



## Assignment A1

Title : Pass I of a 2 Pass Assembler

Date of Completion : 09/01/2020

### Problem Statement :

Design & suitable data structures & implement pass I of a 2-Pass Assembler for a pseudo-machine in Java using object oriented feature. Implementation should consist of a few instructions from each category & few assembler directives.

### Objectives :

- Implement Pass I of a 2 Pass Assembler
- Use Object - Oriented features for implementing pass I
- Implement instructions from each category
- Implement assembler directives as well

### Outcome :

- Student should be able to :
- Understand & implement pass I of a 2 pass assembler
- Implement pass I for instructions from each category

### Software & Hardware Requirements :

- Fedora 64 bit OS
- Eclipse IDE for Java
- i5 processor
- 4 GB RAM
- 500 GB HD

## Theory :

Assembler: A program that converts assembly language into machine code.

An assembler is a program that accepts an assembly language as the input & produces its machine code.

An assembler has to perform the following major functions:

- i. Replace Symbolic Address by Numeric Address
- ii. Replace Symbolic Operation Codes by Machine Operation Codes
- iii. Reserve Storage for instruction & data
- iv. Handle literals

## Two Pass Assembler:

Pass I : Many Assemblers are designed in 2 pass

### Pass I :-

- i. Separate the labels, mnemonics, opcode & operand file
- ii. Determine the storage requirement for every assembly language statement & update the location counter
- iii. Build the symbol table

### Pass II :-

Generate the machine code.

2-Pass Assembler can handle literals

Pass I uses the following data structures:

1. Machine Opcode Table (MOT)
2. Symbol Table (ST)
3. Literal Table (LT)
4. Pool Table (PT)
5. Table for storing intermediate code in internal format (I)

### 1. Machine Opcode Table

Type of Mnemonic Symbol	Value of Class Field
Imperative Statement (IS)	1
Declarative Statement (DL)	2
Assembler Directive (AD)	3
CPU Register (RG)	4
Condition Code (CC)	5

### 2. Symbol Table

i.e A symbol table entry contains :-

- label
- address
- size

### 3. Literal Table

A literal table entry contains :-

- value
- address

### 4. Pool Table

A Pool Table contains the literal number of the starting literal of each literal pool.

### 5 Intermediate Code :

IC Table entry contains :

- Type 1 : Gives the type of mnemonic symbol

- Code 1 : Gives the index of mnemonic symbol

### Algorithm :

- $LC = 0$ ;  $iPT$  also stores address of strings of characters

- $iPT = 0$ ;

- $PT[iPT] = 0$ ;

- $iLT = 0$ ;

- Open the source file;

- While ('end of the source file') {

- Read the next line

Store the line in character array nextline  
after substituting special character with blank

- if (nextline is an END statement) then  
break;

- if nextline contains a label then  
store label in Symbol Table

- Separate opcode & operands

- if nextline is declaration statement then

- Type 1 = type of declaration statement

- Code 1 = code (index in MOT)

- size = size of memory area req. by D

- Generate intermediate code & store in

- $LC = LC + size$

- if nextline is imperative statement then

- Type 1 = type of imperative statement

- Code 1 = code (index in MOT)

- size = size of instruction

- If operand is literal then

literal is stored in literal table

v. Generate intermediate code & store in IC table  
vi.  $LC = LC + \text{size}$

g) If nextline is LTORG statement then

i. Literals of current pool are assigned addresses

ii.  $iPT = iPT + \text{size}$

iii.  $PT[iPT] = iT$

iv.  $LC$  modified accordingly.

h) If nextline is START or ORIGIN then

i.  $LC = \text{address specified in statement}$

iv. Generate intermediate code & store the same in IC table

i) If nextline is EQU statement then

i. Address of variable is converted in symbol table  
New address is extracted from statement

ii. Generate intermediate code & store the same in IC table.

7. Processing of END statement

i. Perform  $G(g)$

ii. Go to pass II.

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<u>Test Cases</u>			
Description	Input	Output	Conclusion
1 General IS statement	MOVER AREG,A	(13,04)   (5,01)	Success
2 START statement	START 200	LC set to vi	Success
3 LTORG statement	LTORG	Address assigned to literals	Success
4 EQU statement	BACK: EQU 100	Back is given address of loop	Success
5 END statement	END	Address assigned to literals & Program moves to Pass II	Success

### Conclusion:

Successful implementation of Pass I of two pass assembler using object oriented features & proper data structures.