

# 1B Paper 6: Communications

## Handout 4: Digital Baseband Modulation

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## Data Transmission

We have seen how analogue sources can be digitised. E.g., An MPEG or QuickTime file is a **stream of bits**



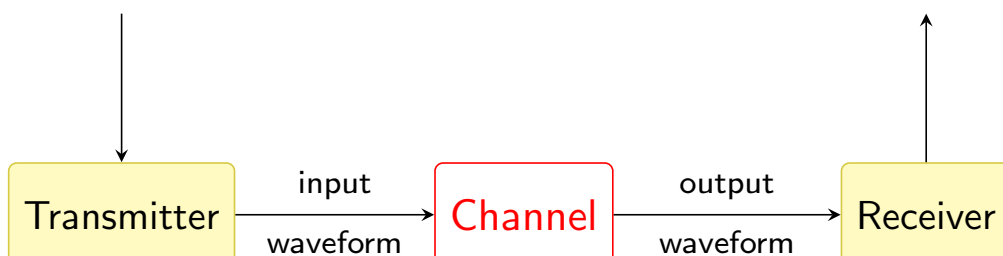
↔ ...10110010001101010...

Now we have to transport those bits across a channel:

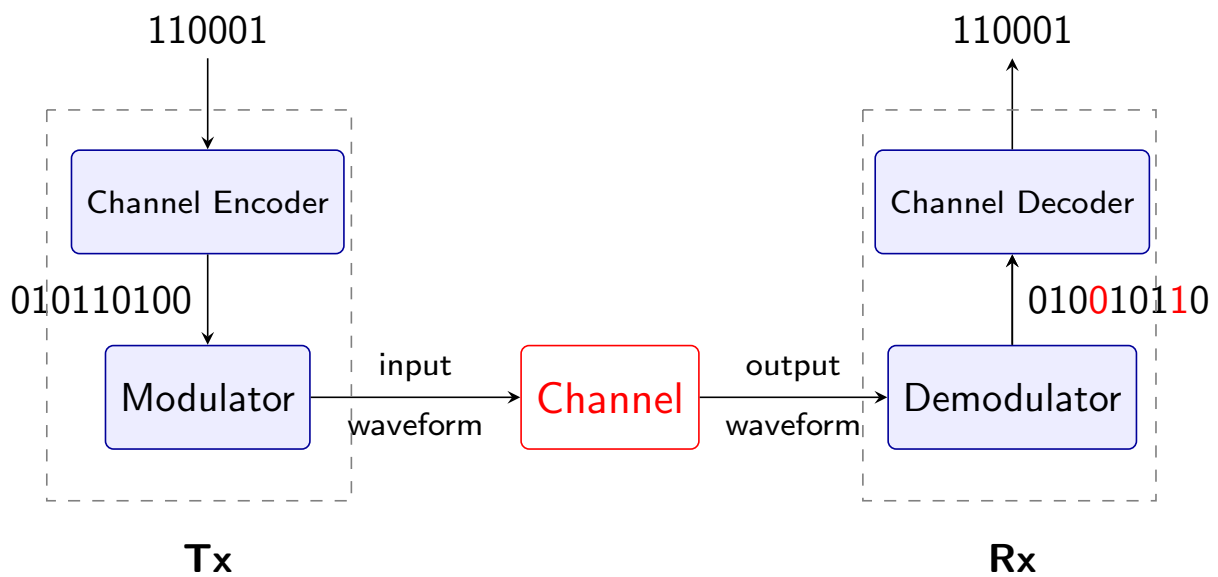
(Digitised source)

110 001 100 111

110 000 100 111



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The transmitter (Tx) does two things:

1. **Encoding:** Adding redundancy to the source bits to protect against noise
2. **Modulation:** Transforming the coded bits into waveforms.

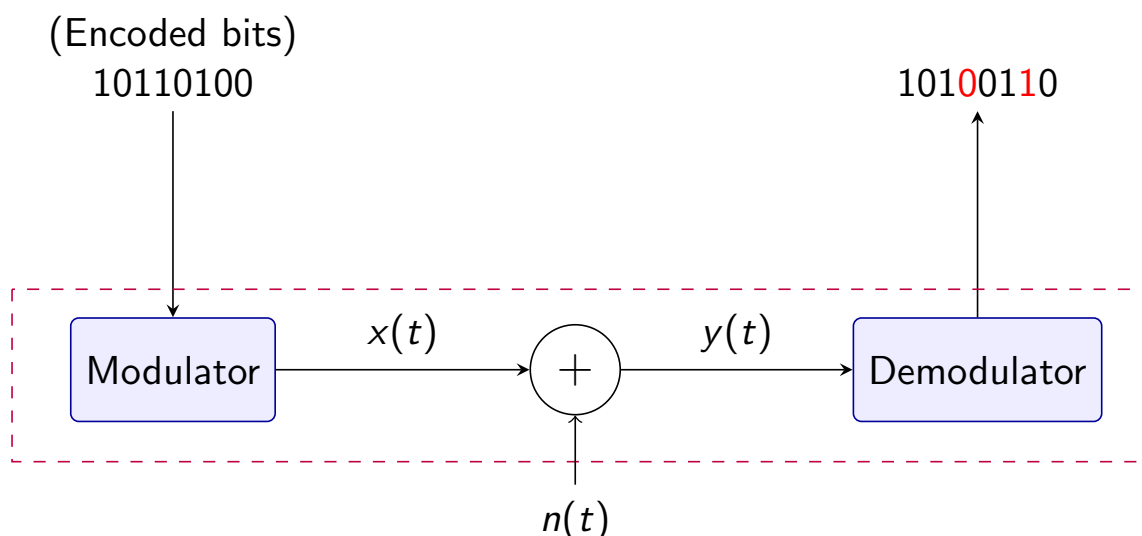
The receiver (Rx) does:

- **Demodulation:** noisy output waveform  $\rightarrow$  output bits
- **Decoding:** Try to correct errors in the output bits and recover the source bits

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## Modulation/Demodulation

We'll first consider the modulation and demodulation blocks assuming that the channel encoder/decoder are fixed, and look at the design of the channel encoder and decoder later.



We now study a digital baseband modulation technique called *Pulse Amplitude Modulation* (PAM) & analyse its performance over an Additive White Gaussian Noise (AWGN) channel

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# The Symbol Constellation

The digital modulation scheme has two basic components.

1. The first is a mapping from bits to real/complex numbers, e.g.

$$0 \rightarrow -A, \quad 1 \rightarrow A \quad (\text{binary symbols})$$

$$00 \rightarrow -3A, \quad 01 \rightarrow -A, \quad 10 \rightarrow A, \quad 11 \rightarrow 3A \quad (\text{4-ary symbols})$$

The set of values the bits are mapped to is called the *constellation*, e.g., the 4-ary constellation above is  $\{-3A, -A, A, 3A\}$ .

Once we fix a constellation, a sequence of bits can be uniquely mapped to constellation symbols. E.g., with constellation  $\{-A, A\}$

$$0 \ 1 \ 0 \ 1 \ 1 \ 1 \ 0 \ 0 \ 1 \ 0 \rightarrow -A, A, -A, A, A, A, -A, -A, A, -A$$

With constellation  $\{-3A, -A, A, 3A\}$ , the same sequence of bits is mapped as 01 01 11 00 10  $\rightarrow -A, -A, 3A, -3A, A$

In a constellation with  $M$  symbols, each symbol represents  $\log_2 M$  bits

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## The Pulse Shape

2. The second component of Pulse Amplitude Modulation is a unit-energy *baseband* waveform denoted  $p(t)$ , called the *pulse shape*. E.g., a sinc pulse or a rect pulse:

$$p(t) = \frac{1}{\sqrt{T}} \operatorname{sinc}\left(\frac{\pi t}{T}\right) \quad \text{or} \quad p(t) = \begin{cases} \frac{1}{\sqrt{T}} & \text{for } t \in \left(-\frac{T}{2}, \frac{T}{2}\right] \\ 0 & \text{otherwise} \end{cases}$$

$T$  is called the *symbol time of the pulse*

A sequence of constellation symbols  $X_0, X_1, X_2, \dots$  is used to generate a *baseband* signal as follows

$$x_b(t) = \sum_k X_k p(t - kT)$$

Thus we have the following **important** steps to associate bits with a baseband signal  $x_b(t)$ :

$$\dots 0 \ 1 \ 0 \ 1 \ 1 \ 1 \ 0 \ 0 \ 1 \ 0 \dots \rightarrow X_0, X_1, X_2, \dots \rightarrow \sum_k X_k p(t - kT)$$

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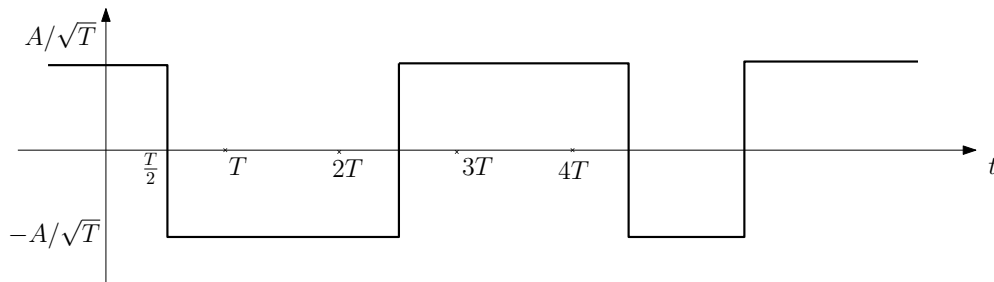
# Rate of Transmission

The modulated baseband signal is  $x_b(t) = \sum_k X_k p(t - kT)$ .

With the rect pulse shape

$$p(t) = \begin{cases} \frac{1}{\sqrt{T}} & \text{for } t \in \left(-\frac{T}{2}, \frac{T}{2}\right] \\ 0 & \text{otherwise} \end{cases}$$

and  $X_k \in \{+A, -A\}$ ,  $x_b(t)$  looks like



Every  $T$  seconds, a new symbol is introduced by shifting the pulse and modulating its amplitude with the symbol.

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The *transmission rate* is  $\frac{1}{T}$  **symbols/sec** or  $\frac{\log_2 M}{T}$  **bits/second**

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## Desirable Properties of the Pulse Shape $p(t)$

$p(t)$  is chosen to satisfy the following important objectives:

1. We want  $p(t)$  to decay quickly in time, i.e., the effect of symbol  $X_k$  should not start much before  $t = kT$  or last much beyond  $t = (k + 1)T$
2. We want  $p(t)$  to be approximately band-limited.  
For a fixed sequence of symbols  $\{X_k\}$ , the spectrum of  $x_b(t)$  is

$$X_b(f) = \mathcal{F} \left[ \sum_k X_k p(t - kT) \right] = P(f) \sum_k X_k e^{-j2\pi f k T}$$

*Hence the bandwidth of  $x_b(t)$  is the same as that of the pulse  $p(t)$*

3. The retrieval of the information sequence from the *noisy* received waveform  $x_b(t) + n(t)$  should be simple and relatively reliable. In the absence of noise, the symbols  $\{X_k\}_{k \in \mathbb{Z}}$  should be recovered perfectly at the receiver.

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# Orthonormality of pulse shifts

Consider the third objective, namely, simple and reliable detection. To achieve this, the pulse is chosen to have the following “**orthonormal shifts**” property:

$$\int_{-\infty}^{\infty} p(t - kT)p(t - mT) dt = \begin{cases} 1 & \text{if } k = m \\ 0 & \text{if } k \neq m \end{cases} \quad (1)$$

We’ll see how this property makes signal detection at the Rx simple

- This property is satisfied by the rect pulse shape

$$p(t) = \begin{cases} \frac{1}{\sqrt{T}} & \text{for } t \in \left(-\frac{T}{2}, \frac{T}{2}\right] \\ 0 & \text{otherwise} \end{cases}$$

- The sinc pulse  $p(t) = \frac{1}{\sqrt{T}} \text{sinc}\left(\frac{\pi t}{T}\right)$  also has orthonormal shifts! (You will show this in Examples Paper 9, Q.2)

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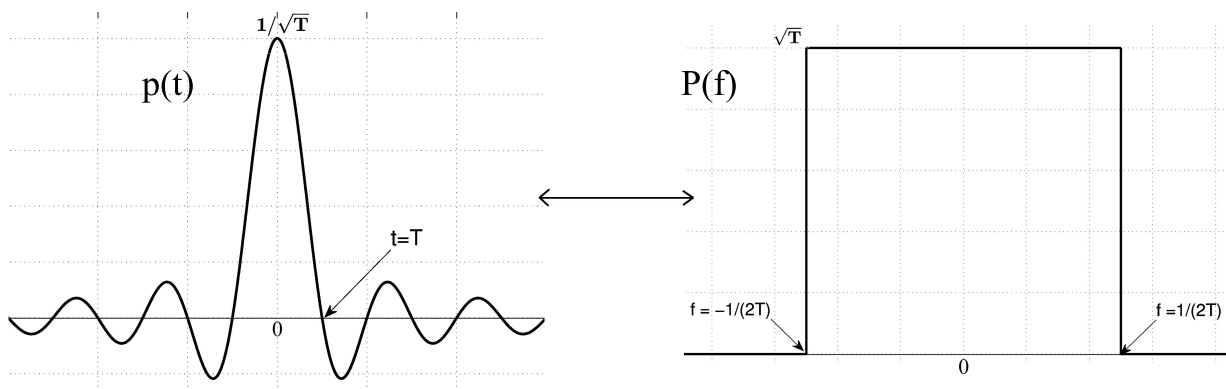
## Time Decay vs. Bandwidth Trade-off

The first two objectives say that we want  $p(t)$  to:

1. Decay quickly in time
2. Be approximately band-limited

But ... faster decay in time  $\Leftrightarrow$  larger bandwidth

Consider the pulse  $p(t) = \frac{1}{\sqrt{T}} \text{sinc}\left(\frac{\pi t}{T}\right)$

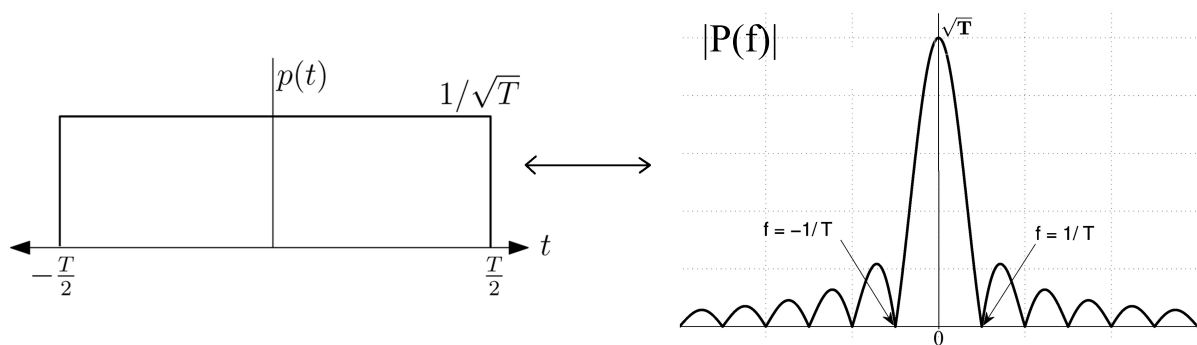


The sinc is perfectly band-limited to  $W = \frac{1}{2T}$   
But decays slowly in time  $|p(t)| \sim \frac{1}{|t|}$

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Next consider the rectangular pulse

$$p(t) = \begin{cases} \frac{1}{\sqrt{T}} & \text{for } t \in \left(-\frac{T}{2}, \frac{T}{2}\right] \\ 0 & \text{otherwise} \end{cases}$$

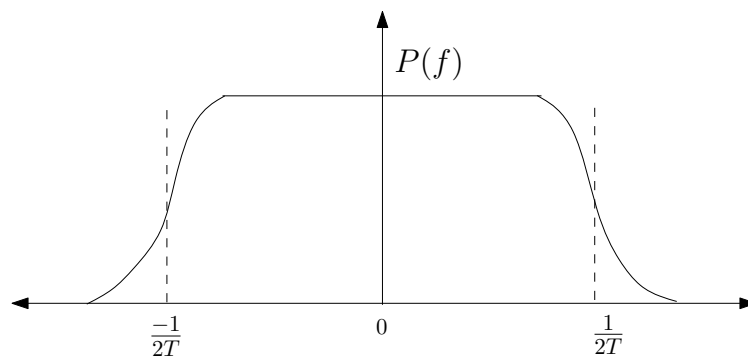


This pulse is perfectly time-limited to the interval  $[-T/2, T/2]$ .  
But ...

- Decays slowly in freq.  $|P(f)| \sim \frac{1}{|f|}$
- Main-lobe bandwidth =  $\frac{1}{T}$

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In practice, the pulse shape is often chosen to have a *root raised cosine* spectrum



Bandwidth slightly larger than  $\frac{1}{2T}$ ; decay in time  $|p(t)| \sim \frac{1}{|t|^2}$

A happy compromise!

- More on raised cosine pulses in 3F4
- For intuition, it often helps to envision  $x_b(t)$  with a rect pulse, though it is never used in practice

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Data  $\rightarrow$  constellation symbols  $\rightarrow$  continuous waveform

Thus we have the following **important** steps to associate bits with a baseband signal  $x_b(t)$ :

$$\dots 0101110010\dots \longrightarrow X_0, X_1, X_2, \dots \longrightarrow \sum_k X_k p(t - kT)$$

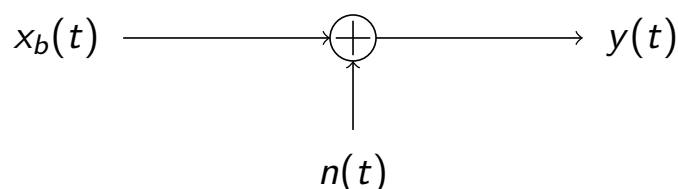
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## PAM Demodulation

Now, assume that we have picked a constellation and a pulse shape satisfying the objectives, and we transmit the baseband waveform

$$x_b(t) = \sum_k X_k p(t - kT)$$

over a *baseband* channel  $y(t) = x_b(t) + n(t)$



How does the receiver recover the information symbols  $\{X_0, X_1, X_2, \dots\}$  from  $y(t)$ ?

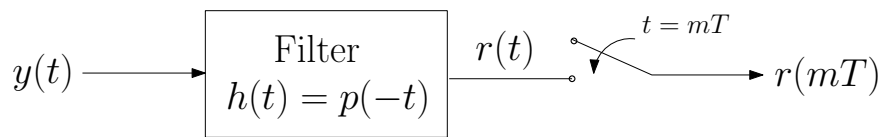
- This process is called *demodulation*
- We will see that the orthonormal shift property of  $p(t)$  leads to a simple and elegant demodulator

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# Matched Filter Demodulator

Let us first understand the operation assuming no noise, i.e.,

$$y(t) = x_b(t) = \sum_k X_k p(t - kT)$$



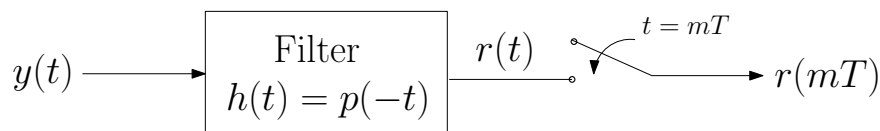
$y(t)$  is passed through a filter with impulse response  $h(t) = p(-t)$

This is called a **matched filter**. The filter output is

$$\begin{aligned} r(t) &= y(t) \star h(t) = x_b(t) \star h(t) \quad (\text{assuming no noise}) \\ &= \int_{-\infty}^{\infty} x_b(\tau) h(t - \tau) d\tau \\ &= \int_{-\infty}^{\infty} x_b(\tau) p(\tau - t) d\tau \\ &= \sum_k X_k \int_{-\infty}^{\infty} p(\tau - kT) p(\tau - t) d\tau \end{aligned}$$

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## Matched filter output



$$\begin{aligned} r(t) &= \sum_k X_k \int_{-\infty}^{\infty} p(\tau - kT) p(\tau - t) d\tau \\ &= \sum_k X_k \int_{-\infty}^{\infty} p(u + t - kT) p(u) du \quad (\text{using } u = \tau - t) \\ &= \sum_k X_k g(t - kT) \end{aligned}$$

where

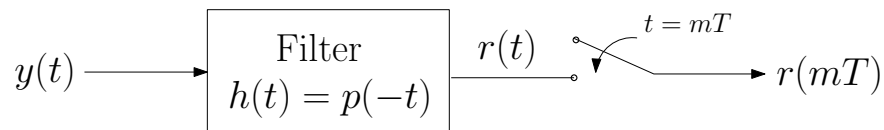
$$g(t) = \int_{-\infty}^{\infty} p(u + t) p(u) du$$

Matched filter output  $r(t)$  is of the form as the PAM signal, except that the 'pulse' is now  $g(t)$

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## Sampled matched filter output



By sampling the filter output at time  $t = mT$ ,  $m = 0, 1, 2, \dots$ , you get

$$r(mT) = \sum_k X_k g((m - k)T)$$

Because of the *orthonormal shifts* property of  $p(t)$

$$g((m - k)T) = \int_{-\infty}^{\infty} p(u + (m - k)T) p(u) du = \begin{cases} 1 & \text{if } k = m \\ 0 & \text{if } k \neq m \end{cases}$$

Therefore,

$$r(mT) = \sum_k X_k g((m - k)T) = X_m$$

*Orthonormal shifts* property is crucial for this demodulator to work!

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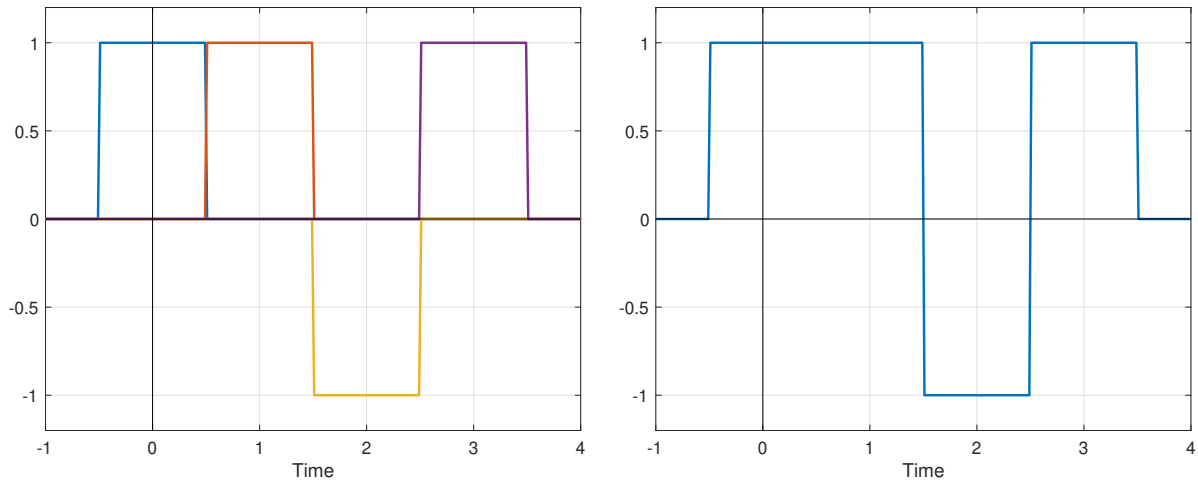
For two different choices of pulse  $p(t)$ , we now visualize

- the transmitted PAM signal  $\sum_k X_k p(t - kT)$ , and
- the matched filter output  $r(t) = \sum_k X_k g(t - kT)$

## Example 1: Transmitted signal

Rectangular pulse:  $p(t) = \begin{cases} \frac{1}{\sqrt{T}} & \text{for } t \in \left(-\frac{T}{2}, \frac{T}{2}\right] \\ 0 & \text{otherwise} \end{cases}$

Assume  $T = 1$  and the symbols  $\{X_0, X_1, X_2, X_3\} = \{1, 1, -1, 1\}$

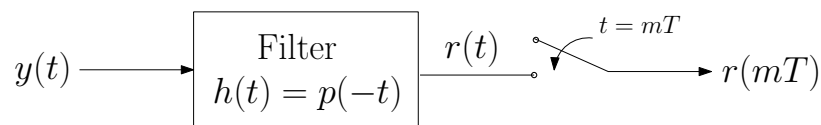


Right panel shows the transmitted PAM signal  $\sum_{k=0}^3 X_k p(t - kT)$

Left panel shows each component of the sum separately

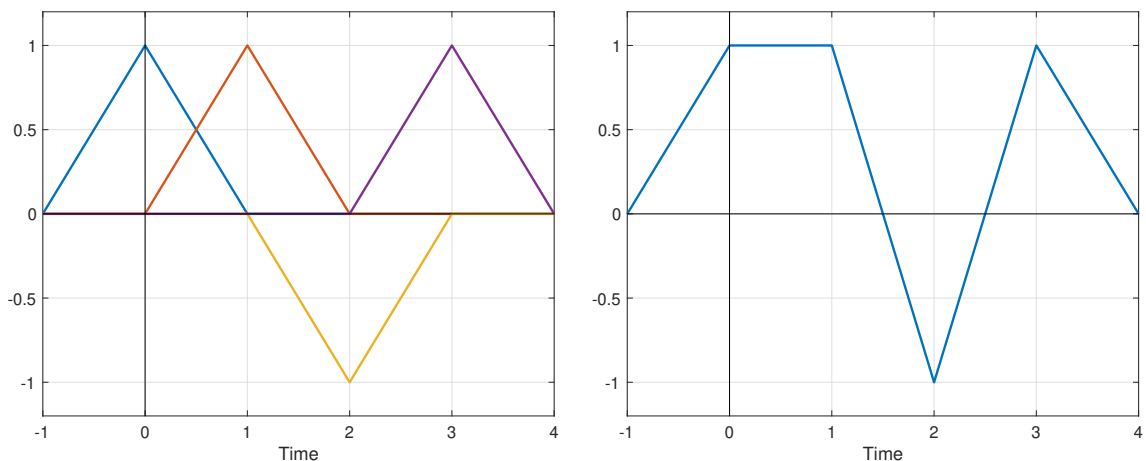
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## Example 1: Matched filter output



Matched filter output  $r(t) = \sum_{k=0}^3 X_k g(t - kT)$

$$g(t) = \int_{-\infty}^{\infty} p(u + t)p(u)du = \begin{cases} 1 + \frac{t}{T}, & -T \leq t \leq 0, \\ 1 - \frac{t}{T}, & 0 \leq t \leq T \end{cases}$$

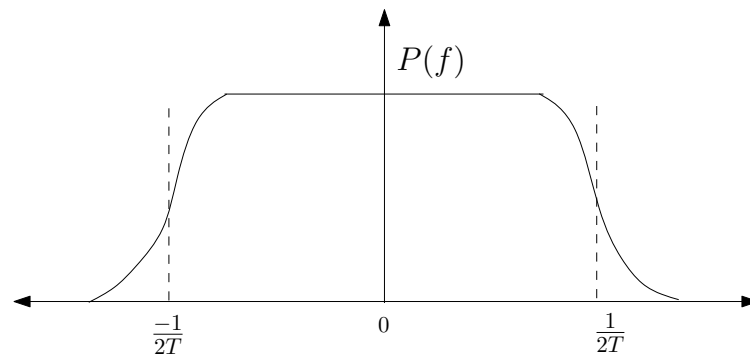


Left: each component of the sum separately, Right:  $r(t)$

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## Example 2

A practical choice: Root-raised cosine pulse  $p(t)$



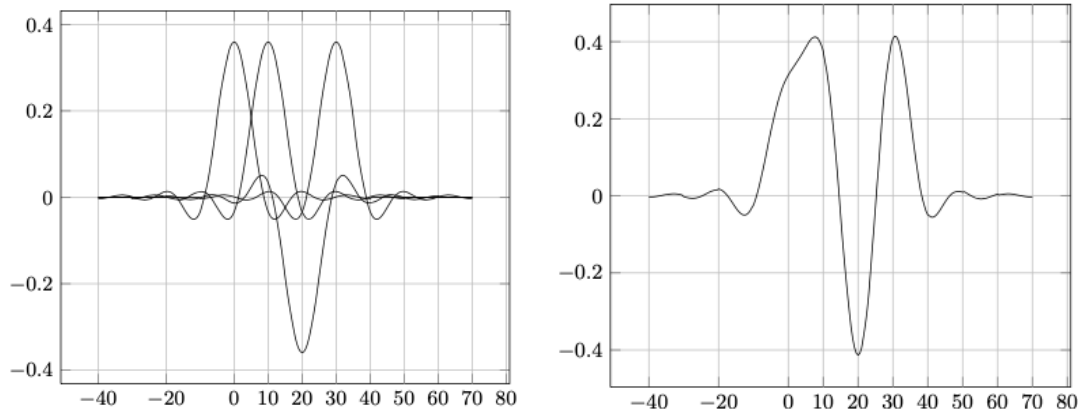
Assume  $T = 10$  and the symbols  $\{X_0, X_1, X_2, X_3\} = \{1, 1, -1, 1\}$

Let us visualise :

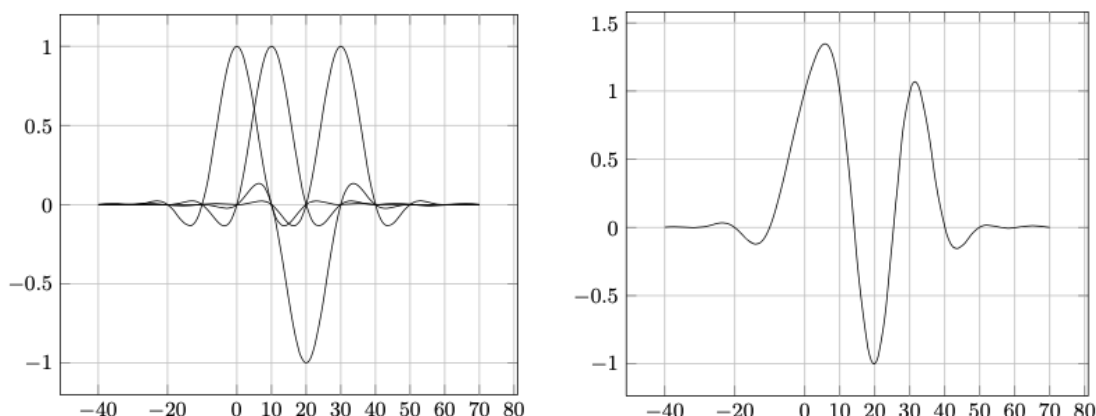
- the transmitted PAM signal  $\sum_{k=0}^3 X_k p(t - kT)$ , and
- the matched filter output  $r(t) = \sum_{k=0}^3 X_k g(t - kT)$

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Transmitted PAM signal  $\sum_{k=0}^3 X_k p(t - kT)$ :



Matched filter output  $r(t) = \sum_{k=0}^3 X_k g(t - kT)$ :

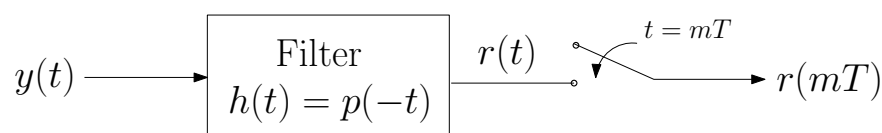


# What happens when there is noise at the receiver?

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## Demodulation with Noisy $y(t)$

Now consider the noisy case. The receiver gets  $y(t) = x(t) + n(t)$



The matched filter output is

$$\begin{aligned} r(t) &= y(t) \star h(t) = x_b(t) \star h(t) + n(t) \star h(t) \\ &= \sum_k X_k g(t - kT) + \int_{-\infty}^{\infty} n(\tau) p(\tau - t) d\tau \end{aligned}$$

Sampling at  $t = mT$ ,  $m = 0, 1, 2, \dots$ , we now get

$$r(mT) = X_m + N_m$$

where  $N_m$  is noise part of the filter output at time  $mT$ :

$$N_m = \int_{-\infty}^{\infty} n(\tau) p(\tau - mT) d\tau$$

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# Properties of the Noise

Let us denote  $r(mT)$ , the sampled output at time  $mT$ , by  $Y_m$ .

$$Y_m = X_m + N_m, \quad m = 0, 1, 2, \dots$$

Note that this is a *discrete-time channel*. We have converted the continuous-time problem into a discrete-time one of detecting the symbols  $X_m$  from the noisy outputs  $Y_m$ .

- To do this, we first need to understand the properties of the noise  $N_m$ . Recall that

$$N_m = \int_{-\infty}^{\infty} n(\tau) p(\tau - mT) d\tau$$

- $N_m$  is a *random variable* whose distribution depends on the statistics of the *random process*  $n(t)$ .

You will learn about random processes and their characterisation in 3F1 & 3F4, but this is outside the scope of this course. For now, we will directly specify the distribution of  $N_m$  and analyse the detection problem.

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$$Y_m = X_m + N_m, \quad m = 0, 1, 2, \dots$$

Modelling  $n(t)$  as a Gaussian process leads to the following **important** characterisation of  $N_m$ :

- For each  $m$ ,  $N_m$  is a *Gaussian random variable* with zero mean, and variance  $\sigma^2$  that can be estimated empirically
- Further  $N_1, N_2, \dots$  are *independent*
- Thus the sequence of random variables  $\{N_m\}, m = 0, 1, \dots$  are *independent* and *identically distributed* as  $\mathcal{N}(0, \sigma^2)$ .

## Detection

- At the Rx, how do we detect the information symbol  $X_m$  from  $Y_m$  for  $m = 0, 1, \dots$ ?
- Remember that each  $X_m$  belongs to the *symbol constellation*

# Detection for Binary PAM

Let's start with a simple binary constellation, then generalise.

Consider a constellation where each  $X_m \in \{-A, A\}$ . This is called *binary PAM* or BPSK ('Binary Phase Shift Keying')

$$Y = X + N$$

The detection problem is now:

*Given  $Y$ , how to decide whether  $X = A$  or  $X = -A$ ?*

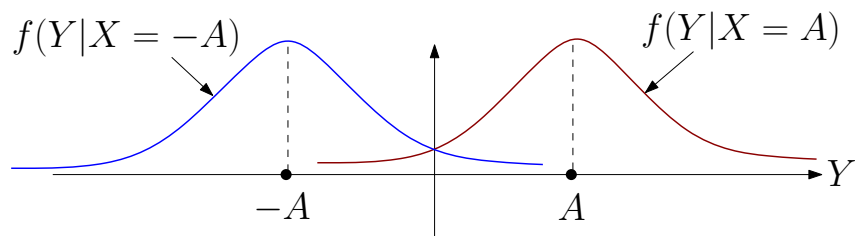
Observe that:

$$Y = A + N \text{ if } X = A \quad \text{and} \quad Y = -A + N \text{ if } X = -A$$

- $N$  is distributed as  $\mathcal{N}(0, \sigma^2)$
- Therefore the pdf  $f(Y|X = A)$  is Gaussian with mean  $A$  and variance  $\sigma^2$
- Similarly the pdf  $f(Y|X = -A)$  is Gaussian with mean  $-A$  and variance  $\sigma^2$

Note: Adding a constant to a random variable just shifts the mean, does not change the shape of the distribution

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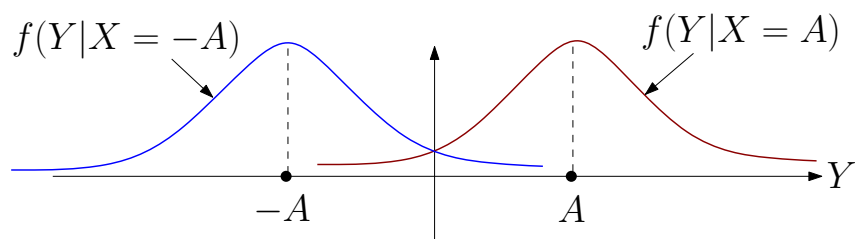
Let  $\hat{X}$  denote the decoded symbol. When the symbols  $A$  and  $-A$  are *a priori* equally likely, the optimal detection rule is:

$$\begin{aligned} \hat{X} &= A \quad \text{if } f(Y | X = A) \geq f(Y | X = -A) \\ \hat{X} &= -A \quad \text{if } f(Y | X = -A) > f(Y | X = A) \end{aligned}$$

*“Choose the symbol from which  $Y$  is most likely to have occurred”*

- This decoder is called the **maximum-likelihood decoder**
- This decoder is intuitive and seems sensible, and is in fact, the optimal detection rule when all the constellation symbols are equally likely (we will not prove this here)
- It is then a special case of the Maximum a Posteriori (MAP) detection rule, which minimises the probability of detection error (discussed in 3F4)

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The detection rule can be compactly written as

$$\hat{X} = \arg \max_{x \in \{A, -A\}} f(Y|X = x)$$

$$\hat{X} = \arg \max_{x \in \{A, -A\}} f(Y|X = x)$$

$$= \arg \max_{x \in \{A, -A\}} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-(Y-x)^2/2\sigma^2} = \arg \min_{x \in \{A, -A\}} (Y - x)^2$$

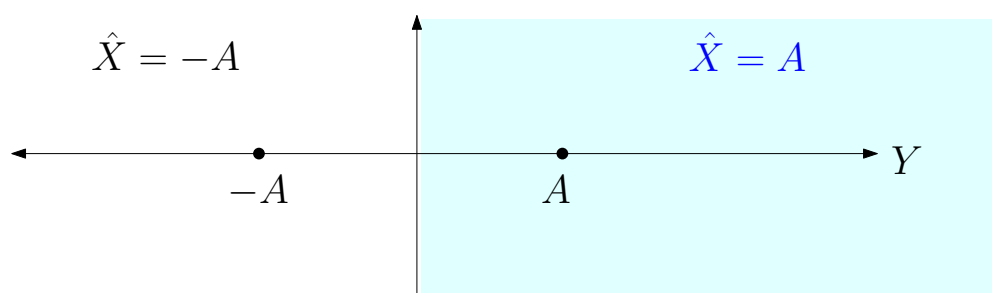
Thus the detection rule is just:  $\hat{X} = A$  if  $Y \geq 0$ ,  $\hat{X} = -A$  if  $Y < 0$   
*“Choose the constellation symbol closest to the output  $Y$ ”*

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## Decision Regions

The detection rule partitions the space of  $Y$  (the real line) into **decision regions**.

For binary PAM, we just derived the following decision regions:



Q: When does the detector make an error?

A: When  $X = A$  and  $Y < 0$ , or When  $X = -A$  and  $Y > 0$

We will calculate the probability of error shortly, but let's first find the detection rule for general PAM constellations

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# Detection for General PAM Constellations

The detection rule can easily be extended to a general constellation  $\mathcal{C}$

- E.g.,  $\mathcal{C}$  may be the 3-ary constellation  $\{-2A, 0, 2A\}$  or a 4-ary constellation  $\{-3A, -A, A, 3A\}$
- The maximum-likelihood principle is the same: “Choose the constellation symbol from which  $y$  is most likely to have occurred”

$$\begin{aligned}\hat{X} &= \arg \max_{x \in \mathcal{C}} f(Y|X=x) \\ &= \arg \max_{x \in \mathcal{C}} \frac{1}{\sqrt{2\pi\sigma^2}} e^{-(Y-x)^2/2\sigma^2} = \arg \min_{x \in \mathcal{C}} (Y-x)^2\end{aligned}$$

Thus, the detection rule for any PAM constellation boils down to:  
“Choose the constellation symbol closest to the output  $Y$ ”

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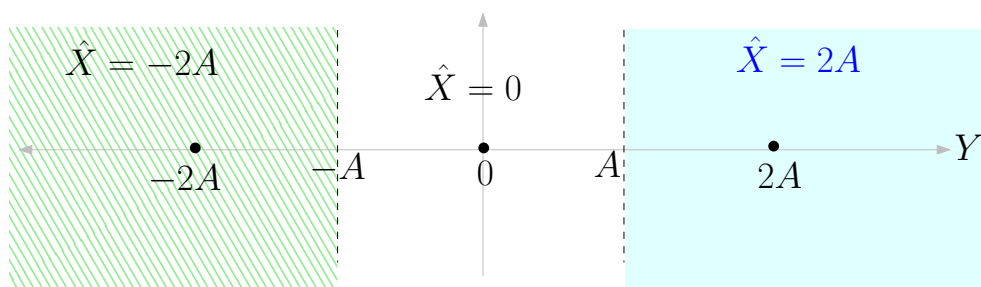
## Example: 3-ary PAM

$$Y = X + N, \quad N \sim \mathcal{N}(0, \sigma^2)$$

What is the optimal detection rule and the associated decision regions if  $X$  belongs to the 3-ary constellation  $\{-2A, 0, 2A\}$ ?

The “nearest symbol to  $Y$ ” decoding rule yields

$$\hat{X} = \begin{cases} -2A & \text{if } Y < -A \\ 0 & \text{if } -A \leq Y < A \\ 2A & \text{if } Y \geq A \end{cases}$$



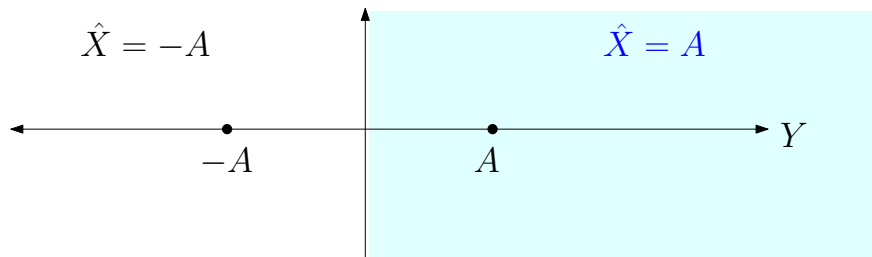
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# Probability of Detection Error

$$Y = X + N$$

Consider binary PAM with  $X \in \{A, -A\}$ . The decision regions are:



The detector makes an error when  $X = A$  and  $Y < 0$ , or when  $X = -A$  and  $Y > 0$

The probability of detection error is

$$\begin{aligned} P_e &= P(\hat{X} \neq X) \\ &= P(X = -A)P(\hat{X} = A | X = -A) + P(X = A)P(\hat{X} = -A | X = A) \\ &= \frac{1}{2}P(\hat{X} = A | X = -A) + \frac{1}{2}P(\hat{X} = -A | X = A) \end{aligned}$$

(The symbols are equally likely  $\Rightarrow P(X = A) = P(X = -A) = \frac{1}{2}$ ) 33 / 42

Let us first examine  $P(\hat{X} = A | X = -A)$

$$\begin{aligned} P(\hat{X} = A | X = -A) &= P(Y > 0 | X = -A) \\ &= P(-A + N > 0 | X = -A) \\ &= P(N > A | X = -A) \stackrel{(a)}{=} P(N > A) \end{aligned}$$

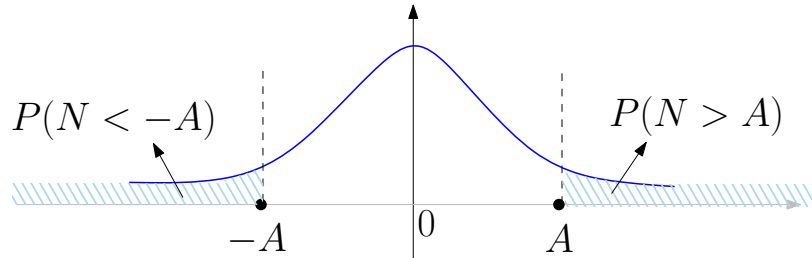
(a) is true because the noise random variable  $N$  is **independent** of the transmitted symbol  $X$ . Similarly,

$$\begin{aligned} P(\hat{X} = -A | X = A) &= P(Y < 0 | X = A) \\ &= P(A + N < 0 | X = A) \\ &= P(N < -A | X = A) = P(N < -A) \end{aligned}$$

The probability of detection error is therefore

$$\begin{aligned} P_e &= \frac{1}{2}P(\hat{X} = A \mid X = -A) + \frac{1}{2}P(\hat{X} = -A \mid X = A) \\ &= \frac{1}{2}P(N > A) + \frac{1}{2}P(N < -A) \\ &\stackrel{(b)}{=} P(N > A) \stackrel{(c)}{=} P\left(\frac{N}{\sigma} > \frac{A}{\sigma}\right) \end{aligned}$$

- (b) holds due to the symmetry of the Gaussian pdf  $\mathcal{N}(0, \sigma^2)$ :



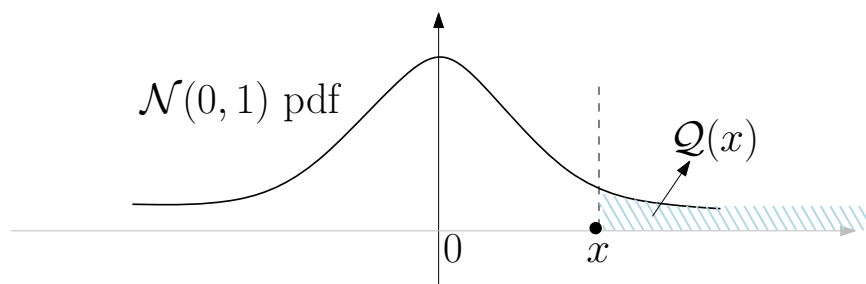
- In (c), we have expressed the probability in terms of a **standard** Gaussian random variable with distribution  $\mathcal{N}(0, 1)$
- Recall from 1B Paper 7 (Probability) that if  $N$  is distributed as  $\mathcal{N}(0, \sigma^2)$  then  $\frac{N}{\sigma}$  is distributed as  $\mathcal{N}(0, 1)$

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## The $Q$ -function

The error probability is usually expressed in terms of the  $Q$ -function, which is defined as:

$$Q(x) = \int_x^{\infty} \frac{1}{\sqrt{2\pi}} e^{-u^2/2} du$$



- $Q(x)$  is the probability that a **standard Gaussian**  $\mathcal{N}(0, 1)$  random variable takes value greater than  $x$
- Also note that  $Q(x) = 1 - \Phi(x)$ , where  $\Phi(\cdot)$  is the cdf of a  $\mathcal{N}(0, 1)$  random variable

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## $P_e$ in terms of the signal-to-noise ratio

The probability of detection error is therefore

$$P_e = P(N > A) = P\left(\frac{N}{\sigma} > \frac{A}{\sigma}\right) = \mathcal{Q}\left(\frac{A}{\sigma}\right) = \mathcal{Q}\left(\sqrt{\frac{E_s}{\sigma^2}}\right)$$

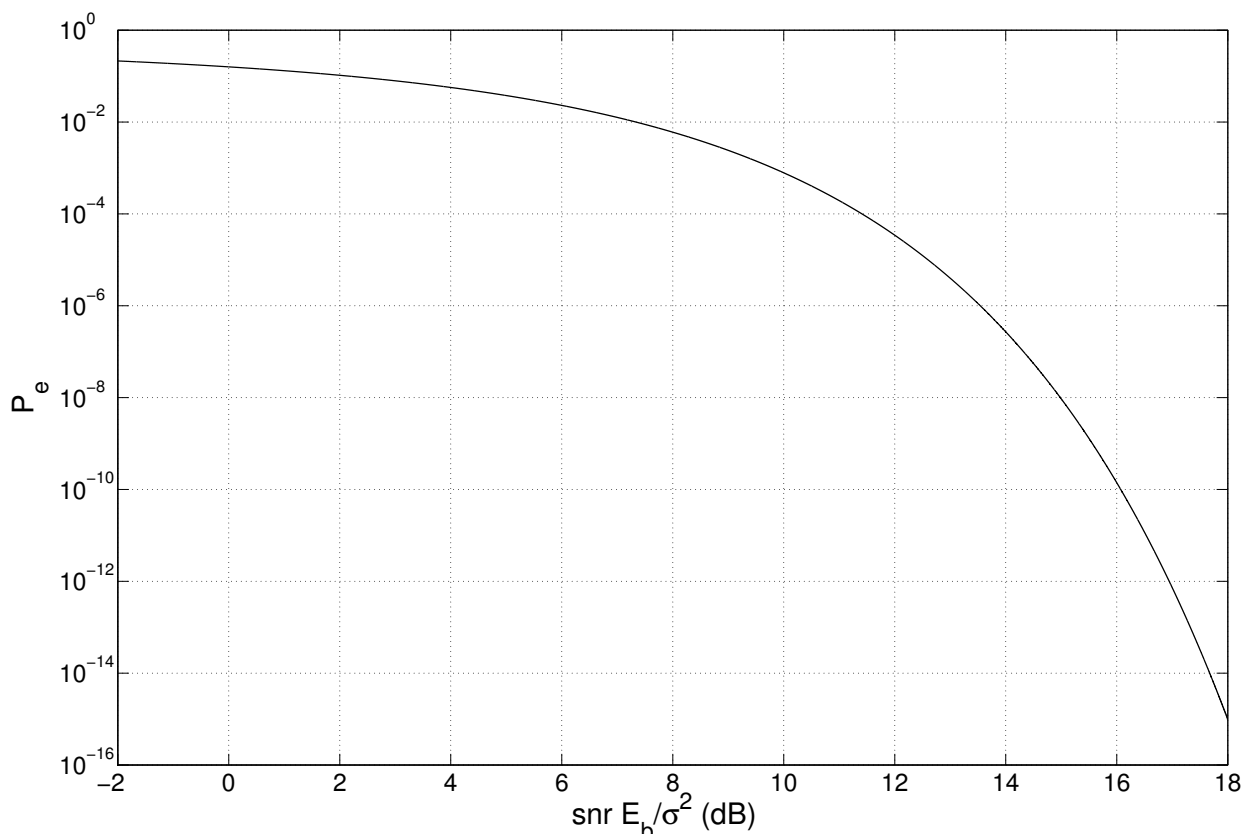
where  $E_s$  is the *average energy per symbol* of the constellation:

$$E_s = \frac{1}{2}(A^2 + (-A)^2) = A^2$$

- For a binary constellation, each symbol corresponds to 1 bit.  
 $\Rightarrow$  the *average energy per bit*  $E_b$  is also equal to  $A^2$  in this case
- For an  $M$ -point constellation, the average energy per symbol  $E_s = E_b \log_2 M$
- $\frac{E_b}{\sigma^2}$  is called the signal-to-noise ratio (snr) of the transmission scheme
- $P_e$  can be plotted as a function of the snr  $\frac{E_b}{\sigma^2} \dots$

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## $P_e$ vs snr for binary PAM



To get  $P_e$  of  $10^{-3}$ , we need  $\text{snr } E_b/\sigma^2 \approx 9$  dB

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# Error Probability vs Transmit Power

The probability of error for binary PAM decays rapidly as  $\text{snr} \uparrow$ :

- $Q(x) \approx e^{-x^2/2}$  for large  $x > 0 \Rightarrow P_e \approx e^{-\text{snr}/2}$

Can we set the  $\text{snr} \frac{E_b}{\sigma^2}$  to be as high as we want, by increasing  $E_b$ ? (i.e., by increasing  $A$  since  $E_b = E_s = A^2$  for binary PAM)

- The problem is that transmitted power also increases!
- Intuition: 1 symbol transmitted every  $T$  seconds with average energy  $E_s \Rightarrow$  transmit power is  $E_s/T$
- Thus as you increase the  $\text{snr}$ , you battery drains faster!

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## Power of PAM signal

$$x_b(t) = \sum_k X_k p(t - kT)$$

With any constellation the power of the baseband PAM signal  $x_b(t)$  is

$$\frac{E_s}{T} = \frac{E_b \log_2 M}{T},$$

where

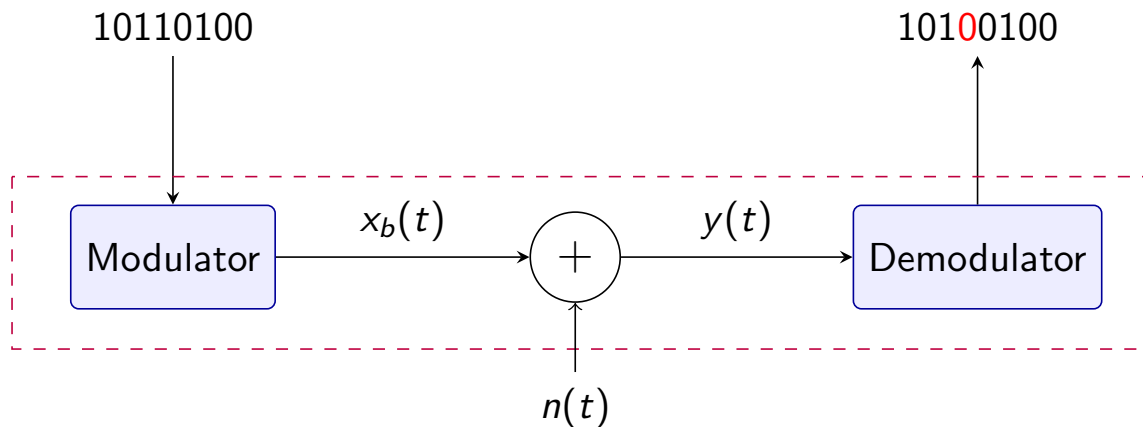
- $E_s$  is the average symbol energy of the constellation.
- $E_b$  is the average energy per bit

**Intuition:**

- In each symbol period of length  $T$ , a symbol with average energy  $E_s$  modulates a unit energy pulse
- A rigorous calculation of power has to take into account the fact that the transmitted symbols  $X_1, X_2, \dots$  are drawn *randomly* from the constellation (done in 3F4)

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# Pulse Amplitude Modulation - The Key Points



PAM is a way to map a sequence of information bits to a *continuous-time baseband* waveform

1. Pick a constellation, map the information bits to symbols  $X_1, X_2, \dots$  in the constellation
2. These symbols then modulate the amplitude of a *pulse shape*  $p(t)$  to generate the baseband waveform  $x_b(t)$

$$x_b(t) = \sum_k X_k p(t - kT)$$

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Desirable properties of the pulse shape  $p(t)$ :

- $p(t)$  should decay quickly in time; its bandwidth  $W$  shouldn't be too large
- *Orthonormal* shifts property for simple and reliable decoding

At the receiver, first **demodulate** then **detect**:

- The demodulator is a *matched filter* with IR  $h(t) = p(-t)$
- Matched filter output is *sampled* at times  $\dots, 0, T, 2T, \dots$   
At time  $mT$ , the output is

$$Y_m = X_m + N_m$$

$N_m$  is Gaussian noise with zero mean and variance  $\sigma^2$  that can be empirically estimated

- Detection rule:  $\hat{X}_m$  = the constellation symbol closest to  $Y_m$

Probability of detection error can be calculated:

- Decays exponentially with snr  $E_s/\sigma^2$
- $E_s$  is average energy/symbol of the constellation; power of PAM waveform  $x_b(t)$  is  $E_s/T$

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