UNIT-I Introduction to Ho-duap

Introducing Hadoop:

- a Big Data Seems to be the buzz word!
- a Enterprises, the world over, are beginning to realize that there is a huge volume of interpred information before them in the form of Structured, Semi-Structured, and unstructured dute.
- a this varied variety of idea is spread across the networks.
- > Few Statistics to get on idea of the amount of data which generated every day, every minute, and every second.

1. Every day:

- · NYSE (New York Stock Exchange) generate 1.5 billion shares and trade date.
- · Facebook stores 2.7 billion comments and Likes.
- · Google processes about 24 petu-bytes of dute.

2. Every minute:

- · Facebook users share nearly 2.5 million pieces of content.
- Instrgram users post nearly e, 20,000 new photos
- · You Tube users upload 72 hours of new video content.
- · Google receives over 4 million Search Queries.

3. Every/ Second:

· Bunking applications process more than 10,000 credit and transactions.

Duty: The Tree-Sure Trove:

- 1. provides business advantages such as generating product recommendations, inventing new products, analyzing the market, and many more.
- e) provides few early key indicators that can turn the fortune of
- 3) provides ruom for precise analysis, If we have more data for analysis, then we have greater precision of analysis.

I am inundated with date"

How to store tempytes of mounting date?

I need this dote to be processed ovicky. My decision is pending, " How to access the information ovickly?

I have date in varied sources.

I have date that is rich in variety - Structured, SemiStructured, unstructured." How to work with date that is so very different?

Fig: challenges with big volume, variety & velocity of data.

Why Hadoop?
This combility to handle massive amount of data, different categories of data - fairly ouickly.

Key considerations of Hadoop:
1) Low cost: Hadoop is an open-source framework and uses commodity

hardware to store enormous open-titles of data.

- e) Computing powers: Hadoop is based on distributed computing model which processes very large volume of data fairly processes very large volume of data fairly puickly. The more the number of computing nodes, the more the processing power at hand.
- 3) Scalability: Simply adding nodes as the system grows and requires much less administration.
- 4) Storing Steribility: unlike the traditional data-bases, in Hadoop data need not to be pre-processed beforing storing it.

 Hadoop provides the convenience of storing as much data as one needs and also, the added flexibility of deciding later as to how to use the stored data.

- 5) Inherent data protection: Ho-doop protects data and executing opplication against hardure failure. If a node fails, it automotically redirets the Jubs that had been assigned to this node to the other functional and available nodes and ensures that distributed computing does not fuil.
- -> Hodoor makes use of commodity hardware, distributed file System, and distributed computing, as shown below.

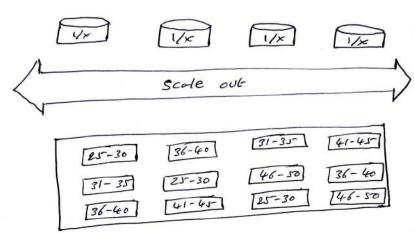


Fig: Hardoop framework (distributed file system)

- In this new design, groups of machine are gothered together; it is Known as a cluster.
- -> The duto can be managed with Hadoop as follows:
 - 1) Distributes the data and duplicate chunks of each data file across several nodes, for example 25-30 is one chunk of
 - 2) Locally available compute resource is used to process each chunk of dato in parallel.
 - 3) Hadoop framework handles failover smoothy and automatically.

Why Not ROBMS:

- -> RDBMS is not suitable for storing and processing large files, images and videos.
- -> ROBMS is not a good choice when it comes to advanced analytics involving machine learning, with respect to cost & Storage.

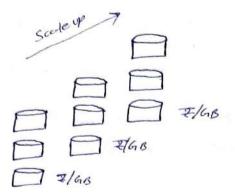


Fig: ROBMS with respect to cost/GB of Storage.

RDBMS Vers	sus Hadoop:	
Parameters U System	ROBMIS Relational database Management system	Mode besed flet structure
e) Data	Suitable for Structured	Ison, text be sed flut file etc.
3) Processing	OL TP	Analytical, Big Dota processing
4) choice	The do-to- needs consistent relationship	Big Deter processing, which does not recuise only consistent relationship b/w date.
s) processor	Need expensive hastwo-re or high-end processors to store huge volumes of deter	In - hadoup cluster, a node require only a
6) cost	cost a sound \$10,000 to \$40,000 per tembytes of slowinge	cost asom \$ 4,000 per wasytes of slosage,

History of Hadoop: - Hadoop are created by Doug cutting", the creator of Apoche Lucene (a commonly used text sourch Library).

- Hadog is a part of Apriche Notch (Yahoo) Project (an open-source web search engines, and also a part of the Lucene Project.

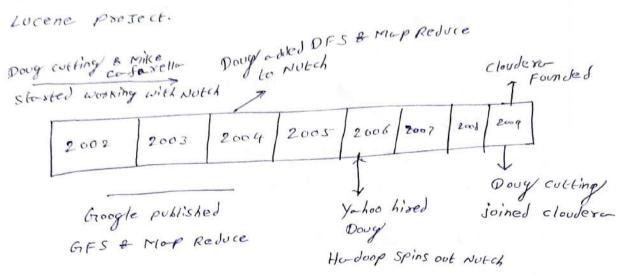


Fig: History of Hadoop

- -> The "Name Hadoop" is not an acronym; it's made-up
- How the Name come about: "The name my kid gave a Stuffed Yellow elephont. Short, relatively easy to spell & pronounce, meaningless, & not used elsewhere: those are my noming criteria. Kids are good at generating such. Googral is a Kid's term".

Hadoop overview:

- -> open Source Suftware framework to store and process massive amount of duta in a distributed fushion on large clusters of commodity hosdowse.
- -> Busicully, Hardoop accomplishes two tasks:
 - 1) Massive data storage
 - 2) Foster duto processing.

Key aspects of Hadoop!

- 1) Open Source Software: It is free to download, use and contribute to.
- E) Framework: Means everything that you will need to develop & execute and application is provided programs, tools, etc.
- 3) Distributed: Divides and Stores data across multiple computers. Processing is time in parallel across multiple connected nodes.
- 4) Massive Storage: Stores large amount of data across nodes of 10w-cost commodity hardware.
- 5) Faster processing: Large amounts of data is processed in parallel.

Hadoop Components:

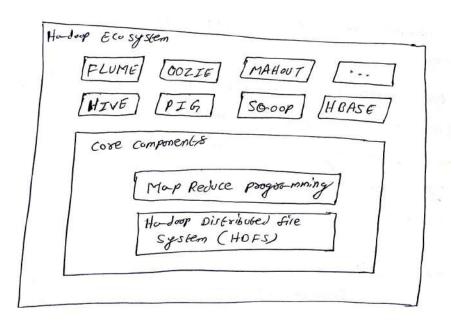


Fig :- Hadrop Compenents

1) Hardoop core components:

a) HOFS !-

- · Storye components
- · Distributes dute across everal modes
- · Natively redundant.

(b) Map Reduce 1.

- · Computational framework
- · Splits a lask across multiple nodes
- · processes duto in parallel.

Plto-doop Eco System:

- Hadoop Eco system are support projects to enhance the functionality of Hudoup case compenents.
- -> The Eco Projects are as follows: 4) HBASE 7) MAHOUT
 - 1) HIVE

- 2) PI 67
- 5) FLOME
- 3) SQ 00P
- () 00ZZE

Hadoop conceptual Layoris

- It is conceptually divided into:
 - Date Storige Loyen
 - 2) Data Processing Layer.

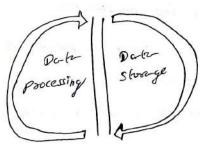


Fig: Hadoop conceptual Lugar

- Duto Storage Loyer: Stores huge volume of doto
- Duta processing Layer: processes duta in parallel to extract sicher & meaningful insights from date.

High - Level Architecture of Hadoop:

- -> Hadoop is a distributed "Moster-slave Architecture"
- -> Moster node is known as "Name wode".
- -> Slave Node is known as Date Node".

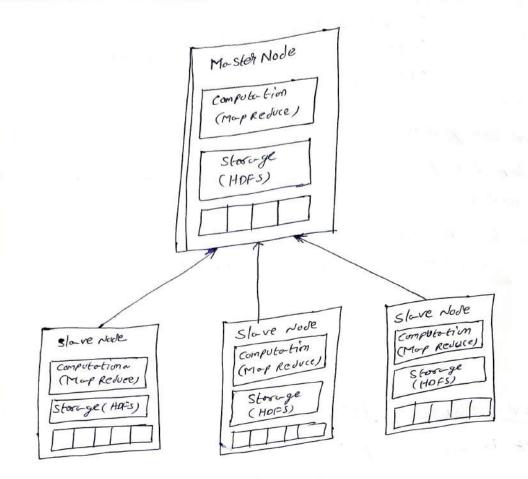
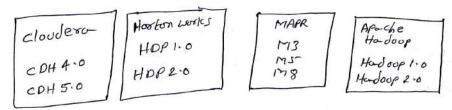


Fig: - Hodoop high-level architecture

- -> The Key components of the Master Mode:
 - 1) Moster HDFS: It's Main responsibility is postitioning the duto storage across the slave nodes. It also keeps track of location of data on Data Modes
 - 2) Master Map Reduce +- It decides and schedules computation task on slave modes.

Ho-doop Distributors:

-> It provides products that include Apache Handoup, Commercial support, and/or tools and utilities related to Hadoop.



Eig: Ho-doop Distributors.

Features of Hadoop:

- It is optimized to handle massive aventities of structured, semi-Structured, & instructured dute, using commodity hardware, that is, relatively inexpensive computers.
- -> Hodoop has a shared nothing architecture.
- -> It replicates its do-to across multiple computers so that, if one goes down, the so-te an still be processed from a-nother machine that stores its replica.
- Hudoop is for high throughput rather than low latency, It is a butch operation handling massive overtities of data. there fore the response time is not immediate.
- -> Its complements "OLTP" and "OLAP". However, it is not a replacement for a "ROBMS".
- It is not good when work connect be porollelized (or) when these are dependencies within the data.
- -> It is not good for processing/ small files. It works best will hoge dota files & sotusets.

use case of Hadoop:

- -> click Stream do-to helps you to understand the purchasing behaviour
- y click stream analysis helps online marketers to optimize their product web pages, promotional content, etc, to improve their business.

21101	date analysis using H	
oin click Stream	Stores years of dote	Hive on Pig
in clickson	asithout much	Script to
to with CRM & Sales data.	incremental cost.	onolyze data.

Fig: click stream do to Anolysis.

Key benefits:

- -> click stream duta analysis using Hadoop.
 - DIT helps to join clickstream do-to with other data
 Sources such as customer Relationship Management
 do-to- (customer demographics data, sales do-to-, &
 Information on Advertising Compaigns). This additional duta
 often provides much information to understand customer
 behaviour.
 - 2) It's Sculo-bility property helps you to store years of data without much incremental cost. This helps you to perform temporal (08) year over year analysis on click stream data which your competitors may miss.
- 3) Business analysts can use "Hive" or pig" for website analysis. with these tools, you can organize click stream data by user session, refine it, and feed it to visualization or analytics tools.

Hadoop Distributed File System (HDFS):-

Key Points of HDFS:

- 1) Slosuge Component
- 2) Distributed File System
- 3) Modeled after google file System
- 4) Optimized for high throughput
- 5) you can replicate a file for a configured number of times, which is tolerant in terms of both s/w & h/w.
- 6) Re-Replicate data blocks outomotivally on nodes that have failed.
- 7) you can realize the power of HDFS when you perform read (30) write an large files.
- 8) Sits on top of notive file system such as ext3 & ext4.

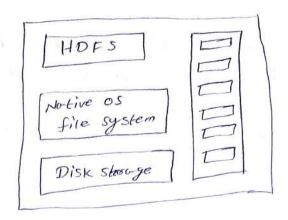
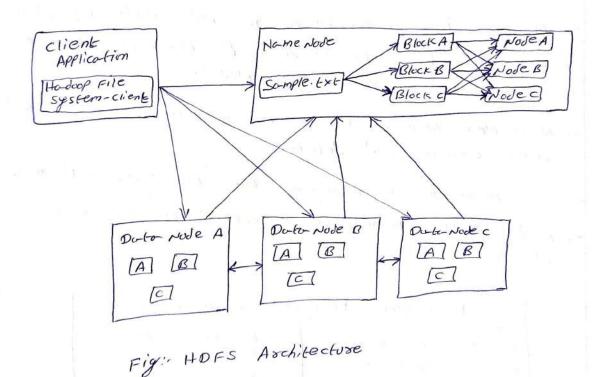


Fig: Hadoop Distributed file system.

Ha-doop Di	stributed file sy	1 2 6 11
Block .	Default Replication	Default Block
st-suctured file	factor : 3	Size: 64 ME

Fig: HDFS Key Points



-> client Application interacts with Name Node for meta data related Activities.

^{-&}gt; client Application communicates with Data Node to read & writer

^{-&}gt; Data-Node converse with each other for pipeline reads & writers.

For example:

Let us assume that the sile "Sample txt" is size of "192178.

As per the default Block size (64 MB), it will be split into three blocks and replicated across the nodes

on the cluster bused on the default replication fuctor.

HDFS Daemons:

is not interactive.

Name Node:

- > HDFS breaks a large file into smaller pices called blocks.
- -> Name node uses "rock ID" to identify "Data rodes" in the rock.
- -> A rack is a collection of Data Nodes within the cluster.
- it is placed on various Duta Modes.
- -> Name Node manages file related operations such as read, write, execte, and delete.
- Its main job is managing the File System Name space.
- A file system no-mespace is collection of files in the cluster.

Hame Node - Man	ages Sile Related operations
FSIMage - Fire, in which	Edit by - Records every transaction that occurs to file system meta-date.

Fig: Nome Node

-> No-me Node uses stores "HDFS no-me space".

File system name space includes mapping of blocks to file file properties and is stored in a file called "FS Impg".

- -> Nome Node uses a Editlog to second every transaction that happens to the file system meta-data.
- -> Nume Node Starts up, it reads FS Image and Editlog from disk & applies all transactions from the Editlog to inmemory representation of the FS Image.
- Then it Flushes out new version of FS Image on disk & truncates the old Edit-Log because the changes are updated in the FS I mage ".

De-tr-Node: -

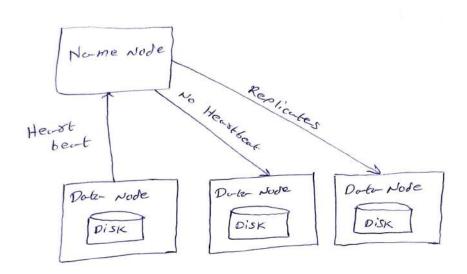


Fig: - Communication with Name Node & Duty Node.

- There are multiple Dute Node per cluster.
- -> During Pipeline read & write Data rudes communicate with euch other.
- A Data Note also continuously sends "heartbeat" message to Name Node to ensure connectivity b/w the Name Node & Duto Node.
- In case there is no "heartbeat" message from date Node, the Nume rose replicates that date node within the cluster & Keeps on running as is nothing had happened.

- Secondary Nome Node:
- The Secondary Nome Node takes a Snopshot of HDFS metadata at intervals specified in the Hadoop configuration.
- T Since the memory requirements of secondary themewade are the same as Namewoode, it is better to run Namewoode & Secondary Namewoode on different machines.
- Nome Node and be consigured monutally to bring up the cluster.
- the secondary Name Node does not second only real-time changes that happens to the HOFS meta-data.

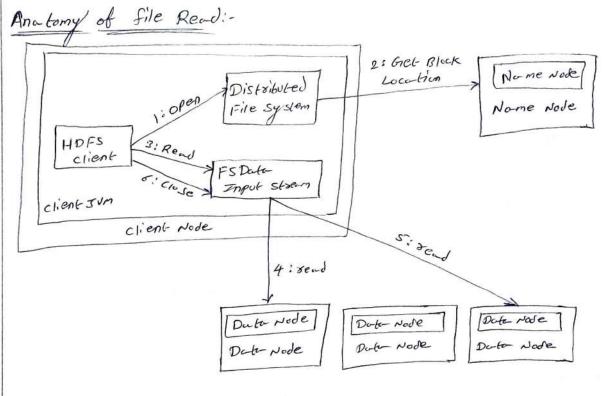
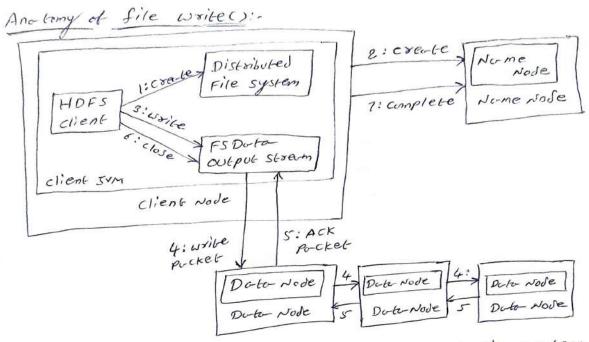


Fig: File Read

- D client opens the file to send from by "culling opens" on the Distributed file system.
- 2) Distributed file system communicates with the Name Node to get the location of data blocks, Name Node return with the address of data nodes that the data blocks are stored on subsequent to this, the "DFS" returns on FSData Input-Straum to client to read from the file.
- 3) client then co-11s read() on the stream DFS Input stream, which has address of the Duta made for the first few blocks of the file, connects to the closest Data mode for the first block in the file.

- 4) client calls reads () repeatedly to stream the data from the Dute Node
- s) when end of the block is reached, DFS Input Stream closes the connection with the Duta-Nove. It repeats the steps to find the best Dute node for the next block & subsequent blocks.
- 1) The client completes the reading of the file, it calls closer) on the FSData-Input Stream to close the connection.



- 1) The client culls & create () on Distributed file system.
- 2) An Remote Procedure calls (RPC) to the Name Mode happens through the Distributed File system to wente a new file. The Name Node performs various checks to create a new file. Initially, the Name Node exectes a sile without associating any duter block to the file. The DFS returns on FS Data -Output Streem to the client to perform write.
- 3) The client writes duty, date is splits into packets by DFS output stream, which is then written to an internal queue called data overe. Data Streamer consumes the data overe. The data streamer request the name node to allocate new block by selecting a list of suitable Data nodes to store replices. The list of puter Hodes makes a pipeline.

- 4) Date Strand strang the provets to the Sinst Detectode in the pipeline. It stores provet a foodereds it to the Second date node in the pipeline. In the same way, the Second date node stores the packet & foodereds it to the third packet & foodereds it to
- 5) The internal overe, Description close manages on "Acknowledgement by extender. I hat are writing for the acknowledgement by extender. A parket is removed from the "Acknowledge many it it is acknowledge by all the detenders in the pipeline.
- en the streem.
- 7) This flushes all the remaining packets to the Dutarlose pipeline & with for relevant acknowledgment before communicating with the Name Node to inform the client that the creation of the file is complete.

Hadoop Replica Placement Strubeyy:

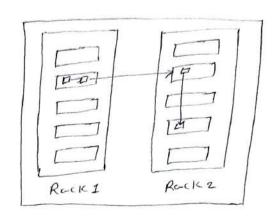


Fig: Replica Placement strategy

- is placed on the same node as the client.
- on different rock.
- -> It places third Replica on the same rack as second, but on a different node in the rack.
- -> once, replice locations have been set, a pipeline is built.

- Working with HDFS Commends: - The get the list of directories & files at the root of HDFS hadoop fs -15/ - To get the list of complete directories and files of HOFS hadoop fs -15 -R/
- -> To create a directory in HOFS (Sample) hadoop &s -mkdis/sample
- -> To display the contents of on HDFS file on console hadoop Is -cot / Sample / file. Ext
- -> To remove a directory from HDFS hadoop fs -sm-x/sample
- -> To copy a file from one directory to another on HDFS hadoop fs - cp /sample/file. txt /sample 1
- -> To copy or file from local file system to HDFS hadoop fs -put / sunt/sumple/test. txt / sample/test. txt -> To copy or file from HDFS to local file system
- hadoop for -get /sample/test. txt /sout/sample/test. txt

Special fectures of HDFS:-

- > Do-to Replication is a-bsolutely no need for a client application to track all blocks. It directs the client to the necrest replice to ensure high performance.
- -> Duta pipeline: client application writes a block to the first Dute Nude in the pipeline. Then this do-to-node takes over & forces the do-to to the next node in the pipeline. This process continues for all the doto blocks, & subsequently all the replicas are written to the disk.

Processing Data with Hadoop:

- -> Morp Reduce programming is a software framework.
- -> Map Reduce Programming helps you to process massive (on large amount of dato in po-80-11el.
- -s In Map Reduce programming, the input data-set is split into independent chunks.
- Map task" process these independent chunks completely in a provilet manner.
- -> The output of the impress are automatically shuffled and stored by the framework.
- -> Map Reduce framework souts the output based on keys.
- -> This sorted output becomes the input to the "yeduce to-sks"
- -> Reduce tousk provides reduced output by combining the output of the verious mappers. Job input and output are stored in a file system.
- -> Mc-P Reduce framework also takes care of other tasks Such as scheduling, monitoring, re-executing) failed tusks etc,
- -> HDFS & MOP REDUCE formework you on the some set of nodes. This configuration allows effective scheduling of tasks on the nodes where data is present Data Locality).

Mo-P Reduce Franework

phoses:

Map: converts input into key value paix.

Reduce: combines output of Mo-ppers & produces o- reduced result set.

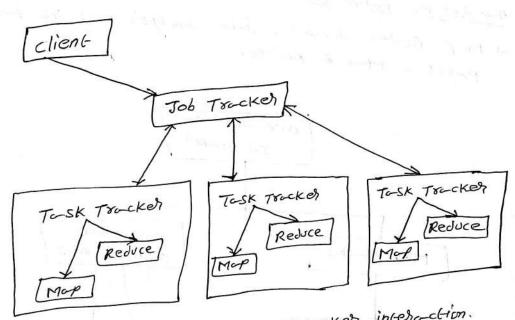
Da-emons:

Job tracker: Master Schedules bask Tosk Truckers slave, executes to-sk.

Fig: Map Reduce programming phases & daemons

- There are two doesnows associated with MapReduce Programming.
 - 1) A single master "Job tracker" per cluster & one slave "Task tracker" per cluster node.
 - 2) The "Job trucker" is responsible for scheduling tousks to the "tosk truckers", monitoring the tousk, & re-executing the trusk just in case the Trusk Trucker" fulls.

Map Reduce Daemons:



Figs Job tracker & Task tracker interaction.

- > once client submits a Jub to the Jub tracker, it partitions

 & assigns diverse Map Reduce tasks for each task Pracker
- In the cluster
 3 Tob tracker: provides commectivity blus Hadoop & application.

 3 Tob tracker: provides commectivity blus Hadoop & application.

 When you submit code to cluster, Job tracker creates, the

 When plan by deciding which task to assign to which

 execution plan by deciding which tasks.

 node. It also monitors all the running tasks.

 Nhen a task fails, automatically re-schedules the task

 when a task fails, automatically re-schedules the task

 when a task fails, automatically re-schedules for

 to a different node after a predefined number of

 to a different node after a master doesnow responsible for

 retries. Job Tracker is a master doesnow responsible for

 retries. Job Tracker is a master doesnow responsible for

 retries.

That k Trucker :- responsible for executing individual test that is assigned by Job Ecocker. There is a signed by Job Ecocker. There is a single Task Trucker per share and spouns multiple (JVMs) to handle multiple map at reduce tousks in partiel. "Task Trucker" continuously tousks in partiel. "Task Trucker" continuously sends "heart beat" message to Job Erecker. Sends "heart beat" message to receive "heart beat" when the Job trucker has failed to receive "heart beat" message from a Task Trucker, the Job tracker assumes that the bash trucker has failed a resubmits the task to another overlieble node in the cluster.

How does map Reduce was ?:

-> map Reduce divides and do conclysis tousk into the parts - Map & Reduce.

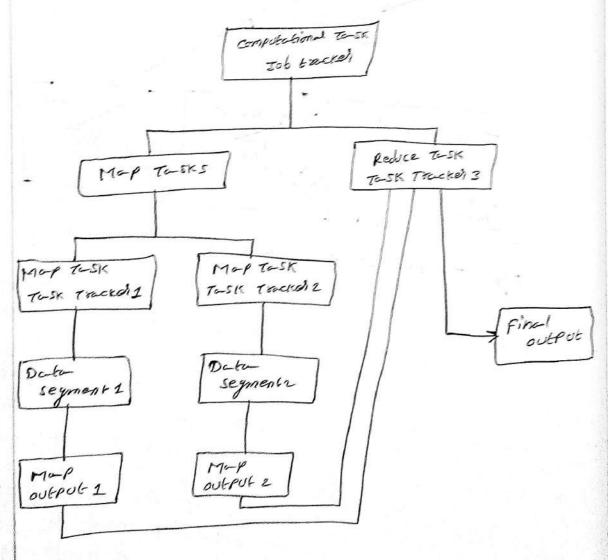


Fig: Map Reduce Programming flow

Example: 140-pper

There are two mappers and one reducer. Each mapper works on the partial dotaset that is stored on that node & the reducer combines the output from the mappens to produce the reduced result set.

working model of mop Reduce Programing:

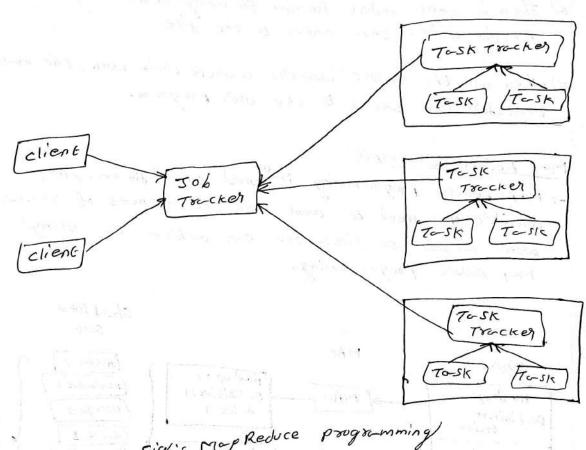


Fig: Map Reduce programming orchitecture

- 1) The input data-set is split into multiple pieces of data. · (small subset).
- 2) The framework creates a moster & several workers processes and executes the worker processes remotely.
- 3) several map to-SKS work Simultaneously and read pieces of dolon that are assigned to each map task. The map worker uses the most function to extract only) those doton that are present on their server a generatest key/volve pair for the extended doto.

- 4) Map worker uses partitioner function to divide the date into regions partitioner decides which reducer should get the output of the specified mapper.
- s) when the map workers complete their work, the master instructs the reduce workers to begin their work, the reduce workers to begin their work, the key/value in turn contact the map workers to get the key/value data for their partitioner.
- 6) Then it calls reduce function for every unique key. This function writes the output to the file.
- 4) when all the reduce workers complete their work, the master transfers the control to the user program.

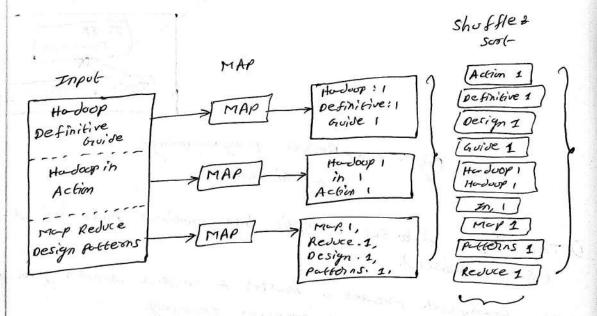
Mor Reduce programming is word count for example,

mor Reduce programming is word count for example,

consider you need to count the occurrences of similar

words o-cross so siles. You can achieve this using

Map Reduce programming.



Reduce

Action 1
pessign 1
pessign 1
puive 1
Handoopp

In 1
Map 1
putterns, 1
Reduce, 1

- NO SOL Stonds for NOE only SOL - No SOL is developed in the year 1998 by corto Strozzi. . No sol is a light - weight, open - source, more la limed data-ba-se.

Features of Nosal duta hase as follows:

- 1) They orse open-source
- e) They are non-relational
- 3) They are distributed
- 4) They are schema-less
- s) They are cluster-friendly
- 6) They are born out of 21st century web application.

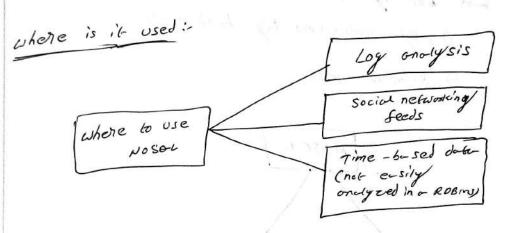


Fig: where to use NOSOL

+ Nosal databases are videly used in big data & other - NOSOL dotabases is used to strock log dota which con real-time web applications.

be pulled for analysis.

- they are hugely popular today owing to their ability to scale out or scale horizontally & the adepthess at dealing with - rich vo-riety of date: structurely semf-structured &

instructioned Jahr.

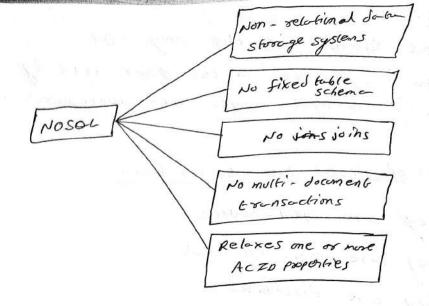


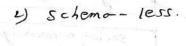
Fig: NOSOL: What it is

Types of NOSOL dolo-bases:

NOSOL duto-bases are non-relational.

They are broadly classified into the following:

Ney -value (00) big hash table



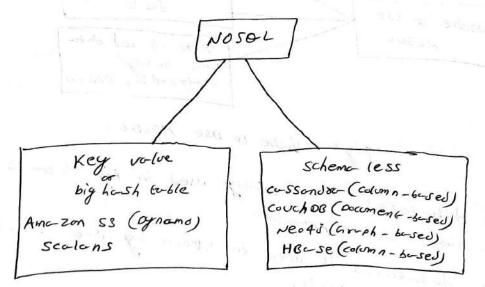


Fig: Types of NOSOL Sambuses

1) Key - Value 1.

It maintains a big hash table of key & values.

For example: Dynamo, Redis, etc.,

Sample key - Value pair in key - value Databases

Last Nome Adam

It maintains data in collections constituted of documents. 2) Document: For example: Mongo DB, couchbase, Apache couch db etc.

> "Book Name": "Big date" year " : " 2007 " " publisher" : " wiley"

- Each storage block has date from only one column. for example: - Co-ssandow, HBuse, etcs
- They are also called network database. A graph stores data in nodes. For example: Neo 45, Hyper Grouph DB, etc.,

- It has scale out a schitecture instead of the monolithic architecture of relational data-bases.

of It can house large volumes of structured, semi-structured, 2

- + NOSOL database allows insertion of data without pre-Setines
- -> It automotically spreed date across on arbitrary. number of servers.
- It offers good suppost of Jute replication which in turns gurantees high availability, fault tolerance, & disaster recovery.

Advantages of NOSOL:

- 1) cheap, easy to implement
- 2) Easy to distribute
- 3) can easily scale up & down
- 4) Relaxes the data consistency requirement
- s) Doesn't requires a pre-defined schema
- 6) Dute can be replicated to multiple nodes & can be partitioned.

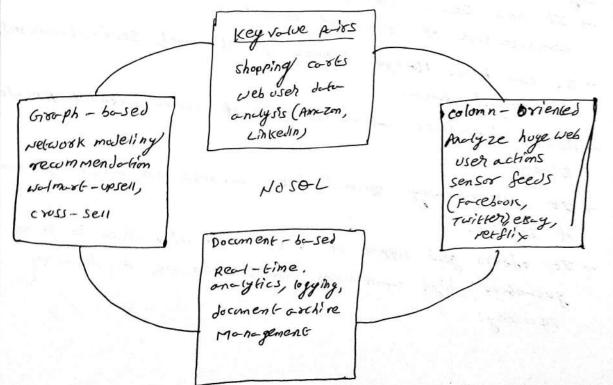
What we miss with NOSOL:

- 1) NOSOL does not support joins.
- 2) It does not have provision for ACID properties of transactions.
- 3) NOSOL does not have a standard SOL Interface but NOSOL data-bases such as Mongo DB & couch DB cassandra have their own rich overy Longuage to compensate for the lack of it.

4) one thing which is dearly missed is the easy integration with other applications that suppost sect.

use of NOSOL in industry:

Figure of Nosot in industry.



company	product	MOSE wisely used by
Amazon	DynonoDB	Linkedzn, Mozina
Fucebook	co-ssandre	Net flix, Twitter, ely
Groogle	Big To-ble	Adobe photoshop

SOL VS NOSOL :-

NO SOL VENDORS:

SO-L:

- 1) Relational detabase, model
- e) pre-defined scheme, toble bused dutubuses
- 3) vertically scalable (by increasing) system resources)
- 4) Not preferred for large dutusets, a best fit for hierarchical do-ta.
- 5) Emphasis on ACID properties
- 6) Excellent Suppost from Vendors
- 7) Supports complex averying & date kepping needs, can be consigured for strang consistency
- & Ex: Mysol, oracle, postgresol, DB2, etc,

- 1) Non-relutional, distributed database, & model-less approach
- 2) Dynamic scheme for unstructured duter
- 3) Document bused, graph bused, wide column store + key- value pair detersets, & it's horizontally scalable.
- 4) uses instructured every Linguize (UnoL), It's preferred for large
- 5) Best fit for hierarchical storage as it follows the key-value pair of storing data similar to ISON.
- c) Follows Brewer's CAP theorem. & relies heavily on community Suppost.
- 7) doesn't have good suppost for complex overying.
- 8) Few suppost strong consistency (e.g: Mongo DB), some others con be consigured for eventual consistency (e.g. cassandow).
- 9) Exir Mongo DB, HBOSE, conssendor, Neo 41, couch DB, couch bose, etcy