

# Practical Functional Programming

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**THIS  
PRESENTATION  
IS NOT ABOUT  
LISP**



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All the examples are written in PHP 5.3  
(So, you can listen and not just pretend)

# Overview

- Functional Programming is...
- How can that possibly work?
- Evaluation Strategies
- FP in "non-FP" Languages
- The Future

# What is Functional Programming?

Based on Church and Kleene's  $\lambda$ -calculus, a formal system designed to investigate function definition, function application and recursion

- Functions are first class citizens
- Passed as parameters
- Returned as results
- $\lambda$ -calculus is capable of expressing **any** algorithm
  - (e.g. it's computationally equivalent to a Turing Machine)

# What is Functional Programming?

- A programming paradigm that treats computation as the evaluation of mathematical functions and avoids state and mutable data.
- In Contrast, Imperative Programming emphasizes changes in state, and mutating data.

# What are "Side Effects"?

A function or expression produces a side effect if it modifies some state in addition to returning a value

- Examples:
  - Writing to the screen
  - Reading from Standard Input/File
  - Modifying a global variable
  - Modifying a variable passed by reference
- Imperative programs are notorious for producing side effects.



# What is Referential Transparency?

A function is Referentially Transparent when its value only relies on its inputs. Therefore calling it twice with the same arguments is equivalent.

- A function that produces side effects is **not** referentially transparent
- Memoization, an optimization that only computes what hasn't been computed before can only be used with referentially transparent functions.
- EXAMPLE

# Why get rid of side effects?

- Makes it easier to reason about code
- More Predictable Behavior
- Much easier to test functionality in automated tests.

How can that possibly work?

# Recursion

- Define functions in terms of themselves
- Get "state" for free, as stored on the call stack
- It's a side effect free looping construct
- **BONUS:** Many algorithms are more easily understood using recursively
  - Dealing with trees
  - "Divide and Conquer"

Recursion is **INEFFICIENT!**

# Recursion is INEFFICIENT!

- Allocates new stack frame for every call
- Pushes arguments on the stack
- Sets up return address
- etc, etc, etc...
- EXAMPLE

# SOLUTION: Tail Calls and Tail Call Optimization

# Tail Calls

- A function call is in tail position if upon its return there is no computation left to perform
- EXAMPLE



# Tail Call Optimization

- Tail calls have no computation left...
- There's still a stack frame, a return address, everything we need, why not reuse it?
- OK! Instead of allocating a new frame and doing all the setup, overwrite the old arguments with the new and jump to the beginning of the code!
- while loops are just guarded gotos!
- TCO == while loops

But, of course...

... not all languages optimize tail  
calls. UGH.

# Most functional languages do... But,

- Perl (sort of does... )
- PHP
- Python
- Java
- ...

Don't!

- They consume stack space, thus consume memory.
- Memory isn't infinite. Thus, we can't recurse infinitely...

... Or can we?

# How do we infinitely recurse?

## Trampolines!

- A bounce on a trampoline calls a function that returns a package with the result and either an indication of "land" or "bounce"
- Requires some work
- EXAMPLE

# Other cool things...

- Continuation Passing Style
  - Instead of returning...
    - Pass as an extra argument a function to call at the end of the computation
    - Need to have first class functions
- Bundle with Trampolining in a language without TCO (or first class continuations) and we can get THREADS
- Can also implement suspend/resume for computation!
  - Great idea for web programming! (Ask me later)

# Recursion: Higher Order Looping

- Map
  - applies a function to each element of a list and returns a list
    - *all\_upcased = map(strtoupper, list\_of\_strings)*
- Fold
  - applies the previous result and the next element of a list to a function and returns an accumulated value (could be a list, or a number, etc..)
    - *sum\_of\_list = fold(+, 0, list\_of\_ints)*



# Higher Order Functions

- Functions treated as "first-class" values
  - (e.g. Passed as parameters, returned from functions, assigned to variables)
- Closures

# Closures

*"A closure is an object that supports exactly one method: 'apply'"*

*-- Guy Steele*

- A function whose free variables are defined in an enclosing scope.
  - Free variable: A variable not bound by either a function parameter or a variable declaration within a function
  - Contrast to: Bound variable: A variable bound by either a function parameter, or a variable declaration.

# Free/Bound Variables

```
function make_counter(n) {  
  // n is "bound" with respect to make_counter  
  return function(increment) {  
    // n is "free" with respect to this function  
    // increment is "bound" with respect to this function  
    n = n + increment;  
    return n;  
  };  
}
```

EXAMPLE

# Currying

Function is applied to its arguments one at a time, with each application returning a new function that accepts the next argument.

- The closure example was an example of currying.
- Allowed us to create "state"

# Partial Application

More general form of currying. A function is applied to some of its arguments and returns a new function that accepts the rest of the arguments.

- Used in a similar fashion to currying.

# Properties of Pure Functions

## No Memory or I/O Side Effects (i.e. Referentially Transparent)

- If a result is not used, the call can be eliminated without affecting other expressions.
- Result of a function call with parameters  $X$  is constant with respect to  $X$ . (e.g. we only need to compute it once!)
- If there is no data dependency between pure functions  $F$  and  $G$ , they can be called in *parallel*, or even in reverse.
  - (e.g.  $F(); G() == G(); F()$ )
  - $G(F()) != F(G())$
- If the whole language is pure (i.e. expressions as well) then any evaluation strategy can be used!

# Strict vs. Non-Strict Evaluation

# Strict Evaluation: "Call by value"

- Evaluated at the call site, at call time.
- (e.g. `$x = some_long_computation();` // `$x` gets the value of `some_long_computation()` at the time of assignment)



# Non-Strict (a.k.a Lazy): "Call by Need"

- Only evaluate something when it's absolutely needed
  - Like, when you have a sink full of dishes and no forks left in the drawer...
- Can avoid unnecessary computation that might happen in call by value
  - (e.g. `$x = some_long_computation()` // `$x` is now a promise to compute `some_long_computation()`)

EXAMPLE

FP in "non-FP" languages

# FP in "non-FP" languages

- Avoid mutation
- Avoid pass by reference, except to avoid copying large objects, but you must never mutate the reference object
- Avoid global state
  - Create referentially transparent functions that "compute" state
    - (e.g. Pacman: The dots blink.)
      - `show_dots(time) -> (time % 2) == 0`
- Accept more arguments and use Currying and partial application to avoid always passing them in
- FP is pretty damn hard to really do without "first-class" functions, but it's techniques are useful for any paradigm

The Future

# Multi-core, Parallel Processing, Concurrency

Hard problems because independent things share the same memory

- Race conditions
  - We modify `account_balance` at the same time. I increment, you decrement... what's our balance?

# Multi-core, Parallel Processing, Concurrency

Key thing was "modify"

- Remember when I said "pure functions" can be run in parallel?

# Functional Programming Makes Sense!

... and it's a solution for the future

# Learn more!

- Languages: Haskell, OCaml, Scheme, Common Lisp, Erlang
- Read:
  - Wikipedia
  - <http://lambda-the-ultimate.org>
  - <http://www.readscheme.org/>
  - SICP: <http://mitpress.mit.edu/sicp/>