Notes on Mathematics

Contents

1 Calculus

1.1 Differentation and Integration

Lemma 1.1 (Simple Calculations).

1. For $1 = xx^{-1}$ the product rule yields $0 = x^{-1} + x(x^{-1})'$. Hence

$$\frac{d}{dx}x^{-1} = -\frac{1}{x^2}$$

2. Similarly $x = \sqrt{x}^2$ and $1 = 2\sqrt{x}\sqrt{x}'$ and so

$$\frac{d}{dx}\sqrt{x} = \frac{1}{2\sqrt{x}}$$

3. It is

$$\frac{d}{dx}x^n = nx^{n-1}$$

   since via induction the product rule yields

$$\frac{d}{dx}x^n = \frac{d}{dx}xx^{n-1} = x^{n-1} + \frac{d}{dx}x^{n-1} = x^{n-1} + (n-1)x^{n-1} = nx^{n-1}$$

4. Again, applying the product rule gives

$$\left(\frac{1}{g}\right)' = \left(\frac{1}{x} \circ g\right)' = -\frac{g'}{g^2}$$

   and the quotient rule

$$\left(\frac{f}{g}\right)' = \frac{f'}{g} + f\left(\frac{1}{g}\right)' = \frac{f'}{g} - \frac{fg'}{g^2} = \frac{gf' - fg'}{g^2}$$

5. Also $x = f \circ f^{-1}$ and $1 = (f^{-1})'f' \circ f^{-1}$. Thus

$$(f^{-1})' = \frac{1}{f' \circ f^{-1}}$$

   where defined. Especially for $x \neq 0$

$$\log'(x) = \frac{1}{\exp'(\log(x))} = \frac{1}{x}$$

6. $(1-q)(1 + q + q^2 + \cdots + q^n) = 1 - q + q - q^2 + q^2 - q^3 + \cdots + q^{n+1}$ gives

$$\sum_{k=0}^{n} q^k = \frac{1 - q^{n+1}}{1 - q} \quad \text{and} \quad \sum_{k=m}^{n} q^k = \frac{q^m - q^{n+1}}{1 - q}$$

Lemma 1.2 (Exponential Function).

1. It is
$$\exp(x + y) = \exp(x)\exp(y)$$

Hence
$$\exp(0) = 1$$
$$\exp(-x) = \exp(x)^{-1}$$
$$\exp(nx) = \exp(x)^{n}$$

2. For the derivative
$$\exp'(x) = \sum_{k=0}^{\infty} \frac{1}{k!}(x^k)' = \sum_{k=0}^{\infty} \frac{1}{k!}kx^{k-1} = \sum_{k=1}^{\infty} \frac{1}{(k-1)!}x^{k-1} = \exp(x)$$

Lemma 1.3 (Sinus and Cosinus)**.**

1. Sinus and Cosinus power series
$$\cos(x) = \sum_{k=0}^{\infty} \frac{(-1)^k}{2k!}x^{2k}$$
$$\sin(x) = \sum_{k=0}^{\infty} \frac{(-1)^k}{(2k+1)!}x^{2k+1}$$

2. Symmetry
$$\cos(-x) = \sum_{k=0}^{\infty} \frac{(-1)^k}{2k!}(-x)^{2k} = \cos(x)$$
$$\sin(x) = \sum_{k=0}^{\infty} \frac{(-1)^k}{(2k+1)!}(-x)^{2k+1} = -\sin(x)$$

3. Derivatives
$$\cos'(x) = \sum_{k=1}^{\infty} \frac{(-1)^k}{(2k-1)!}x^{2k-1} = \sum_{k=0}^{\infty} \frac{(-1)^{k+1}}{(2k+1)!}x^{2k+1} = -\sin(x)$$
$$\sin'(x) = \sum_{k=0}^{\infty} \frac{(-1)^k}{2k!}x^{2k} = \cos(x)$$

Theorem 1.4 (Fermat Stationary Point)**.** *Let $\Omega \subseteq \mathbb{R}$ be open and $f \in C^1(\Omega)$. If $x^* \in \Omega$ is local extremum then $f'(x^*) = 0$.*

Proof. Assume x^* is the minimum of f in Ω and let $f(x^*) > 0$. Since $f \in C^1(\Omega)$ there exist $\varepsilon, \delta > 0$, so that for $|h| \leq \varepsilon$
$$\frac{f(x^* + h) - f(x^*)}{h} > \delta$$
Pick a negative $h \in [-\varepsilon, 0)$. Then
$$f(x^* + h) < f(x^*) + \delta h < f(x^*)$$
and x^* cannot be the minimum. Analog for maximum with a positive h, then apply to $-f$. $\qquad\square$

Theorem 1.5 (Rolle)**.** *Let $f \in C[a, b]$ with $f(a) = f(b)$. If f is differentiable in (a, b) then there exists a $\xi \in (a, b)$ with $f'(\xi) = 0$.*

Proof. Assume f is not constant. Since $[a, b]$ is compact there exists either a global minimum or maximum $\xi \in (a, b)$ and Theorem 1.4 can be applied. □

Theorem 1.6 (Mean Value)**.** *Let $f \in C[a, b]$ be differentiable in (a, b). Then there exists a $\xi \in (a, b)$ with*

$$f'(\xi) = \frac{f(b) - f(a)}{b - a}$$

Proof. Apply Theorem 1.5 to

$$g(x) = f(x) - \frac{f(b) - f(a)}{b - a}(x - a)$$

□

Remark 1.7.

1. *More generally choose any $\varphi \in C^1[a, b]$ with $\varphi(a) = 0$ and $\varphi(b) = f(b) - f(a)$. Set $g(x) = f(x) - \varphi(x)$ to see there is a $\xi \in (a, b)$ with $f'(\xi) = \varphi'(\xi)$.*

2. *Let f be differentiable in (a, b) with $f' = 0$. For $x, y \in (a, b)$*

$$0 = f'(\xi) = \frac{f(y) - f(x)}{y - x}$$

   and f is a constant.

3. *Another useful generalization: let $\Omega \subseteq \mathbb{R}^n$ be open and $f \in C^1(\Omega)$. For $x, y \in \Omega$ define $\varphi(t) = f(tx + (1 - t)y)$ and apply the chain rule for differentiation*

$$\varphi'(\xi) = \nabla f(\xi x + (1 - \xi)y)^T (x - y) = f(x) - f(y)$$

4. *The Cauchy Schwarz inequality then yields*

$$\|f(x) - f(y)\| \leq \|\nabla f(\xi x + (1 - \xi)y)\| \|(x - y)\|$$

Theorem 1.8 (Differentiation Theorem)**.** *Let $f \in C[a, b]$ and define*

$$F(x) = \int_a^x f(t)\, dt$$

Then $F \in C^1[a, b]$ with $F'(x) = f(x)$ for $x \in [a, b]$.

Proof. Applying the Mean Value Theorem of Integration gives

$$F(x + h) - F(x) = \int_x^{x+h} f(t)\, dt = f(\xi)h$$

for some $\xi \in (x, x + h)$. □

Theorem 1.9 (Fundamental Theorem of Calculus)**.** *Let $F \in C^1[a,b]$ with $F' = f$ Then*

$$F(b) - F(a) = \int_a^b f(t)\, dt$$

Lemma 1.10 (Integration by Substitution)**.** *Let $I \subseteq \mathbb{R}$ be an interval and $f \in C(I)$. For $\varphi \in C([a,b], I)$ it follows*

$$\int_{\varphi(a)}^{\varphi(b)} f(x)\, dx = \int_a^b f(\varphi(t))\varphi'(t)\, dt$$

Proof. Let $F \in C^1(I)$ with $F' = f$. Then the chain rule for differentiation yields

$$\int_{\varphi(a)}^{\varphi(b)} f(x)\, dx = F(\varphi(b)) - F(\varphi(a))$$

$$= F \circ \varphi(b) - F \circ \varphi(a)$$

$$= \int_a^b (F \circ \varphi)'(t)\, dt$$

$$= \int_a^b f(\varphi(t))\varphi'(t)\, dt$$

$\square$

Examples 1.11.

1. For $\varphi(x) = x^2 + 1$ it is $\varphi(0) = 1$ and $\varphi(2) = 5$. Thus

$$\int_0^2 x \cos(x^2 + 1)\, dx = \frac{1}{2} \int_0^2 2x \cos(x^2 + 1)\, dx = \frac{1}{2} \int_1^5 \cos(t)\, dt = \frac{1}{2}(\sin(5) - \sin(1))$$

2. Consider $\varphi(x) = \sin(x)$ where $\varphi(0) = 0$ and $\varphi(\pi/2) = 1$. Since $\cos(t) = \sqrt{1 - \sin^2(t)}$ it follows

$$\int_0^1 \sqrt{1 - x^2}\, dx = \int_{\cos(0)}^{\cos(\pi/2)} \sqrt{1 - x^2}\, dx = \int_0^{\pi/2} \sqrt{1 - \sin^2(t)} \cos(t)\, dt = \int_0^{\pi/2} \cos^2(t)\, dt$$

3. Let $f \in C[a,b]$ and $\varphi(x) = a + t(b - a)$. Then

$$\int_a^b f(x)\, dx = (b - a) \int_0^1 f(a + t(b - a))\, dt$$

4. Let $f(x) = x^n$ and $\varphi(x) = t^m$. As expected

$$\int_0^1 x^n\, dx = \int_0^1 t^{nm} m t^{m-1}\, dt = m \int_0^1 t^{m(n+1)-1}\, dt = \left[\frac{m}{m(n+1)} t^{m(n+1)} \right]_0^1 = \frac{1}{n+1}$$

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1.2 Directional Derivative and Gradients

Lemma 1.12 (Directional Derivative). *Let $\Omega \subseteq \mathbb{R}^n$ be open and $f \in C^1(\Omega)$. Then*

$$\frac{\partial f}{\partial d}(x) = \nabla f(x)^T d$$

for any $d \in \mathbb{R}^n$.

Proof. Let $\varphi(t) = f(x + td)$. Then $\varphi \in C^1[-\varepsilon, \varepsilon]$ for some $\varepsilon > 0$ and the chain rule yields

$$\varphi'(t) = \nabla f(x + td)^T d$$

Hence

$$\varphi'(0) = \lim_{t \to 0} \frac{\varphi(x + td) - \varphi(0)}{t} = \lim_{t \to 0} \frac{f(x + td) - f(x)}{t} = \nabla f(x)^T d$$

$\square$

Remarks 1.13.

1. *Note that by definition the directional derivative is invariant under multiplication with any $\lambda \neq 0$.*

2. *A similar proposition holds under the weaker assumption that d is a only feasable direction for f in x*

3. *For $d = \nabla f(x)/\|\nabla f(x)\|$ it follows that*

$$\frac{\partial f}{\partial d}(x) = \|\nabla f(x)\| > 0$$

and for any other $d \in \mathbb{R}^n$ with $\|d\| = 1$ the Cauchy Schwarz inequality yields

$$|\frac{\partial f}{\partial d}(x)| = |\nabla f(x)^T d| \leq \|\nabla f(x)\|\|d\| = \|\nabla f(x)\|$$

Hence $\nabla f(x)$ is the direction of the greatest ascent and respectively, $-\nabla f(x)$ is the direction of the greatest descent.

Theorem 1.14 (First Order Necessary Condition). *Let $\Omega \subseteq \mathbb{R}^n$ be open and $f \in C^1(\Omega)$. If $x^* \in \Omega$ is a local minimzer then $\nabla f(x^*) = 0$.*

Proof. Let $h \in \mathbb{R}^n$ and $\delta > 0$ so that $x^* + th \in \Omega$ for all $t \in (-\delta, \delta)$. Then 0 is local minimizer for $\varphi(t) = f(x^* + th)$ and

$$\varphi'(0) = \nabla f(x^*)^T h = 0$$

Now let $h = \nabla f(x^*)$. $\square$

Theorem 1.15 (Banach Fixed-Point Theorem). *Let X be a Banach space and $f \in C(X, X)$ a contraction*

$$\|f(x) - f(y)\| \leq q\|x - y\| \text{ for all } x, y \in X$$

for some $0 < q < 1$. Then there exists a unique fix point $x^* \in X$ with

$$f(x^*) = x^*$$

Furthermore for any $x_0 \in X$ the sequence defined by

$$x_{n+1} = f(x_n)$$

converges aganist x^*.

Proof. Since $\|x_{n+1} - x_n\| = \|f(x_n) - f(x_{n-1})\| \le q\|x_n - x_{n-1}\|$ it follows, that

$$\|x_{n+1} - x_n\| \le q^n\|x_1 - x_0\|$$

Furthermore

$$\|x_n - x_m\| \le \sum_{k=m}^{n} q^k\|x_1 - x_0\| = \frac{q^m - q^{n+1}}{1 - q}\|x_1 - x_0\|$$

and (x_n) is a Cauchy sequence. For its limit x^* we have

$$x^* = \lim_{n \to \infty} x_{n+1} = \lim_{n \to \infty} f(x_n) = f(x^*)$$

For any other $y^* \in X$ with $f(y^*) = y^*$ it follows, that

$$\|x^* - y^*\| = \|f(x^*) - f(y^*)\| \le q\|x^* - y^*\|$$

and therefore $x^* = y^*$.

$\square$

2   Nonlinear Optimization

2.1   Minimization without Constraints

Lemma 2.1 (Gradient Inequality). *Let $M \subseteq \mathbb{R}^n$ be a convex set and $f \in C^1(M)$. Then f is convex if and only if*

$$f(x) \geq f(y) + \nabla f(y)^T (x - y)$$

for all $x, y \in M$.

Proof. Let f be convex and $x, y \in M$. For $0 \leq \lambda \leq 1$ we have

$$f(\lambda x + (1 - \lambda)y) \leq \lambda f(x) + (1 - \lambda)f(y) = \lambda f(x) - \lambda f(y) + f(y)$$

and

$$f(x) - f(y) \geq \frac{f(\lambda x + (1 - \lambda)y) - f(y)}{\lambda} = \frac{f(y + \lambda(x - y)) - f(y)}{\lambda}$$

For $d = x - y$ and $\lambda \to 0$ the term on the right converges to the direction derivative of f in d

$$\frac{\partial f}{\partial d}(y) = \nabla f(y)^T d = \nabla f(y)^T (x - y)$$

Now let $x, y \in M$ and $0 \leq \lambda \leq 1$. For $z = \lambda x + (1 - \lambda)y \in M$ it follows that

$$\lambda f(x) \geq \lambda f(z) + \lambda \nabla f(z)^T (x - z)$$
$$(1 - \lambda)f(y) \geq (1 - \lambda)f(z) + (1 - \lambda)\nabla f(z)^T (y - z)$$

Adding the two inequalities gives

$$\begin{aligned}
\lambda f(x) + (1 - \lambda)f(y) &\geq f(z) + \nabla f(z)^T (\lambda x - \lambda z + (1 - \lambda)y - (1 - \lambda)z) \\
&= f(z) + \nabla f(z)^T (\lambda x + (1 - \lambda)y - z) \\
&= f(z)
\end{aligned}$$

$\square$

Exercise 2.2 (Facility Locations). *The facilities are located at:*

$$(3, 0), (0, -3), (1, 4)$$

Proof. Let

$$\begin{aligned}
f(x) &= (x - 3)^2 + y^2 + x^2 + (y + 3)^2 + (x - 1)^2 + (y - 4)^2 \\
&= x^2 - 6x + 9 + y^2 + x^2 + y^2 + 6y + 9 + x^2 - 2x + 1 + y^2 - 8y + 16 \\
&= 3x^2 + 3y^2 - 8x - 2y + 35
\end{aligned}$$

Then

$$\nabla f(x, y) = (6x - 8, 6y - 2) \text{ and } \nabla^2 f(x, y) = \begin{pmatrix} 6 & 0 \\ 0 & 6 \end{pmatrix} > 0$$

Hence $(4/3, 1/3)$ is the gobal minimum. $\square$

Exercise 2.3 (Convex Functions). *The sum of convex functions is convex.*

Proof. Let $x, y \in M$. Since $\alpha_i > 0$ we have

$$f(\lambda x + (1 - \lambda)y) = \sum_{i=1}^{m} \alpha_i f_i(\lambda x + (1 - \lambda)y)$$

$$\leq \sum_{i=1}^{m} \alpha_i \lambda f_i(x) + \sum_{i=1}^{m} \alpha_i (1 - \lambda)f_i(y) = \lambda f(x) + (1 - \lambda)f(y)$$

Let $f(x) = x^2$. Then $-f$ is not convex, e.g. $x = 1, y = -1$ and $\lambda = 0.5$.

Exercise 2.4 (Solution of Quadratic Inequality). *Let*

$$f(x) = x^T A x + b^T x + c$$

Proof. The product rule gives

$$\nabla f(x) = x^T A + Ax + b = (A^T + A)x + b = 2Ax + b$$

Thus $\nabla^2 f(x) = 2A > 0$ and f is convex. Hence the level set Γ_{-c} is convex. Since the intersection of convex sets is convex $\Gamma_{-c} \cap \{x \in \mathbb{R}^n : g^T x + h = 0\}$ is convex, too.

Exercise 2.5 (Line Search on Compact Convex Sets). *Let $S \subset \mathbb{R}^n$ be compact and convex. Furthermore let $f \in C^1(S)$ be convex, $x \in S$ and $d \in \mathbb{R}^n$ a descent direction of f in x with $\nabla f(x)^T d < 0$.*

Proof. If $x + \lambda^* d$ is an optimal solution then $\nabla f(x + \lambda^* d)^T d = 0$ according to Theorem 1.14. Let $\nabla f(x + \lambda^* d)^T d = 0$. Then Lemma 2.1 gives

$$f(x + \lambda d) \geq f(x + \lambda^* d) + (\lambda - \lambda^*)\nabla f(x + \lambda^* d)^T d = f(x + \lambda^* d)$$

and $x + \lambda^* d$ is an optimal solution.

Exercise 2.6 (Steepest Descent). *Let*

$$f(x) = \frac{1}{2}x^T A x + b^T x + c$$

where A is symmetrical and positive definite.

Proof. Since $\nabla f(x) = Ax + b$ and $\nabla^2 f(x) = A > 0$ it follows $x^* = -A^{-1}b$. Let v be eigenvector with $Av = \mu v$. For $x_0 = x^* + \theta v$ we have

$$\nabla f(x_0) = Ax^* + \mu\theta v + b = \mu\theta v$$

and for $\lambda \geq 0$

$$\arg\min\{f(x_0 - \lambda\nabla f(x_0)\} = \arg\min\{f(x^* + \theta v - \lambda\mu\theta v)\} = \mu^{-1}$$

Thus

$$x_1 = x_0 - \mu^{-1}\nabla f(x_0) = x^* + \theta v - \mu^{-1}\mu\theta v = x^*$$

and $\nabla f(x_1) = 0$. Hence the algorithm stops after the first iteration. Now let

$$x_0 = x^* + \sum_{i=0}^{m} \theta_i v_i$$

for orthogonal eigenvectors with $Av_i = \mu_i$ and $m \leq n$. Then

$$\nabla f(x_0) = Ax^* + \sum_{i=0}^{m} \mu_i \theta_i v_i + b = \sum_{i=0}^{m} \mu_i \theta_i v_i$$

and

$$x_1 = x_0 - \lambda \sum_{i=0}^{m} \mu_i \theta_i v_i = x^* + \sum_{i=0}^{m} \theta_i v_i - \lambda \sum_{i=0}^{m} \mu_i \theta_i v_i = x^* + \sum_{i=0}^{m} (1 - \lambda \mu_i) \theta_i v_i$$

Since x^* is the minimum we have $\nabla f(x_1) = 0$ iff $\lambda = \mu^{-1}$ for all $0 \leq i \leq m$. $\qquad\square$

2.2 One Dimensional Minimization and Direct Search

Definition 2.7 (Unimodal Function). *A function $f : [a, b] \to \mathbb{R}$ is called unimodal if there exists a $\xi \in [a, b]$, so that f is strictly decreasing in $[a, \xi]$ and strictly increasing in $[\xi, b]$.*

In fact ξ is the unique minimum of f in $[a, b]$. According to the definition, for $a \leq x < y \leq b$ we have

$$f(x) > f(y) \text{ for } x, y \in [a, \xi) \text{ and } f(x) < f(y) \text{ for } x, y \in (\xi, b]$$

Thus

$$\xi \in [a, y] \text{ if } f(x) < f(y) \text{ and } \xi \in [x, b] \text{ if } f(x) \geq f(y)$$

Consider now a symmetrical partioning of the interval $[0, 1]$ where two consecutive partionings hold the same ratio respectively:

$$\sigma = 1 - \tau \text{ and } \frac{1}{\tau} = \frac{\tau}{\sigma}$$

Then $1 - \tau = \tau^2$ and solving the quadratic equation $\tau^2 + \tau = 1$ yields

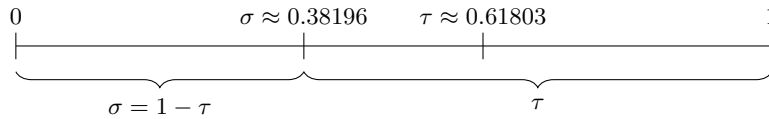$$\tau = \frac{\sqrt{5} - 1}{2} \approx 0.61803$$

Figure 1: Golden Section

Let now $[a_0, b_0] = [a, b]$ and define

$$[a_{k+1}, b_{k+1}] = \begin{cases} [a_k, y_k] & \text{if } f(x_k) < f(y_k) \\ [x_k, b_k] & \text{if } f(x_k) \geq f(y_k) \end{cases}$$

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where

$$x_k = b_k - \tau(b_k - a_k)$$
$$y_k = a_k + \tau(b_k - a_k)$$

It follows that $[a_k, b_k] \supset [a_{k+1}, b_{k+1}]$ is a decreasing series of intervals with

$$(b_{k+1} - a_{k+1}) = \tau(b_k - a_k)$$

where the interval converges to ξ. This leads to the following algorithm:

Algorithm 2.8 (Golden Section Search)**.**

```python
import math


def golden_section_search(f, I, eps=0.00001):
    t = 0.5 * (math.sqrt(5) - 1)
    a, b = I
    while abs(b - a) > eps:
        x, y = b - t * (b - a), a + t * (b - a)
        if f(x) > f(y):
            a = x
        else:
            b = y
    return (a + b) / 2
```

Algorithm 2.9 (Steepest Descent)**.**

Let $f \in C^1(\mathbb{R}^n)$ and $x_0 \in \mathbb{R}^n$. For $0 < \alpha \leq \beta < 1$ and $\gamma < 1$ let

Exercise 2.10 (Surprising Convergence)**.** *Example for $f \in C^2(\mathbb{R})$ with a sequence of strict local minima converging to a strict local maximum.*

Proof. Let $f \in C[a, b]$ and $\xi \in (a, b)$ so that f is strictly increasing in $(a, \xi]$ and strictly decreasing in $[\xi, b)$. Define

$$g(x) = \int_{\xi - x}^{\xi + x} f(t)\, dt$$

$\square$

3 Functional Analysis

3.1 The Hahn-Banach Theorem

Remarks 3.1.

1. Let $f = u + \mathrm{i}v$ be a complex valued linear functional. Then $f(\mathrm{i}x) = \mathrm{i}f(x)$ gives

$$u(\mathrm{i}x) = -v(x) \text{ and } v(\mathrm{i}x) = \mathrm{i}u(x)$$

   Thus

$$f(x) = u(x) - \mathrm{i}u(\mathrm{i}x)$$

2. Assume $|u(x)| \leq C\|x\|$ and choose $t \in \mathbb{R}$ so that $|f(x)| = e^{\mathrm{i}t}f(x)$. Then

$$|f(x)| = e^{\mathrm{i}t}f(x) = f(e^{\mathrm{i}t}x) = u(e^{\mathrm{i}t}x) \leq C\|e^{\mathrm{i}t}x\| = C\|x\|$$

3. See also the *Cauchy-Rieman Eqautions 4.10*.

Theorem 3.2 (Hahn-Banach Theorem). *Let E be a normed space and F a linear subspace. For any linear functional $f \in F^*$ there exists a $\tilde{f} \in E^*$ so that*

$$\tilde{f}|_F = f \text{ and } \|\tilde{f}\| = \|f\|$$

Proof. Assume that $f : E \to \mathbb{R}$ is a real valued linear functional and $x' \in E \setminus F$. Consider the subspace $F' = \{x + \lambda x' | x \in F\}$ and define
$$f'(x + \lambda x') = f(x) + \lambda r$$

for some $r \in \mathbb{R}$. As sum of two linear functionals f' again is a linear functional with $\|f'\| \geq \|f\|$.

Using Zorn's Lemma allows the extension to E while the remarks above can be used to extend to complex valued linear functionals. $\qquad \square$

3.2 The Riesz Representation Theorem

Theorem 3.3 (The Riesz Representation Theorem). *Let X be a locally compact Hausdorff space and $L \in C_c(X)$ a positive linear functional. Then ther exists a σ-algebra M in X and a unique positive measure μ on M so that*

$$L(f) = \int_X f \, d\mu$$

for every $f \in C_c(X)$.

4 The Road to Reality

4.1 Hyperbolic Geometry

The ratio between the area A and A' of two similar shapes is given by

$$A' = k^2 A$$

Theorem 4.1 (Pythagoras).
$$a^2 + b^2 = c^2$$

Proof. Let A, B and C be the areas of the three triangles respectively. All triangles are similar, hence

$$B = \frac{b^2}{a^2} A \text{ and } C = \frac{c^2}{b^2} B$$

Since $A + B = C$ it follows that

$$a^2 + b^2 = \frac{b^2 A}{B} + b^2 = \frac{b^2(A+B)}{B} = \frac{b^2 C}{B} = c^2$$

$\square$

Lemma 4.2 (Conformal and Projective Representation). *The mapping from conformal and projective representation of any point is given by the radial expansion of the following factor*

$$\frac{2R}{R^2 + r^2}$$

Proof. For any point the distance from the origin with regard to the two representations is given by

$$\log \frac{R+r}{R-r} = \frac{1}{2} \log \frac{R+r'}{R-r'} = \log \frac{(R+r')^2}{(R-r')^2}$$

This gives
$$(R-r)^2 (R+r') = (R+r)^2 (R-r') \text{ and } -4R^2 r + 2R^2 r' + 2r^2 r' = 0$$

Hence
$$r' = \frac{2R^2}{R^2 + r^2} r$$

$\square$

4.2   Complex Numbers

Lemma 4.3 (Basic Formulas)**.**

1. *It is*
$$(a + \mathrm{i}b)(c + \mathrm{i}d) = (ac - bd) + \mathrm{i}(ad + bc)$$

2. *Thus*
$$(a + \mathrm{i}b)^2 = (a^2 - b^2) + \mathrm{i}2ab$$

   and
$$(a + \mathrm{i}b)(a - \mathrm{i}b) = a^2 + \mathrm{i}ab - \mathrm{i}ab - \mathrm{i}^2 b^2 = a^2 + b^2$$

3. *Hence*
$$\frac{a + \mathrm{i}b}{c + \mathrm{i}d} = \frac{(a + \mathrm{i}b)(c - \mathrm{i}d)}{c^2 + d^2} = \frac{ac + bd}{c^2 + d^2} + \mathrm{i}\frac{bc - ad}{c^2 + d^2}$$

4. *For*
$$z = \sqrt{\frac{1}{2}(a + \sqrt{a^2 + b^2})} + \mathrm{i}\sqrt{\frac{1}{2}(-a + \sqrt{a^2 + b^2})}$$

   it follows
$$z^2 = \frac{1}{2}(a + \sqrt{a^2 + b^2}) - \frac{1}{2}(-a + \sqrt{a^2 + b^2}) + \mathrm{i}2\sqrt{\frac{1}{4}(\sqrt{a^2 + b^2}^2) - a^2} = a + \mathrm{i}b$$

Lemma 4.4 (Binomial Theorem)**.**

1. *For the binomial coefficient Pascal's identity holds*
$$\binom{n}{k-1} + \binom{n}{k} = \binom{n+1}{k}$$

2. *The following equation states the binomial identity*
$$(a + b)^n = \sum_{k=0}^{n} \binom{n}{k} a^k b^{n-k} = \sum_{k=0}^{n} \binom{n}{k} a^{n-k} b^k$$

3. *For $a = 1$ follows*
$$(1 + x)^n = \sum_{k=0}^{n} \binom{n}{k} x^k$$

Proof. It is
$$\binom{n}{k} + \binom{n}{k-1} = \frac{n!}{k!(n-k)!} + \frac{n!}{(k-1)!(n-k+1)!} = \frac{n!(n+1-k) + n!k!}{k!(n+1-k)!} = \binom{n+1}{k}$$

Furthermore by using induction

$$(a+b)^{n+1} = \sum_{k=0}^{n} \binom{n}{k} a^{k+1} b^{n-k} + \sum_{k=0}^{n} \binom{n}{k} a^k b^{n+1-k}$$

$$= \sum_{k=1}^{n+1} \binom{n}{k-1} a^k b^{n+1-k} + \sum_{k=0}^{n} \binom{n}{k} a^k b^{n+1-k}$$

$$= \sum_{k=0}^{n+1} \binom{n+1}{k} a^k b^{n+1-k}$$

$\square$

4.3   Exponential Function and Logarithms

Exercise 4.5 (Exponential Function). *The Cauchy product yields*

$$\sum_{n=0}^{\infty} a_n \sum_{n=0}^{\infty} b_n = \sum_{n=0}^{\infty} \sum_{k=0}^{n} a_k b_{n-k}$$

if at least one of the series is absolutely convergent. Hence

$$\sum_{n=0}^{\infty} \frac{1}{n!} z^n \sum_{n=0}^{\infty} \frac{1}{n!} w^n = \sum_{n=0}^{\infty} \sum_{k=0}^{n} \frac{1}{k!} z^k \frac{1}{(n-k)!} w^{n-k}$$

$$= \sum_{n=0}^{\infty} \frac{1}{n!} \sum_{k=0}^{n} \binom{n}{k} z^k w^{n-k}$$

$$= \sum_{n=0}^{\infty} \frac{1}{n!} (z+w)^n$$

Let $t \in \mathbb{R}$. Then

$$e^{it} = \sum_{k=0}^{\infty} \frac{1}{k!} (it)^k$$

$$= \sum_{k=0}^{\infty} \frac{1}{2k!} (it)^{2k} + \sum_{k=0}^{\infty} \frac{1}{(2k+1)!} (it)^{2k+1}$$

$$= \sum_{k=0}^{\infty} \frac{(-1)^k}{2k!} t^{2k} + i \sum_{k=0}^{\infty} \frac{(-1)^k}{(2k+1)!} t^{2k+1}$$

$$= \cos t + i \sin t$$

More generally for $z = \log r + it$

$$e^z = e^{\log r + it} = r e^{it} = r(\cos t + i \sin t)$$

For $r = 1$ and $t = 2\pi$ this yields

$$e^{2\pi i} = \cos 2\pi + i \sin 2\pi = 1$$

and for $t = 2\pi$ we get

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Lemma 4.6 (Euler Equation)**.**

$$e^{\pi i} + 1 = 0$$

Exercise 4.7.

1. *If* $e^z = w$ *then* $z + \pi i$ *is a logarithm to* $-w$: $e^{z+\pi i} = e^z e^{\pi i} = -e^z = -w$.

2. *Since* $e^{i(s+t)} = e^{is}e^{it}$ *it follows*

$$\cos(s+t) + i\sin(s+t) = (\cos s + i\sin s)(\cos t + i\sin t)$$
$$= \cos s \cos t - \sin s \sin t + i(\cos s \sin t + \sin s \cos t)$$

   Hence

$$\cos(s+t) = \cos s \cos t - \sin s \sin t$$
$$\sin(s+t) = \cos s \sin t + \sin s \cos t$$

3. *It is* $e^{3it} = \left(e^{it}\right)^3$ *and thus*

$$\cos 3t + i\sin 3t = (\cos t + i\sin t)^3 = \cos^3 t - 3\cos t \sin^2(t) + i(\cos^2 t \sin t - \sin^3 t)$$

4. *Fun facts*

$$e^{1-4\pi^2} = e^{1+(2i\pi)^2} = ee^{2\pi i}e^{2\pi i} = e$$

   and $i = e^{i\pi/2}$ *gives*

$$i^i = e^{i\log i} = e^{ii\pi/2} = e^{-\pi/2} \in \mathbb{R}$$

4.4 Complex Analysis

Definition 4.8 (holomorphic Function)**.** *Let* $\Omega \subseteq \mathbb{C}$ *be open. A function* $f : \Omega \to \mathbb{C}$ *is called* differentiable *at* $z \in \Omega$ *if the limit*

$$f'(z) = \lim_{h \to 0} \frac{f(z+h) - f(z)}{h}$$

exists. f *is called* holomorphic *on* Ω *if* f *is complex differentiable at all points of* Ω *and* $f' : \Omega \to \mathbb{C}$ *is called the* derivative *of* f.

Remarks 4.9.

1. f *is differentiable at* $z_0 \in \Omega$ *iff the limit*

$$f'(z_0) = \lim_{z \to z_0} \frac{f(z) - f(z_0)}{z - z_0}$$

   exists

2. *If* f *is differentiable at* $z_0 \in \Omega$ *and* $\varepsilon > 0$ *then there exists a small enough environment of* z_0, *so that*

$$|f(z) - f(z_0) - f'(z_0)(z - z_0)| < \varepsilon |z - z_0|$$

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Theorem 4.10 (Cauchy Riemann Equations). *Let $f = u + iv$ be holomorphic. Then f satisfies the Cauchy Riemann equations*

$$\frac{\partial u}{\partial x} = \frac{\partial v}{\partial y}$$

$$\frac{\partial u}{\partial y} = -\frac{\partial v}{\partial x}$$

Proof. For $h \in \mathbb{R}$ follows

$$\lim_{h \to 0} \frac{f(z+h) - f(z)}{h} = \frac{\partial u}{\partial x}(z) + i\frac{\partial v}{\partial x}(z)$$

and

$$\lim_{h \to 0} \frac{f(z+ih) - f(z)}{ih} = \frac{\partial u}{i\partial y}(z) + \frac{\partial v}{\partial y}(z) = \frac{\partial v}{\partial y}(z) - i\frac{\partial u}{\partial y}(z)$$

$\square$

Examples 4.11.

1. Let $f(z) = z^3$. Then $u(x,y) + iv(x,y) = x^3 - 3xy^2 + i(3x^2y - y^3)$ and as expected

$$\frac{\partial u}{\partial x}(x,y) = x^3 - 3y^2 \quad and \quad \frac{\partial u}{\partial y}(x,y) = -6xy$$

$$\frac{\partial v}{\partial x}(x,y) = 6xy \quad and \quad \frac{\partial v}{\partial y}(x,y) = x^3 - 3y^2$$

Lemma 4.12. *Let $D \subseteq \mathbb{C}$ be connected. For arbitrary $z, w \in D$ there exists a polygonal path from z to w.*

Proof. For any path from z to w the image is compact, which can be used to define a finite subcover of disks. Use the center points to define the polygonal path. $\square$

Lemma 4.13. *Let $\gamma : [a, b] \to \mathbb{C}$ a smooth path, $\psi : [c, d] \to [a, b]$ a smooth and increasing bijection and f continous.*

$$\int_\gamma f(z) \, dz = \int_{\gamma \circ \psi} f(z) \, dz$$

Proof. It is

$$\int_{\gamma \circ \psi} f(z) \, dz = \int_c^d f(\gamma \circ \psi(t))(\gamma \circ \psi)'(t)) \, dt$$

$$= \int_{\psi(a)}^{\psi(b)} f(\gamma(\psi(t)))\gamma'(\psi(t))\psi'(t) \, dt$$

$$= \int_a^b f(\gamma(s))\gamma'(s) \, ds = \int_\gamma f(z) \, dz$$

$\square$

Lemma 4.14. *For a smooth path* $\gamma : [a, b] \to \mathbb{C}$ *define* $-\gamma(t) = a + b - t$. *Then*

$$\int_{-\gamma} f(z)\, dz = -\int_{\gamma} f(z)\, dz$$

Proof. Using integration by substitution

$$\int_{-\gamma} f(z)\, dz = -\int_a^b f(\gamma(a + b - t))\gamma'(a + b - t))\, dt = \int_b^a f(\gamma(s))\gamma'(s)\, ds = -\int_{\gamma} f(z)\, dz$$

$\square$

In order to use the results from real calculus recall the fact, that for every $z \in \mathbb{C}$ there exists a $t \in [0, 2\pi]$, so that $z = |z|e^{it}$ and hence $|z| = ze^{-it}$.

Lemma 4.15. *Let* $f \in C[a, b]$. *Then*

$$\left| \int_a^b f(x)\, dx \right| \le \int_a^b |f(x)|\, dx$$

Proof. Using the estimation for integrals from real calculus

$$\left| \int_a^b f(x)\, dx \right| = e^{-it} \int_a^b f(x)\, dx \le \int_a^b |e^{-it} f(x)|\, dx = \int_a^b |f(x)|\, dx$$

$\square$

Let $\gamma : [a, b] \to \mathbb{C}$ be a smooth path and $a = t_0 < t_1 < \cdots < t_n = b$ a partioning of $[a, b]$. Then

$$\sum_{k=1}^n |\gamma(t_k) - \gamma(t_{k-1})| = \sum_{k=1}^n \left| \frac{\gamma(t_k) - \gamma(t_{k-1})}{t_k - t_{k-1}} \right| (t_k - t_{k-1}) = \sum_{k=1}^n |\gamma'(\xi_k)|(t_k - t_{k-1})$$

yields a reasonable approximation of the length of the path. Hence

Definition 4.16. *For a smooth path* $\gamma : [a, b] \to \mathbb{C}$

$$L(\gamma) = \int_a^b |\gamma'(t)|\, dt$$

is called the length of γ.

Lemma 4.17 (Estimation Lemma). *Let* $\gamma : [a, b] \to \mathbb{C}$ *be a smooth path. Then*

$$\left| \int_{\gamma} f(z)\, dz \right| \le L(\gamma) \max_{\gamma[a,b]} f$$

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Proof. Using the definition above

$$\left| \int_\gamma f(z)\,dz \right| = \left| \int_a^b f(\gamma(t))\gamma'(t)\,dt \right| \le \int_a^b |f(\gamma(t))\gamma'(t)|\,dt \le \max_{\gamma[a.b]} f \int_a^b |\gamma'(t)|\,dt$$

$\square$

Examples 4.18.

1. Let $\gamma(t) = t + \mathrm{i}t$. Then

$$\int_\gamma z^2\,dz = \int_0^1 (t + \mathrm{i}t)^2(1 + \mathrm{i})\,dt = (1 + \mathrm{i})\int_0^1 2\mathrm{i}t^2\,dt = \left[(-2 + 2\mathrm{i})t^2\right]_0^1 = -\frac{2}{3} + \mathrm{i}\frac{2}{3}$$

2. For $\gamma(t) = t^2 + \mathrm{i}t$

$$\int_\gamma z^2\,dz = \int_0^1 (t^2 + \mathrm{i}t)^2(2 + \mathrm{i}t)\,dt = \int_0^1 (2t^5 - 4t^3) + \mathrm{i}(5t^4 - t^2 t)\,dt$$

$$= \left[\frac{1}{3}t^6 - t^4\right]_0^1 + \mathrm{i}\left[t^5 - \frac{1}{3}t^3\right]_0^1 = -\frac{2}{3} + \mathrm{i}\frac{2}{3}$$

3. And $\gamma(t) = \mathrm{i} + \mathrm{e}^{\mathrm{i}t}$

$$\int_\gamma z^2\,dz = \int_{3/2\pi}^{2\pi} (\mathrm{i} + \mathrm{e}^{\mathrm{i}t})^2 \mathrm{i}\mathrm{e}^{\mathrm{i}t}\,dt = \int_{3/2\pi}^{2\pi} (-1 + 2\mathrm{i}\mathrm{e}^{\mathrm{i}t} + \mathrm{e}^{2\mathrm{i}t})\mathrm{i}\mathrm{e}^{\mathrm{i}t}\,dt$$

$$= \int_{3/2\pi}^{2\pi} -\mathrm{i}\mathrm{e}^{\mathrm{i}t} - 2\mathrm{e}^{2\mathrm{i}t} + \mathrm{i}\mathrm{e}^{3\mathrm{i}t}\,dt = \left[-\mathrm{e}^{\mathrm{i}t} + \mathrm{i}\mathrm{e}^{2\mathrm{i}t} + \frac{1}{3}\mathrm{e}^{3\mathrm{i}t}\right]_{3/2\pi}^{2\pi}$$

$$= \left(-1 + \mathrm{i} + \frac{1}{3}\right) - \left(\mathrm{i} - \mathrm{i} + \frac{1}{3}\mathrm{i}\right) = -\frac{2}{3} + \mathrm{i}\frac{2}{3}$$

4. Let $\gamma(t) = \mathrm{e}^{\mathrm{i}t}$ and $k \ne -1$. Then

$$\int_\gamma z^k\,dz = \int_0^{2\pi} \mathrm{e}^{\mathrm{i}kt}\mathrm{i}\mathrm{e}^{\mathrm{i}t}\,dt = \int_0^{2\pi} \mathrm{i}\mathrm{e}^{\mathrm{i}(k+1)t}\,dt = 0$$

Theorem 4.19. *Let $D \subseteq \mathbb{C}$ be a connected domain and $f \in C(D)$. Then the following assertions are equivalent*

1. *f has an antiderivative*

2. *For every closd path γ*

$$\int_\gamma f(z)\,dz = 0$$

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Proof. Let $F' = f$. Since γ is closed

$$\int_\gamma f(z)\,dz = \int_a^b f(\gamma(t))\gamma'(t)\,dt = \int_a^b (F \circ \gamma)'(t)\,dt = F(\gamma(b)) - F(\gamma(a)) = 0$$

Now fix some arbitrary $a \in D$. For $z \in D$ let γ_z be a path from a to z and define

$$F(z) = \int_{\gamma_z} f(\zeta)\,d\zeta$$

This is well defined since the integral of f vanishes over each closed path. Moreover, since $\gamma_{z+h}+[z+h,z]-\gamma_z$ defines a closed path

$$F(z+h) - F(z) = \int_{\gamma_{z+h}} f(z)\,dz - \int_{\gamma_z} f(z)\,dz = \int_{[z,z+h]} f(z)\,dz = h\int_0^1 f(z+th)\,dt$$

Here the latter integral is continous at 0 with respect to h

$$\left| \int_0^1 f(z+th) - f(z)\,dt \right| \leq \int_0^1 |f(z+th) - f(z)|\,dt \leq \max_{t\in[0,1]} |f(z+th) - f(z)|$$

$\square$

Corollary 4.20. *The second assertion can be weakend to*

$$\int_{\partial\Delta} f(z)\,dz = 0$$

for every triangle $\Delta \subset D$, where e.g. D is convex or star shaped. Here the antiderivative can directly be defined as

$$F(z) = \int_{[a,z]} f(\zeta)\,d\zeta$$

similar to the real calculus approach. Note, that under this conditions f always has a local antiderivative.

Examples 4.21.

1. Let $z_0 \in \mathbb{C}$ and $\gamma(t) = z_0 + e^{it}$ for $t \in [0, 2\pi]$. Then

$$\int_\gamma \frac{1}{z - z_0}\,dz = \int_0^{2\pi} \frac{ie^{it}}{z_0 + e^{it} - z_0}\,dt = \int_0^{2\pi} i\,dt = 2\pi i$$

   and thus $1/(z - z_0)$ has no antiderivative on $\mathbb{C} \setminus \{z_0\}$

2. Let $z_0 \in \mathbb{C}$ and $z \in D = D_r(z_0)$. Applying *Theorem 4.19.* to ∂D and a small enough circle around z gives

$$\int_{\partial D} \frac{1}{\zeta - z}\,d\zeta = \int_{\partial D} \frac{1}{\zeta - z_0}\,d\zeta = 2\pi i$$

Theorem 4.22 (Goursat). *Let $\Omega \subseteq \mathbb{C}$ be open and f holomorphic on Ω. Then*

$$\int_{\partial \triangle} f(z)\, dz = 0$$

for every triangle $\triangle \subset \Omega$.

Proof. Choose a sequence of triangles $\triangle \supset \triangle_0 \supset \triangle_1 \cdots \supset \triangle_k$ as depicted. Since all the triangles are compact with a vainishing diameter there exists a unique $z_0 \in \Omega$ with $\bigcap \triangle_k = \{z_0\}$. Thus

$$\left| \int_{\partial \triangle} f(z)\, dz \right| \leq 4^k \left| \int_{\partial \triangle_k} f(z)\, dz \right| = 4^k \left| \int_{\partial \triangle_k} f(z) - f(z_0) - f'(z_0)(z - z_0)\, dz \right|$$

Furthermore $L(\partial \triangle) = 2^{-k} L(\partial \triangle_k)$ and

$$|z - z_0| < L(\partial \triangle_k) = 2^{-k} L(\partial \triangle)$$

for any $z \in \triangle_k$. Since f is holomorphic at z_0 for any given $\varepsilon > 0$ there exists a sufficiently large enough k, so that

$$\left| \int_{\partial \triangle} f(z)\, dz \right| \leq 4^k L(\partial \triangle_k) \max_{z \in \triangle_k} |f(z) - f(z_0) - f'(z_0)(z - z_0)|$$

$$\leq 4^k L(\partial \triangle_k)\, \varepsilon \max_{z \in \triangle_k} |z - z_0|$$

$$\leq L(\partial \triangle)^2\, \varepsilon$$

$\square$

Corollary 4.23.

1. *A holomorphic function always has a local antiderivative*

2. *A holomorphic function on a star shaped domain has a global antiderivative and*

$$\int_{\gamma} f(z)\, dz = 0$$

   for any closed path

3. *The prerequites of Goursat theorem can be weakened to continous and holomorphic with the exception of a finite number of points: adequate partioning of the original triangle*

Theorem 4.24 (Cauchy's Intergral Formula). *Let $\Omega \subseteq \mathbb{C}$ be open and f holomorphic on Ω. Further let $D \subset \Omega$ be a disc. Then*

$$f(z) = \frac{1}{2\pi i} \int_{\partial D} \frac{f(\zeta)}{\zeta - z}\, dz$$

for $z \in D$.

Proof. For $z \in D$ define
$$h(\zeta) = \frac{f(\zeta) - f(z)}{\zeta - z}$$
for $\zeta \neq z$ and $f'(z)$ for $\zeta = z$. Then h is holomorphic on $D \setminus \{z\}$ and continous at z
$$0 = \int_{\partial D} h(\zeta)\, d\zeta = \int_{\partial D} \frac{f(\zeta)}{\zeta - z}\, d\zeta - f(z) \int_{\partial D} \frac{1}{\zeta - z}\, d\zeta = \int_{\partial D} \frac{f(\zeta)}{\zeta - z}\, d\zeta - 2\pi \mathrm{i} f(z)$$

$\square$

5 Neural Networks

5.1 The Perceptron

Definition 5.1 (Binary Classifiers). *Let $X \subset \mathbb{R}^n$ be the union of two finite disjoint sets $X = M \cup N$.*

1. *A* binary classification problem *is the task to find a mapping $f : X \to \{0, 1\}$ with*

$$f(x) = \begin{cases} 1 & \text{for } x \in M \\ 0 & \text{for } x \in N \end{cases}$$

   f then is called a binary classifier *for X*

2. *X is called* separable *if there exists a* weight vector *$w \in \mathbb{R}^n$ and a* bias *$b \in \mathbb{R}$ so that*

$$wx + b > 0 \quad \text{for } x \in M$$
$$wx + b < 0 \quad \text{for } x \in N$$

3. *The weight w and the bias b are called* solution to the classification problem. *They implicitly define a binary classifier via*

$$f(x) = \begin{cases} 1 & \text{if } wx + b > 0 \\ 0 & \text{if } wx + b < 0 \end{cases}$$

Examples 5.2.

1. *Let $X = \{0, 1\} \times \{0, 1\}$ and consider the* and *operator $f(1, 1) = 1$ and $f(x, y) = 0$ elsewhere. Then $w = (3, 3)$ and $b = -5$ yield a solution to the classification problem $M = f^{-1}(1)$ and $N = f^{-1}(0)$*

2. *Again let $X = \{0, 1\} \times \{0, 1\}$ and $f(1, 0) = f(0, 1) = 1$ and $f(0, 0) = f(1, 1) = 0$, the* xor *operator. Thus for any weight (w_1, w_2) and any bias b*

$$w_1 + b > 0$$
$$w_2 + b > 0$$

$$w_1 + w_2 + b \leq 0$$
$$b \leq 0$$

   Adding two equations respectively shows that there cannot be a solution

3. *The bias can be integrated into the weight vector via $w' = (w, b) \in \mathbb{R}^{n+1}$ and $x' = (x, 1) \in \mathbb{R}^{n+1}$. Separability then reduces to*

$$w'x' > 0$$

Geometrical Interpretation

The idea for the perceptron most likely has its origin in a simple geometrical observation. Recall that for $x, y \in \mathbb{R}^n$ the dot product can be expressed as

$$xy = \|x\| \|y\| \cos \alpha$$

where α is the angle between the two vectors. Hence the product is positive if the angle is less than $90°$ degrees and negative if the angle is between $90°$ and $180°$ degrees

$$xy > 0 \quad \text{for } 0 \leq \alpha < \pi/2$$
$$xy < 0 \quad \text{for } \pi/2 < \alpha \leq \pi$$

Note, that the sign does not depend on the vector lengths, but solely on the angle.

For any two vectors it is easy enough to find a weight that satisfies $wx > 0$ and $wy > 0$. Generally $w = x + y$ is a good guess, but not always correct as shown below in Figure 2.
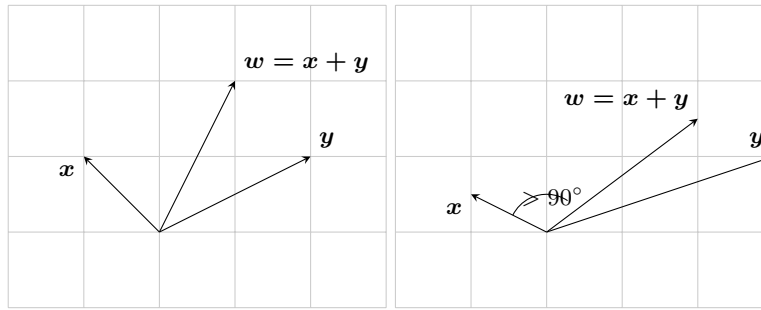

Figure 2: Dot Product and Angle

But, the more similar the lengths of the two vectors are the more likely $x + y$ works. The actual threshold is given by the following

Lemma 5.3. *Let $x, y \in \mathbb{R}^n$ with $\|x\| < \|y\|$ and $xy = \|x\| \|y\| \cos \alpha$. If*

$$\|x\| > -\|y\| \cos \alpha$$

then $x(x + y) > 0$.

Proof. Let $x(x + y) = \|x\| \|x + y\| \cos \beta$. Since

$$x(x + y) = xx + xy = \|x\|^2 + \|x\| \|y\| \cos \alpha$$

it follows

$$\|x + y\| \cos \beta = \|x\| + \|y\| \cos \alpha$$

Hence $\cos \beta > 0$ if the inequality above holds. $\qquad \square$

An iterative approach is to repeatedly increase $w = x + y$ in the direction of the shorter vector aka the one with the angle greater than $90°$ degrees.

$$w' = \begin{cases} w + x & \text{if } wx \leq 0 \\ w + y & \text{if } wy \leq 0 \end{cases}$$

While this seems reasonable it is unclear whether the algorithm always yields a result after a finite number of iterations. The answer to this question will be given later as a special case of the Perceptron Convergence Theorem 5.6.

This approach can be used to to separate two vectors. Finding a common weight for x and $-y$ now yields $wx > 0$ and $-wy > 0$, hence $wy < 0$.

$$w' = \begin{cases} w + x & \text{if } wx \leq 0 \\ w - y & \text{if } wy \geq 0 \end{cases}$$

Algorithm 5.4 (Weight).

```
from vector import add, dotprod


def weight(x, y):
    w = add(x, y)
    while True:
        if dotprod(w, x) <= 0:
            w = add(w, x)
        elif dotprod(w, y) <= 0:
            w = add(w, y)
        else:
            return w
```

Examples 5.5.

1. Let $x = (-1, 1)$ and $y = (6, 1)$. Then

$$\begin{aligned} w_0 &= (5, 2) & w_0 x &= -3 & w_0 y &= 32 \\ w_1 &= (4, 3) & w_1 x &= -1 & w_1 y &= 27 \\ w_2 &= (3, 4) & w_2 x &= 1 & w_2 y &= 22 \end{aligned}$$

2. Let $x = (4, -6)$ and $y = (-10, 5)$. Then

$$\begin{aligned} w_0 &= (-6, -1) & w_0 x &= -18 & w_0 y &= 55 \\ w_1 &= (-2, -7) & w_1 x &= 34 & w_1 y &= -15 \\ w_2 &= (-12, -2) & w_2 x &= -36 & w_2 y &= 110 \\ w_3 &= (-8, -8) & w_3 x &= 16 & w_3 y &= 40 \end{aligned}$$

3. Task: for any given integer k find two vectors so that more than k steps are needed

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4. *Let $w = x/\|x\| + y/\|y\|$. Then*

$$wx = \frac{\|x\|^2}{\|x\|} + \frac{yx}{\|y\|} = \|x\| + \|x\| \cos \alpha = (1 + \cos \alpha)\|x\| > 0$$

and on the other hand $wx = \|w\|\|x\| \cos \beta$. Similarily $wy = (1 + \cos \alpha)\|y\| = \|w\|\|y\| \cos \gamma$. Hence

$$1 + \cos \alpha = \|w\| \cos \beta = \|w\| \cos \gamma$$

and thus $\beta = \gamma$. Also, as expected, $wx > 0$ and $wy > 0$

The following theorem is the generalization of the approach above for binary classifiers. Roughly sepaking the algorithm changes the weight only for misqualified vectors. The proof measures the change of angle against the change of length of the weight vector and provides an estimation for the maximumn number of required steps.

Theorem 5.6 (Perceptron Convergence Theorem). *Let $X = M \cup N$ be separable by $w^* \in \mathbb{R}^n$. Define $w_0 = 0$ and repeat to iterate over all $x \in X$ via*

$$w_{k+1} = \begin{cases} w_k + x & \text{if } x \in M \text{ and } w_k x \leq 0 \\ w_k - x & \text{if } x \in N \text{ and } w_k x \geq 0 \\ w_k & \text{else} \end{cases}$$

until no further changes occurr. Suppose $\|x\| \leq r$ and $|w^* x| \geq \delta > 0$ for $x \in X$. Then the number of iterations before the algorithm stops is limited by

$$k \leq \frac{r^2}{\delta^2}$$

Proof. Let $x \in M$ with $w_k x \leq 0$. Then

$$w^* w_{k+1} = w^* w_k + w^* x \geq w^* w_k + \delta$$

Furthermore

$$\|w_{k+1}\|^2 = \|w_k + x\|^2 = \|w_k\|^2 + 2w_k x + \|x\|^2 \leq \|w_k\|^2 + r^2$$

The same estimations holds for $x \in N$ with $w_k x \geq 0$ and using induction yields

$$w^* w_k \geq k\delta \text{ and } \|w_k\|^2 \leq kr^2$$

Assuming $\|w^*\| = 1$ now gives

$$k^2 \delta^2 \leq \|w_k\|^2 \leq kr^2$$

which proves the initial inequality. $\qquad \square$

Algorithm 5.7 (Perceptron).

```python
from weight import dotprod


def perceptron(T, t=0.01, epochs=50):
    n = len(T[0])
    w, b = n * (0.0,), 0.0
    for _ in range(epochs):
        done = True
        for x, y in T:
            if dotprod(w, x) + b >= 0:
                z = 1
            else:
                z = 0
            w = tuple(w[i] - t * (z - y) * x[i] for i in range(n))
            b = b - t * (z - y)
            if not y == z:
                done = False
        if done:
            break
    return w, b
```

Remark 5.8. *Check Banach Fixed-Point Theorem 1.15 for an alternative proof?!*

5.2 Single Layer Feedforward Neural Networks and the Delta Rule

Definition 5.9 (Activation Function).

1. *There is no valuable formal definition for an* activation function $\alpha : \mathbb{R} \to Y \subseteq \mathbb{R}$. *However, the behaviour of an activation function should somehow relate to the fireing behaviour of an actual biological neuron. Continuity, differentiability, monotony and the behaviour at infinity are of special interest in order to apply mathematical concepts*

2. *An activation function $\alpha : \mathbb{R} \to \mathbb{R}$ with $\alpha(x) \to 0$ for $x \to -\infty$ and $\alpha(x) \to 1$ for $x \to \infty$ is called a* sigmoidal activation function

3. *The* Heaviside *function $H : \mathbb{R} \to \{0,1\}$ is defined as*

$$H(x) = \begin{cases} 1 & \text{for } x > 0 \\ 0 & \text{for } x \leq 0 \end{cases}$$

4. *The* sigmoid *function $\sigma \in C^{\infty}(\mathbb{R})$ is defined as*

$$\sigma(x) = \frac{1}{1 + e^{-x}}$$

Remarks 5.10.

1. *The heaviside function is not continuous and therefore not differentiable at 0. In that sense the sigmoid function can be considered its smooth counterpart*

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2. *The definition of the sigmoid function yields $0 < \sigma(x) < 1$ as well as $\sigma(x) \to 0$ for $x \to -\infty$ and $\sigma(x) \to 1$ for $x \to \infty$*

3. *The quotient rule yields*

$$\sigma'(x) = -\frac{-e^{-x}}{(1 + e^{-x})^2} = \sigma(x)\frac{1 + e^{-x} - 1}{1 + e^{-x}} = \sigma(x)(1 - \sigma(x))$$

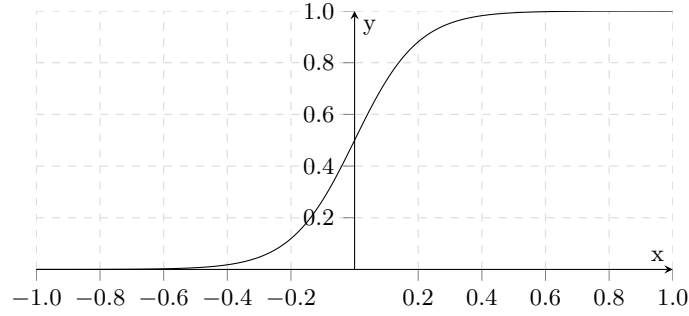*and σ is monotonically increasing over its domain*

Figure 3: The sigmoid function $\sigma(x) = \frac{1}{1+e^{-x}}$

Remarks 5.11.

1. *Let $f \in C^1(\mathbb{R}^n)$ and $g(x) = f(x|^2$. The multivariable chain rule yields*

$$\frac{\partial g}{\partial x_i}(x) = \frac{\partial f^2}{\partial x_i}(x) = 2f(x)\frac{\partial f}{\partial x_i}(x)$$

and

$$\nabla g(x) = 2f(x)\nabla f(x)$$

2. *More generally for $f = (f_1, f_2, \ldots f_n) \in C^1(\mathbb{R}^n, \mathbb{R}^m)$ and*

$$g(x) = \|f(x)\|^2 = \sum_{j=1}^{m} f_j(x)^2$$

it follows that

$$\frac{\partial g_j}{\partial x_i}(x) = \sum_{j=1}^{m} \frac{\partial f_j^2}{\partial x_i}(x) = 2\sum_{j=1}^{m} \frac{\partial f_j}{\partial x_i}(x)f_j(x)$$

and

$$\nabla g(x) = 2\sum_{j=1}^{m} \nabla f_j(x)f_j(x)$$

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3. Now let $x \in \mathbb{R}^n$ and $y \in R$ be fixed. For $q(w, b) = wx + b - y$ follows

$$\frac{\partial q}{\partial w_i}(w, b) = \frac{\partial}{\partial w_i} \sum_{i=0}^{n} w_i x_i + b - y = x_i$$

$$\frac{\partial q}{\partial b}(w, b) = 1$$

Hence

$$\nabla q(w, b) = (x_1, x_2, \ldots, x_n, 1)$$

4. Let $\alpha \in C^1(\mathbb{R})$ be an activation function and define $p(w, b) = (y - \alpha(wx + b))^2$. Then

$$\frac{\partial p}{\partial w_i}(w, b) = -2(y - \alpha(wx + b))\alpha'(wx + b)x_i$$

$$\frac{\partial p}{\partial b}(w, b) = -2(y - \alpha(wx + b))\alpha'(wx + b)$$

and

$$\nabla p(w, b) = -2(y - \alpha(wx + b))\alpha'(wx + b)\nabla q(w, b)$$

Definition 5.12 (Single Layer Feedforward Neural Network). *Let $\alpha \in C^1(\mathbb{R})$ be an activation function, $w \in \mathbb{R}^n$ and $b \in \mathbb{R}$ be weight and bias. Define $f \in C^1(\mathbb{R}^n)$ as*

$$f(x) = \alpha(wx + b)$$

Then $f = \alpha(wx + b)$ is called a single layer feedforward neural network.

Definition 5.13 (Error Function). *Let $f = \alpha(wx + b)$ be a single layer feedforward neural network.*

1. *A finite set $T \subset \mathbb{R}^n \times \mathbb{R}$ is called* test set

2. *The* error function *is defined as the mean squared error over the samples*

$$E(w, b) = \frac{1}{2} \sum_{(x,y) \in T} (y - f(x)))^2$$

Lemma 5.14. *It is*

$$\frac{\partial E}{\partial w_i}(w, b) = \sum_{(x,y) \in T} -(y - f(x))\alpha'(wx + b)x_i$$

$$\frac{\partial E}{\partial b}(w, b) = \sum_{(x,y) \in T} -(y - f(x))\alpha'(wx + b)$$

Proof. This follows from the remarks above. $\square$

The Delta Rule applies the method of the Steepest Descent to the error function of a neural network.

Algorithm 5.15 (Delta Rule)**.** *Let* $f = \alpha(wx + b)$ *be a single layer feedforward neural network and* T *a test set. The delta rule modifies weight and bias in the direction of the greatest descent* $-\eta\nabla E(w, b)$ *at a given learning rate* $\eta > 0$. *For* $(x, y) \in T$ *this is*

$$\Delta w_i = \eta(y - f(x))\alpha'(wx + b)x_i$$
$$\Delta b = \eta(y - f(x))\alpha'(wx + b)$$

```python
from random import random
from perceptron import dotprod


def delta_rule(T, a, da, t=0.01, eps=0.05, epochs=50):
    n = len(T[0])
    w, b = n * (random(),), random()
    for _ in range(epochs):
        err = 0.0
        for x, y in T:
            z = dotprod(w, x) + b
            az, daz = a(z), da(z)
            d = y - az
            w = tuple(w[i] + t * d * daz * x[i] for i in range(n))
            b = b + t * d * daz
            err = err + d * d
        if 0.5 * err < eps:
            break
    return w, b
```

5.3   The Univeral Approximation Theorem

Definition 5.16 (Single Layer Neural Network)**.** *Let* $w_k \in \mathbb{R}^n$ *and* $b_k \in \mathbb{R}$. *For* $\alpha_k \in \mathbb{R}$ *define* $f : \mathbb{R}^n \to \mathbb{R}$

$$f(x) = \sum_{k=0}^{m} \alpha_k \sigma(w_k x + b_k)$$

Then f *is called a* single layer feedforward neural network *of* m *units.*

Theorem 5.17 (The Universal Approximation Theorem)**.** *The subspace of all single layer feedforward neural networks is dense in* $C(I^n)$.