

CS 367 Homework 2

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1 Question One

We are starting with a `LinkedList` class implemented as a doubly-linked chain of nodes with a dummy header node. The class uses the class `DblListNode` and has the following fields:

```
private DblListNode<E> items;
private int numItems;
```

We have been given the following code to reverse the order of the objects in a specified sublist:

```
public void reverse(int pos1, int pos2){
    // If pos1 == pos2, reversing a single list element does nothing
    // If pos1 > pos2, reversing an expty sublist does nothing
    if (pos1 >= pos2)
        return;

    // We swap the 1st and last items in the sublist,
    // then recursively reverse the remaining sublist
    // We stop when the remaining sublist has size 0 or 1

    // Swap list items at pos1 and pos2
    E temp = remove(pos2);
    add(pos2, get(pos1));
    remove(pos1);
    add(pos1, temp);

    // Now recursively reverse remainder of sublist (if any)
    // The remaining sublist is from pos1+1 to pos2-1
    reverse(pos1+1, pos2-1);
}
```

And we will now complete a second version of this method that directly changes the chain of nodes by unlinking the nodes to be swapped and re-links them into the chain in the appropriate way using only `DblListNode` methods.

```
/**
 * Reverses the order of the items from pos1 to pos2
 * NOTE: The Double Linked List will have a header node that is null.
 *
 * @param pos1 The start position of the reversal group
 * @param pos2 The end position of the reversal group
 */
public void reverse(int pos1, int pos2){
```

```

if(pos1 < 0 || pos2 < 0 || pos1 > numItems - 1 || pos2 > numItems - 1) {
    throw new IndexOutOfBoundsException();
}

if (pos1 >= pos2 ) {
    return; //return if input doesn't make sense
}

//Access the node at pos1
Dbllistnode<E> pos1node = items.getNext();
// items is the header node so this is the first node in the list
for (int k = 0; k < pos1; k++) {
    pos1node = pos1node.getNext();
}

//Access the node at pos2
Dbllistnode<E> pos2node = pos1node.getNext();
//To get the second node, start from the first
for (int k = pos1 + 1; k < pos2; k++ ) {
    pos2node = pos2node.getNext();
}

//If pos1 is the beginning of the list, it still has a next and prev
//If pos2 is at the end of the list, it only has a prev
Dbllistnode<E> node1Prev = pos1node.getPrev();
Dbllistnode<E> node1Next = pos1node.getNext();
Dbllistnode<E> node2Prev = pos2node.getPrev();

//treat pos2 at the end of the list as a special case
if (pos2 == numItems - 1) {
    if(pos2-pos1 == 1) {
        //deal with nodes next to each other
        pos1node.setNext(null);
        pos1node.setPrev(pos2node);
        node1Prev.setNext(pos2node);
        pos2node.setNext(pos1node);
        pos2node.setPrev(node1Prev);
    } else {
        pos1node.setNext(null);
        pos1node.setPrev(node2Prev);
        node1Prev.setNext(pos2node);
        node1Next.setPrev(pos2node);
        pos2node.setNext(node1Next);
        pos2node.setPrev(node1Prev);
        node2Prev.setNext(pos1node);
    }
}

//Deal with all other cases
else {
    Dbllistnode<E> node2Next = pos2node.getNext();
    if(pos2-pos1 == 1) {
        //deal with nodes next to each other
        pos1node.setNext(node2Next);
        pos1node.setPrev(pos2node);
    }
}

```

```

        node1Prev.setNext(pos2node);
        pos2node.setNext(pos1node);
        pos2node.setPrev(node1Prev);
        node2Next.setPrev(pos1node);
    } else {
        pos1node.setNext(node2Next);
        pos1node.setPrev(node2Prev);
        pos2node.setNext(node1Next);
        pos2node.setPrev(node1Prev);
        node1Next.setPrev(pos2node);
        node1Prev.setNext(pos2node);
        node2Next.setPrev(pos1node);
        node2Prev.setNext(pos1node);
    }
    reverse(pos1+1, pos2-1);
}

```

2 Question Two

We can find the worst-case time complexity for the method from the previous question in terms of N , the list size.

All of the test statements for positions in the beginning are clearly $O(1)$. When we obtain `pos1node` and `pos2node`, we could potentially have to traverse the entire list, so this step is $O(N)$. From there, all we are doing is constant time operations. However, that is only one of the recursive iterations. Worst case, we could have `pos1` reference the first node in the list and `pos2` reference the last node in the list. In this case, we must perform $N/2$ iterations. Although each successive iteration requires 2 less operations to obtain the original nodes, we can still bound the time complexity by N iterations of $O(N)$ operations, thus it is quadratic in time, or $O(N^2)$.