Braids and the Jones polynomial

Thesis presentation

Apoorv Potnis

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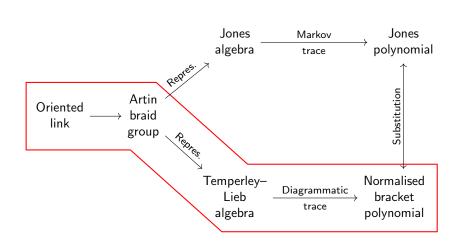


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Three dimensional representation

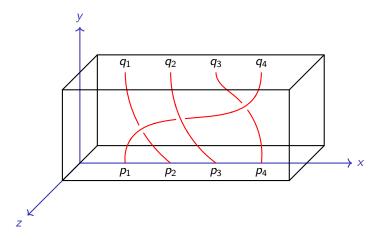


Figure: Three dimensional geometric representation of a braid

Two dimensional representation

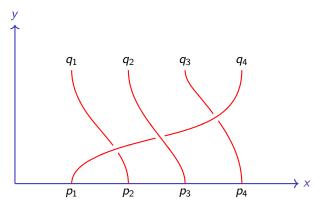


Figure: A projection of the braid

Multiplication of braids

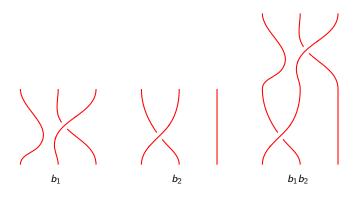


Figure: Multiplication of two braids

The identity braid \mathbf{I}_n

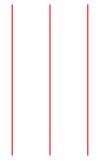


Figure: The identity I₃

Inverse of braids

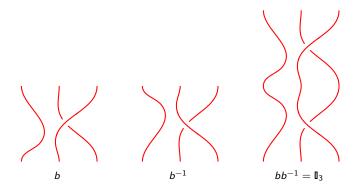


Figure: Inverse of a braid



Subsection 2

Generators and relations

Generators of the braid group

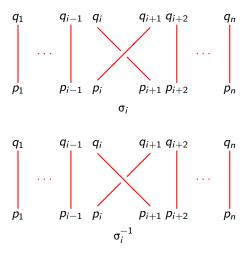


Figure: Generators σ_i and σ_i^{-1}

Type II move: $\sigma_i \sigma_i^{-1} = \mathbf{I}_n$

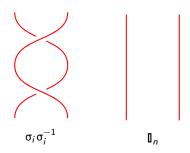


Figure: A type II move illustrating $\sigma_i \sigma_i^{-1} = \mathbb{I}_n$

Type III move: $\sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1}$

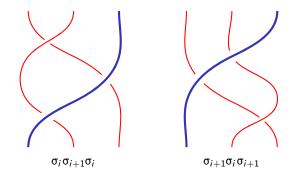


Figure: A type III move illustrating $\sigma_i \sigma_{i+1} \sigma_i = \sigma_{i+1} \sigma_i \sigma_{i+1}$

Sliding of crossings: $\sigma_i \sigma_j = \sigma_j \sigma_i$

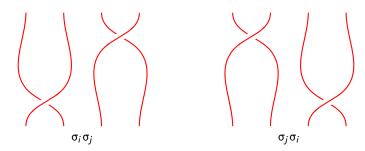


Figure: Sliding of crossings illustrating $\sigma_i \sigma_j = \sigma_j \sigma_i$

Subsection 3

Algebraic definition

Presentation of the braid group

The Artin braid group B_n admits the following presentation on the generators σ_i , for $1 \le i \le n-1$.

$$\mathsf{B}_n = \left\langle \begin{array}{ccc} \sigma_1, \dots, \sigma_{n-1} & \sigma_i \sigma_i^{-1} & = & \mathbf{I}_n \\ \sigma_i \sigma_{i+1} \sigma_i & = & \sigma_{i+1} \sigma_i \sigma_{i+1} & \text{if } i+1 \leq n-1 \\ \sigma_i \sigma_i & = & \sigma_i \sigma_i & \text{if } |i-j| \geq 2 \end{array} \right\rangle$$

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Closure of a braid

Closure of a braid \overline{b}

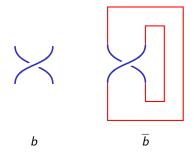


Figure: Closure of a braid

Braids and links

Every closure of a braid is a link.

Theorem (Alexander)

Every link is ambient isotopic to a closure of a braid.

Subsection 2

Equivalence of closures of braids

Conjugation

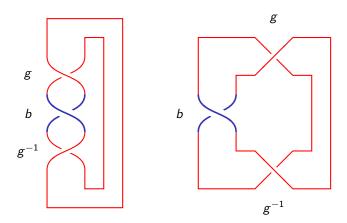


Figure: Conjugation process illustrating the link equivalence of $\overline{gbg^{-1}}$ and \overline{b} (part 1)

Conjugation

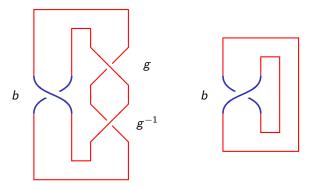


Figure: Conjugation process illustrating the link equivalence of $\overline{gbg^{-1}}$ and \overline{b} (part 2)

Markov theorem

Theorem (Markov)

Two braids whose closures are ambient isotopic to each other are related by a finite sequence of the following operations.

- Braid equivalences, i.e. equivalences resulting due to the braid relations.
- 2. Conjugation.
- 3. Markov moves.

Markov move

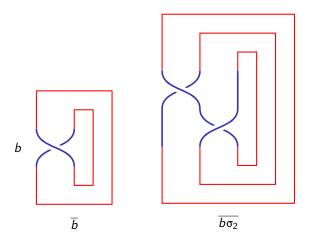


Figure: Markov move with $b = \sigma_1^{-1}$.

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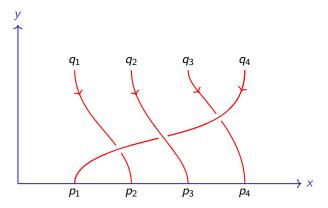
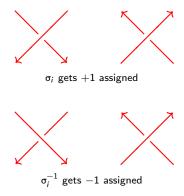


Figure: A projection of the braid with downward orientation

For the consistency of the orientation, it must be either upwards or downwards for all strands.



Writhe is the sum of the assigned numbers.

Bracket polynomial of a braid:

$$\langle \cdot \rangle \colon \mathsf{B}_n \to \mathbb{Z}[A, A^{-1}]$$

 $\langle \cdot \rangle \colon b \mapsto \langle \overline{b} \rangle$

 $\langle\cdot\rangle$ is well defined and invariant under conjugation.

Normalisation using writhe:

$$L(K) := (-A^3)^{-w(b)} \langle \overline{b} \rangle$$

$$\left\langle \left| \cdots \right| \right\rangle = A \left\langle \left| \cdots \right| \right\rangle + A^{-1} \left\langle \left| \cdots \right| \right\rangle \left| \cdots \right| \right\rangle$$

$$\left| \cdots \right| \left| \bigcirc \left(\left| \cdots \right| \right| = \mathbf{I}_n \quad \text{and} \quad \mathsf{U}_i \coloneqq \left| \cdots \right| \middle| \bigcirc \left| \cdots \right|$$

 $\langle \sigma_i^{-1} \rangle = A \langle \mathsf{U}_i \rangle + A^{-1} \langle \mathbf{I}_n \rangle$

$$\langle \sigma_i \rangle = A \langle \mathbf{I}_n \rangle + A^{-1} \langle \mathsf{U}_i \rangle$$

$$\rangle + A^{-1}\langle 0 \rangle$$