

Midterm Report : Single Cycle v. Pipeline  
ECE 437 : Computer Design and Prototyping  
Katie Bauschka  
Apoorv Gaur  
Chuan Yean Tan  
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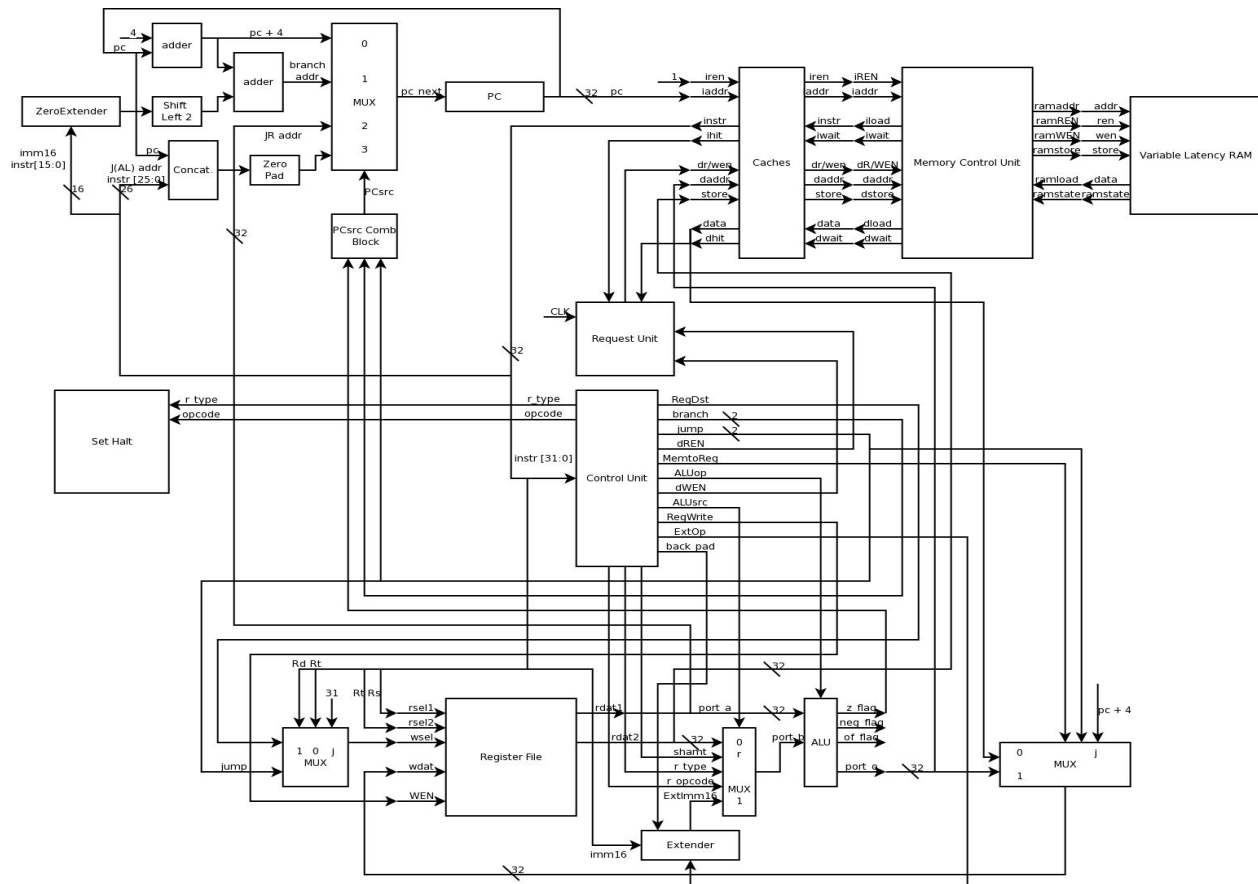
## Overview

The throughput of each design is the main feature being compared between the two designs. The biggest benefit of the pipelined processor is that the throughput of any given program will be much higher than that of the single cycle. More instructions can be executed at a time, rather than one instruction at a time. On the other hand, the single cycle design is comparatively simplistic. There is no need for hazard or forwarding units.

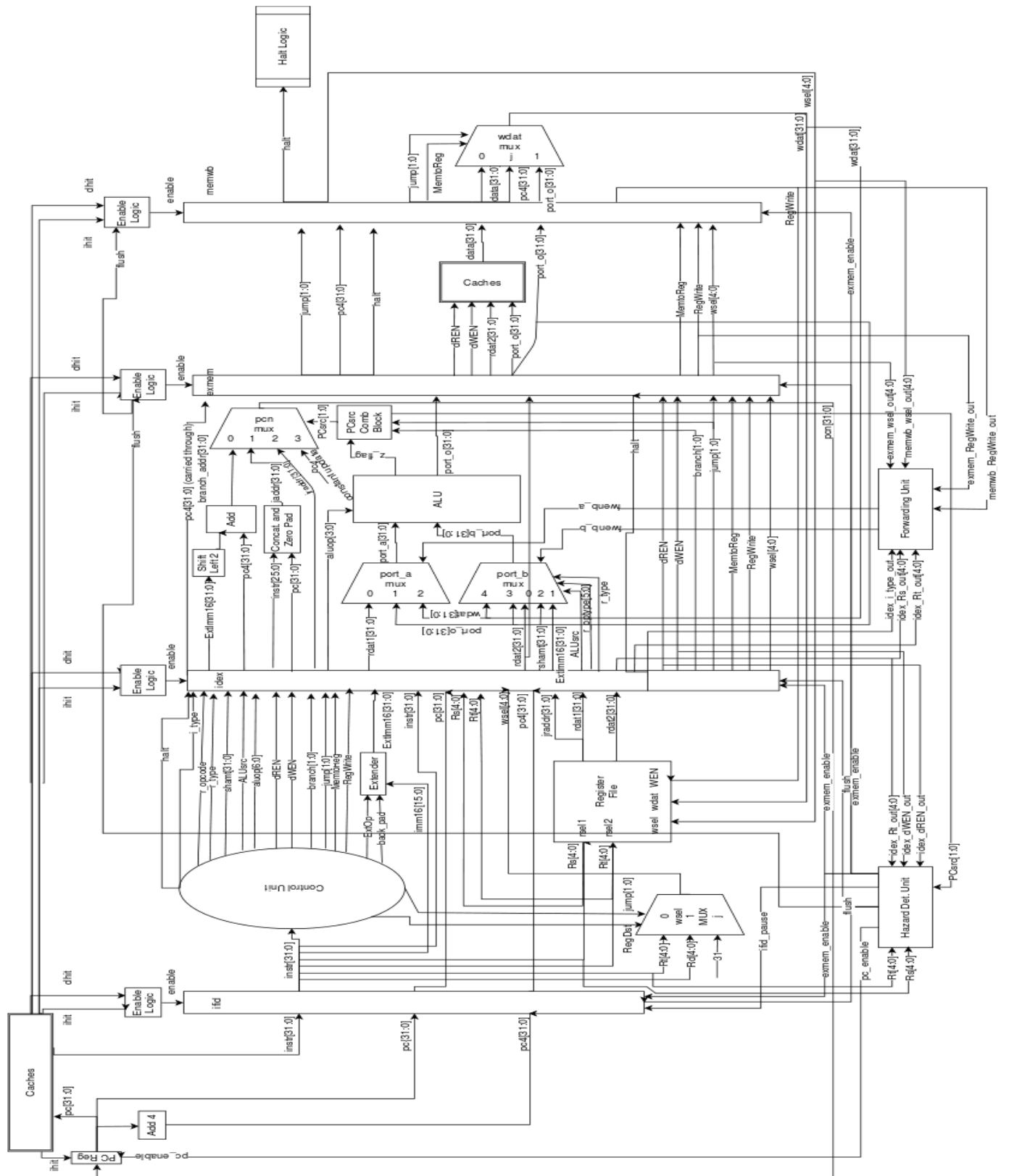
We are using mergesort.asm to test our two processor designs. This program will test all the types of instructions. For the pipeline design, it thoroughly tests the effectiveness of the hazard and forwarding units. At the end of the testing, we expect to see that the pipeline design has a higher throughput than the single cycle design.

## Processor Design

## Single Cycle



# Pipeline



## **Results**

### Single Cycle Processor

	Single Cycle Processor
Frequency (maximum possible)	$\min(36.46, 34.395) = 34.395 \text{ MHz (CLK)}$
Frequency (highest achieved)	1 / (10ns)
Latency (of one instruction)	$1 / 34.395\text{MHz} = 0.029\text{ms}$
Performance in MIPS*	0.0135 (13791 cycles)
Length of critical path	27.360ns
FPGA resources	
Total Combinational functions	2781
Total Registers	1279

### Pipeline Processor

	Pipeline Processor
Frequency (maximum possible)	$\min(76.91, 50.77) = 50.77 \text{ MHz (CLK)}$
Frequency (highest achieved)	1/ (5ns)
Latency (of one instruction)	$5 * (1/50.77\text{MHz}) = 0.0984\text{ms}$
Performance in MIPS*	0.00271 (20221 cycles)
Length of critical path	19.615ns
FPGA Resources	
Total Combinational functions	3384
Total Registers	1725

\*Formula Used:

$$MIPS = \frac{n \cdot 10^{-6}}{latency \cdot cycles}$$

where n is total number of instructions used by mergesort (5399).

## **Conclusions**

As seen from the tables above, a pipelined processor can operate at a higher frequency because fewer instructions need to be executed to per clock cycle.

The latency, however, increases for a pipelined processor as compared to a single cycled processor because while trying to execute several instructions at once, we introduce stalls and flushes to avoid overlaps, forwarding and hazarding issues. This is expressed in performance in MIPS, too, as an increase in latency leads to fewer number of instructions per cycle.

The length of the critical path decreases, which is explained by the increase in number of registers that are used in a pipelined processor. This is also depicted by the increase in number of registers used for a synthesized simulation (FPGA Resources) of the pipeline processor as compared to a single cycle processor.

## **Contributions**

### Single Cycle Processor

- Block Diagram - Katie
- Final Design - Katie

### Pipeline Processor

- Block Diagram - Katie
- Final Design
  - Hazard Unit - Katie and Apoorv
  - Forwarding Unit - Katie and Apoorv
  - Datapath - Apoorv and Katie

### Midterm Report Result Tables - Apoorv