



Theory:

Deadlock:

In multiprogramming environment, several processes may compete for a fixed number of resources. A process requests a resource B if the resource is not available at that time, it enters wait state. It may happen that it will never gain access to a resource, if that resource is being held by other waiting processes. A set of processes is deadlocked if each process is waiting for a resource that is held by a deadlocked process.

Conditions for Deadlock:

1. Mutual Exclusion
2. Hold & Wait
3. No preemption
4. Circular Wait Condition.

Strategies to deal with deadlock:

1. Ignore the problem altogether.
2. Detection & Recovery
3. Dynamic Avoidance by Careful Resource allocation
4. Prevention by negating above 3 conditions.

more the problem (Ostrich Algorithm)

### Detection & Recovery

Detection - A deadlock detection algorithm tries to detect the presence of a cycle in a resource graph.

The algorithm terminates if

1. Cycle is found
2. Cycle doesn't exist

Method - Methods to recover from deadlock are:

1. Preemption
2. Rollback
3. Killing Processes

### Deadlock Avoidance

Deadlock avoidance approach ensures that deadlock does not arise in a system.

Approaches include:

1. Do not start a process if its demands might lead to a deadlock.
2. Do not grant incremental resources request to a process if this allocation might lead to a deadlock.

Example - Banker's Algorithm.

### Bankers Algorithm :

Bankers Algorithm is the algorithm of resource allocation derived.

Tables used by Bankers Algorithm :

1. Resources in Existence (E)
2. Resources Available (A)
3. Maximum Claim Matrix (C)
4. Current Allocation Matrix (R)

Bankers Algorithm states that a request should be granted iff granting it leads to a safe state.

State of a system - is the current allocation of resources to processes.

Safe State - is a state in which there is a way to satisfy all the currently pending requests by running the process in some order.

### Deadlock Avoidance Steps :-

When a process makes a request for a set of resources, allocation matrix is updated assuming that request is granted.

Step - 1. If System is in safe state, then

2.1. grant request for resources

2.2. reverse current allocation

2.2. block the process until safe state.

## Program Outcomes

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### Databases for Banker's Algorithm

Maximum Claim Matrix

int cm [10][10]

Initial Allocation Matrix

int am [10][10]

Resource vector for resources in existence

int rv [10]

Vector for available resources after allocation

int av [10]

Vector to read the request for grant of

int vector [10]

resource for any process

Number of requesting process

int pno

Number of resources

int nr

Number of processes

int np.

### Test Cases

Input	Expected O/P	Actual O/P	Result
Allocation Matrix	No. of Safe sequences = 6	No. of safe sequences = 6	Success
A B C 1 0 0	P2 → P1 → P3 → P4	P2 → P1 → P3 → P4	Success
6 1 2	P2 → P1 → P4 → P3	P2 → R1 → P4 → P3	
2 1 1	P2 → P3 → P1 → P4	P2 → P3 → P1 → P4	
0 0 2	P2 → P3 → P4 → P1	P2 → P3 → P4 → P1	
Claim Matrix	P2 → P4 → P1 → P3 P2 → P4 → P3 → P1	P2 → P4 → P1 → P3 P2 → P4 → P3 → P1	
A B C 3 2 2			
6 1 3			
3 1 4			
4 2 2			
Resource Vector	9 3 6		

### Conclusion

Hence Barker's Algorithm has been successfully implemented using concepts of OOP.

Thing → no back → no listing

Result: answer for 3x3 example of first and second

row: 1st row: 1 2 3 4 5 6 7 8 9

2nd row: 1 2 3 4 5 6 7 8 9

3rd row: 1 2 3 4 5 6 7 8 9

4th row: 1 2 3 4 5 6 7 8 9

5th row: 1 2 3 4 5 6 7 8 9

6th row: 1 2 3 4 5 6 7 8 9

7th row: 1 2 3 4 5 6 7 8 9

8th row: 1 2 3 4 5 6 7 8 9

9th row: 1 2 3 4 5 6 7 8 9