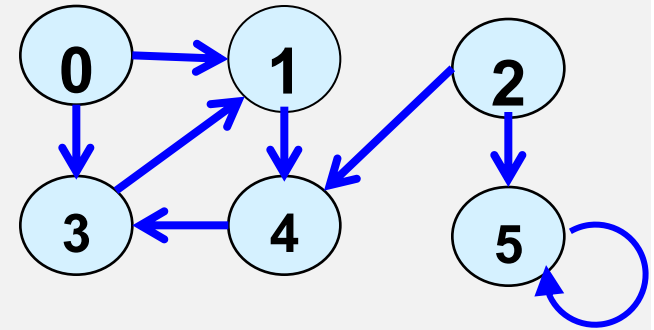




graph traversal

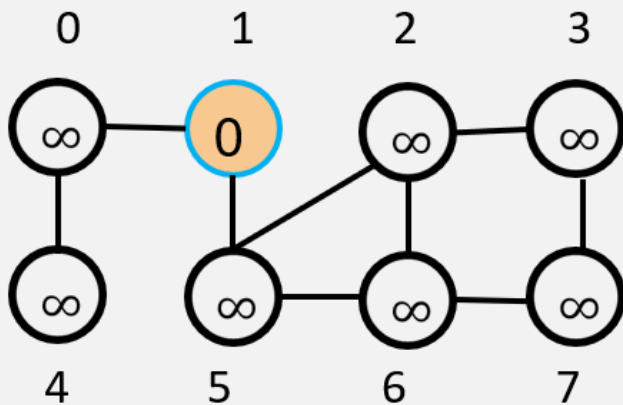
- กระบวนการเข้าไปเยือนโหนดในกราฟ
- แต่ละโหนดจะถูกเยือนเพียงครั้งเดียว
- ในกราฟระหว่างโหนดอาจจะมีหลายเส้นทาง ดังนั้นเพื่อป้องกันการท่องไปในเส้นทางที่ซ้ำเดิมจึงจำเป็นต้องทำเครื่องหมายบริเวณที่ได้เยือนเสร็จเรียบร้อยแล้วเพื่อไม่ให้เข้าไปเยือนอีก
- วิธีการท่องกราฟมี 2 แบบดังนี้
 - BFS
 - DFS





4.3 Breadth-first search (BFS) : Algorithm for searching a graph

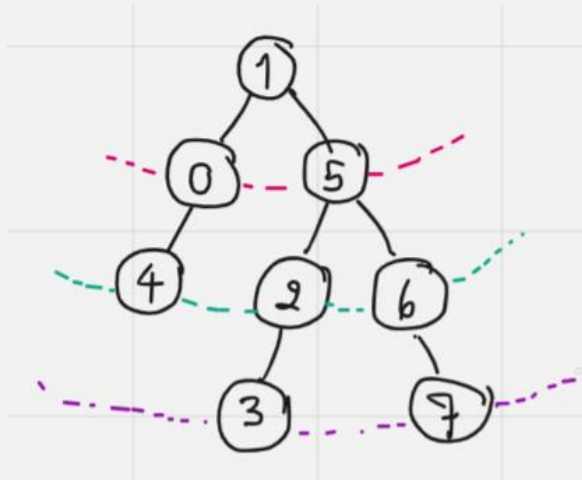
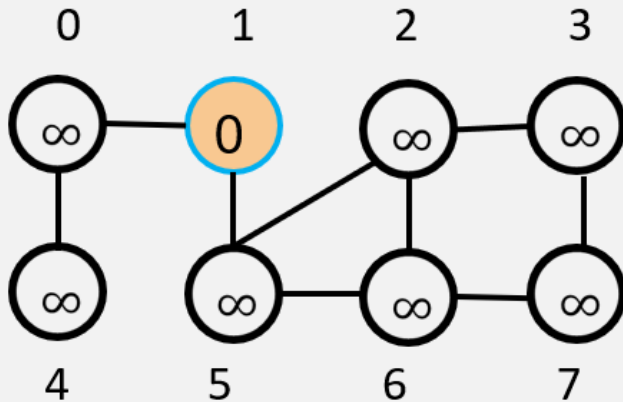
Definition : Given a graph $G = (V, E)$ and distinguished source vertex s , bfs systematically explores the edges of G to discover every vertex that is reachable from s to all such reachable vertices.



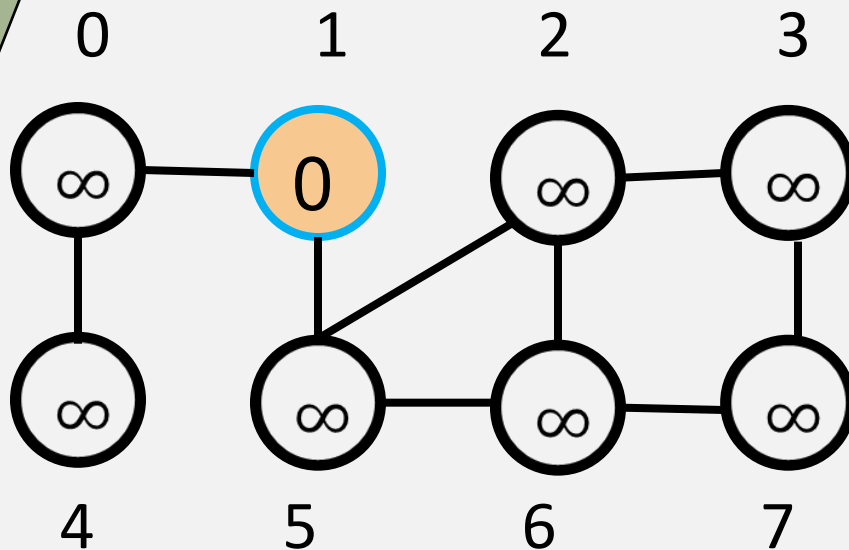


Output

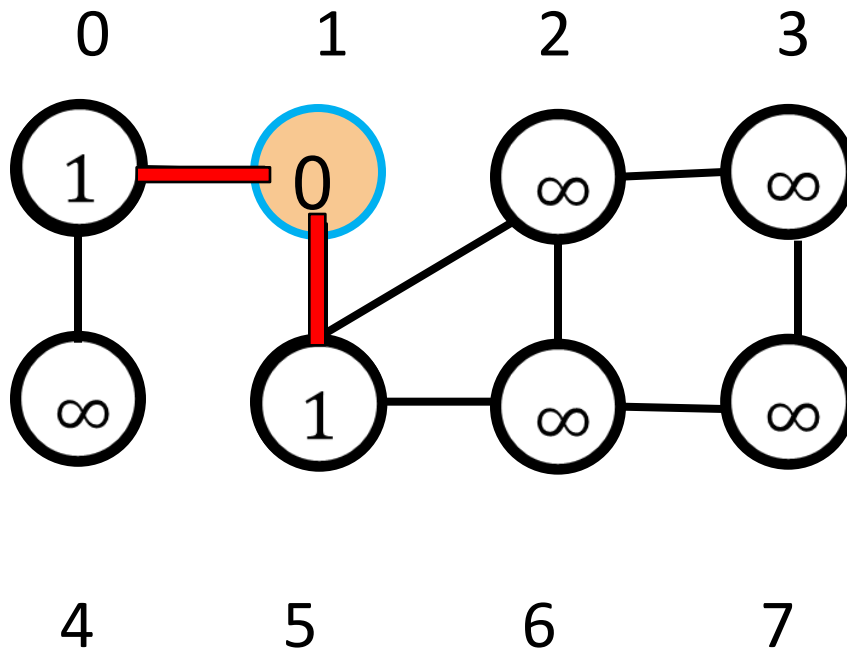
- Compute distance
- Produces a bfs tree.



ทุกเส้นทาง
มีระยะ
เท่ากัน

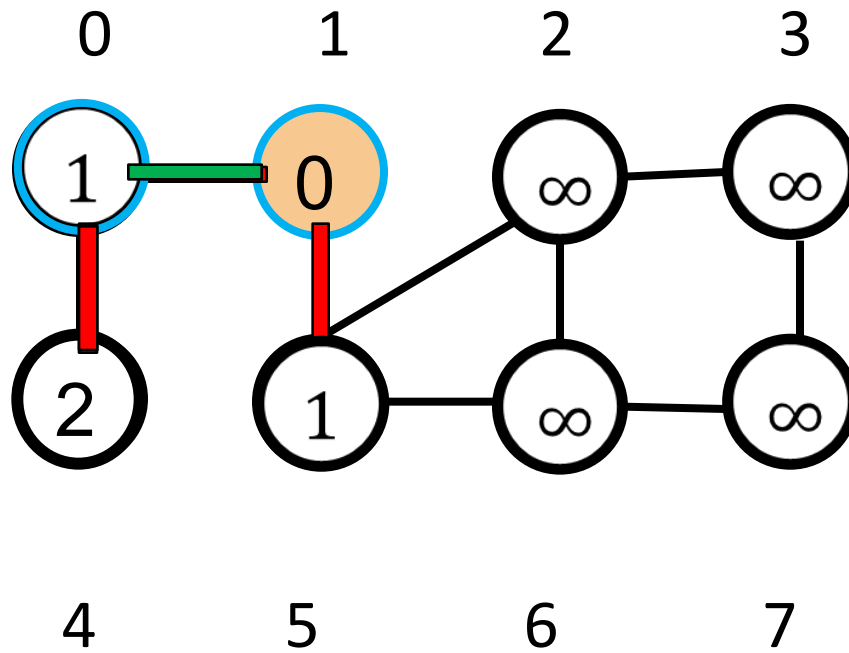


Q 1



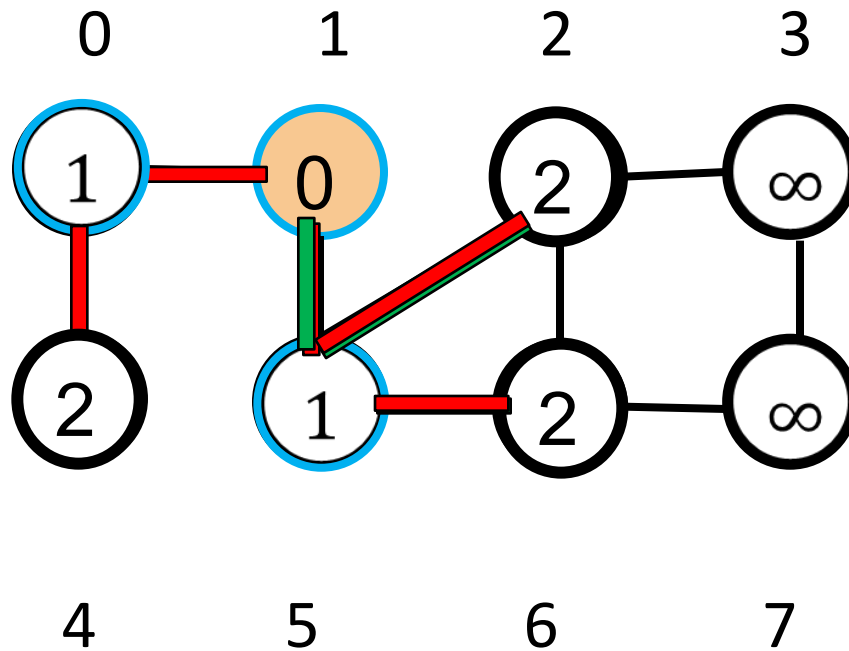
Q

1	0	5
---	---	---



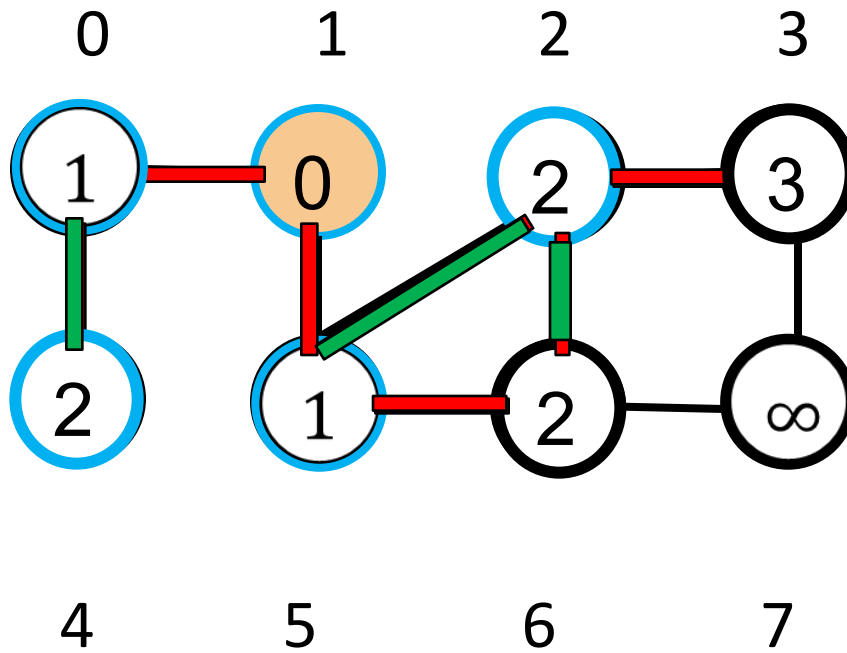
Q

0	5	4
---	---	---



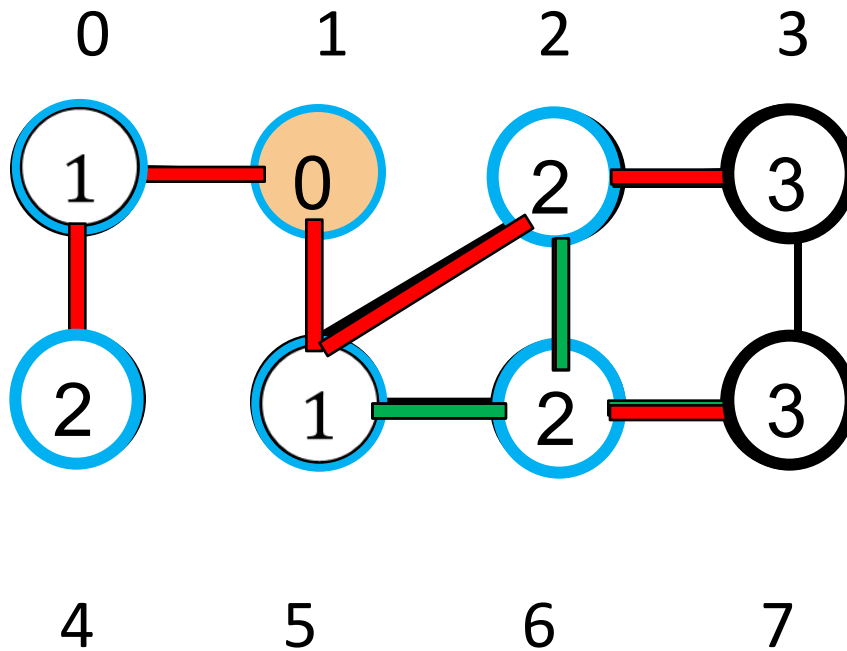
Q

5	4	2	6
---	---	---	---



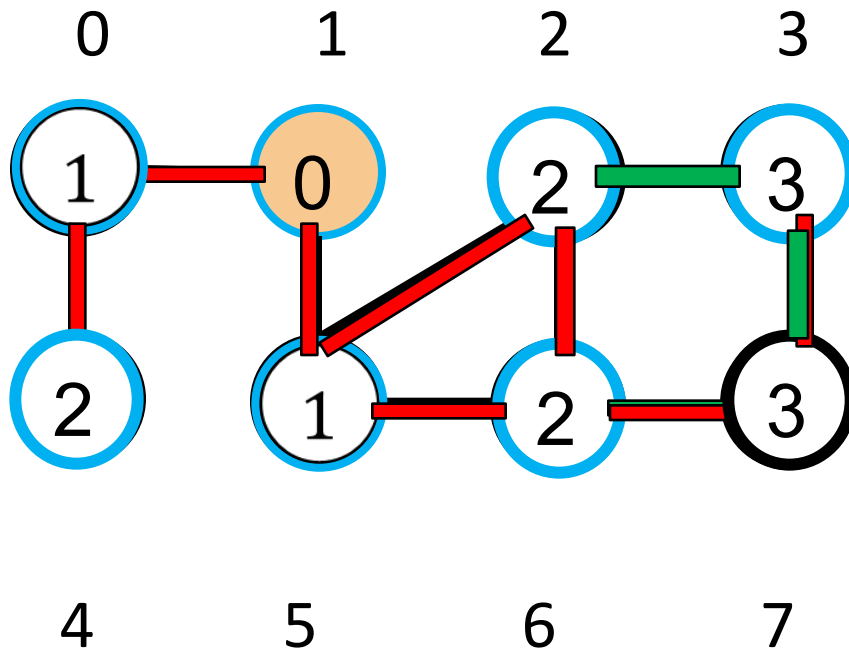
Q

4	2	6	3
---	---	---	---



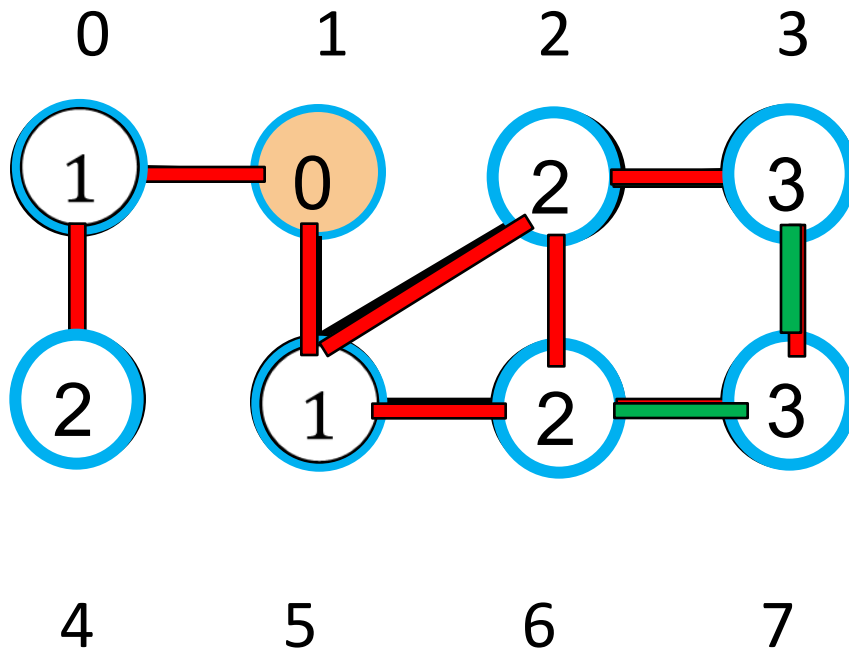
Q

6	3	7
---	---	---



Q

3	7
---	---



Q

7



4.3.1 BFS Algorithm

BFS(G, s)

for each vertex $u \in V[G] - \{s\}$

do $\text{pass}[u] = 0$

$d[u] = -1$

$\text{pred}[u] = -1$

$\text{pass}[s] = 1$

$d[s] = 0$

$\text{pred}[s] = s$



4.3.1 BFS Algorithm

BFS(G, s)

for each vertex $u \in V[G] - \{s\}$

do $\text{pass}[u] = 0$

$d[u] = -1$

$\text{pred}[u] = -1$

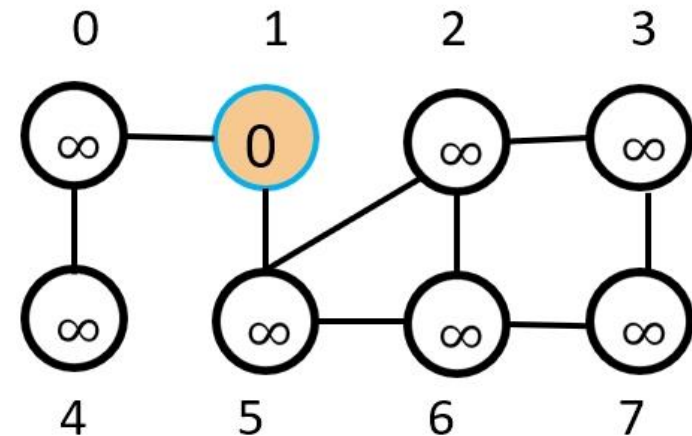
$\text{pass}[s] = 1$

$d[s] = 0$

$\text{pred}[s] = s$



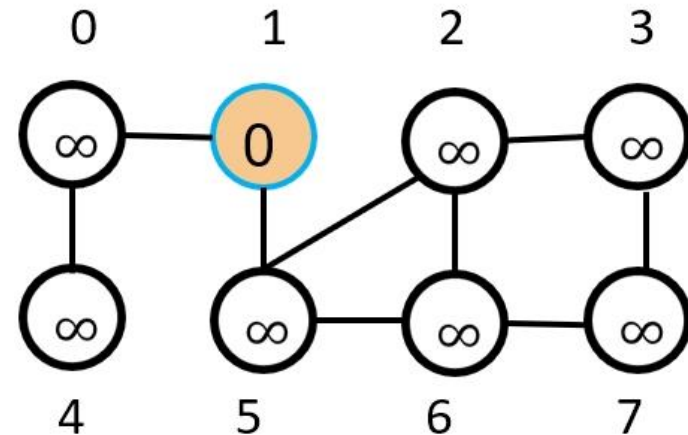
```
Enqueue(s);
while (queueSize() != 0)
{
    u = head(q)
    node = adj[u];
    while (node != NULL)
    {
        v = node->value;
        if (pass[v] == 0)
        {
            pass[v] = 1
            d[v] = d[u] + 1
            pred[v] = u
            enqueue(v)
        }
        node = node->next;
    }
    dequeue()
}
```



	pass	d	pred
0			
1			
2			
3			
4			
5			
6			
7			



```
Enqueue(s);
while (queuesize() != 0)
{
    u = head(q)
    node=adj[u];
    while( node!= NULL)
    {
        v= node->value;
        if (pass[v]=0)
        {
            pass[v] = 1
            d[v] = d[u]+1
            pred[v] = u
            enqueue(v)
        }
        node=node->next;
    }
    dequeue()
}
```



			pass	d	pred
0	<div><div></div><div>→ 1 → 4</div></div>	0	<div></div>	<div></div>	<div></div>
1	<div><div></div><div>→ 0 → 5</div></div>	1	<div></div>	<div></div>	<div></div>
2	<div><div></div><div>→ 3 → 5 → 6</div></div>	2	<div></div>	<div></div>	<div></div>
3	<div><div></div><div>→ 2 → 7</div></div>	3	<div></div>	<div></div>	<div></div>
4	<div><div></div><div>→ 0</div></div>	4	<div></div>	<div></div>	<div></div>
5	<div><div></div><div>→ 1 → 2 → 6</div></div>	5	<div></div>	<div></div>	<div></div>
6	<div><div></div><div>→ 2 → 5 → 7</div></div>	6	<div></div>	<div></div>	<div></div>
7	<div><div></div><div>→ 3 → 6</div></div>	7	<div></div>	<div></div>	<div></div>



1. Edges are explored out of the most recently discovered vertex v that still has unexplored edge leaving it.
2. When all of v 's edges have been explored, the search backtracks to explore edges leaving the vertex from which v has discovered. Until we have discovered all the vertices that are reachable from the original source vertex