y}')
    i, w = 0, add(x, y)
    while True:
        a, b = dotproduct(w, x), dotproduct(w, y)
        print(f'    w[{i}] = {w}, w[{i}] * x = {a}, w[{i}] * y = {b}')
        if a < 0:
            w = add(w, x)
        elif b < 0:
            w = add(w, y)
        else:
            break
        i = i + 1

if __name__ == '__main__':
    n = 10
    for _ in range(3):
        x, y = (randint(-n, n), randint(-n, n)), (randint(-n, n), randint(-n, n))
        test_weight(x, y)

```

Algorithm 4.5 (Perceptron).

```

def heaviside(x):
    return 0 if x < 0.0 else 1

def dotproduct(v, w):
    return sum(x * y for x, y in zip(v, w))

def perceptron(M, t=0.1, max_iterations=50):
    n = len(M[0])
    w, b = n * [0.0], 0.0
    for _ in range(max_iterations):
        done = True
        for x, y in M:
            if dotproduct(w, x) + b >= 0:
                z = 1
            else:
                z = 0
            w = [w[i] - t * (z - y) * x[i] for i in range(n)]
            b = b - t * (z - y)
            if not y == z:
                done = False
        if done:
            break
    return lambda x: heaviside(dotproduct(w, x) + b)

def and_(x, y):
    return int(x and y)

def or_(x, y):
    return int(x or y)

def xor(x, y):
    return int(x and not y or not x and y)

def nand(x, y):
    return int(not x and not y)

def test_perceptron(f, X):
    M = [(x, f(*x)) for x in X]
    p = perceptron(M)
    for x, y in M:
        print(f'{x} -> {p(x)}: {y}')
    print(f'{f.__name__}: {all(p(x) == y for x, y in M)}')

if __name__ == '__main__':
    X = [(0, 0), (0, 1), (1, 0), (1, 1)]
    test_perceptron(or_, X)
    test_perceptron(and_, X)
    test_perceptron(nand, X)
    test_perceptron(xor, X)

```

4.2 The Backtracking Algorithm

Definition 4.6 (Activation Functions).

1. The Heaviside function $H : \mathbb{R} \rightarrow \{0, 1\}$ is defined as

$$H(x) = \begin{cases} 1 & \text{for } x > 0 \\ 0 & \text{for } x \leq 0 \end{cases}$$

2. The sigmoid function $\sigma \in C^\infty(\mathbb{R})$ is defined as

$$\sigma(x) = \frac{1}{1 + e^{-x}}$$

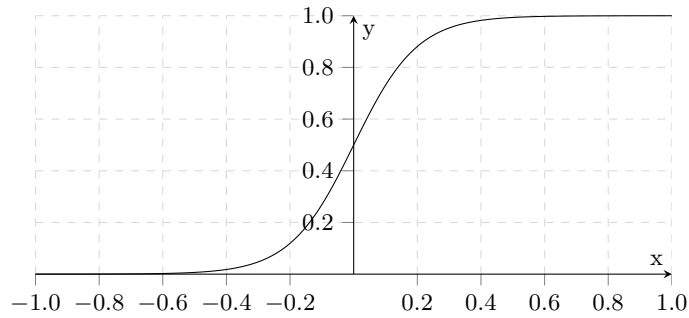


Figure 3: The sigmoid function $\sigma(x) = \frac{1}{1+e^{-x}}$

Remarks 4.7.

1. Since the heaviside function is not continuous and therefore not differentiable at 0 the sigmoid function is often considered its smooth counterpart
2. The definition of the sigmoid function yields $0 < \sigma(x) < 1$ as well as $\sigma(x) \rightarrow 0$ for $x \rightarrow -\infty$ and $\sigma(x) \rightarrow 1$ for $x \rightarrow \infty$
3. The quotient rule yields

$$\sigma'(x) = -\frac{-e^{-x}}{(1 + e^{-x})^2} = \sigma(x) \frac{1 + e^{-x} - 1}{1 + e^{-x}} = \sigma(x)(1 - \sigma(x))$$

and σ is monotonically increasing over its domain