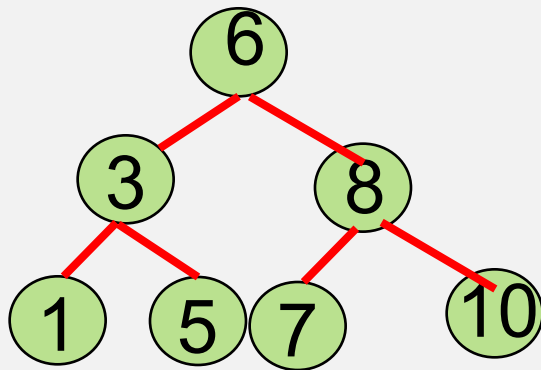




การ insert Binary Search Trees

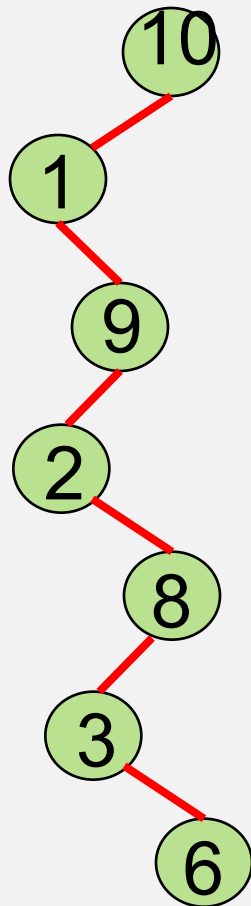
6 3 8 1 5 7 10



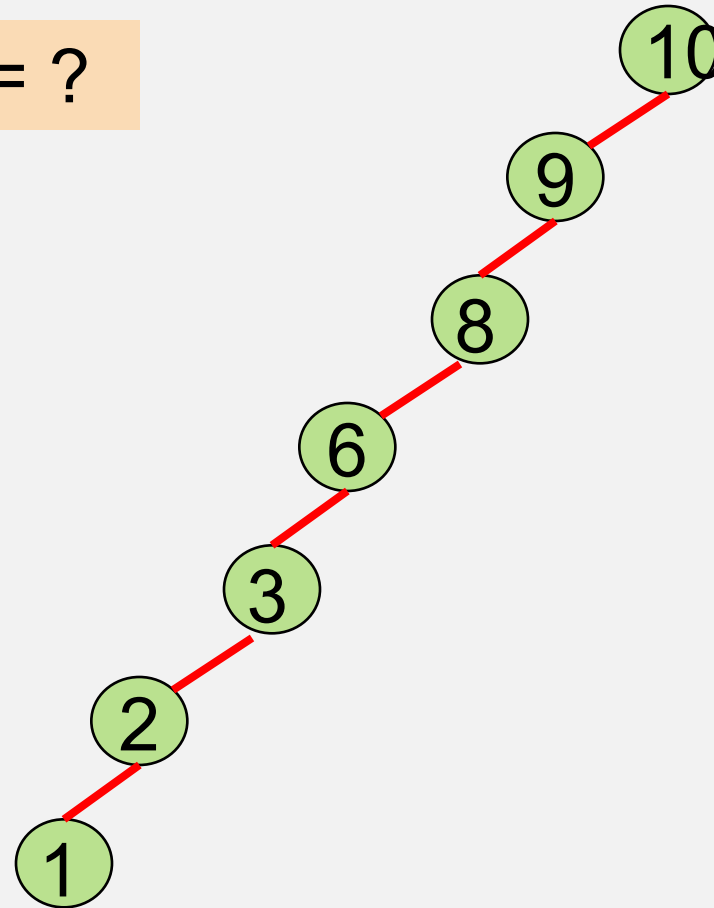


Insert : 10 1 9 2 8 3 6

Insert : 10 9 8 6 3 2 1



Big O = ?





ปัญหา Big O ของ Binary Search Trees

worse case $O(n)$

แก้ไข

2.1 ต้องทำให้หรือ Perfect Balance Trees : ทำยาก

2.2 AVL Trees

AVL : Binary Search Trees ที่มี Balance condition
(ความสูงต่างกันไม่เกิน 1)

- Single Rotation
- Double Rotation



3.4 AVL Trees

A binary search tree with **balance condition**, it ensures that the depth of the tree is $O(\log_n)$.

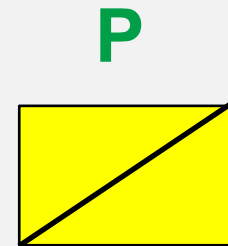
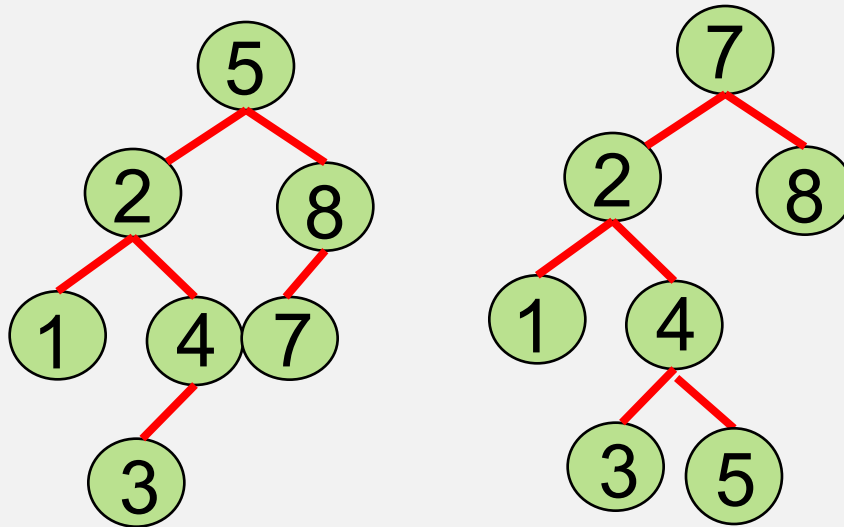
Balance condition : Every node must have left and right subtree of the same height.

Problem : Only **perfect balance trees** of $2^k - 1$ node would satisfy this criterion.



Definition : An AVL(Adelson-Velskii and Landis) tree is identical to a binary search tree, except that for every node in the tree, the height of the left and the right subtree **can differ by at most 1**.

AVL Operation All = $O(\log_N)$ except insertion



The height of an empty subtree is defined to be -1.



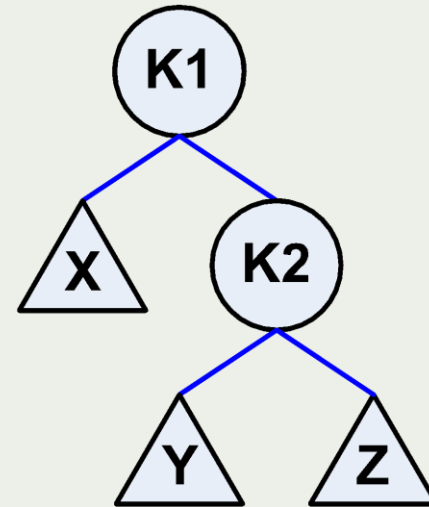
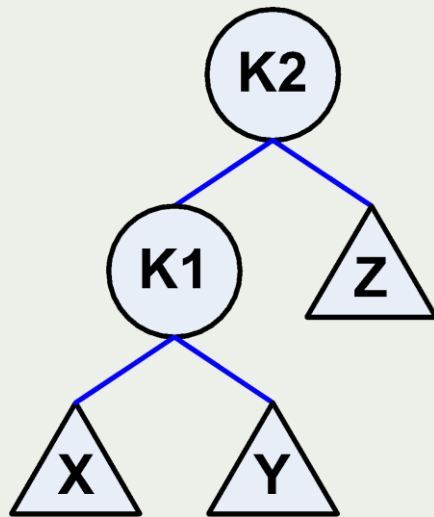
3.4.1 Insertion operation :

1. Update all the balancing information for the nodes on the path back to the root
2. The insertion a node could violate the AVL tree property, then property has to be restored before the insertion step. Called a *rotation*
 - **Single rotation**
 - **Double rotation**



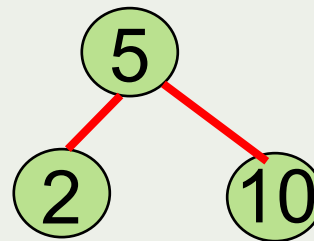
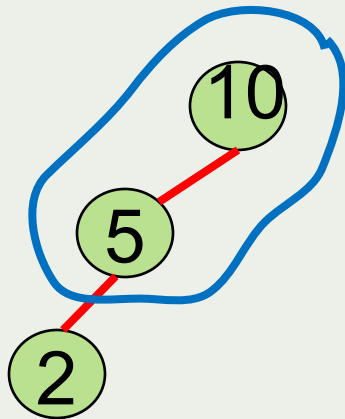
3.4.2 Single Rotation

- A rotation involves only a few pointer changes, and changes the structure of the tree while preserving tree property.





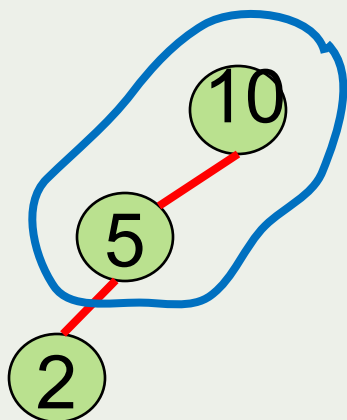
Example insert 10, 5, 2



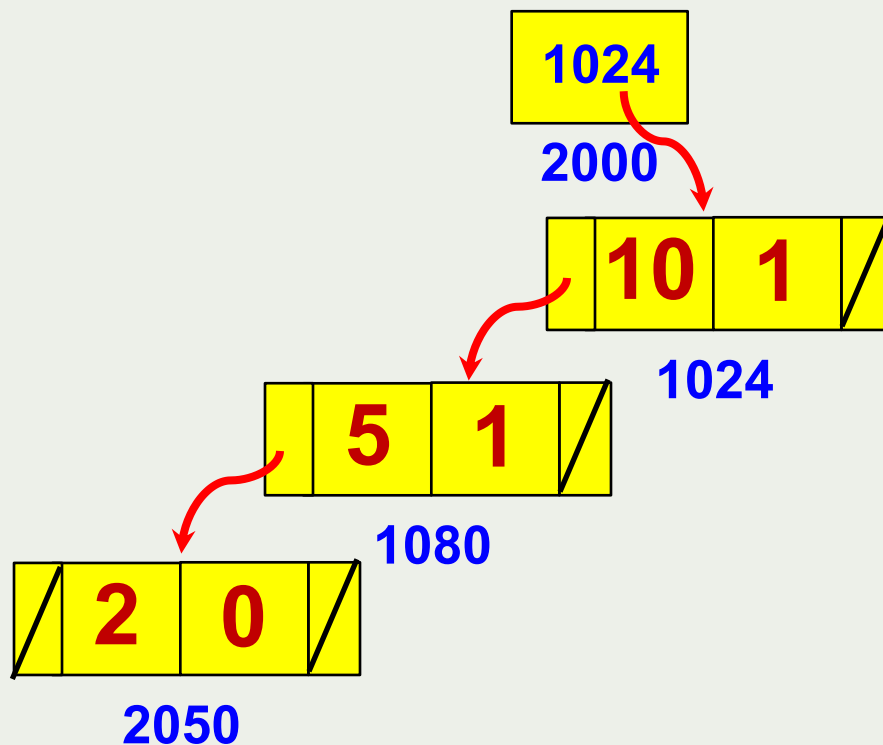


Example insert 10, 5, 2

นอกช่วง

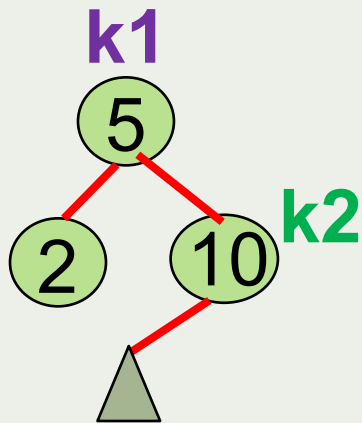
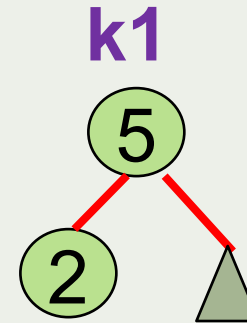
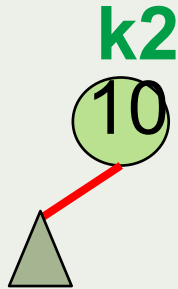
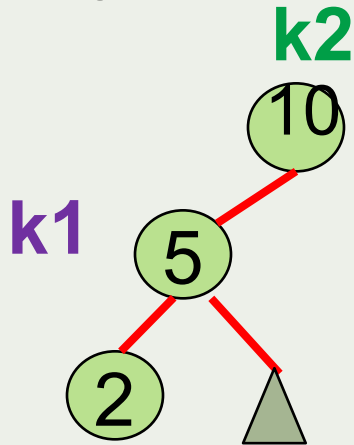


tree





Example insert 10, 5, 2



```
struct node *srleft(struct node *k2)
{
    struct node *k1;
    k1=k2->left;
    k2->left = k1->right;
    k1->right = k2;
}
```



```
struct node *srleft(struct node *k2)
```

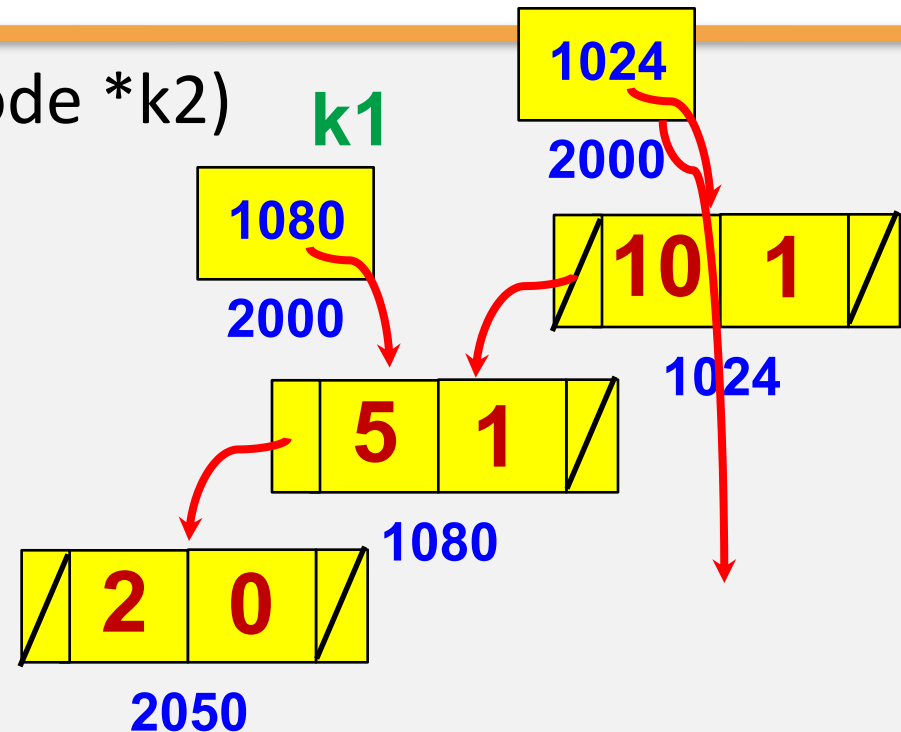
```
{ struct node *k1;
```

```
  k1=k2->left;
```

```
  k2->left = k1->right;
```

```
  k1->right = k2;
```

```
}
```



Example insert 1, 2, 3, 4, 5, 6, 7

นอกช่วง

Example insert 15, 14

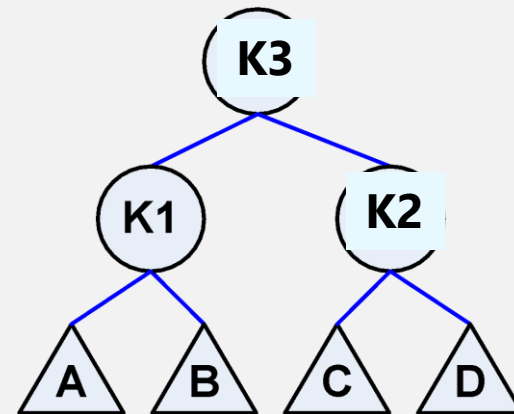
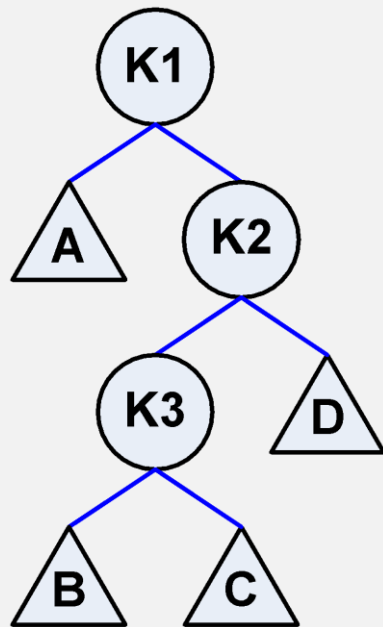
insert 13, 12, 11, 10, 9, 8, 8.5

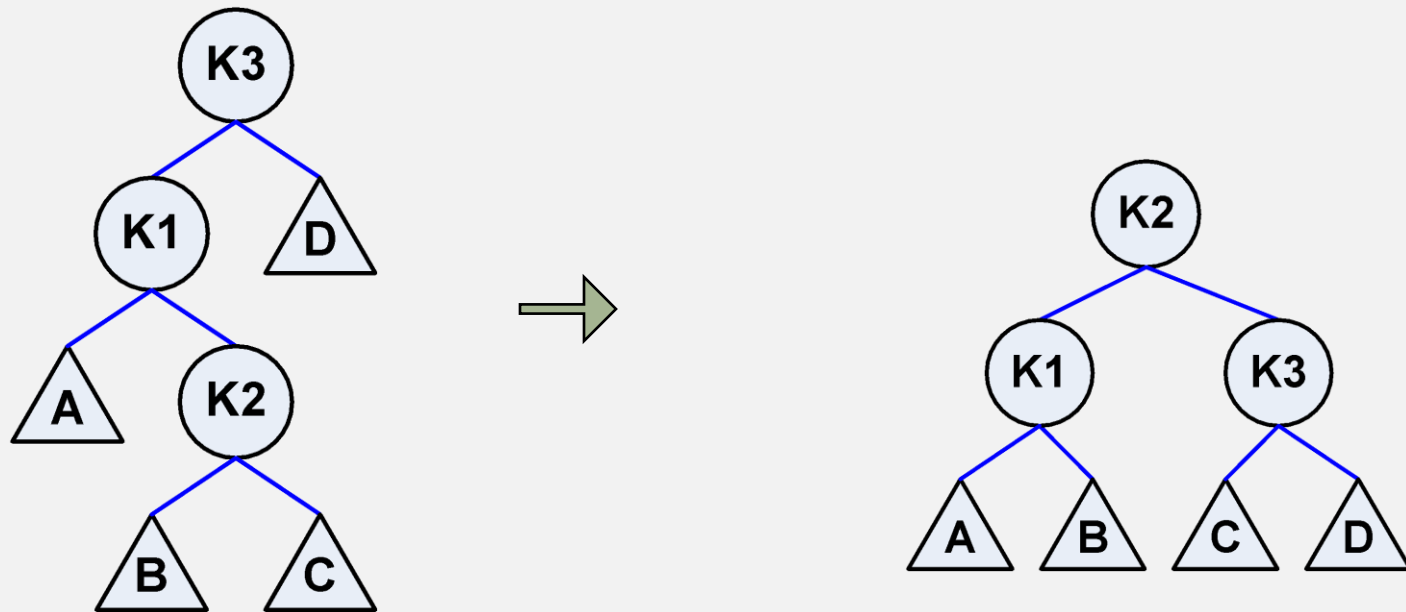
นอกช่วง



3.4.3 Double rotation

Problem : Single rotation has not fixed the height imbalance by a node inserted into the tree containing the middle elements.





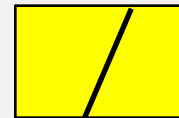


```
struct node
{
    int value;
    int height;
    struct node *left;
    struct node *right;
};
```

```
1. int fheight( struct node *P)
2. { if ( P==NULL)
3.     return -1;
4.     else
5.     return P->height;
6. }
```



P



P



1024



Code Insert

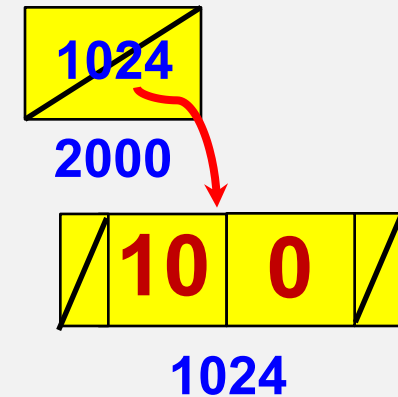
```
1 ถ้า tree == NULL
2   สร้าง node ใหม่ ใส่ค่า
3 else
4   ถ้า x น้อยกว่า tree->value ให้ recursive ลงไปทางซ้าย
5   tree->left = insert (tree->left)
6   ถ้าความสูงของทรีทางซ้ายและทางขวาต่างกันเกิน 2
7     ถ้า x น้อยกว่า tree->left->value
8       single rotation (อยู่นอกช่วง)
9     else
10      double rotation (อยู่ในช่วง)
11 update ความสูง
12 return tree
```

เมื่อ insert แล้ว return
กลับมาเช็คความสูง



```
struct node *insert(int x, struct node *T)
1{   if(T == NULL)
2   {   T=new struct node;
3       T->value=x;
4       T->left=T->right=NULL;
5       T->height=0;
6   }
   else
```

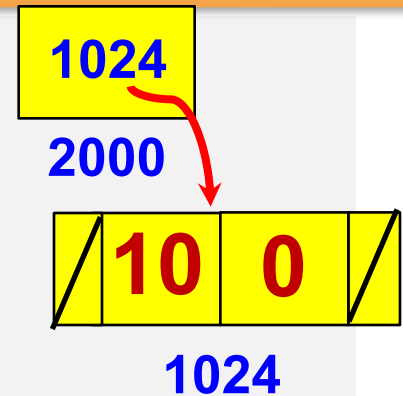
tree





```
8 if(x < T->value )
9     { T->left = insert(x,T->left);
10         if(fheight(T->left) - fheight(T->right) == 2)
11             if( x < T->left->value)
12                 T=srleft(T);
13             else
14                 T=drleft(T);
15     }
16 else
17 ถึง 29 if(x > T->value )
18     { ข้างขวาบ้าง ให้เขียนเอง ... }
```

```
30 T->height =  $\max(-1, -1) = -1$  +1;
31 return T;
}
```

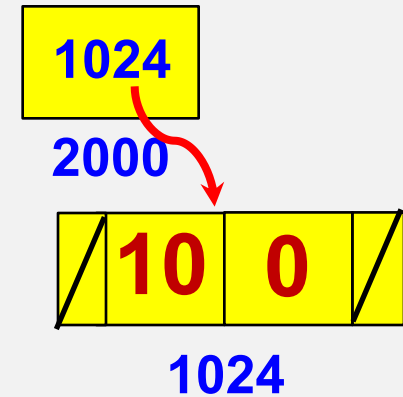




```
struct node *insert(int x, struct node *T)
1{  if(T == NULL)
2    {  T=new struct node;
3        T->value=x;
4        T->left=T->right=NULL;
5        T->height=0;
        }
7  else
```



tree



ถ้า x น้อยกว่า tree->value
ให้ recursive ลงไปทางซ้าย

– Tree

tree

21

5, NULL

5

```
8 if(x < T->value )
9   { T->left = insert(x,T->left);
10     if(fheight(T->left) - fheight(T->right) == 2)
11       if( x < T->left->value)
12         T=srleft(T);
13       else
14         T=drleft(T);
15   }
```

16else

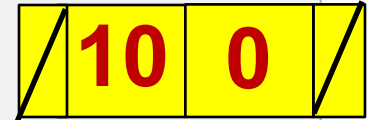
```
17ถึง 29 if(x > T->value )
    { ข้างขวาบ้าง ให้เขียนเอง ... }
```

```
30 T->height = max(fheight(T->left), fheight(T->right)) + 1;
```

```
31 return T;
}
```

1024

2000



1024

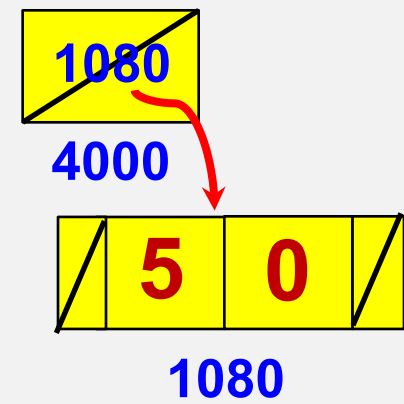


Recursive 1

5, NULL

```
struct node *insert(int x, struct node *T)
1{   if(T == NULL)
2   {   T=new struct node;
3       T->value=x;
4       T->left=T->right=NULL;
5       T->height=0;
        }
7   else
        ....
```

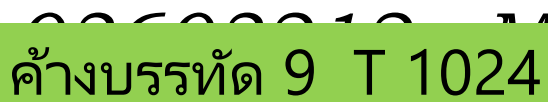
tree



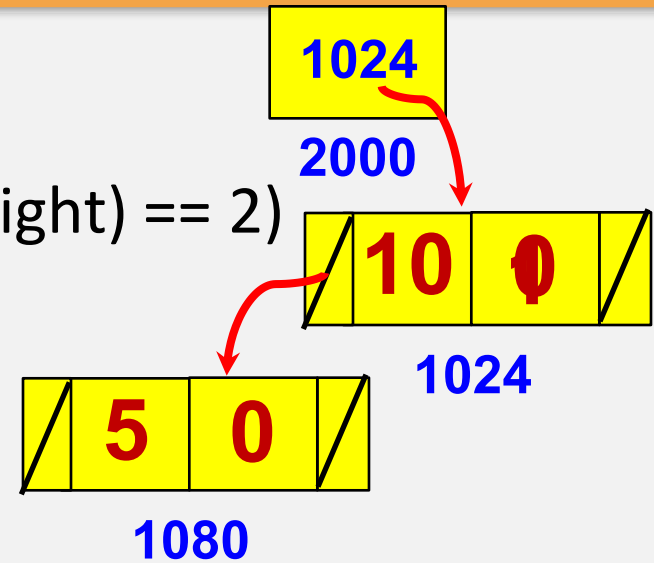
$$\text{max}(-1, -1) = -1 \quad +1$$

```
30 T->height = max(fheight(T->left), fheight(T->right)) + 1;
31 return T;
    }
```

1080

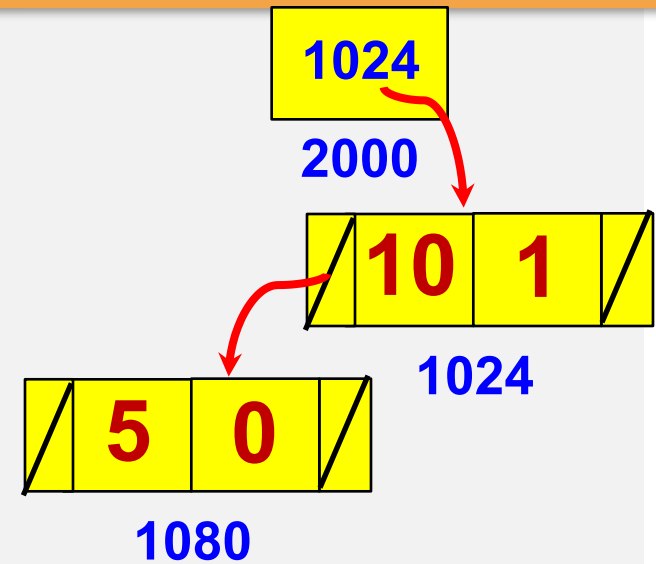


24 **1080**





```
main()
{
    .....
    tree=insert(2,tree);
}
```



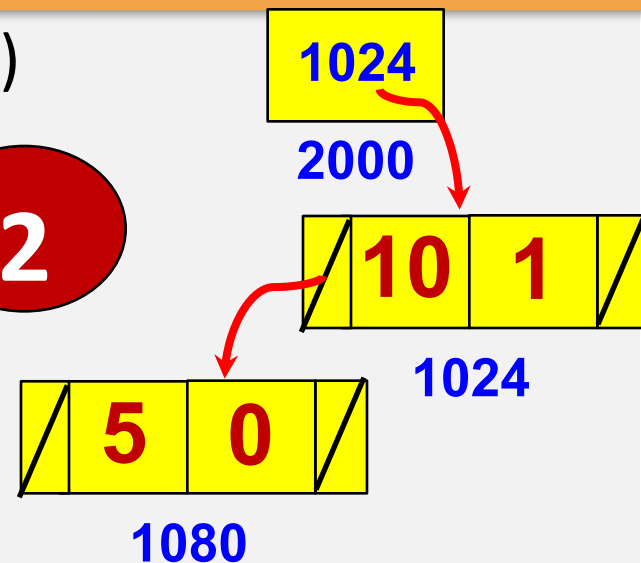


จาก main

2, 1024

```
struct node *insert(int x, struct node *T)
```

```
1{   if(T == NULL)
2   {   T=new struct node;
3       T->value=x;
4       T->left=T->right=NULL;
5       T->height=0;
        }
7   else
        ....
```



ถ้า x น้อยกว่า tree->value
ให้ recursive ลงไปทางซ้าย

Tree

tree

26

8 if(x < T->value)

2, 1080

9 { T->left = insert(x,T->left);

10 if(fheight(T->left) - fheight(T->right) == 2)

11 if(x < T->left->value)

12 T=srleft(T);

13 else

14 T=drleft(T);

15 }

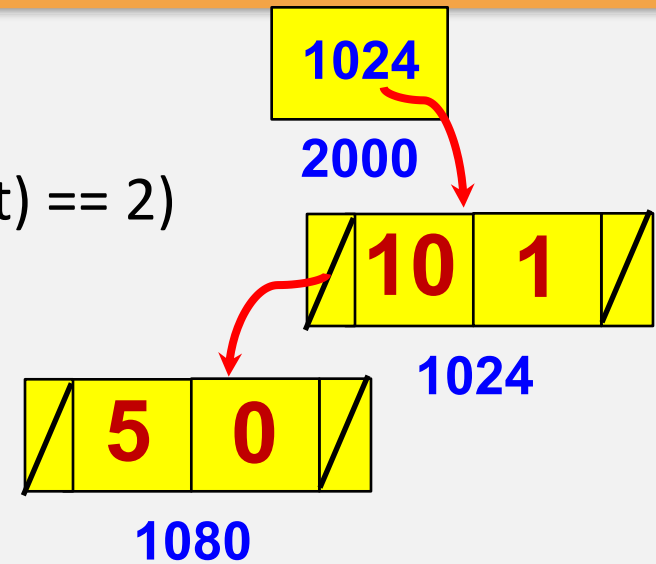
16ถึง 29 if(x > T->value)

{ ข้างขวาบ้าง ให้เขียนเอง ... }

30 T->height = max(fheight(T->left), fheight(T->right)) + 1;

31 return T;

}



(A) T = 1024 ค้าง 9 หน้า 26

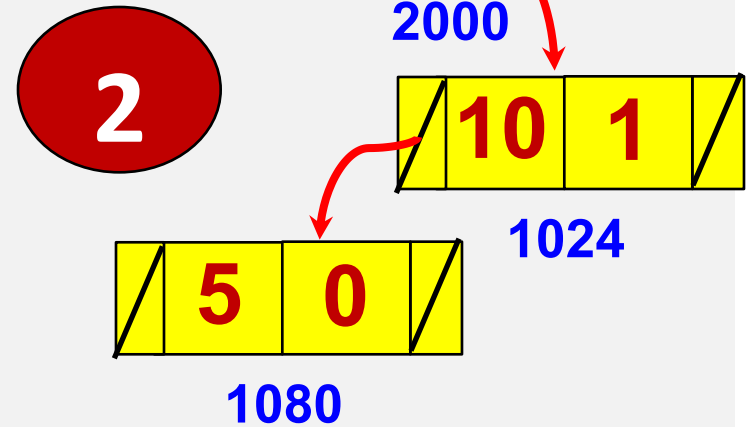


Recursive 1

2, 1080

```
struct node *insert(int x, struct node *T)
```

```
1{   if(T == NULL)
2   {   T=new struct node;
3       T->value=x;
4       T->left=T->right=NULL;
5       T->height=0;
6   }
7   else
    ....
```

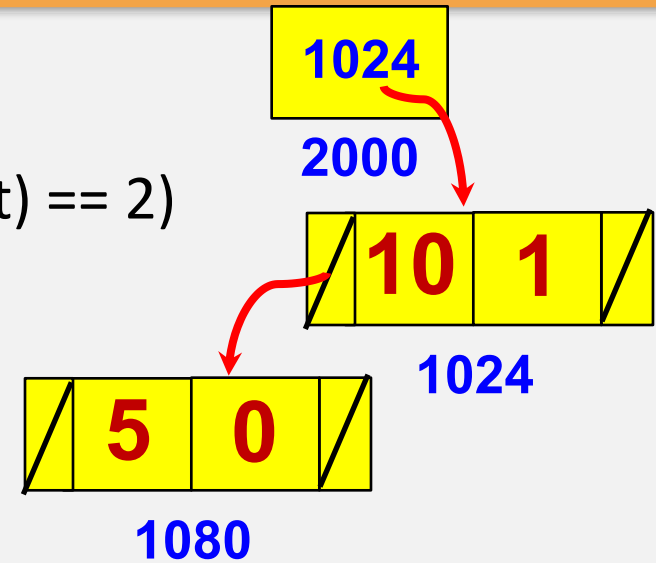




tree

```
8 if(x < T->value )
9   { T->left = insert(x,T->left);
10     if(fheight(T->left) - fheight(T->right) == 2)
11       if( x < T->left->value)
12         T=srleft(T);
13       else
14         T=drleft(T);
15   }
16ถึง 29 if(x > T->value )
17   { ข้างขวาบ้าง ให้เขียนเอง ... }
30 T->height = max(fheight(T->left), fheight(T->right)) + 1;
31 return T;
}
```

2, NULL



(B) T = 1080 ค้างบรรทัด 9 หน้า 28



Recursive 2

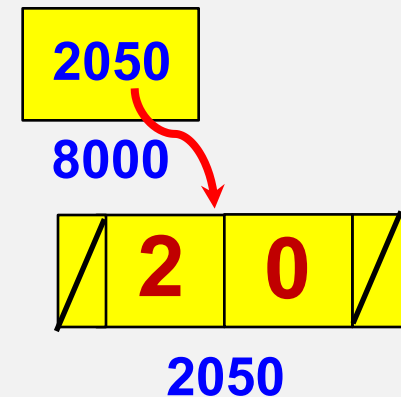
2, NULL

```
struct node *insert(int x, struct node *T)
```

```
1{   if(T == NULL)
2   {   T=new struct node;
3       T->value=x;
4       T->left=T->right=NULL;
5       T->height=0;
        }
7   else
        ....
```



tree



```
30   T->height =  $\max(-1, -1) = -1$   $+1$  fheight(T->right)) + 1;
31   return T;
}
```

กลับไปที

(B) T = 1080 ค้างบรรทัด 9

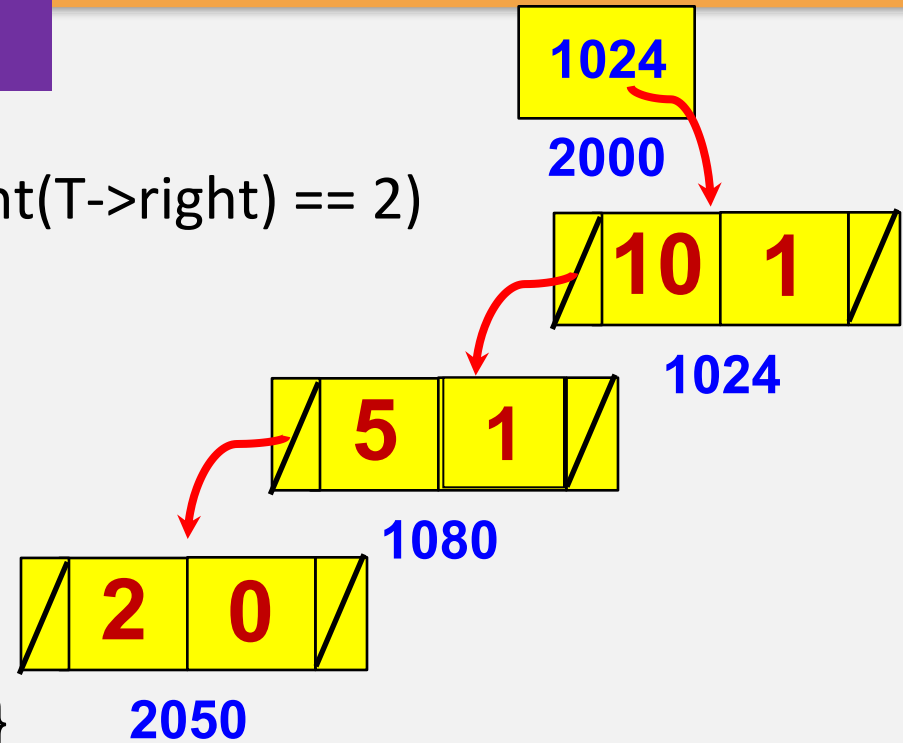
(B) T = 1080 ค้างบรรทัด 9

tree

30

```
8 if(x < T->value )
9     { T->left = insert(x,T->left);
10         if(fheight(T->left) - fheight(T->right) == 2)
11             if( x < T->left->value)
12                 T=srleft(T);
13             else
14                 T=drleft(T);
15     }
16ถึง 29 if(x > T->value )
17     { ช้างขวาบ้าง ให้เขียนเอง ... }
30 T->height = max(fheight(T->left), fheight(T->right)) + 1;
31 return T;
}
```

2050





(A) ค้างบรรทัด 9 1024

– Tree

tree

31

8 if(x < T->value)

1080

9 { T->left = insert(x,T->left);

10 if(fheight(T->left) - fheight(T->right) == 2)

11 if(x < T->left->value)

12 T=srleft(T);

13 else

14 T=drleft(T);

15 }

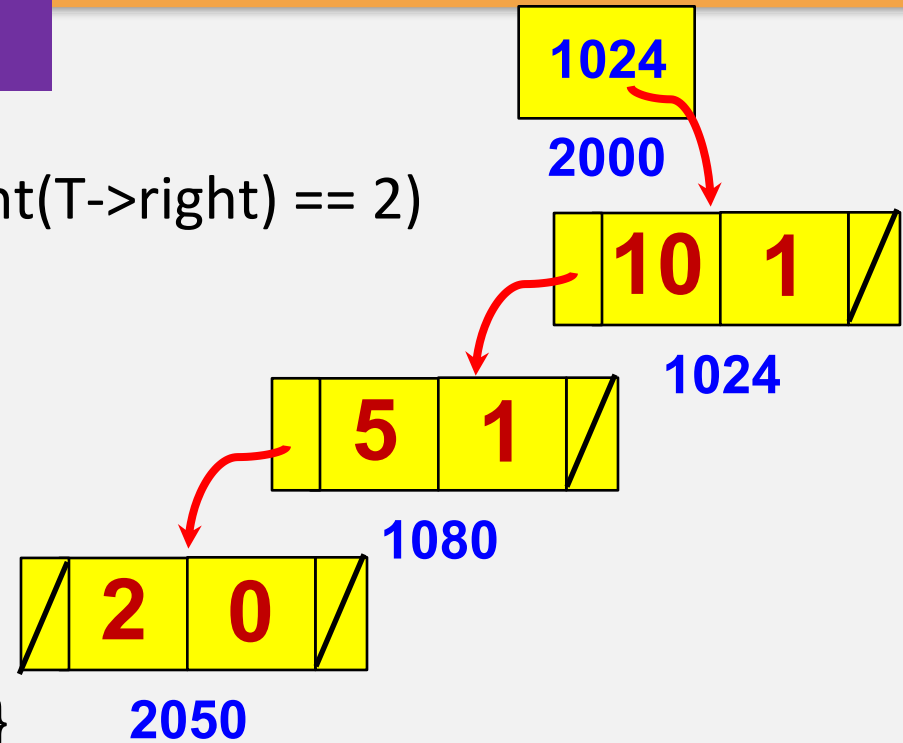
16ถึง 29 if(x > T->value)

{ ข้างขวาบ้าง ให้เขียนเอง ... }

30 T->height = max(fheight(T->left), fheight(T->right)) + 1;

31 return T;

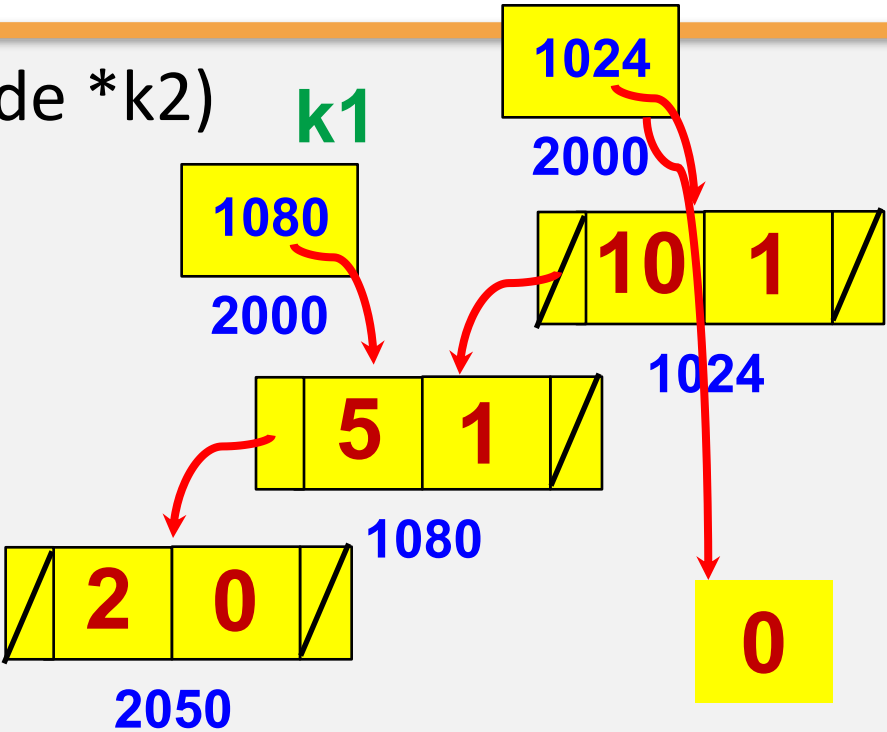
}



```

struct node *srleft(struct node *k2)
{
    struct node *k1;
    k1=k2->left;
    k2->left = k1->right;
    k1->right = k2;

```



```

    k2->height = max(fheight(k2->left), fheight(k2->right)) + 1;
    k1->height = max(fheight(k1->left), k2->height) + 1;
    return k1;
}

```



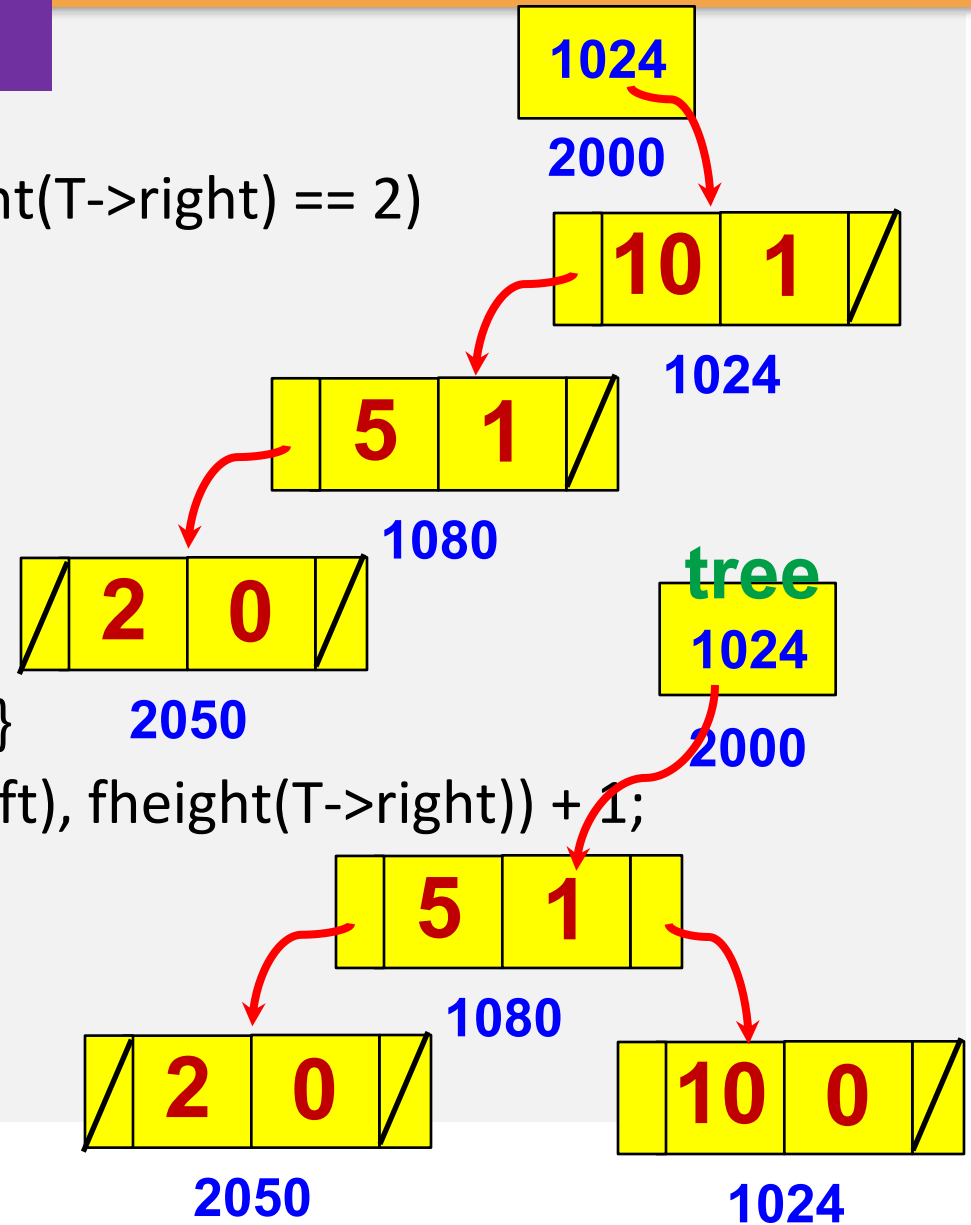

A ค้างบรรทัด 9 1024 le3 - Tree

33

tree

```
8 if(x < T->value )
9   { T->left = insert(x,T->left);
10     if(fheight(T->left) - fheight(T->right) == 2)
11       if( x < T->left->value)
12         T=srleft(T);
13       else
14         T=drleft(T);
15   }
16ถึง 29 if(x > T->value )
17   { ช้างขวาบ้าง ให้เขียนเอง ... }
30 T->height = max(fheight(T->left), fheight(T->right)) + 1;
31 return T;
}
```

1080





```
struct node *drleft(struct node *k3)
{
    k3->left = srright(k3->left);
    return srleft(k3);
}
```

3.5 Splay Trees

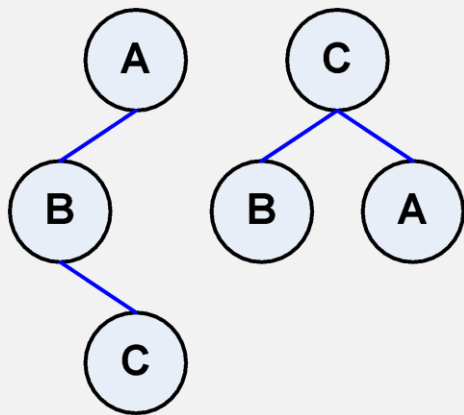
- $O(N)$ worst-case time per operation for binary search trees.
- After a node is accessed, it is pushed to the root



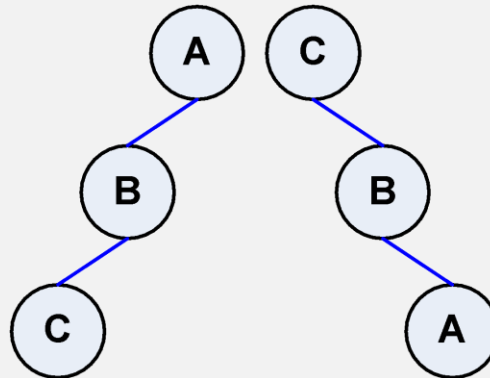
Rotation idea

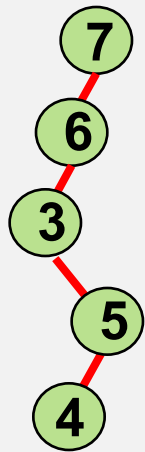
Transform

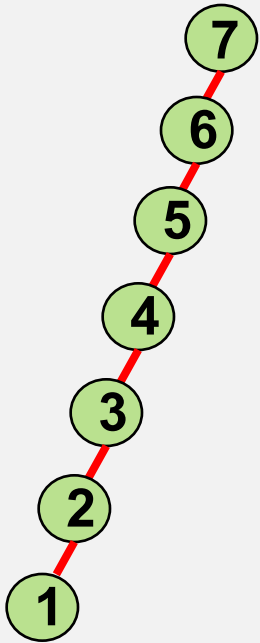
1. zig-zag



2. zig-zig

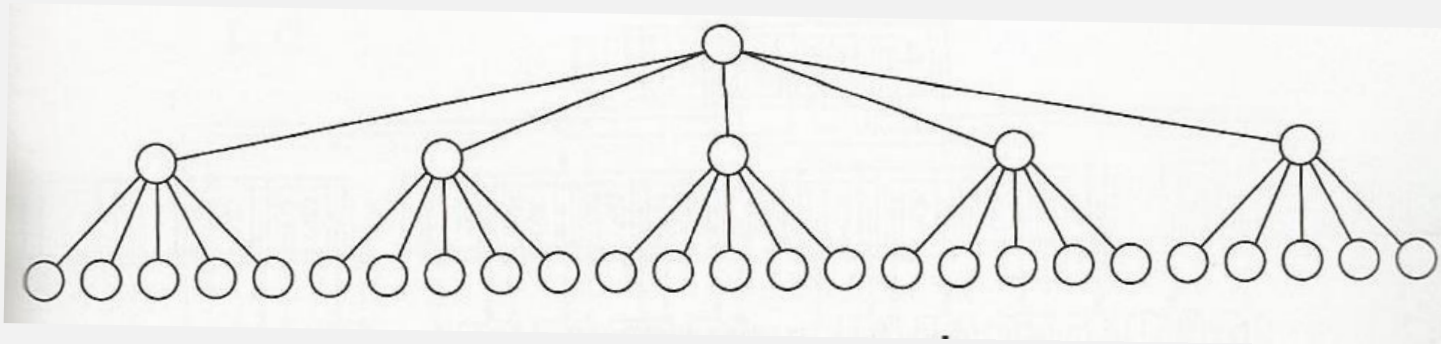
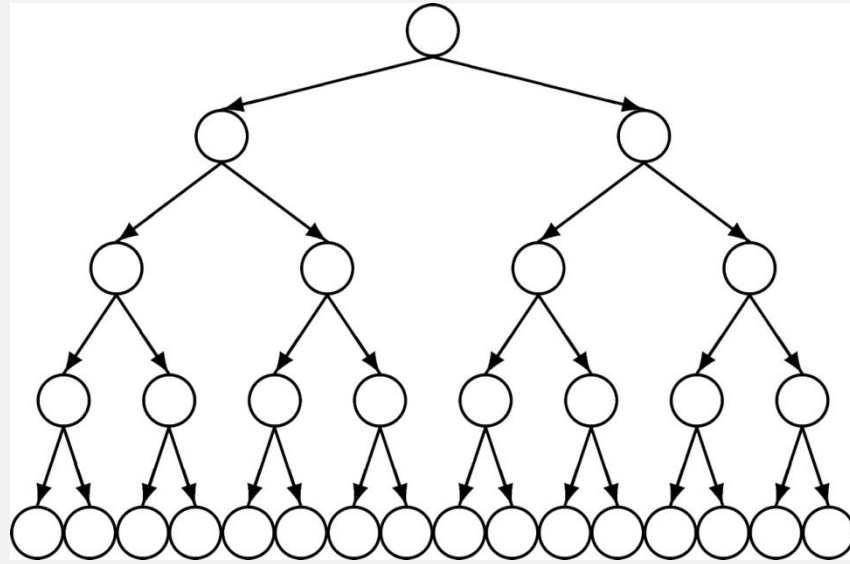






3.6 B-Trees

Complete binary
Tree 31 nodes





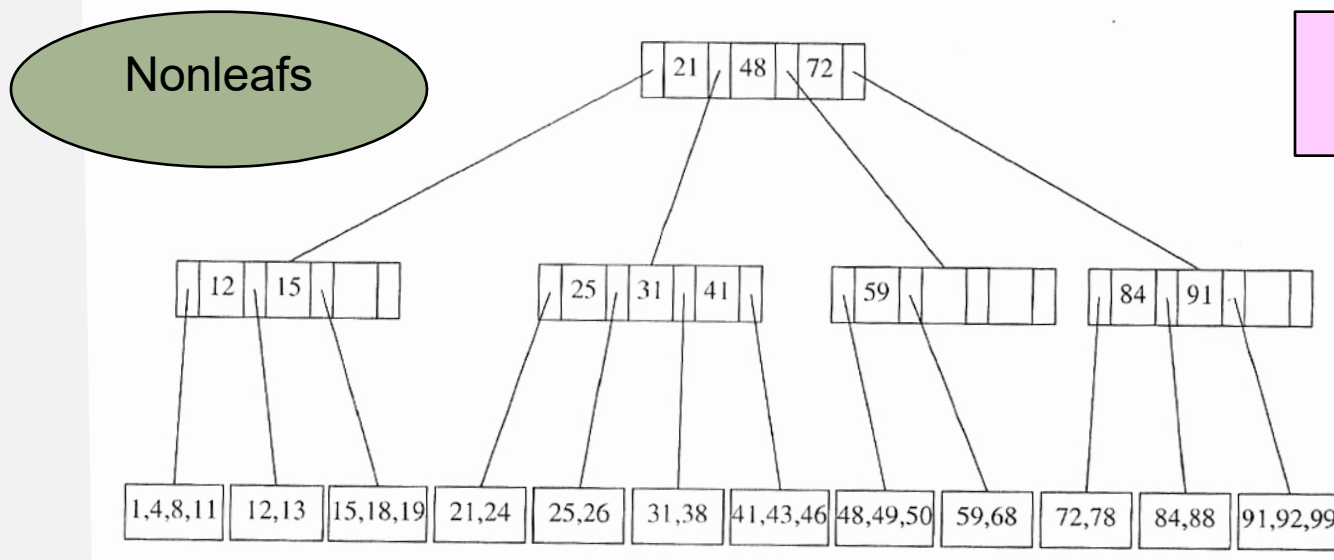
- สำหรับข้อมูลขนาดใหญ่
- ส่วนใหญ่ใช้ในระบบฐานข้อมูลและระบบไฟล์



3.6 B-Trees

A B-tree of **order m** is a tree with the following structural properties:

- 1) The data items are stores at leaves.
- 2) The nonleafs nodes store up to $M-1$ keys to guide the searching; key i represents the smallest key in subtrees $i+1$

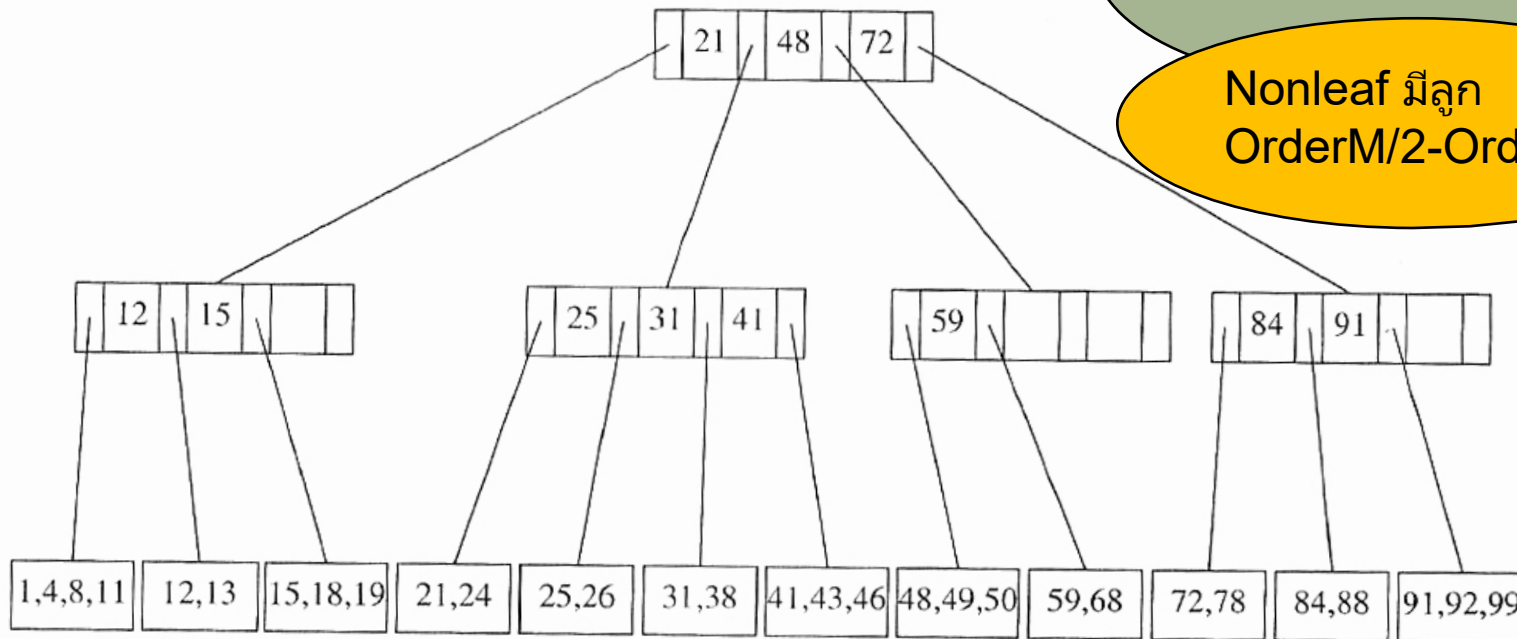




- 3) The root is either a leaf or has between two and M children.
- 4) All nonleaf nodes (Except the root) have between $\lceil M/2 \rceil$ and M children

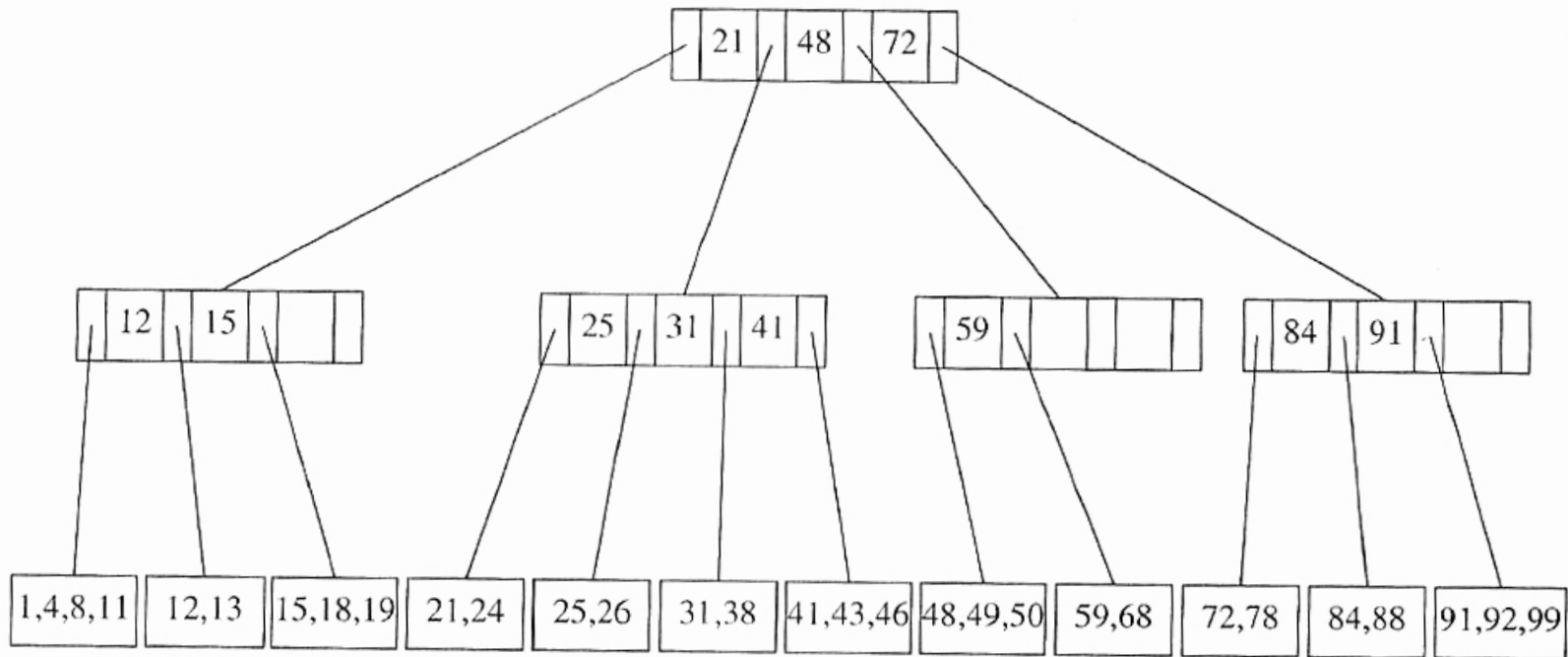
Root มีลูกจำนวน
2-OrderM

Nonleaf มีลูก
OrderM/2-OrderM





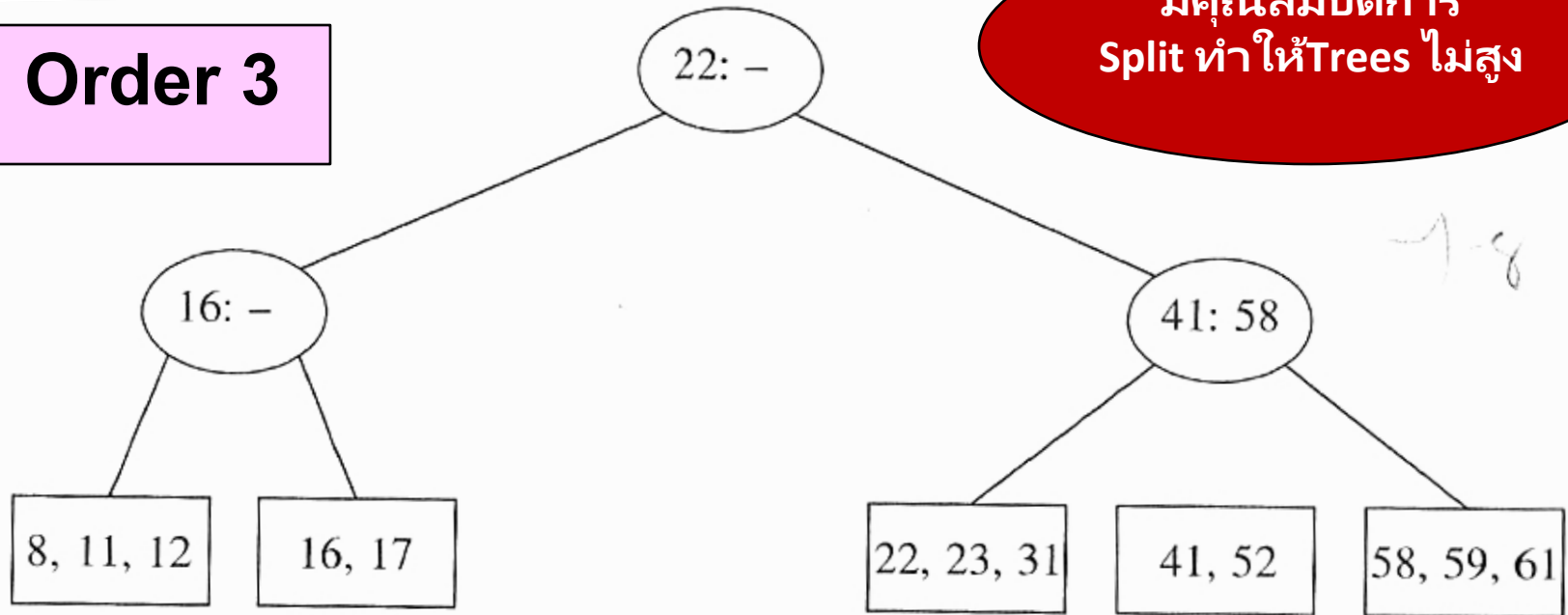
5) All leaves are at the same depth and have between $\lceil L/2 \rceil$ and L children.



leaf มีลูกจำนวน
 $L/2$ -OrderM



Order 3



Insert 18, 1, 19, 28



Order 5

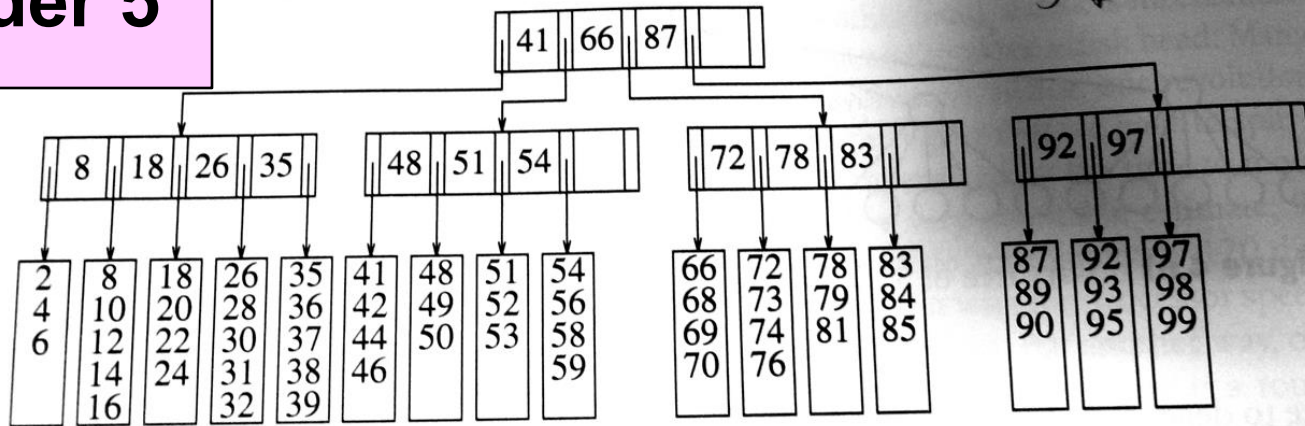
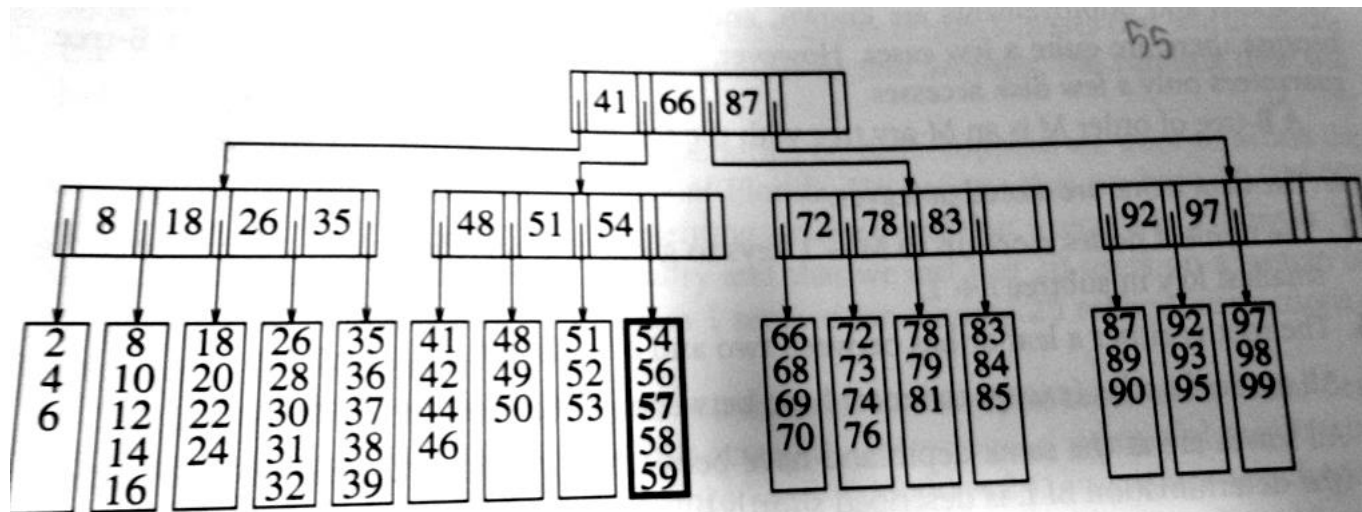


Figure 4.62 B-tree of order 5



57 is the tree in Figure 4.62

